

NANYANG TECHNOLOGICAL UNIVERSITY
School of Electrical & Electronic Engineering

EE/IM4152 Digital Communications

Tutorial No. 12 (Sem 1, AY2016-2017)

1. For linear block codes, we adopt the maximum-likelihood (ML) strategy for decoding. We decide the codeword \mathbf{c}_i if $\Pr(\mathbf{r}|\mathbf{c}_i)$ has the largest value among all codewords, where \mathbf{r} is the received word. Suppose the Hamming distance between \mathbf{r} and \mathbf{c}_i is d_i . That is, the channel noise causes errors in d_i positions and the Hamming weight of $w(\mathbf{r} \oplus \mathbf{c}_i) = d_i$.
 - (a) Determine $\Pr(\mathbf{r}|\mathbf{c}_i)$ if the channel bit-error probability is ε over a binary symmetric channel (BSC), where $\varepsilon < 0.5$.
 - (b) Taking the logarithm of $\Pr(\mathbf{r}|\mathbf{c}_i)$ will not affect the comparisons. Show that $\max_i \Pr(\mathbf{r}|\mathbf{c}_i)$ is equivalent to $\min_i [d_i]$.
 - (c) Show that the result in part (b) is equivalent to choosing the error pattern with the minimum Hamming weight.

2. For a t -error-correcting (n, k) linear block code, the received word is in error if more than t errors occur in n coded bits (i.e., the worst case).
 - (a) Suppose the probability of channel bit error is q_c . Determine the probability of word error, P_{ew} , in terms of n , t and q_c .
 - (b) Given that a received word is wrong, show that the average probability of bit error is

$$P_{ec} = \sum_{j=t+1}^n \binom{n-1}{j-1} (1-q_c)^{n-j} q_c^j.$$

3. For the convolutional encoder with $(n, k, N) = (2, 1, 3)$, as shown in Fig. 1, design a trellis diagram to generate the encoded sequence corresponding to the information sequence $\mathbf{x} = [1101011]$. As observed from Fig. 1, the output bits are not only determined by the present $k=1$ information bit but also by the previous information bits. This dependence on the previous information bits causes the encoder to be a finite state machine. Here we use the contents of (d_{k-2}, d_{k-1}) to define the state of the encoder: $a = (0, 0)$, $b = (1, 0)$, $c = (0, 1)$ and $d = (1, 1)$. Remember to append $N-1=2$ zero bits to end the transmission of the current block $\mathbf{x} = [1101011]$, which is different from the truncation approach used in the lecture notes. In this way, the trellis path will be forced to return to state $a = (0, 0)$ due to those appended zeros.

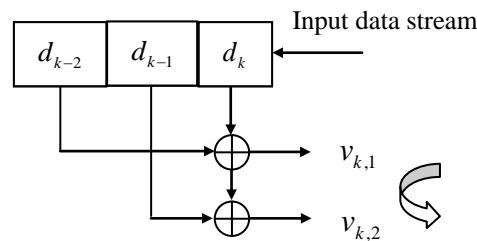


Figure 1: Convolutional encoder with $(n, k, N) = (2, 1, 3)$

Assuming hard-decision decoding, the received sequence is $\mathbf{y} = [01101111010001]$. Find the maximum-likelihood information sequence and the corresponding number of errors contained in \mathbf{y} .