

CMPUT 333

Security in a Networked World

Lab Assignment 1

Part 1

Goal: To decrypt “ciphertext1”

Available Data

- **Plaintext: Printable text ASCII file**
- **Key: Printable and non-printable (control) ASCII characters**

Lookup Table Map[x][k]

- **Map [16] [16]: Will ultimately give, the ciphertext's 4 higher or lower bits (ch,cl)**
 - Rows: x, 4 bit quantity
 - Columns: k, 4 bit quantity

Part 1

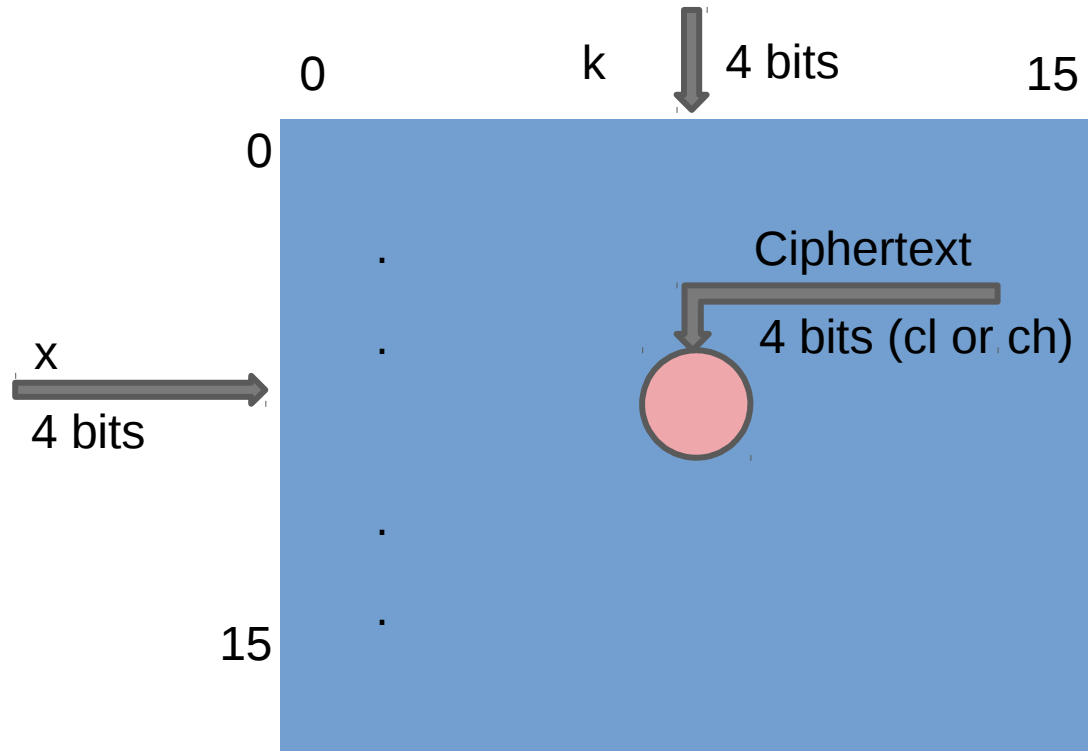


Figure 1: Encryption table

Part 1

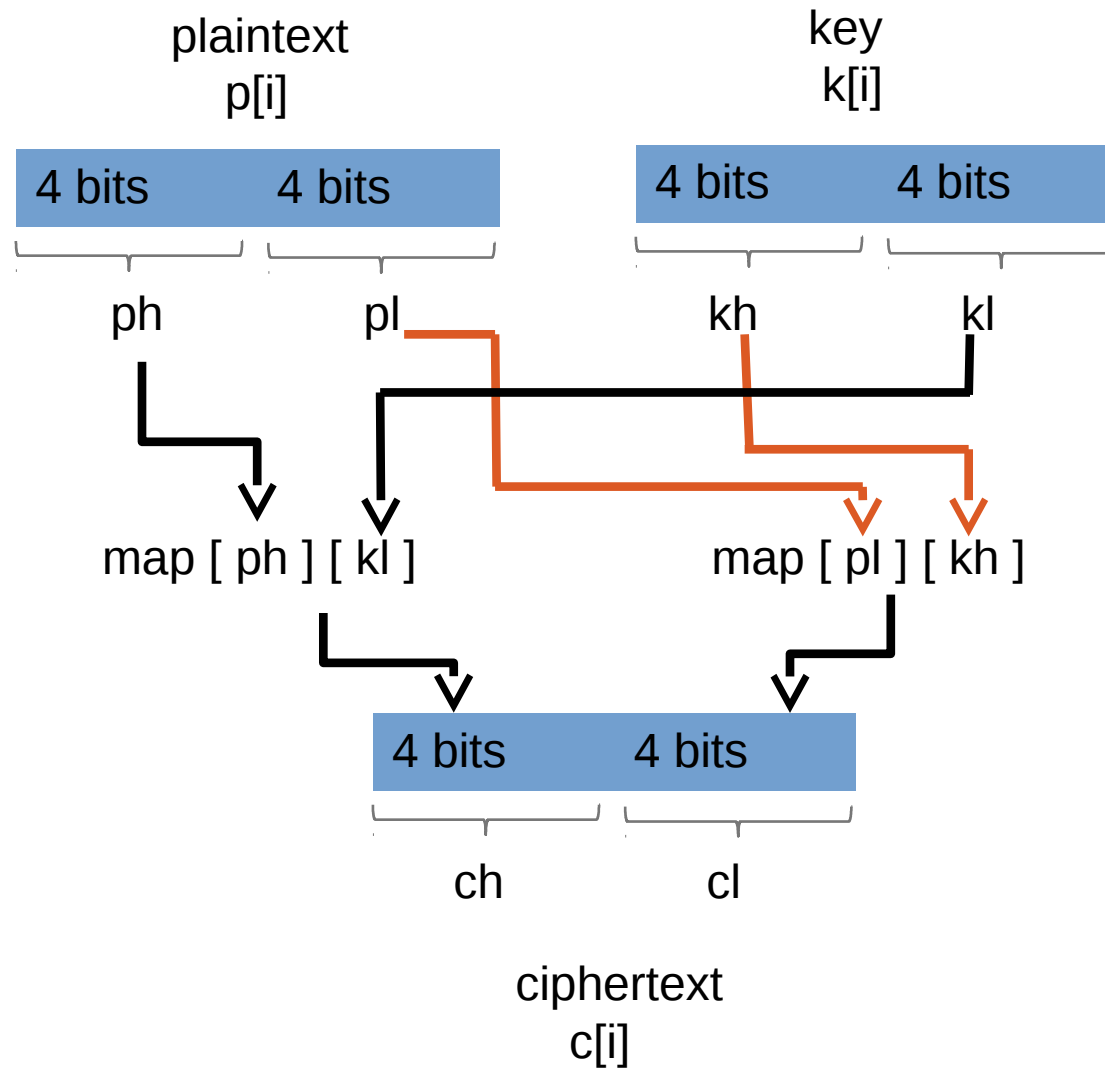


Figure 2: Encryption diagram

Part 1

- **How is encryption performed?**

Example: Suppose that the plaintext to be encrypted is “hello!” and the keyword is “key”.

Step 1: Extract p_h , p_l from the byte p and k_h , k_l from the byte k .

Step 2: Use these formulas to find the values of ch and cl in the map.

$ch \leftarrow \text{map} [p_h] [k_l]$

$cl \leftarrow \text{map} [p_l] [k_h]$

Step 3: Combine ch and cl into the byte c .

Part 1

Dec	Hx	Oct	Char	Dec	Hx	Oct	Html	Chr	Dec	Hx	Oct	Html	Chr	Dec	Hx	Oct	Html	Chr
0	0	000	NUL (null)	32	20	040	 	Space	64	40	100	@	@	96	60	140	`	`
1	1	001	SOH (start of heading)	33	21	041	!	!	65	41	101	A	A	97	61	141	a	a
2	2	002	STX (start of text)	34	22	042	"	"	66	42	102	B	B	98	62	142	b	b
3	3	003	ETX (end of text)	35	23	043	#	#	67	43	103	C	C	99	63	143	c	c
4	4	004	EOT (end of transmission)	36	24	044	$	\$	68	44	104	D	D	100	64	144	d	d
5	5	005	ENQ (enquiry)	37	25	045	%	%	69	45	105	E	E	101	65	145	e	e
6	6	006	ACK (acknowledge)	38	26	046	&	&	70	46	106	F	F	102	66	146	f	f
7	7	007	BEL (bell)	39	27	047	'	'	71	47	107	G	G	103	67	147	g	g
8	8	010	BS (backspace)	40	28	050	((72	48	110	H	H	104	68	150	h	h
