

## Lecture 10

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Today, we will continue our approach to factoring bivariate polynomials. We will first focus on the tool of Hensel's Lifting; and then describe how we perform the factoring.

## 1 Hensel's Lifting

Suppose  $f(x, y) = g(x)h(x) \pmod{y}$ . We wanted a factorization for higher powers of  $y$ :  $f(x, y) = \tilde{g}(x, y)\tilde{h}(x, y) \pmod{y^{2k}}$ . Hensel's Lifting says we can obtain this if  $g$  and  $h$  are “relatively prime,” and that the  $\tilde{g}$  and  $\tilde{h}$  we obtain are essentially unique. Formally:

**Lemma 1 (Hensel's Lifting)** *If  $R$  is a ring,  $I \subseteq R$  is an ideal, and there exist  $f, g, h, a, b$  in  $R$  such that*

$$(H1) \quad f = gh \pmod{I}$$

$$(H2) \quad ag + bh = 1 \pmod{I}.$$

*Then, for every positive integer  $s$  that is a power of 2, there exist  $\tilde{g}, \tilde{h}$  and  $\tilde{a}, \tilde{b}$  in  $R$ , such that*

$$(C1) \quad f = \tilde{g}\tilde{h} \pmod{I^s}$$

$$(C2) \quad \tilde{g}\tilde{a} + \tilde{b}\tilde{h} = 1 \pmod{I^s}$$

$$(C3) \quad g = \tilde{g} \pmod{I} \text{ and } h = \tilde{h} \pmod{I}$$

*Furthermore, the solution satisfying the above three conditions is “unique” in the following sense:*

**Uniqueness:** *We say that an ideal  $J$  is special if for all all integer  $k$ , and  $a, b$  such that  $ab \in J^k$ , there is an integer  $l$  such that  $a \in J^l$  and  $b \in J^{k-l}$ .*

*Assume that  $I$  is special. Then, for every two solutions  $g_1, h_1$  and  $g_2, h_2$  satisfying conditions C1 to C3, there exists  $u \in I^t$ , such that:*

$$g_2 = g_1(1+u) \pmod{I^s}, \quad h_2 = h_1(1-u) \pmod{I^s}$$

As we will see later, towards factoring a bivariate polynomial  $f \in \mathbb{F}_q[x, y]$ , we will apply the Hensel's lifting to  $R = F_q[x, y]$ ,  $I = (y)$  and a factorization of  $f \pmod{I}$ . When  $I = (y)$ ,  $f \pmod{y}$  is simply a univariate polynomial on  $x$  over  $\mathbb{F}_q$ , which we know how to factor from previous lectures. Next we proceed to prove Lemma 1.

**Proof** We prove this lemma by induction.

**Base Case  $s = 2$ :** We first show that there exists  $\tilde{g}, \tilde{h}$  and  $\tilde{a}, \tilde{b}$  satisfying C1 to C3, and then establish the uniqueness of the solution. By condition H1 and H2, we have

$$f = gh + q, \text{ for some } q \in I$$

$$ag + bh = 1 + r \text{ for some } r \in I$$

Write  $\tilde{g} = g + g_1$ ,  $\tilde{h} = h + h_1$ , for some  $g_1, h_1 \in I$  to be set. Then

$$\tilde{g}\tilde{h} = gh + g_1h + h_1g + h_1g_1$$

Since  $h_1g_1 \in I^2$ , in order to satisfy condition C1, we want  $g_1h + h_1g + h_1 = q \pmod{I^2}$ . To satisfy this, set  $g_1 = bq$ ,  $h_1 = aq$ , and get that  $g_1h + h_1g = q(bh + ag) = q(1+r)$ , which equals to  $q$  modulo  $I^2$  as required. By construction  $\tilde{g} = g \pmod{I}$  and  $\tilde{h} = h \pmod{I}$ , satisfying condition C3. To show that  $\tilde{g}$  and  $\tilde{h}$  are also relatively prime, observe that  $a\tilde{g} + b\tilde{h} = ag + bh + r' = 1 + r + r'$ , for some  $r' \in I$ . Let  $r'' = r + r' \in I$ . Now we can take  $\tilde{a} = a(1 - r'')$  and  $\tilde{b} = b(1 - r'')$ , and get that:

$$\tilde{a}\tilde{g} + \tilde{b}\tilde{h} = (1 - r'')(a\tilde{g} + b\tilde{h}) = (1 - r'')(1 + r'') = 1 - r''^2 = 1 \pmod{I^2}$$

