

Efficacy of noisy dynamical decoupling

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(To be edited.) This paper study dynamical decoupling (DD), a family of pulse schemes aimed at removing noise in quantum systems. We first review the formalism of DD with an emphasis on the error phase, a figure-of-merit characterizing the performance of DD. We then develop noise suppression criteria from the error phase perspective for the periodic and concatenated DD schemes. A rigorous proof for the noise decoupling order of concatenated DD is also presented.

I. INTRODUCTION

The adverse effects of noise remains the single biggest obstacle to the realisation of large-scale quantum technologies. The very quantum effects that give quantum technologies their edge over classical devices are extremely fragile and easily destroyed by the presence of unwanted interactions with the environment—noise. Much of the current research and technological push in the community today are centered around implementing methods to remove the effects of noise in quantum devices.

A key noise-removal technique is dynamical decoupling (DD) [1–11], a low-resource-cost approach that only requires application of fast pulse sequences on individual qubits in the quantum device to average away the effects of noise processes slow compared to the pulse time. With its roots in spin-echo techniques in NMR systems, DD has been used in many different types of experiments as a simple way to reduce noise in quantum information processing systems. Such active noise removal scheme is in direct contrast with the widely studied approach of quantum error correction (QEC) [12]. Compared with QEC, dynamical decoupling is typically more economical as it require no encoding of logical qubits using multiple physical qubits, nor real-time close-loop control through periodic syndrome measurement. All that is needed are regular single-qubit fast pulses that are usually easy to implement. DD be used by itself, or as the first layer of defense against noise In practice, one can choose to only implement DD or to incorporate DD within a standard QEC scheme as the low-level noise-reduction approach.

The use of DD does not, however, come at no cost. To implement it, one has to apply multiple pulses, which can be imperfect. On the one hand, if the DD pulses are perfect, the slow noise will be averaged away, leaving weaker residual noise on the system; on the other hand, if the DD pulses are imperfect, those imperfections can add errors to the system, and if those errors happen often enough, they can eliminate the benefit of having DD in the first place. Like the cost-benefit analysis familiar to fault-tolerant quantum computing, one has to ask whether there is an error threshold for DD, limiting the level of imperfections allowed in the DD pulses above

which DD simply cannot offer any benefit.

More to come.

II. PRELIMINARIES

A. Basics of DD

DD involves the repeated application of a fixed sequence of short pulses (or fast gates) to individual quantum registers—usually qubits—that averages away the effect of any noise with a time scale slow compared to the sequence time. For a given DD scheme, let L be the number of pulses in the sequence. We denote the i th pulse of the sequence as P_i , and write the sequence as $P_L \dots P_2 P_1$ (I moved the equation here to inline), proceeding from right to left in time. Pulse P_i is applied at time $t_i = t_{i-1} + \tau_i$ where t_{i-1} is the time of the application of P_{i-1} , and t_0 is the start time of the sequence. We illustrate a typical sequence of DD control gates in Figure 1.

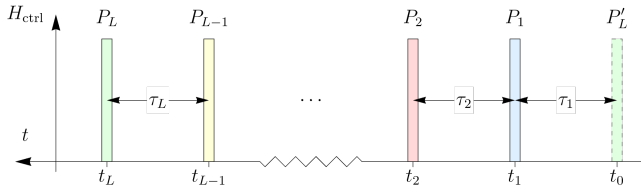


FIG. 1. Illustration of the external control pulses for a typical DD sequence. Different gates are filled with different colors. The dashed gate at t_0 corresponds to the previous DD period.

All DD schemes share some fundamental similarities. First, the pulses are chosen from a specified group of transformations $P_i \in \mathbf{G}$, for $i = 1, 2, \dots, L$. Second, any DD sequence satisfies the constraint

$$P_L P_{L-1} \dots P_1 \doteq I, \quad (1)$$

where the dot above the equal sign means “up to an overall phase factor”. This constraint ensures that there is no net transformation on the quantum register at the completion of the sequence. Different strategies could differ in the following aspects: (i) the transformation group \mathbf{G} .—For a qubit register, an obvious choice is perhaps the Pauli group generated by the Pauli operators X, Y , and Z . Simpler schemes such as spin echo [13] and the CPMG sequence [14, 15] use only the subgroup $\{I, Z\}$. (ii) the sequence length L and the specific sequence of pulses. (iii) the pulse times.—One can have regular-interval pulses, with $\tau_i \equiv \tau \forall i$, as is the case in the periodic DD (PDD) [1] and concatenated DD (CDD) [4] schemes. One could, however, have variable-interval schemes, such as the Uhrig DD (UDD) sequence [7], nested UDD [9], and quadratic DD [11]. In this work, we focus our analysis on the regular-interval schemes of PDD and CDD, differing in sequence length and the specific pulse sequence, but which both employ pulses drawn from the Pauli group.

The physical model we are considering is one where gates, including DD pulses, can be applied at some operating frequency, with time τ in between consecutive

gates. A nonzero τ can be simply the finite switch time between gates, or there may be other practical reasons for a synchronized clock cycle time. We model the DD pulses as instantaneous gates taking no time. There are two sources of errors: (i) the system—the quantum register—interacts with a bath, and that interaction is the source of noise on the system, present even in the absence of the DD pulses; (ii) the pulse imperfections. In past work (see, for example, Ref. [4]), imperfect DD pulses were modeled as finite-duration pulses during which the always-on system-bath interaction acts and leads to errors in the pulses. Here, we allow for the more general, and practically relevant, situation of additional control errors in the pulses. We model noisy finite-duration pulses as instantaneous ideal pulses followed by a noise map that captures the background noise as well as any additional control errors. DD sequences are ideally designed to average away the background noise arising from the system-bath interaction. We refer to the original DD sequences with perfect instantaneous pulses as ideal DD, and to the case with imperfect pulses as noisy DD.

In the absence of any DD pulses, the system and bath evolve jointly according to the Hamilton operator H , assumed to be time-independent as is appropriate for standard DD analysis (put differently, any parametric time dependence in the joint dynamics occurs because of degrees of freedom excluded from the bath; here we think of all such degrees of freedom as part of the bath). H here can be written as $H = H_B + H_{SB}$, with H_B as the bath-only Hamilton operator, and H_{SB} the interaction. No system-only term appears in H as we assume no nontrivial dynamics (other than that arising from H_{SB}) occur in the system during the DD period. With this, we can write the evolution operator for ideal DD as

$$U_{DD} \equiv P_L e^{-i\tau H} P_{L-1} \dots P_2 e^{-i\tau H} P_1 e^{-i\tau H} \equiv e^{-i\Omega_{DD}}. \quad (2)$$

Here, we have defined $\Omega_{DD} \equiv T H_{\text{eff}}$ as the dimensionless (with the T factor) effective Hamilton operator appropriate for describing evolution for the time of the DD sequence. Ω_{DD} can be written down formally using the Magnus expansion,

$$\Omega_{DD} = \sum_{m=1}^{\infty} \Omega_{DD}^{(m)}, \quad (3)$$

where the m th term consists of products of m copies of τH , and hence of order $\|\tau H\|^m$. We refer the reader to past analyses of DD (cite) for a detailed derivation; here we provide only the basic expressions needed for our discussion below. (Add back a bit of the derivation of the $\Omega_{DD}^{(1)}$ and $\Omega_{DD}^{(2)}$ terms.) Within the radius of convergence of the Magnus series, $\Omega_{DD}^{(m)}$ decreases in importance as m increases. A DD scheme such that $\Omega_{DD}^{(m)}$ acts trivially on the system (i.e., acts as I on system) for all $m \leq n$ is said to achieve n th-order decoupling—the system sees weakened noise, of strength $\|\Omega_{DD}^{(n)}\| \sim \|\tau H\|^n$, compared to $\|\Omega_{DD}^{(1)}\| \sim \|\tau H\|$ without DD.

B. Quantifying the efficacy of DD

(More to come. To put back parts of the fidelity, etc. discussion, as needed in the remainder of the text.)

1. Decoupling Order

According to Eq. (??), the DD protocol can be regarded as a transformation $\Omega \rightarrow \Omega_{\text{DD}}$. Since the noise map \mathcal{E} is a function of Ω and ρ_B , the transformation on Ω takes the original quantum channel into a new one $\mathcal{E} \rightarrow \mathcal{E}_{\text{DD}}$. As is argued in Sec. ??, error phase is a convenient characterization of noise strength, we apply the two-component norm to the dimensionless Hamiltonian before and after DD, leading to

$$\|\Omega\| \equiv (\alpha, \beta) \quad \text{and} \quad \|\Omega_{\text{DD}}\| \equiv (\alpha_{\text{DD}}, \beta_{\text{DD}}). \quad (4)$$

The “decoupling” reference in the name of DD is understood in terms of the coupling part: i.e., the error phase of the new channel gets reduced through this transformation. We say noise is suppressed if the ratio $\beta_{\text{DD}}/\beta < 1$, and a lower ratio would indicate better noise suppression.