9	9	011	TAB (horizontal tab)	41	29	051))	73	49	111	I	I	105	69	151	i	i
10	A	012	LF (NL line feed, new line)	42	2A	052	*	*	74	4A	112	J	J	106	6A	152	j	j
11	B	013	VT (vertical tab)	43	2B	053	+	+	75	4B	113	K	K	107	6B	153	k	k
12	C	014	FF (NP form feed, new page)	44	2C	054	,	,	76	4C	114	L	L	108	6C	154	l	l
13	D	015	CR (carriage return)	45	2D	055	-	-	77	4D	115	M	M	109	6D	155	m	m
14	E	016	SO (shift out)	46	2E	056	.	.	78	4E	116	N	N	110	6E	156	n	n
15	F	017	SI (shift in)	47	2F	057	/	/	79	4F	117	O	O	111	6F	157	o	o
16	10	020	DLE (data link escape)	48	30	060	0	0	80	50	120	P	P	112	70	160	p	p
17	11	021	DC1 (device control 1)	49	31	061	1	1	81	51	121	Q	Q	113	71	161	q	q
18	12	022	DC2 (device control 2)	50	32	062	2	2	82	52	122	R	R	114	72	162	r	r
19	13	023	DC3 (device control 3)	51	33	063	3	3	83	53	123	S	S	115	73	163	s	s
20	14	024	DC4 (device control 4)	52	34	064	4	4	84	54	124	T	T	116	74	164	t	t
21	15	025	NAK (negative acknowledge)	53	35	065	5	5	85	55	125	U	U	117	75	165	u	u
22	16	026	SYN (synchronous idle)	54	36	066	6	6	86	56	126	V	V	118	76	166	v	v
23	17	027	ETB (end of trans. block)	55	37	067	7	7	87	57	127	W	W	119	77	167	w	w
24	18	030	CAN (cancel)	56	38	070	8	8	88	58	130	X	X	120	78	170	x	x
25	19	031	EM (end of medium)	57	39	071	9	9	89	59	131	Y	Y	121	79	171	y	y
26	1A	032	SUB (substitute)	58	3A	072	:	:	90	5A	132	Z	Z	122	7A	172	z	z
27	1B	033	ESC (escape)	59	3B	073	;	;	91	5B	133	[[123	7B	173	{	{
28	1C	034	FS (file separator)	60	3C	074	<	<	92	5C	134	\	\	124	7C	174	|	
29	1D	035	GS (group separator)	61	3D	075	=	=	93	5D	135]]	125	7D	175	}	}
30	1E	036	RS (record separator)	62	3E	076	>	>	94	5E	136	^	^	126	7E	176	~	~
31	1F	037	US (unit separator)	63	3F	077	?	?	95	5F	137	_	_	127	7F	177		DEL