Now it remains to show that  $\tilde{g}, \tilde{h}$  is the unique solution satisfying C1 to C3. That is, if  $g^*, h^*$  is a different solution satisfying C1 to C3, then there is  $u \in I$  such that  $g^* = \tilde{g}(1+u)$  and  $h^* = \tilde{h}(1-u)$ . By condition C3, we have  $g^* = \tilde{g} + g_2$  and  $h^* = \tilde{h} + h_2$  for some  $g_2, h_2 \in I$  (because, modulo  $I$ , we know that  $g^* = g = \tilde{g}$  and  $h^* = h = \tilde{h}$ ). Therefore, we have:

$$g^*h^* = \tilde{g}\tilde{h} + g_2\tilde{h} + h_2\tilde{g} + g_2h_2$$

By condition C1, we know that  $g^*h^* = f = \tilde{g}\tilde{h} \pmod{I^2}$ . Thus, the above equation modulo  $I^2$  gives,

$$g_2\tilde{h} + h_2\tilde{g} = 0 \pmod{I^2}$$

By condition C2, we have  $\tilde{a}\tilde{g} + \tilde{b}\tilde{h} = 1 \pmod{I^2}$ . Therefore,

$$\begin{aligned} \tilde{b}(g_2\tilde{h} + h_2\tilde{g}) &= 0 \pmod{I^2} \\ g_2\tilde{b}\tilde{h} + \tilde{b}h_2\tilde{g} &= 0 \pmod{I^2} \\ g_2(1 - \tilde{a}\tilde{g}) + \tilde{b}h_2\tilde{g} &= 0 \pmod{I^2} \\ g_2 &= (\tilde{a}g_2 - \tilde{b}h_2)\tilde{g} \pmod{I^2} \end{aligned}$$

Let  $u = \tilde{a}g_2 - \tilde{b}h_2$ . Since  $g_2$  and  $h_2$  are all elements in  $I$ , so is  $u$ . Furthermore we have  $g^* = \tilde{g} + g_2 = \tilde{g}(1+u)$ . Similarly, by symmetry, we obtain that

$$h_2 = (\tilde{b}h_2 - \tilde{a}g_2)\tilde{h} \pmod{I^2}$$

Therefore,  $h^* = \tilde{h}(1-u)$ . This concludes the proof for the base case.

**Induction Step:** Assume that for the case of  $s = t$ , there exist  $g_0, h_0 \in R[x]$  and  $a_0, b_0 \in I^t$  satisfying conditions C1 to C3, and the solution to the three conditions is “unique”. We show that for the case of  $s = 2t$ , we can construct  $g_1, h_1 \in R[x]$  and  $a_1, b_1 \in I^{2t}$  satisfying conditions C1 to C3, and the solution is also unique.

The existence of  $g_1, h_1, a_1, b_1$  satisfying conditions C1 to C3 follows exactly the same proof as in the base case. Therefore, we focus on the proof of uniqueness. Let  $g_1, h_1$  be a different solution from  $g_0, h_0$ . Then both  $g_1 \pmod{I^t}, h_1 \pmod{I^t}$  and  $g_0 \pmod{I^t}, h_0 \pmod{I^t}$  are solutions satisfying C1 to C3 for the case of  $s = t$ . Then by the induction hypothesis, we have that there is a  $u_0 \in I^{t/2}$  such that,

$$\begin{aligned} (g_2 \pmod{I^t}) &= (g_1 \pmod{I^t})(1 + u_0) \pmod{I^t} \\ (h_2 \pmod{I^t}) &= (h_1 \pmod{I^t})(1 - u_0) \pmod{I^t} \end{aligned}$$