A widely used indicator to characterize the performance of DD schemes is the *decoupling order*. A DD scheme is said to achieve n th-order decoupling if the first n terms in the series (3) have vanishing coupling part:

$$\mathcal{P}_c(\Omega_{\text{DD}}^{(k)}) = 0, \quad \text{for } 1 \leq k \leq n. \quad (5)$$

The decoupling order is thus a single positive integer to grade different DD schemes. Greater decoupling order is usually linked with better noise suppression rate. This can be seen by evaluating the norm of the Magnus series projected on to the coupling part. For an n th-order decoupling sequence, we have

$$\beta_{\text{DD}} = \left\| \sum_{k=n+1}^{\infty} \mathcal{P}_c(\Omega_{\text{DD}}^{(k)}) \right\| \simeq (\alpha, \beta)^{n+1}, \quad (6)$$

where we use the notation $(\alpha, \beta)^{n+1}$ to represent a polynomial that is $(n+1)$ th order in α and β . Provided that the Magnus series converges absolutely, β_{DD} is bounded by a power series in $\|\Omega\|$ whose leading term is on the order of $n+1$. Therefore Eq. (6) provides an alternative characterization for the n th-order decoupling condition, which can be useful when directly calculating the first n Magnus terms is unfeasible.

2. Error Phase

To gauge the efficacy of DD, we need to quantify the deviation of the actual state of the quantum system, with and without DD, from the ideal, no-noise state. Following Ref. [5], we make use of the *error phase*, which measures the strength of the system-bath interaction,

the source of the noise on the system. The system and bath evolve jointly for some specified time T according to the evolution operator $U(0, T)$. The underlying joint Hamilton operator generating the dynamics can be time-dependent, and can include—or not—the DD pulses on the system. We write $U(0, T) \equiv e^{-iTH_{\text{eff}}}$, for some effective time-independent system-bath Hamilton operator H_{eff} . H_{eff} can be split into two pieces: $H_{\text{eff}} \equiv H_{\text{eff},B} + H_{\text{eff},SB}$, where $H_{\text{eff},B} \equiv \frac{I_S}{d_S} \otimes \text{tr}_S(H_{\text{eff}})$ acts on the bath alone, while $H_{\text{eff},SB} \equiv H_{\text{eff}} - H_{\text{eff},B}$ contains all the pieces that act nontrivially on the system. $H_{\text{eff},SB}$ can be thought of as the effective system-bath interaction over this time T . We define the error phase as the norm of that interaction, multiplied by the time interval,

$$\Phi_{\text{SB}} \equiv T\|H_{\text{eff},SB}\|. \quad (7)$$

We use a norm such that $\|S \otimes B\| = \|S\|\|B\|$, and $\|I\| = 1 = \|X\| = \|Y\| = \|Z\|$ for single-qubit I, X, Y , and Z .

As we will see, the norm of $H_{\text{eff},B}$ will also enter our analysis. We define the corresponding error phase for the bath-only term as $\Phi_B \equiv \|H_{\text{eff},B}\|T$, and also “two-component” error phase as

$$\Phi \equiv T\|H_{\text{eff}}\| \equiv (\Phi_B, \Phi_{\text{SB}}) \quad (8)$$

(To merge with new text.) Following the properties of operator norm, thus-defined map is absolutely-scalable $\|a\Omega\| = |a| \cdot \|\Omega\|, \forall a \in \mathbb{C}$, and satisfies the inequality, $\|A+B\| \leq \|A\| + \|B\|$, where the “ \leq ” sign between number pairs is defined by each of the components: $(a_1, b_1) \leq (a_2, b_2)$ if both $a_1 \leq a_2$ and $b_1 \leq b_2$; vice versa for the “ \geq ” sign. It can also be easily verified that the two-component norm is invariant under arbitrary disjoint unitary transformation on $\mathcal{H}_S \otimes \mathcal{H}_B$:

$$\|(U_S \otimes U_B)\Omega(U_S^\dagger \otimes U_B^\dagger)\| = \|\Omega\|, \quad \forall U_S(U_B) \in \mathcal{U}(d_{S(B)}), \quad (9)$$

where $\mathcal{U}(d_{S(B)})$ is set of unitary operators on \mathcal{H}_S or \mathcal{H}_B . This invariance suggests that the error phase is invariant under arbitrary interaction picture defined by a system or bath evolution.

3. Channel fidelity measures

C. PDD and its concatenation

We specialize here to the case of interest, that of the single-qubit PDD scheme. The single qubit interacts with a bath, with a joint Hamilton operator (in the absence of DD) that can be written, without loss of generality, as

$$H \equiv I \otimes B_I + X \otimes B_X + Y \otimes B_Y + Z \otimes B_Z, \quad (10)$$

where I is the identity operator on the qubit, X, Y , and Z are the Pauli operators on the qubit, and B_i , for $i = I, X, Y, Z$, are operators on the bath. We identify $I \otimes B_I$

as the bath-only H_B and $X \otimes B_X + Y \otimes B_Y + Z \otimes B_Z$ as the interaction Hamilton operator H_{SB} , and define

$$\|\tau H\| = (\|\tau H_B\|, \|\tau H_{SB}\|) \equiv (\phi_B, \phi_{SB}). \quad (11)$$

(I moved τ inside the norm operator to conform the convention that norm should be taken on a dimensionless operator, not on Hamiltonian alone. But in the coming texts there are many violations.)

The PDD scheme uses a simple 4-pulse sequence $ZXZX$. The corresponding Ω_{DD} of Eqs. (2) and (3), now re-labeled as Ω_{PDD} , can be worked out to be

$$\Omega_{PDD} = \Omega_{PDD}^{(1)} + \Omega_{PDD}^{(2)} + \dots \quad (12)$$

$$\text{with } \Omega_{PDD}^{(1)} = (4\tau)I \otimes B_I \quad (13)$$

$$\Omega_{PDD}^{(2)} = -(2\tau^2)\{X \otimes 2i[B_I, B_X] + Y \otimes (i[B_I, B_Y] + \{B_X, B_Z\})\}, \quad (14)$$

where $[\cdot, \cdot]$ and $\{\cdot, \cdot\}$ are the commutator and anti-commutator, respectively. Ω_{PDD} should be compared against the evolution operator (over the same time period of 4τ) in the absence of DD: $U = e^{-i4\tau H} \equiv e^{-i\Omega}$, with

$$\Omega = (4\tau)H = (4\tau)(I \otimes B_I + H_{SB}) \quad (15)$$

Comparing $\Omega_{PDD}^{(1)}$ —usually the dominant term—with Ω , we see that the H_{SB} in Ω no longer appears in $\Omega_{PDD}^{(1)}$ and $\Omega_{PDD}^{(1)}$ is trivial on the system. This corresponds to the fact that PDD is able to remove the lowest-order noise and that it achieves first-order decoupling. (To revisit to see if a re-phrasing is needed. I'm worried about the $[B_I, B_X]$ term in $\Omega_{PDD}^{(2)}$, which is of order $\|H_B H_{SB}\|$, not $\|H_{SB}\|^2$. Also the last sentence of the paragraph following Eq. (3).)

[I am changing most of the Ω -like quantities into τH s (except when we talk about the Magnus series) – it is important to keep the time parameter visible. Sometimes, τ is the relevant time—for the no DD case—sometimes $K\tau$ is the relevant one—for the DD case. These should just be visible in the analysis, not kept hidden within Ω . This becomes especially important when considering CDD.]

A given DD sequence yields a given decoupling order, setting a limit on the scheme's ability to reduce noise in the system. To increase the power of the DD scheme, one can employ the method of concatenation introduced in Ref. [4]. In that work, concatenated DD (CDD) was built upon the basic PDD scheme; the same procedure of concatenation, however, can be applied to other basic DD sequences. The idea is to make use of concatenation to increase the decoupling order of the resulting DD sequence, hence scaling up the protection of the system against noise.

CDD can be described in a recursive manner. We begin with the bare evolution, without any DD sequence,

writing the evolution operator over some time interval t as

$$U_0(t) \equiv e^{-itH_0}, \quad (16)$$

where $H_0 \equiv H$, with the subscript 0 added here in preparation for concatenation to higher levels. With a DD scheme of L pulses, applied at time intervals (assuming regular-interval DD) τ_0 , the evolution operator is

$$U_1(\tau_1) \equiv P_L U_0(\tau_0) \dots P_2 U_0(\tau_0) P_1 U_0(\tau_0), \quad (17)$$

where $\tau_1 \equiv L\tau_0$. The subscript 1 is to be understood as indicating that this is concatenation level 1. To concatenate further, $U_k(\tau_k)$ is defined recursively,

$$U_{k+1}(\tau_{k+1}) \equiv P_L U_k(\tau_k) \dots P_2 U_k(\tau_k) P_1 U_k(\tau_k), \quad (18)$$

increasing the concatenation level by 1 each time, and with $\tau_{k+1} \equiv L\tau_k$. With each $U_k(\tau_k)$, we associate an effective Hamilton operator H_k , and a dimensionless Hamilton operator $\Omega_k \equiv \tau_k H_k$, such that

$$U_k(\tau_k) \equiv e^{-i\tau_k H_k} = e^{-i\Omega_k}. \quad (19)$$

Each CDD scheme is determined by specifying the maximal concatenation level n , and either the value of τ_0 or τ_n . CDD at level n , denoted as CDD_n , is then a sequence of $L_n \equiv L^n$ pulses separated by time interval τ_0 and taking total time $\tau_n = L^n \tau_0$ to complete.

In the remainder of the paper, we will restrict our discussion to CDD built upon the basic PDD scheme. The authors of Ref. [4] showed that CDD_n achieves n th order decoupling. This quantifies the benefit of scaling up the noise protection by concatenation. Appendix A re-derives this conclusion with a different and what we believe to be clearer argument than that in the original reference.

