# Intermediate Programming Day 15

## Outline

- Exercise 14
- Numerical representation
- Casting
- Random Numbers
- Review questions

Convert char \* message to int number

```
encrypt.c
int str_to_int( char msg[] , int len )
     int num = 0;
    if(len>32)
          fprintf( stderr , "[WARNING] Insufficient bits\n" );
          len = 32;
     for(int i=0; i<len; i++) if( msg[len-i-1]=='1') num += pow(2, i);
     return num;
```

Convert char \* message to int number

• Note that  $2^i = 1 << i$ 

```
encrypt.c
int str_to_int( char msg[] , int len )
     int num = 0;
     if(len>32)
          fprintf( stderr , "[WARNING] Insufficient bits\n" );
          len = 32;
     for( int i=0 ; i<len ; i++ ) if( msg[len-i-1]=='1' ) num += 1<<i;
     return num;
```

# Convert char \* message to int number

- Note that 2<sup>i</sup> = 1<<i
- Note that if i and j are variables with no common bits turned on (i&j==0) then i+j = i|j

```
encrypt.c
int str_to_int( char msg[] , int len )
     int num = 0:
     if(len>32)
          fprintf( stderr , "[WARNING] Insufficient bits\n" );
          len = 32;
     for( int i=0 ; i<len ; i++ ) if( msg[len-i-1]=='1' ) num |= 1<<i;
     return num;
```

Convert int number to char \* message

```
encrypt.c
void int_to_str( int num_encrypted , char msg_encrypted[] , int len )
     for( int i=0 ; i<len ; i++ )
          if( num_encrypted&1 ) msg_encrypted[len-i-1] = '1';
                                 msg_encrypted[len-i-1] = '0';
          else
         num_encrypted >>= 1;
     if( num_encrypted )
          fprintf( stderr , "[WARNING] Insufficient bits\n" );
```

Compute the encrypted message by repeatedly left-shifting the message by 1 and XORing.

```
encrypt.c
int main( void )
     int num_msg = 0;
     char msg[33] = {' \ 0'};
     int n = -1;
     int num_encrypted = 0;
     for( int i=0; i<n; i++) num_encrypted ^= num_msg<<i;
     return 0;
```

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#### Arithmetic

- The integers are a set (of numbers)
- There is an addition operator, +, that takes a pair of integers and returns an integer
  - There is a zero element, 0, with the property that adding zero to any integer gives back that integer:

$$a + 0 = a$$

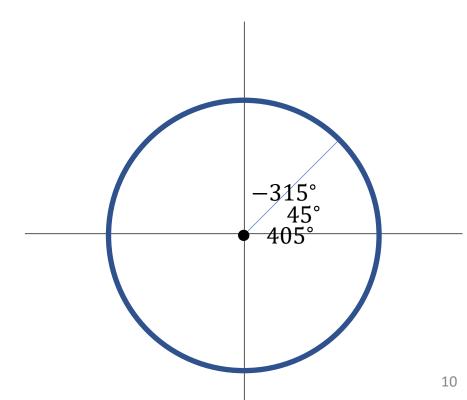
• Every integer a has an inverse -a such that the sum of the two is zero:

$$a + (-a) = a - a = 0$$

• Given a positive integer, M, we say that two integers  $\alpha$  and b are equivalent modulo M, if there is exists some integer k such that:

$$a \equiv b + k \cdot M$$

- Degrees in a circle (mod 360°)
- Hours on a clock (mod 12)



• Given a positive integer, M, we say that two integers a and b are equivalent modulo a, if there is exists some integer a such that:  $a \equiv b + k \cdot M$ 

- We can represent integers mod M using values in the range [0, M)
  - While an integer is bigger than or equal to M, repeatedly subtract M
  - While an integer is less than zero, repeatedly add M

• Given a positive integer, M, we say that two integers a and b are equivalent modulo M, if there is exists some integer k such that:  $a \equiv b + k \cdot M$ 

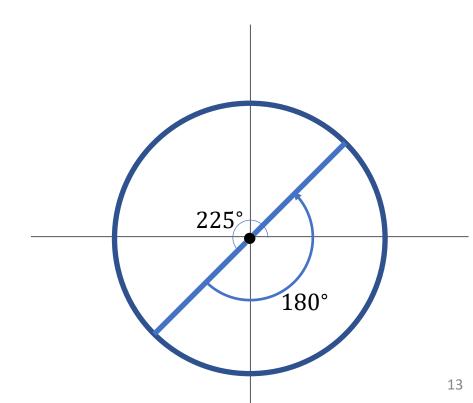
- We can represent integers mod M using values in the range [0, M)
- Or, we can represent integers mod M using the range  $\left[-\frac{M}{2}, \frac{M}{2}\right]$ 
  - While an integer is bigger than or equal to  $\frac{M}{2}$ , repeatedly subtract M
  - While an integer is less than  $-\frac{M}{2}$ , repeatedly add M

