

(proof of unique interpolation - Newton form)

Let k be the number of nodes.

$$(k=1) \quad p_0 = y_0 = C_0$$

$$(k=2) \quad p_1 = C_0 + C_1(x-x_0)$$

$$\text{From } y_1 = p_1(x_1) = C_0 + C_1(x_1 - x_0)$$

$$C_1 = [y_1 - C_0] / (x_1 - x_0)$$

$$(k=3) \quad p_2 = p_1(x) + C_2(x-x_1)(x-x_0)$$

This construction allows us to only consider the new condition

$$p_2(x_2) = y_2 \quad \text{b/c we already know}$$

$$p_1(x_0) = y_0 \quad \text{and} \quad p_1(x_1) = y_1 \quad \text{while}$$

$C_2(x-x_1)(x-x_0)$ vanishes if $x=x_0$ or x_1 .

$$\text{From } y_2 = p_2(x_2) = p_1(x_2) + C_2(x_2-x_1)(x_2-x_0)$$

$$C_2 = [y_2 - p_1(x_2)] / (x_2-x_1)(x_2-x_0)$$

Repeat this until $k=n$. Since the process is constructive, the resulting $p_n(x)$ must be unique.

Example: Find $p \in \Pi_3$ such that

x	0	1	2	3
y	1	1	2	3

$$p_0 = 1 = C_0$$

$$p_1(x) = C_0 + \frac{y_1 - C_0}{x_1 - x_0} = 1 + \frac{1 - 1}{1 - 0} (x - 0) = 1$$

$$p_2(x) = p_1(x) + \frac{y_2 - p_1(x_2)}{(x_2 - x_1)(x_2 - x_0)} (x - x_1)(x - x_0)$$

$$= 1 + \frac{2 - 1}{2 \cdot 1} (x - 1)x = 1 + \frac{x}{2} (x - 1)$$

$$p_3(x) = p_2(x) + \frac{y_3 - p_2(x_3)}{(x_3 - x_2)(x_3 - x_1)(x_3 - x_0)} (x - x_2)(x - x_1)(x - x_0)$$

$$= 1 + \frac{x}{2} (x - 1) + \frac{3 - 1 - \frac{3}{2} \cdot 2}{1 \cdot 2 \cdot 3} (x - 2)(x - 1)x$$

$$= 1 + \frac{x}{2} (x - 1) - \frac{x}{6} (x - 2)(x - 1)$$

p_0

p_1

p_2

p_3

Horner's algorithm

number of
multiplications

$$c_0 + c_1 d_1 + c_2 \underline{d_1 d_2} + c_3 \underline{d_1 d_2 d_3} \quad 1+2+3$$

$$= c_0 + c_1 \underline{d_1} + (c_2 + c_3 d_3) \underline{d_1 d_2}$$

$$= c_0 + (c_1 + (c_2 + c_3 d_3) d_2) d_1$$

nested $\frac{\alpha}{c's} + \underbrace{\beta r}_{\substack{\text{previous} \\ \text{result}}} \text{ form}$

$$= c_0 + (c_1 + \underbrace{(c_2 + c_3 d_3)}_{\alpha_1 + \beta_1 r_1}) d_2 \quad 1+1+1$$

$$\underbrace{\alpha_2 + \beta_2 \cdot r_2}$$

$$\underbrace{\alpha_3 + \beta_3 \cdot r_3}$$

example : Lagrange basis subordinate to
nodes $x_0 = 1, x_1 = 2, x_2 = 3, x_3 = 4$

$$l_0(x) = \frac{(x-x_1)(x-x_2)(x-x_3)}{(x_0-x_1)(x_0-x_2)(x_0-x_3)} = -\frac{1}{6} (x-2)(x-3)(x-4)$$

$\begin{matrix} -1 & -2 & -3 \end{matrix}$

$$l_1(x) = \frac{(x-x_0)(x-x_2)(x-x_3)}{(x_1-x_0)(x_1-x_2)(x_1-x_3)} = \frac{1}{2} (x-1)(x-3)(x-4)$$

$\begin{matrix} 1 & -1 & -2 \end{matrix}$

$$l_2(x) = \frac{(x-x_0)(x-x_1)(x-x_3)}{(x_2-x_0)(x_2-x_1)(x_2-x_3)} = -\frac{1}{2} (x-1)(x-2)(x-4)$$

$\begin{matrix} 2 & 1 & -1 \end{matrix}$

$$l_3(x) = \frac{(x-x_0)(x-x_1)(x-x_2)}{(x_3-x_0)(x_3-x_1)(x_3-x_2)} = \frac{1}{6} (x-1)(x-2)(x-3)$$

