Private Cloud Compute

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2024-10-08

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1.1 Requirements

Stateless computation

Private Cloud Compute must use the personal user data that it receives exclusively for the purpose of fulfilling the user's request. This data must never be available to anyone other than the user, not even to Apple staff, not even during active processing. And **this data must not be retained**, including via logging or for debugging, after the response is returned to the user. In other words, we want a strong form of stateless data processing where **personal data leaves no trace** in the PCC system.

1.1 Requirements

Enforceable guarantees

Security and privacy guarantees are strongest when they are entirely technically enforceable, which means it must be possible to **constrain and analyze all the components** that critically contribute to the guarantees of the overall Private Cloud Compute system. To use our example from earlier, it's very difficult to reason about what a TLS-terminating load balancer may do with user data during a debugging session. Therefore, PCC must not depend on such external components for its core security and privacy guarantees. Similarly, operational requirements such as collecting server metrics and error logs must be supported with mechanisms that do not undermine privacy protections.











Appendix

1.1 Requirements

No privileged runtime access

Private Cloud Compute **must not contain privileged interfaces** that would enable Apple's site reliability staff to bypass PCC privacy guarantees, even when working to resolve an outage or other severe incident. This also means that PCC must not support a mechanism by which the privileged access envelope could be enlarged at runtime, such as by loading additional software.

1.1 Requirements

Non-targetability

An attacker should not be able to attempt to compromise personal data that belongs to specific, targeted Private Cloud Compute users without attempting a broad compromise of the entire PCC system. This must hold true even for exceptionally sophisticated attackers who can attempt physical attacks on PCC nodes in the supply chain or attempt to obtain malicious access to PCC data centers. In other words, a limited PCC compromise must not allow the attacker to steer requests from specific users to compromised nodes; targeting users should require a wide attack that's likely to be detected. To understand this more intuitively, contrast it with a traditional cloud service design where every application server is provisioned with database credentials for the entire application database, so a compromise of a single application server is sufficient to access any user's data, even if that user doesn't have any active sessions with the compromised application server.

1.1 Requirements

Verifiable transparency

Security researchers need to be able to verify, with a high degree of confidence, that our privacy and security guarantees for Private Cloud Compute match our public promises. We already have an earlier requirement for our guarantees to be enforceable. Hypothetically, then, if security researchers had sufficient access to the system, they would be able to verify the guarantees. But this last requirement, verifiable transparency, goes one step further and does away with the hypothetical: **security researchers must be able to verify the security and privacy guarantees of Private Cloud Compute**, and they must be able to verify that the software that's running in the PCC production environment is the same as the software they inspected when verifying the guarantees.

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1.2 Taxonomy

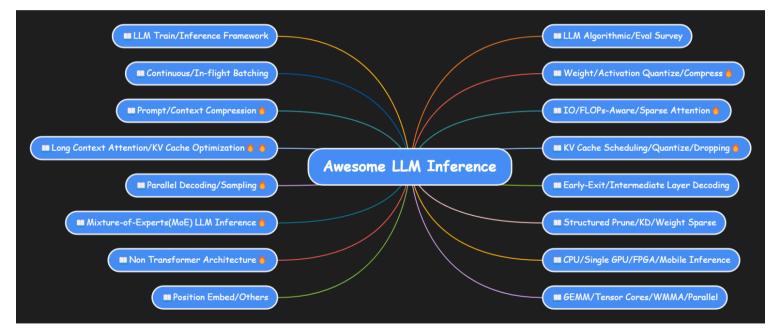
Category	Requirements	Threats	Guarantees	
Technical	Stateless computation	Trace of data after processing Example: Logging, debugging	(Purpose) Only use user data to perform requested operations (Transient) Delete the data after fulfilling the request (Scope) Not available to even Apple staff	
	Non-targetability Targeted attack Example: Steer request to compromised nodes		(Hardware) Hardened supply chain (Scheduler) Requests cannot be user/content-specific routed (Anonymity) OHTTP Relay, RSA Blind Signature (Scope) No system-wide encryption	
Ecosystem -	Enforceable guarantees	Technical enforceability Example: External TLS-terminating load balancer	(Hardware) Secure Enclave, Secure Boot (System) Signed System Volume, Swift on Server (Software) Code Signing, Sandboxing	
	No privileged runtime access	Privileged interfaces Example: Shell access by SREs	No remote shell. Only pre-specified, structured, and audited logs/metrics can leave the node User data is reviewed by multiple indepedent layers	
	Verifiable transparency	Uninspected code	Every production build of PCC publicly available	

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2.1 LLM Inference

Most of the popular decoder-only LLMs (GPT-3, for example) are pretrained on the causal modeling objective, essentially as next-word predictors. These LLMs take a series of tokens as inputs, and generate subsequent tokens autoregressively until they meet a stopping criteria.