III. THE BREAK-EVEN CONDITIONS FOR IDEAL AND NOISY PDD

Now we have introduced the basic principles of DD, a natural question will be how

Any noise mitigation strategy is effective only if its costs—the increased complexity of carrying out the quantum computational task—are lower than its benefits—the increased ability to remove adverse effects of the noise. The costs enter not just because any noise mitigation approach requires the use of more resources, e.g., more gates, more qubits, etc., but more that those added resources are themselves imperfect in practice, such that the unmitigated noise of the resulting larger system is unavoidably larger than if no noise mitigation strategy was adopted. There is then a net benefit only if the added noise is low enough to not overwhelm the added noise-removal capabilities. Such is the content of any fault-tolerant analysis, usually applied to quantum computing tasks protected by quantum error correction. Here, we apply the same logic to DD, and ask when the added benefit of averaging away part of the noise outweighs the added cost of having to do additional DD pulses.

A. Ideal case

Let us first discuss when ideal PDD is useful—we expect to recover the usual statement of when the idea of PDD works at all. We choose the error phase as our figure of merit: $\epsilon \equiv \Phi$. For PDD to be useful, we require, under the computational setting,

$$\Phi_{\text{SB}} \leq \phi_{\text{SB}}. \quad (20)$$

From our earlier discussion of PDD, it is clear that $\Phi_{\text{SB}} = \|\Omega_{\text{PDD}}^{(2)} + \Omega_{\text{PDD}}^{(3)} + \dots\|$. The full Magnus series is difficult to write down, but we can employ the dominant term, $\Omega_{\text{PDD}}^{(2)}$ which acts nontrivially on the system, for the approximate condition:

$$\Phi_{\text{SB}} \simeq \|\Omega_{\text{PDD}}^{(2)}\| \leq \phi_{\text{SB}}. \quad (21)$$

Using the expression for $\Omega_{\text{PDD}}^{(2)}$ from Eq. (14), we have

$$\begin{aligned} \|\Omega_{\text{PDD}}^{(2)}\| &\leq 2\tau^2(2\|[B_I, B_X]\| + \|[B_I, B_Y]\| + \|[B_X, B_Z]\|) \\ &\leq 4\phi_B\tau(2\|B_X\| + \|B_Y\|) + 4\tau^2\|B_X\|\|B_Z\| \\ &\leq 12\phi_B\phi_{\text{SB}} + 4\phi_{\text{SB}}^2 \end{aligned} \quad (22)$$

noting that $\phi_B \equiv \tau\|H_B\| = \tau\|B_I\|$, and that $\phi_{\text{SB}} \equiv \tau\|H_{\text{SB}}\| = \tau\left[\sum_{i=X,Y,Z}\|B_i\|^2\right]^{1/2} \geq \tau\|B_i\|$ for any $i = X, Y, Z$ (The first equality is true under the H-S induced norm, but missing a factor of 2, as written; can be remedied by using normalized Pauli operators. Need to fix this throughout the paper). To get condition (21), it then suffices to require

$$12\phi_B + 4\phi_{\text{SB}} \leq 1. \quad (23)$$

This amounts to a requirement that both ϕ_{SB} and ϕ_B be small for PDD to work well. That ϕ_{SB} has to be small comes as no surprise— ϕ_{SB} quantifies the strength of the noise on the system, and DD is expected to work well as long as the noise is weak enough so that the remnant noise is weak. The requirement that ϕ_B be small is perhaps a bit more surprising – after all, it bounds the bath-only term which does not directly lead to noise on the system. Nevertheless, H_B determines the evolution rate of the bath. The requirement that ϕ_B also has to be small should be understood as the requirement that any DD scheme requires the noise to remain more or less constant throughout the sequence, for good averaging by the pulses. ϕ_B quantifies how rapidly the bath evolves, and hence how quickly the noise seen by the system changes (through the evolution of the bath operators B_i in the Heisenberg picture). A slowly evolving noise requires then a small ϕ_B . Tracing back to the original calculation, we see that part of the pre-factor of 12 of ϕ_B is the “4” for the length of the PDD sequence; a longer sequence will have a larger pre-factor and hence a more stringent requirement that ϕ_B is small, consistent with our interpretation here and our standard understanding of DD.

(To be edited and merged with the new text.) On the other hand, the inequalities (??) inspires use to use the following approximation for the error phase: (The following is exactly true, not just approximately so, up to these factors of 2 that can be fixed – but don't matter for the final inequality for the threshold condition – for the Hilbert-Schmidt (H-S) induced norm. What is the norm used in producing Fig. 2? If that was the H-S norm, that would explain why this version of the inequality works well.) (That is the infinite norm, i.e., the largest eigenvalue, and consistently used across the original script. The reason was explained in the original error phase section. Basically the norm need to satisfy $\|I\| = 1$ for whatever dimensionality of the identity. This is necessary for attaching an ancilla system without changing the error phase.—Jiaan)

$$\left\|\sum_{i=1}^3 \sigma_i \otimes B_i\right\| \approx \sqrt{\sum_{i=1}^3 \|B_i\|^2}. \quad (24)$$

Such approximation is valid up to a factor $\in [1/\sqrt{3}, \sqrt{3}]$. This approximation allows us to derive an “approximate condition”. We estimate

$$\begin{aligned} \beta_{\text{PDD}}^2 &\approx -4i[B_0, B_1]\|^2 + \|-2i[B_0, B_2] - 2\{B_1, B_3\}\|^2 \\ &\leq (8\alpha\|B_1\|)^2 + (4\alpha\|B_2\| + 4\|B_1\|\|B_3\|)^2, \end{aligned} \quad (25)$$

Defining $x_i \equiv \|B_i\|/\beta$, then upper-bounding the r.h.s. by β^2 requires

$$\max_{x_1^2+x_2^2+x_3^2=1} (2\alpha x_1)^2 + (\alpha x_2 + \beta x_1 x_3)^2 \leq \frac{1}{16}. \quad (26)$$

Solution of this constrained optimization problem leads

to:

$$\left(\frac{\beta}{\sqrt{3}} \leq \alpha \leq \frac{1}{8}\right) \text{ or } \left(10\alpha^2 + \frac{9\alpha^4}{\beta^2} + \beta^2 \leq \frac{1}{4}\right). \quad (27)$$

This set of inequalities automatically implies (??), it gives a broader but less confident condition due to the approximation used for the error phase.

The criteria (??) and (27) each determines a region in the (α, β) parameter space where noise suppressing is achieved using the leading order Magnus term. But in order for the calculations to make sense, we have assumed that the Magnus series converge absolutely. According to (??), the necessary condition here is $4\|\Omega\|_\infty < \xi$. Since $\|\Omega\| \leq \alpha + \beta$, we may use a more stringent convergence bound:

$$\alpha + \beta \leq \frac{1}{4}. \quad (28)$$

To examine the accuracy of the two theoretical bounds, we perform numeric tests as a verification. We first randomly generate four 2×2 Hermitian matrices for the bath operators satisfying the normalization condition in Eq. (??), then simulate the PDD protocol and use matrix logarithm to find the actual Ω_{PDD} and error phase β_{PDD} . For each fixed (α, β) point, we numerically maximize the error phase reduction ratio β_{PDD}/β over 1000 random instances. The result is presented in Figure 2, where a point in blue color represents a greater-than-one ratio and yellow color represents smaller-than-one ratio. Thus the graph reflects the true noise suppressing condition by the continuous yellow-colored region. We also plot the boundaries for the two theoretical noise suppression region from (??) and (27), together with the absolute convergence region (28) for comparison.

As long as the two-component norm of the free evolution Hamiltonian falls in the yellow region on the lower-left corner, the idealist PDD scheme is guaranteed to work, in the sense of reducing the error phase. This reflects the fundamental “small noise” assumption of DD. From the numerical results, we learn that the bound by the “strict condition” from (??) is too loose compared with the “approximate condition” (27), which works surprisingly well in reflecting the true boundary of the noise-reduction region. The faithfulness of the latter bound is particularly striking considering the points beyond the absolute-convergence boundary (above the green dotted line), where our derivation should be unreliable in principle. An explanation is that the absolute-convergence criteria (28), sufficient but not necessary, is a loose one and the leading order Magnus series works well beyond the region predicted by this inequality. This piece of evidence boosts our confidence in using the leading order Magnus term as approximation in the following parts of this work. But for safety, the absolute convergence region will also be indicated.

B. Noisy pulses

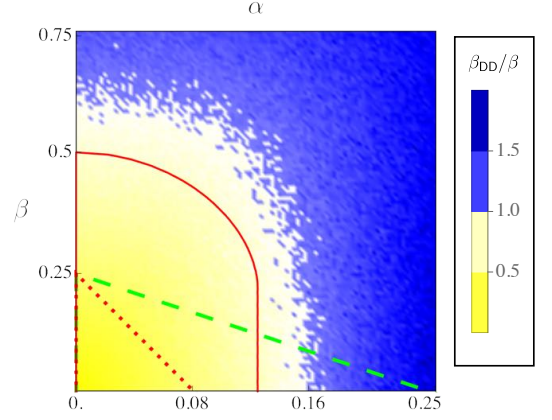


FIG. 2. The noise reduction ratio β_{DD}/β obtained through maximization over random bath operators plotted in the (α, β) parameter plane. In comparison, the theoretical boundaries for noise reduction are also plotted, with dotted red line for (??) and solid red line for (27), also with the boundary for the absolute-convergence condition (28) in dashed green line. Note the ranges for α and β are chosen differently for better visual effect.