Figure 3: ASCII table

Part 1

	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	{0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe},															
1	{0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0},															
2	{0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7},															
3	{0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa},															
4	{0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4},															
5	{0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3},															
6	{0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1},															
7	{0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf },															
8	{0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2},															
9	{0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5},															
A	{0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb},															
B	{0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6},															
C	{0x9, 0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8},															
D	{0xd, 0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9},															
E	{0xc, 0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd},															
F	{0xe, 0xf, 0x7, 0x6, 0x4, 0x5, 0x1, 0x0, 0x2, 0x3, 0xb, 0xa, 0x8, 0x9, 0xd, 0xc}};															

Figure 4: The map table

Part 1

Step 1: Extract ph, pl from the byte p and kh, kl from the byte k
The first plaintext letter is 'h' and the first key letter is 'k'

ASCII code, in HEX, for 'h' is 0x68 and for 'k' is 0x6B

So, ph = 0x06, pl = 0x08, and
kh = 0x06, kl = 0x0B

Step 2: Define ch and cl as:

ch <- map [ph] [kl]

cl <- map [pl] [kh] so by substituting we have,

ch <- map [0x06] [0x0B]

cl <- map [0x08] [0x06] ,

ch=0x07

cl =0x0C

Step 3: Combine ch and cl into the byte c, so c=0x7C

Part 1

What about the decryption?

- How can we use the information provided by the map table and encryption diagram for decryption?
- Looking for p_h and p_l given c_h, c_l and k_h, k_l
 - $c_h = \text{map}[?][k_l]$
 - $c_l = \text{map}[?][k_h]$

Part 1

▪ Decryption Example

From the encryption example we have:

Ciphertext byte: $c=0x7C$ (as extracted from the encryption example)

Key byte: $k=0x6B$

Step 1: Extract ch , cl from the ciphertext byte and kh , kl from the key character

So, $ch = 0x07$, $cl = 0x0C$, and

$kh = 0x06$, $kl = 0x0B$

Step 2: We have the formulas

$ch \leftarrow \text{map}[ph][kl]$,

$cl \leftarrow \text{map}[pl][kh]$,

How to find ph and pl using the map?

- $\text{map}[?][0x0B] = 0x07$
 - Which row in column B has value of 0x07?
 - Row 0x06
 - So, $ph=0x06$
- $\text{map}[?][0x06] = 0x0C$
 - Which row in column 6 has value of 0x0C?
 - Row 0x08
 - So, $pl=0x08$

Step 3: Combine ph and pl into the byte p , so $p=0x68$, ASCII for 'h'

Part 1

Time for some HINTS ...

Hint 1: The plaintext is printable ASCII, and the key is a combination of **printable and non-printable (control) ASCII characters**. Use these facts when searching for the key.

Hint 2: How could you use the frequencies of character occurrence in ASCII text of a language to automate the process of recognizing the right key?

Part 2

Goal: To decrypt “ciphertext2”

Available Data

- **Plaintext:** Not a regular text file, but some other commonly used file format
- **Key:** Any combination of printable ASCII characters

Lookup Table Map[x][k]

- **Map [16] [16]:** Will ultimately give, the ciphertext's 4 higher or lower bits (ch,cl)
 - Rows: x, 4 bit quantity
 - Columns: k, 4 bit quantity

Part 2

Hint 1: Consider the possibility that the file format corresponding to *ciphertext2* may not observe standard frequency characteristics.

Hint 2: The key for *ciphertext2* can be any combination of **printable** ASCII characters. Thus, the key used to encrypt *ciphertext2* is **not** restricted the same way as the key that encrypted *ciphertext1*.

Hint 3: The key is substantially longer than the one for *ciphertext1*.

QUESTIONS ?