

**SYNERGISTIC GEOMETRY PROCESSING:
FROM ROBUST GEOMETRIC MODELING TO EFFICIENT PHYSICAL
SIMULATIONS**

by

Zhongshi Jiang

A DISSERTATION SUBMITTED IN PARTIAL FULFILLMENT
OF THE REQUIREMENTS FOR THE DEGREE OF
DOCTOR OF PHILOSOPHY
DEPARTMENT OF COMPUTER SCIENCE
NEW YORK UNIVERSITY
SEPTEMBER, 2022

Professor Daniele Panozzo

© ZHONGSHI JIANG
ALL RIGHTS RESERVED, 2022

ACKNOWLEDGEMENTS

ABSTRACT

The numerical solution of partial differential equations (PDEs) portrays how a craft would behave even before it is created. In practice, a typical product begins its lifecycle from a digital specification of the shape and material. Computer simulation would then test the behavior to match the intent: the car hood's deformation under pressure or a pan's temperature when heated. In this case, a robust and accurate simulation could lift many iterations of manufacturing to the digital platform. In addition, as it is safer to conduct more stress experiments in a virtual environment, a reliable simulation also promotes safety, especially in designing biomechanical equipment.

There are various approaches to specify a shape, depending on the different geometry characteristics. Shape modeling evolves beyond smooth and regular objects like sculpture or furniture. Algorithm deduced structures have gained popularity since they deliver the promise of maximal durability with limited material. Downstream, computer simulation requires a discrete tessellation of the object. A notable example is a mesh: it contains a list of vertices for the spatial locations and a list of simple polygons or polyhedra. However, it is rarely the case where one mesh fits all. More degrees of freedom give a more accurate description of the desired shape, but it inevitably means more computational resources are needed to handle the computation.

Different representations between design methods and within various adaptations of simulations, even for the same shape, are not reliably correlated. In the general case, the material, attribute, or experiment setting cannot be robustly re-used on another representation without laborious manually fixing. Such a pipeline demands the engineer or designer to understand precisely the properties of the specific numerical methods of choice. It also prohibits an automatic and robust pipeline to compute optimal structure for the desired physical behavior.

My Ph.D. research's central thesis is to develop algorithms that reliably associate different representations of discrete geometry entities at diverse design and simulation stages. Instead of re-assigning properties at various phases of the design-simulation pipeline, or whenever the required budget is different, a good association between geometry at different stages gives more freedom to the design and development of automatic and adaptive solvers. The adaptation can be bi-directional: on the one hand, with coarse and concise tessellation, the algorithm utilizes the information associated with the original design to obtain more faithful simulations. On the other hand, in early-stage prototyping, and automatic data generation for deep learning, the principles I established allows for simplifying the shape and retaining the settings, providing a fast answer. In addition, I use the same principle to design the first large-scale and robust approach for curved mesh representation, which provides accurate representation as well as efficient physical simulations.

Besides the algorithm design, I also explore the underpinning theory to provide guarantees. Theoretical inquiry not only gives us a better understanding of digital geometry and shapes but also lays the foundation for suitable applications. On the other hand, we perform extensive numerical validations with the implemented software, involving tens of thousands of complex different geometry shapes. Finally, I also release the program implementation and detailed specifications to be open source and accessible.

CONTENTS

Acknowledgments	iii
Abstract	iv
List of Figures	ix
List of Tables	xvii
1 Introduction	1
1.1 Preliminaries	1
1.1.1 bijective maps	1
1.1.2 mathematical optimization	1
2 Bijective Maps with Triangulation Augmentation	2
2.1 Introduction	2
2.2 Previous Work	4
2.3 Method	6
2.3.1 General Formulation	6
2.3.2 Surface Parametrization	9
2.3.3 Extension to 3D	11
2.4 Results	11
2.5 Limitations and Concluding Remarks	16
3 Bijective Projection in the Shell	21
3.1 Introduction	21
3.2 Related Works	24
3.2.1 Attribute Transfer	24
3.2.2 Shell Generation	25
3.2.3 Robust Geometry Processing	26
3.2.4 Isotopy between surfaces	27
3.3 Method	27
3.3.1 Shell and Projection	28
3.3.2 Validity Condition	31

3.3.3	Shell Initialization	32
3.3.4	Shell Optimization	35
3.3.5	Singularities	38
3.3.6	Boundaries	40
3.4	Results	40
3.5	Applications	45
3.6	Variants	48
3.7	Limitations	53
3.8	Concluding Remarks	54
4	Bijective and Coarse High-Order Tetrahedral Meshes	55
4.1	Introduction	55
4.2	Related Works	56
4.2.1	Curved Tetrahedral Mesh Generation	57
4.2.2	Curved Structured Mesh Generation	58
4.2.3	Boundary Preserving Tetrahedral Meshing	59
4.2.4	Curved Surface Fitting	60
4.3	Shell Preliminaries	60
4.3.1	Variation from the Original Algorithm	62
4.4	Curved Tetrahedral Mesh Generation	62
4.4.1	High-order Shells	66
4.4.2	Distance Bound	66
4.4.3	Tetrahedral Meshing	67
4.5	Feature Preserving Curved Shell	69
4.5.1	Feature straightening.	72
4.6	Results	73
4.6.1	Large Scale Validation.	73
4.6.2	Comparisons	75
4.6.3	Flexibility	77
4.6.4	Applications	78
4.7	Limitations and Concluding Remarks	80
5	Declarative Specification for Unstructured Mesh Editing Algorithms	81
5.1	Introduction	81
5.2	Related Work	83
5.2.1	Mesh Data Structures	83
5.2.2	Domain specific languages in graphics	84
5.2.3	Parallel Meshing	84
5.2.4	Scope of Mesh Editing	85
5.3	Method	87
5.3.1	Mesh editing components	87
5.3.2	Declarative Specification	89
5.3.3	Implementation.	92

5.3.4	Example: Shortest Edge Collapse	93
5.4	Applications	94
5.4.1	Parallelization	99
5.4.2	Algorithm Modifications	100
5.4.3	Large-scale dataset validation	100
5.5	Concluding Remarks	103
6	Conclusion	104
6.1	Overview	104
6.2	Future Directions	104
A	Shell Specifics	105
A.1	Proofs	105
A.1.1	Theorem 3.2	105
A.1.2	Proposition 3.4	105
A.1.3	Theorem 3.5	106
A.1.4	Theorem 3.6	106
A.1.5	Theorem 3.7	107
A.2	Extension to meshes with Singularities	108
A.3	Reduction to a Standard QP	108
A.4	Bijectivity of the Nonlinear Prismatic Transformation	109
A.5	Double Slab	109
A.6	Unpublished Material	109
A.6.1	Pointless Variant	110
A.6.2	Dynamic Intersection Check	110
A.6.3	Feature Snapping Shell	110
A.6.4	Positive Determinant for Natural Prism Map	110
B	Curve Meshing related	111
B.1	The geometric map is bijective	111
B.2	Local Operations	112
B.3	Boundary preserving TetGen comparison	113
B.4	TetWild Overview	113

LIST OF FIGURES

2.1	The Nefertiti model with prescribed seams is UV mapped by our algorithm. Each chart is bijective mapped into a circle or ring with Tutte's embedding and achieves minimal distortion in less than a second. The layout is further improved interactively and the final parametrized model is shown on the right. Our approach guarantees a valid UV map with no inverted elements or overlapping triangles. See the attaching video for the optimization and manual interaction.	2
2.2	The initial mesh \mathcal{M} (in green, left), is embedded in another mesh D (in gray, middle) that covers a box in the ambient space and contains the same triangles as \mathcal{M} . D might contain additional points (orange). We denote the triangles that are in D but not in \mathcal{M} as the scaffold S . Our algorithm deforms D , inducing a corresponding deformation on \mathcal{M} (right), while keeping the boundary (blue vertices) fixed and preventing changes in the triangle orientation.	7
2.3	Different values of λ do not affect the result, but they change the number of iterations needed. From left to right: we used a large weight (100x ours), our weight, and a small weight (0.01x ours). The optimization took 9,7, and 8 iterations, respectively, to reach the same energy level.	9
2.4	A bijective map from a circle (left) to a spiral (right) is computed without (top) and with (bottom) the iterative remeshing step. The slivers in the triangulation locks the optimization (top), preventing it from reaching the target shape.	10
2.5	Two models are cut using [Bommes et al. 2009] and bijectively parametrized using our algorithm. See the additional material for more examples.	12
2.6	We compare the distortion energy with respect to the number of iterations on a set of Lucy's meshes with different resolutions (from 1 to 12 million faces). In the center of the plot, we show the 1M Lucy model parametrized by our algorithm.	13
2.7	A mesh is cut by an artist into a single chart and parametrized using SLIM [Rabinovich et al. 2017] (left) and with our algorithm (right). Note that local-injectivity is not sufficient for this model, since the global overlaps in the highlighted region prevent this parametrization from being a UV texture map. Our result (right) is guaranteed to be bijective.	14
2.8	A model with multiple chart (left) is automatically parametrized in a texture atlas (bottom-right) by first mapping each component to a circle (top-right) and then minimizing the distortion.	15

2.9	We remove the self-intersections from a genus 0 model using the conformalized flow [Kazhdan et al. 2012; Sacht et al. 2013]. The flow is inverted, while using our algorithm to compute a bijective volumetric map, to recover a self-intersection free version of the original surface. The final model can now be meshed using TetGen, since it is free from self-intersections.	16
2.10	We grow a bunny inside a box, while preventing self-intersections. We show the result after 0,10,20,30,40, and 50 iterations.	16
2.11	We repeat the challenging test in [Smith and Schaefer 2015] with a subdivided version of their Hilbert curve to increase the triangle count. Our method starts from a disc (upper left), gracefully extends (upper right), and reaches the same minimum (lower left) in 39 minutes whereas [Smith and Schaefer 2015] didn't terminate more than 5 days (lower right), highlighting our performance boost of over 200 times.	17
2.12	We apply our algorithm on 4 models used in [Smith and Schaefer 2015] (using the same stopping criteria) obtaining visually identical results. Distortion errors produced by our algorithm (outer) and theirs (inner) are shown in black.	18
2.13	A single iteration of our algorithm (from left to right) drastically reduces the distortion. The black vector in the center is 150 times longer than the average edge length of its 1-ring. Iterative methods would need thousands of iterations to achieve a similar progress.	19
2.14	Our algorithm is independent to the initial orientation. We rotate the initializing Tutte's mapping of the camel model and obtain results with similar isometric distortion.	20
3.1	A low-quality mesh with boundary conditions (a) is remeshed using our shell (b) to maintain a bijection between the input and the remeshed output. The boundary conditions (arrows in (a)) are then transferred to the high-quality surface (c), and a non-linear elastic deformation is computed on a volumetric mesh created with TetGen (e). The solution is finally transferred back to the original geometry (d). Note that in this application setting both surface and volumetric meshing can be hidden from the user, who directly specifies boundary conditions and analyses the result on the input geometry.	22
3.2	Overview of our algorithm. We start from a triangle mesh, find directions of extrusion, build the shell, and optimize to simplify it.	27
3.3	Example of the top (left, outer) and bottom (right, inner) surface of the prismatic shell.	28
3.4	A prism Δ (left) is decomposed into 6 tetrahedra (middle, for clarity, we only draw the 3 tetrahedra of the top slab). Each tetrahedron has a constant vector field in its interior (pointing toward the top surface), which is parallel to the only pillar of the prism that contains the point.	29

3.5	A point p (left) is traced through \mathcal{V} inside the top part of the shell. A ray with p as origin and \mathcal{V} as direction is cast inside the orange tetrahedron (middle). The procedure is repeated (on the blue tetrahedron) until the ray hits a point in the middle surface (right).	30
3.6	A 2D illustration for the normal dot product condition. The blue arrows agrees with the background vector field (white arrows), while the red arrows do not agree.	30
3.7	The composition $\mathcal{P}_{\mathcal{S}_2}^{-1}(\mathcal{P}_{\mathcal{S}_1}(x))$ with $x \in \mathcal{S}_1$, of a direct and an inverse projection operator defines a bijection between two sections \mathcal{S}_1 and \mathcal{S}_2	31
3.8	The vector field aligned with the pillar edge (orange) has a negative dot product with the green triangle normal; the tetrahedron with this field direction meets the green triangle at the red vertex. As a consequence, the mid-surface is not a section. After topological beveling, the shell becomes valid since the dot product between the green normal and the new pillar (purple) is positive.	34
3.9	The beveling patterns used to decompose prisms for which \mathcal{T} is not a section.	34
3.10	Our algorithm refines the input model (left) with beveling patterns (middle) to ensure the generation of a valid shell. The subsequent shell optimizations gracefully remove the unnecessary vertices (right).	35
3.11	Different local operations used to optimize the shell. Mesh editing operations translate naturally to the shell setting. For vertex smoothing, we decompose the operation into 3 intermediate steps: pan, zoom, and rotate.	36
3.12	A model with 48 singularities and a close up around one (left). Our shell is pinched around each of them without affecting other regions (right).	38
3.13	Left to right: examples of a singularity as a feature point and as a meshing artefact; and illustration of the degenerate prism with a singularity (red) and its tetrahedral decomposition made of only four tetrahedra.	39
3.14	AA' is a direction with positive dot product with respect to all its neighboring faces. However, no valid shell can be built following that direction.	40
3.15	An example of a mesh with boundary.	41
3.16	The effect of different target thickness on the number of prisms $ F $ of the final shell, and distribution of final thickness (shown as the box plots on the top).	41
3.17	Gallery of shells built around models from Thingi10k [Zhou and Jacobson 2016] and ABC [Koch et al. 2019].	42
3.18	Statistics of 5018 shells in Thingi10k dataset [Zhou and Jacobson 2016] (left) and 5545 in ABC Dataset [Koch et al. 2019] (right).	43
3.19	The <i>UV based method</i> cannot simplify the prescribed seams, and introduces self-intersections. The projection induced by the <i>Naive Cage</i> method is not continuous and not bijective; it leads to visible spikes in the reconstructed geometry.	44
3.20	Our projection is three orders of magnitude more accurate than the baseline method, and bijectively reconstructs the input vertex coordinates.	45

3.21 We attempt to simplify <i>rockerm</i> (top left) from 20088 triangles to 100. QSLim [Garland and Heckbert 1998] succeeds in reaching the target triangle count (top right) but generates an output with a self-intersection (red) and flipped triangles (purple). With our shell constraints (bottom left), the simplification stagnates at 136 triangles, but the output is free from undesirable geometric configurations. Note that both examples use the same quadratic error metric based scheduling [Garland and Heckbert 1998].	46
3.22 The heat method (right top) produces inaccurate results due to a poor triangulation (left top). We remesh the input with our method (left bottom), compute the solution of the heat method [Crane et al. 2013] on the high-quality mesh (middle bottom) and transfer the solution back to the input mesh using the bijective projection (bottom right). This process produces a result closer to the exact discrete geodesic distance [Mitchell et al. 1987] (top middle, error shown in the histograms).	47
3.23 <i>Knot</i> simplified with our method and tetrahedralized with TetGen. The original model is converted to 1,503,428 tetrahedra (left) while the simplified surface is converted to only 176,190 tetrahedra (right).	48
3.24 The union of two meshes is coarsened through our algorithm, while preserving the exact correspondence, as shown through color transfer.	49
3.25 Top: an input mesh decimated to create a coarse base mesh, and the details are encoded in a displacement map (along the normal) with correspondences computed with Phong projection [Kobbelt et al. 1998]. Bottom: the decimation is done within our shell while using the same projection as above. With our construction, this projection becomes bijective (Appendix A.4), avoiding the artifacts visible on the ears of the bunny in the top row.	50
3.26 Two volumetric chainmail textures (right) are applied to a shell (bottom left) constructed from the <i>Animal</i> mesh (top left). The original UV coordinates are transferred using our projection operator.	50
3.27 An example of a shell built around a self-intersecting mesh.	51
3.28 The <i>Armadillo</i> model with four nested cages. We create a shell from the original mesh, and then rerun our algorithm on the outer shell to create the other three layers. Note that all layers are free of self-intersections, and we have an explicit bijective map between them.	51
3.29 After optimization, the shell may self-intersect (left). Our postprocessing can be used to extract a non-selfintersecting shell, which is easier to use in downstream applications (right).	52
3.30 We generate a pinched shell for a model with a singularity (left). Optionally, we can complete the shell using our Boolean construction.	52

4.1	Our pipeline starts from a dense <i>linear</i> mesh with annotated features (green), which is converted in a curved shell filled with a high-order mesh. The region bounded by the shell is then tetrahedralized with linear elements, which are then optimized. Our output is a coarse, yet accurate, curved tetrahedral mesh ready to be used in FEM based simulation. Our construction provides a bijective map between the input surface and the boundary of the output tetrahedral mesh, which can be used to transfer attributes and boundary conditions.	55
4.2	Input triangle mesh \mathcal{M} and points \mathcal{P} . Output curved tetrahedral mesh \mathcal{T}^k and bijective map ϕ^k	62
4.3	Lagrange nodes on the reference element $\hat{\tau}$ for different $k = 1, 2, 3$ and example of geometric mapping g	63
4.4	Effect of the choice of the set \mathcal{P} on the output.	63
4.5	Our algorithm maintains free of intersection even on challenging models, without the need of setting adaptive threshold.	64
4.6	Overview of curved mesh generation pipeline.	65
4.7	A model simplified with different distances.	67
4.8	Two dimensional overview of the five steps of our boundary preserving tetrahedral meshing algorithm.	67
4.9	Input triangle mesh with features and output curved mesh with feature preserved equipped with bijective map ϕ^k	69
4.10	A sphere with different marked features (green). As we increase the number of features our algorithm will preserve them all but the quality of the surface suffers.	70
4.11	Input feature (green) is not preserved after traditional shell simplification.	70
4.12	Overview of the construction of the first stage of our pipeline.	71
4.13	The input mesh has feature edges snapped, to create a valid shell, as well as the curved mesh.	71
4.14	The input edges feature (green) are grouped together in poly-lines and categorized in graph (left) and loops (right). For every graph we add the nodes (blue) to the set of feature vertices.	72
4.15	Illustration of smoothing on a feature.	72
4.16	Curved meshes of different order. The additional degrees of freedom allows for more coarsening.	73
4.17	Relative average edge length (with respect to longest bounding box edge of each model) of our curved meshes versus number of input vertices.	74
4.18	Within the same distance bound (10^{-3} of the longest bounding box side), our method generates a coarser high order mesh, compared to the linear counterpart generated by fTetWild.	74
4.19	Surface and volume average MIPS energy of the output of our method (the CAD volume energy is truncated at 100, excluding 6 models).	75
4.20	Timing of our algorithm versus the input number of vertices.	75
4.21	Compared with Curved ODT, our method does not rely on setting vertex number and sizing field, and can generate coarse valid results.	76

4.22	Example of a BRep meshed with Gmsh where the optimization fails to untangle elements when fixing the surface. By allowing the surface to be modified, the mesh becomes “wiggly”. Our method successfully generate a positive curved mesh.	77
4.23	Example of a STEP file meshed with Gmsh where, due to the low mesh density, the tetrahedralization is not positive. Gmsh manages to generate a positive mesh by using a denser initial tessellation. Since our method starts from a dense mesh and coarsen it, it can successfully resolve the geometry.	77
4.24	Our algorithm processes triangle meshes that can be extracted from different formats: an implicit microstructure geometry from [Tozoni et al. 2020] or a subdivision surface from [Crane 2013]. The bijective map preserved on the surface allows taking advantage of the plethora of surface algorithms including polyhedral geodesic computation [Mitchell et al. 1987] and texture mapping.	78
4.25	L^2 error of the solution of the Poisson equation with respect to model size on our three datasets.	78
4.26	By meshing the region between a box and a complicated obstacle, we are able to perform non-linear fluid simulation on our curved mesh.	79
5.1	Overview of the components behind our specification. A mesh is represented through its topology, implemented by our library, and a list of user provided attributes. Before an operation is attempted, we explicitly perform a pre-check, and, if successful, we generate a mesh (attributes and topology). At this point, we trigger the after check to validate the operation (e.g., check if the newly generated mesh has positive volume). In case the after check fails, we <i>automatically</i> rollback the operation and restore the mesh to its previous <i>valid</i> state.	88
5.2	Example of the locking region for two edges. In the example the operation requires locking a two-ring neighborhood (e.g., for the edge collapse operation). If the two edges are sufficiently far (right) both operations can be safely executed in parallel. When the two edges are close (left) the operations might fail acquiring the mutexes in the shared area.	93
5.3	Comparison of our parallel implementation (32 threads) of shortest edge collapse (scalability plot on the right, from 1 to 32 threads) of a model with 281,724 faces with the serial version in libigl. Both libigl and our output have 28,168 faces and comparable edge length (1.058 for libigl versus 1.061 for ours). Our serial method runs in 4.9s (5.84s on a single thread, 0.52s with 32 thread, leading to a speedup of 11 \times), while libigl runs in 2.74s.	96
5.4	Comparison of our parallel implementation of QSlim with the serial version in libigl for a model with 1,909,755 faces. Top right, the libigl output has 17,891 faces and takes 41.26s. Bottom left, our output has 17,906 faces and runs in 306.59s. Our implementation scales well: 347.59s with one thread and 13.88s with 32 (25 \times speedup).	97

5.5	Example of uniform remeshing a model with 2,529,744 triangles (left) with the same target edge length. Middle, [Möbius and Kobbelt 2010] remeshes it to 78,322 faces in 31.96 seconds. On the right is our 32-thread implementation, which generates 71,640 triangles in 8.2 seconds (78.34s serial, 95.0s for a single thread, leading to a speedup of 11 \times). The difference in density of the meshes is due to differences in the detail of the implementation, which makes the two methods reach an average vertex valence of 5.999, and a similar target edge length (differ 0.01% of the bounding box diagonal length) with a different element budget.	98
5.6	Example of <i>Harmonic Triangulations</i> starting with one million Gaussian distributed random points. Both our and the reference implementation reach a similar target number of tetrahedra (5.9 million for the reference and 6.1 million for ours, due to a difference in operation ordering) and a similar Mean Harmonic Index (0.547 for the reference and 0.554 for ours). Our method takes 3.82s with 32 threads (15.49s serial, 40.49s on a single thread, speedup of 11 \times), while the reference serial implementation takes 6.37s.	99
5.7	Tetrahedralize a surface with 856,294 faces. Original TetWild (top right) generates a mesh with 56,761 tetrahedra in 287.58s; our reimplementation (bottom left) generates a mesh with 44,866 tetrahedra in 153.33s with 8 threads (452.42s serial, 521.47s on a single thread, speedup of 3.4 \times). The difference in number of tetrahedra is likely due to the different order of scheduling of operations due to the partitioning.	100
5.8	Example of splitting the longest edge on the bottom of the triangle. On the left, the edge can be split to generate the dark blue edge. In the next iteration, the left edge is split (by the light blue edge) leading to a decreasing maximum edge-length. On the left, the bottom edge is locked illustrated by a dashed line. In this case the next iteration splits the left edge (dark blue edge) and so on. As long as the dashed edge is locked it will never be split preventing the maximum edge length to decrease.	101
5.9	Shortest edge collapse with envelope containment of a model with 857,976 faces. Our method successfully generates a mesh with 71,298 faces in 37.49s with 32 threads (731.32s serial, 725.34s on a single thread, speedup of 20 \times).	101
5.10	Uniform remeshing with envelope containment check of a model with 198,918 faces. Our method produces a mesh with 68,202 faces in 29.11s with 32 threads and 493.54s for a single thread (483.68s for the serial version) leading to a speedup of 16 \times	102
5.11	Timings, target edge length ratio, and valence for every model in the dataset. Most models finish within a minute. The target edge length ratio measure how well our algorithm simplifies the meshes to reach the desired edge length, with an optimal value of 1. Since uniform remeshing strives to generate regular meshes, for most model our algorithm is able to obtain the optimal valence of 6.	102
5.12	Timing, max and average AMIPS energy (capped at 20) for maximum 25 iterations of tetrahedral meshing. Most models finish withing 20 minutes with only a few taking up to a day. Even by limiting the iterations to 25, most models reach an average AMIPS energy lower than 10, with optimal value at 3.	103

A.1	Bevel patterns used in the proofs in Appendix A.1.4 and Appendix A.2	107
B.1	Histogram of mean and maximum conformal AMIPS energy [Rabinovich et al. 2017] of the output of our method and TetGen.	112
B.2	Histogram of output tetrahedra number for TetGen and our method.	113

LIST OF TABLES

- | | |
|---|----|
| 2.1 Timings and statistics for the models shown in the paper. From left to right: number of input vertices and simplices, number of initial/final scaffold vertices and simplices, number of iterations, running time in seconds. The numbers in parenthesis refer to the Newton optimization. Note that our timings are considerably higher than those reported in the SLIM paper for the Lucy model since we used the reference implementation in [Jacobson et al. 2014], which does not use a multi-threaded solver. | 19 |
|---|----|

1 | INTRODUCTION

Large scale processing: ABC – Robust tetwild –

Mathematical Contributions:

Algorithmic contributions: Robust algorithms through experimental and theory.

Educational: auxiliary line or helping line is a useful tool in elementary geometry proofs. It is usually overlooked in more advanced textbooks. However, as I have showcased in several applications, educations with such tools would make sense computationally.

Outreach and impact:

1.1 PRELIMINARIES

1.1.1 BIJECTIVE MAPS

1.1.2 MATHEMATICAL OPTIMIZATION

2 | BIJECTIVE MAPS WITH TRIANGULATION AUGMENTATION

2.1 INTRODUCTION

The computation of discrete maps is a fundamental problem in computer graphics that has been extensively studied in the last three decades. The problem is challenging due to the large solution space and the non-linearity of the desired properties (in both distortion measures and constraints). Algorithms for robustly and efficiently computing locally injective (i.e. non-flipping) maps have been only recently introduced [Lipman 2012; Kovalevsky et al. 2016; Rabinovich et al. 2017] and are now having a major impact in many research areas outside of traditional texture mapping, including remeshing [Bommes et al. 2013a], image editing [Poranne and Lipman 2014], and cultural heritage [Pal et al. 2014].

In this chapter, we consider the problem of generating bijective maps, i.e. locally injective maps with non-intersecting boundaries. This is a difficult problem, exacerbated by the fact that any pair of boundary elements could overlap, leading to non-linear constraints whose number is quadratic in the size of the boundary. This problem is usually tackled by iteratively deforming an existing map, checking for overlaps after each step, and then preventing the overlap using

scaf-tex/figs/teaser.pdf

Figure 2.1: The Nefertiti model with prescribed seams is UV mapped by our algorithm. Each chart is bijective mapped into a circle or ring with Tutte’s embedding and achieves minimal distortion in less than a second. The layout is further improved interactively and the final parametrized model is shown on the right. Our approach guarantees a valid UV map with no inverted elements or overlapping triangles. See the attaching video for the optimization and manual interaction.

constraints [Harmon et al. 2011] or penalty forces [Harmon 2010]. These methods require a spatial acceleration structure to find the candidate pairs of overlapping elements and a resolution strategy that updates the map while avoiding the detected overlaps. However, the newly computed displacement might in turn lead to new overlaps, and this process has to be performed iteratively in the hope that no collisions are left. Difficult cases with many collisions might require tens or hundreds of iterations before all the candidate intersecting pairs are detected.

Our approach sidesteps the need to find candidate self-intersections based on a simple observation: if the entire ambient space (with a fixed simple boundary) is tessellated, then local injectivity implies global bijectivity [Zhang et al. 2005; Lipman 2013; Müller et al. 2015]. We thus propose an optimization framework based on this idea making it possible to leverage recent techniques for locally injective maps; our algorithm is simple to implement, robust, and two orders of magnitude faster than competing methods.

OVERVIEW. Given an input bijective map represented by a discrete triangle mesh and its mapped vertex locations, we create a new scaffold mesh for a bounding box that contains the initial, mapped triangle mesh (Figure 2.1) and conforms to its boundary. We then optimize for the desired property of the map (such as distortion, positional constraints, etc.) while ensuring that no triangle will flip. This property is achieved using a variational formulation that combines a user-defined energy for the map with a regularization term that allows the scaffold to freely deform without hindering the optimization of the map properties. During the optimization, we refine and optimize the connectivity of the scaffold mesh to prevent possible locking situations.

While a scaffold mesh has been already used in previous works [Zhang et al. 2005; Misztal and Bærentzen 2012; Müller et al. 2015], we propose to use an isometric distortion energy on the scaffold mesh, with the reference reset to the current rest pose at each iteration, and an online remeshing strategy. Our scaffold energy aims for "isometry to the current iteration", which resembles how plasticity is usually modeled in elasto-plastic simulations — this leads to a global and natural deformation of the scaffold elements that opens up space for the evolving boundary and allows for an efficient optimization using recent numerical methods for locally injective maps [Rabinovich et al. 2017].

We demonstrate the practical utility of our algorithm in the context of single patch mesh parametrization by producing distortion minimizing bijective maps for a collection of 119 challenging models. Our algorithm is ideal to compute tight UV maps for models with multiple connected components and seams as we demonstrate in our interactive texture packing experiments. Our algorithm can also be easily extended to 3D by replacing the triangular scaffold with one composed of tetrahedra. For the 3D case, we show that our method can be used to deform surfaces preventing self-intersections and to remove self-intersection from existing genus-0 surfaces when paired with a mean conformalized flow [Kazhdan et al. 2012; Sacht et al. 2013].

In the additional material, we provide a video (showing the optimization iterations) and the input/output meshes for each figure in the paper. To foster replicability of results, we will release an open-source reference implementation of our algorithm.

2.2 PREVIOUS WORK

Bijective maps find a host of applications in a variety of fields including physical simulation, surface deformation, and parametrization. We review only the most relevant prior works here and refer to the following surveys for more details [Floater and Hormann 2005; Sheffer et al. 2006; Hormann et al. 2007].

LOCALLY INJECTIVE MAPS.

There are many methods that focus on creating locally injective maps, which amounts to requiring that triangles maintain their orientation (i.e. they do not flip). In mesh parameterization, many flip-preventing metrics have been developed: the idea is to force the metric to diverge to infinity as triangles become degenerate, inhibiting flips. These metrics optimize various geometric properties such as angle [Hormann and Greiner 2000a; Degener et al. 2003] or length [Sander et al. 2001; Sorkine et al. 2002; Aigerman et al. 2014; Poranne and Lipman 2014; Smith and Schaefer 2015] preservation. Similar techniques in the context of deformation have been used to add barrier functions to enforce local injectivity in deformations [Schüller et al. 2013]. Our method uses these techniques to prevent flips in the scaffold.

