# Math 104A - Intro to Numerical Analysis

Numerical Differentiation and Integration

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# Numerical Differentiation and Integration

# **Numerical Differentiation and Integration**

Intro

# Numerical Differentiation

# Problem of interest

Given the access to function values can we suggest something close to f'(x)?

# **Numerical Differentiation and Integration**

**Numerical differentiation** 

# DIFFERENCE QUOTIENTS

The definition of the derivative,  $f'(x) = \lim_{h\to 0} \frac{f(x+h)-f(x)}{h}$ , motivates the following approximation.

# Forward/Backward difference quotient

For a small h > 0,

$$f'(x) \approx \frac{f(x+h) - f(x)}{h}$$
 (forward difference quotient)  
 $f'(x) \approx \frac{f(x) - f(x-h)}{h}$  (backward difference quotient)

■ Backward difference quotient is obtained by setting  $h \leftarrow -h$ .

# DIFFERENCE QUOTIENTS

# Theorem (Order of forward/backward difference quotient)

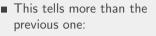
If f is smooth enough,

$$\frac{f(x+h) - f(x)}{h} = f'(x) + \frac{f''(\xi)}{2}h$$
$$\frac{f(x) - f(x-h)}{h} = f'(x) - \frac{f''(\xi)}{2}h,$$

where  $\xi \in (x, x + h)$ , and  $\hat{\xi} \in (x - h, x)$ .

# Proof.

Board work.



$$\frac{f(x+h) - f(x)}{h} - f'(x)$$
$$= \frac{f''(\xi)}{2}h$$

- "How fast" does not tell you "how close." For "how close," you also need  $f''(\xi)$ , which is usually not available.
- Usually, a faster method (i.e., higher order) is considered better if other factors are the same

# CENTERED DIFFERENCE QUOTIENT

# Theorem (Centered difference quotient)

If f is smooth enough and h > 0 is small enough,

$$\frac{f(x+h) - f(x-h)}{2h} = f'(x) + \frac{f'''(\xi)}{6}h^2$$

where  $\xi \in (x - h, x + h)$ .

# Numerical differentiation using interpolation

Given the nodes  $x_0, x_1, \dots, x_n$ , we know from Lagrange interpolation theorem

$$f(x) = \sum_{i=0}^{n} f(x_i) \ell_i(x) + \frac{1}{(n+1)!} f^{(n+1)}(\xi_x) w(x),$$

where  $w(x) = \prod_{i=0}^{n} (x - x_i)$  and  $\ell_i = \prod_{j \neq i}^{n} (x - x_j)/(x_i - x_j)$ . If we are interested in one of the nodes, say  $x_k$ , differentiating this and evaluating at  $x_k$ ,

$$f'(x_k) = \sum_{i=0}^{n} f(x_i) \ell'_i(x_k) + \frac{1}{(n+1)!} f^{(n+1)}(\xi_{x_k}) \prod_{\substack{j=0\\j\neq k}}^{n} (x_k - x_j)$$

■ This argument is not rigorously true though may be acceptable for a practical reason: see the next slide.

# Numerical differentiation using interpolation

If  $x \neq x_0, x_1, \dots, x_n$ , the previous argument doesn't work even in the practical sense.

$$f'(x) = \sum_{i=0}^{n} f(x_i) \ell'_i(x) + \frac{1}{(n+1)!} f^{(n+1)}(\xi_x) w'(x) + \frac{1}{(n+1)!} w(x) \frac{d}{dx} f^{(n+1)}(\xi_x)$$

Fortunately, there is a way to obtain a similar result but it take a bit more work. In particular, we have a good result for smooth function.

# Theorem (Suli and Mayers (2003))

If  $f \in C^{\infty}$ , letting  $M_{n+1} = \max_{x \in [a,b]} |f^{(n+1)}(x)|$ , for all  $x \in [a,b]$ ,

$$|f'(x)-p'_n(x)|\leq \frac{(b-a)^nM_{n+1}}{n!},$$

where  $p_n(x)$  is the Lagrange interpolation at  $x_0, x_1, \dots, x_n$ .

