

Lecture 18: Low-level code generation

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November 2, 2022

601.428/628 Compilers and Interpreters



High-level code generation

These slides present a few thoughts/recommendations about low-level (x86-64) code generation

Print comments as debugging output

Printing C-style `/* comments */` is a really useful way to emit debugging information in a way that won't interfere with the generated code being assembled

For example:

- ▶ Storage allocation decisions
- ▶ Computations involving memory layout

This technique is also useful for high-level code generation

Accessing memory in the stack frame

- ▶ Provided translations of `enter` and `leave` will (hopefully!) create ABI-compliant stack frames
- ▶ Assume that N bytes of memory are reserved in the stack frame
- ▶ `%rbp - N` points to the “bottom” of the local memory area
- ▶ Assume that i is the offset of a local variable: its displacement from `%rbp` should be $N - i$
- ▶ For example, if $N = 32$ and $i = 0$, use `-32(%rbp)` to access the memory location

Allocating storage for virtual registers

- ▶ Each function will use a certain number of virtual registers as
 - ▶ Storage for temporary (computed) values
 - ▶ Storage for (some) scalar local variables
- ▶ Note that `vr0` really means `%rax` and `vr1`, `vr2`, etc. are argument registers (`%rdi`, `%rsi`, etc.)
- ▶ For Assignment 4: allocate each `vreg` (other than `vr0` through `vr9`) in memory in the stack frame
 - ▶ This is in addition to the memory needed for local variables whose storage is in memory
- ▶ Assignment 5: you can do local register allocation to promote some virtual registers to CPU registers

Machine register sizes

- ▶ Each machine register has “subregisters” of various sizes
- ▶ These are specified as `Operand::Kind` values
- ▶ E.g., `Operand(Operand::MREG32, MREG_RAX)` represents the `%eax` register (i.e., the 32-bit sub-register of `%rax`)
- ▶ The `select_mreg_kind` helper function assists in selecting the correct machine register size

Instruction variants

- ▶ For instructions which move, compute, or compare values, there are different variants for different operand sizes
- ▶ The `select_ll_opcode` assists in selecting the correct low-level opcode

Temporary machine registers

- ▶ You can use `%r10` and `%r11` (and sub-registers of `%r10` and `%r11`) to store temporary values
- ▶ Use for dealing with situations such as
 - ▶ An x86-64 instruction can have at most one memory operand
 - ▶ Some instructions doesn't allow an immediate operand and a memory operand

What if a virtual register is used as a pointer?

Your high-level code will probably have operands like `(vr10)`, where a virtual register (in this case `vr10`) is being used as a pointer to access a data value in memory

Since virtual register values will be stored in memory, just referring to the contents of the virtual register requires a memory reference (e.g., `-24(%rbp)`). How to dereference a pointer if the pointer is in memory?

Solution: copy the pointer to a machine register, e.g.:

```
movq -24(%rbp), %r11
```

...code can now use `(%r11)` to dereference the pointer...

Conditions/decisions

- ▶ The comparison instructions provided high-level opcodes yield a boolean data value
- ▶ The `cjmp_t` and `cjmp_f` instructions consume this computed boolean data value
- ▶ How to generate code?
 - ▶ Use `setxx` instruction to use condition codes to set a boolean value in an 8 bit register

Example of evaluating a condition, control flow

...code for lhs subexpression...
...code for rhs subexpression...

```
cmpl rhsval, lhsval  
setl %r10b  
cmpb $0, %r10b  
je .Lsome_label
```