Many methods have also been developed to optimize these distortion energies including moving one vertex at a time [Hormann and Greiner 2000a], parallel gradient descent [Fu et al. 2015a], as well as other quasi-newton approaches [Smith and Schaefer 2015; Kovalsky et al. 2016; Rabinovich et al. 2017]. Other approaches construct such maps by performing a change of basis, projecting to an inversion free space, and then constructing a parametrization from the result [Fu and Liu 2016]. While our method could potentially use any of these optimization methods, we use [Rabinovich et al. 2017] for its large step sizes. We elaborate on this choice in Section 2.3.1.

BIJECTIVE MAPS.

In addition to injective constraints, bijective maps have the additional requirement that the boundary does not intersect. One simple method for creating a bijective map in 2D involves constraining the boundary to a convex shape such as a circle [Tutte 1963; Floater 1997]. Such parametrizations guarantee a bijective map in 2D but create significant distortion. Even so, these methods are commonly used to create a valid starting point for further optimization [Schüller et al. 2013; Smith and Schaefer 2015; Rabinovich et al. 2017]. While methods that produce bijective maps with fixed boundaries exist [Weber and Zorin 2014; Campen et al. 2016], we aim to produce maps where the boundary is free to move to reduce the distortion of the map.

[Gotsman and Surazhsky 2001; Surazhsky and Gotsman 2001] introduced the concept of scaffolding where the free space is triangulated for the purpose of morphing without self-intersection. In [Zhang et al. 2005], the scaffold triangles are given a step function for their error: zero if not flipped, otherwise infinity. Hence, the bijective condition becomes local in that the shape can evolve until a scaffold triangle flips, in which case the free space is retriangulated and the optimization continues. The main limitation of this work is the lack of an evolving triangulation

during the line search and the absence of a rotationally invariant metric for the scaffold triangles, which lead to very small steps and an inefficient optimization.

The Deformable Simplicial Complex (DSC) method [Misztal and Bærentzen 2012] utilize a triangulation of both the free space and the interior of an object to track the interface between the two volumes. Similar to [Zhang et al. 2005], the DSC retriangulates at degeneracies but also performs operations to improve the shape of the triangles. This method changes the triangulation of the interface that it tracks, which works well for simulation, but it is not allowed in many other applications such as UV mapping.

Air meshes [Müller et al. 2015] extends the technique of Zhang et al. [Zhang et al. 2005] to add the concept of triangle flipping based on a quality measure during the optimization instead of simply retriangulating at the first sign of a degeneracy. However, this method does not maintain bijective maps as boundaries are allowed to inter-penetrate during optimization: the scaffold is only used to efficiently detect problematic regions, and the local injectivity requirement is a soft constraint in the optimization. The problem tackled in this chapter is much harder, because we do not allow any overlap during any stage of the optimization to guarantee that the resulting maps will be bijective.

[Smith and Schaefer 2015] take a different approach: instead of using a scaffold triangulation, the authors introduce a locally supported barrier function for the boundary to prevent intersection and explicitly limit the line search by computing the singularities of both the distortion energy and the boundary barrier function. Such an approach is inspired by traditional collision detection and response methods that are discussed below. Given a bijective starting point, this approach never leaves the space of bijective maps during optimization. Its main limitation is that it is computationally expensive, especially for large models. Our method is two order of magnitude faster (Figure 2.11).

COLLISION DETECTION AND RESPONSE.

While not directly related to our approach, bijective maps inherently involve some form of collision detection and response to avoid overlaps. The field on collision detection is vast, and we refer the reader to a survey [Jimenez et al. 2001]. In terms of simulations, methods such as asynchronous contact mechanics [Harmon et al. 2009; Harmon 2010; Ainsley et al. 2012] ensure the bijective property but are very expensive and designed to operate as part of a simulation. Differently, our approach is specialized for geometric optimization, where we are interested in a quasi-static solution (i.e. we do not want to explicitly simulate a dynamic system, but only find an equilibrium solution).

The work that is closer to ours in term of application (but very different in term of formulation) is [Harmon et al. 2011], where collision detection and response is used to interactively deform shapes while avoiding self-intersections. Similarly to the previous methods, the explicit detection and iterative response is expensive when many collisions happen at the same time. Our work avoids these expensive computations, and can robustly handle hundreds of simultaneous collisions while still making large steps in the optimization.

SEAM CREATION.

In the context of parametrization, some approaches optimize the connectivity of the charts of the surface during parametrization to obtain a bijective map. [Lévy et al. 2002; Zhou et al. 2004] parameterize the surface and then split charts based on whether they intersect [Lévy et al. 2002] or based on a level of distortion [Zhou et al. 2004]. Sorkine et al. [Sorkine et al. 2002] employ a bottom-up approach and add triangles to a parametrization chart until bijectivity would be violated. The problem we are solving is more general (seams are only useful for texture mapping applications) and constrained (we preserve the prescribed seams). Our algorithm could be used by these algorithms to parametrize single charts, which could reduce the number of additional seams.

2.3 METHOD

Our method, Simplicial Complex Augmentation Framework (SCAF) utilizes a scaffold structure to robustly compute a bijective map between a pair of simplicial meshes with the same connectivity. SCAF is specialized for the context of geometric optimization, i.e. we are interested in maps with low distortion and optionally satisfying a set of geometric constraints. We assume our maps are continuous and piecewise affine, i.e. the map deforms every simplex with an affine deformation. Thus, we can fully define the map using the image of its vertices.

Our algorithm uses a discrete, bijective identity map (encoded as a non-overlapping and non-flipping triangle/tetrahedral mesh) as initialization and then iteratively refines it, displacing the vertices while always ensuring that it remains bijective. Our method is composed of three stages: (1) augment the initial mesh with a scaffold, filling a bounding box around the initial map image; (2) optimize the extended mesh (scaffold included), reducing the geometric distortion of the map; (3) update the vertices and scaffold, enlarging the bounding box if necessary, to improve the quality of the triangulation. Steps (2) and (3) are iterated until the quality of the map is deemed sufficient.

2.3.1 GENERAL FORMULATION

Denote the input simplicial mesh by $\mathcal{M} = (\mathcal{V}, \mathcal{F})$ with a single, simple boundary representing a compact d -dimensional manifold embedded in the d -dimensional Euclidean space, where \mathcal{V} is the set of n -vertices and \mathcal{F} is the set of m -simplices. Our goal is to compute a continuous and piecewise affine mapping $\Phi : \mathcal{M} \rightarrow \mathbb{R}^d$ with $\Omega := \Phi(\mathcal{M}) = (\mathcal{V}', \mathcal{F}')$ the resulting simplicial mesh with the same connectivity as \mathcal{M} .

We are interested in the bijective map that minimizes a given type of geometric energy:

$$\begin{aligned} \min_{\mathcal{V}'} E_{\mathcal{M}}(\Phi) \\ \text{s.t. } \Phi \text{ is bijective,} \end{aligned} \tag{2.1}$$

where $E_{\mathcal{M}}$ is a user-defined geometric energy.

REDUCTION TO LOCAL ORIENTATION PRESERVATION. A sufficient condition for the simplicial map $\Phi : M \rightarrow \Omega$ to be bijective is that the map preserves orientation and its restriction to the boundary $\Phi|_{\partial M} : \partial M \rightarrow \partial\Omega$ is bijective [Lipman 2013]. In this light, we are able to take advantage of the following simple construction (Figure 2.2): the algorithm extends the axis aligned bounding box of M to get a d -orthotope and fills the enclosed region with a scaffold simplicial complex to form a simplicial complex mesh D that includes and conforms to $M \subset D$. We can now define the continuous and piecewise affine map $\Psi : D \rightarrow D'$ with $\Phi = \Psi|_M$ where $\Psi|_{\partial D}$ is the identity map, and denote the scaffold region as $S = D \setminus M$. Now Φ is guaranteed to be bijective if Ψ preserves orientation. Using this observation, we can translate the bijectivity into a local, orientation-preserving requirement defined per simplex.

scaf-tex/figs/scaffold_illustration.pdf

Figure 2.2: The initial mesh M (in green, left), is embedded in another mesh D (in gray, middle) that covers a box in the ambient space and contains the same triangles as M . D might contain additional points (orange). We denote the triangles that are in D but not in M as the scaffold S . Our algorithm deforms D , inducing a corresponding deformation on M (right), while keeping the boundary (blue vertices) fixed and preventing changes in the triangle orientation.

VARIATIONAL FORMULATION Minimizing the distortion of Φ using the augmented map Ψ poses an interesting challenge: what is the desired shape of the scaffold S ? Ideally we would like the simplices in the scaffold to maintain their orientation and not affect the optimization in any other way. Such a requirement is difficult to model directly, since it is a discontinuous condition that is not well-suited for the variational framework that we would like to use to minimize E_M .

We propose a regularized version of this condition modeled with an energy $E_M(\Psi|_S)$ that still diverges when elements change orientation and that mildly penalizes any non-rigid distortion.

We choose a reweighted version of the symmetric Dirichlet energy \mathcal{D} [Smith and Schaefer 2015], measured w.r.t. the Jacobian of the map Ψ computed from the rest pose of S for each simplex f ,

$$J_f := \nabla \Psi_f \quad (2.2)$$

where Ψ_f is the restriction of Ψ over the simplex f , which is an affine map. We divide the energy of each scaffold simplex f by its area A_f , and sum them up to obtain the final energy that favors

an equal contribution regardless of the size of each scaffold simplex:

$$\begin{aligned} E_S(\Psi|_S) &= \sum_{f \in S} \frac{1}{A_f} \mathcal{D}(J_f) \\ &= \sum_{f \in S} (\|J_f\|_F^2 + \|J_f^{-1}\|_F^2 - 2d). \end{aligned} \quad (2.3)$$

The $-2d$ term ensures that the energy is 0 when $J_f = \mathbb{I}$. The map is then computed by summing the two terms:

$$\begin{aligned} \min_D \quad & E_M(\Psi|_M) + \lambda E_S(\Psi|_S) \\ \text{s.t.} \quad & \Psi|_{\partial D} \text{ is Identity} \\ & \Psi \text{ preserves orientation.} \end{aligned} \quad (2.4)$$

where $\lambda > 0$ is balancing the contribution of the two energies, decreasing as the optimization proceeds.

ITERATIVE REGULARIZATION Solving this problem leads to a bijective and distortion minimizing map, but the regularizer will affect the stationary points of E_M , which is problematic, especially for large deformations. To address this problem we iteratively minimize this energy, regenerating the scaffold at each iteration, and use the new scaffold as a rest pose for the regularization term $E_S(\Psi|_S)$. This iterative procedure has two positive effects: (1) it acts as a proximal regularization term without inhibiting movement since the rest pose is updated at each iteration; (2) the meshing quality of the scaffold is high, which avoids locking configurations.

INTERPOLATION COEFFICIENT. We experimentally observed that our algorithm is robust to different choices of λ , generating indistinguishable results in most cases. However, λ affects the convergence speed (Figure 2.3). We used $\lambda = \frac{1}{100} \frac{E_M(\Psi|_M)}{|S|}$ for all our experiments, where $|S|$ is the number of scaffold simplices.

SOLVER. Since our energy is rotational invariant, we can minimize the energy with the same quadratic proxy proposed in SLIM [Rabinovich et al. 2017], enriching the approach with the equality constraint needed to fix the boundary ∂D . We also employ the orientation-preserving line search [Smith and Schaefer 2015] (with exact predicates [Shewchuk 1996a]) to ensure that no triangles can change orientation. Alternatively, other methods such as AQP [Kovalsky et al. 2016] could be used to minimize this energy. Since our approach changes the mesh connectivity at every iteration, AQP loses much of its advantages as the approximate Hessian must be recomputed at each iteration and not prefactored. Therefore, our approach is an ideal fit for [Rabinovich et al. 2017], which takes large steps at every iteration without relying on a constant, prefactored matrix at each iteration. For practical applications, SLIM iterations are sufficient to minimize the energy to acceptable levels. For two stress tests (figures 2.4 and 2.11) we used the result of SLIM as a warm start for a Newton optimization (as suggested by [Rabinovich et al. 2017]), which quickly converges to a numerical minimum.

scaf-tex/figs/hand_weight.pdf

Figure 2.3: Different values of λ do not affect the result, but they change the number of iterations needed. From left to right: we used a large weight (100x ours), our weight, and a small weight (0.01x ours). The optimization took 9, 7, and 8 iterations, respectively, to reach the same energy level.

2.3.2 SURFACE PARAMETRIZATION

Our framework can be used for many applications, one of which is computing a bijective surface parametrization from a 3D surface into the UV plane. We follow [Liu et al. 2008] and assume that each 3D triangle f^{3D} is equipped with a rigid transformation R_f such that applying R_f to f^{3D} maps f^{3D} to the plane. Given this transformed triangle f , we can now measure the distortion of the map using the Jacobian of the affine transformation (a 2×2 matrix) from $R_f(f^{3D})$ to f , the location of the triangle in the parametrization.

INITIALIZATION Our method is initialized with Tutte’s embedding algorithm [Tutte 1963]:

$$\Phi^0 : \mathcal{M}^{3D} \rightarrow \Omega^0,$$

where Ω^0 is a simplicial disk domain. Then we construct a larger rectangular domain $D^0 \supset \Omega^0$, where ∂D^0 is an axis-aligned rectangle, and use Triangle [Shewchuk 1996a] to triangulate the region in between. We enforce a quality bound of 20° to obtain a graded mesh that is coarse on the boundary and conforming the boundary of Ω^0 . This grading implicitly produces an approximate inverse distance weighting of the scaffold space with respect to the error function, which enables a larger deformation per iteration. Then we define $\Psi : D^0 \rightarrow D \subset \mathbb{R}^2$ and restrict $\Psi|_{\partial D^0}$ to be the identity.

MESH IMPROVEMENT At the end of each iteration, we improve the quality of the scaffold. The reason for maintaining a good mesh quality is two-fold. First, as observed in [Zhang et al. 2005; Müller et al. 2015], fixing the scaffold will potentially prevent movement. Secondly, the quality of the scaffold affects the condition of the linear system in SLIM [Rabinovich et al. 2017]; a higher quality leads to larger and more efficient iterations.

We resort to Triangle [Shewchuk 1996a] to create the initial scaffold and to regenerate the scaffold mesh in the improvement step. Our experiments show that, in 2D, it is faster to generate the mesh from scratch at every iteration instead of trying to optimize the scaffold using local operations as suggested in [Müller et al. 2015]. Since our solver makes large steps in each iteration, the scaffold requires significant connectivity changes each iteration, which explains why regenerating the triangulation is faster than local operations. This is in stark contrast with physical simulation scenarios where each iteration represents a small time step and, thus, a minor change in the vertex positions.

We demonstrate the effectiveness of the remeshing strategy in Figure 2.4 where our method recovers from a large rotation — note that the scaffold is updated during the iterations and always leaves space for the map to move freely.

scaf-tex/figs/coil-remeshing.pdf

Figure 2.4: A bijective map from a circle (left) to a spiral (right) is computed without (top) and with (bottom) the iterative remeshing step. The slivers in the triangulation locks the optimization (top), preventing it from reaching the target shape.

SLIDING & DEGENERACY PREVENTION As the optimization proceeds and some of the boundary elements get closer, some of the scaffold triangles might (and often will) get smaller and smaller, restricting the amount of sliding that is allowed in one iteration as well as introducing numerical difficulties in computing the corresponding Jacobian J_f whose singular values will approach infinity.

To avoid this issue, we replace the degenerating target when computing its Jacobian. For the triangles with an area smaller than ϵ , we use an equilateral triangle with area ϵ to compute the local Jacobian. In our experiment, we traverse through the boundary of the interior of the uv domain at the current iteration to find the minimum edge length l and set $\epsilon = \frac{l^2}{4}$.

A theoretical downside of this modification is that it affects the distortion energy. We experimentally observed that the changes are negligible, and we thus used it for all our experiments. However, on the practical side, it discourages fully degenerate elements — this change, coupled with the orientation preserving line search [Smith and Schaefer 2015], makes our algorithm robust enough for the challenging stress tests shown in Figure 2.5.

2.3.3 EXTENSION TO 3D

Our formulation naturally unifies bijective geometric optimization problems in 2D and 3D, so the algorithm readily extends to the 3D case with only one major difference: the scaffold becomes a tetrahedral mesh, which is computationally more challenging to create and update.

MESH IMPROVEMENT We use TetGen [Si 2015a] to generate the initial scaffold and the local operations proposed in [Klingner 2009] to optimize the scaffold’s quality in the subsequent iterations. It is unfortunately not possible to directly use TetGen at every iteration as we did with Triangle in the 2D case, since TetGen fails when boundaries get too close, which is common in our experiments.

GUARANTEE Similarly to the 2D case, we are guaranteed that no flipping or self-intersecting tetrahedra will occur since we are following an interior point strategy. Notice the contrast with Air Mesh [Müller et al. 2015] where only if penetration happens can the constraints be of effect. However, as pointed out in [Dougherty et al. 2004], the local operations we are performing may not be sufficient to explore the entire space of possible tetrahedralizations. Therefore, we cannot guarantee that the algorithm achieves the globally optimal solution.

2.4 RESULTS

We implemented our algorithm in C++ using Eigen for linear algebra routines. We ran our experiments on a desktop with a 4-core Intel i7 processor clocked at 4 GHz and 32 GB of memory but using only one thread on a single core. For all experiments, the scaffold bounding box is computed by uniformly scaling by three times the bounding box of the image of the current map.

ROBUSTNESS. To demonstrate the robustness of our algorithm, we computed bijective maps for all the 102 meshes parametrized by the MIQ algorithm [Bommes et al. 2009] and for the 17 meshes parametrized by [Myles et al. 2014] in the dataset proposed by [Myles et al. 2014]. The cuts in these meshes have been designed for locally injective parametrization that usually have major self-overlap. We use them as a stress test for the effectiveness and robustness of our method: the

scaf-tex/figs/database_demo.pdf

Figure 2.5: Two models are cut using [Bommes et al. 2009] and bijectively parametrized using our algorithm. See the additional material for more examples.

cuts introduce a massive distortion in the Tutte initialization and lead to boundaries that are prone to overlap in hundreds of locations. Our method successfully creates bijective parametrizations for all these models with default parameters. We attach all the parametrized models in the additional material and show two examples in Figure 2.5.

SCALABILITY. Our methods scales gracefully to large datasets, similarly to [Rabinovich et al. 2017]. We repeat their scalability experiment, but producing bijective maps instead of just locally injective maps (Figure 2.6). The behaviour is remarkably similar — the density of the model (and consequently of the scaffold) does not affect the number of required iterations.

TEXTURE ATLAS GENERATION UV mapping is a time consuming procedure required in most geometric modeling pipelines. Existing commercial tools provide the ability to flatten single patches and arrange them in UV layouts where multiple patches are tightly packed inside a rectangular domain, which is then loaded in the texture memory of a GPU.

Our algorithm can bijectively parameterize a single patch (Figure 2.7), avoiding the typical manual UV postprocessing required with traditional tools. Our algorithm can also be used to create automatic UV charts of models with multiple connected components (or predefined cut edges). We show an example in Figure 2.8 where we detected the connected components, bijectively map the patches into a set of circles (using a grid layout), and reduce their distortion using our algorithm. The result is a tight and automatic packing without resorting to any user-interaction. Additional interactive tools can further improve the atlas by dragging&dropping regions or translating islands

scaf-tex/figs/lucy-scalability.pdf

Figure 2.6: We compare the distortion energy with respect to the number of iterations on a set of Lucy’s meshes with different resolutions (from 1 to 12 million faces). In the center of the plot, we show the 1M Lucy model parametrized by our algorithm.

while ensuring that no overlaps are introduced (Figure 2.1). We show interactive sessions using our packing tool in the additional material.

PREVENTING SELF-INTERSECTIONS Our algorithm can be generalized to handle mixed dimension problems, such as the deformation of 2D surface in 3D space, while preventing self-intersections. In Figure 2.9, we demonstrate the use of our method to resolve self-intersections of surfaces. First we perform a conformalized flow [Kazhdan et al. 2012] using the algorithm proposed in [Sacht et al. 2013] to resolve any self-intersections. While [Sacht et al. 2013] will resolve the intersections, the resulting surface may be geometrically far from the initial shape (see Figure 2.9). Next we tetrahedralize the ambient space while conforming to the deformed surface mesh and minimize Equation 2.4 with an additional energy term that strives to restore the rest pose geometry of the surface, using the surface ARAP energy proposed in [Sorkine and Alexa 2007]. The result

scaf-tex/figs/the_animal.pdf

Figure 2.7: A mesh is cut by an artist into a single chart and parametrized using SLIM [Rabinovich et al. 2017] (left) and with our algorithm (right). Note that local-injectivity is not sufficient for this model, since the global overlaps in the highlighted region prevent this parametrization from being a UV texture map. Our result (right) is guaranteed to be bijective.

is a surface similar to the original mesh, but without self-intersections. In this example, it is possible to observe that even dramatic changes of scale (on the foot) can be robustly handled by our parametrization algorithm.

A more challenging stress test is shown in Figure 2.10, where the bunny model is scaled up inside a box, to 30 times its original size. No self-intersections are introduced, despite the extreme, constrained deformation.

COMPARISON WITH [SMITH AND SCHAEFER 2015] The algorithm closest to ours is [Smith and Schaefer 2015], which tackles a similar problem (restricted to the 2D case). We replicated the space filling curve experiment and obtained remarkably similar results, where our running time is 96s, compared with 8,472s for [Smith and Schaefer 2015] (88 times faster). We show in Figure 2.11 a more challenging experiment with a subdivided version of the space filling curve to emphasize the performance difference: our algorithm converges in 39 minutes, while [Smith and Schaefer 2015] did not converge after 5 days and 21 hours. For this example, we used the procedure suggested in [Rabinovich et al. 2017]: we performed a few iterations minimizing the quadratic proxy and then switch to a traditional newton method until numerical convergence. A video of the optimization is provided in the additional material.

Our method produces results that are visually identical to [Smith and Schaefer 2015]. In Figure 2.12 we repeat the experiments shown in [Smith and Schaefer 2015], stopping our optimization at the same energy value.

LOCAL VS GLOBAL OPTIMIZATION Both [Zhang et al. 2005] and [Misztal and Bærentzen 2012] use a construction similar to ours to generate bijective maps (Section 2.2). Both methods explicitly prevent changes of orientation using a local approach: they optimize the map using coordinate descent iterations [Solomon 2015] allowing only one vertex at a time to move in its 1-ring and thus ensuring that no triangle flip. This strategy severely limits the maximal displacement per iteration and restricts the step to the size of the 1-rings. Such a restriction makes these methods impractical for parametrization applications since the difference in scale between the Tutte’s

scaf-tex/figs/maneki_neko_colorful.pdf

Figure 2.8: A model with multiple chart (left) is automatically parametrized in a texture atlas (bottom-right) by first mapping each component to a circle (top-right) and then minimizing the distortion.

embedding and the final result is extreme (the ratio of min and max triangle area is 10^{-6} in Figure 2.13). We show an example of one of our iterations in Figure 2.13, where the highlighted vertex traversed a distance of 150 times the size of the average edge length of its 1-ring in one single step. Using coordinate descent would have required hundreds of iterations to achieve the same effect.

Despite the orientation-dependent box used as scaffold boundary, our optimization produces results that are, in practice, independent of orientation. We show this effect in Figure 2.14 where we initialize the optimization with 1000 randomly rotated Tutte's mappings of the same camel model and run our optimization. The isometric distortion of the model after 50 iterations is quite similar in all trials (the minimum, maximum, average, standard deviation of distortion errors in all 1000 runs are 0.1086, 0.1107, 0.1095, 3.2698e-4 resp.) indicating very little change based on the initial orientation.

scaf-tex/figs/leg-flow.pdf

Figure 2.9: We remove the self-intersections from a genus 0 model using the conformalized flow [Kazhdan et al. 2012; Sacht et al. 2013]. The flow is inverted, while using our algorithm to compute a bijective volumetric map, to recover a self-intersection free version of the original surface. The final model can now be meshed using TetGen, since it is free from self-intersections.

scaf-tex/figs/rabbit_grow.pdf

Figure 2.10: We grow a bunny inside a box, while preventing self-intersections. We show the result after 0,10,20,30,40, and 50 iterations.

TIMINGS The timings for all the results in the paper are reported in Table 2.1.

2.5 LIMITATIONS AND CONCLUDING REMARKS

We proposed a simple and robust algorithm to generate bijective maps, both in 2D and 3D. We demonstrated the practical value of the algorithm in UV mapping and deformation applications, and its robustness with extensive stress tests.

One major venue for future work is the support of hard positional constraints, which are favored over soft constraints in many practical applications. Our current algorithm only supports

scaf-tex/figs/dense_hilbert.pdf

Figure 2.11: We repeat the challenging test in [Smith and Schaefer 2015] with a subdivided version of their Hilbert curve to increase the triangle count. Our method starts from a disc (upper left), gracefully extends (upper right), and reaches the same minimum (lower left) in 39 minutes whereas [Smith and Schaefer 2015] didn't terminate more than 5 days (lower right), highlighting our performance boost of over 200 times.

soft constraints as geometric energy [Schüller et al. 2013]. To support hard constraints we would need to generate a bijective starting point that guarantees those constraints, and then preserve them in our optimization. While bijective maps with hard constraints can be constructed for a 2D patch homeomorphic to a disk [Weber and Zorin 2014] and for a 3D volume homeomorphic to a ball [Campen et al. 2016], the generic solution is still elusive.

In 3D cases, the generation of the initial scaffold is not as robust as in 2D, since TetGen fails for geometries with self-intersections and other imperfections. Our algorithm is also slower in 3D due to larger and denser linear systems, as well as the need for local mesh refinement operations instead of regenerating the entire tetrahedralization. We believe a more optimized and parallel implementation could reduce this overhead, and plan to explore this in the future.

scaf-tex/figs/compare_smith.pdf

Figure 2.12: We apply our algorithm on 4 models used in [Smith and Schaefer 2015] (using the same stopping criteria) obtaining visually identical results. Distortion errors produced by our algorithm (outer) and theirs (inner) are shown in black.

scaf-tex/figs/camel_step.pdf

Figure 2.13: A single iteration of our algorithm (from left to right) drastically reduces the distortion. The black vector in the center is 150 times longer than the average edge length of its 1-ring. Iterative methods would need thousands of iterations to achieve a similar progress.

Type	Model	#V	#F	#Vs	#Fs	It.	Total Time (s)	It. Time (s)
Atlas	Nefertiti (Fig. 2.1)	1697	2823	983/ 247	1945/ 728	50	0.71	0.01
	Maneki-Neko (Fig. 2.8)	23025	43648	2427/ 725	7174/ 3770	50	16.81	0.34
	Hand (Fig. 2.3)	2239	4046	347/280	1104/970	7	0.14	0.02
	Spiral (Fig. 2.4)	54	52	78/36	190/106	50(50)	0.04(0.21)	0.01
	Thai Statue (Fig. 2.5, left)	42405	79970	3665/1593	12148/8004	50	28.28	0.56
	Filigree (Fig. 2.5, right)	56062	100000	9160/2627	30422/17356	100	75.99	0.76
	Lucy (Fig. 2.6)	501105	1000000	1856/ 3470	5900/ 5674	100	2524.22	25.24
	Lucy (Fig. 2.6)	1001375	1999999	2284/ 4400	7297/ 7133	100	7251.00	72.51
	Lucy (Fig. 2.6)	2002031	3999999	3587/ 6930	11215/ 10985	100	22500.07	225.00
	Lucy (Fig. 2.6)	3002899	5999999	5135/ 9859	16047/ 15601	100	52235.31	522.35
2D	Lucy (Fig. 2.6)	4002816	8000000	5140/ 10288	15890/ 15918	100	59413.14	594.13
	Lucy (Fig. 2.6)	5003408	10000000	6194/ 12231	19182/ 19040	100	95247.59	952.47
	Lucy (Fig. 2.6)	6004111	12000000	7357/6418	2291/21036	50	78726.05	1574.52
	Animal (Fig. 2.7)	19937	39040	747/593	2306/1998	50	15.36	0.31
	Space Filling (Fig. 2.11)	79545	146832	90815/88237	181608/176452	200(250)	547.13(1836.58)	5.30
	Horse (Fig. 2.12)	20636	39698	1343/984	4238/3520	30(10)	8.26(12.03)	0.28(1.20)
	Camel (Fig. 2.12)	2032	3576	384/272	1234/1010	30(10)	0.52(1.13)	0.02(0.11)
	Cow (Fig. 2.12)	3195	5804	491/277	1546/1118	30(10)	0.81(1.74)	0.03(0.17)
	Tricera (Fig. 2.12)	3163	5660	544/329	1732/1302	30(10)	0.83(1.77)	0.03(0.18)
	3D	Leg (Fig. 2.9)	6617	13230	5016/5021	68521/68544	500	3251.17
	Bunny (Fig. 2.10)	568	1132	683/706	6209/6289	50	7.16	0.14

Table 2.1: Timings and statistics for the models shown in the paper. From left to right: number of input vertices and simplices, number of initial/final scaffold vertices and simplices, number of iterations, running time in seconds. The numbers in parenthesis refer to the Newton optimization. Note that our timings are considerably higher than those reported in the SLIM paper for the Lucy model since we used the reference implementation in [Jacobson et al. 2014], which does not use a multi-threaded solver.

scaf-tex/figs/random-rotate.pdf

Figure 2.14: Our algorithm is independent to the initial orientation. We rotate the initializing Tutte's mapping of the camel model and obtain results with similar isometric distortion.

3 | BIJECTIVE PROJECTION IN THE SHELL

3.1 INTRODUCTION

Triangular meshes are the most popular representation for discrete surfaces, due to their flexibility, efficiency, and direct support in rasterization hardware. Different applications demand different meshes, ranging from extremely coarse for collision proxies, to high-resolution and high-quality for accurate physical simulation. For this reason, the adaptation of a triangle mesh to a specific set of criteria (surface remeshing) is a core building block in geometry processing, graphics, physical simulation, and scientific computing.