- The factor in red is problematic  $\xi_x$  may not smoothly depend on x.
- If interested, see Süli and Mayers, *An introduction to numerical analysis*, Theorem 6.5 and Corollary 6.1.

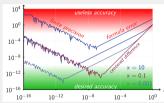
# CAUTIONS

**Caveat**: Rounding error may destroy everything if you pursue too much precision.

**Example**: (Centered difference quotient)

- $f'(x) = \frac{f(x+h)-f(x-h)}{2h} + \mathcal{O}(h^2)$  tells you the smaller h gets the more accurate the approximation is.
- Every time you store something you have rounding error. Thus, what you really compute is  $\frac{f(x+h)+e_1-f(x-h)-e_2}{2h}=\frac{f(x+h)-f(x-h)}{2h}+\frac{e_1-e_2}{2h}.$
- 2nd term may amplify the error as *h* gets smaller.
- Or, the numerator may just be zero to the computer since they are so close: leading to f'(x) = 0.

■ See Matlab example.



# RICHARDSON'S EXTRAPOLATION

If you have a numerical method in the following form, you can accelerate it using **Richardson's extrapolation**.

**Example**: Centered difference quotient (Board work) Start with Taylor expansion so that

$$f'(x) = \varphi(h) - a_2h^2 - a_4h^4 - \cdots,$$

where  $\varphi(h) = \frac{f(x+h) - f(x-h)}{2h}$ ,  $a_2 = \frac{f^{(3)}(x)}{3!}$  and  $a_4 = \frac{f^{(5)}(x)}{5!}$  and so on. Plug in  $h \leftarrow h/2$  and do something similar to something in middle school.

Then we get (Board work)

$$f'(x) = \frac{4}{3}\varphi(h/2) - \frac{1}{3}\varphi(h) - a_4h^4/4 - 5a_6h^6/16 - \cdots$$

- This process can be repeated to achieve higher accuracy, say put  $\psi(h) := \frac{4}{3}\varphi(h/2) \frac{1}{3}\varphi(h)$  and cancel  $h^4$ -term.
- This process can be applied to other methods that have the same structure.
- Subjective question: Does this look like a magic to you? Or, not exactly? Give the reason.

# **Numerical Differentiation and Integration**

Quadrature based on interpolation

# Numerical Integration (Quadrature)

#### Problem of interest

Given the access to function values can we suggest something close to  $\int_a^b f(x)dx$ ?

■ As you know, functions with integrals in closed form are very limited. E.g.,  $e^{-x^2/2}$ ,  $\cos(\sin(x^2))$ , etc. have no definite integral in a simple formula.

#### NEWTON-COTES FORMULA

# Definition (Newton-Cotes)

Given a < b, let  $x_i = a + hi$   $(i = 0, 1, \cdots, n)$ , where h = (b - a)/n, be equally spaced nodes on [a, b]. Then, the *Newton-Cotes* formula for the approximate integral of  $\int_a^b f(x) dx$  is given by

$$\sum_{i=0}^n A_i f(x_i),$$

where

$$A_i = \int_a^b \ell_i(x) \mathrm{d}x$$

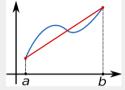
and

$$\ell_i(x) = \prod_{\substack{j=0\\j\neq i}}^n \frac{x - x_j}{x_i - x_j}$$

- Idea: replace *f* with something similar.
- In words, Newton-Cotes is a quadrature obtained from equally spaced Lagrange interpolation.
- Note that, once f is fixed,  $f(x_i)$ 's are data (they acts more like fixed numbers) while the "real" functions are  $\ell_i(x)$ 's.
- Notice that  $f \mapsto \sum_{i=0}^{n} A_i f(x_i)$  is a linear mapping.

