CompSci131

Parallel and Distributed Systems

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Today's topics

More MPI

- Reading assignment:
 - Today: 4.3 and lecture notes
 - Next lecture: Ch 5.
 - » Complete the assignment before next class

Last Lecture Covered

- Berkeley sockets
- MPI



MPI Summary

- Communicator = set of tasks
- MPI_world = a communicator with all the tasks
- Rank = logical id within a communicator
 - 0 to N-1
- Blocking comm. Primitives : Send, Recv
 - Assumes communication is reliable and ordered!
 - » 2 message sent from task I to J are received by J in the same order
- Blocking means
 - Send blocks until its buffer is available for reuse
 - Recv blocks until its buffer has the message

MPI Summary

- Non-blocking comm. Primitives: Isend, Irecv
 - Async. Primitives, continue immediately
 - Return a request id for later tracking
 - » MPI_Wait takes the request id and will block until done
 - Same conditions as for Send, Recv to return
 - MPI_Test is a non-blocking test for a request
- Collective primitives: Barrier, Reduce

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Example C program (LLNL tutorial)

Ring data exchange

```
#include "mpi.h"
#include <stdio.h>
main(int argc, char *argv[]) {
int numtasks, rank, next, prev, buf[2], tag1=1, tag2=2;
MPI_Request reqs[4]; // required variable for non-blocking calls
MPI_Status stats[4]; // required variable for Waitall routine
MPI_Init(&argc,&argv);
MPI Comm_size(MPI_COMM_WORLD, &numtasks);
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
prev = rank-1; next = rank+1; // compute neighbours ranks
if (rank == 0) prev = numtasks - 1;
if (rank == (numtasks - 1)) next = 0;
```

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Example program continued

After MPI_Init each process becomes part of

```
// post non-blocking receives and sends for neighbors
MPI_Irecv(&buf[0], 1, MPI_INT, prev, tag1, MPI_COMM_WORLD,
&reqs[0]);
MPI_Irecv(&buf[1], 1, MPI_INT, next, tag2, MPI_COMM_WORLD,
&reqs[1]);
MPI_Isend(&rank, 1, MPI_INT, prev, tag2, MPI_COMM_WORLD,
&reqs[2]);
MPI_Isend(&rank, 1, MPI_INT, next, tag1, MPI_COMM_WORLD,
&reqs[3]);
// do some work while sends/receives progress in background
// wait for all non-blocking operations to complete
MPI Waitall(4, reqs, stats);
// do more work
MPI_Finalize(); } own id (rank), total number of nodes
```

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More useful MPI primitives

- MPI_Bcast (&buffer,count,datatype,root,comm)
 - Root task sends message to all others
- MPI_Scatter(&sendbuf,sendcnt,sendtype,&recv buf, recvcnt,recvtype,root,comm)
 - Task with data distributes to all
- MPI_Gather = the opposite of MPI_Scatter
- MPI_Barrier
 - Wait for all
- MPI_Reduce

Other types of MPI primitives

Has a number of communication modes, primitives

Primitive	Meaning
MPI_bsend	Buffered send - append outgoing message to a local send buffer
MPI_send	Send a message and wait until copied to local or remote buffer
MPI_ssend	Synchronous send a message and wait until receipt starts
MPI_sendrecv	Send a message and wait for reply
MPI_isend	Pass reference to outgoing message, and continue
MPI_issend	Pass reference to outgoing message, and wait until receipt starts
MPI_recv	Receive a message; block if there are none
MPI_irecv	Check if there is an incoming message, but do not block

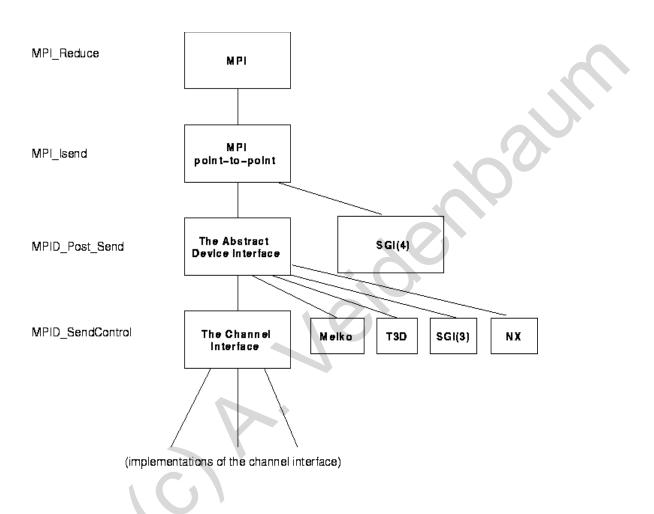
- MPI_Irecv (&buf, count, datatype, dest, tag, comm, request)
- MPI_Wait(&request, &status)
 - Now block until finished

- MPI_Test(&request, &flag, &status)
 - Flag can be set when finished
- MPI_Get_count(&status, datatype, &count)

Portable MPI implementation (MPICH)

- From ANL, initially for MPI-1
- Goals: portability and good performance
- Defined ADI, a set of low-level primitives to
 - Implement all major MPI functions
 - Abstract away details of communication
 - » And possibly buffering
- One implementation defines an even lower interface, that requires only 5 functions

Abstract Device Interface (MPICH)



ADI functionality

Implements 4 sets of functions

- 1. specifying a message to be sent or received
- 2. moving data between the ADI and the messagepassing hardware
- 3. managing lists of pending messages
 - » both sent and received
- 4. providing basic information about the execution environment (e.g., how many tasks are there).

ADI layer contains the code for

- packetizing messages and attaching header information
- managing multiple buffering policies
- matching posted receives with incoming messages (or queuing them)
- etc

Channel Interface

- Five required functions
 - MPID_SendControl or MPID_SendControlBlock» The latter is blocking
 - MPID_RecvAnyControl
 - MPID_ControlMsgAvail
 - MPID_SendChannel
 - MPID_RecvFromChannel
- Notice separation of control and data
- Can have additional functions

How to implement these functions?

Can be done with standard Unix I/O operations

-Select

- » wait until the file descriptor becomes "ready" for some class of I/O operation (e.g., input possible)
- » A file descriptor is ready if it is possible to perform the corresponding I/O operation without blocking
 - e.g., read(2)

-Read

- » from fd to the specified buffer, up to the specified # of bytes
 - "Up to" because of the EOF

- Write

» to fd from the specified buffer, the specified # of bytes

Buffering

- Where is the message buffered?
- A critical issue for good performance
- The channel interface implements three data exchange mechanisms
 - Eager (MPICH default)
 - » data is sent to the destination immediately
 - Rendezvous
 - » When a receive is posted
 - Control is always sent
 - Get
 - » Receiver reads data directly
 - Best on PGAS systems

Sending Modes

- Synchronous mode (MPI_Ssend):
 - the send does not complete until a matching receive has begun
- Buffered mode (MPI_Bsend)
 - the user supplies the buffer to system for its use
- Ready mode (MPI_Rsend)
 - user guarantees that matching receive has been posted.
 - » undefined behavior if the matching receive is not posted
- Non-blocking versions
 - MPI_Issend, MPI_Irsend, MPI_Ibsend
- Note that an MPI_Recv may receive messages sent with any send mode.