• Given a positive integer, M, we say that two integers  $\alpha$  and b are equivalent modulo M, if there is exists some integer k such that:

$$a \equiv b + k \cdot M$$

• We can add numbers modulo *M*:

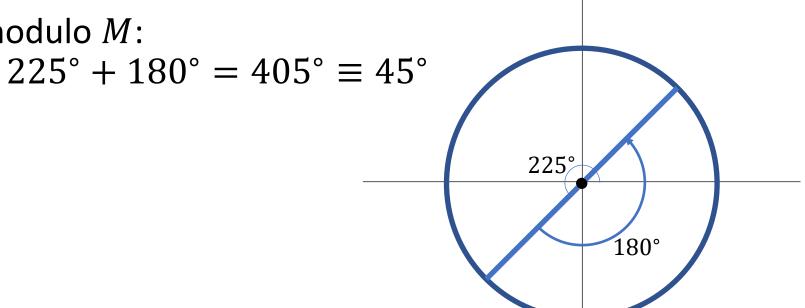
$$225^{\circ} + 180^{\circ} = 405^{\circ}$$



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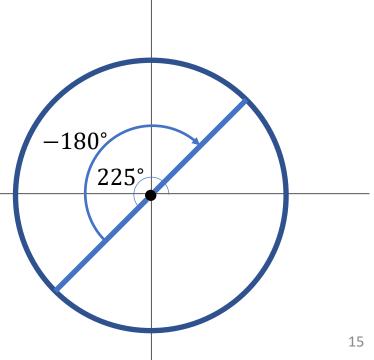


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$$a \equiv b + k \cdot M$$

• We can add numbers modulo *M*:

$$225^{\circ} - 180^{\circ} \equiv 225^{\circ} + 180^{\circ} = 405^{\circ} \equiv 45^{\circ}$$

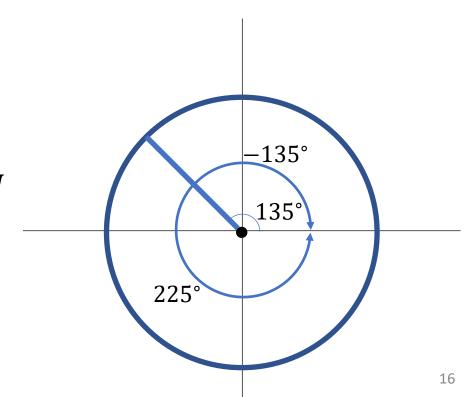


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We can add numbers modulo M

• For any integer a, the negative of a modulo M can be represented by M-a



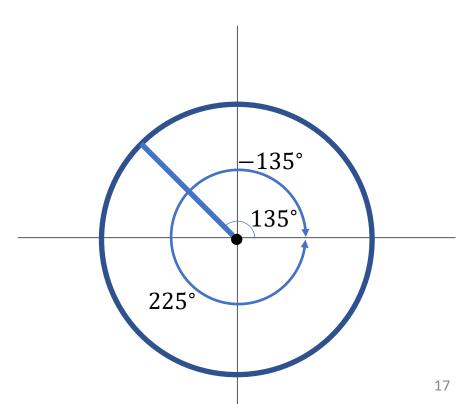
• Given a positive integer, M, we say that two integers  $\alpha$  and b are equivalent modulo M, if there is exists some integer k such that:

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• We can add numbers modulo M

• For any integer a, the negative of a modulo M can be represented by M-a:

$$a + (M - a) = (a - a) + M$$



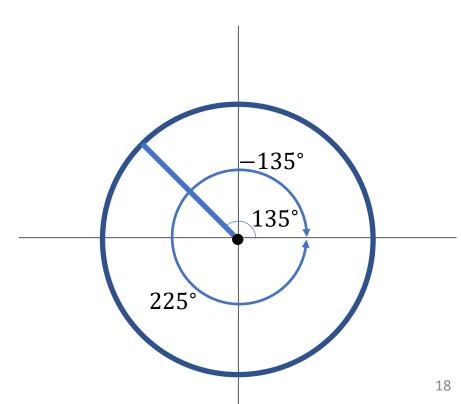
• Given a positive integer, M, we say that two integers  $\alpha$  and b are equivalent modulo M, if there is exists some integer k such that:

$$a \equiv b + k \cdot M$$

• We can add numbers modulo M

• For any integer a, the negative of a modulo M can be represented by M-a:

$$a + (M - a) = (a - a) + M = M$$



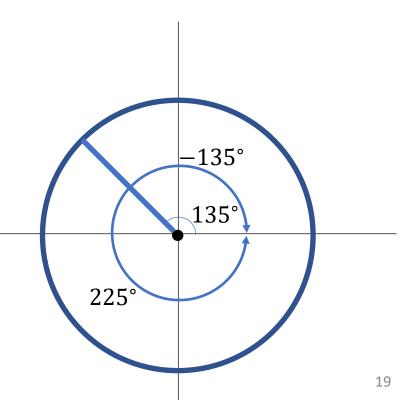
• Given a positive integer, M, we say that two integers  $\alpha$  and b are equivalent modulo M, if there is exists some integer k such that:

$$a \equiv b + k \cdot M$$

We can add numbers modulo M

• For any integer a, the negative of a modulo M can be represented by M-a:

$$a + (M - a) = (a - a) + M = M \equiv 0$$



• When we write out an integer in decimal notation, we are representing it as a sum of "one"s, "ten"s, "hundred"s, etc.

