

# **Abstract Algebra Note**

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### **Chapter 1 Group Theory**

#### 1.1 Monoid Group

#### **Definition 1.1 (monoid)**

We say that (S,\*) is a monoid if the binary operation satisfies the associative law and has an identity element. That is.

$$\forall x, y, z \in S, \quad x * (y * z) = (x * y) * z$$

and

$$\exists e \in S, \forall x \in S, \quad e * x = x * e = x$$

#### **Definition 1.2 (commutative monoid)**

We say that (S, \*) is a commutative monoid if it is a monoid and the operation satisfies the commutative law. That is,

$$\forall x, y \in S, \quad x * y = y * x$$

#### **Proposition 1.1 (unique of identity element)**

Let  $(S, \cdot)$  be a monoid. Then the identity element is unique.

**Proof** Suppose that e and e' are both identity elements of S. Then

$$e = e \cdot e' = e'$$

so e = e'.

#### **Proposition 1.2 (expand of associative law)**

Let  $x_1, \ldots, x_n, y_1, \ldots, y_m \in S$ . Then

$$x_1 \cdot x_2 \cdot \ldots \cdot x_n \cdot y_1 \cdot y_2 \cdot \ldots \cdot y_m = (x_1 \cdot x_2 \cdot \ldots \cdot x_n) \cdot (y_1 \cdot y_2 \cdot \ldots \cdot y_m)$$

**Proof** We prove this by induction on n.

**Base Case** (n = 1): When n = 1, the statement simplifies to:

$$x_1 \cdot y_1 \cdot y_2 \cdot \ldots \cdot y_m = x_1 \cdot (y_1 \cdot y_2 \cdot \ldots \cdot y_m)$$

This is clearly true by the associative property of multiplication.

**Inductive Step:** Assume the statement holds for n = k, that is:

$$x_1 \cdot x_2 \cdot \ldots \cdot x_k \cdot y_1 \cdot y_2 \cdot \ldots \cdot y_m = (x_1 \cdot x_2 \cdot \ldots \cdot x_k) \cdot (y_1 \cdot y_2 \cdot \ldots \cdot y_m)$$

We need to show that the statement holds for n = k + 1. Consider:

$$x_1 \cdot x_2 \cdot \ldots \cdot x_k \cdot x_{k+1} \cdot y_1 \cdot y_2 \cdot \ldots \cdot y_m$$

By the associative property, we can regroup the terms as:

$$(x_1 \cdot x_2 \cdot \ldots \cdot x_k) \cdot (x_{k+1} \cdot y_1 \cdot y_2 \cdot \ldots \cdot y_m)$$

Using the inductive hypothesis on the first k terms, we have:

$$(x_1 \cdot x_2 \cdot \ldots \cdot x_k) \cdot x_{k+1} \cdot (y_1 \cdot y_2 \cdot \ldots \cdot y_m) = (x_1 \cdot x_2 \cdot \ldots \cdot x_k \cdot x_{k+1}) \cdot (y_1 \cdot y_2 \cdot \ldots \cdot y_m)$$

Thus, the statement holds for n = k + 1.

#### **Proposition 1.3**

Let  $x \in S$  and  $m, n \in \mathbb{N}$ . Then

$$x^{m+n} = x^m \cdot x^n$$

**Proof** We will prove this in three steps:

**Step 1:** First, recall from Proposition 1.2 that for any elements in S:

$$x_1 \cdot x_2 \cdot \ldots \cdot x_n \cdot y_1 \cdot y_2 \cdot \ldots \cdot y_m = (x_1 \cdot x_2 \cdot \ldots \cdot x_n) \cdot (y_1 \cdot y_2 \cdot \ldots \cdot y_m)$$

**Step 2:** Now, consider the special case where all elements are equal to x:

- Let  $x_1 = x_2 = \ldots = x_m = x$
- Let  $y_1 = y_2 = \ldots = y_n = x$

**Step 3:** By definition of exponentiation in a monoid:

$$x^{m+n} = \underbrace{x \cdot x \cdot \dots \cdot x}_{m+n \text{ times}}$$

$$= \underbrace{(x \cdot x \cdot \dots \cdot x)}_{m \text{ times}} \cdot \underbrace{(x \cdot x \cdot \dots \cdot x)}_{n \text{ times}}$$

$$= x^m \cdot x^n$$

Therefore, we have proved that  $x^{m+n} = x^m \cdot x^n$  for all  $x \in S$  and  $m, n \in \mathbb{N}$ .

#### **Definition 1.3 (Submonoid)**

Let  $(S, \cdot)$  be a monoid. If  $T \subset S$ , we say that  $(T, \cdot)$  is a submonoid of  $(S, \cdot)$  if:

- 1. The identity element  $e \in T$
- 2. T is closed under multiplication, that is:

$$\forall x, y \in T, \quad x \cdot y \in T$$

#### **Proposition 1.4**

If  $(T, \cdot)$  is a submonoid of  $(S, \cdot)$ , then  $(T, \cdot)$  is a monoid.

**Proof** We need to verify two properties:

1. The operation is associative in T:

Since  $T \subset S$  and  $\cdot$  is associative in S, it is also associative in T.