In most applications, the triangular mesh is equipped with attributes, such as textures, displacements, physical properties, and boundary conditions (Figure 3.1). Whenever remeshing is needed, these properties must be transferred on the new mesh, a task which has been extensively studied in the literature and for which robust and generic solutions are still lacking (Section 3.2). Defining a continuous bijective map, more precisely, a homeomorphism where the inverse is also continuous, between two geometrically close piecewise-linear meshes of the same topology is a difficult problem, even in its basic form, when one of these meshes is obtained by adapting the other in some way (e.g., coarsening, refining, or improving triangle shape). Common approaches to this problem are Euclidean projection [Jiao and Heath 2004], parametrization on a common domain [Praun et al. 2001; Kraevoy and Sheffer 2004; Lee et al. 1998], functional maps [Ovsjanikov et al. 2012], and generalized barycentric coordinates [Hormann and Sukumar 2017]. However, the problem is not fully solved, as all existing methods, as we discuss in greater detail in Section 3.2, often fail to achieve bijectivity and/or sufficient quality of the resulting maps when applied to complex geometries. Our focus is on correspondences between meshes obtained during a remeshing procedure, instead of solving the more general problem of processing arbitrary mesh pairs.

In this work, we propose a general construction designed to enable attribute mapping between geometrically close (in a well-defined sense) meshes by jointly constructing: (1) a shell \mathcal{S} around triangle mesh \mathcal{T} spanned by a set of prisms, inducing a volumetric vector field \mathcal{V} in its interior and (2) a projection operator \mathcal{P} that bijectively maps surfaces inside the shell to \mathcal{T} , as long as the dot product of the surface face normals and \mathcal{V} is positive (we call such a surface a *section* of \mathcal{S}). Given a surface mesh \mathcal{T} and its shell \mathcal{S} , it is now possible to exploit the bijection induced by \mathcal{P} in many existing remeshing algorithms by adding to them an additional constraint ensuring that the generated surface is a section of a given shell.

As long as the generated mesh is a section, the projection operator \mathcal{P} can be used to transfer application-specific attributes. At a higher level, the middle surface of our shell can be seen as

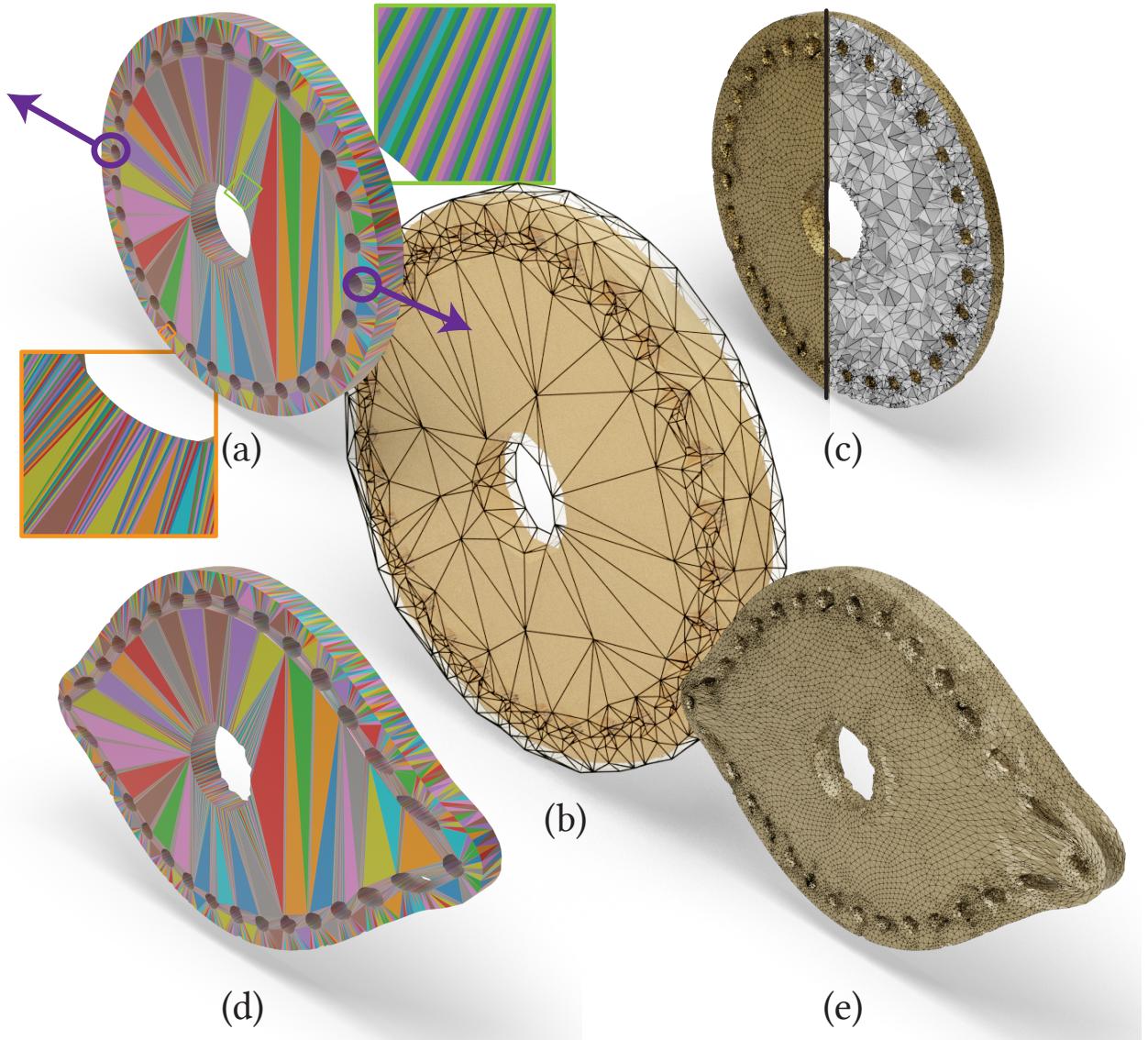


Figure 3.1: A low-quality mesh with boundary conditions (a) is remeshed using our shell (b) to maintain a bijection between the input and the remeshed output. The boundary conditions (arrows in (a)) are then transferred to the high-quality surface (c), and a non-linear elastic deformation is computed on a volumetric mesh created with TetGen (e). The solution is finally transferred back to the original geometry (d). Note that in this application setting both surface and volumetric meshing can be hidden from the user, who directly specifies boundary conditions and analyses the result on the input geometry.

a common parametrization domain shared by sections within the shell: differently from other methods that map the triangle meshes to a disk, region of a plane, a canonical polyhedron, or orbifolds, our construction uses an explicit triangle mesh embedded in ambient space as the common parametrization domain. This provides additional flexibility since it is adaptive to the density of the mesh and naturally handles models with high genus, while being numerically stable under floating-point representation (and exact if evaluated with rational arithmetic). The downside is that it is defined only for sections contained within the shell. The construction and optimization of our shell, and corresponding bijective mapping, is computationally more expensive than remeshing-only methods: our algorithms takes seconds to minutes on small and medium sized models, and might take hours on the large models in our tests.

We evaluate the robustness of the proposed approach by constructing shells for a subset of the models in Thingi10k [Zhou and Jacobson 2016] and in ABC [Koch et al. 2019] (Section 3.4). We also integrate it in six common geometry processing algorithms to demonstrate its practical applicability (Section 3.5):

1. *Proxy*. The creation of a proxy, high-quality remeshed surface to solve PDEs (e.g., to compute geodesic distances or deformations), avoiding the numerical problems caused by a low-quality input in commonly used codes. Bijective projection operators associated with a shell enable us to transfer boundary conditions to the proxy mesh, compute the solution on the proxy, and then transfer the solution back to the original geometry.
2. *Boolean operations*. The remeshing of intermediate results of Boolean operations, to ensure high-quality intermediate meshes while preserving a bijection to transfer properties between them.
3. *Displacement Mapping*. The automatic conversion of a dense mesh into a coarse approximation and a regularly sampled displacement height map. Our method generates a bijection that allows us to bake the geometric details in a displacement map.
4. *Tetrahedral Meshing*. The conversion of a surface mesh of low quality into a high-quality tetrahedral mesh, with bijective correspondence.
5. *Geometric Textures*. Generation of complex topological structures using volumetric textures mapped to the volumetric parametrization of a simplified shell defined by \mathcal{P} . Our analysis on the initial shell also complements the literature on shell maps.
6. *Nested Cages*. A robust approach to generate a coarse approximation of a surface for collision checking, cage-based deformation, or multigrid approaches.

Our contributions are:

1. An algorithm to build a prismatic shell and the corresponding projection operator around an orientable, manifold, self-intersection free triangle mesh with arbitrary quality;
2. A new definition of bijective maps between *close-by* discrete surfaces;
3. A reusable, reference implementation provided at <https://github.com/jiangzhongshi/bijective-projection-shell>.

3.2 RELATED WORKS

We review works in computer graphics spanning both the realization of maps for attribute transfer (Section 3.2.1), and the explicit generation of boundary cages (Section 3.2.2), which are closest to our work.

3.2.1 ATTRIBUTE TRANSFER

Transferring attributes is a common task in computer graphics to map colors, normals, or displacements on discrete geometries. The problem is deeply connected with the generation of UV maps, which are piecewise maps that allow to transfer attributes from the planes to surfaces (and composition of a UV map with an inverse may allow transfer between surfaces). We refer to [Floater and Hormann 2005; Sheffer et al. 2006; Hormann et al. 2007] for a complete overview, and we review here only the most relevant works.

PROJECTION Modifying the normal field of a surface has roots in computer graphics for Phong illumination [Phong 1975], and tessellation [Boubekeur and Alexa 2008]. Orthographic, spherical, and cage based projections are commonly used to transfer attributes, even if they often leads to artifacts, due to their simplicity [Community 2018; Nguyen 2007]. Projections along a continuously-varying normal field has been used to define correspondences between neighbouring surfaces [Kobbelt et al. 1998; Lee et al. 2000; Panizzo et al. 2013; Ezuz et al. 2019], but it is often discontinuous and non-bijective. While the discontinuities are tolerable for certain graphics applications (and they can be reduced by manually editing the cage), these approaches are not usable in cases where the procedure needs to be automated (batch processing of datasets) or when bijectivity is required (e.g., transfer of boundary conditions for finite element simulation). These types of projection may be useful for some remeshing applications to eliminate surface details [Ebke et al. 2014], but it makes these approaches not practical for reliably transferring attributes. Our shell construction, algorithms, and associated projection operator, can be viewed as guaranteed continuous bijective projection along a field.

COMMON DOMAINS A different approach to transfer attributes is to map both the source and the target to a common parametrization domain, and to compose the parametrization of the source domain with the inverse parametrization of the target domain to define a map from source to target. In the literature, there are methods that map triangular meshes to disks [Tutte 1963; Floater 1997], region of a plane [Maron et al. 2017; Aigerman et al. 2015, 2014; Schüller et al. 2013; Smith and Schaefer 2015; Rabinovich et al. 2017; Jiang et al. 2017; Weber and Zorin 2014; Campen et al. 2016; Müller et al. 2015; Gotsman and Surazhsky 2001; Surazhsky and Gotsman 2001; Zhang et al. 2005; Fu and Liu 2016; Litke et al. 2005; Schmidt et al. 2019], a canonical coarse polyhedra [Kraevoy and Sheffer 2004; Praun et al. 2001], orbifolds [Aigerman and Lipman 2015; Aigerman et al. 2017; Aigerman and Lipman 2016], Poincare disk [Springborn et al. 2008; Stephenson 2005; Kharevych et al. 2006; Jin et al. 2008], spectral basis [Ovsjanikov et al. 2012; Shoham et al. 2019; Ovsjanikov et al. 2017], and abstract domains [Kraevoy and Sheffer 2004; Schreiner et al. 2004; Pietroni et al.

2010]. While these approaches allow mappings between completely different surfaces, this is a hard problem to tackle in full generality fully automatically, with guarantees on the output (even some instances of the problem of global parametrization, i.e. maps from a specific type of almost everywhere flat domains to surfaces, lack a fully robust automatic solution).

Our approach uses a coarse triangular domain embedded in ambient space as the parametrization domain, and uses a vector field-aligned projection within an envelope to parametrize close-by surfaces bijectively to the coarse triangular domain. Compared to the methods listed above, our approach has both pros and cons. Its limitation is that it can only bijectively map surfaces that are similar to the domain, but on the positive side, it: (1) is efficient to evaluate, (2) guarantees an exact bijection (it is closed under rational computation), (3) works on complex, high-genus models, even with low-quality triangulations, (4) less likely to suffer from high distortion (and the related numerical problems associated with it), often introduced by the above methods. We see our method not as a replacement for the fully general surface-to-surface maps (since it cannot map surfaces with large geometric differences), but as a complement designed to work robustly and automatically for the specific case of close surfaces, which is common in many geometry processing algorithms, as well as serve as a foundation for generating such close surfaces (e.g., surface simplification and improvement, see Section 3.5)

ATTRIBUTE TRACKING In the specific context of remeshing or mesh optimization, algorithms have been proposed to explicitly track properties defined on the surface [Garland and Heckbert 1997; Cohen et al. 1997; Dunyach et al. 2013] after every local operation. By following the operations in reverse order, it is possible to resample the attributes defined on the input surface. These methods are algorithm specific, and provide limited control over the distortion introduced in the mapping. Our algorithm provides a generic tool that enables any remeshing technique to obtain such a map with minimal modifications.

3.2.2 SHELL GENERATION

The generation of shells (boundary layer meshes) around triangle meshes has been studied in graphics and scientific computing.

ENVELOPES Explicit [Cohen et al. 1996, 1997] or implicit [Hu et al. 2016] envelopes have been used to control geometric error in planar [Hu et al. 2019b], surface [Guéziec 1996; Hu et al. 2017; Cheng et al. 2019], and volumetric [Hu et al. 2018, 2019d] remeshing algorithms. Our shells can be similarly used to control the geometric error introduced during remeshing, but they offer the advantage of providing a bijection between the two surfaces, enabling to transfer attributes between them without explicit tracking [Cohen et al. 1997]. We show examples of both surface and volumetric remeshing in Section 3.5. Also, [Barnhill et al. 1992; Bajaj et al. 2002] utilize envelopes for function interpolation and reconstruction, where our optimized shells can be used for similar purposes.

SHELL MAPS 2.5D geometric textures, defined around a surface, are commonly used in rendering applications [Wang et al. 2003, 2004; Porumbescu et al. 2005; Peng et al. 2004; Lengyel et al. 2001; Chen et al. 2004; Huang et al. 2007; Jin et al. 2019]. The requirement is to have a thin shell around the surface that can be used to map an explicit mesh copied from a texture, or a volumetric density field used for ray marching. Our shells are naturally equipped with a 2.5D parametrization that can be used for these purposes, and have the advantage of allowing users to generate coarse shells which are efficient to evaluate in real-time. The bijectivity of our map ensures that the volumetric texture is mapped in the shell without discontinuities. We show one example in Section 3.5.

BOUNDARY LAYER Boundary layers are commonly used in computational fluid dynamics simulations requiring highly anisotropic meshing close to the boundary of objects. Their generation is considered challenging [Aubry et al. 2015, 2017; Garimella and Shephard 2000]. These methods generate shells around a given surface, but do not provide a bijective map suitable for attribute transfer.

COLLISION AND ANIMATION Converting triangle meshes into coarse cages is useful for many applications in graphics [Sacht et al. 2015], including proxies for collision detection [Calderon and Boubekeur 2017] and animation cages [Thiery et al. 2012]. While not designed for this application, our shells can be computed recursively to create increasingly coarse nested cages. We hypothesize that a bijective map defined between all surfaces of the nested cages could be used to transfer forces from the cages to the object (for collision proxies), or to transfer handle selections (for animation cages). [Botsch et al. 2006; Botsch and Kobbelt 2003] uses a prismatic layer to define volumetric deformation energy, however their prisms are disconnected and only used to measure distortion. Our prisms could be used for a similar purpose since they explicitly tessellate a shell around the input surface.

3.2.3 ROBUST GEOMETRY PROCESSING

The closest works, in terms of applications, to our contribution are the recent algorithms enabling black-box geometry processing pipelines to solve PDEs on meshes *in the wild*.

[Dyer et al. 2007; Liu et al. 2015] refines arbitrary triangle meshes to satisfy the Delaunay mesh condition, benefiting the numerical stability of some surface based geometry processing algorithms. These algorithms are orders of magnitude faster than our pipeline, but, since they are refinement methods, cannot coarsen dense input models. While targeting a different application, [Sharp et al. 2019] offers an alternative solution, which is more efficient than the extrinsic techniques [Liu et al. 2015] since it avoids the realization of the extrinsic mesh (thus naturally maintaining the correspondence to the input, but limiting its applicability to non-volumetric problems) and it alleviates the introduction of additional degrees of freedom. [Sharp and Crane 2020] further generalizes [Sharp et al. 2019] to handle non-manifold and non-orientable inputs, which our approach currently does not support.

TetWild [Hu et al. 2018, 2019d] can robustly convert triangle soups into high-quality tetrahedral meshes, suitable for FEM analysis. Their approach does not provide a way to transfer boundary

prism-tex/figs/pipeline.pdf

Figure 3.2: Overview of our algorithm. We start from a triangle mesh, find directions of extrusion, build the shell, and optimize to simplify it.

conditions from the input surface to the boundary of the tetrahedral mesh. Our approach, when combined with a tetrahedral mesher that does not modify the boundary, enables to remesh low-quality surface, create a tetrahedral mesh, solve a PDE, and transfer back the solution (Figure 3.1). However, our method does not support triangle soups, and it is limited to manifold and orientable surfaces.

3.2.4 ISOTOPY BETWEEN SURFACES

[Chazal and Cohen-Steiner 2005; Chazal et al. 2010] presents conditions for two sufficiently smooth surfaces to be isotopic. Specifically, the projection operator is a homeomorphism. [Mandad et al. 2015a] extends this idea to make an approximation mesh that is isotopic to a region. However, they did not realize a map suitable for transferring attributes.

3.3 METHOD

Our algorithm (Figure 3.2) converts a self-intersection free, orientable, manifold triangle mesh $\mathcal{T} = \{V_{\mathcal{T}}, F_{\mathcal{T}}\}$, where $V_{\mathcal{T}}$ are the vertex coordinates and $F_{\mathcal{T}}$ the connectivity of the mesh, into a shell composed of generalized prisms $\mathcal{S} = \{(B_S, M_S, T_S), F_S\}$, where B_S, M_S, T_S are bottom, middle, and top surfaces of the shell, consisting of bottom, middle, and top triangles of the prisms, and F_S is the connectivity of the prisms (Figure 3.3). The algorithm initially generates a shell \mathcal{S} whose middle surface M_S has the same geometry as the input surface \mathcal{T} (possibly with refined connectivity), and then optimizes it while ensuring that \mathcal{T} is contained inside and projects bijectively to M_S . The shell induces a volumetric vector field \mathcal{V} and a projection operator \mathcal{P} in the interior of each of its prisms (Section 3.3.1). This output can be used directly in many geometry processing tasks, as we discuss in detail in Section 3.5.

We first introduce the definition of our projection operator \mathcal{P} and the conditions required for

prism-tex/figs/corona.pdf

Figure 3.3: Example of the top (left, outer) and bottom (right, inner) surface of the prismatic shell.

bijectivity of its restrictions to sections of the shell (Section 3.3.1). We then define shell validity (Section 3.3.2), present our algorithm for creating an initial shell (Section 3.3.3) and optimizing it to decrease the number of prisms (Section 3.3.4). To simplify the exposition, we initially assume that our input triangle mesh does not contain *singular points* (defined in Section 3.3.3) and boundary vertices, and we explain how to modify the algorithm to account for these cases in sections 3.3.5 and 3.3.6.

3.3.1 SHELL AND PROJECTION

Let us consider a single generalized prism Δ in a prismatic layer S (Figure 3.4 left). The generalized prism Δ is defined by the position of the vertices of three triangles, one at the top, with coordinates t_1, t_2, t_3 , one at the bottom, with coordinates b_1, b_2, b_3 , and one in the middle, implicitly defined by a per-vertex parameter $\alpha_i \in [0, 1]$, with coordinates $m_i = \alpha_i t_i + (1 - \alpha_i)b_i, i = 1, 2, 3$. We will call the top (bottom) “half” of the prism *top (bottom) slab* (we refer to Appendix A.5 for an explanation on why we need two slabs). For brevity, we will refer to a generalized prism as a prism.

DECOMPOSITION IN TETRAHEDRA We decompose each prism Δ into 6 tetrahedra (3 in the top slab and 3 in the bottom one, Figure 3.4 middle), using one of the patterns in [Dompierre et al. 1999, Figure 4]. The patterns are identified by the orientation (rising/falling) of the two edges cutting the side faces of the prism. While, for a single prism, any decomposition would be sufficient for our purposes, we need a consistent tetrahedralization between neighboring prisms to avoid inconsistencies in the projection operator. To resolve this ambiguity, we use the technique proposed in [Garimella and Shephard 2000]: we define a total ordering over the vertices of the

prism-tex/figs/prism_projection.pdf

Figure 3.4: A prism Δ (left) is decomposed into 6 tetrahedra (middle, for clarity, we only draw the 3 tetrahedra of the top slab). Each tetrahedron has a constant vector field in its interior (pointing toward the top surface), which is parallel to the only pillar of the prism that contains the point.

middle surface of Δ (naturally, we use the vertex id) and split (for each half of the prism) the face connecting vertices v_1 and v_2 with a rising edge if $v_1 < v_2$ and a falling edge otherwise.

FORWARD AND INVERSE PROJECTION We define a piecewise constant vector field \mathcal{V} inside the decomposed prism, by assigning to each tetrahedron T_j^Δ , $j = 1, \dots, 6$, the constant vector field defined by the only edge of T_j^Δ which is a *oriented pillar* of Δ connecting the bottom surface to the top surface passing through the middle surface). That is, for any $p \in T_j^\Delta$

$$\mathcal{V}(p) = t_i - b_i, \quad (3.1)$$

where i is the index of the vertex corresponding to the pillar edge of T_j^Δ . Note that \mathcal{V} is constant on each tetrahedron and might be discontinuous on the boundary: we formally define the value of \mathcal{V} on the boundary as any of the values of the incident tetrahedra. This choice does not affect our construction. There is exactly one integral (poly-)line passing through each point of the prism if all the decomposed tetrahedra have positive volumes (Theorem 3.2). This allows us to define the *projection operator* $\mathcal{P}(p)$ for a point $p \in \Delta$ as the intersection of the integral line $f_p(t)$ of the vector field \mathcal{V} passing through p , with the middle surface of Δ (Figure 3.5). Intuitively, we can project any mesh that does not fold in each prism (Figure 3.6) to the middle surface. Formally, we introduce the following definition, to describe this property in terms of the triangle normals of the mesh.

With a slight abuse of the notation, for meshes and collections of prisms A and B , we use $A \cap B$ to denote the intersection of their corresponding geometry.

Figure 3.5: A point p (left) is traced through \mathcal{V} inside the top part of the shell. A ray with p as origin and \mathcal{V} as direction is cast inside the orange tetrahedron (middle). The procedure is repeated (on the blue tetrahedron) until the ray hits a point in the middle surface (right).

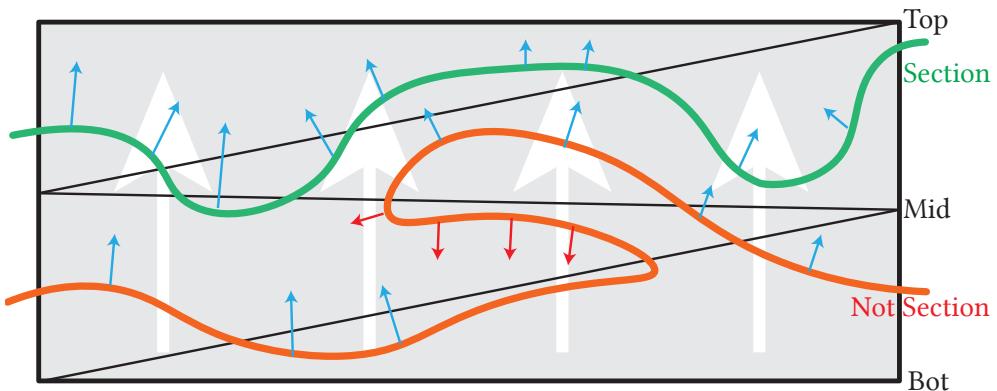


Figure 3.6: A 2D illustration for the normal dot product condition. The blue arrows agree with the background vector field (white arrows), while the red arrows do not agree.

Definition 3.1. A section $\hat{\mathcal{T}}$ of a prism Δ is a manifold triangle mesh whose intersection with Δ is a simply connected submesh $\hat{\mathcal{T}} \cap \Delta$ whose single boundary loop is contained in the boundary of Δ , excluding its top and bottom surface, and such that for every point $p \in \hat{\mathcal{T}} \cap \Delta$ the dot product between the face normal $n(p)$ and the vector field $\mathcal{V}(p)$ is strictly positive. Similarly, a triangle mesh $\hat{\mathcal{T}}$ is a section of a shell \mathcal{S} , if it is a section of all the prisms of \mathcal{S} .

Note that this definition implies that all sections are contained inside the shell. Additionally, the definition implies that the section does not intersect with either bottom or top surface. However, our definition allows for the bottom or top surface to self-intersect. The intersection of the shell does not invalidate the *local* definition of projection since it is defined per prism. Allowing intersections is crucial to an efficient implementation of our algorithm since it allows us to take advantage of a static spatial data structure in later stages of the algorithm (Section 3.3.4).

Theorem 3.2. If all 6 tetrahedra T_j^Δ in a decomposition of a prism Δ have positive volume, then the projection operator \mathcal{P} defines a bijection between any section $\hat{\mathcal{T}}$ of Δ and the middle triangle of the prism (M in Figure 3.7).

Proof. We prove this theorem in Appendix A.1.1. □

prism-tex/figs/composition.pdf

Figure 3.7: The composition $\mathcal{P}_{S_2}^{-1}(\mathcal{P}_{S_1}(x))$ with $x \in S_1$, of a direct and an inverse projection operator defines a bijection between two sections S_1 and S_2 .

The *inverse projection operator* \mathcal{P}^{-1} is defined for a section $\hat{\mathcal{T}}$ as the inverse of the forward projection restricted to $\hat{\mathcal{T}}$. It can be similarly computed by tracing the vector field in the opposite direction, starting from a point in the middle surface of the prism. Note that, differently from the inverse Phong projection [Panizzo et al. 2013; Kobbelt et al. 1998], whose solution depends on the root-finding of a quadric surface, our shell has an explicit form for the inverse and does not require a numerical solve. The combination of forward and inverse projection operators allows to bijectively map between any pair of sections, independently of their connectivity (Figure 3.7). An interesting property of our forward and inverse projection algorithm, which might be useful for applications requiring a provably bijective map, is that our projection could be evaluated exactly using rational arithmetic.

3.3.2 VALIDITY CONDITION

Shell, projection operator, section definitions, and the bijectivity condition (Theorem 3.2) are dependent on a specific tetrahedral decomposition, which depends on vertex numbering.

To ensure that our shell construction is independent from the vertex and face order, we define the validity of a shell by accounting for all 6 possible tetrahedral decompositions [Dompierre et al. 1999, Figure 4].

Definition 3.3. We say that a prismatic shell \mathcal{S} is *valid with respect to a mesh* $\hat{\mathcal{T}}$ if it satisfies two conditions for each prism.

- I1 *Positivity.* The volumes of 24 tetrahedra (Appendix A.1.2) corresponding to 6 tetrahedral decompositions are positive.
- I2 *Section.* $\hat{\mathcal{T}}$ is a section of \mathcal{S} for all 6 decompositions.

If a shell is valid, from I1, I2, then by Theorem 3.2, it follows that any map between sections induced by the projection operator \mathcal{P} is bijective.

I2 ensures that the input mesh is a *valid section* independently from the decomposition, that is, we require the dot product to be positive with respect to all *three* pillars of a prism inside the convex hull. An interesting and useful side effect of this validity condition is that it ensures the bijectivity of a natural nonlinear parametrization of the prism interior (Appendix A.4, Figure 3.25).

3.3.3 SHELL INITIALIZATION

We now introduce an algorithm to compute a *valid* prismatic shell $\mathcal{S} = \{(B_{\mathcal{S}}, M_{\mathcal{S}}, T_{\mathcal{S}}), F_{\mathcal{S}}\}$ with respect to a given triangle mesh $\mathcal{T} = \{V_{\mathcal{T}}, F_{\mathcal{T}}\}$ such that \mathcal{T} is geometrically identical to the middle surface of \mathcal{S} . We assume that the faces of the triangle mesh are consistently oriented.

EXTRUSION DIRECTION. The first step of the algorithm is the computation of an extrusion direction for every vertex of \mathcal{T} . These directions are optimized to be pointing towards the *outside* of the triangle mesh (which we assume to be orientable), that is, they must have a positive dot product with the normals of all incident faces. More precisely, for a vertex v , we are looking for a direction d_v such that $d_v \cdot n_f > 0$ for each adjacent face f with normal n_f . We can formulate this as the optimization problem

$$\begin{aligned} & \max_{d_v} \min_{f \in N_v} n_f \cdot d_v, \\ \text{s.t. } & n_f \cdot d_v \geq \epsilon, \quad \forall f \in N_v \\ & \|d_v\|^2 = 1. \end{aligned} \tag{3.2}$$

A solution, if it exists, can be found solving the following quadratic programming problem (Appendix A.3)

$$\begin{aligned} & \min \|x\|^2, \\ \text{s.t. } & Cx \geq 1, \end{aligned} \tag{3.3}$$

with $d_v = x/\|x\|$ and C the matrix whose rows are the normals n_f of the faces in the 1-ring N_v of vertex v . Solutions not satisfying $\|x\| \leq 1/\epsilon$ needs to be discarded (Appendix A.3). The QP can be solved with an off-the-shelf solver [Stellato et al. 2017; Cheshmi et al. 2020], and in particular it can be solved exactly [Gärtner and Schönherr 2000] to avoid numerical problems. Note that the Problem (3.2) is studied in a similar formulation in [Aubry and Löhner 2008] but their solution requires tolerances in multiple stages of the algorithm to handle cospherical point configurations.