(I have rewritten this subsection to address the problems raised by Hui Khoon.) Previously, we treated the DD control gates in an ideal manner, which implies instantaneous and noise-free pulses. In practice, of course, the control gates are noisy quantum channels that need not be unitary. However, one may always include an ancillary system to transform a quantum channel into a unitary one over the extended Hilbert space. For the analytical treatment below, we will always assume this dilation for the actual gates and marked the extended unitary operators with overtilde, e.g., \tilde{P}_i . A general DD sequence under those noisy pulses are then

$$\tilde{U}_{\text{DD}} = \tilde{P}_L e^{-i\Omega} \dots \tilde{P}_1 e^{-i\Omega} \equiv e^{-i\tilde{\Omega}_{\text{DD}}}, \quad (29)$$

where we also put tildes over U_{DD} and Ω_{DD} to remind that the DD control gates are noisy. The overall goal is to figure out the conditions required by the noisy gates such that the noise-inflicted DD error phase $\tilde{\Phi}_{\text{SB}} \equiv \|\tilde{\Omega}_{\text{DD}, \text{SB}}\|$ is still reduced compared to the bare Hamiltonian error phase $\phi_{\text{SB}} = \|\Omega_{\text{DD}, \text{SB}}\|$. For simplicity, let us employ a single parameter η , which is determined later, to quantify the noise strength associated with the gates. The break-even condition is, in essence, to invert the inequality

$$\tilde{\Phi}_{\text{SB}}(\phi_{\text{B}}, \phi_{\text{SB}}, \eta) < \phi_{\text{SB}}, \quad (30)$$

and obtain an upper bound on the noise level $\eta < \eta_{\text{max}}(\phi_{\text{B}}, \phi_{\text{SB}})$. To achieve this, we need to correctly estimate the DD error phase function $\tilde{\Phi}_{\text{SB}}(\phi_{\text{B}}, \phi_{\text{SB}}, \eta)$ while also looking for an appropriate definition for the noise strength η associated with the gates.

1. Gate-independent noise

For each noisy gate \tilde{P}_i , we may separate out a “pure error” map, $P_i^\dagger \tilde{P}_i$, by applying the inverse gate. As the resulting operator is unitary and close to identity, we may associate to it a Hermitian generator Γ_i to write the noisy gate as

$$\tilde{P}_i = P_i e^{-i\Gamma_i}. \quad (31)$$

The evolution of the system and bath under noisy pulses thus becomes

$$e^{-i\tilde{\Omega}_{\text{DD}}} = P_L e^{-i\Gamma_L} e^{-i\Omega} \dots P_1 e^{-i\Gamma_1} e^{-i\Omega}. \quad (32)$$

The already established map $\Omega \rightarrow \Omega_{\text{DD}}$ for the idea-pulse DD sequence can no longer be reused, unless if the noise is gate-independent, i.e.,

$$\tilde{P}_i = P_i e^{-i\Gamma}, \quad \forall i. \quad (33)$$

Let us for now focus on such gate-independent noise. The problem readily reduces to the ideal DD case by replacing $\Omega = \tau H$ in our earlier analysis by $\tilde{\Omega} \equiv \tau \tilde{H}$ such that

$$e^{-i\tilde{\Omega}} = e^{-i\Gamma} e^{-i\Omega}. \quad (34)$$

Since Γ and Ω are both small, as required for DD to work, one may invoke the BCH formula to get $\tilde{\Omega}$ as a series of nested commutators. Similar to $\Omega = \Omega_B + \Omega_{\text{SB}}$, let us also split Γ into the pure bath part Γ_B plus the interaction part Γ_{SB} and define $(\eta_0, \eta) \equiv (\|\Gamma_B\|, \|\Gamma_{\text{SB}}\|)$. Assuming Γ and Ω not particularly correlated, we may retain to the second order terms, and then apply triangular inequality to derive the upper bounds for $(\tilde{\phi}_B, \tilde{\phi}_{\text{SB}}) \equiv (\|\tilde{\Omega}_B\|, \|\tilde{\Omega}_{\text{SB}}\|)$:

$$\begin{cases} \tilde{\phi}_B \lesssim (\eta_0 + \phi_B)(1 + \eta + \phi_{\text{SB}}) \equiv \tilde{\phi}_B^{(\text{ub})} \\ \tilde{\phi}_{\text{SB}} \lesssim (\eta + \phi_{\text{SB}})(1 + \eta_0 + \phi_B) \equiv \tilde{\phi}_{\text{SB}}^{(\text{ub})} \end{cases} \quad (35)$$

In Fig. 3, we numerically verify this approximate bounds by randomly generating small Γ and Ω operators and compute $(\tilde{\phi}_B, \tilde{\phi}_{\text{SB}})$ for the combined $\tilde{\Omega}$. The result seems positive with no violation found. With this, we can bound the error phase relevant for this noisy-pulse situation under the different schemes by replacing ϕ_B and ϕ_{SB} for the ideal case with their noisy upper bounds:

$$\tilde{\Phi}_{\text{SB}} \leq \Phi_{\text{SB}}(\tilde{\phi}_B^{(\text{ub})}, \tilde{\phi}_{\text{SB}}^{(\text{ub})}) < \phi_{\text{SB}}. \quad (36)$$

Further bounding the right hand side by ϕ_{SB} then gives us the condition for the break-even points.

The above analysis is generic and applies to all DD schemes with gate-independent noise. The resulting bound is general, but could be suboptimal for a specific DD scheme and a specific type of pulse imperfection. Below, we examine a concrete examples of imperfections for

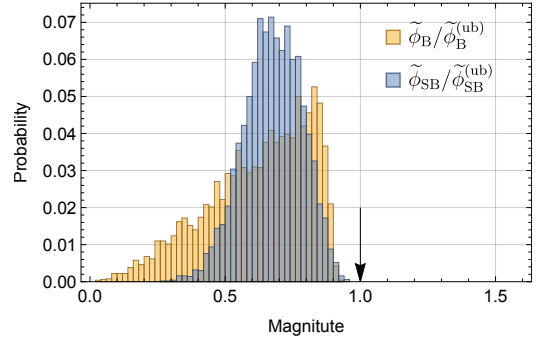


FIG. 3. The histogram for the ratios $\tilde{\phi}_B/\tilde{\phi}_B^{(\text{ub})}$ and $\tilde{\phi}_{\text{SB}}/\tilde{\phi}_{\text{SB}}^{(\text{ub})}$ over 5000 randomly generated Γ and Ω operators (as 4×4 Hermitian matrices). The random Γ s and Ω s are set to follow the normal distribution with norms around 0.01. Our bounds in Eq. (35) is indicated with an arrow.

PDD: global unitary error in the control gates, which results in the noisy pulse

$$\tilde{P}_i = P_i e^{-i\theta \cdot \sigma}, \quad \forall i. \quad (37)$$

We use the notation $\sigma \equiv (X, Y, Z) = (\sigma_1, \sigma_2, \sigma_3)$, and $\theta \equiv (\theta_1, \theta_2, \theta_3)$ with θ_i real constants—taken as small for weak noise—that parameterize the unitary error. This can arise from, for example, a systematic calibration error in the pulse control leading to a consistent over or under rotation. ~~One could also include in this model a random over/under-rotation by drawing θ from a specified distribution that describes the randomness in the control. Here, we restrict to the systematic error case and note only that our analysis applies to this random error situation simply by doing an average over the distribution of θ . For simplicity, we assume gate-independent noise, so the same θ applies to all pulses. (Random unitary rotation does not fit into the gate-independent noise model. So it is not suitable for discussion here.)~~ In this situation, $\Gamma \equiv \theta \cdot \sigma$ acts only on the system and does not have a pure bath term. In the appendix, we derive an tighter bound:

$$\begin{cases} \tilde{\phi}_B \leq \phi_B + \frac{1}{3}\theta^2\phi_{\text{SB}}^2, \\ \tilde{\phi}_{\text{SB}} \leq \phi_{\text{SB}} + \theta^2\phi_{\text{SB}} \end{cases} \quad (38)$$

where $\theta \equiv \|\theta\|$ is the rotation angle associated with the gate error. Assuming the noise to be small, we may only use the leading order approximation, which is by itself accurate to the second order.

Where do the $\tilde{\alpha}$ and β' expressions come from? I see that the $\tilde{\alpha}$ one comes from setting $\Gamma_{\text{SA}} = \tau\tilde{\theta} \cdot \tilde{\sigma}$ and hence $\eta \equiv \tau\|\Gamma_{\text{SA}}\| \propto \theta$, and then plugging that into Eq. (69) of the earlier draft, with $\eta_0 = 0$. But why is the proportionality constant 1? What is the norm used here? Same question for β' : Where does the coefficient for the $\theta^2\beta$ come from? I'm understanding that we don't have a linear-in- θ term because of the rotation argument

above, but that does not explain the coefficient of 1 for the quadratic term. Is it just a matter of algebra using the BCH formula for the next-order terms? Also, how is θ related to the vector $\vec{\theta}$? What norm is being used for this? (I'm moving all these details to the appendix. For now you can just use the conclusion. The following paragraph should briefly describe the comparison between theoretical and numerical results. —Jiaan)

2. Gate-dependent noise

For the more generic case of gate-dependent noise, we can no longer reuse the ideal DD transformation formula $\Omega \rightarrow \Omega_{\text{DD}}$ as there is no well defined Ω . But we can still rewrite Eq. (32) to allow derivation of a converging series for Ω_{DD} :

$$e^{-i\tilde{\Omega}_{\text{DD}}} = e^{-iP_L\Gamma_L P_L^\dagger} e^{-i\Omega_L} \dots e^{-i\Omega_2} e^{-i\Gamma_1} e^{-i\Omega_1}. \quad (39)$$

As all the matrix exponents are small, we may use the Magnus formula to derive the leading order Magnus series as good approximation to $\tilde{\Omega}_{\text{DD}}$.