$\begin{matrix} 3 & 2 & 1 \end{matrix}$

(proof of linear independence)

Assume

equality as fn's.


$$a_0 f_0(x) + a_1 f_1(x) + \dots + a_n f_n(x) = 0$$

plug in $x = x_0$, then

$$\begin{aligned} 0 &= a_0 \underbrace{f_0(x_0)}_{\delta_{00} \quad 1} + a_1 \underbrace{f_1(x_0)}_{\delta_{10} \quad 0} + \dots + a_n \underbrace{f_n(x_0)}_{\delta_{n0} \quad 0} \\ &= a_0. \end{aligned}$$

plug in $x = x_1$, then

$$\begin{aligned} 0 &= a_0 \underbrace{f_0(x_1)}_{\delta_{01} \quad 0} + a_1 \underbrace{f_1(x_1)}_{\delta_{11} \quad 1} + \dots + a_n \underbrace{f_n(x_1)}_{\delta_{n1} \quad 0} \\ &= a_1. \end{aligned}$$

Repeat this to conclude $a_0 = a_1 = \dots = a_n = 0$ 

(proof of interpolation error)

① If x is one of the nodes we have

$$0 = 0. \checkmark$$

② Assume $x \neq x_i$ ($i=0,1,2,\dots,n$).

** The trick is (due to Cauchy) to think of x as a new node. **

Put $w(t) = \prod_{i=1}^n (t - x_i)$, then $w(x) \neq 0$.

Let $\lambda = (f(x) - p(x))/w(x)$ and introduce

$$\varphi(t) = f(t) - p(t) - \lambda w(t) \in C^{n+1}[a,b]$$

Observe that $\varphi(x_i) = 0$ ($i=0,1,2,\dots,n$)

and $\varphi(x) = 0$ (by the construction of λ).

Thus, use Rolle's theorem $(n+1)$ times to argue $\exists \xi_j$ ($j=0,1,\dots,n$) s.t. $\varphi'(\xi_j) = 0$.

Next, do the similar to argue $\exists \zeta_k$ ($k=0,1,\dots,n-1$)

s.t. $\varphi''(\zeta_k) = 0$. Repeat this to show that

$\exists \xi_x$ s.t. $\varphi^{(n+1)}(\xi_x) = 0$. (see picture

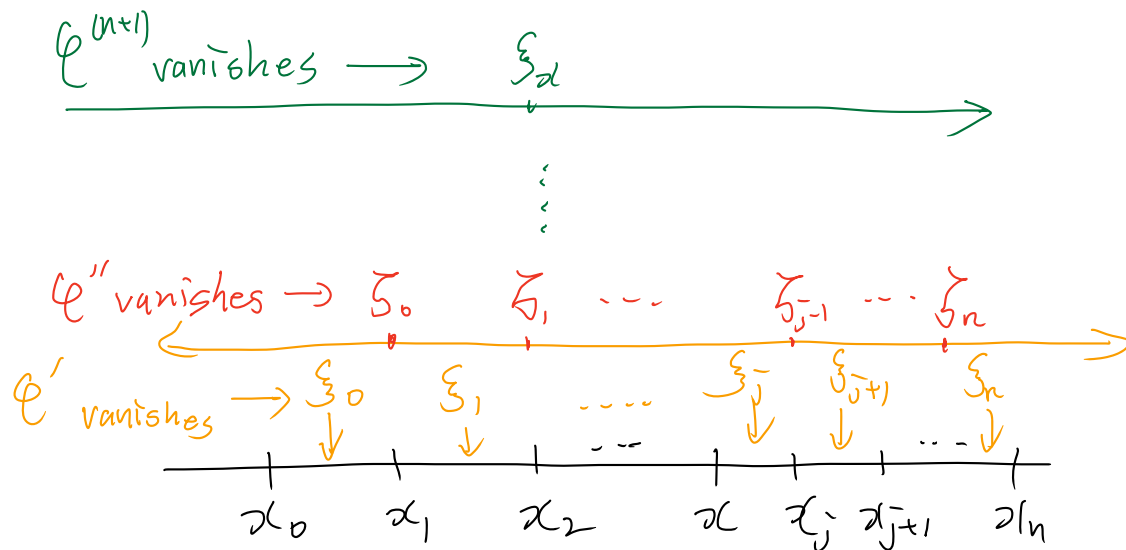
below). But

$$0 = \varphi^{(n+1)}(\xi_x) = f^{(n+1)}(\xi_x) - \lambda (n+1)! \quad (\text{why?})$$

$$= f^{(n+1)}(\xi_x) - \frac{(f(x) - p(x))}{\omega(x)} \cdot (n+1)!$$

Rearranging, we obtain

$$f(x) - p(x) = \frac{f^{(n+1)}(\xi_x)}{(n+1)!} \prod_{i=0}^n (x - x_i)$$



example of error bound on sine interpolation.