The admissible set of (3.2) might be empty for a vertex v , that is, no vector d_v satisfies $Cd_v \geq \epsilon$. In this case we call v a *singularity*. For example, Figure 3.12 shows a triangle mesh containing a singularity: there exist no direction whose dot product with the adjacent face normals is positive. To simplify the explanation, we assume for the remainder of this section that \mathcal{T} does not contain singularities and also that it does not contain boundaries: we postpone their handling to sections 3.3.5 and 3.3.6.

Proposition 3.4. *Let \mathcal{T} be a closed (without boundary) triangle mesh without singularities and \mathcal{N} be a per-vertex displacement field satisfying $CN_i > 0$ for every vertex v_i of \mathcal{T} . Then there exist a*

strictly positive per-vertex thickness δ_i such that vertices t_i and b_i obtained by displacing v_i by δ_i in the direction of \mathcal{N}_i and in the opposite direction, define a shell that satisfies invariant I1.

Proof. We provide a proof in Appendix A.1.2. \square

INITIAL THICKNESS We first show that a strictly positive per-vertex thickness δ exists for a shell \mathcal{S} with \mathcal{T} as its middle surface, and then discuss a practical algorithm to realize it.

Theorem 3.5. *Given a closed, orientable, self-intersection free triangle mesh \mathcal{T} such that for all its vertices Problem (3.2) has a solution, a shell \mathcal{S} exists such that \mathcal{T} is the middle surface and there exist a strictly positive per-vertex thickness δ .*

Proof. We provide a proof in Appendix A.1.3. \square

To find a valid per-vertex thickness δ_i to construct the top surface, we initially cast a ray in the direction of \mathcal{N} for each vertex and measure the distance to the first collision with \mathcal{T} (and cap it to a user-defined parameter δ_{max} if no collision is found). An initial top mesh is built with this extrusion thickness; then we test whether any triangle in the top surface intersects \mathcal{T} , through triangle-triangle overlap test [Guigue and Devillers 2003], and iteratively shrink δ_i in this triangle by 20% until we find a thickness that prevents intersections between the input and the top surface. Analogously, we build the bottom shell along the opposite direction. Note that the thickness of a vertex for the top and bottom surface can be different.

VALIDITY OF THE INITIAL SHELL Proposition 3.4 and Theorem 3.5 ensure that the initial shell constructed using a displacement field \mathcal{N} obtained by solving (3.3), satisfies property I1. However, \mathcal{T} might not formally satisfy the conditions for being a section of \mathcal{S} , despite being identical to the middle surface of \mathcal{S} , due to our definition of the projection operator \mathcal{P} . The reason for this can be seen in Figure 3.8. After the initialization, the middle surface is the input mesh. Thus every prism Δ corresponds to a triangle T_Δ of \mathcal{T} . The intersection $\Delta \cap \mathcal{T}$, required to check if \mathcal{T} is a section of Δ (Definition 3.1) contains points from both T_Δ and its 1-ring neighborhood.

The projection operator (and thus the definition of the section) is based on all tetrahedral decompositions of Δ ; it is possible that multiple tetrahedra (with different vector field values in \mathcal{V}) overlap on the boundary of Δ .

Depending on the dihedral angles in the mesh \mathcal{T} , it is possible that the dot product between one of the pillars (orange in Figure 3.8) and the face normals of some of the triangles from 1-ring neighborhood (green in Figure 3.8) is negative. While it may seem that this problem could be addressed by changing the definition so that each edge is assigned to one of the incident triangles, so that the field direction only in one incident tetrahedron needs to be considered, this problem is more significant than it may seem, as it leads to instability under small perturbations (e.g., due to floating-point rounding of coordinates). Such small perturbations can change the set of triangles intersecting a prism and thus violate the validity of the shell and, consequently, the bijectivity of \mathcal{P} . We propose instead to refine \mathcal{T} , without changing its geometry, so that the shell corresponding to the refined mesh satisfies I2 (i.e., its middle surface is a section).

prism-tex/figs/need_to_bevel.pdf

Figure 3.8: The vector field aligned with the pillar edge (orange) has a negative dot product with the green triangle normal; the tetrahedron with this field direction meets the green triangle at the red vertex. As a consequence, the mid-surface is not a section. After topological beveling, the shell becomes valid since the dot product between the green normal and the new pillar (purple) is positive.

prism-tex/figs/bevel.pdf

Figure 3.9: The beveling patterns used to decompose prisms for which \mathcal{T} is not a section.

TOPOLOGICAL BEVELING. We identify a prism Δ_v for which I2 does not hold and use a beveling pattern [Coxeter 1973; Conway et al. 2016; Hart 2018] to decompose Δ_v in a way that \mathcal{T} becomes a section for all 6 decompositions (I2). We refer to this operation as *topological beveling*, as it does not change the geometry of the mesh, only its connectivity (Figure 3.10). We use the pattern in Figure 3.9a for Δ_v , and we use the other two patterns (b) and (c) on the adjacent prisms to ensure valid mesh connectivity. The positions of the vertices are computed using barycentric coordinates (we used $t = 0.2$, i.e., the orange dot is at 1/5 of the horizontal edge), and the normals of the newly inserted vertices are copied from the closest vertex (in Figure 3.9, the internal vertices have the normal of the triangle corner with the same color).

Theorem 3.6. *Suppose \mathcal{T} is the middle surface of \mathcal{S} , and neither $T_{\mathcal{S}}$ or $B_{\mathcal{S}}$ intersects with \mathcal{T} . After topological beveling, I2 holds, that is, \mathcal{T} is a section of the shell \mathcal{S} for all 6 decompositions.*

Proof. We provide a proof in Appendix A.1.4. □

OUTPUT The output of this stage is a valid shell with respect to \mathcal{T} (Section 3.3.2), that is, it satisfies I1 and I2.

prism-tex/figs/screwdriver_bevel.pdf

Figure 3.10: Our algorithm refines the input model (left) with beveling patterns (middle) to ensure the generation of a valid shell. The subsequent shell optimizations gracefully remove the unnecessary vertices (right).

3.3.4 SHELL OPTIMIZATION

During shell optimization, we perform local operations (Figure 3.11) on a valid shell to reduce its complexity and increase the quality. Before applying every operation, we check the validity of the operation to ensure that: (1) the resulting middle surface will be manifold [Dey et al. 1999b] and (2) the shell will be valid with respect to \mathcal{T} (to ensure a bijective projection). We forbid any operation that does not pass these checks. We would like to remark that, while there are different choices to guide the shell modification, we experimentally discovered that allowing shell simplification and optimization consistently leads to thicker shells with a richer space of sections.

Theorem 3.7. *Let \mathcal{S} be a valid shell with respect to a mesh \mathcal{T} and let $C = \{\Delta_i\}_{i \in I} \subset \mathcal{S}$ be a collection of prisms such that the middle surface M_C of C is a simply connected topological disk.*

Let \mathbb{O} be an operation replacing C with a new collection of prisms $C' = \{\Delta'_i\}_{i \in I'} \subset \mathcal{S}$, preserving both geometry and connectivity of the sides of the prism collection C , and ensuring that $M_{C'}$ is a simply connected topological disk.

If these three assumptions hold:

1. *property I1 holds for C' ,*

prism-tex/figs/local-operations.pdf

Figure 3.11: Different local operations used to optimize the shell. Mesh editing operations translate naturally to the shell setting. For vertex smoothing, we decompose the operation into 3 intermediate steps: pan, zoom, and rotate.

2. *the top and bottom surfaces of C' do not intersect \mathcal{T} ($T_{C'} \cap \mathcal{T} = B_{C'} \cap \mathcal{T} = \emptyset$),*
3. *the dot product condition $n(p) \cdot \mathcal{V}(p) > 0$ is satisfied for all points $p \in \mathcal{T} \cap \Delta'_i$ for all pillars of every prism Δ'_i of C' ,*

then, $\forall i \in I'$, $\mathcal{T} \cap \Delta'_i$ is a simply connected topological disk. In other words, \mathcal{T} is a section of the new shell \mathcal{S}' obtained by applying the operation \odot to \mathcal{S} .

Proof. We prove this theorem in Appendix A.1.5. □

We note that assumption (2) in Theorem 3.7 prevents the input surface from crossing the bottom/top surface, thus avoiding it to move in the interior of a region covered by more than one prism.

Our local operations (satisfying the definition of \odot in Theorem 3.7) are translated from surface remeshing methods [Dunyach et al. 2013] since our shell can be regarded as a triangle mesh (middle surface) extruded through a displacement field \mathcal{N} . All the local operations described below directly change the middle surface, and consequently affect the extruded shell. After every operation, the middle surface is recomputed by intersecting \mathcal{T} with the edges of the prisms in \mathcal{S} .

SHELL QUALITY We measure the quality of the shell \mathcal{S} using the MIPS energy [Hormann and Greiner 2000b] of its middle surface M_S . For each triangle T of the middle surface, we build a local reference frame, and compute the affine map J_T transforming the triangle into an equilateral reference triangle in the same reference frame. The energy is then measured by

$$\sum_{T \in M_S} \frac{\text{tr}(J_T^T J_T)}{\det(J_T)}.$$

This energy is invariant to scaling, thus allowing the local operations to coarsen the shell whenever possible while encouraging the optimization to create well-shaped triangles. Good quality of the middle surface decreases the chances, for the subsequent operations, to violate the shell invariants.

SHELL CONNECTIVITY MODIFICATIONS. We translate three operations for triangular meshes to the shell settings (Figure 3.11 top). Edge collapse, split, and flip operations can be performed by simultaneously modifying the top and bottom surfaces and retrieve the positions for the middle surface through the intersection. We only accept the operations if they pass the invariant check.

VERTEX SMOOTHING. Due to the additional degree of freedom on vertex-pairs (position, direction, and thickness), we decompose the smoothing operations into three components (Figure 3.11 bottom). *Pan* moves the positions of the top and bottom vertex at the same time, minimizing the MIPS quality of the middle surface. Neither the thickness or direction will be changed. *Rotate* re-aligns the local direction to be the average of the neighboring ones while keeping the position of the middle vertex fixed. *Zoom* keeps the direction and position of the middle vertex, and set the thickness of both top and bottom to be 1.5 times of the neighbor average, capped by the input target thickness.

INVARIANT CHECK We use exact orientation predicates [Shewchuk 1997] to make sure all the prisms satisfy positivity (I1). Further, we ensure that the original surface \mathcal{T} is not intersecting with the bottom and top surface, except at the prescribed singularities. The check is done using the triangle-triangle overlap test [Guigue and Devillers 2003], accelerated using a static axis-aligned bounding box tree constructed from \mathcal{T} . To accelerate the checks for normal condition, for each prism Δ_i , we maintain a list triangles overlapping with its convex hull (an octahedron), and check their respective normals against all the three pillars of Δ_i . These three checks ensure that the three conditions in Theorem 3.7 are satisfied. Note that the vertex smoothing operation is continuous, in the sense that any point between the current position and the optimal one improves the shell. We, however, handle it as a discrete operation to check our conditions: we attempt a full step, and if I1 is not satisfied, we perform a bisection search for a displacement that does. We avoid bisection for the other two conditions since they are expensive to evaluate.

PROJECTION DISTORTION An optional invariant to maintain (not necessary for guaranteeing bijectivity, but useful for applications), is a bound on the maximal distortion $\mathcal{D}_{\mathcal{P}}(\Delta)$ of \mathcal{P} for a prism Δ . We measure it as the maximal angle between the normals of the set C containing the faces of \mathcal{T} intersecting Δ and \mathcal{V} :

$$\mathcal{D}_{\mathcal{P}}(\Delta) = \max_{p \in C} \angle(n_p, \mathcal{V}(p)),$$

where \angle is the unsigned angle in degrees. This quantity is bounded from below by the smallest dihedral angle of \mathcal{T} , making it impossible to control exactly. However, we can prevent it from increasing by measuring it and discarding the operations that increase it. In our experiments, we use a threshold of 89.95 degrees.

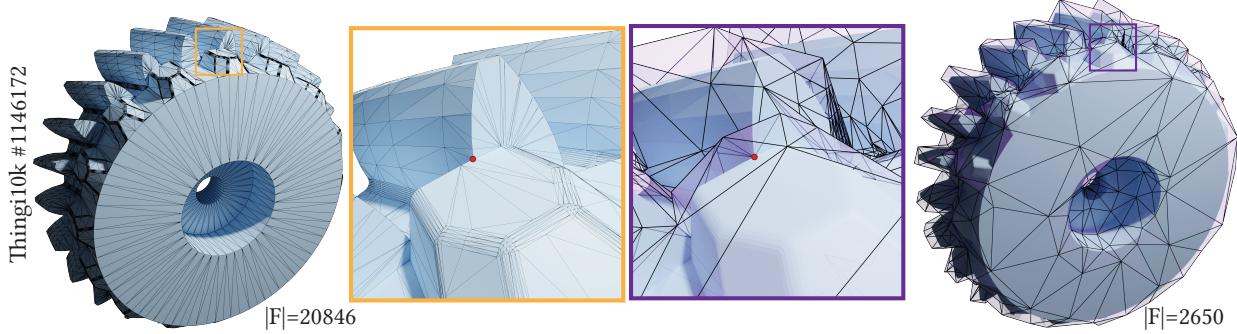


Figure 3.12: A model with 48 singularities and a close up around one (left). Our shell is pinched around each of them without affecting other regions (right).

SCHEDULING AND TERMINATION Our optimization algorithm is composed of two nested loops. The outer loop repeats a set of local operations until the face count between two successive iteration decreases by less than 0.01%. In the inner loop we: (1) flip every edge of \mathcal{S} decreasing the MIPS energy and avoiding high and low vertex valences [Dunyach et al. 2013]; (2) smooth all vertices which include pan, zoom, and rotate; and (3) collapse every edge of \mathcal{S} not increasing the MIPS energy over 30. Note that for every operation, we check the invariants, the projection distortion, and manifold preservation and reject any operation violating them. After the outer iteration terminates (i.e. the shell cannot be coarsened anymore), we further optimize the shell with 20 additional iterations of flips and vertex smoothing.

3.3.5 SINGULARITIES

Singularities, i.e. vertices of \mathcal{T} for which the constraints set of problem (3.2) is empty, are surprisingly common in large datasets (Figure 3.12 shows an example). For instance, in our subset of Thingi10k [Zhou and Jacobson 2016], although only 0.01% vertices are singular, 8% of the models have at least one singular point. This has been recently observed as a limitation for the construction of nested cages [Sacht et al. 2015, Appendix A], and it is a well-known issue when building boundary layers [Aubry et al. 2015, 2017; Garimella and Shephard 2000]. There are two main situations that give rise to singular points. The first one naturally generates a singular point when more than two ridge-lines meet (e.g., figures 3.12 and 3.30), thus making the point a feature point. The second one is a pocket-like mesh artifact, often produced as a result of mesh simplification (Figure 3.13).

While singularities might seem fixable by applying local smoothing or subdivision as a pre-process, it is not desirable in the case of a feature point, and is likely to introduce self-intersection or more serious geometric inconsistency. Therefore, due to the above reasons and observing that they are very uncommon, we propose to extend our theory (Appendix A.2) and algorithm to handle isolated singularities by *pinching* the thickness of the shell. Note that in the rare case where two singular points are sharing the same edge, they are *automatically* separated by our topological beveling.

prism-tex/figs/degenerate_prism_decompose.pdf

Figure 3.13: Left to right: examples of a singularity as a feature point and as a meshing artefact; and illustration of the degenerate prism with a singularity (red) and its tetrahedral decomposition made of only four tetrahedra.

PINCHING. We extend our definition of the shell by allowing it to have zero thickness on singularities, thus tessellating the degenerate prism with 4 instead of 6 tetrahedra (Figure 3.13). We further remark that these isolated points must be excluded from Definition 3.1. In the implementation, this requires to change the intersection predicates to skip the singular vertices. With this change, the singularity becomes a trivial point of the projection operator \mathcal{P} , and the rest of our shell can still be used in applications without further changes. Since singularities tends to be isolated (they are usually located at the juncture of multiple sharp features), this solution has minimal effects on applications: for example, when our shell is used for remeshing, pinching the shell at singularities will freeze the corresponding isolated vertices while allowing the rest of the mesh to be freely optimized.

The topological beveling algorithm is changed most significantly: for singularities, there is no pillar to copy from. In this case, we apply an additional edge split, to use the pattern in the inset (with the singularity marked by a white dot) in the one-ring neighborhood of the singularity. The newly inserted vertices lie either inside a triangle (uncircled red and orange dots), or in the interior of an edge (circled red and orange dots). Therefore, we assign to the orange vertices the average normal of the two adjacent triangles, and to red the pillar of the connected orange one.

Additionally, the edges connecting singularities will always be beveled/split after beveling. Therefore no prism will contain more than one singular point. We discuss the technical extensions for our proofs to shells with pinched prisms in Appendix A.2.

prism-tex/figs/singular_bevel.pdf

prism-tex/figs/boundary_singular.pdf

Figure 3.14: AA' is a direction with positive dot product with respect to all its neighboring faces. However, no valid shell can be built following that direction.

3.3.6 BOUNDARIES.

We introduced our algorithm, assuming that the input mesh does not have boundaries. We will now extend our construction to handle this case, which requires minor variations to our algorithm.

For some vertices on the boundary, it might be impossible to extrude a valid shell (Figure 3.14), even if problem (3.2) has a solution, as Theorem 3.5 does not apply in its original form. We identify such cases by connecting every edge in the 1-ring neighborhood of the boundary vertex to the extruded point and check if they collide with the existing 1-ring triangles (e.g., the triangle $A'AB$ intersects the existing input triangle in Figure 3.14). If it is the case, we consider this vertex as a singularity, and we pinch the shell. Note that this is an extremely rare case and, in our experiments, we detected it only for models where the loss of precision in the STL export introduces rounding noise on the boundary.

Once we pinch all boundary singularities, our construction extends naturally to the boundary. The only necessary modification is in the shell optimization (Section 3.3.4), where we skip all operations acting on boundary vertices to maintain the bijectivity of the induced projection operator (Figure 3.15). We thus freeze these vertices and never allow them to move or be affected by any other modification of the shell. Note that, in certain applications, it might also be useful to freeze additional non-boundary vertices to ensure that these remain on the middle surface during optimization (e.g., to exactly represent a corner of a CAD model).

3.4 RESULTS

Our algorithm is implemented in C++ and uses Eigen [Guennebaud et al. 2010] for the linear algebra routines, CGAL [The CGAL Project 2020] and Geogram [Lévy 2015] for predicates and spatial searching, and libigl [Jacobson et al. 2016] for basic geometry processing routines. We run our experiments on cluster nodes with a Xeon E5-2690 v2 @ 3.00GHz. The reference implementation used to generate the results is attached to the submission and will be released as an open-source project.

prism-tex/figs/open-hand.pdf

Figure 3.15: An example of a mesh with boundary.

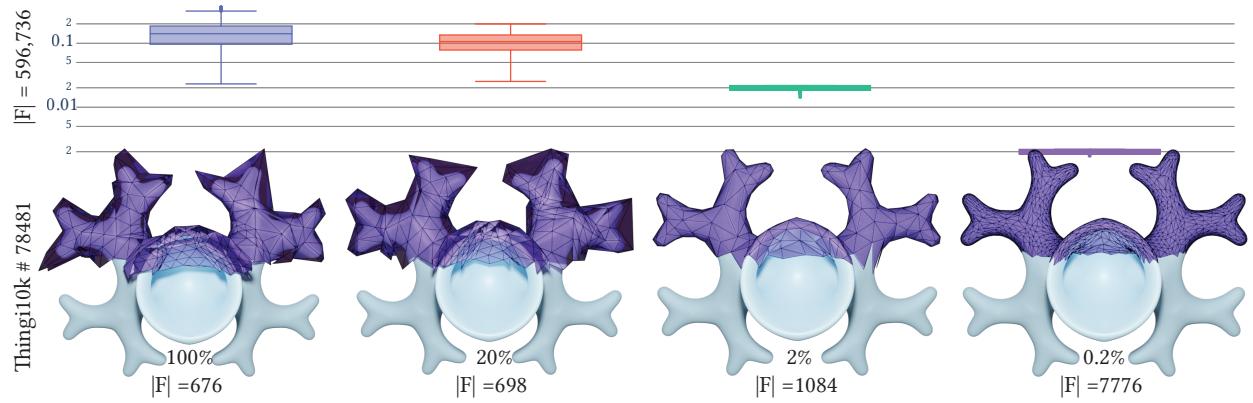


Figure 3.16: The effect of different target thickness on the number of prisms $|F|$ of the final shell, and distribution of final thickness (shown as the box plots on the top).

ROBUSTNESS For each dataset, we selected the subset of meshes satisfying our input assumptions: intersection-free, orientable, manifold triangle meshes without zero area triangles (tested using a numerical tolerance 10^{-16}). We test self-intersections by two criteria: a ball of radius 10^{-10} around each vertex does not contain non-adjacent triangles; and all the dihedral angles are larger than 0.1 degrees.

We tested our algorithm on two datasets: (1) Thingi10k dataset [Zhou and Jacobson 2016] containing, after the filtering due to our input assumptions, 5018 models; and (2) the first chunk of for the ABC dataset [Koch et al. 2019] with 5545 models. The only user-controlled parameter of our algorithm is the target thickness of our shell; in all our experiments (unless stated otherwise), we use 10% of the longest edge of the bounding box. In Figure 3.16, we show how the target thickness influences the usage of the shell: a thicker shell provides a larger class of sections, thus accommodates more processing algorithms, while a thinner one offers a natural bound on the geometric fidelity of the sections.

Our algorithm successfully creates shells for all 5018 models for Thingi10k and 5545 for ABC. We show a few representative examples of challenging models for both datasets in Figure 3.17,

Thingi10k

prism-tex/figs/10k_gallery.pdf

ABC

prism-tex/figs/abc_gallery.pdf

Figure 3.17: Gallery of shells built around models from Thingi10k [Zhou and Jacobson 2016] and ABC [Koch et al. 2019].

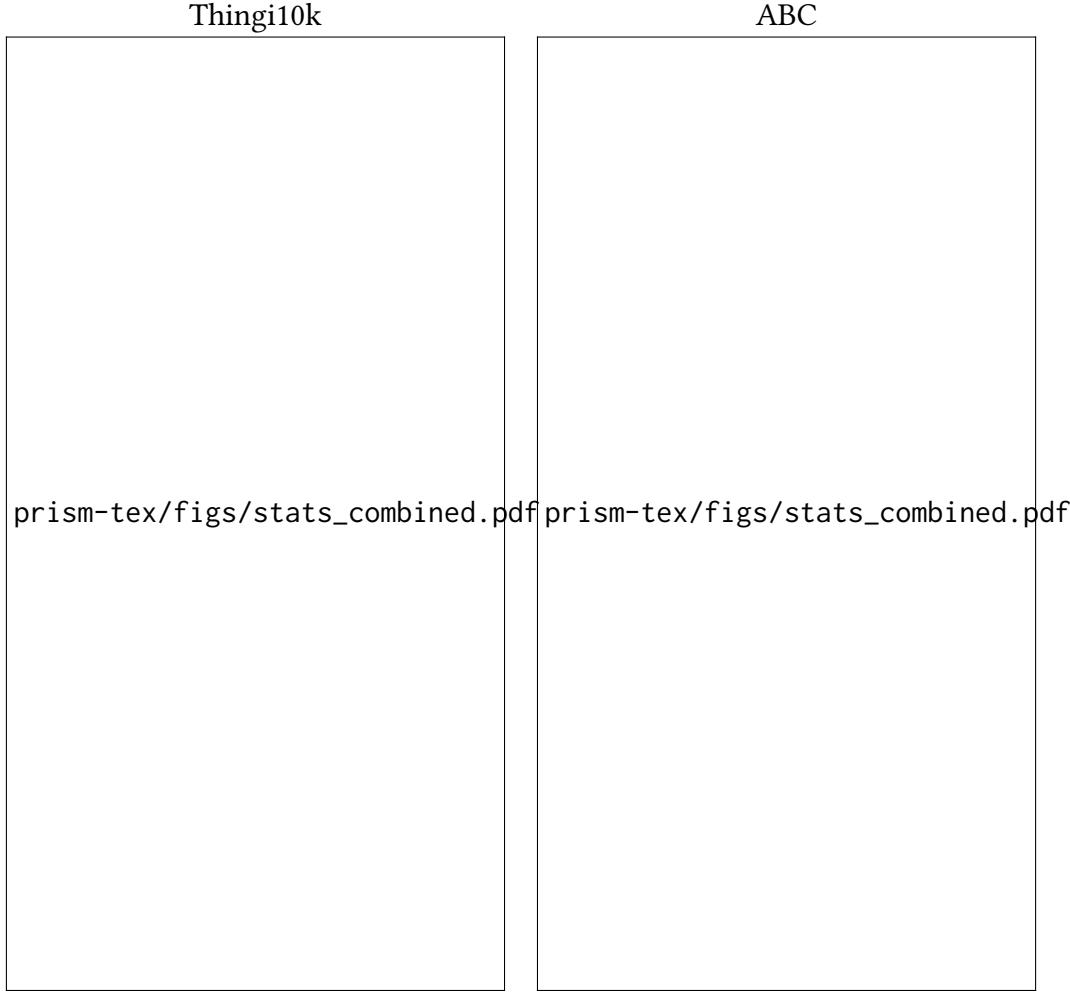


Figure 3.18: Statistics of 5018 shells in Thingi10k dataset [Zhou and Jacobson 2016] (left) and 5545 in ABC Dataset [Koch et al. 2019] (right).

including models with complicated geometric and topological details. In all cases, our algorithm produces coarse and thick cages, with a bijective projection field defined.

We report as a scatter plot the number of output faces, the timing, and the memory used by our algorithm (Figure 3.18). In total, the number of prisms generated by our algorithm is 7% and 2% of the number of input triangles for the Thingi10k and ABC dataset respectively and runs with no more than 4.7 GB of RAM. The generation and optimization of the shell takes 5min and 59s in average and up to 8.6 hours for the largest model. 50% of the meshes finish in 3 minutes and 75% in 6 minutes and 15 seconds.

COMPARISON TO SIMPLE BASELINES In Figure 3.19 we compare to two baseline methods based on [Garland and Heckbert 1998]. For each method, we generate a coarse mesh, uniformly subdivide it for visualization purposes, and query the corresponding spatial position on the original input

prism-tex/figs/thai_comparision.pdf

Figure 3.19: The *UV based method* cannot simplify the prescribed seams, and introduces self-intersections. The projection induced by the *Naive Cage* method is not continuous and not bijective; it leads to visible spikes in the reconstructed geometry.

to form the subdivided mesh. The *UV based* method is a conventional way of establishing correspondence in the context of texture mapping. However, the robust generation of a UV atlas satisfying a variety of user constraints is still an open problem. We use the state-of-the-art methods [Li et al. 2018; Jiang et al. 2017] to generate a low-distortion bijective parametrization, and use seam-aware decimation technique [Liu et al. 2017] to generate the coarse mesh. Due to the complex geometry and the length of the seam (Figure 3.19 second figure), the simplification is not able to proceed beyond the prescribed seam while maintaining the bijectivity, making the pipeline inadequate especially for building computational domains.

We also set up a baseline of the *Naive Cage* method by creating a simplified coarse mesh with [Garland and Heckbert 1998] and use Phong projection to establish the correspondence [Kobbelt et al. 1998; Panozzo et al. 2013]. Such attribute transfer is not guaranteed to be bijective; some face may not be projected (Figure 3.19 third image). With our method, we can generate a coarse mesh while having a low-distortion bijective projection (Figure 3.19 fourth image).

prism-tex/figs/shark_accuracy.pdf

Figure 3.20: Our projection is three orders of magnitude more accurate than the baseline method, and bijectively reconstructs the input vertex coordinates.

NUMERICAL ACCURACY To evaluate the numerical error introduced by our projection operator when implemented with floating-point arithmetic, we *transfer* the vertices of the input mesh to the middle surface and *inverse transfer* them from the middle surface back to the input mesh. We measure the Euclidean distance with respect to the source vertices (Figure 3.20). We compare the same experiment with the Phong projection [Kobbelt et al. 1998]. This alternative approach exhibits distance errors up to 10^{-5} even after ruling out the outliers for which the method fails due to its lack of bijectivity. The maximal error of our projection is on the order of 10^{-8} ; this error could be completely eliminated (for applications requiring an exact bijection) by implementing the projection operator and its inverse using rational arithmetic.

3.5 APPLICATIONS

Using our shell \mathcal{S} , we implement the following predicates and functions:

- $\text{is_inside}(p)$: returns true if the point $p \in \mathbb{R}^3$ is inside \mathcal{S} .
- $\text{is_section}(\mathcal{T})$: returns true if the triangle mesh \mathcal{T} is a section of \mathcal{S} .
- $\mathcal{P}(p)$: returns the prism id (pid), the barycentric coordinates (α, β) in the corresponding triangle of the middle surface, and the relative offset distance from the middle surface (h , which is -1 for the bottom surface, and 1 for the top surface) of the projection of the point p .
- $\mathcal{P}^{-1}(\text{pid}, \alpha, \beta, h)$ is the inverse of $\mathcal{P}(p)$.
- $\mathcal{P}_{\mathcal{T}}(\text{tid}, \alpha, \beta) = \mathcal{P}(q)$, where q is the point in the triangle tid of the mesh \mathcal{T} , with barycentric coordinates α, β .
- $\mathcal{P}_{\mathcal{T}}^{-1}(\text{pid}, \alpha, \beta)$ is the inverse of $\mathcal{P}_{\mathcal{T}}$.

As explained in Section 3.3, our shell may self-intersect and we opted to simply exclude the overlapping regions. In practice this affects only the function $\text{is_inside}(p)$ which needs to check if p is contained in two or more non-adjacent prisms.

prism-tex/figs/rockerarm_qslim.pdf

Figure 3.21: We attempt to simplify *rockerarm* (top left) from 20088 triangles to 100. QSlim [Garland and Heckbert 1998] succeeds in reaching the target triangle count (top right) but generates an output with a self-intersection (red) and flipped triangles (purple). With our shell constraints (bottom left), the simplification stagnates at 136 triangles, but the output is free from undesirable geometric configurations. Note that both examples use the same quadratic error metric based scheduling [Garland and Heckbert 1998].

These functions are sufficient to implement all applications below, demonstrating the flexibility of our construction and how easy it is to integrate in existing geometry processing workflows.