This implies that

$$\begin{aligned} g_2 &= g_1(1 + u_0) \pmod{I^t} \\ h_2 &= h_1(1 - u_0) \pmod{I^t} \end{aligned}$$

Notice that this is different from the condition in the base case where any two solutions must equal modulo  $I$ . Nevertheless, following the same argument, we can derive that there is an element  $u \in I^t$  such that,

$$g_2 = g_1(1 + u_0)(1 + u) \pmod{I^{2t}}$$

$$h_2 = h_1(1 - u_0)(1 - u) \pmod{I^{2t}}$$

Below we show that  $u_0$  is in fact an element in  $I^t$ , then

$$g_2 = g_1(1 + u_0 + u + u_0u) = g_1(1 + u_0 + u) \pmod{I^{2t}}$$

$$h_2 = h_1(1 - u_0 - u + u_0u) = h_1(1 - u_0 - u) \pmod{I^{2t}}$$

Thus  $g_2 = g_1(1 + u')$  and  $h_2 = h_1(1 - u')$  for  $u' = u_0 + u \in I^t$  as desired.

To show that  $u_0 \in I^t$ , consider:

$$g_2h_2 = g_1h_1(1 - u_0^2)(1 - u^2) \pmod{I^{2t}}$$

$$g_2h_2 = g_1h_1(1 - u_0^2) \pmod{I^{2t}} \quad [\text{as } u \in I^t]$$

Since  $g_1$  and  $h_1$  are not elements in  $I$ , for the last equation to hold, it must be the case that  $u_0 \in I^t$ . Therefore, we conclude the lemma. ■

## 2 Outline of Factoring, revisited

We now give a more complete outline for factoring bivariate polynomials.

Given a monic  $f(x, y) \in \mathbb{F}_q[x, y]$ , with total degree  $d$ , the factoring algorithm *SPLIT* proceeds as follows:

1. If  $g = \gcd(f, \frac{\partial f}{\partial x}) \neq 1$ , then output  $(g, f/g)$  and stop. Otherwise, continue the following steps.
2. Find  $y_0 \in \mathbb{F}$  such that  $f(x, y_0)$  has no repeated factors. This can be done by computing  $\text{Res}\left(f, \frac{\partial f}{\partial x}\right)$ , and plugging in  $y_0 = 1, 2, \dots$  until we find one that makes the resultant non-zero.  
We claim that this will terminate in at most  $d^2$  iterations, as  $\text{Res}\left(f, \frac{\partial f}{\partial x}\right)$  is a polynomial in  $y$  with degree at most  $d^2$ . Furthermore, the first step ensures that  $f$  does not have repeated roots; therefore,  $\text{Res}\left(f, \frac{\partial f}{\partial x}\right)$  is not a zero polynomial. Hence it has at most  $d^2$  roots.
3. Put  $f_{y_0}(x) = f(x, y_0) \pmod{(y - y_0)} = f(x, y_0)$  and factor it. This can be done by using the factoring algorithm for univariate polynomial over  $\mathbb{F}$ . Let  $g$  be an irreducible factor of  $f_{y_0}(x)$ , and  $h$  such that  $f = gh \pmod{(y - y_0)}$ .
4. Now we apply Hensel's Lifting to obtain  $f = g_1h_1 \pmod{(y - y_0)^t}$  for a  $t \approx d^2$
5. Next, from  $g_1$  we ask if we can find a nontrivial factor  $\tilde{g}$  of  $f$ . This is done through the "Jump" step, which tries to find polynomials  $\tilde{g}$  and  $\tilde{h}$  such that  $\tilde{g} = g_1\tilde{h} \pmod{(y - y_0)^t}$ , and  $\tilde{g}$  has small degrees in  $y$  (smaller than  $d$ ) and minimal degree in  $x$ .
6. Finally, return  $\tilde{g}$  and  $f/\tilde{g}$ .