$$365 = 3 \times 100 + 6 \times 10 + 5 \times 1$$
  
=  $3 \times 10^{2} + 6 \times 10^{1} + 5 \times 10^{0}$ 

• This is unique because each digit is in the range 0 to 9, written [0,10)

- We add two numbers by adding the digits from smallest to largest
  - If the sum of digits falls outside the range [0,10) we carry

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Q: If we use three digits, how many numbers can we represent?

A:  $1000 = 10^3$  (including zero)

#### Note:

 The sum of two numbers represented using three digits may require four digits to store:

Q: If we use three digits, how many numbers can we represent?

A:  $1000 = 10^3$  (including zero)

#### Note:

 The sum of two numbers represented using three digits may require four digits to store:

- If we only use three digits, we lose the leading digit to overflow
- This is the same as the number mod  $10^3$

We can also write out numbers in base two

$$(s_3s_2s_1s_0)_2 = s_3 \times 8 + s_2 \times 4 + s_1 \times 2 + s_0 \times 1$$
  
=  $s_3 \times 2^3 + s_2 \times 2^2 + s_1 \times 2^1 + s_0 \times 2^0$ 

where  $s_0, s_1, s_2, s_3$  are either 0 or 1.

$$\begin{array}{ccccc}
 & (1 & 1 & 0)_2 \\
 + & (0 & 1 & 1)_2 \\
\hline
 & & & )_2
\end{array}$$

$$\begin{array}{cccc} & (1 & 1 & 0)_2 \\ + & (0 & 1 & 1)_2 \\ \hline ( & & 1)_2 \end{array}$$

$$\begin{array}{ccccc}
 & 1 & & \\
 & (1 & 1 & 0)_2 \\
 + & (0 & 1 & 1)_2 \\
\hline
 & & 0 & 1)_2
\end{array}$$

Q: Using three digits in base two, how many numbers can we represent?

A: 8 (including zero)

#### Note:

• As before, the sum of two numbers represented using three digits may require four digits to store: 1 1

$$\begin{array}{cccc} & (1 & 1 & 0)_2 \\ + & (0 & 1 & 1)_2 \\ \hline (1 & 0 & 0 & 1)_2 \end{array}$$

Q: Using three digits in base two, how many numbers can we represent?

A: 8 (including zero)

#### Note:

• As before, the sum of two numbers represented using three digits may require four digits to store:

$$\begin{array}{c|cccc} & (1 & 1 & 0)_2 \\ + & (0 & 1 & 1)_2 \\ \hline & (0 & 0 & 1)_2 \end{array}$$

- If we only use three digits, we lose the leading digit to overflow
- This is the same as the number mod 8.

Given a number in base 10:

$$16,384 = 1 \times 10^4 + 6 \times 10^3 + 3 \times 10^2 + 8 \times 10^1 + 4 \times 10^0$$

we can get an expression of the number in base 100 by grouping digits:

$$16,384 = 1 \times 100^2 + 63 \times 100^1 + 84 \times 100^0$$

Similarly, we can get an expression of the number in base 1000, etc.

### Bases (two)

• Similarly, given a number in base two:

$$(s_3s_2s_1s_0)_2 = s_3 \times 8 + s_2 \times 4 + s_1 \times 2 + s_0 \times 1$$

we can get an expression of the number in base 4 by grouping digits:

$$(s_3s_2s_1s_0)_2 = (s_3 \times 2 + s_2) \times 4 + (s_1 \times 2 + s_0) \times 1$$

Similarly, we can get an expression of the number in base 8, or base 16, or ...

What is the value in base 10?

• 
$$(1101)_2 =$$

What is the value in base 10?

• 
$$(1101)_2 = 1 \times 8 + 1 \times 4 + 0 \times 2 + 1 \times 1$$
  
=  $8 + 4 + 1$   
=  $13$ 

What is the value in base 2?

• 27 =

What is the value in base 2?

• 
$$27 = 16 + 8 + 4 + 2 + 1$$
  
=  $1 \times 16 + 1 \times 8 + 0 \times 4 + 1 \times 2 + 1 \times 1$   
=  $(11011)_2$ 

What is the value in base 2?

• 
$$27 = 16 + 8 + 4 + 2 + 1$$
  
=  $1 \times 16 + 1 \times 8 + 0 \times 4 + 1 \times 2 + 1 \times 1$   
=  $(11011)_2$ 

What is the value in base 4?