## Trapezoidal rule

Trapezoidal rule  $\longrightarrow$  Newton-Cotes with only two nodes a,b



$$\int_a^b f(x)dx \approx \frac{b-a}{2}[f(a)+f(b)]$$

The error term is

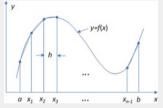
$$\frac{b-a}{2}[f(a)+f(b)]-\int_a^b f(x)dx=\frac{1}{12}(b-a)^3f''(\xi)$$

We are not going to prove this version of error, but a more general version involving absolute value of error.

- Trapezoidal rule is exact for  $f \in \Pi_1$
- For quadrature rules, the standard measure of accuracy is the degree of polynomials a formula is exact for.
- Most books has minus sign in the error. But I am sticking to the convention (error) = (estimate) (true): so negative error indicates underestimation (e.g., concave) and the positive overestimation (e.g., convex).

# Trapezoidal rule and its composite version

Composite trapezoidal rule  $\longrightarrow$  refine the interval into n subintervals and apply trapezoidal rule to each subinterval



$$\sum_{i=1}^{n} \int_{x_{i-1}}^{x_{i}} f(x) dx \approx \frac{1}{2} \sum_{i=1}^{n} (x_{i} - x_{i-1}) [f(x_{i-1}) + f(x_{i})]$$

If the refinement is uniform with h = (b - a)/n,

$$\int_a^b f(x)dx \approx \frac{h}{2} \left[ f(a) + 2 \sum_{i=1}^{n-1} f(a+ih) + f(b) \right]$$

**Exercise**) Show that the error term for the uniform refined case is  $\frac{1}{12}(b-a)h^2f''(\xi)$ . (Hint: You may want to use the discrete version of mean value theorem for integral)

# SIMPSON'S RULE

**Exercise**: Derive a Newton-Cotes formula for functions defined on [a,b] that is exact for  $p \in \Pi_2$ .

#### Guidelines

- Does this makes sense?
- Definition of Newton-Cotes?
- How many nodes to be chosen?
- What tool to use?

# SIMPSON'S RULE

**Exercise**: write a program that implement the composite Simpson's rule and test it.

- (minimum) input: a, b (a < b) (interval; two numbers), f (integrand; function handle), n (number of subintervals; an even integer)
- output: I (integral approximated by Simpson's rule)
- test: output  $\int_0^{\pi/2} \sin x dx$ ,  $\int_{-1}^1 f(x) dx$ , and  $\int_0^2 e^{-x^2/2} dx$ , where  $f(x) = \begin{cases} x^2/2 & (x \ge 0) \\ -x^2/2 & (x < 0) \end{cases}$ .
- Confirm the convergence rate of the composite Simpson's rule when possible.
  - ► For each function, repeat the test while doubling n (number of subintervals) (e.g.  $n = 16, 32, 64, 128, \cdots$ )
  - ▶ Report a table of n, approximate integrals, error,  $\log_2 e_k \log_2 e_{k+1}$ .
  - ► You need true value to compute the errors. You may want to use Wolfram alpha for such cases.

■ This will be one of the options for computational HW2.

#### Before we begin

- $\prod_{i \neq j} \frac{x x_j}{x_i x_i}$  and  $\prod_{i \neq j} \frac{x x_j}{x_i x_i}$  are totally different.
- From iclicker (last Tue): "happy upcoming thanksgiving," "HELLO!!" "thanks for today's great class!"
- Some answers to "Any patterns in error estimates of derivatives and integrals?"
  - ► Both depend on h?
  - well, they both look complicate but follow the logic
  - ▶ i think they all approach to zero when h is approaching to zero
  - $\,\blacktriangleright\,$  Both errors are dependent on a derivative of f which they are
  - accurate up to the number of nodes the is the degree of h?
  - ► I'm having a really hard time understanding this chapter. I can't find any pattern. But I am here and trying to learn. (not sure: I don't recognize a pattern; and more)
  - This error may looks like double integration of that of derivatives.
  - ► the LHS of equation is understandable with the difference, but the Max with h(n+2) has me lost

# CHANGE OF INTERVAL

**Goal**: A quadrature formula for on [a, b] for  $\int_a^b f(x) dx$ .

What we have:  $\int_c^d f(t) dt \approx \sum_{i=0}^n A_i \ell_i$  on [c, d].