$$|f^{(n)}(x)| \leq 1 \quad \text{for all } x \in [0, 1].$$

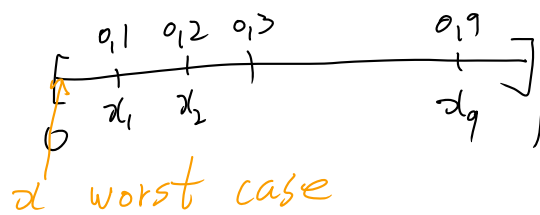
(an answer) Let us use the theorem we just proved.

$$\begin{aligned} |f(x) - p(x)| &= \left| \frac{f^{(11)}(x)}{11!} \prod_{i=1}^{10} (x - x_i) \right| \\ &\leq \frac{1}{11!} \prod_{i=0}^{10} \underbrace{|x - x_i|}_{\leq 1} \longrightarrow (*) \\ &= \frac{1}{11!} \approx 2.5 \times 10^{-8} \end{aligned}$$

If x_i 's are equally spaced, (*) continues

$$\leq \frac{1}{11!} \prod_{i=1}^{10} \frac{1}{10} = \frac{1}{10^{10} \cdot 11} \approx 9.1 \times 10^{-12}$$

very pessimistic bound $|x - x_0| \leq 1$ is used.



/* If you don't like this the pessimistic bound $|x - x_0| \leq 1$ (even though you that must be very small), you can do the following. For the first two nodes, you can model $|x - x_0| |x - x_1|$ as $|x(x - 0.1)|$ on $[0, 0.1]$

But we know

$$|x(x - 0.1)| \leq 0.5^2 = 0.0025$$

Then, the error bound reads

$$\begin{aligned} &\leq \frac{1}{11!} 0.0025 \sum_{i=2}^{10} \frac{1}{i!} \frac{1}{10} \\ &= \frac{0.0025}{10^9 \cdot 11} \simeq 2.27 \times 10^{-13} \end{aligned}$$

But when we bound errors, too precise calculations are not what we are after. Find a good balance. */

(proof of characterization of Chebyshev poly's.)
We argue by an induction.

$$\textcircled{1} T_0(x) = \cos(0 \cdot \cos^{-1}x) = \cos 0 = 1$$

$$\textcircled{2} T_1(x) = \cos(\cos^{-1}x) = x$$

$$\textcircled{3} \text{ Suppose } T_k(x) = \cos(k \cos^{-1}x) \text{ for } k=0,1,\dots,n.$$

$$/* \cos(A+B) = \cos A \cos B - \sin A \sin B$$

$$\Rightarrow \left| \begin{array}{l} \cos((n+1)\theta) = \cos(n\theta)\cos\theta - \sin(n\theta)\sin\theta \\ \cos((n-1)\theta) = \cos(n\theta)\cos\theta + \sin(n\theta)\sin\theta \end{array} \right.$$

$$+ \quad \underline{\hspace{10cm}}$$

$$\cos((n+1)\theta) + \cos((n-1)\theta) = 2\cos(n\theta)\cos\theta$$

$$\text{Take } \theta = \cos^{-1}x \text{ for } x \in [-1,1]:$$

$$\begin{aligned} \cos((n+1)\cos^{-1}x) &= 2\cos(n\cos^{-1}x) \overbrace{\cos(\cos^{-1}x)}^x \\ &\quad - \cos((n-1)\cos^{-1}x) \end{aligned}$$

$$= 2xT_n(x) - T_{n-1}(x)$$

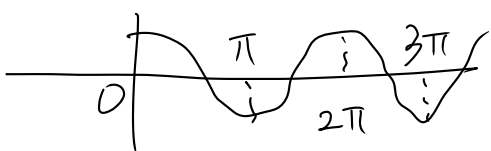
$$= T_{n+1}(x) \quad \text{by recursive} \\ \text{def. of Chebyshev}$$

(proof of basic properties of Chebyshev poly's.)