For PDD, the sequence comprises only X and Z gates. This amounts to writing the noisy X and Z pulses as

$$\tilde{X} = X e^{-i\Gamma'} \quad \text{and} \quad \tilde{Z} = Z e^{-i\Gamma''}. \quad (40)$$

Here, Γ' and Γ'' can generally act on the system and bath (including possibly an ancillary system). We define

$$\Gamma' \equiv \sum_i \sigma_i \otimes B'_i \quad \text{and} \quad \Gamma'' \equiv \sum_i \sigma_i \otimes B''_i. \quad (41)$$

The time evolution over the noisy PDD sequence then becomes

$$e^{-i\tilde{\Omega}_{\text{PDD}}} = e^{-iZ\Gamma''Z} e^{-i\Omega_3} e^{-iY\Gamma'Y} e^{-i\Omega_2} e^{-iX\Gamma''X} e^{-i\Omega_1} e^{-i\Gamma'} e^{-i\Omega}. \quad (42)$$

Every term on the exponent is small and we can apply the Magnus formula to evaluate $\tilde{\Omega}_{\text{PDD}}$. The first order Magnus term reads:

$$\begin{aligned} \tilde{\Omega}_{\text{PDD}}^{(1)} &= \sigma_0 \otimes (4B_0 + 2B'_0 + 2B''_0) \\ &\quad + \sigma_2 \otimes (2B'_2 - 2B''_2). \end{aligned} \quad (43)$$

The term that acts nontrivially on the system in $\tilde{\Omega}_{\text{PDD}}^{(1)}$ is proportional to $B'_2 - B''_2$. When the noise is gate-independent, i.e., when $\Gamma' = \Gamma''$, the expression reduce back to the ideal PDD situation, consistent with our earlier discussion for the gate-independent imperfections. With gate-dependent noise, however, exact first order decoupling is lost, as $B'_2 - B''_2 \neq 0$. This has significant implication for the concatenate DD scheme as well, as the first order coupling would persist throughout all concatenation levels.

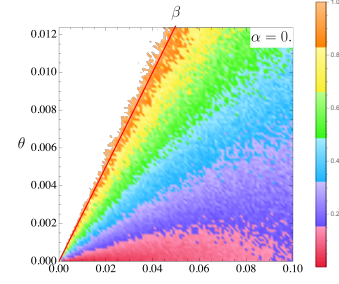


FIG. 4. (Which one is for PDD? Why do some points violate the threshold condition? Is it because of numerical—sampling?—inaccuracies?) The noise threshold compared for PDD and CDD₃ simulated at $\alpha = 0$. The red solid line represents the noise threshold prediction $\eta \leq \beta/4$. different colors represents different levels of noise reduction ration β_{DD}/β after numerically maximized over random bath operator instances.

Analogous to the gate-independent noise situation, we take $\eta \equiv \max\{\|\Gamma'_{\text{SB}}\|, \|\Gamma''_{\text{SB}}\|\}$. Assuming no particular relation between the X and Z noise, we can upper-bound the size of the first order coupling by:

$$\|\tilde{\Omega}_{\text{PDD:SB}}^{(1)}\| \leq 4\eta. \quad (44)$$

When the gate noise is at least of similar strength as the noise from H_{SB} —the gate noise is in fact stronger than the background noise in many experiments— $\tilde{\Omega}_{\text{PDD}}$ is dominated by the first order term. So that the error phase for PDD can be estimated as $\Phi_{\text{SB}} \simeq \|\tilde{\Omega}_{\text{PDD:SB}}^{(1)}\| \leq 4\eta$. With this estimate, the condition for the break-even point for PDD under such gate-dependent noise can be written as

$$4\eta \leq \phi_{\text{SB}}. \quad (45)$$

(Please provide more details for the figure, on how the numerics are done. For example, what random bath operators are we referring to? The B 's and B ''s? What do you mean by “random” here? Uniform distribution over what? Initial system and bath states? etc.) In Figure 4, we simulate both PDD and CDD₃ with the gate-dependent noise model. We find that our prediction works pretty well.

3. Gate-independence version 2

The decomposition $\tilde{P}_i = P_i e^{-i\Gamma'_i}$ previously considered for the noisy gate is by construction and not fundamental. Another possible decomposition is to include both left-error and right-error maps, as opposed to only one-sided error map. That is,

$$\tilde{P}_i = e^{-iL_i} P_i e^{-iR_i}, \quad (46)$$

where L_i and R_i are gate-dependent Hermitian operators responsible for generating the left and right error maps.

We expect both operators to be small as the noise is weak. Using this decomposition, we get another definition for gate-independent noise. We say that the noise associated with the gate is gate-independent if there exists Hermitian L and R operators such that $L_i = L$ and $R_i = R$ for all noisy pulses \tilde{P}_i . This definition automatically covers the previous one by setting $L_i = 0$. After plugging the gate-independent noise model into the DD sequence, one finds that

$$e^{-i\tilde{\Omega}_{\text{DD}}} = e^{-iL} \left(P_L e^{-i\tilde{\Omega}} \dots P_1 e^{-i\tilde{\Omega}} \right) e^{iL}, \quad (47)$$

where we have defined the effective free-evolution Hamiltonian $\tilde{\Omega}$ by

$$e^{-i\tilde{\Omega}} = e^{-iR} e^{-i\Omega} e^{-iL}. \quad (48)$$

In the bracket of Eq. (47), the sequence once again reduces to that of ideal DD gates. Therefore, the decoupling order will be unaffected within the bracket. The edge terms $e^{\pm iL}$ in Eq. (47) seems to be problematic, as they will break the exact decoupling enjoyed by previously considered gate-independent noise. In fact, one need not be worried about the edge terms if DD is periodically performed in a long run. This is commonly assumed for both quantum memory and quantum computation tasks. Within the typical decoherence time for a NISQ era qubit, thousands of gates, if not more, can be performed. Naturally, a DD-protected task will involve in way more than one DD cycle. After joining these cycles, the edge terms will travel to the edges of the whole sequence, where they are usually overwhelmed by the state preparation and measurement errors. We do acknowledge these potential slight overheads, but the scaling benefit provided by DD are more meaningful. Hence it is safe to ignore these edge terms.

A gate-independent noise with both left and right errors may be viewed as gate-dependent in if only right noise is considered. This difference is only a superficial matter on the definitions. But how about the decoupling behavior? From previous discussion we know that gate-dependent noise will result in inexact decoupling, but gate-independent noise will have exact decoupling similar to the ideal case. To investigate this contraction, we consider PDD as an example. After equating both ways of decomposing the noisy gate, we have:

$$\begin{aligned} e^{-iL} X e^{-iR} &= X e^{-i\Gamma'}, \\ e^{-iL} Z e^{-iR} &= Z e^{-i\Gamma''}. \end{aligned} \quad (49)$$

For given L and R , we can solve for Γ' and Γ'' by basic algebra and the BCH formula. But the reverse is generally unsolvable. Since both L and R are small, we may truncate to the second order expansion:

$$\begin{aligned} \Gamma' &= XLX + R + \frac{1}{2i}[XLX, R] + \dots, \\ \Gamma'' &= ZLZ + R + \frac{1}{2i}[ZLZ, R] + \dots. \end{aligned} \quad (50)$$

According to Eq. (43), the DD effective Hamiltonian has a non-vanishing first order coupling, which is proportional to $B'_2 - B''_2$. Let us introduce a smallness parameter η by shifting $L \rightarrow \eta L$ and $R \rightarrow \eta R$. We can use the formula for Γ' and Γ'' to show that

$$B'_2 - B''_2 = \mathcal{O}(\eta^2). \quad (51)$$

Hence one can still argue that the first order decoupling is exact despite it is superficially not so. Exact decoupling or not is only a matter of perspective. We say PDD achieves first order decoupling because the first order smallness in the effective Hamiltonian is removed. This settles the decoupling order disagreement between the two views of gate-independence.

How useful is the gate-independent noise model anyway? One may scorn it off as being too contrived and does not bear any practical relevance. To counter this doubt we consider the finite-width pulse noise model from a new perspective. The noisy gates are approximately:

$$e^{-i(\frac{\pi}{2}\sigma_\alpha + \delta H)} \approx e^{-i\frac{1}{2}\delta H} e^{-i\frac{\pi}{2}\sigma_\alpha} e^{-i\frac{1}{2}\delta H}, \quad \text{for } \alpha = X, Z. \quad (52)$$

This is an example of the Trotter-Suzuki type of approximation. And our approximation is valid to the second order as one may check. For PDD and CDD, all we will be concerned is the first two orders of Magnus term. Hence it suffices to use this approximation.