**Strategy**: Change of variable:  $x = \alpha t + \beta$  so that  $c \mapsto a$  and

 $d\mapsto b$ . Call such a map  $x=\lambda(t)$ .

#### Quadrature formula for different intervals

$$\int_{a}^{b} f(x) dx \approx \frac{b-a}{d-c} \sum_{i=0}^{n} A_{i} f\left(\lambda\left(t_{i}\right)\right) = \frac{b-a}{d-c} \sum_{i=0}^{n} A_{i} f\left(x_{i}\right),$$

where  $x_i = \lambda(t_i)$  are the new nodes on [a, b] and  $t_i$  are the old nodes on [c, d].

#### Proof.

Board work.

In words.

- You can use a quadrature formula on an interval, say [0, 1], for any other intervals.
- You change only the function values  $f(x_i) = f(\lambda(t_i))$  while use the same weight  $A_i$ .
- You also need to adjust the total sum according to stretch or shrink (i.e., Jacobian (b-a)/(d-c)).

# A NAIVE BUT GENERAL ERROR ESTIMATE

**Intuition**: If functions are similar, so are their integrals.

# Quadrature error – single interval

For h > 0, Newton-Cotes formula of degree n on [a, a + h] satisfies

$$\left| \int_a^{a+h} f(x) dx - \sum_{i=0}^n A_i f(x_i) \right| \leq \frac{\hat{M}}{(n+1)!} h^{n+2},$$

where  $\hat{M} := \max_{a \le x \le a+h} |f^{(n+1)}(x)|$  and  $x_i = a + ih/n$  (i-th node;  $i = 0, 1, 2, \dots, n$ ).

Board work.

- Recall:  $f(x) - p(x) = \frac{1}{(n+1)!}$  $f^{(n+1)}(\xi_x) \prod_{i=0}^{n} (x - x_i)$
- Subjective question: Compare this result on the error with that of derivatives. Do you see any patterns or something notable?

# A NAIVE BUT GENERAL ERROR ESTIMATE

# Quadrature error – composite rules

Composite Newton-Cotes formula of degree n on [a, b] with k uniform subintervals satisfies

$$\left| \int_{a}^{b} f(x) dx - \sum_{j=1}^{k} \sum_{i=0}^{n} A_{i} f(x_{j,i}) \right| \leq \frac{M(b-a)}{(n+1)!} h^{n+1}$$

where  $M:=\max_{a\leq x\leq b}|f^{(n+1)}(x)|$  and  $x_{j,i}=a+(j-1)h+ih/n$  (i-th node in j-th subinterval;  $i=0,1,2,\cdots,n$  and  $j=1,2,\cdots,k$ ).

Board work.

 Draw a picture to better decipher indices.

# **Numerical Differentiation and Integration**

**Gaussian Quadrature** 

# WEIGHTED QUADRATURE

Sometimes, we need to integrate things w.r.t. some *density* or weight w(t)dt.

Fortunately, what we have developed works similarly.

# Weighted quadrature based on interpolation

Given  $x_0, x_1, \dots, x_n$ , the following quadrature formula  $\sum_{i=0}^n \tilde{A}_i f(x_i)$  is  $\Pi_n$ —exact, where

$$\tilde{A}_i = \int_a^b \ell_i(x) \mathbf{w}(x) dx,$$

Here, we are using only (n+1) degrees of freedom (for sensible  $\tilde{A}_i$ ). **Question**: can we do better by exploiting more degrees of freedom? If so, what to tune?