REMESHING We integrated our shell in the meshing algorithm proposed in [Dunyach et al. 2013] by adding envelope checks ensuring that the surface is a section after every operation. After simplification, we can use the projection operator to transfer properties between the original and remeshed surface (e.g., figures 3.22, 3.23). Since the remeshed surface is a section, a very practical side effect of our construction is that the remeshed surface is *guaranteed to be free of self-intersections*. As shown in Figure 3.21, the constraints enforced through our shell prevents undesirable geometric configurations (intersections, pockets, or triangle flips).

PROXY A particularly useful application of our shell is the construction of proxy domains for the solutions of PDEs on low-quality meshes. Additionally, by specifying the target thickness parameter (Figure 3.16), we are able to bound the geometry approximation error to the input as well. Using our method, we can (1) convert a low-quality mesh to a proxy mesh with higher quality and desired density, (2) map the boundary conditions from the input to the proxy using

prism-tex/figs/heat-in-nut.pdf

Figure 3.22: The heat method (right top) produces inaccurate results due to a poor triangulation (left top). We remesh the input with our method (left bottom), compute the solution of the heat method [Crane et al. 2013] on the high-quality mesh (middle bottom) and transfer the solution back to the input mesh using the bijective projection (bottom right). This process produces a result closer to the exact discrete geodesic distance [Mitchell et al. 1987] (top middle, error shown in the histograms).

the bijective projection map, (3) solve the PDE on the proxy (which is a standard mesh), and (4) transfer back the solution on the input surface (Figure 3.22). Our algorithm can be directly used to solve volumetric PDEs. by calling an existing tetrahedral meshing algorithm between steps 2 and 3 (figures 3.1, 3.23). In this case, we control the geometric error by setting the target thickness (we use 2% of the longest edge of the bounding box).

BOOLEAN OPERATIONS The mesh arrangements algorithm enables the robust and exact (up to a final floating-point rounding) computation of Boolean operations on PWN meshes [Zhou et al. 2016]. However, the produced meshes tend to have low triangle quality that might hinder the performance of downstream algorithms. By interleaving a remeshing step performed with our algorithm after every operation, we ensure high final quality and more stable runtime. The

prism-tex/figs/tet-knot.pdf

Figure 3.23: Knot simplified with our method and tetrahedralized with TetGen. The original model is converted to 1,503,428 tetrahedra (left) while the simplified surface is converted to only 176,190 tetrahedra (right).

composition of the bijections enables us to transfer properties between different nodes of the CSG tree (Figure 3.24).

DISPLACEMENT MAPPING The middle surface of the shell is a coarse triangular mesh that can be directly used to compress the geometry of the input mesh, storing only the coarse mesh connectivity and adding the details using normal and displacement maps (Figure 3.25). A common way to build such displacement is to project [Kobbelt et al. 1998; Collins and Hilton 2002] the dense mesh on the coarser version. As shown in Appendix A.4, our method guarantees that this natural projection is also bijective as long as the coarse mesh is a section. This alleviates the loss of information even on challenging geometry configurations, and our shell can thus be used to automate the creation of projection cages and displacement maps.

GEOMETRIC TEXTURES. The inverse projection operator provides a 2.5D parametrization around a given mesh and can be used to apply a volumetric texture (Figure 3.26). Note that we build the volumetric texture on the simplified shell, while still being able to bijectively transfer the texture coordinates.

3.6 VARIANTS

INPUT WITH SELF-INTERSECTIONS Up to this point, we assumed that our input meshes are without self-intersections. This requirement is necessary to guarantee a bijection between any section (e.g.,

prism-tex/figs/birdengine.pdf

Figure 3.24: The union of two meshes is coarsened through our algorithm, while preserving the exact correspondence, as shown through color transfer.

the input mesh) and the middle surface. Such bijection is essential for a key target application, the transfer of boundary conditions for solving PDEs on meshes or mesh-bounded domains.

However, our method can be easily extended to meshes containing self-intersections, broadening the class of meshes it can be applied to, at the cost of making the resulting shell usable in fewer application scenarios: for example, if it is used for remeshing, it will likely generate a new surface that still contains self-intersections.

If \mathcal{T} contains self-intersections, our algorithm can be trivially extended to generate a shell which will be *locally injective*, and the bijectivity of the mapping between sections still holds but with respect to the immersion. The only change required is to modify the invariance checks (Section 3.3.4): we have to replace the global intersection check with checking whether local triangles overlap with the current prisms. Figure 3.27 shows an example of a mesh \mathcal{T} with self-intersections, the generated shell, and the isolines of geodesic transferred on the coarser middle surface.

prism-tex/figs/bunny-displacement.pdf

Figure 3.25: Top: an input mesh decimated to create a coarse base mesh, and the details are encoded in a displacement map (along the normal) with correspondences computed with Phong projection [Kobbelt et al. 1998]. Bottom: the decimation is done within our shell while using the same projection as above. With our construction, this projection becomes bijective (Appendix A.4), avoiding the artifacts visible on the ears of the bunny in the top row.

prism-tex/figs/chain-lizard.pdf

Figure 3.26: Two volumetric chainmail textures (right) are applied to a shell (bottom left) constructed from the *Animal* mesh (top left). The original UV coordinates are transferred using our projection operator.

prism-tex/figs/leg-intersect.pdf

Figure 3.27: An example of a shell built around a self-intersecting mesh.

prism-tex/figs/armadillo_cage.pdf

Figure 3.28: The *Armadillo* model with four nested cages. We create a shell from the original mesh, and then rerun our algorithm on the outer shell to create the other three layers. Note that all layers are free of self-intersections, and we have an explicit bijective map between them.

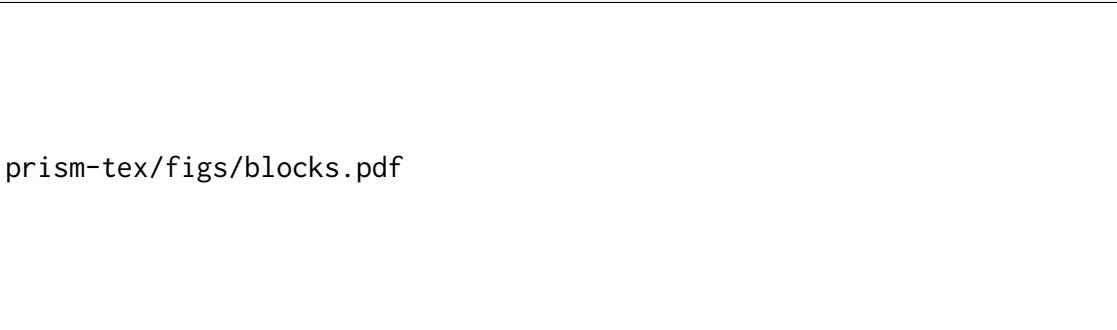
RESOLVING SHELL SELF-INTERSECTIONS For certain applications it might be preferable to have a shell whose top and bottom surfaces do not self-intersect: for example, in the construction of nested cages [Sacht et al. 2015] (useful for collision proxies and animation cages), we want to iteratively build nested shells while ensuring no intersections between them (Figure 3.28). With a small modification, our algorithm can be used to generate nested cage automatically and robustly, with the additional advantage of being able to map any quantity bijectively across the layers and to the input mesh. In contrast, [Sacht et al. 2015] does not provide guarantees on the success (e.g., the reference implementation of [Sacht et al. 2015] fails on Figure 3.19, probably due to the presence of a singularity).

To resolve the self-intersections of the top (bottom) surface, we identify the regions covered by more than one prism by explicitly testing intersections between the tetrahedralized prisms, accelerated using [Zomorodian and Edelsbrunner 2000]. For every detected prism, we reduce the thickness by 20%, and iterate until no more intersections are found. Differently from the procedure in Section 3.3.3, where reducing the thickness of the shell always maintains the validity of the shell, at this stage, the shrinking of the shell may make the shell invalid, since \mathcal{T} may not be contained



prism-tex/figs/camel.pdf

Figure 3.29: After optimization, the shell may self-intersect (left). Our postprocessing can be used to extract a non-selfintersecting shell, which is easier to use in downstream applications (right).



prism-tex/figs/blocks.pdf

Figure 3.30: We generate a pinched shell for a model with a singularity (left). Optionally, we can complete the shell using our Boolean construction.

anymore in \mathcal{S} . Whenever this happens, we perform one step of red-green refinement [Bank et al. 1983] on the regions we wish to thin, and we iterate until we succeed. This procedure is guaranteed to terminate since, on the limit of the refinement, the middle surface will be geometrically identical to \mathcal{T} , and thus Theorem 3.5 holds. In the worst case, the procedure terminates when the size of triangles on the middle surface is comparable to the input; then no refinement is required to shrink below the minimum separation of the input. Figure 3.29 shows how the intersecting shell between the legs of the camel can be shrunk to generate an intersection-free shell.

PINCHING ALTERNATIVE. For certain applications, such as boundary layer meshing, it is necessary to have a shell with non-zero thickness everywhere, including at singularities, and it is tolerable to lose bijectivity at the vicinity of a singularity. For these cases, we propose a Boolean construction to *fill* the shell around singularities, and to extend the projection operator \mathcal{P} inside these regions. That is, every point in the filled region will project to the singularity.

Without loss of generality, let us assume that \mathcal{T} has a single singularity (Figure 3.30). We initially construct a pinched shell, with zero thickness at the singularity, construct a valid shell (Section 3.3), and then perform a corefinement [Loriot et al. 2020] between a tetrahedron (centered at the singularity and whose size is smaller than the minimal thickness of the neighboring vertices) and the shell. The result of the corefinement operation (Figure 3.30 middle) consists of triangles belong to the tetrahedron, or the shell surface. The remaining part of the tetrahedron is a star-shaped polyhedron with the singularity in its kernel, and sharing a part of its boundary with the shell. This polyhedron can be easily tetrahedralized by connecting its triangulated boundary faces (one of them is highlighted in red in Figure 3.30 middle) with the singularity. For every point p in these tetrahedra, the projection operator \mathcal{P} projects p to the singularity. The remaining triangles are divided into two groups: the triangles with only one new vertex complete the degenerate prisms (one of them is highlighted in blue in Figure 3.30 middle), while the others map to the edges they are attached to.

3.7 LIMITATIONS

Currently, our algorithm is limited to manifold and orientable surfaces: its extension to non-manifold and/or non-orientable meshes is a potential venue for future work. With such an extension, the integration of shells with robust tetrahedral meshers [Hu et al. 2018, 2019d] would allow to solve PDEs on imperfect triangle meshes without ever exposing the user to the volumetric mesh, allowing them to directly work on the boundary representation to specify boundary conditions and to analyze the solution of the desired PDE. Since we rely on the additional checks for the bijective constraints, our method is slower than classical surface mesh adaptation algorithms and it is not suitable for interactive applications.

Integrating our approach into existing mesh processing algorithms might lose some of their guarantees or properties since our shell might prevent some local operations. Other surface processing algorithms guarantee some properties under some regularity assumptions on the input, which might not hold when our bijective constraints are used. For example, QSLim [Garland and Heckbert 1997] might not be able to reach the desired target number of vertices and [Dey and Ray 2010] might not be able to achieve the bounded aspect ratio. A practical limitation is that integrating our approach into existing remeshing or simplification implementations requires code level access.

Our shell is ideal for triangle remeshing algorithms employing incremental changes: not every geometry processing algorithm requiring a bijective map can use our construction. For example, it is unclear how isosurface-extraction methods [Hass and Trnkova 2020] could use our shell or how global parametrization algorithms [Kraevoy and Sheffer 2004; Schreiner et al. 2004; Alliez et al. 2003; Bommes et al. 2013b] could benefit from our method since they already compute a

map to a common domain.

3.8 CONCLUDING REMARKS

We introduce an algorithm to construct shells around triangular meshes and define bijections between surfaces inside the shell. We proposed a robust algorithm to compute the shell, validated it on a large collection of models, and demonstrated its practical applicability in common applications in graphics and geometry processing.

We believe that many applications in geometry processing could benefit from bijectively mapping spatially close surfaces, and that the idea of using an explicit mesh as a common parametrization domain could be extended to the more general case of computing cross-parametrizations between arbitrary surfaces. To foster research in this direction, we will release our reference implementation as an open-source project.

4 | BIJECTIVE AND COARSE HIGH-ORDER TETRAHEDRAL MESHES

4.1 INTRODUCTION

Piecewise linear approximations of surfaces are a popular representation for 3D geometry due to their simplicity and wide availability of libraries and algorithms to process them. However, dense sampling is required to faithfully approximate smooth surfaces. Curved meshes, that is, meshes whose element's geometry is described as a high-order polynomial, are an attractive alternative for many applications, as they require fewer elements to achieve the same representation accuracy of linear meshes. In particular, curved meshes have been shown to be effective in a variety of simulation settings in mechanical engineering, computational fluid dynamics, and graphics. Despite their major benefits, they are not as popular as linear meshes: We believe that one of the main reason for their limited usage is the lack of an automatic, robust way of constructing them.

While robust meshing algorithm exists for volumetric linear tetrahedral meshing, there are few algorithms for curved meshes, and even fewer of them having either a commercial or open-source implementation (Section 4.2). Only a few algorithms work directly on arbitrary triangle meshes (most of them require the input geometry to be either a CAD file or an implicit function), and

curve_meshing_in_shell_tex/figs/teaser.pdf

Figure 4.1: Our pipeline starts from a dense *linear* mesh with annotated features (green), which is converted in a curved shell filled with a high-order mesh. The region bounded by the shell is then tetrahedralized with linear elements, which are then optimized. Our output is a coarse, yet accurate, curved tetrahedral mesh ready to be used in FEM based simulation. Our construction provides a bijective map between the input surface and the boundary of the output tetrahedral mesh, which can be used to transfer attributes and boundary conditions.

none of them can reliably process a large collection of 3D models.

We propose the first robust and automatic algorithm to convert dense piecewise linear triangle meshes (which can be extracted from scanned data, volumetric imaging, or CAD models) into coarse, curved tetrahedral meshes equipped with a bijective map between the input triangle mesh and the boundary of the tetrahedral mesh. Our algorithm takes advantage of the recently proposed bijective shell construction [Jiang et al. 2020] to allow joint coarsening, remeshing, and curving of the dense input mesh, which we extend to support feature annotations. Our outputs are guaranteed to have a self-intersection free boundary and the geometric map of every element is guaranteed to be bijective, an important requirement for FEM applications.

The key ingredient of our algorithm is the separation of the curved volumetric meshing problem into a near-surface, surface curving problem, followed by a restricted type of linear volumetric meshing.

We believe that our curved meshing algorithm will enable wider adoption of curved meshes, as it will provide a way to automatically convert geometric data in multiple formats into a coarse tetrahedral mesh readily usable in finite element applications. To showcase the benefits of our approach we study 2 settings: (1) we show that a coarse proxy mesh can be used to compute non-linear deformations efficiently and transfer them onto a high-resolution geometry, targeting real-time simulation (Figure 4.1), and (2) we show that our meshes are ready to use in downstream FEM simulations.

We validate the reference implementation of our approach on a large collection of more than 8000 geometrical models, which will be released as an open-source project to foster the adoption of curved meshes in academia and industry.

Our contributions are:

- An algorithm to convert dense piecewise linear meshes into coarse curved volumetric meshes while preserving a bijective map with the input and bounding the approximation error.
- An extension of the bijective shell construction algorithm to support feature annotations.
- An algorithm for conforming tetrahedral meshing without allowing refinement on the boundary, but allowing internal Steiner points.
- A large-scale dataset of high-order tetrahedral meshes.

4.2 RELATED WORKS

We review the literature on the generation of unstructured and structured curved meshes. We also review the literature on boundary-preserving linear meshing, as it is an intermediate step of our algorithm.

4.2.1 CURVED TETRAHEDRAL MESH GENERATION

High-order meshes are used in applications in graphics [Bargteil and Cohen 2014; Mezger et al. 2009; Suwelack et al. 2013] and engineering analysis [Jameson et al. 2002] where it is important to reduce the geometric discretization error [Babuska and Guo 1988; Babuška and Guo 1992; Bassi and Rebay 1997; Luo et al. 2001; Oden 1994], while using a low number of degrees of freedom. The creation of high-order meshes is typically divided into three steps: (1) linear meshing of the smooth input surface, (2) curving of the linear elements to fit the surface, and (3) optimization to heal the elements inverted during curving. We first cover steps 2 and 3, and postpone the overview of linear tetrahedral algorithms to Section 4.2.3.

DIRECT METHODS. Direct methods are the simplest family of curving algorithms, as they explicitly interpolate a few points of the target curved surface or project the high-order nodes on the curved boundary [Dey et al. 1999a; Ghasemi et al. 2016; Moxey et al. 2015; Abgrall et al. 2012; Sherwin and Peiró 2002; Turner 2017; Marcon et al. 2019]. The curved elements are represented using Lagrange polynomials, [Dey et al. 1999a; Peiró et al. 2008], quadratic or cubic Bézier polynomials [George and Borouchaki 2012; Lu et al. 2013; Luo et al. 2002], or NURBS [Engvall and Evans 2016, 2017]. [Shephard et al. 2005; Sherwin and Peiró 2002] further optimizes the high-order node distribution according to geometric quantities of interest, such as length, geodesic distance, and curvature. In the case where no CAD information is available, [Wang et al. 2016], [Jiao and Wang 2012] use smooth reconstruction to compute high-order nodes and perform curving.

DEFORMATION METHODS Deformation methods consider the input linear mesh as a deformable, elastic body, and use controlled forces to deform it to fit the curved boundary. Different physical models have been employed such as linear, [Abgrall et al. 2014, 2012; Dobrzynski and El Jannoun 2017; Xie et al. 2013], and (variants of) non-linear elasticity [Persson and Peraire 2009; Moxey et al. 2016; Fortunato and Persson 2016]. A comparison between different elasticity and distortion energies is presented in [Poya et al. 2016; Turner et al. 2016; Dobrev et al. 2019].

Direct and deformation methods have been tested on small collections of simple models, and, to the best of our knowledge, none of them can provide guarantees on the validity of the output or has been tested on large collections of models. There are also no reference implementations we could compare against.

INVERSIONS AND INTERSECTIONS. Most of these methods introduce inverted elements during the curving of the high-order elements. Inverted elements can be identified by extending Jacobian metrics for linear elements [Knupp 2002, 2000] to high-order ones [Engvall and Evans 2018; Peiró et al. 2014; Johnen et al. 2013; Poya et al. 2016; Roca et al. 2012]. Various untangling strategies have been proposed, including geometric smoothing and connectivity modifications [Cardoze et al. 2004; Dey et al. 1999a; Gargallo Peiró et al. 2013; George and Borouchaki 2012; Lu et al. 2013; Luo et al. 2002; Peiró et al. 2008; Shephard et al. 2005; Dobrzynski and El Jannoun 2017; Gargallo-Peiró et al. 2015; Geuzaine et al. 2015; Roca et al. 2012; Ruiz-Gironés et al. 2017, 2016a,b; Stees and Shontz 2017; Steve L. Karman and Stefanski 2016; Toulonge et al. 2013, 2016; Ziel et al.

2017; Dobrev et al. 2019; Turner 2017; Luo et al. 2008; Lu et al. 2014]. None of these techniques can guarantee to remove the inverted elements.

An alternative approach is to start from an inversion-free mesh and slowly deform it [Persson and Peraire 2009; Ruiz-Gironés et al. 2017], explicitly avoiding inversions at the cost of possibly inaccurate boundary reproductions. These methods cannot however guarantee that the boundary will not self-intersect. Our approach follows a similar approach but uses a geometric shell to ensure element validity and prevention of boundary self-intersections.

CURVED OPTIMAL DELAUNAY TRIANGULATION [Feng et al. 2018a] generalize optimal delaunay triangulation paradigm to the high-order setting, through iteratively update vertices and connectivity. Their algorithm starts with a point cloud sampled from triangle meshes. However, the success of the method depends on the choice of final vertex number and sizing field, where insufficient vertices may result in broken topology or invalid tetrahedral meshes.

SOFTWARE IMPLEMENTATION. Despite the large literature on curved meshing generation, there are very few implementations available.

Nektar [Moxey et al. 2018] is a finite element software with a meshing component, which can generate high-order elements. We do not explicitly compare as their documentation (Section 4.5.1.5 Mesh Correction¹) states that the algorithm is not fully automatic and not designed to process robustly large collections of models. Gmsh [Geuzaine and Remacle 2009] is an open source software that supports the curved meshing of CAD models, but it does not support dense linear meshes as input. Despite the difference in the input type, we provide a comparison with Gmsh in Section 4.6.2, as it is the only method that we could run on a large collection of shapes.

To the best of our knowledge, the commercial software that support curved meshing (Pointwise [Pointwise 2018; Steve L. Karman and Stefanski 2016]) are also requiring a CAD model as input.

ANIMATION Curved tetrahedral meshes have also gained popularity in the context of fast animation. With fewer degrees of freedom and preserved geometric fidelity, [Mezger et al. 2007] observe the benefit of quadratic tetrahedra in the pipeline of physically based animation. [Suwelack et al. 2013] further investigate the transfer problem when using curved meshes as a proxy.

4.2.2 CURVED STRUCTURED MESH GENERATION

The use of a hexahedral mesh as a discretization for a volume, allows to naturally define C^k splines over the domain, which can be used as basis functions for finite element methods: this idea has been pioneered by IsoGeometric Analysis (IGA), and it is an active research area. The generation of volumetric, high-order parametrizations that conform to a given input geometry is an extremely challenging problem [Sorger et al. 2014; Peiró et al. 2015]. Most of the existing methods rely on linear hexahedral mesh generation, which is on its own a really hard problem for which automatic and robust solutions to generate coarse meshes are still elusive [Li et al. 2012; Gao et al. 2019; Guo

¹<https://doc.nektar.info/userguide/5.0.0/user-guidese17.html>

et al. 2020; Palmer et al. 2020; Zhang et al. 2020; Marschner et al. 2020] due to the inherently global nature of the problem. The current state of the art for IGA meshing is a combination of manual decomposition of the volume and semi-automated geometrical fitting [Yu et al. 2020; Coreform 2020].

In contrast, our approach is automatic, i.e. can automatically process thousands of models without any manual intervention, while providing explicit guarantees on both the validity of the elements and the maximal geometric error. Its downside is the C^0 continuity of the basis on the elements' interfaces. However, it is unclear to us if this is a real limitation: in many common settings in computer graphics and mechanical engineering (Poisson problems, linear and non-linear elasticity) the higher smoothness offered by spline functions does not make noticeable difference experimentally [Schneider et al. 2019] (and it actually leads to much worse conditioning of the system matrix, which is problematic for iterative solvers), and we thus believe that curved tetrahedral meshing is a very exciting alternative as it dramatically simplifies both the meshing and fitting of high-order elements.

4.2.3 BOUNDARY PRESERVING TETRAHEDRAL MESHING

We refer to [Hu et al. 2018] for a detailed overview of linear tetrahedral meshing, and we focus here only on the techniques that target boundary preserving tetrahedral meshing.

The most popular linear tetrahedral meshing methods are based on *Delaunay refinement* [Chew 1993; Shewchuk 1998a; Ruppert 1995a], i.e. the insertion of new vertices at the center of the circumscribed sphere of the worst tetrahedron in term of radius-to-edge ratio. This approach is used in the most popular tetrahedral meshing implementations [Si 2015b; Jamin et al. 2015a], and, in our experiments, proved to be consistently successful as long as the boundary is allowed to be refined. A downside of these approaches is that a 3D Delaunay mesh, unlike the 2D case, might still contain “sliver” tetrahedra, thus requiring mesh improvement heuristics [Cheng et al. 2000; Du and Wang 2003a; Alliez et al. 2005a; Tournois et al. 2009a]. [Alexa 2019] discusses this issue in detail and provides a different formulation to avoid it without the use of a postprocessing. [Alexa 2020] introduces an approach that does not allow insertion of Steiner points, making it not suitable for generic polyhedra domains.

There are many variants of Delaunay-based meshing algorithms including *Conforming Delaunay tetrahedralization* [Murphy et al. 2001a; Cohen-Steiner et al. 2002a], *constrained Delaunay tetrahedralization* [Chew 1989a; Si and Gärtner 2005; Shewchuk 2002a; Si and Shewchuk 2014a], and *Restricted Delaunay tetrahedralization* [Cheng et al. 2008a; Boissonnat and Oudot 2005a; Engwirda 2016].

To the best of our knowledge, all these methods are designed to allow some modifications of the input surface (either refinement, resampling, or approximation). One exception is the constrained Delaunay implementation in TetGen [Si 2015b] that allows disabling any modification to the boundary. However, this comes at the cost of much lower quality and potential robustness issues, as we show in Appendix B.3.

A different tetrahedral meshing approach has been proposed in [Hu et al. 2018], and its variants [Hu et al. 2020, 2019c], where the problem is relaxed to generate a mesh that is close to the input

to increase robustness. However, these approaches are not directly usable in our setting, as we require boundary preservation.

Due to these issues, we propose a novel boundary-preserving tetrahedral meshing algorithm specifically tailored for the shell mesh generated by our curved meshing algorithm.

4.2.4 CURVED SURFACE FITTING

There are many algorithms for fitting curved *surfaces* to dense 3D triangle meshes. The most popular approaches fit spline patches, usually on top of a quadrangular grid. Since generating quadrilateral meshes is a challenging problem for which robust solutions do not exist yet, we refer to [Bommes et al. 2012] for an overview, and only review in this section algorithms for unstructured curved mesh generation, which are more similar to our algorithm. We note that the focus of our paper is volumetric meshing: while we generate an intermediate curved surface mesh, this is not the goal of our algorithm, especially since the generated surface is only C^0 on edges.

[Hoppe et al. 1994] fits a smooth surface represented by a point cloud to a curved triangle mesh based on a subdivision surface scheme and an interleaving mesh simplification and fitting pipeline that preserves sharp features. The algorithm does not provide explicit correspondence to the input: they are defined using distance closest point, which is not bijective far from the surface.

[Krishnamurthy and Levoy 1996] converts dense irregular polygon meshes of arbitrary topology into coarse tensor product B-spline surface patches with accompanying displacement maps. Based on the work [Lin et al. 2007] that fits triangle surface meshes with Bézier patches, [Zhang et al. 2011] fits triangle surface meshes with high-order B-spline quadrilateral patches and adaptively subdivide the patches to reduce the fitting error. These methods produce smooth surfaces but do not have feature preservation.

Another related topic is the definition of smooth parametric surfaces interpolating triangle meshes. We refer to [Zorin 2000] for an overview of subdivision methods and discuss here the approaches closer to our contribution.

[Hahmann and Bonneau 2003] proposed to use triangular Bézier patches to define smooth surfaces over arbitrary triangle meshes ensuring tangent plane continuity by relaxing the constraint of the first derivatives at the input vertices. Following Hahmann’s work, [Yvert et al. 2005b] presents a more complete pipeline: perform QEM simplifications, trace the coarse mesh onto the dense one and perform parameterization relaxing and smoothing. Then it fits a hierarchical triangular spline [Yvert et al. 2005a] to the surface. More recent work [Tong and Kim 2009] approximates the triangulation of an implicit surface with a G^1 surface. These schemes are usually designed to interpolate existing meshes rather than simplifying a dense linear mesh into a coarse curved mesh and are thus orthogonal to our contribution.

4.3 SHELL PRELIMINARIES

We briefly overview [Jiang et al. 2020] as our work uses and extends it. [Jiang et al. 2020] introduces bijective projection shells (which we will abbreviate as shells in the rest of the paper), a new

geometry processing tool to perform mesh editing while preserving a close-by bijective map to the input surface.

SHELL. The shell is defined by three triangulated surface meshes $\bar{\mathcal{S}} = \{(B_s, V_s, T_s), F_s\}$ sharing the same mesh connectivity F_s , where B_s, V_s , and T_s are the vertices of the bottom, middle, and top surface respectively. Each triangle in F_s corresponds to a *generalized prism*, defined by connecting the corresponding triangles in the three surfaces with straight edges, called *pillars*. Each prism P has three vertices $v_i \in V_s, t_i \in T_s, b_i \in B_s, i = 1, 2, 3$ on the middle, top, and bottom surface, respectively. Each pillar decomposes into a top $h_i^T = t_i - v_i, i = 1, 2, 3$ and bottom $h_i^B = b_i - v_i, i = 1, 2, 3$ slab, and each slab can be canonically decomposed into 3 tetrahedra, and each tetrahedron contains a constant vector field aligned with the pillars it is connected with [Jiang et al. 2020, Figure 4]. The vector field is used to define a *projection operator* Π within the shell.

PROJECTION OPERATOR. For every prism P , the projection operator Π_P is defined as the tracing of the piece-wise constant vector field V inside the decomposed prism, by assigning to each tetrahedron $T_j^P, j = 1, \dots, 6$, the constant vector field defined by the only edge of T_j^P which is one of the oriented pillars $h_i^T, h_i^B, i = 1, 2, 3$ ([Jiang et al. 2020, Figure 4]). That is, $\forall p \in T_j^P$

$$\Pi_P(p) = h_i^k,$$

where i is the index of the vertex corresponding to the pillar edge of T_j^P , and k is either the top or bottom surface. The shell projection operator Π is defined as the operator whose restriction to P , $\Pi|_P$ is Π_P , and it defines a bijective map between every pair of specific triangle meshes contained in the shell, called *sections*.

SECTION. A triangle mesh is a section if it is contained within the shell, and if the dot product of the normal of each of its triangles with the vector field in each of the overlapping tetrahedra is positive. The projection operator Π defines a bijective map between any pair of sections if the shell is *valid*.

SHELL VALIDITY. [Jiang et al. 2020] defines a shell $\bar{\mathcal{S}}$ to be *valid* with respect to an input mesh \mathcal{M} if it satisfies two conditions:

1. The volumes of all possible tetrahedral decomposition of a prism (24 of them) are positive.
2. \mathcal{M} is a section for all possible tetrahedral decompositions. That is, the input mesh is contained within the shell, and the dot product between the mesh's normals and the shell's pillar is positive.