LLM Inference

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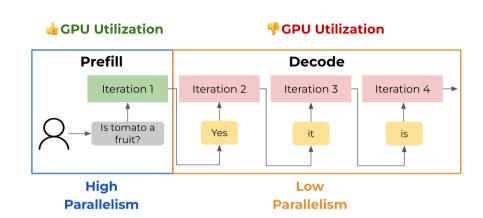
2.2 Phases

Prefill: Processing the input

In the prefill phase, the LLM processes the input tokens to compute the intermediate states (keys and values), which are used to generate the "first" new token. Each new token depends on all the previous tokens, but because the full extent of the input is known, at a high level this is a matrix-matrix operation that's **highly parallelized**. It effectively **saturates GPU utilization**.

Decode: Generating the output

In the decode phase, the LLM generates output tokens autoregressively one at a time, until a stopping criteria is met. Each sequential output token needs to know all the previous iterations' output states (keys and values). This is like a matrix-vector operation that underutilizes the GPU compute ability compared to the prefill phase. The speed at which the data (weights, keys, values, activations) is **transferred to the GPU from memory** dominates the latency, not how fast the computation actually happens. In other words, this is a **memory-bound operation**.



2.3 Challenges

Workload Heterogeneity

Universality and application diversity lead to heterogeneity of the inference requests, in terms of input lengths, output lengths, expected latencies, etc

• Queuing Delays, Preemptions, Interference

Execution Unpredictability

Unknown a priori how many tokens will be generated before the stopping criteria is met. As such, the execution time and the resource demand of a request are both unpredictable.

Multi-Tenant and Dynamic Environment

The system must scale to support multiple users and adapt to the dynamic nature of the environment.

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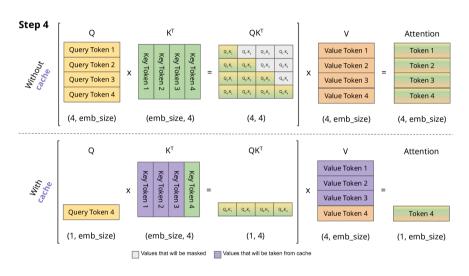


 $\underset{\bigcirc}{\mathsf{Summary}}$

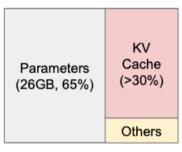
Appendix

3.1 Memory Management

KV Cache



Transformers use attention mechanisms that compute attention scores between tokens. The KV Cache helps by storing previously computed key-value pairs, allowing the model to quickly access and reuse them for new tokens, avoiding redundant calculations.



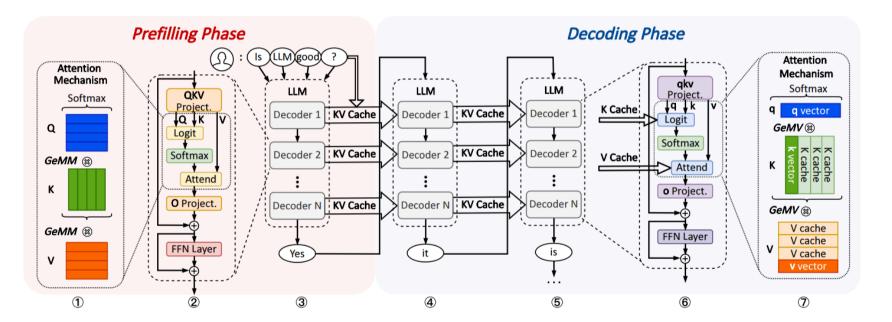
NVIDIA A100 40GB

Memory layout when serving an LLM with 13B parameters on NVIDIA A100. The parameters (gray) persist in GPU memory throughout serving. The memory for the KV cache (red) is (de)allocated per serving request. A small amount of memory (yellow) is used ephemerally for activation.

3.1 Memory Management

LLM inference architecture primarily comprises multiple stacked decoder blocks, each consisting of a self-attention module and a Feed-Forward Neural Network (FFN) module.

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3.1 Memory Management

Paged Attention

Paged Attention* is a technique that divides the attention matrix into smaller pages. This approach provides a near-perfect solution for mitigating fragmentation and hence, PagedAttention has become the de facto standard for dynamic memory allocation in LLM serving systems.

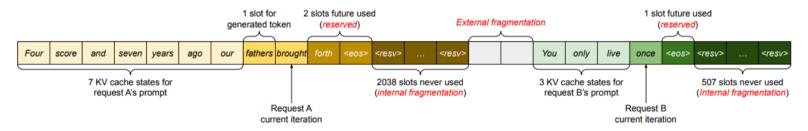


Figure 3. KV cache memory management in existing systems. Three types of memory wastes – reserved, internal fragmentation, and external fragmentation – exist that prevent other requests from fitting into the memory. The token in each memory slot represents its KV cache. Note the same tokens can have different KV cache when at different positions.