$$\textcircled{1} |T_n(x)| = |\cos(n \cos^{-1}x)| \leq 1$$

$$\textcircled{2} T_n(\cos \frac{j\pi}{n}) = \cos(\cancel{n} \cos^{-1}(\cancel{\cos \frac{j\pi}{n}}))$$

$$= \cos(j\pi)$$

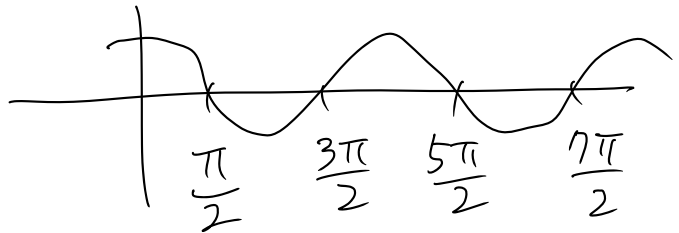
$$= (-1)^j$$


$$\textcircled{3} T_n(\cos \frac{2j-1}{2n} \pi)$$

$$= \cos(\cancel{n} \cos^{-1}(\cancel{\cos \frac{2j-1}{2n} \pi}))$$

$$= \cos(j\pi - \frac{\pi}{2})$$

$$= 0$$



(proof of maximum bound of monic poly.)
 (equality) Let $g(x) = 2^{1-n} T_n(x)$. This is monic
 and of degree n :

$$T_2(x) = 2x T_1(x) - T_0(x) = \underline{2x^2} - 1 \in \Pi_2$$

$$T_3(x) = \underline{2x} \underline{T_2(x)} - T_1(x) = \underline{2^2 x^3} - 3x \in \Pi_3$$

$$\vdots$$

and so on

$$T_n(x) = 2x(T_{n-1}(x)) - T_{n-2}(x) = 2^{1-n} x^n + \dots \in \Pi_n$$

Also, $|T_n(x)| \leq 1$ on $[-1, 1]$ by previous
 corollary.

Thus, for all $x \in [-1, 1]$, we have

$$|2^{1-n} \cdot T_n(x)| = 2^{1-n} |T_n(x)| \leq 2^{1-n}$$

(Inequality) Suppose $\exists p \in \Pi_n$, monic such that

$$\max_{|x| \leq 1} |p(x)| < 2^{1-n}$$

Let $x_j = \cos\left(\frac{j\pi}{n}\right)$. Then, observe

$$(*) \quad (-1)^{\bar{j}} \cdot p(x_{\bar{j}}) \leq |p(x_{\bar{j}})| < 2^{1-n} = (-1)^{\bar{j}} \bar{g}(x_{\bar{j}})$$

The last equality holds b/c

$$T_n(x_{\bar{j}}) = (-1)^{\bar{j}} \quad (\text{previous corollary})$$

$$\Rightarrow \bar{g}(x_{\bar{j}}) = 2^{1-n} \cdot T_n(x_{\bar{j}}) = 2^{1-n} (-1)^{\bar{j}}$$

$$\Rightarrow (-1)^{\bar{j}} \bar{g}(x_{\bar{j}}) = (-1)^{\bar{j}} \cdot (-1)^{\bar{j}} \cdot 2^{1-n} = 2^{1-n}$$

Rearranging (*), we have

$$(**) \quad (-1)^{\bar{j}} (\bar{g}(x_{\bar{j}}) - p(x_{\bar{j}})) > 0 \quad (\bar{j} = 0, 1, \dots, n)$$

$\Rightarrow g(x) - p(x)$ changes sign n times $\Rightarrow n$ roots

But $\deg(g(x) - p(x)) = n-1$ since x^n terms have been cancelled out. Besides,

$g(x) - p(x)$ is not zero function

since, e.g., $\bar{g}(x_0) - p(x_0) > 0$ from (**)

Thus, we have contradiction. \square

(proof of error bound with Chebyshev nodes)

$$|f(x) - p(x)| = \left| \frac{f^{(n+1)}(\xi)}{(n+1)!} \prod_{i=0}^n (x - x_i) \right|$$

$$\leq \frac{1}{(n+1)!} \max_{|t| \leq 1} |f^{(n+1)}(t)|$$

$$\cdot \max_{|x| \leq 1} \left| \prod_{i=0}^n (x - x_i) \right| \xrightarrow{\in \Pi_{n+1}}$$

minimized by $2^{-n} T_{n+1}$

, which is achieved by
choosing x_i 's to be the
zeros of T_{n+1} .

$$= \frac{1}{2^n (n+1)!} \max_{|t| \leq 1} |f^{(n+1)}(t)|$$