What is the value in base 2?

• 
$$27 = 16 + 8 + 4 + 2 + 1$$
  
=  $1 \times 16 + 1 \times 8 + 0 \times 4 + 1 \times 2 + 1 \times 1$   
=  $(11011)_2$ 

What is the value in base 4?

- Decimal (base 10)
  - We have ten fingers
- Sexagesimal (base 60):
  - Minutes / seconds
  - Easy to tell if a number is divisible by 2, 3, 4, 5, 6, 10, 12, 15, or 30
  - Dates back to the Babylonians

- Binary (base 2)
  - Numbers in a computer

- Hexadecimal a.k.a. hex (base 16)
  - Numbers in a computer  $(16 = 2^4)$ 
    - We can easily convert binary to hex by grouping sets of four digits
    - We get a more compact representation, replacing 4 digits with 1

- Binary (base 2)
  - Numbers in a computer

Hexadecimal a.k.a. hex (base 16)

Q: How should we separate the digits?  $(115)_{16}$ 

• 
$$(115)_{16} = 1 \times 16^2 + 1 \times 16^1 + 5 \times 16^0$$

• 
$$(115)_{16} = 1 \times 16^1 + 15 \times 16^0$$

• 
$$(115)_{16} = 11 \times 16^1 + 5 \times 16^0$$

- Binary (base 2)
  - Numbers in a computer

Hexadecimal a.k.a. hex (base 16)

Q: How should we separate the digits?  $(115)_{16}$ 

A: Use numbers and letters:

- {0,1,2,3,4,5,6,7,8,9} to represent numbers in the range [0,10)
- $\{a, b, c, d, e, f\}$  to represent values in the range [10,16):
  - $(115)_{16} = 1 \times 16^2 + 1 \times 16^1 + 5 \times 16^0$
  - $(1f)_{16} = 1 \times 16^1 + 15 \times 16^0$
  - $(b5)_{16} = 11 \times 16^1 + 5 \times 16^0$

- On most machines, [unsigned] ints are represented using 4 bytes\*
  - Each byte is composed of 8 bits
  - ⇒ An [unsigned] int is represented by 32 bits
  - Each bit can be either "on" or "off"
  - ⇒ An [unsigned] int is represented in binary using 32 digits with values 0 or 1
  - $\Rightarrow$  An [unsigned] int can have one of  $2^{32}$  values

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  - $\Rightarrow$  An [unsigned] int can have one of  $2^{32}$  values

On the machine,  $\alpha$  is assigned the value:

```
a \leftarrow (00\ 00\ 00\ 1e)_{16}
```

```
#include <stdio.h>
int main( void )
{
    int a = 30;
    printf( "%d\n" , a );
    return 0;
}

>> ./a.out
30
>>
```

- On most machines, [unsigned] ints are represented using 4 bytes
  - Each byte is composed of 8 bits
  - ⇒ An [unsigned] int is represented by 32 bits
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On the machine,  $\alpha$  is assigned the value:

```
a \leftarrow (00\ 00\ 00\ 1e)_{16}
```

• You can assign using base 16 by preceding the number with 0x to indicate hex

```
#include <stdio.h>
int main( void )
{
    int a = Ox1e;
    printf( "%d\n" , a );
    return 0;
}
    >> ./a.out
30
```

- On most machines, [unsigned] ints are represented using 4 bytes
  - Each byte is composed of 8 bits
  - ⇒ An [unsigned] int is represented by 32 bits
  - Each bit can be either "on" or "off"
  - ⇒ An [unsigned] int is represented in binary using 32 digits with values 0 or 1
  - $\Rightarrow$  An [unsigned] int can have one of  $2^{32}$  values

On the machine,  $\alpha$  is assigned the value:

```
a \leftarrow (00\ 00\ 00\ 1e)_{16}
```

- You can assign using base 16 by preceding the number with 0x to indicate hex
- You can print the base 16 representation by using %x for formatting

```
#include <stdio.h>
int main( void )
{
    int a = 30;
    printf( "%x\n" , a );
    return 0;
}
    >> ./a.out
1e
```

- On most machines, [unsigned] chars are represented using 1 byte ⇒ A [unsigned] char can have one of 2<sup>8</sup> values
- On most machines, [unsigned] long ints are represented using 8 bytes
  - A [unsigned] long int can have one of 2<sup>64</sup> values

- $\Rightarrow$  An [unsigned] char can have one of  $2^8 = 256$  values
- $\Rightarrow$  [unsigned] chars are integer values mod  $2^8$ 
  - unsigned char: We will use the range [0,256) to represent integers
  - char: We will use the range [-128,128) to represent integers

Q: What's the difference? Integers mod  $2^8$ , regardless of the representation!!!

- unsigned char: We will use the range [0,256) to represent integers
- char: We will use the range [-128,128) to represent integers

Q: What's the difference?