Efficient Memory Management for Large Language Model Serving with PagedAttention



3.1 Memory Management

Paged Attention

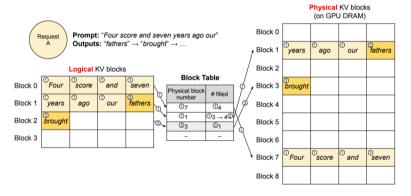


Figure 6. Block table translation in vLLM.

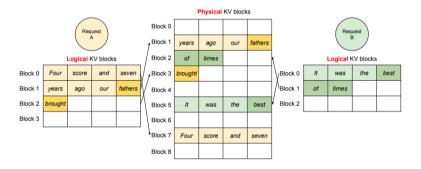


Figure 7. Storing the KV cache of two requests at the same time in vLLM.

Pitfalls*

- Requires re-writing the attention kernel.
- Adds software complexity and redundancy (CPU code), can degrade throughput by 11%.
- \bullet Introduces performance overhead. 20-26% slower than original FasterTransformer kernel.

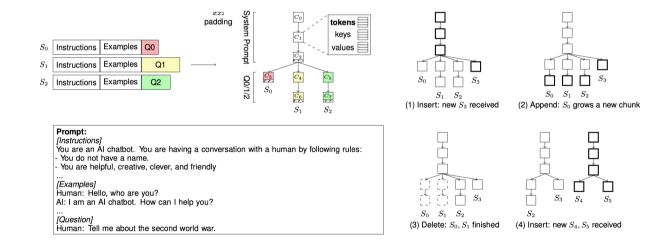
vAttention: Dynamic Memory Management for Serving LLMs without PagedAttention. 2024



3.1 Memory Management

Prefix Caching

Prefix Caching* is a technique that caches the intermediate states of the model during the prefill phase. These states are then reused during the decode phase to speed up inference.



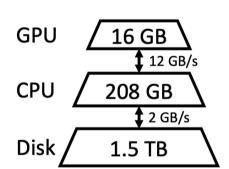
ChunkAttention: Efficient Self-Attention with Prefix-Aware KV Cache and Two-Phase Partition. 2024

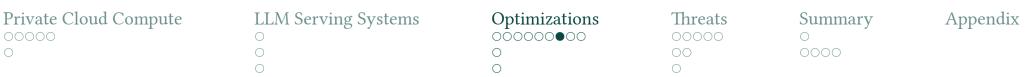
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3.1 Memory Management

KV Cache Offloading

The KV Cache Offloading technique moves the KV cache from the GPU to the CPU to free up GPU memory for other tasks.





3.1 Memory Management

Flash Attention

GPU: One kind of computation done on the input data at a time in sequence

Fusing: Fusing multiple layers together during the actual computation can enable minimizing the data access by GPUs.

FlashAttention* uses **tiling** to fully compute and write out a small part of the final matrix at once

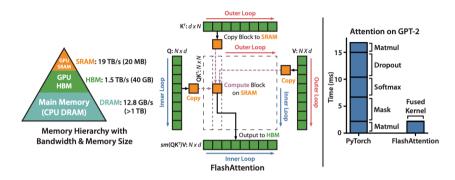


Figure 1: **Left:** FLASHATTENTION uses tiling to prevent materialization of the large $N \times N$ attention matrix (dotted box) on (relatively) slow GPU HBM. In the outer loop (red arrows), FLASHATTENTION loops through blocks of the **K** and **V** matrices and loads them to fast on-chip SRAM. In each block, FLASHATTENTION loops over blocks of **Q** matrix (blue arrows), loading them to SRAM, and writing the output of the attention computation back to HBM. **Right:** Speedup over the PyTorch implementation of attention on GPT-2. FLASHATTENTION does not read and write the large $N \times N$ attention matrix to HBM, resulting in an 7.6× speedup on the attention computation.

FlashAttention: Fast and Memory-Efficient Exact Attention with IO-Awareness

3.1 Memory Management

Group-Query Attention

- Standard Attention: Compute attention for each query separately. Complexity is $O(n^2)$.
- **Multi-Query Attention**: Reuse the same attention matrix for multiple queries. Queries are similar enough to share the same attention distribution.
- **Group-Query Attention***: Divide queries into groups and compute attention for each group separately.

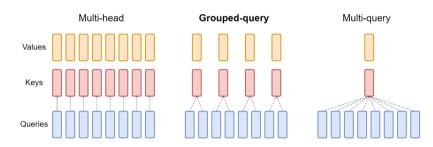


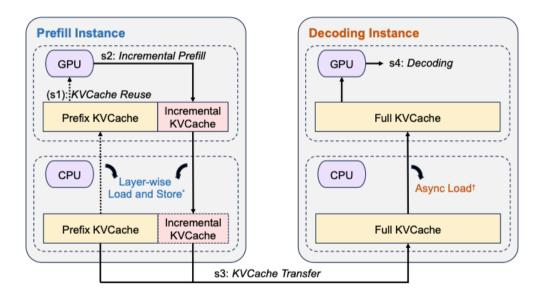
Figure 2: Overview of grouped-query method. Multi-head attention has H query, key, and value heads. Multi-query attention shares single key and value heads across all query heads. Grouped-query attention instead shares single key and value heads for each *group* of query heads, interpolating between multi-head and multi-query attention.