2. T has an identity element:

By definition of submonoid, the identity element  $e \in T$ .

Therefore,  $(T, \cdot)$  satisfies all properties of a monoid.

#### **Definition 1.4 (Monoid Homomorphism)**

Let  $(S, \cdot)$  and (T, \*) be monoids, and let  $f: S \to T$  be a mapping. We say f is a monoid homomorphism if f preserves multiplication and maps the identity element to the identity element. That is:

1. For all  $x, y \in S$ :

$$f(x \cdot y) = f(x) * f(y)$$

2. For the identity elements  $e \in S$  and  $e' \in T$ :

$$f(e) = e'$$

**Remark** While a homomorphism preserves operations, an isomorphism represents complete structural equivalence. An isomorphism is first a **bijective mapping**, meaning it establishes a one-to-one correspondence between elements - essentially "relabeling" elements uniquely. Beyond being bijective, an isomorphism preserves operations under this relabeling, implying that the only difference between two structures (like monoids) is their labeling.

**Example 1.1 Different Types of Monoid Maps** Let's examine several maps between monoids:

- 1. A homomorphism that is not an isomorphism: Consider  $f:(\mathbb{Z},+)\to(\mathbb{Z},+)$  defined by f(n)=2n
  - Preserves operation: f(a + b) = 2(a + b) = 2a + 2b = f(a) + f(b)
  - Is injective:  $f(a) = f(b) \implies 2a = 2b \implies a = b$
  - Not surjective: odd numbers are not in the image
  - Therefore: homomorphism but not isomorphism
- 2. Non-isomorphic homomorphism: Consider  $h: (\mathbb{Z}, +) \to (\mathbb{Z}_2, +)$  defined by  $h(n) = n \mod 2$ 
  - Preserves operation:  $h(a+b)=(a+b) \bmod 2=(a \bmod 2+b \bmod 2) \bmod 2=h(a)+h(b)$
  - Not injective: h(0) = h(2) = 0
  - Surjective: image is all of  $\mathbb{Z}_2$
  - Therefore: homomorphism but not isomorphism

#### **Definition 1.5 (Generated Submonoid)**

Let  $(S, \cdot)$  be a monoid and  $A \subset S$  be a subset. The submonoid generated by A, denoted by  $\langle A \rangle$ , is defined as the intersection of all submonoids of S containing A. That is:

$$\langle A \rangle = \bigcap \{ T \subset S : T \supset A, \ T \text{ is a submonoid} \}$$



#### **Proposition 1.5**

Let  $(S, \cdot)$  be a monoid and  $A \subset S$  be a subset. Then  $\langle A \rangle$  is also a submonoid. Therefore, it is the smallest submonoid containing A.

**Proof** We will prove this in two steps:

**Step 1:** Show  $\langle A \rangle$  contains the identity element

Let  $\{T_{\alpha}\}_{{\alpha}\in I}$  be the collection of all submonoids containing A, Each  $T_{\alpha}$  contains the identity e (by definition of submonoid), Therefore  $e\in\bigcap_{{\alpha}\in I}T_{\alpha}=\langle A\rangle$ 

Step 2: Show closure under multiplication

Let  $x, y \in \langle A \rangle = \bigcap_{\alpha \in I} T_{\alpha}$ , Then  $x, y \in T_{\alpha}$  for all  $\alpha \in I$  Since each  $T_{\alpha}$  is a submonoid,  $x \cdot y \in T_{\alpha}$  for all  $\alpha \in I$ , Therefore  $x \cdot y \in \bigcap_{\alpha \in I} T_{\alpha} = \langle A \rangle$ .

#### **Definition 1.6 (Monoid Isomorphism)**

Let  $(S, \cdot)$  and (T, \*) be monoids, and let  $f: S \to T$  be a mapping. We say f is a monoid isomorphism if f is bijective and a homomorphism. That is:

1. f is bijective (one-to-one and onto)

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2. For all  $x, y \in S$ :

$$f(x \cdot y) = f(x) * f(y)$$

3. For the identity elements  $e \in S$  and  $e' \in T$ :

$$f(e) = e'$$

#### **Proposition 1.6**

If  $f:(S,\cdot)\to (T,*)$  is a monoid isomorphism, then  $f^{-1}:T\to S$  is a monoid homomorphism. Therefore,  $f^{-1}$  is also a monoid isomorphism.

**Proof** Since f is an isomorphism,  $f^{-1}$  exists and is bijective. We need to show:

1.  $f^{-1}$  preserves operation:

$$f^{-1}(a * b) = f^{-1}(f(f^{-1}(a)) * f(f^{-1}(b)))$$
$$= f^{-1}(f(f^{-1}(a) \cdot f^{-1}(b)))$$
$$= f^{-1}(a) \cdot f^{-1}(b)$$

2.  $f^{-1}$  preserves identity:

$$f^{-1}(e') = e$$
 where  $e'$  and  $e$  are identity elements

Therefore,  $f^{-1}$  is both a homomorphism and bijective, making it an isomorphism.

## **Chapter 2** Ring Theory

## **Chapter 3** Ring Theory

## **Chapter 4 Polynomial Theory**

## **Chapter 5** Field Theory