A: Is  $125 < 129 \mod 256$ ? Since  $129 \equiv -127 \mod 256$ , it depends on the range we use

```
#include <stdio.h>
int main( void )
{
    unsigned char c1 = 125 , c2 = 129;
    printf( "%d\n" , c1<c2 );
    return 0;
}

>> ./a.out
1
>>
```

- unsigned char: We will use the range [0,256) to represent integers
- char: We will use the range [-128,128) to represent integers

Q: What's the difference?

A: Is  $125 < 129 \mod 256$ ? Since  $129 \equiv -127 \mod 256$ , it depends on the range we use

#### • Addition:

We add two numbers, a + b, by adding the digits from smallest to largest

- We carry as necessary
- And we cut off at 8 bits

```
\begin{array}{r}
11 & 11 \\
(11010011)_2 \\
+(01000110)_2 \\
\hline
(100011001)_2 \\
= (00011001)_2
\end{array}
```

#### • Addition:

We add two numbers, a + b. by adding the digits from smallest to largest

- We carry as necessary
- And we cut off at 8 bits

$$\begin{array}{r}
11 & 11 \\
(11010011)_2 \\
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\hline
(00011001)_2
\end{array}$$

Q: What about subtraction, a - b?

#### • Addition:

We add two numbers, a + b. by adding the digits from smallest to largest

- We carry as necessary
- And we cut off at 8 bits

$$\begin{array}{r}
11 & 11 \\
(11010011)_2 \\
+(01000110)_2 \\
\hline
(00011001)_2
\end{array}$$

Q: What about subtraction, a - b = a + (-b)? Equivalently, how do we define the negative of a number?

#### Negation

#### • Recall:

The negative of an integer is the number we would have to add to get back zero.

- <u>Defining negative one</u>:
  - Mod 256, we have  $-1 \equiv 255 = (111111111)_2$

#### Negation

#### Recall:

The negative of an integer is the number we would have to add to get back zero.

- <u>Defining negatives in general</u>:
  - 1. Given a binary value in 8 bits:

$$(10011101)_2$$

2. We can flip the bits:

$$(01100010)_2$$

- 3. Adding the two values we get  $255 \equiv -1$ :  $(111111111)_2$
- 4. Adding one to that we get 0

#### Negation

#### • Recall:

The negative of an integer is the number we would have to add to get back zero.

#### • 2's complement:

To get the binary representation of the negative of a number

- 1. Flip the bits
- 2. Add 1

### Floating point value representation

$$\pm s \times 2^e$$

- On most machines, floats are represented using 4 bytes (32 bits)
  - These are (roughly) used to encode:
    - The sign  $(\pm)$ : 1 bit
    - The signed (integer) exponent (e): 8 bits\*
    - The unsigned (integer) significand/mantissa/coefficient (s): 23 bits

### Floating point value representation

$$\pm s \times 2^e$$

- On most machines, doubles are represented using 8 bytes (64 bits)
  - These are (roughly) used to encode:
    - The sign  $(\pm)$ : 1 bit
    - The signed (integer) exponent (e): 11 bits
    - The unsigned (integer) significand/mantissa/coefficient (s): 52 bits

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When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

```
<type-1> lhs;
<type-2> rhs;
lhs = rhs;
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

• If both are integers and sizeof(LHS)>=sizeof(RHS)

⇒ the conversion happens without loss of information

```
#include <stdio.h>
int main( void )
{
    char c = 'a';
    int i = c;
    printf( "%d -> %d\n" , c , i );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

• If both are integers and sizeof(LHS)  $\prec$ sizeof(RHS)  $\Rightarrow$  an implicit "modulo" operation is performed (modulo  $2^b$  where b is the

number of bits in the LHS)

```
#include <stdio.h>
int main( void )
{
    int i = 511;
    char c = i;
    printf( "%d -> %d\n" , i , c );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

• If both are floats and sizeof(LHS)>=sizeof(RHS)

⇒ the conversion happens without loss of information

```
#include <stdio.h>
int main( void )
{
    float f = 1.5;
    double d = f;
    printf( "%.8f -> %.8f\n" , f , d );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

If both are floats and sizeof(LHS) \sizeof(RHS)
 ⇒ rounding is performed

```
#include <stdio.h>
int main( void )
{
    double d = 1.7;
    float f = d;
    printf( "%.8f -> %.8f\n" , d , f );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

• If the LHS is an integer and the RHS is a floating point value ⇒ the fractional part is discarded

```
#include <stdio.h>
int main( void )
{
    double d = -3.6;
    int i = d;
    printf( "%.8f -> %d\n" , d , i );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

• If the LHS is a floating point value and the RHS is an integer ⇒ the closest floating point representation is used

```
#include <stdio.h>
int main( void )
{
    int i = 123456789;
    float f = i;
    printf( "%d -> %.Of\n" , i , f );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS). The same rules apply when passing values to/from a function

```
#include <stdio.h>
char foo( unsigned char c ){ return c; }
int main( void )
{
    double d = 511.5;
    float f = foo( d );
    printf( "%g -> %g\n" , d , f );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