PCC

GQA: Training Generalized Multi-Query Transformer Models from Multi-Head Checkpoints

3.1 Memory Management

Real-World System: Mooncake



3.2 Batch Processing

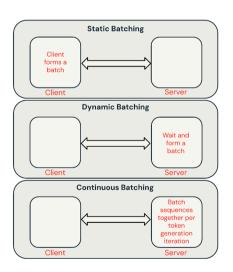
Static batching: Client packs multiple prompts into requests and a response is returned after all sequences in the batch have been completed. Our inference servers support this but do not require it.

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Dynamic batching: Prompts are batched together on the fly inside the server. Typically, this method performs worse than static batching but can get close to optimal if responses are short or of uniform length. Does not work well when requests have different parameters.

Continous Batching: A batch that is continuously processed, leveraging the opportunity by batching new requests once some old requests are finished



Algorithm 1 LLM serving with Continuous batching

- 1: Initialize current batch $B \leftarrow \emptyset$, waiting queue $Q \leftarrow \emptyset$
- 2: > with monitoring stream:
- 3: while True do
- 4: **if** new request *r* arrived **then**
- 5: $Q \leftarrow Q + r$
- 6: ▷ with execution stream:
- 7: while True do
- 8: **if** can_add_new_request() **then**
- 9: $B_{new} \leftarrow select_new_requests(\mathbf{Q})$
- 10: $\operatorname{prefill}(B_{new})$
- 11: $B \leftarrow B + B_{new}$
- 12: decode(B)
- 13: $B \leftarrow \text{filter_finished_requests}(B)$

3.3 Parallel Processing

Data Parallelism (DP)

Each device processes a different batch of data. This allows multiple devices to work on independent data batches in parallel, improving throughput.

Sequence Parallelism (SP)

The input sequence is divided into chunks, with each chunk processed by a separate device. This enables parallel processing of sequence segments across devices.

Pipeline Parallelism (PP)

The model is split vertically into chunks, where each chunk contains a subset of layers. Each device handles a different set of layers, allowing different stages of the model to be executed in parallel across devices.

Tensor Parallelism (TP)

The model's parameters are sharded horizontally into chunks, with each chunk distributed across devices. This allows the computation within each layer to be split and processed in parallel.

3.4 Speculative Inference

Standard inference

Sequence generation is strictly sequential. Each token must be generated based on the previously generated token, which leads to high latency, especially for long-sequence tasks.

Speculative inference*

- **Predict multiple tokens ahead**: When generating the first token, the model simultaneously makes speculative predictions about the next several tokens.
- **Parallel processing**: These speculative predictions allow the model to process multiple possible outcomes in parallel, speeding up the inference.
- Validate predicted paths: If the speculative predictions are correct, the model can continue with these results, avoiding the need to recalculate. If the predictions are incorrect, the model adjusts and corrects the path.

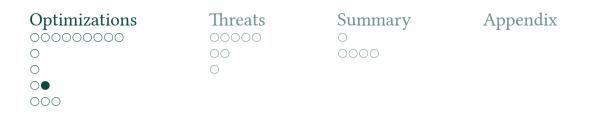
Blockwise Parallel Decoding for Deep Autoregressive Models