We now make the above observation rigorous.

Proposition 1: For PDD if the Hermitian generator responsible for the noise is gate-independent. We have exact first order decoupling.

Proposition 2: For CDDn if the Hermitian generator responsible for the noise is gate-independent. We have exact n-th order decoupling.

4. Unitary control noise with finite-width pulses

(To add)

IV. SCALING UP PROTECTION - CONCATENATION AND THE ACCURACY THRESHOLD

(Generalization of the previous PDD section; CDD content from Secs. IVA-D of draft by Jiaan & Xiansong. The main goal here is to show that an accuracy threshold for DD, under the CDD scheme for scaling up, fails to exist. This means that there is a maximum level for CDD, beyond which further levels only make things worse. To discuss the differences from using different measures.) We can distinguish between two operational considerations: to preserve states in a quantum memory (no computational gates), or to reduce noise in the course of carrying out a quantum computation. In the memory setting, the goal is to preserve the quantum informational state for some time T without damage. During this time T , we can do DD pulses, and ask if the noise at the end of each complete DD sequence is lower than the noise over the same time period if there were no DD pulses. In other words, for DD to work well, we require

$$\text{Memory setting: } \epsilon_{\text{DD}} < \epsilon(L), \quad (53)$$

where ϵ_{DD} quantifies the noise after the DD sequence, while $\epsilon(L)$ is the noise after bare evolution for the time taken for the L DD pulses to be applied. In the computational setting, the comparison of with and without DD is different. With DD, computational gates can only be applied at the completion of each iteration of the DD sequence, as opposed to every time step without DD. The condition for DD to be effective in this case is then

$$\text{Computational setting: } \epsilon_{\text{DD}} < \epsilon(1), \quad (54)$$

where $\epsilon(1)$ is the noise after bare evolution for a single time step.

These two perspectives result in very different overall decoupling strategies. Especially if one is using CDD to “scale up protection”. As illustrated in Fig. 5, for the quantum memory setting, as total time is fixed T , one only need to come up with finer and finer slicing of the time evolution pieces; but for the quantum computation setting, one need to join the DD control gates to form a dilated time evolution map. The past papers in favor of CDD had taken the former perspective. It is undeniable that with finer and finer slicing of time evolution, one gets better and better performance. But this picture has some series drawbacks. First the total evolution time T is typically unknown, even for quantum memory. Hence there is no way of knowing in advance which gate to perform at what time. Second, there exists a minimal switch/cool-down time between consecutive gates due to technological limitations. Hence it is impractical to arbitrarily suppress noise by increasing concatenation level.

In our work, we focus on the computational setting, the more stringent one among the two, though it is straightforward to modify our analysis for the memory case. As

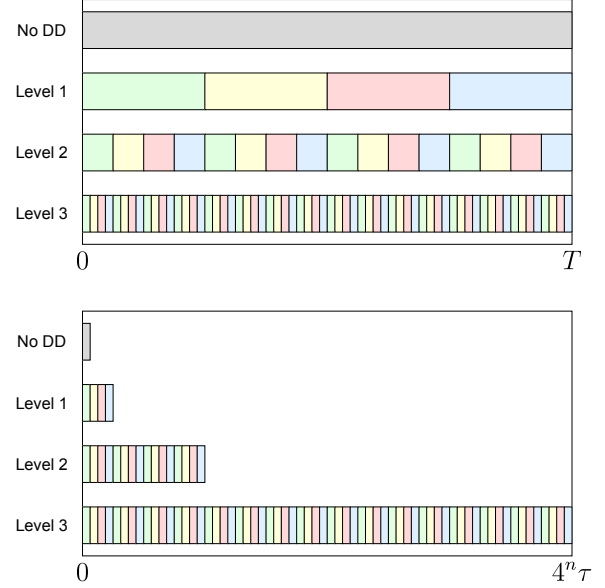


FIG. 5. An illustration of the CDD time evolution maps under increasing concatenation levels. We compare the quantum memory setting (above) to the quantum computation setting (under).

we will see, Condition (54) will yield a requirement on the noise parameters characterizing the noise in the system and the DD pulses. We refer to that requirement as the “break-even point” for DD, borrowing terminology from quantum error correction and fault-tolerant quantum computing. (There was a comment in the .tex file that some past papers have analyzed the memory case. What are those papers? Do you just mean the Khodjasteh-Lidar 2007 paper? That paper looks only at imperfections due to finite-width pulses, not control errors as we do here. Their paper hence only has a single-parameter Φ_{SB} , not our two-parameter Φ_{SB} and η situation. In fact, in current experiments η is likely the larger problem; memory noise is usually much smaller than gate noise. Or are there other earlier papers?) (Yes, that paper in particular. Their conclusion that CDD decouples to order n and dramatically outperforms PDD come from this perspective. Our paper is focused on the quantum computation case which is more practical. See my new added picture 5). And let me quote their words (bottom left page 10) on how they do the comparison: “The higher order data for PDD are obtained by simply repeating the basic sequence and shrinking the pulse interval. This is necessary to get the long-term behavior of decoupling. ...” I move this memory/computation argument from the previous section as it is more appropriate when we discuss the scaling behavior. —Jiaan)

(To be merged with the new text.) The conditions for achieving noise suppression are already derived previously for the idealistic PDD and CDD case. In general,

we find that smaller (α, β) values would lead to better noise suppression rate. Therefore the gates should be performed as frequently as possible. But this may lead to suboptimal outcomes in practice given the extra noise introduced by imperfect gate operations. To study the fault-tolerance of DD, we must account for the additional noise associated with the control gates.

Next, we turn to the situation of CDD. We can certainly discuss the break-even point for a particular CDD_n sequence, treating it simply as a “flattened” sequence of pulses, and ask when the error phase after CDD_n is smaller than ϕ_{SB} as was done for PDD in the previous section. However, what is more interesting is the question of the so-called accuracy threshold for CDD. As described in the Preliminaries section, CDD tries to achieve better and better noise removal by concatenating the same basic DD scheme—PDD for our discussion here—to higher and higher levels. Here, we are interested in understanding how the noise removal improves as the DD scheme scales up. In fault-tolerant quantum computing, there is the concept of a fault tolerance noise threshold, a strength of the noise below which scaling up the QEC code leads to improved protection against noise, and hence more accurate quantum computation. This threshold is referred to as the accuracy threshold. Here, we ask the analogous question of CDD, whether there is a condition on the noise parameter of the problem such that increasing the CDD level always leads to improved noise removal capabilities. An accuracy threshold for CDD exists if the parameters governing the noise of the situation satisfying a condition $\eta < \eta_0$, such that

$$\Phi_{\text{SB},n+1} < \Phi_{\text{SB},n} \quad \forall n = 1, 2, \dots \quad (55)$$

Below, we examine this for both the ideal CDD case where pulses are perfect, and the noisy CDD case where imperfections in the pulses are allowed.

A. Ideal case

If a noise threshold indeed exist, then ideal case, where the noise is non-existent, would definitely apply. We first consider CDD with perfect pulses. Khodjasteh and Lidar propose in their paper to truncate to the second order term gave in their paper an estimation of the CDD_n generator as:

$$\begin{aligned} \Omega_{(n)} &\approx \sigma_0 \otimes 4^n B_0 \\ &+ \sigma_1 \otimes (-i)^n 2^{n(n+1)} \text{ad}_{B_0}^n(B_1) \\ &+ \sigma_2 \otimes (-i)^n 2^{n^2} \text{ad}_{B_0}^{n-1}([B_0, B_2] - i\{B_1, B_3\}). \end{aligned} \quad (56)$$

It is shown in the appendix A that the above expression indeed captures the leading order behavior of $\Omega_{(n)}$. Different from the leading order term in the Magnus series, the leading order “behavior” is manifested separately for the pure bath part and the interaction part. As mentioned in the Preliminaries section, CDD_n achieves n th-order decoupling. From the leading order expression, we see it is indeed so as the interaction part is of $(n+1)$ th order. From Eq. (56), we can estimate $\Phi_{\text{B},n} \simeq 4^n \phi_{\text{B}}$ and derive an upper bound for the error phase:

$$\Phi_{\text{SB},n} \lesssim 2^{(n+1)^2} \phi_{\text{B}}^n \phi_{\text{SB}} f_n\left(\frac{\phi_{\text{SB}}}{\phi_{\text{B}}}\right), \quad (57)$$

where we define the piece-wise function:

$$f_n(x) = \begin{cases} 1, & x \leq \sqrt{4^n - 1}, \\ \frac{1}{2^n} \sqrt{\left(\frac{4^n - 1}{2x} + \frac{x}{2}\right)^2 + 1}, & x > \sqrt{4^n - 1}. \end{cases} \quad (58)$$

to remind the reader, $\phi_{\text{B}} \equiv \phi_{\text{B}}(\tau_0) \equiv \tau_0 \|H_{\text{B}}\|$, and $\phi_{\text{SB}} \equiv \phi_{\text{SB}}(\tau_0) \equiv \tau_0 \|H_{\text{SB}}\|$, noting the dependence of both quantities on τ_0 , the time interval between two consecutive pulses in the CDD sequence. (To reconsider: We need both an upper and lower bound for $\Phi_{\text{SB},n}$, which we have from the estimator [Eq. (A5)]. For use in Eq. (55), we should have the lower bound on the RHS of that condition, and upper bound expression on the LHS, for a sufficient condition.) As mentioned in the Preliminaries section, CDD_n achieves n th-order decoupling. This can be understood from the error phase expression: $\Phi_{\text{SB},n}$ after CDD_n is of order ϕ^{n+1} , for $\phi \sim \phi_{\text{B}}, \phi_{\text{SB}}$. That ϕ_{B} enters the error phase should again be of no surprise—as in the PDD case, ϕ_{B} determines how quickly the noise seen by the system evolves, and hence affects the efficacy of DD which does a good elimination of the noise only when the noise remains nearly unchanged for the full DD sequence.