SINGULARITY. Singular vertices are a special geometric configuration, where the neighboring triangles of a specific vertex admits a conflicting set of normals. [Jiang et al. 2020] extends the shell construction to allow such cases. Around the (isolated) singular vertices, the prisms are pinched to become generalized pyramids, composed of two tetrahedra instead of three.

curve_meshing_in_shell_tex/figs/illustrations/input-output.pdf

Figure 4.2: Input triangle mesh \mathcal{M} and points \mathcal{P} . Output curved tetrahedral mesh \mathcal{T}^k and bijective map ϕ^k .

The algorithm to build the shell creates an initial valid extrusion, potentially thin and dense, and then iteratively uses the shell local operations (i.e., vertex smoothing, edge collapse, edge split, and edge flip) [Jiang et al. 2020, Section 3.4] to improve its quality while preserving the validity.

4.3.1 VARIATION FROM THE ORIGINAL ALGORITHM

To extend the shell formulation in [Jiang et al. 2020] to accommodate for feature preservation (Section 4.5), we modify the definition of a valid section [Jiang et al. 2020, Definition 3.1], by relaxing the intersection between a triangle and a prism in the discrete case. That is, we do not consider the prism to be intersecting a triangle if they share only one vertex of the triangle; we also ignore the intersection if they intersect on a feature edge on opposite sides. The bijectivity and validity condition of the shell projection trivially holds.

4.4 CURVED TETRAHEDRAL MESH GENERATION

INPUT. The input of our algorithm is a collection of oriented manifold, watertight, self-intersection-free triangle mesh $\mathcal{M} = (V, F)$, and a set of points $p_i \in \mathcal{P}$ (possibly empty) on the surface of \mathcal{M} (Figure 4.2, left) where the distance bound ε is prescribed. The collection \mathcal{M} must be consistently oriented such that it is the boundary of an oriented 3-manifold. A set of edges can also be optionally provided as annotated features (Section 4.5).

OUTPUT. The output of our algorithm is a tetrahedral mesh $\mathcal{T}^k = (V^k, T^k)$ of order k . Formally, each tetrahedron $\tau \in T^k$ is defined through the *geometric map* from the reference tetrahedron $\hat{\tau}$,

$$g^\tau = \sum_{j=1}^n c_j^\tau L_j(\hat{u}, \hat{v}, \hat{w}), \quad (4.1)$$

where $\hat{u}, \hat{v}, \hat{w}$ are the local coordinates of a point in $\hat{\tau}$, c_j^τ are the control points for a tetrahedron τ , and L_j are polynomial bases (typically Lagrange bases). For two tetrahedra τ_1 and τ_2 of T^k sharing

curve_meshing_in_shell_tex/figs/illustrations/high-order.pdf

Figure 4.3: Lagrange nodes on the reference element $\hat{\tau}$ for different $k = 1, 2, 3$ and example of geometric mapping g .

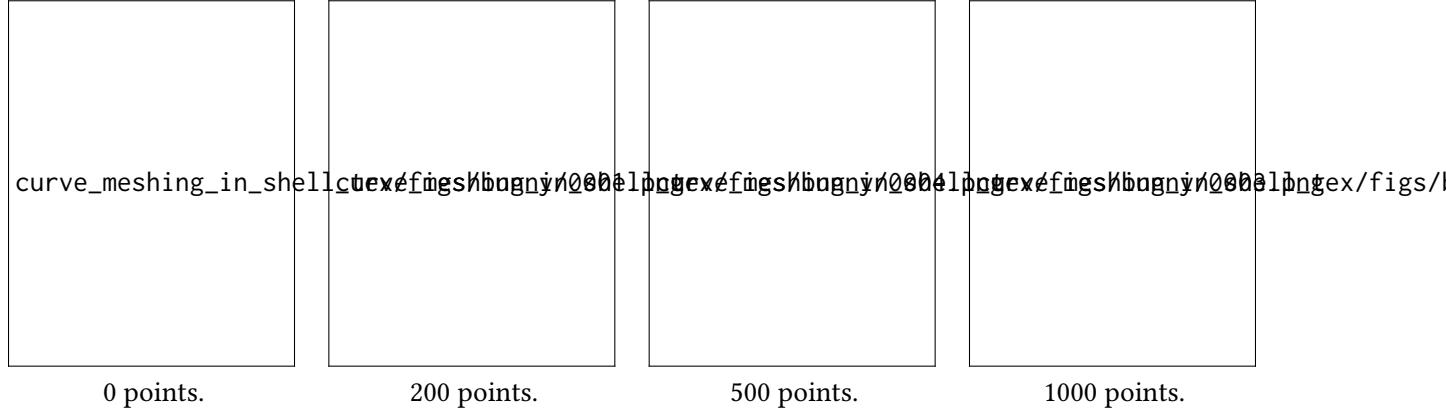


Figure 4.4: Effect of the choice of the set \mathcal{P} on the output.

a face F , the restriction of the maps g^{τ_i} , $i = 1, 2$, to F coincide. Figure 4.3 shows the position of the control points c_j on the reference element for $k = 1, 2, 3$ for the Lagrange bases. We call the tetrahedralization of a curved mesh \mathcal{T}^k *positive* if $\det(J_{g^\tau}) > 0$ everywhere on every τ . In particular, for $k = 1$, since g is affine, J_{g^τ} is constant and the positivity reduces to the positive orientation of the vertices [Shewchuk 1997].

Note that, while bijectivity of the geometric map g^τ implies positivity, the reverse is not true. Therefore, our algorithm not only checks for $\det(J_{g^\tau}) > 0$, but also ensures that the boundary $\partial\mathcal{T}^k$ does not intersect; we show in Appendix B.1 that these two conditions guarantee the bijectivity of g^τ . Furthermore, our algorithm ensures that the distance from any point in \mathcal{P} to $\partial\mathcal{T}^k$ (the surface of \mathcal{T}^k) is smaller than an user-controlled parameter ε .

We are not assuming anything on \mathcal{P} : a sparse set of points will generate a mesh less faithful to the input geometry, while a dense sampling computed for instance with Poisson disk sampling [Bowers et al. 2010] will prevent the surface from deviating too much (Figure 4.4).

Our algorithm guarantees that the tetrahedralization is positive and that $\partial\mathcal{T}^k$ does not have self-intersections. It also aims at coarsening \mathcal{T}^k as much as possible while striving to obtain a good geometric quality. To reliable fit $\partial\mathcal{T}^k$ to \mathcal{M} we require a bijective map

$$\phi^k: \mathcal{M} \rightarrow \partial\mathcal{T}^k$$

from the input \mathcal{M} to the surface of \mathcal{T}^k (Figure 4.2 right). Our algorithm also generates this map and exposes it as an output for additional uses, such as attribute transfer. Note that, since we build

curve_meshing_in_shell_tex/figs/intersection-free.pdf

Figure 4.5: Our algorithm maintains free of intersection even on challenging models, without the need of setting adaptive threshold.

upon the shell construction in [Jiang et al. 2020], we also guarantee $\partial\mathcal{T}^k$ is homeomorphic and topology-preserving with respect to \mathcal{M} (Figure 4.5).

To simplify the explanation, we use the bar $\bar{\cdot}$ to represent quantities on the *straight* linear shell, the tilde $\tilde{\cdot}$ for the *curved* shell, and hat $\hat{\cdot}$ for the reference elements (e.g., \bar{P} is the prism on the straight coarse shell, \tilde{P} is the curved prism, and \hat{P} is the prism on the reference configuration).

Definition 4.1. We call a curved mesh \mathcal{T}^k and its mapping ϕ^k to \mathcal{M} *valid* if it satisfies the following conditions:

1. ϕ^k is bijective;
2. the distance between any $p \in \mathcal{P}$ and $\partial\mathcal{T}^k$ is less than ε ;
3. \mathcal{T}^k is positive (i.e., every geometric map g^τ has positive Jacobian's determinant).

OVERVIEW. Our algorithm starts by creating a valid mesh (i.e., it satisfies 4.1), then it performs local operations (Appendix B.2) to improve \mathcal{T}^k (i.e., coarsen it and improve its quality) while ensuring all the conditions remain valid with respect to local modification. To achieve this goal, our algorithm uses two stages: (1) curved shell generation and (2) tetrahedral mesh generation and optimization.

curve_meshing_in_shell_tex/figs/illustrations/pipeline.pdf

Figure 4.6: Overview of curved mesh generation pipeline.

STAGE 1: CURVED SHELL CONSTRUCTION. In the first stage (Section 4.4.1) we extend the shell construction of [Jiang et al. 2020] by combining the shell projection Π with a high-order mapping

$$\psi^k: \tilde{\mathcal{S}} \rightarrow \bar{\mathcal{S}}.$$

We start from the *valid* shell $\bar{\mathcal{S}}$ constructed from the input mesh \mathcal{M} (i.e., \mathcal{M} is a section $\bar{\mathcal{S}}$). We call $\bar{\mathcal{S}}$ a *projection shell* and call the prismatic projection Π . Together with the construction of $\bar{\mathcal{S}}$, we build an order k *curved prismatic shell* $\tilde{\mathcal{S}}$ that defines a curved layer around \mathcal{M} and a *bijective* map ψ^k between $\tilde{\mathcal{S}}$ and $\bar{\mathcal{S}}$ (Figure 4.6, first three figures) that ensures that the distance between \mathcal{M} and $\partial\mathcal{T}^k$ is smaller than ε (Section 4.4.2). That is, $\phi^k(p) < \varepsilon$ for any $p \in \mathcal{P}$. (Note that we do not require \mathcal{M} to be a section of $\tilde{\mathcal{S}}$.)

To facilitate the volumetric meshing in the next stage (Section 4.4.3), we restrict the top and bottom surface of $\tilde{\mathcal{S}}$ to be linear (independently from the order of ψ^k). The final output of this first stage is a high-order volumetric shell, a bijective mapping $\phi^k = \Pi \circ \psi^k$, and a *positive* tetrahedralization of $\tilde{\mathcal{S}}$ with flat boundary. In other words, the tetrahedralization of $\tilde{\mathcal{S}}$ satisfies 4.1.

STAGE 2: TETRAHEDRAL MESH GENERATION. In the second stage (Section 4.4.3) we use boundary-conforming tetrahedralization to connect the top and bottom surface of $\tilde{\mathcal{S}}$ with a background tetrahedral mesh, thus generating a *positive* order k tetrahedralization \mathcal{T}^k of a bounding box around the input, which we can further optimize with local operations to improve its quality (Figure 4.6, last two figures).

To ensure that our first condition is satisfied, we define the mapping ϕ^k as a composition of several mappings, which we ensure are bijective. For the second condition, we initialize our construction with ϕ^k as the identity, and thus, the distance at the sample points is zero. After every operation, we recompute the distance and “undo” the operation if the distance becomes larger than ε . To ensure that the last condition holds, we rely on checking if all prisms (linear and curved) have positive geometric mapping, which ensures that they can be tetrahedralized with a positive tetrahedralization. Ensuring the condition while coarsening \mathcal{M} allows us to generate a coarse *curved* tetrahedral mesh \mathcal{T}^k and the *bijective* map ϕ^k to the input mesh \mathcal{M} .

4.4.1 HIGH-ORDER SHELLS

To simplify the explanation, we first focus on the case where $\varepsilon = \infty$, that is, we aim at generating an as-coarse-as-possible curved mesh. Note that, the trivial solution (i.e., a single tetrahedron) is not necessarily a valid \mathcal{T}^k since it would be impossible to build the bijective mapping ϕ^k .

The output of [Jiang et al. 2020] is a coarse shell $\bar{\mathcal{S}}$ with a piecewise linear middle surface. To curve it, we construct a shell $\tilde{\mathcal{S}}$ and the bijective map ψ^k while constructing $\bar{\mathcal{S}}$. The shell $\tilde{\mathcal{S}}$ is constructed warping every prism \tilde{P} of $\tilde{\mathcal{S}}$ with ψ^k . Since we define ϕ^k as $\Pi \circ \psi^k$, and Π satisfies the first two conditions 4.1, we only need to ensure that ψ^k is bijective, for ϕ^k to be bijective. We define the mapping $\psi^k = \bar{\omega} \circ (\tilde{\omega}^k)^{-1}$ through two cross-parametrization maps from the reference prism \hat{P} :

$$\bar{\omega}: \hat{P} \rightarrow \bar{P}, \quad \tilde{\omega}^k: \hat{P} \rightarrow \tilde{P}.$$

Both mappings $\bar{\omega}$ and $\tilde{\omega}^k$ are defined as the tensor product between the base triangular mapping (high-order for $\tilde{\omega}^k$) and pillar's barycentric heights. That is, for a prism \tilde{P}

$$\tilde{\omega}^k(\hat{u}, \hat{v}, \hat{h}) = \sum_{j=1}^n c_j^P L_j(\hat{u}, \hat{v}) ((1 - \hat{u} - \hat{v}) h_1^{\tilde{P}} + \hat{u} h_2^{\tilde{P}} + \hat{v} h_3^{\tilde{P}}) \hat{h},$$

where $\hat{u}, \hat{v}, \hat{h}$ are the barycentric coordinates in the reference prism, c_j the control points of a triangle, $h_i^{\tilde{P}}$ the three pillars heights, and L_j a triangle polynomial basis. By ensuring that ψ^k is bijective, we guarantee that any curved tetrahedralization of a prism \tilde{P} will be a valid tetrahedralization of $\tilde{\mathcal{S}}$.

We note that to decouple the following tetrahedral mesh generation and the curved shell generation, we ensure that $\tilde{\omega}^k$ maps the top and bottom face of the curved prism \tilde{P} to a linear triangle.

After each local operation, we generate samples $\hat{s}_i, i = 1, \dots, m$ on the parametric base of the prism \hat{P} and use $\Pi^{-1} \circ \bar{\omega}$ to map \hat{s}_i back to \mathcal{M} and $\tilde{\omega}^k$ to map them to $\partial\mathcal{T}$. Using the mapped point we solve

$$\min_{c_i} \sum_{i=1}^m \|(\Pi^{-1} \circ \bar{\omega})(\hat{s}_i) - \tilde{\omega}[c_i]^k(\hat{s}_i)\|_2^2,$$

where c_i are the control points of $\tilde{\omega}^k$. As $\tilde{\omega}[c_i]^k(\hat{s}_i)$ is a linear function of c_i . This is a quadratic optimization problem. The control points of the top and bottom surface are fixed to ensure that $\tilde{\omega}^k$ maintains the two surfaces as linear. We validate the bijectivity of $\tilde{\omega}^k$ by checking positivity of the determinant [Johnen et al. 2013] and that the top and bottom surfaces are intersection free. The intersection is simplified in our case since fast and exact algorithms[Guigue and Devillers 2003] are available since the top and bottom surfaces stay linear.

4.4.2 DISTANCE BOUND

In the previous section, we explained how to generate a curved shell $\tilde{\mathcal{S}}$ that satisfies 4.1. To ensure that the middle surface of $\tilde{\mathcal{S}}$ has a controlled distance from the points in \mathcal{P} , we interleave

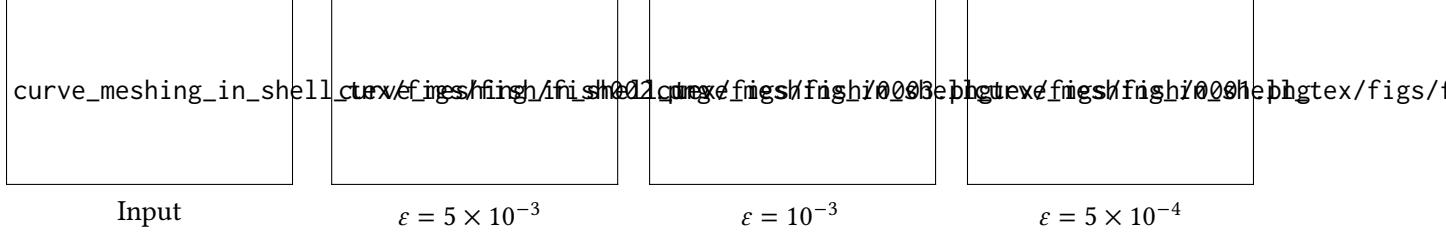


Figure 4.7: A model simplified with different distances.



Figure 4.8: Two dimensional overview of the five steps of our boundary preserving tetrahedral meshing algorithm.

a distance check in the construction of $\tilde{\omega}^k$ after each local operation. Formally, after every local operation we use the mapping ϕ^k to map every point $p_i \in \mathcal{P}$ to $\tilde{p}_i = \phi^k(p_i)$ a point on the coarse curved middle surface of $\tilde{\mathcal{S}}$ and, if $\|p_i - \tilde{p}_i\| \geq \varepsilon$ we reject the operation. The initial shell is trivially a valid initialization as ϕ^k is identity and thus, the distance is zero. Note that, \tilde{p}_i is not necessarily the closest point to p_i on $\partial\mathcal{T}$, thus $\|p_i - \tilde{p}_i\|$ is an upper bound on the actual pointwise distance. Figure 4.7 shows the effect of the distance bound on the surface; a small distance will lead to a denser mesh with more details, while a large one will allow for more coarsening.

4.4.3 TETRAHEDRAL MESHING

The outcome of the previous stage is a curved tetrahedralization of $\tilde{\mathcal{S}}$ that closely approximates \mathcal{M} with linear (“flat”) boundaries. We now consider the problem of filling its interior (and optionally its exterior) with a tetrahedral mesh, a problem known as *conforming boundary preserving* tetrahedralization.

Several solution exists for this problem (Section 4.2.3) and the most common implementation is TetGen [Si 2015b]. Most algorithms refine the boundary, which allows deriving bounds on the quality of the tetrahedral mesh. However, in our setting, this is problematic, as any change will have to be propagated to the curved shell. To avoid coupling the volumetric meshing problem with the curved shell coarsening, while technically possible it is very challenging to implement robustly, we opt for using a tetrahedral meshing algorithm that preserves the boundary exactly. Not many algorithms support this additional constraint, the only one with a public implementation is the widely used TetGen algorithm. However, we discovered that, when this option is used, it

suffers from robustness issues, which we detail in Appendix B.3. To solve this problem in our specific setting, we propose in the following five step algorithm (Figure 4.8) taking advantage of the availability of a shell, based on the TetWild [Hu et al. 2018] algorithm.

STEP 1. To generate a boundary preserving linear mesh, we first exploit the shell to extrude the bottom surface B (and top T) further by a positive (potentially small) constant δ such that the newly extruded bottom surface B_e (and top T_e) does not self-intersect. The space between B and B_e (and between T and T_e) consists of prisms divided into positive tetrahedra.

STEP 2. Then we insert B_e and T_e in a background mesh \mathcal{B} generated following TetWild algorithm ([Hu et al. 2018, Section 3.1]), that is, we use the triangle of B_e and T_e as the input triangle meshes for the first stage of the TetWild algorithm, which inserts them into a background mesh \mathcal{B} , so that each input triangle is a union of faces of refinement of \mathcal{B} .

We interrupt the algorithm after the binary space partitioning (BSP) subdivision (and before the TetWild mesh optimization [Hu et al. 2018, Section 3.2]) to obtain a positive tetrahedral mesh in rational coordinates with a surface with the same geometry of B_e and T_e , but possibly different connectivity as TetWild might refine it during the BSP stage.

STEP 3. Our original goal was to compute a mesh conforming to B and T , but we could not do it directly with TetWild as they might be refined. We now replace the mesh generated by TetWild between B_e and T_e with another one conforming to B and T . To achieve this, we delete all tetrahedra between B_e and T_e , and insert the surfaces B and T , which will “float” in the empty space between B_e and T_e . We now want to fill the space between B and B_e with positive tetrahedra conforming to the surfaces B and T .

STEP 4. Every prism P , made by a bottom triangle B^T and a bottom extruded triangle B'_e and its corresponding bottom extruded refined triangle $B'^T \in \mathcal{B}$, can be tetrahedralized without refining B : That is, we first decompose the prism B^T, B'_e in tetrahedra (always possible by construction), then refine every tetrahedron touching B'^T . By repeating the same operation on the space between T and T_e we will have a positive linear boundary conforming tetrahedral mesh of B and T .

STEP 5. The tetrahedra generated in the previous step will have rational coordinates and will also likely have low quality. To round the coordinates to floating-point representation and to improve their quality, we use the mesh optimization stage of TetWild, with the minor variant of keeping the vertices and edges on B and T frozen. Note that the vertices in B and T are already roundable to floating-point representation, as they were part of the input.

CURVED TETRAHEDRAL MESH OPTIMIZATION After generating the conforming linear tetrahedral mesh, we stitch it with the tetrahedralized $\tilde{\mathcal{S}}$ to obtain a valid output mesh \mathcal{T}^k (Definition 4.1). However, its quality might be low, in particular in the curved region, as $\tilde{\mathcal{S}}$ can be thin with large triangles. To improve the quality of \mathcal{T}^k we adapt the local operation of a linear pipeline to our



curve_meshing_in_shell_tex/figs/illustrations/input-output-feature.pdf

Figure 4.9: Input triangle mesh with features and output curved mesh with feature preserved equipped with bijective map ϕ^k .

curved settings. We propose three local operations: smoothing, collapse, and flip. Since the surface of \mathcal{T}^k is already coarse and of high quality, as part of the definition of ϕ^k , we prevent any local operation from changing it. We validate every local operation (i.e., check the positivity of \mathcal{T}^k) using the convex-hull property [Johnen et al. 2013] and reject the operation if it is violated. Our local operations are prototypical, and we leave as future work a more comprehensive study of curved mesh optimization.

SMOOTHING. As for the linear case, we compute the total energy of a vertex v by summing up the energies of the tetrahedra adjacent to it, which we compute on 56 regularly sampled points. We then perform gradient descent for all high-order nodes in the one-ring neighborhood of v . That is, we collect all edge nodes, face nodes, and cell nodes of the one-ring neighborhood of v . Differently from the linear case, the optimization is expensive since the nodes neighborhood typically contains hundreds of nodes.

COLLAPSE AND SWAP The collapse and swap are the same as in a linear mesh, and we place the high-order nodes of the newly created face on the linear flat face.

4.5 FEATURE PRESERVING CURVED SHELL

INPUT We enhance the input to additionally include a set of feature edges $f_i \in \mathcal{F}$ and feature vertices $v_i \in \mathcal{V}$ such that no triangle in F has more than one feature edge (Figure 4.9 left). (This property can be satisfied on any generic mesh by performing 1-to-3 refinement on every triangle with more than one feature edge).

OUTPUT Since the input has features, the output curved mesh \mathcal{T}^k will also have curved feature edges $f_i^k \in \mathcal{T}^k$, feature vertices $v_i^k \in \mathcal{V}^k$, and the bijective map ϕ^k preserves features by bijectively mapping \mathcal{F} to \mathcal{F}^k and \mathcal{V} to \mathcal{V}^k . Our method makes no assumption on the topology and “quality” of the features. If the features are reasonable, it will produce a high-quality mesh, while if the

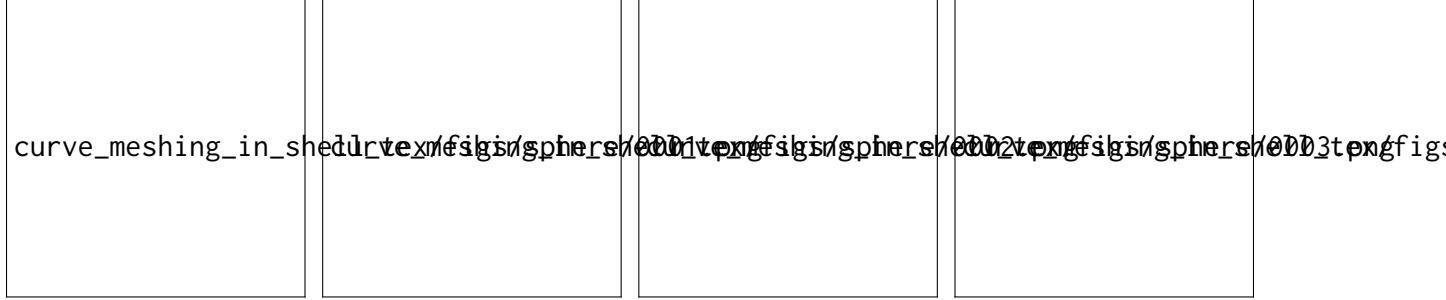


Figure 4.10: A sphere with different marked features (green). As we increase the number of features our algorithm will preserve them all but the quality of the surface suffers.



Figure 4.11: Input feature (green) is not preserved after traditional shell simplification.

features are close our algorithm will preserve them and result in smaller triangles on the surface. (Figure 4.10).

The previous construction generates valid curved tetrahedral meshes and the bijective map ϕ^k based on the construction of [Jiang et al. 2020]. However, the shell construction cannot coarsen features: the authors suggest freezing them. For instance, when performing an edge collapse on the feature, the new coarse edge (orange) will not map to the feature (green) anymore (Figure 4.11). To ensure feature preservation we extend Definition 4.1.

Definition 4.2. We call a curved mesh \mathcal{T}^K and its mapping ϕ^k from \mathcal{M} *valid and feature preserving* if they are valid (Definition 4.1) and ϕ^k bijectively maps \mathcal{F} to \mathcal{F}^k and \mathcal{V} to \mathcal{V}^k

As for the non-feature preserving case, we always aim to maintain a valid feature preserving \mathcal{T}^k .

To account for features, we propose to change the prismatic map Π , that is, we only need to change the first stage. This is done by snapping the input features (Section 4.5.1). That is, we modify \mathcal{M} to “straighten” the feature to ensure that the coarse prismatic projection preserves them and construct a mapping β between the straight mesh $\overline{\mathcal{M}}$ and \mathcal{M} (Figure 4.13).

The outcome is a *valid* shell $\overline{\mathcal{S}}$ with respect to the straight surface $\overline{\mathcal{M}}$ (i.e., $\overline{\mathcal{M}}$ is a section $\overline{\mathcal{S}}$) that preserve features, the prismatic projection Π , and the bijective map β that can be directly used in the curved pipeline (Section 4.4). That is, the mapping ϕ^k will be defined as $\phi^k = \beta \circ \Pi \circ \psi^k$.

As for the non-preserving feature pipeline we ensure that our conditions are always met, starting from a trivial input and rejecting operations violating them. Our goal is to modify the

curve_meshing_in_shell_tex/figs/illustrations/stage-1.pdf

Figure 4.12: Overview of the construction of the first stage of our pipeline.



Figure 4.13: The input mesh has feature edges snapped, to create a valid shell, as well as the curved mesh.

input mesh \mathcal{M} and create $\overline{\mathcal{M}}$ by moving its vertices. In such a way, the mapping β is simply barycentric. To guarantee bijectivity of ϕ^k we need to ensure that all mappings composing it are bijective, in particular β . To ensure that β is bijective it is enough that $\overline{\mathcal{M}}$ is self-intersection free (guaranteed by the shell construction) and that all its triangles have positive area. By straightening the features of \mathcal{M} we ensure that the edges of the prism will map to the feature. Thus, ϕ^k will be feature preserving.

FEATURE GROUPING. The first step of our pipeline consists of grouping successive edges $f_i \in \mathcal{F}$ into poly-lines and identifying two categories: loops and graphs (Figure 4.14). For every graph, we identify its nodes and add them as feature vertices. In other words, we add to \mathcal{V} all the end-points and junction of poly-lines.

curve_meshing_in_shell_tex/figs/illustrations/feat-gr.pdf

Figure 4.14: The input edges feature (green) are grouped together in poly-lines and categorized in graph (left) and loops (right). For every graph we add the nodes (blue) to the set of feature vertices.

curve_meshing_in_shell_tex/figs/illustrations/loc-op.pdf

Figure 4.15: Illustration of smoothing on a feature.

4.5.1 FEATURE STRAIGHTENING.

To allow feature coarsening, we propose to straighten \mathcal{M} to ensure that all features are collinear. In other words, we build, together with the shell $\bar{\mathcal{S}}$, a mesh $\bar{\mathcal{M}} = (\bar{V}, \bar{F})$ (i.e., a mesh with the same connectivity F of \mathcal{M}) and features $\bar{\mathcal{F}}$ such that every triangle of $\bar{\mathcal{M}}$ has a positive area, $\bar{\mathcal{M}}$ is a section of $\bar{\mathcal{S}}$, and the features in $\bar{\mathcal{F}}$ are collinear. In such a way, the mapping β is trivially defined as piecewise affine ($\bar{\mathcal{M}}$ and \mathcal{M} share the same connectivity) and is locally injective as long as all the triangles on $\bar{\mathcal{M}}$ have positive areas. Note that the bijectivity of β follows from the fact that the shell prevents self-intersections of $\bar{\mathcal{M}}$.

To construct a mesh $\bar{\mathcal{M}}$ with straight features, we start with $\bar{\mathcal{M}} = \mathcal{M}$ (in the beginning all prisms of $\bar{\mathcal{S}}$ cover at most one feature edge). Let $\bar{f}^1 = \{\bar{f}_i^1\}, i = 1, \dots, n$ and $\bar{f}^2 = \{\bar{f}_i^2\}, i = 1, \dots, k$ two chains of feature edge belonging to the same feature $\bar{f} \in \bar{\mathcal{F}}$. For every local operation acting on a feature \bar{f}_1 and \bar{f}_2 we first construct the new feature $\bar{f}^n = \{\bar{f}_i^n\}, i = 1, \dots, k$ such that the segments $(\bar{f}_i^n, \bar{f}_{i+1}^n)$ are collinear and their length is proportional to (f_i, f_{i+1}) (the feature vertices in the input mesh \mathcal{M}), that this we use arc-length cross parameterization from f to \bar{f}^n (Figure 4.15 show an example of smoothing a feature). Moving vertices of \bar{f}^n will also move the vertices of $\bar{\mathcal{M}}$ thus, straighten the mesh as the local operations proceed. After the construction of \bar{f}^n we check if the newly constructed $\bar{\mathcal{M}}$ is still a section of $\bar{\mathcal{S}}$ and if the triangles modified by the straightening have areas larger than ϵ . In practice, we choose $\epsilon = 10^{-10}$: a smaller value would lead to numerical instabilities and a larger one to less straightening.