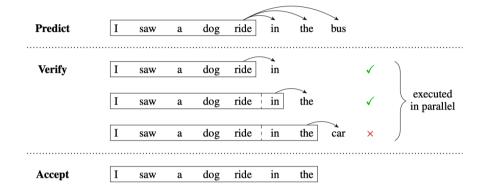
3.4 Speculative Inference

Algorithm*

- p is the smaller draft model, q is the larger target model.

```
Algorithm 2 Speculative Sampling (SpS) with Auto-Regressive Target and Draft Models
  Given lookahead K and minimum target sequence length T.
   Given auto-regressive target model q(.|.), and auto-regressive draft model p(.|.), initial prompt
   sequence x_0, \ldots, x_t.
   Initialise n \leftarrow t.
   while n < T do
      for t = 1 : K do
        Sample draft auto-regressively \tilde{x}_t \sim p(x|, x_1, \dots, x_n, \tilde{x}_1, \dots, \tilde{x}_{t-1})
      end for
      In parallel, compute K + 1 sets of logits from drafts \tilde{x}_1, \dots, \tilde{x}_K:
                       q(x|, x_1, ..., x_n), q(x|, x_1, ..., x_n, \tilde{x}_1), ..., q(x|, x_1, ..., x_n, \tilde{x}_1, ..., \tilde{x}_K)
      for t = 1 : K do
        Sample r \sim U[0, 1] from a uniform distribution.
        if r < \min\left(1, \frac{q(x|x_1,...,x_{n+t-1})}{p(x|x_1,...,x_{n+t-1})}\right), then
           Set x_{n+t} \leftarrow \tilde{x}_t and n \leftarrow n+1.
         else
           sample x_{n+t} \sim (q(x|x_1,...,x_{n+t-1}) - p(x|x_1,...,x_{n+t-1}))_+ and exit for loop.
         end if
      end for
     If all tokens x_{n+1}, \ldots, x_{n+K} are accepted, sample extra token x_{n+K+1} \sim q(x|, x_1, \ldots, x_n, x_{n+K}) and
      set n \leftarrow n + 1.
   end while
```





Accelerating Large Language Model Decoding with Speculative Sampling, 2023

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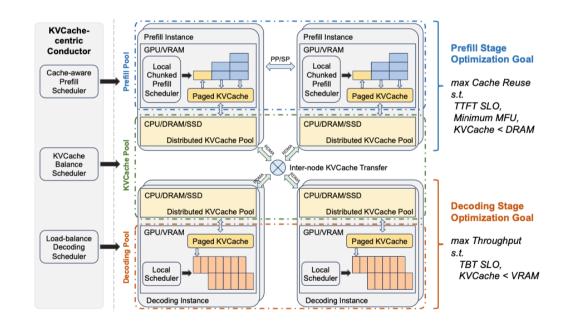
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3.5 Summary

Real-World System: Mooncake

Gray: Control plane



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3.5 Summary

Catagony	Optimization		GPU Resources			Optimization Goal		
Category	Optimization	Compute	Memory	Transmission	Throughput	TTFT	TBT	
	Paging		#		•			
Memory	Prefix Caching		#		•			
Wiemor y	Disk Offloading		H		•			
	Group-Query Attention					+	+	
	Duplication			+		+	+	
Tranmission	Pulling			+	•	+	+	
Traininssion	Request Migration	#	+		#	+	+	
	Disaggregated Arch		+	+			•	
	Continuous Batch			+			•	
Batch	Chunked Prefill					+	+	
	Prepack Prefill							
	Pipeline Parallelism	+			+	-	?	
Parallelism	Tensor Parallelism	+			+		+	
Parallelism	Sequence Parallelism	+				+	?	
	Speculative Inference		-		-		+	

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Optimizations

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3.5 Summary

Trends

Category	Trend	Examples	Conflict
Mamany	Improve reachility of VV Cooks with finer granularity	Paging	S
Memory	Improve reusability of KV Cache with finer granularity	Token-Level Optimization	3
		Data Duplication	
Transmission	Minimizing transmission latency	Prefetching/Pulling	T
		PD Disaggregation	
	Request-level Predictions		
Scheduling	Customized scheduling for specific scenarios	Machine-Level Scheduling	STP
_		Global profiling	
		Pipeline Parallelism	,
Parallelism	Ontimizing navallalism for resource rouse and efficiency	Tensor Parallelism	S
	Optimizing parallelism for resource reuse and efficiency	Sequence Parallelism	3
		Speculative Inference	

S: Stateless computation E: Enforceable guarantees T: Non-targetability P: No privileged runtime access V: Verifiable transparency

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4.1 Stateless Computation

Requirement

The system does not maintain any state between requests and the system does not store any information about previous requests.

Attacker's Capabilities

- **Weak**: An attacker gains access to the system's storage mechanism, potentially compromising databases or disk storage. They can also query the system to infer the presence of sensitive information.
- **Strong**: An attacker gains control over specific nodes, allowing them to request or intercept data within the system. However, they are unable to directly access the original prompt due to model sharding or other security measures.

Overview of Attacker's Capabilities

Capabilities	Goal	Query	Access to Storage	Knowledge of Model	Access to Specific Nodes	Control over Node
Weak	Reconstructing User Inputs and Contexts	/	✓	✓		_
Strong	Reconstructing User Inputs and Contexts	/	√	√	√	<u> </u>

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4.1 Stateless Computation

Prompt & KV Cache

While the KV cache is related to the input prompt, it is not possible to directly infer the original prompt from it due to the complex and non-reversible nature of the transformations involved in generating the cache.

Precompute KV Cache

Precompute the KV cache for a set of known sensitive prefixes. For instance, if the system is used for medical queries, precompute the KV cache for common medical terms.

State

In LLM inference, once KVs are computed for a sequence, their values do not change.