What is perhaps more surprising is the appearance of factors of 2^n and 2^{n^2} in $\Phi_{\text{SB},n}$, a new feature in the CDD case (Do these factors also appear in the Khodjasteh-Lidar analysis? Please check and put in a sentence about this.). For fixed $\tau_0 \equiv \tau$ as n increases (e.g., if τ is the experimental limit for the switch time between consecutive

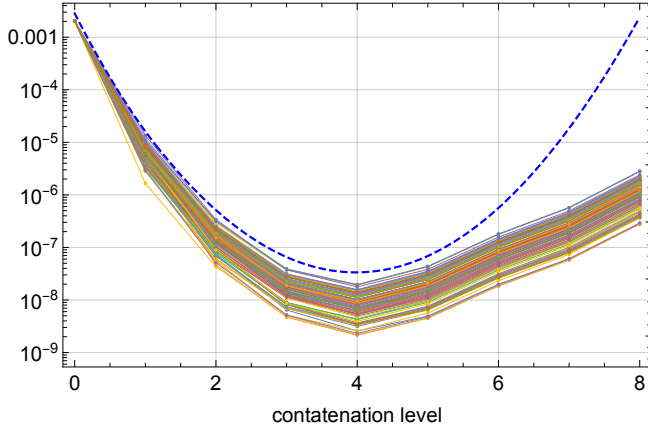


FIG. 6. The error phase $\phi_{SB,n}$ as a function of concatenation level n . The thin color lines are obtained from numerically simulating CDD $_n$ with randomly generated bare Hamiltonian operators satisfying $\phi_B = 1 \times 10^{-3}$ and $\phi_{SB} = 2 \times 10^{-3}$. The blue dashed line is our theoretical upper bound. For this configuration, we set the maximal concatenation level at 4.

pulses), such factors mean that the error phase eventually increases for large enough n : The exponentially decreasing ϕ^n factor is eventually overcome by the super-exponentially increasing 2^{n^2} factor. This means there is no accuracy threshold, i.e., there is no level of noise that is weak enough such that $\Phi_{SB,n+1} \leq \Phi_{SB,n}$ for all n . Instead, there is an maximal useful level of concatenation level, beyond which further concatenation actually increases the noise seen by the system. Fig. 6 plots this situation of fixed τ_0 , for increasing concatenation level n . We observe the initial decrease of $\Phi_{SB,n}$ as n increases, but this turns around eventually.

(To reconsider with my lower/upper bound comment above.) For higher concatenation level to make sense, increasing concatenation level should decrease the error phase. To derive a sufficient condition of it, we examine the recursive relation between concatenation levels, which suggests

$$\begin{aligned} B_{n+1,0} &= 4B_{n,0} \\ B_{n+1,1} &= -4i[B_{n,0}, B_{n,1}] \\ B_{n+1,2} &= -2i[B_{n,0}, B_{n,1}] \\ B_{n+1,3} &= 0 \end{aligned} \quad (59)$$

For both B_1 and B_2 to decrease in size, it suffice to require the map $-4i[B_{n,0}, \cdot]$ a contraction map. By requiring its norm to be smaller than 1, we obtain the bound for the maximal concatenation level:

$$n_{\max} \leq \lceil -\log_4(\phi_B) - 3/2 \rceil, \quad (60)$$

where $\lceil \cdot \rceil$ is the ceiling operator.

Note that the existence of a maximal useful concatenation level does not technically contradict the statement that higher-order decoupling is achieved by higher level CDD. Using decoupling order to quantify the degree of

noise removal requires, in the first place, the convergence of the Magnus series, so that the $(n+1)$ th- and higher-order terms in the series are a small correction to the first n terms. However, if the DD scheme is designed such that the total time for the sequence grows exponentially with n , as is the case for CDD if τ_0 is fixed as n increases, eventually, we exceed the convergence criterion and the decoupling order stops being a reasonable indicator of successful noise removal.

We can, however, analyze a different situation: to have fixed $\tau_n \equiv T$, so that the CDD sequence takes the same amount of time, regardless of n . This requires $\tau_0 = \frac{T}{4^n}$ for each n , so that the pulses are applied at shorter and shorter time intervals as the concatenation level increases. This can be the practical choice if one is doing computation where computational gates, which can only be applied at the end of a complete DD sequence in order to not interfere with the noise averaging process, have to be applied at a particular clock rate. Increasing n to increase the noise removal capabilities must not increase the total time taken for the DD sequence. This requires, of course, the ability to do faster and faster pulse switching.

In this case, it is more illuminating to make use of $\tau_0 \|H_B\|$ and $\tau_0 \|H_{SB}\|$, in place of ϕ_B and ϕ_{SB} , in our estimate of the CDD error phases, so that that the n dependence in τ_0 is explicit. (I suggest we analyze this case as well, to see if an accuracy threshold exists. Ji-aan - could you take a quick look at this?) Eventually, of course, this becomes limited by the technological—not fundamental—constraint of how fast we can switch between pulses, and thus how short τ_0 can be in an experiment.

B. Noisy pulses

(To add.)
(To add.)

V. CONCLUSION

(To come.)

ACKNOWLEDGMENTS

(Acknowledge Gu Yanwu, Jonas Tan, and Ryan Tiew, who did some of the first calculations of the breakeven points for PDD. Yanwu should be asked if he wants to be a co-author.)

Appendix A: Analyzing CDD $_n$

(To be edited.) It was suggested by the authors of CDD that the decoupling order for CDD $_n$ is n [4]. This

statement is correct, but their proof is logically flawed. Given its importance, here we explain why the original proof is insufficient and provide a rigorous proof for the decoupling order.

For the generic CDD_n, it is infeasible to flatten out the gates and calculate the combined Magnus series directly. Instead it is easier to study the iteration map

$$\Omega_{(n)} = \Omega_{\text{PDD}}(\Omega_{(n-1)}) = \sum_{m=1}^{\infty} \Omega_{\text{PDD}}^{(m)}(\Omega_{(n-1)}), \quad (\text{A1})$$

where the update map Ω_{PDD} is defined in Eq. (??). Backtracking the map from level n to level 0, we have $\Omega_{(n)} = \Omega_{\text{PDD}} \circ \dots \circ \Omega_{\text{PDD}}(\Omega) \equiv (\Omega_{\text{PDD}})^n(\Omega)$. In the end, the effective Hamiltonian $\Omega_{(n)}$ can be expanded as:

$$\Omega_{(n)} = \sum_{m=1}^{\infty} \Omega_{(n)}^{(m)}(\Omega), \quad (\text{A2})$$

where $\Omega_{(n)}^{(m)}$, with the size of $\mathcal{O}(\|\Omega\|^m)$, is the m th order Magnus term in Ω . Notice that this series is different from the series in Eq. (A1), where each $\Omega_{\text{PDD}}^{(m)}$ generate the m th order Magnus term in $\Omega_{(k-1)}$.

To calculate the leading order Magnus term in Ω , we do not directly solve $\Omega_{(n)}$ but instead constructing an estimator $\hat{\Omega}_{(n)}$ which is easier to solve. Formally we have an estimator-error pair $(\hat{\Omega}_{(n)}, \delta\Omega_{(n)})$ such that

$$\Omega_{(n)} = \hat{\Omega}_{(n)} + \delta\Omega_{(n)}. \quad (\text{A3})$$

The estimator $\hat{\Omega}_{(n)}$ is faithful if it captures the leading order behavior of $\Omega_{(n)}$. This is defined separately for the pure-bath part and the coupling part, as they are of different orders. In other words, a faithful estimator should contain the first nonzero pure-bath part and the first nonzero coupling part of the series (A2). An equivalent statement is that the error must be of higher order than the estimator. To quantify this statement, we take the two-component norm for both $\delta\Omega_{(n)}$ and $\hat{\Omega}_{(n)}$, and demanding

$$\|\delta\Omega_{(n)}\| \equiv \begin{pmatrix} \delta\alpha_{(n)} \\ \delta\beta_{(n)} \end{pmatrix} \ll \|\hat{\Omega}_{(n)}\| \equiv \begin{pmatrix} \hat{\alpha}_{(n)} \\ \hat{\beta}_{(n)} \end{pmatrix}, \quad (\text{A4})$$

where the comparison is implied for both components and should be understood by comparing the leading powers of the polynomials in α and β . In essence, a “smaller” polynomial should have a larger leading power and thus is of higher order smallness. We refer readers to Appendix ?? for more rigorous descriptions on ordering polynomials by their leading powers. Furthermore, after taking the two-component norm and applying triangular inequality for Eq. (A3), we find,

$$\hat{\beta}_{(n)} - \delta\beta_{(n)} \leq \beta_{(n)} \leq \hat{\beta}_{(n)} + \delta\beta_{(n)}. \quad (\text{A5})$$

According to Eq. (6), the n th-order decoupling condition requires $\beta_{(n)}$ to be an $(n+1)$ th order polynomial in α

and β . It now suffices to show that ① $\hat{\beta}_{(n)}$ is an $(n+1)$ th order polynomial and ② the estimator is faithful.