The same rules apply when passing values to/from a function

• double  $\rightarrow$  unsigned char:  $511.5 \rightarrow 511 \rightarrow 255$ 

```
#include <stdio.h>
char foo( unsigned char c ){ return c; }
int main( void )
{
    double d = 511.5;
    float f = foo( d );
    printf( "%g -> %g\n" , d , f );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

The same rules apply when passing values to/from a function

- double  $\rightarrow$  unsigned char:  $511.5 \rightarrow 511 \rightarrow 255$
- unsigned char  $\rightarrow$  char:  $255 \rightarrow -1$

```
#include <stdio.h>
char foo( unsigned char c ){ return c; }
int main( void )
{
    double d = 511.5;
    float f = foo( d );
    printf( "%g -> %g\n" , d , f );
    return 0;
}
```

When you assign a value to a variable, the right-hand-side (RHS) is implicitly converted (a.k.a. cast) to the type of the left-hand-side (LHS)

The same rules apply when passing values to/from a function

- double  $\rightarrow$  unsigned char:  $511.5 \rightarrow 511 \rightarrow 255$
- unsigned char  $\rightarrow$  char:  $255 \rightarrow -1$
- char  $\rightarrow$  float:  $-1 \rightarrow -1.f$

```
#include <stdio.h>
char foo( unsigned char c ){ return c; }
int main( void )
{
    double d = 511.5;
    float f = foo( d );
    printf( "%g -> %g\n" , d , f );
    return 0;
}
```

#### When casting, the types are ranked:

 Larger size integers/floats are "higher rank" char < int < long</li>

```
unsigned char < unsigned int < unsigned long float < double
```

- Unsigned integers are "higher rank" than signed integers\*
   char < unsigned char < int < unsigned int < long < unsigned long
   float < double</li>
- Floating point values are "higher rank" than integers char < unsigned char < int < unsigned int < long < unsigned long < float < double

When casting, the types are ranked:

char < unsigned char < int < unsigned int < long < unsigned long < float < double

When we cast from lower rank to higher rank, we are **promoting**.

When we cast from higher rank to lower rank, we are narrowing.

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted:

char < unsigned char < int < unsigned int < long < unsigned long < float < double

```
#include <stdio.h>
int main( void )
{
    int i = -1;
    unsigned int ui = 1;
    printf( "d\n" , i<ui );
    return 0;
}</pre>
```

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted: char < unsigned char < int < unsigned int < long < unsigned long < float < double

```
#include <stdio.h>
int main( void )
{
    int i = 2;
    double d = 2.5;
    i = i * d;
    printf( "%d\n" , i );
    return 0;
}
```

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted: char < unsigned char < int < unsigned int < long < unsigned long < float < double

```
#include <stdio.h>
int main( void )
{
    int i = 2;
    double d = 2.5;
    i *= d;
    printf( "%d\n" , i );
    return 0;
}
```

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted: char < unsigned char < int < unsigned int < long < unsigned long < float < double

```
#include <stdio.h>
int main( void )
{
    int one = 1;
    int four = 4;
    int i = one / four * four;
    printf( "%d\n" , i );
    return 0;
}
```

The arithmetic operators \*, /, and % have the same precedence and are evaluated left-to-right.

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted: char < unsigned char < int < unsigned int < long < unsigned long < float < double

```
#include <stdio.h>
int main( void )
{
    double one = 1;
    int four = 4;
    int i = one / four * four;
    printf( "%d\n" , i );
    return 0;
}
```

The arithmetic operators \*, /, and % have the same precedence and are evaluated left-to-right.

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted:

char < unsigned char < int < unsigned int < long < unsigned long < float < double

Since evaluation precedes assignment, we get truncated results even though the LHS doesn't require it.

```
#include <stdio.h>
int main( void )
{
    int one = 1 , four = 4;
    float f = one / four;
    printf( "%g\n" , f );
    return 0;
}
```

When performing a binary operation (arithmetic or comparison) with different types the "lower rank" operand is implicitly promoted:

char < unsigned char < int < unsigned int < long < unsigned long < float < double

The desired behavior can be forced with casting:

- Preceding the variable name with (<type-name>)
  converts the variable to type <type-name>
- Since casting takes precedence over arithmetic operations:
  - 1. We convert one to a float
  - 2. And then divide a float by an int
    - a. This implicitly promotes four to a float
    - o. And then performs float by float division

```
#include <stdio.h>
int main( void )
{
    int one = 1 , four = 4;
    float f = (float)one / four;
    printf( "%g\n" , f );
    return 0;
}
```

>> ./a.out

## Casting between types (pointers)

- Since pointers represent locations in memory (independent of type)
  - We can cast between pointer types (though not necessarily a good idea)
    - This needs to be done explicitly