Input

$$X = [x_1, x_2, ..., x_n].$$

Embedded Input

$$E = \text{Embed}(X) + \text{Positional Encoding}(X)$$

= $[E_1, E_2, ..., E_n].$

K Cache

$$K = [E_1 W^k, E_2 W^k, ..., E_n W^k].$$

V Cache

$$V = [E_1 W^v, E_2 W^v, ..., E_n W^v].$$

4.1 Stateless Computation

Weak Attacker

Inference Attack

An attacker could analyze cached data to infer patterns or user behaviors. Even without full query access, understanding what prefixes are frequently cached might reveal the types of queries being made to the system.

Example: If the cache contains **frequent prefixes related to medical inquiries**, the attacker could infer that the LLM is being used in a healthcare context.

Replay Attack

Attacker can craft new queries that match existing keys in the cache. This allows them to replay or trigger cached computations, potentially extracting sensitive information based on model completions.

Example: If the key in the KV cache is "My bank account number is...", the attacker can craft similar queries to attempt to elicit a continuation that reveals more about the original input.

4.1 Stateless Computation

Strong Attacker

Targeted Data Extraction

An attacker controlling a node could craft specific inputs designed to trigger the retrieval of cached prefixes. By doing this repeatedly, they can extract sensitive information from the cache based on the model's responses.

Example: By submitting inputs like "My Social Security Number is..." and monitoring responses, they could **infer** whether such a prefix was previously cached and what context or continuation it triggers.

Cache Content Mapping

The attacker maps key tensors to actual input tokens using controlled nodes, thereby obtaining the full sequence of user inputs. The attacker can reconstruct the complete input sequence for high-frequency queries or commonly used phrases, directly exposing user input such as personal queries or confidential data.

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4.1 Stateless Computation

Optimization	Stored States	Location	Mitigation	Weak At	tacker	Strong A	ttacker
Optimization	Stored States	Location	Mitigation	Inference	Replay	Extraction	Mapping
D. C. O. L. W. G. I. GI		GPU Memory	Cache Expiry		,	,	
Prefix Caching	KV Cache	CPU Memory	Isolation			/	√
Diale Offlooding	KV Cache	Disk Storage	Encryption	√	,	/	
Disk Offloading		(SSD, Hard Drive)			√		
Dulling	VV Caalaa	GPU Memory	Randomized Scheduler	,	,	,	
Pulling	KV Cache	CPU Memory				/	/
Database-based	Talsas	GPU Memory	Differential Drivers			,	
Speculative Inference	Token	CPU Memory	Differential Priavacy				

4.2 Non-Targetability

Non-Targetability: An attacker should not be able to attempt to compromise personal data that belongs to specific, targeted Private Cloud Compute users without attempting a broad compromise of the entire PCC system.

Definition: Let $S = \{S_1, S_2, ..., S_n\}$ denote the set of all servers in the system, with the capability of each server S_i represented by $C(S_i)$. The set of requests handled by these servers is denoted as $R(S) = \{R(S_1), R(S_2), ..., R(S_n)\}$. The system is considered non-targetable if, for any subset $T = \{T_1, T_2, ..., T_m\} \subseteq S$ of servers, the probability of compromising the data of a specific user u is given by:

$$P(u \in R(T)) = \frac{\sum_{i=1}^{m} C(T_i)}{\sum_{i=1}^{n} C(S_i)}$$

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4.2 Non-Targetability

$$P(u \in R(T)) = \frac{\sum_{i=1}^{m} C(T_i)}{\sum_{i=1}^{n} C(S_i)}$$

Goal

- Increase Hit Count: $\sum C(T_i)$
- Decrease Miss Count: $\sum C(S_i) \sum C(T_i)$

Optimization	Increase Hit Rate	Decrease Miss Rate			
Dunlingtion	✓	✓			
Duplication	Duplicate targeted cache to other victims	Duplicate untargeted cache to non-victims			
Dulling	✓				
Pulling	Pull targeted cache from other victims				
Priority-Based Scheduling	✓	✓			
Priority-based Scheduling	Prioritize targeted requests	Deprioritize untargeted requests			
Degreet Level Drediction	✓	✓			
Request-Level Prediction	Prioritize targeted requests	Deprioritize untargeted requests			
Instance Elin	/				
Instance Flip	Flip victim(s) to prefill instance				

4.3 Ecosystem

No Privileged Runtime Access Definition

The system must not contain privileged interfaces that would enable site reliability staff to bypass PCC privacy.

Violations

• Global Profiling

Enforceable Guarantees Definition

The system must provide guarantees that can be enforced by the system itself. These guarantees must be technically enforceable and not rely on external components or human intervention.

Violations

- Disk Offloading
- Pulling
- Machine-level Scheduler

Verifiable Transparency Definition

Security researchers must be able to verify the security and privacy guarantees of Private Cloud Compute, and they must be able to verify that the software that's running in the PCC production environment is the same as the software they inspected when verifying the guarantees.