- w(t) is positive in the interior of the interval.
- Recall the fundamental interpolating polynomials

$$\ell_i(x) = \prod_{\substack{j=0\\j\neq i}}^n \frac{x-x_j}{x_i-x_j}.$$

They satisfy  $\ell_i(x_j) = \delta_{ij}$  (Kronecker delta)

# Gaussian quadrature

#### Theorem

Let w be a positive weight function and  $q \in \Pi_{n+1}$  be w-orthogonal to  $\Pi_n$ , that is,  $\int_a^b q(x)p(x)w(x)dx = 0$  ( $\forall p \in \Pi_n$ ). If  $x_0, x_1, \cdots, x_n$  are chosen to be the zeros of q, then the quadrature  $\sum_{i=0}^n \tilde{A}_i f(x_i)$  is  $\Pi_{2n+1}$ -exact, where  $\tilde{A}_i = \int_a^b \ell_i(x)w(x)dx$ .

#### Proof.

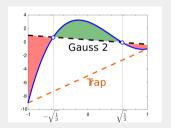
Board work.

■ From the degrees of freedom point of view, we've gained (n+1) higher accuracy by making (n+1) more sensible tuning.

# Gaussian quadrature

#### Example of Gaussian quadrature

setting	quadrature formula
$w(x) \equiv 1$	
[a,b] = [-1,1]	$x_0 = -1/\sqrt{3}$ , $x_1 = 1/\sqrt{3}$
n = 1	$ ilde{\mathcal{A}}_0 =  ilde{\mathcal{A}}_1 = 1$



- In theory, we can find arbitrarily accurate
   Gaussian formula.
- Once [a, b], n, and w(x) are chosen, everything is determined:  $\Pi_n \to$  orthogonal polynomial  $\to$  its roots  $\to$  nodes  $x_i$   $\to$  Lagrange basis  $\to$  weights  $\tilde{A}_i \to$  formula  $\sum_{i=0}^n \tilde{A}_i f(x_i)$
- Exercise: confirm that this is accurate up to degree 3.

# Properties of Orthogonal Polynomials

**Worries**: What if the zeros are repeated? What if we have to plug in numbers outside [a,b]?

#### Theorem

Let  $w \in C[a, b]$  be a positive weight function and  $f \in C[a, b]$  be w-orthogonal to  $\Pi_n$ . Then, f changes sign at least (n + 1) times on (a, b).

#### Proof.

Board work.

#### L

# Corollary

The nodes of Gaussian quadrature are all distinct and inside the interval of integral.

# Error

#### Lemma

The coefficients of the Gaussian quadrature are positive and sum up to  $\int_a^b w(x) dx$ .

#### Proof.

Board work.

# Theorem (Convergence of Gaussian quadrature)

If  $f \in C[a,b]$ , then the Gaussian quadrature converges to the true integral as  $n \to \infty$  (n+1) is the number of nodes. That is,  $\sum_{i=0}^n \tilde{A}_{n,i} f(x_{n,i}) \to \int_a^b f(x) w(x) \mathrm{d}x$ .

# Proof.

Board work.

Convergence theorem is not of practical interest. But this shows why mathematicians care 'dense' subset and is a great exercise.

# Error.

# Theorem (Error of the Gaussian quadrature)

If  $f \in C^{2n+2}[a, b]$ , then there is  $\xi \in (a, b)$  such that

$$\int_{a}^{b} f(x)w(x)dx - \sum_{i=0}^{n} \tilde{A}_{i}f(x_{i})$$

$$= \frac{1}{(2n+2)!} f^{(2n+2)}(\xi) \int_{a}^{b} \prod_{i=0}^{n} (x - x_{i})^{2} dx$$

#### Proof.

We will not prove this.



- We omit the proof. But it is not very hard if we studied Hermite interpolation.
- Absolute error will be of (2n+2)—th order.