```
#include <stdio.h>
int main( void )
{
    ...
    int i = 1;
    int *ip = &i;
    float *fp= (float*)ip;
    ...
}
```

# Casting between types (pointers)

- Since pointers represent locations in memory (independent of type)
  - We can cast between pointer types
    - This needs to be done explicitly
    - Unless one of them has type void\*

```
...
void *malloc( size_t );
...
```

```
#include <stdio.h>
int main( void )
{
    ...
    float *a = malloc( 10 * sizeof( float ) );
    ...
}
```

# Casting between types (pointers)

- Since pointers represent locations in memory (independent of type)
  - We can cast between pointer types
    - This needs to be done explicitly
    - Unless one of them has type void\*
  - We can also explicitly cast between pointers and integers
    - This needs to be done with care since a pointer can have different sizes on different machines:
      - 4 bytes on a 32-bit machine
      - 8 bytes on a 64-bit machine
    - The size\_† type is guaranteed to always have the size of a pointer

```
#include <stdio.h>
int main( void )
{
    int i = 100;
    int *ip = &i;
    size_t addr = (size_t)ip;
    printf( "Address is: %zu\n" , addr );
    return 0;
}
```

- 1. Nothing changes in the binary representation
  - pointers ↔ pointers
    - A memory address is a memory address
    - The compiler needs to know the type to transform element offsets into byte offsets
  - unsigned integers 
     ⇔ signed integers
    - Different representations of numbers modulo *M* still represent the same number
    - The compiler needs to know the type for comparisons

- 1. Nothing changes in the b #include <stdio.h>
  - pointers ↔ pointers
    - A memory address is a memory
    - The compiler needs to know
  - unsigned integers ↔ signed
    - Different representations of r
    - The compiler needs to know

```
void PrintBinary( const void *mem , size_t sz ){ ... }
int main(void)
      int iArray[] = { 1, 2, 3, 4 };
      int *iPtr = iArray;
      char *cPtr = (char *)iPtr;
      PrintBinary( iPtr , sizeof(iPtr) );
      PrintBinary( cPtr , sizeof(cPtr) );
      return 0;
```

- 1. Nothing changes in the b #include <stdio.h>
  - pointers ↔ pointers
    - A memory address is a memory
    - The compiler needs to know
  - unsigned integers ↔ signed
    - Different representations of r
    - The compiler needs to know

```
void PrintBinary( const void *mem , size_t sz ){ ... }
int main(void)
      unsigned int ui =(1 << 31)|1;
      int i = ui;
      printf( " %u = " , ui ); PrintBinary( &ui , sizeof(ui) );
      printf( "%d = " , i ); PrintBinary( &i , sizeof(i) );
      return 0;
       >> ./a.out
        2147483649 = 10000000 00000000 00000000 00000001
       -2147483647 = 10000000 00000000 00000000 00000001
```

- 1. Nothing changes in the binary representation
- 2. Binary representations are truncated/expanded
  - integers ↔ integers (of different sizes)

#### Three types of casting:

- Nothing changes in the bin #include <stdio.h>
- - integers ↔ integers (of differe int main(void)

```
2. Binary representations are void PrintBinary (const void mem, size_t sz) { ... }
                                            int i = 254;
                                            unsigned char c = i;
                                            printf( "%d = " , i ); PrintBinary( &i , sizeof(i) );
                                            printf( "%d = ", c); PrintBinary( &c, sizeof(c));
                                            return 0;
                                               >> ./a.out
```

= 00000000 00000000 00000000 11111110

254 = 111111110

>>

- 1. Nothing changes in the binary representation
- 2. Binary representations are truncated/expanded
- 3. Binary representations are completely different
  - integers ↔ floating point values
  - floating point values ↔ floating point values (of different sizes)

#### Three types of casting:

- 1. Nothing changes in the bir #include <stdio.h>
- 2. Binary representations are
- 3. Binary representations are
  - integers ↔ floating point valu
  - floating point values ↔ floatil

```
void PrintBinary( const void* mem , size_t sz ){ ... }
int main(void)
      int i = 1;
      float f = i;
      printf("%d = ",i); PrintBinary(&i, sizeof(i));
      printf( "%.1f = ", f); PrintBinary(&f, sizeof(f));
      return 0;
```

00000000 00000000 00000001

= 00111111 10000000 00000000 00000000

>>

101

### Outline

- Exercise 14
- Numerical representation
- Casting
- Random numbers
- Review questions

stdlib.h declares two functions for generating random numbers int rand(void);

- Returns a random integer value between 0 and RAND\_MAX
- **RAND\_MAX** is a constant (at least 32,767=2<sup>15</sup>-1)
- Each call to rand creates a new random number

```
#include <stdio.h>
#include <stdlib.h>

int main( void )
{
    printf( "%d<=%d\n" , rand() , RAND_MAX );
    printf( "%d<=%d\n" , rand() , RAND_MAX );
    return 0;
}</pre>
```

stdlib.h declares two functions for generating random numbers int rand(void);

- Returns a random integer value between 0 and RAND\_MAX
- **RAND\_MAX** is a constant (at least 32,767=2<sup>15</sup>-1)
- Each call to rand creates a new random number