Violations

• Non open-source systems

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5.1 Academic Systems

Category	Optimization	Threat	Orca 2206	FlexGen 2303	FastServe 2305	SpecInfer 2305	vLLM 2309	REST 2311	Splitwise 2311	SGLang 2312	Lookahead 2312	Sarathi 23-24	InfiniteLLM 2401	DistServe 2401	Medusa 2401	TetriInfer 2401	AttentionStore 2403	LoongServe 2404	Andes 2405	Llumnix 2406	Preble 2407	MInference 2407	TR 2408
	Paging						Initial			Z		V				V							
Memory	Prefix Caching	SE								V											V		
	Disk Offloading	SE		V													V						
	Duplication	T																					
Tranmission	Pulling	SET												V									
114111111551011	Request Migration																	V		V			
	Disaggregated Arch								V					V		V							
	Continuous Batch		Initial		Z	Z	V					V		N		V							
Batch	Chunked Prefill											Initial				V					V		
	Prepack Prefill													v		v							
	Speculation					V		V		V	Z				V								
	Context-Based Speculation	S						V															
Parallelism	Database-Based Speculation	S									V												V
	Tensor Parallelism					Z																	
	Sequence Parallelism																	V					
	Priority-Based	T			V					V		✓				V			V	/	V		
	Request-Level Prediction	T			V	V										V							
Scheduling	Machine-level Scheduler	E			V				V				V			V		V			V		
	Instance Flip	T							V							V							
	Global Profiling	P		V					V					V								V	
Verification	Non Open-Source	v														V							V

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Private	Cloud	Compute
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LLM Serving Systems

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Appendix

Miscellaneous

Title	Keywords	Contributions
Prompt Cache	Prefill, Memory	Reuse attention states across different LLM prompts. Parse the prompt and use reusable text segments(snippet)
Layer-wise Transmission	Transmission	Transmit each layer's output to the next layer in the pipeline, instead of transmitting the entire model's output
LightLLM	Interface	Use http as the interface to the system
SkyPilot Cross Region & Cloud		Given a job and its resource requirements (CPU/GPU/TPU), SkyPilot automatically figures out which locations (zone/region/cloud) have the compute to run the job, then sends it to the cheapest one to execute
MLC LLM Efficient Execution		Enable efficient execution of large language models across a wide range of hardware platforms, including mobile devices, edge devices, and even web browsers
vAttention	Virtual Memory	stores KV-cache in contiguous virtual memory and leverages OS support for on-demand allocation of physical memory
MemServe	API, Framework	an elastic memory pool API managing distributed memory and KV caches across serving instances
CacheGen	Network, Streaming	CacheGen uses a custom tensor encoder, leveraging KV cache's distributional properties to encode a KV cache into more compact bitstream representations
DynamoLLM	Energy	It exploits heterogeneity in inference compute properties and fluctuations in inference workloads to save energy
MInference	Prefill, Long Context	Addresses the expensive computational cost and the unacceptable latency of the attention calculations in the pre-filling stage of long-context LLMs by leveraging dynamic sparse attention with spatial aggregation patterns
Shared Attention	Attention	directly sharing pre-computed attention weights across multiple layers in LLMs
SnapKV	Compression	Observing that specific tokens within prompts gain consistent attention from each head during generation, our methodology not only retrieve crucial information but also enhances processing efficiency

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5.2 Industrial Systems

Category	Optimization	Threat	vLLM Open Source	LightLLM Open Source	FlexFlow Open Source	SGLang Open Source	Mooncake Moonshot	DeepSpeed Microsoft	TensorRT NVIDIA	TGI Hugging Face	Llama Intel	LMDeploy Shanghai AI lab	fastllm Open Source	rtp-llm Alibaba	MindIE Huawei
	Paging		V			V	V	V		V				/	
	Token Attention			V											
Memory	Prefix Caching	s	V						V						
	Disk Offloading	SE				V	V		V						✓
	Multi-Query Attention								V						
	Duplication	T					V								
Tranmission	Pulling	SET													
Tranmission	Request Migration														
	Disaggregated Arch		V				V								/
	Continuous Batch		V		N		V	V	V	V	V	V	V	✓	
Batch	Chunked Prefill		V				V	V							
	Prepack Prefill														
	Speculation		V		N	N			V	N				V	V
Parallelism	Tensor Parallelism									V					V
	Sequence Parallelism														
	Priority-Based	T				V	V	V							
	Request-Level Prediction	T		V		V									
Scheduling	Machine-level Scheduler	E				V	V								
	Instance Flip	Т													
	Global Profiling	P					V								
Verification	Non Open-Source	v					Z		V						

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5.2 Industrial Systems

The roadmap of the vLLM project includes the following features:

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Version	Date	Memory	Transmission	Batch	Parallelism	Scheduling	Model
v0.1	2306	Paging		Continuous Batching			MQA, GQA
v0.2	2309				Better TP & EP Support		AWQ
v0.3	2401	Prefix Caching					GPTQ
v0.4	2404		Optimize Distributed Communication	Chunked Prefill	Speculative Inference		
v0.5	2407	CPU Offloading			Support PP	Schedule multiple GPU steps in advances	FP8
v0.6	2409					Asynchronous output processor	

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5.2 Industrial Systems

Trends

Private Cloud Compute

Category	Trend	Examples	Conflict			
Memory	Improve reusability of KV Cache with finer granularity	Paging	S			
Wiemory	improve reusability of KV Cache with infer granularity	Token-Level Optimization	S			
		Data Duplication				
Transmission	Minimizing transmission latency	Prefetching/Pulling	T			
		PD Disaggregation				
	Request-level Predictions					
Scheduling	Customized scheduling for specific scenarios	Machine-Level Scheduling	STP			
		Global profiling				
		Pipeline Parallelism				
Parallelism	Tensor Parallelism					
r ai aileiisiii	Optimizing parallelism for resource reuse and efficiency	S				
		Speculative Inference				

LLM Serving Systems

S: Stateless computation E: Enforceable guarantees T: Non-targetability P: No privileged runtime access V: Verifiable transparency



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Attention

https://zh.d2l.ai/chapter_attention-mechanisms/attention-cues.html

Inference Optimization

https://github.com/DefTruth/Awesome-LLM-Inference

https://developer.nvidia.com/blog/mastering-llm-techniques-inference-optimization/

Parallelism

https://developer.nvidia.com/blog/demystifying-ai-inference-deployments-for-trillion-parameter-large-language-models/

Utilities

https://github.com/Trusted-AI/adversarial-robustness-toolbox

6 Appendix

Confidential Computing on nVIDIA H100 GPU: A Performance Benchmark Study

Model	Т	PS (tokens	s/s)	$ ext{QPS (req/s)}$					
	TEE-on	TEE-off	Overhead	TEE-on	TEE-off	Overhead			
LLama-3.1-8B	123.2985	132.3618	6.85%	18.2141	18.8208	3.22%			
Phi3-14B-128k	66.5845	69.7787	4.58%	7.1760	7.3456	2.31%			
Llama-3.1-70B	2.4822	2.4789	$-0.13\%^{1}$	0.8325	0.8295	$-0.36\%^{2}$			

Table 1: Performance comparison of TEE-on and TEE-off modes for various models in terms of TPS (tokens per second) and QPS (queries per second).

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Confidential Computing on nVIDIA H100 GPU: A Performance Benchmark Study

Model		TTFT (s)	ITL (s)					
	TEE-on	TEE-off	Overhead	TEE-on	TEE-off	Overhead			
LLama-3.1-8B	0.0288	0.0242	19.03%	1.6743	1.5549	7.67%			
Phi3-14B-128k	0.0546	0.0463	18.02%	3.7676	3.5784	5.29%			
Llama-3.1-70B	0.5108	0.5129	$-0.41\%^{3}$	94.8714	95.2395	$-0.39\%^{4}$			

Table 2: Comparison of TTFT (Time to First Token) and ITL (Inter Output Token Latency) for TEE-on and TEE-off modes across models.

Model	TPS	- short (to	kens/s)	TPS -	medium (t	okens/s)	TPS - long (tokens/s)			
	TEE-on	TEE-off	Overhead	TEE-on	TEE-off	Overhead	TEE-on	TEE-off	Overhead	
LLama-3.1-8B	127.0310	136.8282	7.16%	122.9356	132.0464	6.90%	122.9705	131.7333	6.65%	
Phi3-14B-128k	70.9799	74.7556	5.05%	66.1690	69.3104	4.53%	66.2987	69.4176	4.49%	
Llama-3.1-70B	2.5983	2.6073	0.34%	2.4413	2.4374	$-0.16\%^{5}$	2.5245	2.5168	$-0.30\%^{6}$	

Table 3: Performance comparison of TEE-on and TEE-off modes across different sequence lengths in terms of TPS (tokens per second). Short sequences are no longer than 100 tokens. Medium sequences are no longer than 500 tokens. Long sequences are between 501 and 1500 tokens.

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Stateful Inference Systems

Static state States in traditional systems can be modified after creation and require various consistency and coherence mechanisms to support parallelism. In LLM inference, once KVs are computed for a sequence, their values do not change.

Regular computation patterns LLMs' transformer computation is regular. Its computing and memory consumption is determined by the model size, the prompt length, and the output generation length. The model size and a request's prompt length are known before execution, and output is generated one token per iteration. Thus, we can estimate the computing and memory consumption for every iteration.

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Quantization

Quantization is the process of reducing the precision of a model's weights and activations.

Sparsity

Sparsity is the process of setting a portion of the model's weights to zero. Then the model can be expressed as a sparse matrix.

Distillation

Distillation is the process of training a smaller model to mimic the behavior of a larger model.