Note that not all features can be straight; for instance, if a triangle has three feature vertices (the snapped feature will result in a degenerate triangle, thus, β will not be bijective) or if the snapping flips the normal ($\bar{\mathcal{M}}$ will no longer be a section of $\bar{\mathcal{S}}$). Both are extremely rare cases in

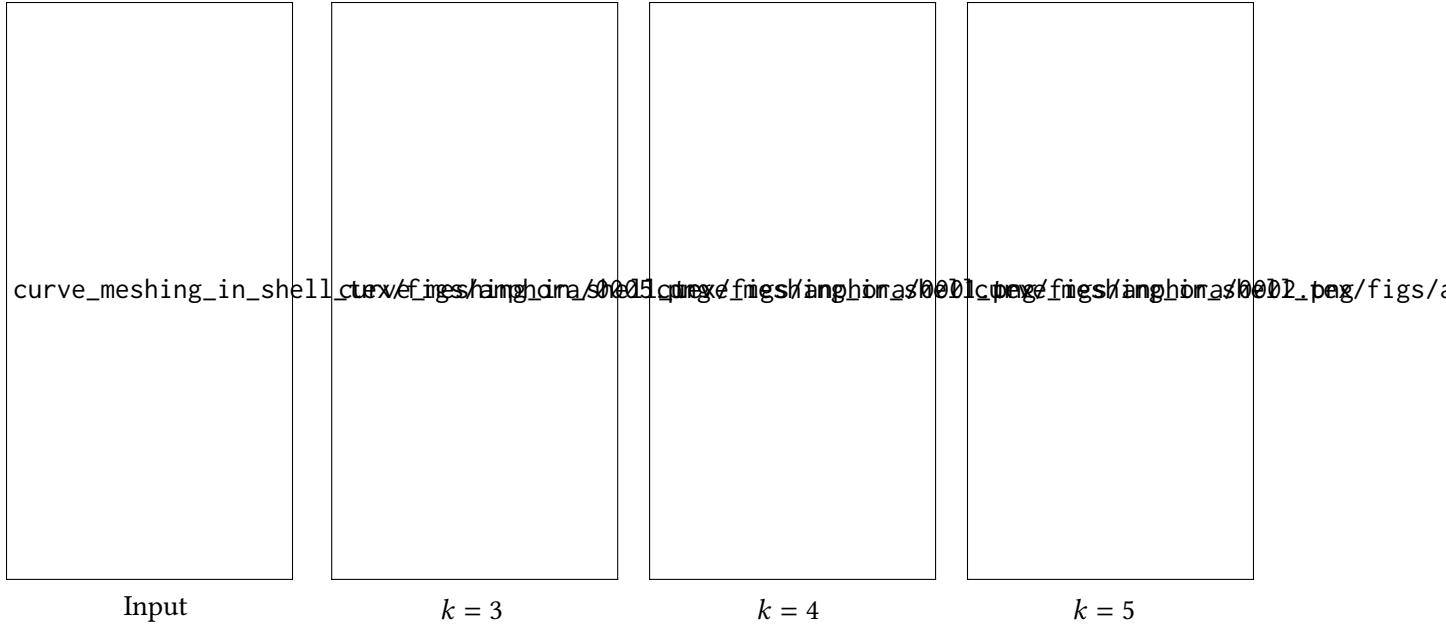


Figure 4.16: Curved meshes of different order. The additional degrees of freedom allows for more coarsening.

our dataset.

4.6 RESULTS

Our algorithm is implemented in C++, using Eigen [Guennebaud et al. 2010] for the linear algebra routines, CGAL [The CGAL Project 2020] and Geogram [Lévy 2015] for predicates and geometric kernel, libigl [Jacobson et al. 2016] for basic geometry processing routines, and meshio [Schlömer 2020] for converting across the different formats. We run our experiments on cluster nodes with Intel Xeon Platinum 8268 CPU 2.90GHz. The reference implementation and the data used to generate the results will be released as an open-source project.

To simplify the exposition, all meshes presented in this section are quartic meshes ($k = 4$). Our method is flexible and, for lower k , it will generate denser meshes (Figure 4.16).

4.6.1 LARGE SCALE VALIDATION.

We tested the robustness and quality of the result produced by our algorithm on three datasets: (1) Thingi10k dataset [Zhou and Jacobson 2016] containing 3574 models without features; (2) the first chunk of the ABC dataset [Koch et al. 2019] with 5328 models with features marked from the STEP file and (3) the CAD dataset [Gao et al. 2019] containing 106 models with semi-manual features.

Note that the original datasets contain more models since, for each of them, we selected meshes satisfying our assumptions: intersection-free (using the same strategy as in [Jiang et al. 2020] with

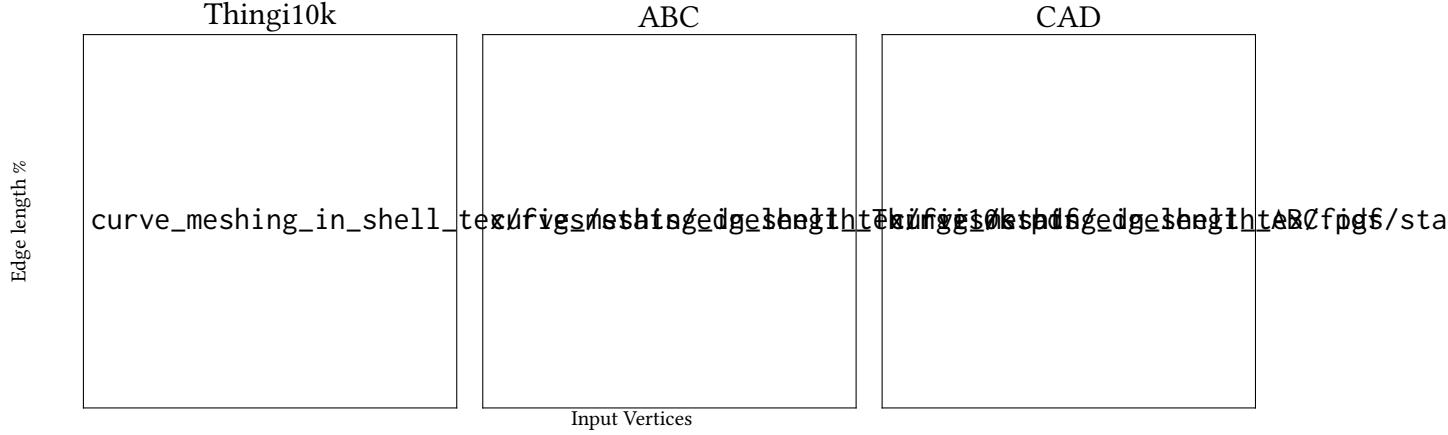


Figure 4.17: Relative average edge length (with respect to longest bounding box edge of each model) of our curved meshes versus number of input vertices.



Figure 4.18: Within the same distance bound (10^{-3} of the longest bounding box side), our method generates a coarser high order mesh, compared to the linear counterpart generated by fTetWild.

a distance tolerance of 10^{-6} and dihedral angle of 2°) oriented, manifold triangle meshes, smallest triangle area larger than 10^{-8} .

Our method has only the geometry accuracy parameter ε , which we set to 1% of the longest bounding box edge, and the point set \mathcal{P} which we set as the input vertices V . With this basic setup, our algorithm aims to produce the coarsest possible mesh while preserving features and striving to generate high-quality meshes. Our algorithm successfully generates curved meshes for 3527 for Thingi10k, 5268 for the ABC, and all for CAD dataset within 12 hours; by allowing more time, all models but 3 can be successfully processed. The 3 failures are due to models with a small one-tetrahedra component that “move” inside the shell as it grows. This is an implementation choice: we use collision detection instead of continuous collision detection for efficiency reasons.

Our method successfully generates coarse meshes whose average edge length is 10% of the model size while preserving features (Figure 4.17). Figure 4.18 shows how our method successfully

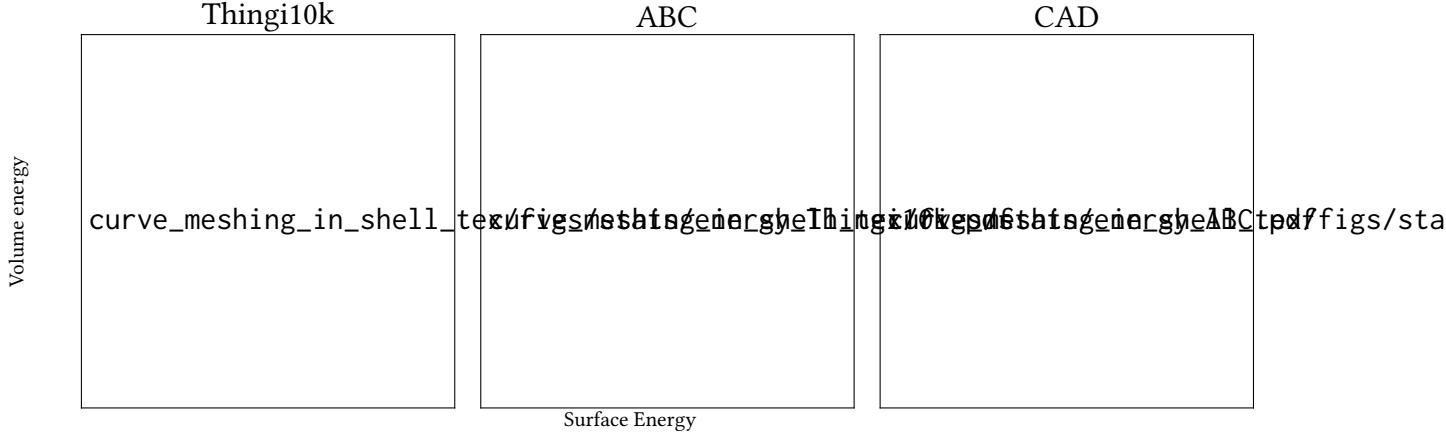


Figure 4.19: Surface and volume average MIPS energy of the output of our method (the CAD volume energy is truncated at 100, excluding 6 models).

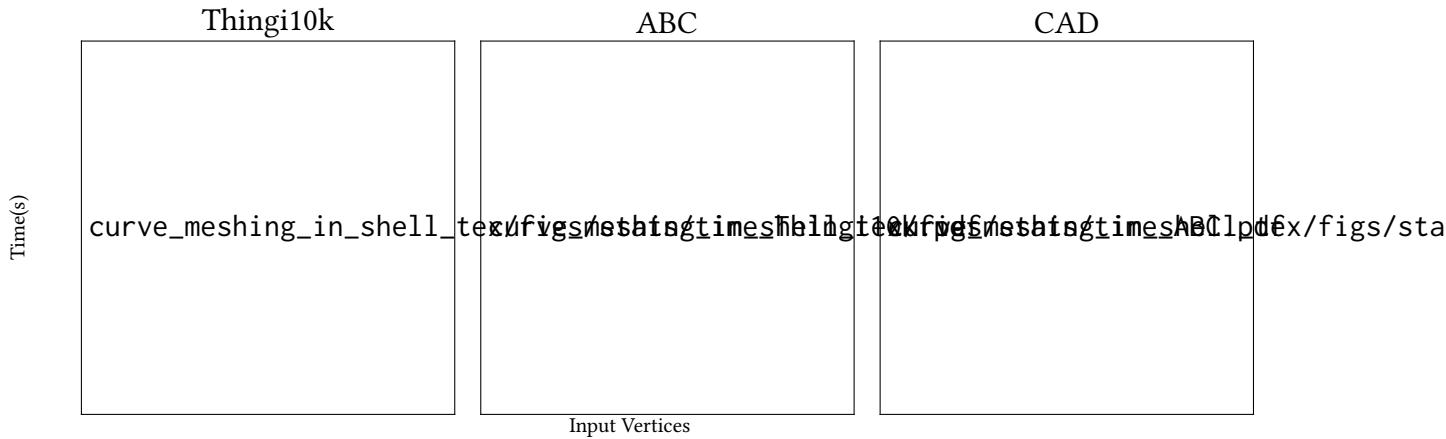


Figure 4.20: Timing of our algorithm versus the input number of vertices.

captures the features and coarsen the surface with curved elements, while many linear elements (generated with fTetWild [Hu et al. 2020]) are required to closely approximate the surface.

The output of our algorithm can be directly used in the simulation (Section 4.6.4) since we guarantee that the geometric mapping g is positive. To ensure good conditioning and performance of the numerical solver, we measure the MIPS energy [Hormann and Greiner 2000b; Fu et al. 2015b] of our output meshes (Figure 4.19).

Figure 4.20 shows the running time of our method with respect to the number of vertices. The running time of our algorithm is linear with respect to the number of input vertices; it takes around an hour for a model with around 10 thousand vertices.

4.6.2 COMPARISONS

CURVED ODT [Feng et al. 2018a] is, to the best of our knowledge, the only existing algorithm designed to convert dense triangle meshes into coarse, curved approximations. The input and

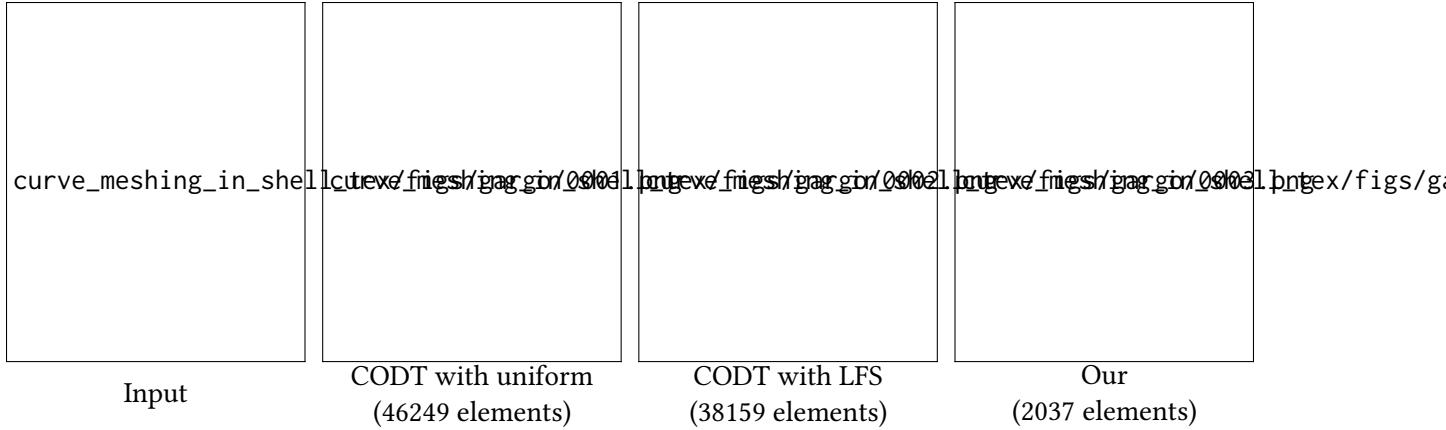


Figure 4.21: Compared with Curved ODT, our method does not rely on setting vertex number and sizing field, and can generate coarse valid results.

output are the same as in our algorithm. However, their method does not provide a bijective map between the input and output, does not guarantee to preserve features, has no bound on the distance to the input surface, and does not guarantee that the elements are positive. While our algorithm has been designed to process large collections of data automatically, exposing only a few intuitive options to control the faithfulness to the input, the reference implementation of the curved ODT method provided to us by the authors requires the user to choose multiple *per-model* parameters to achieve good results, and the parameters have a strong effect on the quality and validity of the result (as shown in [Feng et al. 2018a, Figure 16]). We thus restricted our comparison to only a small selection of models (see additional material) that the authors of [Feng et al. 2018a] processed for us.

From our discussions with the authors, we observed that curved ODT generates a *valid* output when we provide (1) a sufficiently large number of vertices and (2) a good local feature size sizing field (LFS) [Alliez et al. 2005a] to efficiently spend the vertex budget in the regions with more geometrical details. Figure 4.21 shows an example of a model for which [Feng et al. 2018a] fails to converge when using a uniform sizing field, while it succeeds when the sizing field is used.

In contrast, our algorithm can be run automatically on a large collection of geometrical models, it is guaranteed to have positive Jacobian (up to the use of floating-point predicates, Section 4.7), it preserves features, it automatically controls the density of the output depending on the desired user-provided distance threshold, and it provides a bijective map between the input mesh the boundary of the curved surface. For the model in Figure 4.21, our result contains 2037 elements, 18 times less than the curved ODT algorithm.

GMSH [Geuzaine and Remacle 2009] can only generate curved meshes from boundary representation (BRep), that is the input is not exactly the same as ours. To compare both algorithms, we start from the BRep and we generate a dense *linear* mesh that we use for our input. The Gmsh algorithm first constructs a curved mesh by fitting the high-order nodes to the BRep (possibly inverting elements) then performs mesh optimization to untangle them [Remacle et al. 2013];

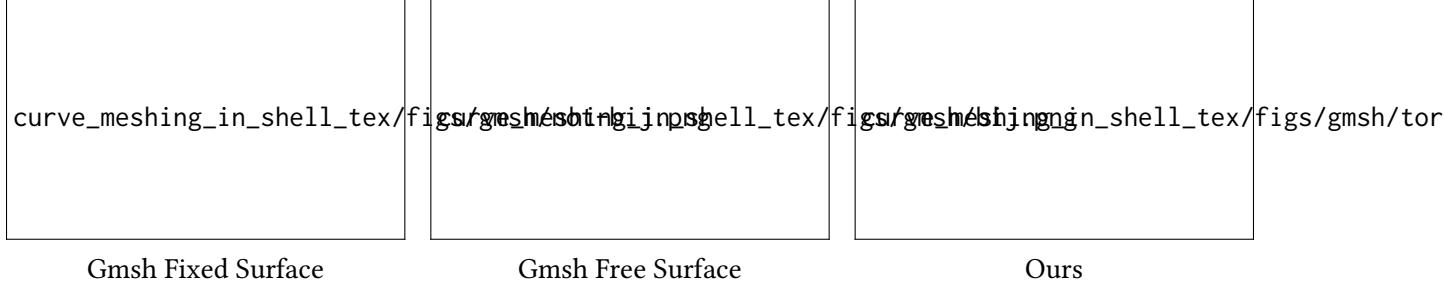


Figure 4.22: Example of a BRep meshed with Gmsh where the optimization fails to untangle elements when fixing the surface. By allowing the surface to be modified, the mesh becomes “wiggly”. Our method successfully generate a positive curved mesh.



Figure 4.23: Example of a STEP file meshed with Gmsh where, due to the low mesh density, the tetrahedralization is not positive. Gmsh manages to generate a positive mesh by using a denser initial tessellation. Since our method starts from a dense mesh and coarsen it, it can successfully resolve the geometry.

thus has no guarantee to generate positive meshes while preserving the surface (Figure 4.22, left). Additionally, Gmsh algorithm cannot control the distance from the input when the untangling allows the surface to move, and thus the surface is “wiggly” and denser than our result (Figure 4.22, center). We also observed that if the initial surface mesh is not dense enough, Gmsh closes holes and cannot generate a valid tetrahedral mesh (Figure 4.23).

4.6.3 FLEXIBILITY

Since the input to our method is a triangle mesh, our method naturally supports a variety of input that can be easily converted into triangle meshes. For instance, \mathcal{M} can be generated from marching an implicit surface or Catmull-Clark subdivision of a hand-made quad mesh (Figure 4.24). The bijective map ϕ^k is used, for instance, to transfer the geodesic distance field or color information from the input triangle mesh to the coarse curved mesh.



Implicit microstructure

subdivision surface

Figure 4.24: Our algorithm processes triangle meshes that can be extracted from different formats: an implicit microstructure geometry from [Tozoni et al. 2020] or a subdivision surface from [Crane 2013]. The bijective map preserved on the surface allows taking advantage of the plethora of surface algorithms including polyhedral geodesic computation [Mitchell et al. 1987] and texture mapping.

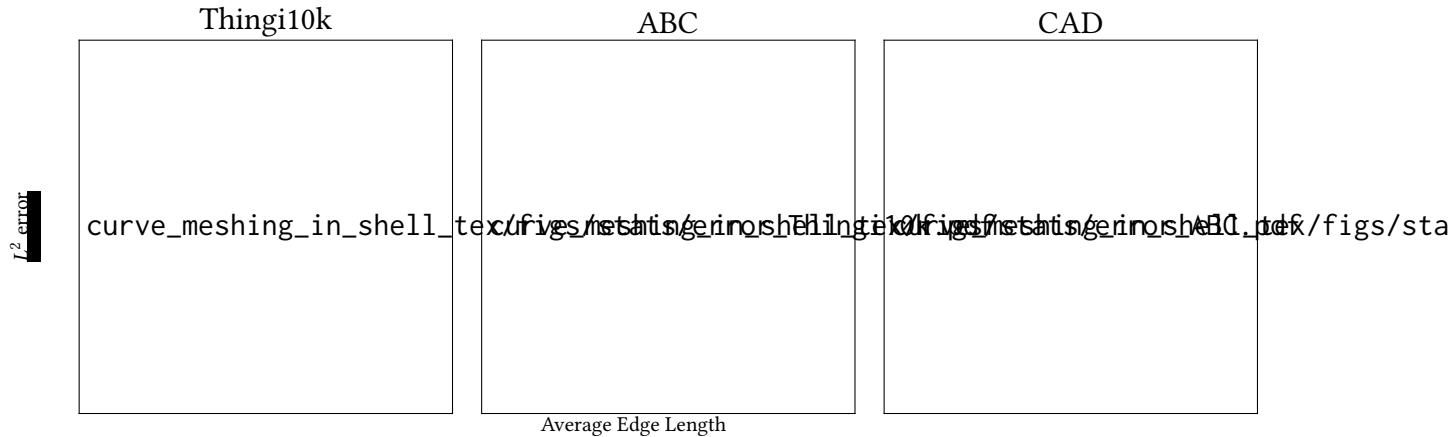
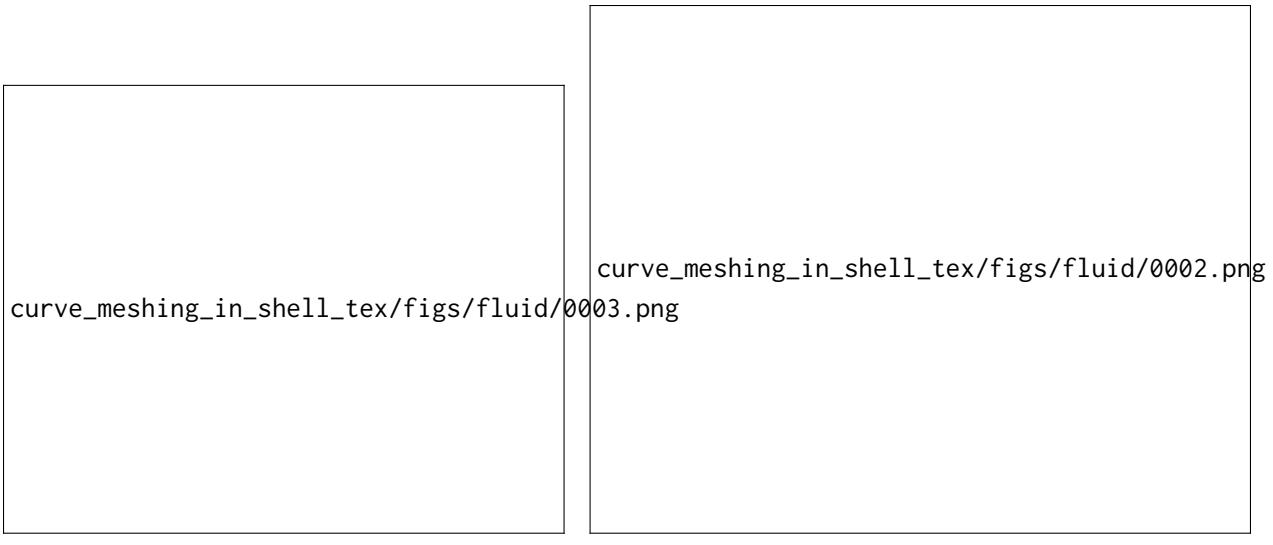


Figure 4.25: L^2 error of the solution of the Poisson equation with respect to model size on our three datasets.

4.6.4 APPLICATIONS

LARGE SCALE POISSON To show that our meshes are ready for simulation, we solve for Poisson equation

$$\Delta u = f, \quad u|_{\partial\Omega} = g,$$



Tetrahedral mesh

Curved simulation

Figure 4.26: By meshing the region between a box and a complicated obstacle, we are able to perform non-linear fluid simulation on our curved mesh.

where Ω is the domain (i.e., the mesh), f is the right-hand side, and g are the Dirichlet boundary conditions. To simplify the setup and the error measurements, we use fabricated solutions [Salari and Knupp 2000]. That is, we choose the function u_{exact} to be

$$\begin{aligned} u_{\text{exact}}(x_1, x_2, x_3) = & \\ & 3/4 e^{-(9x_1-2)^2+(9x_2-2)^2+(9x_3-2)^2)/4} + 3/4 e^{-(9x_1+1)^2/49-(9x_2+1)/10-(9x_3+1)/10} \\ & + 1/2 e^{-(9x_1-7)^2+(9x_2-3)^2+(9x_3-5)^2)/4} - 1/5 e^{-(9x_1-4)^2-(9x_2-7)^2-(9x_3-5)^2}, \end{aligned}$$

then we plug it in the equation to obtain f (g is simply u_{exact}). Figure 4.25 shows the L^2 error (average) distribution across our three datasets using our quartic meshes with quadratic approximation of u (i.e., we use superparametric elements).

HIGH ACCURACY FLUIDS Our curved meshes can be directly used to solve different partial differential equations (PDEs). For instance, by meshing the part outside the top shell and discarding the rest, we can generate a curved background mesh for the Navier-Stokes equation (Figure 4.26).

FAST ANIMATION Our coarse curved meshes can be used as animation proxies as in [Mezger et al. 2007; Suwelack et al. 2013]. We first compute an as-coarse-as-possible curved mesh (i.e., we set ϵ to infinity). Then we apply the boundary condition to simulate an elastic distortion of the curved mesh using linear elements. Finally, we use our bijective map φ^4 to map the displacement back to the input high-detailed surface mesh (Figure 4.1). The results are almost indistinguishable to a classical pipeline (i.e., mesh the input mesh), but the runtime is 400 times faster (8s versus over 50 minutes).

4.7 LIMITATIONS AND CONCLUDING REMARKS

We introduce an automatic algorithm to convert dense triangle meshes in coarse, curved tetrahedral meshes whose boundary is within a user-controlled distance from the input mesh. Our algorithm supports feature preservation, and generates meshes with positive Jacobians and high quality, which are directly usable for FEM simulations.

LIMITATIONS. Our algorithm generates meshes with a C^0 geometric map. For most FEM applications, this is not an issue. However, for geometric modeling applications, where only the mesh boundary is used, the C^0 geometric map introduces normal discontinuities, which are undesirable. While the surface looks smooth from far away, plotting the reflection lines shows the discontinuity between the normals. We believe an exciting extension of our work would be to study the feasibility of using geometric maps that are C^1 [Lyche and Muntingh 2015] or C^2 [Xia and Qian 2017]. A second limitation is that, in our implementation, the validity conditions (Definition 4.1) are currently checked using floating point arithmetic, using numerical tolerances to account for rounding errors. While our implementation works on a large collection of models, it is possible that it will fail on others due to the inexact validity predicates. We are not aware of exact predicates for these conditions, and we believe that developing them is an interesting, and challenging, venue for future work.

FUTURE WORK. Our current high order mesh optimization pipeline is preliminary, as it only supports vertex smoothing, collapse, and swap. Adding additional operators, allow them to exploit the curved geometric map, and optimizing the boundary could lead to a further increase in mesh quality. While simple at a high-level, this change will require to merge the two parts of our algorithm, a major implementation effort.

CONCLUSIONS. We believe that our work will foster adoption of curved meshes, and open the door to a new family of geometry processing algorithms able to take advantage of this highly compact, yet accurate, shape representation.

5 | DECLARATIVE SPECIFICATION FOR UNSTRUCTURED MESH EDITING ALGORITHMS

5.1 INTRODUCTION

Unstructured triangular and tetrahedral meshes are widely used in graphics, engineering, and scientific computing due to their flexibility to represent objects with complex boundaries. Such unstructured meshes find their usage in modeling and rendering 3D objects and scenes, discretizing partial differential equations for physical simulation, collisions detection and response, path planning in robotics, and many other applications.

An unstructured mesh is usually stored in a custom data-structure supporting a set of local operations to add, remove, or change its elements and their properties. A major research effort has been invested in exploring different data-structures and evaluating their generality and efficiency (Section 5.2), which led to the development of mesh libraries such as CGAL [The CGAL Project 2020], VCG/meshlab [Cignoni et al. 2021], OpenMesh [Botsch et al. 2002], libigl [Jacobson et al. 2016], PMP [Sieger and Botsch 2019], and OpenVolumeMesh [Kremer et al. 2013]. Commonly, mesh-editing algorithms are tightly coupled with a data-structure and its API, and porting an algorithm from one library to another is a major engineering effort. Code relying on local operations is also inherently error prone, as it usually involves keeping track of properties attached to mesh elements as the mesh itself changes due to the local operations. Parallelizing code using a mesh data structure is also challenging, due to race conditions when multiple threads attempt to change the same region of the mesh.

At a high-level, it is common practice to describe a mesh editing algorithm as a sequence of topological and geometrical editing operations. We argue that this approach is unnecessarily low-level, as it exposes the algorithm designer to technical problems that can be handled automatically by changing the abstraction level. It also makes it challenging to use or customize mesh editing algorithms in larger projects (such as their use in physical simulation for adaptive refinement), as low-level data structure details percolate in the entire code-base. A particularly difficult challenge in these algorithms is to ensure that a set of conditions (such as manifoldness, being free from self-intersections, minimal quality, maximal geometrical approximation) hold after each operation is applied. This is usually tackled by simulating each operation for the purpose of checking these

conditions, an error-prone process that needs to be carefully designed for each pair of operation and condition. The additional presence of attributes attached to vertices, edges, or faces further complicates these problems.

We propose a different way to describe mesh-editing algorithms on simplicial manifold meshes, using a declarative specification instead of a more traditional procedural approach. Instead of focusing on what the algorithm does, we ask the user to specify what are the requirements that the desired mesh should have. We divide these requirements into two groups: invariants and desiderata. The former is a description of hard requirements on the mesh (for example, no inverted elements or no self-intersections) and the latter is a set desirable properties (such as good quality). A mesh editing algorithm is then described as: (1) a set of per-element invariants (for example, all elements should have correct orientation), (2) a measure for the desiderata (for example, element quality), (3) a set of application-specific attributes attached to mesh elements, and how they are affected by local operations, and (4) a schedule of operation types. We show that many existing algorithms for mesh generation, remeshing, and parametrization, can be concisely expressed in this form (Section 5.4), which we denote IDAS (Invariant-Desiderata-Attributes-Schedule).

