

Time Complexity

Define hypothetical model machine

- Single vs multi processor
- Read / write speed to memory
- 32-bit vs 64-bit architecture
- ...

Running time / rate of growth of time

- Asymptotic notation: Big Oh, Omega, Theta
 - $f(n) = O(g(n))$ means $c * g(n)$ is an upper bound on $f(n)$.
 - $f(n) = \Omega(g(n))$ means $c * g(n)$ is a lower bound on $f(n)$.
 - $f(n) = \Theta(g(n))$ means $c1 * g(n)$ is an upper bound on $f(n)$ and $c2 * g(n)$ is a lower bound on $f(n)$.
- We analyse time complexity of a algorithm for very large input size.
 - Drop lower order terms
 - With Big Oh we discard multiplication of constants: n^2 and $1000 * n^2$ are treated identically.
 - Measure all fragments
 - Pick complexity of condition which is the worst case
 - $O(f(n)) + O(g(n)) \Rightarrow O(\max(f(n), g(n))) \Rightarrow n^3 + n^2 + n + 1 = O(n^3)$
 - $O(f(n)) * O(g(n)) \Rightarrow O(f(n) * g(n))$

Space Complexity

- It is a measure of how efficient your code is in terms of the largest memory use by the program when it runs.
- It analysis happens almost in the same way time complexity analysis happens.

Examples

Growth rates

- $1 < \log(n) < n < n * \log(n) < n^2 < n^3 < 2^n < n!$
 - Constant functions
 - Logarithmic: binary search, arises in any process where things are repeatedly halved or doubled
 - Linear: traversing every item in an array
 - Super-linear: quick sort, merge sort
 - Quadratic: going thru all pairs of elements, insertion sort, selection sort
 - Cubic: enumerating all triples of items
 - Exponential: enumerating all subsets of n items
 - Factorial: generating all permutations or orderings of n items

Recursion Complexity Analysis

- Let $T(n)$ be the recursive function.

- Bellow are examples of inner calls of $T(n)$, their complexity, and example algorithm that uses them:
 - $T(n) = T(n/2) + O(1)$, $O(\log(n))$, binary search
 - $T(n) = T(n-1) + O(1)$, $O(n)$, sequential search
 - $T(n) = 2 * T(n/2) + O(1)$, $O(n)$, tree traversal
 - $T(n) = 2 * T(n/2) + O(n)$, $O(n * \log(n))$, merge sort, quick sort
 - $T(n) = T(n-1) + O(n)$, $O(n^2)$, selection sort
- Geometric progression:

$$a(n) = a(1) * q^{(n-1)}$$

$$S(n) = a(1) * (q^n - 1) / (q - 1)$$

If it converges:

$$S(\infty) = a(1) / (1 - q)$$

Combinatorics

- All pairs: $O(n^2)$
- All triples: $O(n^3)$
- Number of all permutations: $n!$
- n over k : number of combinations for choosing k items from a set of n
 - $(n \text{ over } k) = n! / (k! * (n-k)!)$
- Number of subsets from a set of n : 2^n
 - Subset = any unique set of k elements from n , including the empty set).
 - For example: $\text{set}=\{1, 2, 3\}$, $\text{subsets}=\{\}, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}$ (ordering in a set doesn't matter).

Handy Formulas

- $1 + 2 + \dots + n = n * (n + 1) / 2$
- $x + x/2 + x/4 + \dots = 2x$