```
#include <stdio.h>
#include <stdib.h>

int main( void )
{
    printf( "%d<=%d\n" , rand() , RAND_MAX );
    printf( "%d<=%d\n" , rand() , RAND_MAX );
    return 0;
}
</pre>
```

stdlib.h declares two functions for generating random numbers

void srand( unsigned int );

- Seeds the random number generator
- Calling rand after the random number has been seeded will consistently generate the same set of random numbers.
- Useful for debugging (for consistency)
- Useful for trying different values

```
#include <stdio.h>
#include <stdlib.h>
int main(void)
    srand(1);
    printf( "%d , %d\n" , rand() , rand() );
    srand(2);
    printf( "%d , %d\n" , rand() , rand() );
    srand(1);
    printf( "%d , %d\n" , rand() , rand() );
    return 0;
```

stdlib.h declares two functions for generating random numbers

void srand( unsigned int );

- Seeds the random number generator
- Calling rand after the random number has been seeded will consistently generate the same set of random numbers.
- Useful for debugging (for consistency)
- Useful for trying different values

```
#include <stdio.h>
#include <stdlib.h>
int main(void)
   srand( 1 );
    printf( "%d , %d\n" , rand() , rand() );
    srand(2);
    printf( "%d , %d\n" , rand() , rand() );
   srand(1);
                   >> ./a.out
    printf("%d,%
                   846930886 , 1804289383
    return 0;
                    846930886 , 1804289383
```

#### Warning:

Seeding the random number with a value of 1 has a special meaning. It re-seeds the random number generator to its original state.

On some machines the "original state" is seeded with 0, so that srand(0) and srand(1) do the same thing.

- Seeds the random number generator
- Calling rand after the random number has been seeded will consistently generate the same set of random numbers.
- Useful for debugging (for consistency)
- Useful for trying different values

```
int main(void)
   srand(1);
    printf( "%d , %d\n" , rand() , rand() );
   srand(2);
   printf( "%d , %d\n" , rand() , rand() );
   srand(1);
                   >> ./a.out
   printf( "%d , % 846930886 , 1804289383
                    1738766719 , 1505335290
    return 0;
                    846930886 , 1804289383
```

We can use rand to generate random numbers in an integer range

```
#include <stdio.h>
#include <stdlib.h>
int myRand(int low, int high)
    return low + rand() % ( high - low );
int main(void)
    printf( "%d , %d\n" , myRand(2,6) , myRand(2,6) );
    printf( "%d , %d\n" , myRand(16,26) , myRand(16,26) );
    return 0;
                                                           >> ./a.out
```

Note: We car ge This creates random numbers in the range [low,high). #include <stdio.h> #include <stdlib.h> int myRand(int low, int high) return low + rand() % ( high - low ); int main(void) printf( "%d , %d\n" , myRand(2,6) , myRand(2,6) ); printf( "%d , %d\n" , myRand(16,26) , myRand(16,26)); return 0; >> ./a.out

We can use rand to generate random numbers in a floating point range

```
#include <stdio.h>
#include <stdlib.h>
float myRand(float low, float high)
   return low + (float)rand() / RAND_MAX * ( high - low );
int main(void)
   printf( "%f , %f\n" , myRand(2,6) , myRand(2,6) );
   printf( "%f , %f\n" , myRand(16,26) , myRand(16,26));
   return 0;
                                                         >> ./a.out
                                                         23.984400 , 23.830992
```

### Outline

- Exercise 14
- Numerical representation
- Casting
- Review questions

1. What is *two's complement* representation?

It is a signed integer representation. The negative of a number is obtained by flipping the bits and adding one.

2. How does representation of integers and floating-point values differ in C?

The bits of an integer correspond to its representation in base two.

The bits of a floating-point value are split into three parts – the sign, the mantissa, and the exponent.

3. What is *type narrowing*?

Converting a "higher rank" data type into a "lower rank" one char < unsigned char < int < unsigned int < long < unsigned long < float < double

4. What is *type promotion*?

Converting a "lower rank" data type into a "higher rank" one char < unsigned char < int < unsigned int < long < unsigned long < float < double

5. What is *type casting*?

Explicitly or implicitly converting a value from one type to another

6. What is the output of:

```
int n = 32065;
float x = 24.79;
printf( "int n = %d but (char)n = %c\n" , n , (char)n );
printf( "float x = %f but (long)x = %ld\n" , x , (long)x );
```

In binary, we have:

 $32065 = (00000000 00000000 01111101 01000001)_2$ 

Casting to a char we get:

$$(01000001)_2 = 65 \rightarrow 'A'$$

```
int n = 32065 but (char)n = A
float x = 24.790001 but (long)x = 24
```

### Exercise 14

• Website -> Course Materials -> Exercise 14