The original CDD paper (Ref.[4]) constructed an estimator by keeping the first two Magnus terms for each iteration step. In our notation:

$$\hat{\Omega}_{(n)} \equiv (\Omega_{\text{PDD}}^{(1)} + \Omega_{\text{PDD}}^{(2)})(\hat{\Omega}_{(n-1)}) = (\Omega_{\text{PDD}}^{(1)} + \Omega_{\text{PDD}}^{(2)})^n(\Omega). \quad (\text{A6})$$

With $\Omega_{\text{PDD}}^{(1)}, \Omega_{\text{PDD}}^{(2)}$ explicitly defined in Eq. (??) and Eq. (14), and the starting point $\hat{\Omega}_{(0)} \equiv \Omega$, the estimator can shown to be specified by the formula:

$$\begin{aligned} \hat{\Omega}_{(n)} &= \sigma_0 \otimes 4^n B_0 \\ &+ \sigma_1 \otimes (-i)^n 2^{n(n+1)} \text{ad}_{B_0}^n(B_1) \\ &+ \sigma_2 \otimes (-i)^n 2^{n^2} \text{ad}_{B_0}^{n-1}([B_0, B_2] - i\{B_1, B_3\}). \end{aligned} \quad (\text{A7})$$

After taking the two-component norm of the estimator, we have,

$$\begin{pmatrix} \hat{\alpha}_{(n)} \\ \hat{\beta}_{(n)} \end{pmatrix} \simeq \begin{pmatrix} \alpha \\ \alpha^n \beta + \alpha^{n-1} \beta^2 \end{pmatrix}, \quad (\text{A8})$$

where \simeq represents equal the leading powers as described in Appendix ???. This verifies that $\hat{\beta}_{(n)}$ is a an $(n+1)$ th order polynomial.

To complete the theory, it now remains to show that $\hat{\Omega}_{(n)}$ is indeed faithful. For simplicity, we make the recognition $\alpha \simeq \beta \simeq \tau_0$. This change does not affect the criteria for n th order decoupling as the polynomial order is unchanged. To show $\delta\alpha_{(n)}$ and $\delta\beta_{(n)}$ are of higher order in τ_0 compared to $\hat{\alpha}_{(n)}$ and $\hat{\beta}_{(n)}$, we claim that:

$$\begin{pmatrix} \delta\alpha_{(n)} \\ \delta\beta_{(n)} \end{pmatrix} \lesssim \begin{pmatrix} \tau_0^3 \\ \tau_0^{n+2} \end{pmatrix}. \quad (\text{A9})$$

We prove this assertion by mathematical induction. At $n = 1$, i.e. the PDD sequence, the error comes solely from the Magnus series truncation. Taking the two-component norm and applying triangular inequality:

$$\|\delta\Omega_{(1)}\| \leq \sum_{m=3}^{\infty} \|\Omega_{(1)}^{(m)}\| \simeq \|\Omega_{(1)}^{(3)}\| \lesssim \begin{pmatrix} \tau_0^3 \\ \tau_0^3 \end{pmatrix}. \quad (\text{A10})$$

By definition, the higher level error term $\delta\Omega_{(n+1)}$ can be expressed with the lower level functions as:

$$\begin{aligned} \delta\Omega_{(n+1)} &= \Omega_{\text{PDD}}(\Omega_{(n)}) - (\Omega_{\text{PDD}}^{(1)} + \Omega_{\text{PDD}}^{(2)})(\hat{\Omega}_{(n)}) \\ &= \Omega_{\text{PDD}}^{(1)}(\delta\Omega_{(n)}) + \sum_{m=2}^{\infty} \Omega_{\text{PDD}}^{(m)}(\Omega_{(n)}) - \Omega_{\text{PDD}}^{(2)}(\hat{\Omega}_{(n)}), \end{aligned} \quad (\text{A11})$$

where we have used the fact that $\Omega_{\text{PDD}}^{(1)}$ is a linear map. Assuming Eq. (A9) holds for level n , we need to show $\|\delta\Omega_{(n+1)}\| \lesssim (\tau_0^3, \tau_0^{n+3})$. The argument in the original CDD paper by Khodjasteh and Lidar for truncating the

Magnus series have only showed that the higher order Magnus term is smaller by proving

$$\left\| \sum_{m=3}^{\infty} \Omega_{\text{PDD}}^{(m)}(\hat{\Omega}_{(n)}) \right\| \ll \left\| \hat{\Omega}_{(n+1)} \right\|.$$

If we take $\delta\Omega_{(n)} = 0$, Eq. (A11) would identity the l.h.s. of above as $\left\| \delta\Omega_{n+1} \right\|$, then this inequality would indeed indicate that the error term is of higher order. However Eq. (A11) suggests that this argument is not sufficient: as the estimate error $\delta\Omega_{(n+1)}$ comes not only from the truncation error of the same level, but also accumulates from the lower level error $\delta\Omega_{(n)}$. To properly bound the size of the error term, we take the two-component norm on both sides of Eq. (A11):

$$\begin{aligned} \left\| \delta\Omega_{(n+1)} \right\| &\lesssim \left\| \Omega_{\text{PDD}}^{(1)}(\delta\Omega_{(n)}) \right\| + \left\| \Omega_{\text{PDD}}^{(3)}(\Omega_{(n)}) \right\| \\ &\quad + \left\| \Omega_{\text{PDD}}^{(2)}(\hat{\Omega}_{(n)} + \delta\Omega_{(n)}) - \Omega_{\text{PDD}}^{(2)}(\hat{\Omega}_{(n)}) \right\|, \end{aligned} \quad (\text{A12})$$

where the infinite sum of Magnus terms higher than third order is lead by the non-vanishing third order. We now examine the size of these terms. The first term in (A12) is easy,

$$\left\| \Omega_{\text{PDD}}^{(1)}(\delta\Omega_{(n)}) \right\| = \begin{pmatrix} 4\delta\alpha_{(n)} \\ 0 \end{pmatrix} \simeq \begin{pmatrix} \tau_0^3 \\ 0 \end{pmatrix}.$$

To bound the size of the second term, we make two observations for the PDD Magnus series greater than the third order:

1. At least one $B_{i \neq 0}$ is involved in the coupling part.
2. At least two $B_{i \neq 0}$ is involved in pure bath part.

Therefore, the second term can be bounded by

$$\left\| \Omega_{\text{PDD}}^{(3)}(\Omega_{(n)}) \right\| \lesssim \begin{pmatrix} \beta_{(n)}^2(\alpha_{(n)} + \beta_{(n)}) \\ \beta_{(n)}(\alpha_{(n)} + \beta_{(n)})^2 \end{pmatrix} \simeq \begin{pmatrix} \tau_0^{2n+2} \\ \tau_0^{n+3} \end{pmatrix},$$

where we have used $\alpha_{(n)} \lesssim \hat{\alpha}_{(n)}$ and $\beta_{(n)} \lesssim \hat{\beta}_{(n)}$ according to the induction hypothesis. This can be alternatively shown by explicitly calculating the third order Magnus term using Eq. (??). To bound the third term in (A12), we can calculate the difference $\Omega_{\text{PDD}}^{(2)}(\hat{\Omega} + \delta\Omega) - \Omega_{\text{PDD}}^{(2)}(\hat{\Omega})$ using the expression for $\Omega_{\text{PDD}}^{(2)}$ from Eq. (14). In the end,

$$\begin{aligned} \left\| \Omega_{\text{PDD}}^{(2)}(\hat{\Omega} + \delta\Omega) - \Omega_{\text{PDD}}^{(2)}(\hat{\Omega}) \right\| &\lesssim \begin{pmatrix} 0 \\ \delta\alpha_{(n)}\delta\beta_{(n)} \end{pmatrix} \\ &\quad + \begin{pmatrix} 0 \\ \hat{\beta}_{(n)}\delta\alpha_{(n)} + \hat{\alpha}_{(n)}\delta\beta_{(n)} + \hat{\beta}_{(n)}\delta\beta_{(n)} \end{pmatrix} \simeq \begin{pmatrix} 0 \\ \tau_0^{n+3} \end{pmatrix}. \end{aligned}$$

Since all three terms are bounded by $(\tau_0^3, \tau_0^{n+3})^\top$, so is their sum. This proves our assertion Eq. (A9). \square

Having showed its faithfulness, we are now confident in using the estimator Eq. (A2) to study the full $\Omega_{(n)}$ of CDD_n. In particular, we have $\alpha_{(n)} \approx \hat{\alpha}_{(n)} = 4^n\alpha$ and $\beta_{(n)} \approx \hat{\beta}_{(n)}$. Similar as the PDD case, we can derive a strict bound and an approximate bound. For the strict bound, we have

$$\begin{aligned} \hat{\beta}_{(n)} &\leq 2^{n(n+1)} \|\text{ad}_{B_0}^n(B_1)\| \\ &\quad + 2^{n^2} \|\text{ad}_{B_0}^{n-1}([B_0, B_2] - i\{B_1, B_3\})\| \quad (\text{A13}) \\ &\leq 2^{n(n+1)} \alpha^{n-1} \beta (2^n \alpha + \alpha + \beta). \end{aligned}$$

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