The IDAS specification has been designed with four goals:

1. **Modularity:** The connectivity of the mesh is abstracted from the user, which can only navigate the mesh using a high-level abstraction based on a cell tuple [Brisson 1989]. This reduces the learning curve for a new user, as they only need to learn a navigation API to implement algorithms in IDAS. It will also allow IDAS programs to benefit from continuous progress in data-structure design, as the data structure will be swappable without requiring downstream code changes in the high level IDAS code. This is in stark contrast with existing mesh libraries, which tend to be very invasive in the user code relying on them due to the close connection between navigation, mesh editing, and property management.
2. **Usability:** The user code sees, at all times, a valid mesh: the library simulates each operation transparently allowing the user to navigate on a valid mesh before and after every operation, dramatically simplifying the logic required to define invariants and desiderata. Properties on the mesh are also similarly abstracted, allowing to attach attributes on every simplex independently on the data structure used for implementing the specification.
3. **Efficiency:** The specification purposely requires only definitions of properties on individual elements. This feature allows runtimes for IDAS program to parallelize the computation (Section 5.4) without requiring special attention from an user. We demonstrate that automatic parallelization of mesh editing algorithms is possible on multi-core architectures.
4. **Robustness:** The IDAS specification moves the majority of the robustness issues typical of meshing algorithm on the runtime used to execute a IDAS program instead of the IDAS code itself. This simplifies the development of robust algorithms: for example, the user invariants are guaranteed to be enforced during processing, as the runtime will automatically check them on every modified element. As long as the user provides correct code for the invariant (for example to check for area positivity of an element using a predicate), then the runtime ensures that the invariants will be satisfied for all elements.

Given an algorithm in IDAS form, we design an algorithm and runtime library to realize it, with guarantees on satisfying the invariant and a best effort to maximize the desiderata. Our library exploits shared memory parallelism without any additional effort required from users in the algorithm specification.

To demonstrate the generality and effectiveness of our approach, we provide IDAS formulations for five popular mesh editing algorithms (Section 5.4, Figure ??): (1) shortest edge collapse [Hoppe 1996] (decimation for triangle meshes), (2) QSLim [Garland and Heckbert 1997], (3) isotropic triangle meshing [Botsch and Kobbelt 2004] (remeshing for triangle meshes), (4) harmonic triangulations [Alexa 2019] (quality improvement for 3D volumes), and (5) robust tetrahedral mesh generation [Hu et al. 2019d] (conversion of surface meshes to volumetric meshes). The IDAS formulation closely resemble the textual description of the algorithms in the corresponding papers: it is compact, readable, and easy to adapt for requirements of specific applications. As an example, we show that modifying (1) and (2) to guarantee a maximal geometric error requires is straightforward. Despite its generality, IDAS implementations executed using our library are comparable or faster than state of the art implementations in open-source software: the overhead due to the framework generality is more than compensated by the automatic parallelization (Section 5.4).

We believe our contribution is an important step to allow researchers and practitioners to effectively develop new mesh-editing algorithms, shielding the designer of mesh editing algorithms from many of the robustness and correctness challenges plaguing previous low-level approaches, by moving these components inside the runtime environment. It will also allow mesh editing algorithms to be used more easily in larger systems, as they can be tailored to requirements of a specific application with minimal programming effort.

We provide an open-source implementation of our library and of the five mesh editing algorithms as additional material.

5.2 RELATED WORK

5.2.1 MESH DATA STRUCTURES

Efficient data structures for representing solid geometry have been an intriguing research topic since the early days of computer graphics [Requicha 1980]. As a result, there is a large variety of mesh data structure designs, where they are each optimized for different usage scenarios. Index-array-based mesh data structure encodes each element as a list of vertex indices on its boundary. It is simple and memory efficient, but neighborhood query and local operations are not directly supported. Graph-based mesh data structures, including half-edge [Mäntylä 1987], winged-edge [Baumgart 1972], quad-edge [Guibas and Stolfi 1985], cell-tuple [Brisson 1989], etc., view meshes as graphs, where each element contains links to its adjacent elements. This design allows for efficient local query and update, making it ideal for algorithms like mesh simplification [Garland and Heckbert 1997]. Linear-algebra-based mesh data structures, such as [DiCarlo et al. 2014; Zayer et al. 2017; Mahmoud et al. 2021], encode adjacency information as sparse matrices. This design elegantly reduces neighborhood query and local operations to sparse matrix computations, which are highly optimized for modern parallel computing architecture. Closely related, is the concept of

generalized combinatorial maps [Lienhardt 1994; Dufourd 1991], and the CGoGN library [Kraemer et al. 2014] provide an efficient implementation which provide parallel traversal of the mesh. By design, mesh data structures provide a low-level interface to manipulate vertices, edges, faces, and tetrahedra. Different designs differ vastly in API and implementation details, making it hard to port algorithm from one data structure to another. In contrast, our framework decouples mesh data structure choice from algorithm specification, providing the flexibility of switching the underlying data structure in a seamless manner.

5.2.2 DOMAIN SPECIFIC LANGUAGES IN GRAPHICS

Our abstraction model of mesh processing algorithms draw inspiration from domain specific languages (DSL) in graphics. For dense regularly structured data such as images, Halide [Ragan-Kelley et al. 2013] popularized the idea of decoupling image processing operations from low level scheduling tasks. Similar abstraction that separates algorithm description from low level data structure and/or parallel architecture can also be found in other DSLs such as Simit [Kjolstad et al. 2016] for simulation over triangle meshes, Taco [Kjolstad et al. 2017] for dense and sparse tensor algebra, Taichi [Hu et al. 2019a] for simulation over sparse volumetric data, and Penrose [Ye et al. 2020] for generating diagrams from math notation.

5.2.3 PARALLEL MESHING

To meet the demand of generating large meshes, a number of popular mesh generation algorithms have been redesigned to leverage modern parallel computing hardware, both in a shared memory and distributed memory setting. Typically a divide-and-conquer strategy is adopted where a mesh is partitioned to run local processing operations on each subdomain in parallel. There are two key challenges involved: (1) how to handle operations involving elements shared by multiple partitions; (2) how to ensure load stay balanced across different processors as the mesh evolves.

One way to mitigate both challenges is to ensure mesh is partitioned into similar sized patches with high area to boundary ratio. A large number of partitioning strategies are available, including clustering-based approaches [Mahmoud et al. 2021], spacial-hierarchy-based approach [Loseille et al. 2017; Lo 2012], space-filling-curve-based approach [Marot et al. 2019; Borrell et al. 2018], and general purpose graph partitioning [Karypis and Kumar 1998]. Many variations of space-filling curves have also been used to construct mesh partitions [Chrisochoides 2006; Aluru and Sevilgen 1997]. To handle potential conflicts that may arise at partition boundaries, various synchronization strategies have been proposed [Okusanya and Peraire 1996; Chrisochoides and Nave 2003; Chrisochoides 2006] to minimize the amount of communication.

After generating the submeshes, some methods allow each compute node to work on them independently without synchronization. Once all threads are done, the meshes are merged [Cignoni et al. 1993; Chen 2010; Funke and Sanders 2017; Blelloch et al. 1999]. However these methods require complicated merge steps since the tetrahedra in the intermediate boundaries may not align. There are some techniques that compromise the Delaunay condition in some cases, so that the merging operation can be simpler [Lachat et al. 2014]. To avoid the tricky merge operations, other parallel strategies maintain a single complete Delaunay tetrahedralization and

use synchronization techniques to avoid race conditions when working on a partition boundary [Okusanya and Peraire 1997; Chrisochoides and Nave 2003]. The parallel constrained Delaunay meshing algorithm [Chew et al. 1997] cleverly defines the boundary and edge constraints to reduce the variable and unpredictable communication patterns. Some other techniques use locks for handling conflicts and data races [Blandford et al. 2006; Batista et al. 2010; Foteinos and Chrisochoides 2011].

Another set of methods use recursive divide-and-conquer techniques for parallel implementation on shared memory machines [Marot and Remacle 2020]. All threads independently work on the internal parts of the mesh and skip the operations at the boundary. After this phase, processing of only the boundary elements becomes the new problem. This technique is then recursively used until all the mesh elements are processed. A similar set of techniques use clever space-filling curves for re-partitioning the mesh boundaries after each recursive phase [Chrisochoides 2006; Aluru and Sevilgen 1997].

Since the submesh boundaries are the main areas of concern, some methods entirely avoid any operations on these boundaries while ensuring the correctness of the result [Galtier and George 1996; Linardakis and Chrisochoides 2006]. These methods precompute the domain separators such that their facets are Delaunay admissible. This completely eliminates synchronization overheads, but only applies for Delaunay meshing.

Another conflict handling strategy is to simply reject the offending operations and try executing them later with a new domain partitioning [Marot et al. 2019]. This reject-and-repartition strategy may not guarantee algorithm termination, thus special care is needed to handle this case.

As the domain mesh evolves, keeping load balanced across processors becomes critical. Typically, this is done by periodically repartitioning the updated mesh. Zhou et al. [2012] proposes a predictive load balancing method to keep partitions balanced. Marot et al. [2019] uses simple rescaling of the space-filling curve to repartition the domain.

In this work, we are targeting only shared-memory parallelism, thus making the problem of reducing communications between processors less relevant. We use a graph-based space partitioning technique [Karypis and Kumar 1998] due to its simplicity and availability as open-source code (METIS), but we use it only to reduce the risk of conflicts. To avoid conflicts, we use a shared memory locking mechanism. This approach is only possible for shared-memory parallelism but has the major advantage of not requiring rebalancing and to respect, to a certain degree, the execution order prescribed by the user-code. This approach is possible thanks to the availability of efficient parallel atomic instructions, and parallel libraries based on them (oneTBB).

5.2.4 SCOPE OF MESH EDITING

MESH GENERATION Tetrahedral meshing algorithms heavily rely on mesh editing operations. The most common approaches are Delaunay methods [Shewchuk 1998b; Ruppert 1995b; Remacle 2017; Du and Wang 2003b; Alliez et al. 2005b; Tournois et al. 2009b; Murphy et al. 2001b; Cohen-Steiner et al. 2002b; Chew 1989b; Si and Gartner 2005; Shewchuk 2002b; Si and Shewchuk 2014b; Cheng et al. 2008b; Boissonnat and Oudot 2005b; Jamin et al. 2015b; Dey and Levine 2008; Chen and Xu 2004; Shewchuk 1996b; Cheng et al. 2012; Bishop 2016; Busaryev et al. 2009; Boissonnat et al. 2002; Si 2015a], which strive to generate meshes satisfying the Delaunay condition, grid

methods [Yerry and Shephard 1983; Bern et al. 1994; Molino et al. 2003; Bronson et al. 2013; Labelle and Shewchuk 2007; Doran et al. 2013; Bridson and Doran 2014], which start from a regular lattice or with a hierarchical space partitioning and optionally intersect the background mesh with the input surface, and front-advancing methods [Sadek 1980; Cuilliere et al. 2013; Alauzet and Marcum 2014; Haimes 2014], which insert one element at a time, growing the volumetric mesh (i.e. marching in space), until the entire volume is filled .

These algorithms rely on local operations on mesh data-structures, and benefit from our framework to simplify the implementation and gain automatic parallelization. We discuss an implementation of one the more recent algorithms [Hu et al. 2018, 2019d] in Section 5.4. Note that some of these algorithms use local operation that are not implemented yet (such as 5-6 swap), but they could be added to our framework.

CONSTRAINED MESHING. Downstream applications often require meshes to satisfy either quality (avoidance of zero volume elements) or geometric (distance to the input surface) constraints. For example, Mandad et al. [2015b] creates a surface approximation within a tolerance volume, the TetWild algorithms [Hu et al. 2018, 2019d] use an envelope [Wang et al. 2021] to restricts the geometry of the boundary of the tetrahedral mesh, [Brochu et al. 2012] adds constraints to local remeshing to avoid interpenetrations in simulations, and [Gumhold et al. 2003] extends mesh simplification [Garland and Heckbert 1999; Popović and Hoppe 1997] to ensure a non self-intersecting result.

These criteria are explicitly modeled as invariants in our framework, and they can be easily swapped in and out existing implementations, as we demonstrate in Section 5.4.

MESH IMPROVEMENT. Mesh improvements modifies an existing mesh by changing its connectivity and position of the vertices to improve the quality of its elements [Canann et al. 1996, 1993; A. Freitag and Ollivier-Gooch 1998; Lipman 2012; Chen and Xu 2004; Alliez et al. 2005b; Feng et al. 2018b; Hu et al. 2018; Alexa 2019; Klingner and Shewchuk 2007]. We show in Section 5.4 a reimplementation of [Alexa 2019] in IDAS form.

DYNAMIC REMESHING AND ADAPTIVE MESH REFINEMENT (AMR) Simulations involving large deformations are common in computer graphics, and if the surface or volume deformed is represented by a mesh, it is inevitable that after a large deformation the quality of the elements will deteriorate, and the mesh will have to be updated. Additionally, it is often required to concentrate more elements in regions of interest whose location is changing during the simulation, for example to capture a fold in a cloth simulation, or a fracture in a brittle material. These two challenges are tackled in elastoplastic and viscoplastic simulations [Hutchinson et al. 1996; Bargteil et al. 2007; Wicke et al. 2010; Wojtan and Turk 2008], in fluid simulations [Misztal and Bærentzen 2012; Klingner et al. 2006; Ando et al. 2013; Chentanez et al. 2007; Clausen et al. 2013], in cloth simulation [Villard and Borouchaki 2002; Bender and Deul 2013; Li and Volkov 2005; Narain et al. 2012, 2013; Pfaff et al. 2014; Simnett et al. 2009], and fracture simulation [Busaryev et al. 2013]. All these algorithms could benefit from our contribution, to simplify their implementation and obtaining speedup due to the automatic parallelization offered by our approach.

A different approach is discussed in [Grinspun et al. 2002], where the refinement is performed on the basis to avoid the difficulties with explicit remeshing. However, this approach cannot coarsen a dense input, and also cannot increase the quality of elements, making it usable only for specific scenarios [Grinspun et al. 2002]. Our approach aims at lowering the barrier for integrating explicit remeshing algorithms in simulation applications, thus allowing to directly use standard simulation methods on adaptive meshes without having to pay the high implementation cost for the mesh generation.

When remeshing is paired with algorithms simulating contacts that do not tolerate interpenetrations (for example [Li et al. 2020]), it is necessary to ensure that adaptive remeshing does not break this invariant. This can be achieved adding non-penetration constraints to each local mesh editing operations, as proposed in [Brochu et al. 2012]. Our framework is ideal for developing such methods, as additional constraints can be added to existing mesh editing algorithms with minimal modifications, as we demonstrate in Section 5.4.

PARAMETRIZATION Conformal mesh parametrization algorithms adapt the mesh during optimization, as a fixed triangulation restricts the space of metrics realizable [Luo 2004; Campen and Zorin 2017; Campen et al. 2019; Gu et al. 2018b,a; Springborn 2020; Sun et al. 2015]. Two very recent works [Gillespie et al. 2021; Campen et al. 2021] introduce robust algorithms based on Ptolemy flips to compute conformal maps satisfying a prescribed metric.

All these methods require changing the mesh connectivity of a triangle mesh, and could thus benefit from our framework to simplify their implementation and parallelize the mesh editing operations.

MESH ARRANGEMENTS/BOOLEAN OPERATIONS Boolean operations are basic algorithms often used in geometry processing applications. Recently, [Zhou et al. 2016] proposed a robust way to compute them by constructing a space arrangement, and then filtering the result using the generalized winding number [Jacobson et al. 2013]. A similar approach, using an approximated meshing algorithm, has been extended in [Hu et al. 2019d], using a tetrahedral mesher to create the initial arrangement. The reimplementation of TetWild introduced in this paper (Section 5.4) can be extended for a similar purpose.

5.3 METHOD

Our declarative specification is designed to remove the burden of low-level management of the mesh connectivity and attributes, allowing an algorithm designer to focus only on high-level requirements. The design consists of five components.

5.3.1 MESH EDITING COMPONENTS

OPERATION ROLLBACK. It is common to perform mesh editing to improve a given energy functional, such as mesh quality or element size. However, due to the discrete nature of the operations, it is not possible to use standard smooth optimization techniques, and instead the effect of the

wmtk-tex/figs/pipeline_illustration.pdf

Figure 5.1: Overview of the components behind our specification. A mesh is represented through its topology, implemented by our library, and a list of user provided attributes. Before an operation is attempted, we explicitly perform a pre-check, and, if successful, we generate a mesh (attributes and topology). At this point, we trigger the after check to validate the operation (e.g., check if the newly generated mesh has positive volume). In case the after check fails, we *automatically* rollback the operation and restore the mesh to its previous *valid* state.

energy is evaluated before and after every operation to measure its effect on the energy. This is commonly implemented using an ad-hoc energy evaluation that “simulates” the operation only for the purpose of measuring the energy change. This simulation is complex (especially in 3D), and error-prone, as not only the connectivity changes, but the energy likely depends on properties attached to mesh elements, which needs to be updated accordingly.

We propose instead to make this process opaque to the user, providing to the user-code an explicit copy of the mesh (and up to date attributes) *before and after* the operation is performed to allow an easy and reliable energy evaluation. The correctness and efficiency of this process is handled by the runtime. This reduces the complexity of mesh editing considerably in our experience, as it makes them more similar to traditional finite difference approaches where the energy is evaluated on different points on the domain to approximate its derivative.

EXPLICIT INVARIANTS. It is common to have a set of desiderata on the mesh that needs to be satisfied, such as avoiding triangle insertions or self-intersections. Given the complexity of a mesh editing algorithm it is difficult to ensure that they are satisfied, as these conditions need to be checked after every operation is applied (and they often depend on attributes too, such as vertex positions).

We propose to make these invariants explicit, and delegate to the library the task of ensuring that they are checked after every mesh modifications, and after the input is loaded. In this way,

not only the code is simpler, but it is much easier to ensure correctness, as the checks are handled transparently by the library.

EXPLICIT ATTRIBUTE UPDATE. Mesh attributes are usually handled by low-level meshing libraries, allowing to attach them to the desired mesh element (vertex, edge, face, triangle, or tetrahedra). However, the handling of attributes after a local operation is performed is usually a responsibility of the user code, as it is dependent on the application.

We propose to make this process more explicit, requiring the user to provide the rules on how to update attributes after operation in a high-level specifications, and delegating the actual update to the library. This makes the specification more direct and less error-prone, and allows users to write algorithms without having to know the low-level details on how the local mesh operations work.

PARALLEL SCHEDULING. The type and scheduling of local operations is crucial in mesh editing algorithms. It usually involves maintaining a priority queue of operations, which is updated after every local operation.

We provide a direct way of controlling the operations performed and how the queue is updated. In the library, we can then distribute the work automatically on multiple threads, hiding from the user code the complexity of performing mesh editing operations in parallel and ensuring that race conditions are avoided.

ABSTRACT MESH NAVIGATION. Both invariant and attribute updates require navigating a mesh. Instead of relying on data-structure specific navigation, we favor the use of the cell tuple abstraction [Brisson 1989]. This allows the specification to be independent of the mesh data structure used in the library. The Tuple stores four indices (three for surface meshes), vertex, edge, face, and tetrahedron and provides a single function per index, called `switch`, to change one index while keeping the other indices fixed. For instance `switch_vertex` changes the vertex index while keeping edge, face, and tetrahedron fixed which has the effect of selecting the opposite vertex on an edge.

5.3.2 DECLARATIVE SPECIFICATION

Our API provides two abstractions: a `TetMesh` (and `TriMesh` class for 2D) (Algorithm 1), and a `Scheduler` (Algorithm 2).

MESH CLASSES Both the `TetMesh` and `TriMesh` classes provide the basic local operations (e.g., edge split or collapse) and, for each operation, their corresponding *before* and *after* methods. The mesh class is responsible of implementing the operations changing the topology, and the application code must *only* override the before and after methods to update attributes. The before method has a view of the mesh before the operation, and can thus navigate it to cache local attributes, while the after method has a view after the operation is performed, and it is responsible for updating attributes. In the simple case of regular subdivision of a triangle mesh, the

```

class TetMesh
{
public:

    bool split_edge(const Tuple& t,
                    std::vector<Tuple>& new_tets);
    bool collapse_edge(const Tuple& t,
                       std::vector<Tuple>& new_tets);
    bool swap_edge(const Tuple& t,
                   std::vector<Tuple>& new_tets);
    bool swap_face(const Tuple& t,
                   std::vector<Tuple>& new_tets);
    bool smooth_vertex(const Tuple& t);

    bool insert_triangle(
        const std::vector<Tuple>& intersected_tets,
        const std::vector<Tuple>& intersected_edges);

protected:

    bool invariants(const std::vector<Tuple>& tets);

    bool split_before(const Tuple& t);
    bool split_after(const Tuple& t);

    bool collapse_before(const Tuple& t);
    bool collapse_after(const Tuple& t);

    bool swap_edge_before(const Tuple& t);
    bool swap_edge_after(const Tuple& t);

    bool swap_face_before(const Tuple& t);
    bool swap_face_after(const Tuple& t);

    bool smooth_before(const Tuple& t);
    bool smooth_after(const Tuple& t);

    bool insert_triangle_before(
        const std::vector<Tuple>& faces);
    bool insert_triangle_after(
        const std::vector<Tuple>& faces,
        const std::vector<std::vector<Tuple>>& new_f);
};


```

Algorithm 1: API of our TetMesh class.

```

template <class Mesh>
struct Scheduler
{
    function<double
        (const Mesh&, Op op, const Tuple&)>
    priority = ...;

    function<vector<pair<Op, Tuple>>>
        (const Mesh&, Op, const vector<Tuple>&)
    renew_neighbor_tuples = ...;

    function<vector<size_t>
        (Mesh&, const Tuple&)>
    lock_vertices = ...;

    function<bool(const Mesh&)>
    stopping_criterion = ...;

    function<bool
        (const Mesh&, tuple<double, Op, Tuple>& t)>
    should_process = ...;

    size_t num_threads = ...;
    size_t max_retry_limit = ...;
    size_t stopping_criterion_checking_frequency = ...;

    bool operator()
        (Mesh& m,
        const vector<pair<Op, Tuple>>& ops);
};


```

Algorithm 2: API of our Scheduler.

`split_before` caches the coordinates of the two edge endpoints, and the `split_after` computes the position of the newly inserted vertex by averaging them.

In addition, the mesh class provides a method, which can be overridden by the user code, that *automatically* verifies user-provided invariants (e.g., maintain positive elements' volume). All user-provided methods return a Boolean status to notify the mesh classes if the operation fails; in case it does, our API rolls back the operation and restores the topology to the previous valid state. As the connectivity and attributes management is handled by the class, this ensures that, in case of failure of the operation, the mesh will go back to a valid state.

Our API provides the standard local operations: edge collapsing, edge/face swapping, edge splitting, and smoothing. We also provide an additional, non standard, operation: triangle insertion. This operation is an enhanced version of splitting where multiple edges, faces, and tetrahedra are subdivided to represent an input triangle provided as input. This operation is useful to compute mesh arrangements, and it is also used in meshing algorithms [Hu et al. 2019d].

Since the `TetMesh` class only handles topology, the operation requires the list of edges and tetrahedron the input triangle intersects. Internally it subdivides all of them, and generates a valid tetrahedral mesh using the connectivity table in [Hu et al. 2019d]. The before operation provides the user the list of faces that will be changed by the operation, allowing the user code to explore the mesh and cache attributes, while the after provides a mapping between the old faces and any newly inserted face in the mesh.

SCHEDULER The second part of our API is the Scheduler that is responsible to control the order of the individual operations and then execute a list of operations. The main purpose of the scheduler is to abstract the operation order and hide parallelization details from the user. Our scheduler provides customizable callbacks, including,

- *Priority* to order the local mesh edit operations.
- *Renew neighbor tuples* that is invoked after a successful operation, to add newly created tuples and operations into the queue.
- *Lock vertices* that provides information on the affected region for the operation, and avoiding conflicts.
- *Stopping criterion* that is checked periodically to terminate the program if certain criterion is met. For example, number of vertices, or quality criterion.

5.3.3 IMPLEMENTATION.

We implement a runtime for our specification in C++, using Intel oneTBB for parallelization.

DATA STRUCTURE. We opt for an indexed data structure, where we explicitly represent the vertices and the simplex of higher dimension (triangle for 2D, tetrahedra for 3D). Each vertex explicitly stores a list of incident simplices, and each simplex stores a sorted list of its vertices.

```
wmtk-tex/figs/lock_illustration.pdf
```

Figure 5.2: Example of the locking region for two edges. In the example the operation requires locking a two-ring neighborhood (e.g., for the edge collapse operation). If the two edges are sufficiently far (right) both operations can be safely executed in parallel. When the two edges are close (left) the operations might fail acquiring the mutexes in the shared area.

While not the most efficient option for navigation, this data structure makes the implementation of local operation much simpler.

PARALLELIZATION. To avoid conflicts between local operations working on the same part of the mesh, we introduce a synchronization mechanism using locks.

Each mesh vertex is associated with a mutex. Whenever a thread wants to access (read/write) any attribute stored in a vertex, edge, triangle, or tetrahedron it must first acquire a lock on *all* the vertices of the tetrahedron containing the element(s) storing the attribute (Figure 5.2). For example, if a thread wants to read a value on an edge of a 3D mesh, it first needs to acquire a lock on all vertices of all the tetrahedra containing that edge. This mechanism is used also for mesh navigation, and for updating the mesh connectivity.

At a first look, this mechanism might seem cumbersome and expensive. However, we rely on asynchronous, tentative lock acquisition operations. We *try* to acquire the lock, and give up if the lock is already taken by another thread. These operations are efficient on modern hardware and dramatically improve the performance, while avoiding deadlocks. A downside is that an operation might be skipped due to impossibility of acquiring a mutex. These operations are retried for several times (by default 10 times) and run serially if they still do not succeed. Before performing any local operation, we try to acquire the lock on vertices in the 1-ring or 2-ring of the vertex involved in the operation. For example, a vertex smoothing operation requires acquiring the 1-ring vertex neighborhood of the smoothed vertex, while an edge-collapse operation on an edge (v_1, v_2) requires acquiring the lock on the 2-ring vertex neighborhood of both v_1 and v_2 (Figure 5.2).

Finally, since we partition the input mesh using Morton Encoding [Karras 2012], the amount of conflicts (and skipped operations) is low.

5.3.4 EXAMPLE: SHORTEST EDGE COLLAPSE

We show how the library is used in a classical example, shortest edge collapse. In this case, we add a 3D position to every vertex as a vertex attribute (by default, there are no attributes attached to mesh elements). For every attribute and for every operation we plan to use in the scheduler, we

need to provide a function that updates such attribute (Algorithm 3). In the `collapse_before` function, we cache the two vertex coordinates associated with the collapsed edge represented by `Tuple t`. In the `collapse_after` function, we generate a new vertex in the middle of the two endpoints of the collapsed edge.

```

//Save two vertices attached to edge t
bool collapse_before(const Tuple& t)
{
    cache.v1p = verts[t.vid()];
    cache.v2p = verts[switch_vertex(t).vid()];
    return true;
}

//Generate a new point
bool collapse_after(const Tuple& t)
{
    verts[t.vid()] = (cache.v1p + cache.v2p) / 2.0;
    return true;
}

```

Algorithm 3: Overridden methods in `TriMesh` sub-class to implement shortest edge collapse.

Equipped with the 3D position attribute, which at this point will be automatically kept up to date by the library, we can now schedule the collapse operation (Algorithm 4). For shortest edge collapse, we want to attempt to collapse all edges, prioritizing the shortest ones, until we reach a fixed number of collapses `n_collapse`: in the code we registering the operation type (`ops`), specify how to update the queue after an operation (`renew`) by adding all the neighbouring edges, and specifying the edge length as a priority (`priority`). Note that the outdated elements in the queue that are affected by a local operation are automatically invalidated using a tagging mechanism on the tuples which is opaque to the user.

5.4 APPLICATIONS

To showcase the generality and effectiveness of our approach, we implement five popular mesh editing algorithms in our framework, and compare them with reference implementations. Overall, the performance of our method are competitive for surface applications, but the overhead due to the approach generality is higher in 3D, leading to higher running time.

SHORTEST EDGE COLLAPSING The simplest algorithm for simplifying a triangle edge is shortest edge collapse [Hoppe 1996], which performs a series of collapse operations prioritizing the shorter edges. The algorithm requires only one local operation, edge collapse. A common criteria for termination is reaching a desired number of mesh elements. We compare our implementation with

```

//Collect edges attached to tris
vector<Tuple> new_edges_after(const vector<Tuple>& tris)
{
    vector<Tuple> new_edges;
    for (auto t : tris) {
        for (auto j = 0; j < 3; j++) {
            new_edges.push_back(
                tuple_from_edge(t.fid(), j));
        }
    }
    return new_edges;
}

bool collapse_shortest(int n_collapses)
{
    //Register operations
    auto ops = vector<pair<Op, Tuple>>();
    for (auto& l : get_edges())
        ops.emplace_back("edgeCollapse", l);

    //After a successful operation,
    //we append all new edges
    auto renew = [] (auto& m, auto op, auto& tris) {
        auto edges = m.new_edges_after(tris);
        auto optup = vector<pair<Op, Tuple>>();
        for (auto& e : edges)
            optup.emplace_back("edgeCollapse", e);
        return optup;
    };

    //priority in which we collapse
    auto priority = [] (auto& m, auto op, const Tuple& e) {
        const auto v1 = m.verts[e.vid()];
        const auto v2 = m.verts[e.switch_vertex(m).vid()];
        auto len2 = (v1 - v2).squaredNorm();
        return -len2;
    };

    //Set the functions to the scheduler
    Scheduler executor;
    executor.renew_neighbor_tuples = renew;
    executor.priority = priority;
    executor.stopping_criterion_checking_frequency =
        n_collapses;
    //We stop only when we perform 95 collapses
    executor.stopping_criterion =
        [] (auto& m) { return true; };
    //Run the executor
    executor(*this, ops);
}

```

wmtk-tex/figs/2d-sec-statue.pdf

Input	libigl	Output	Timing
-------	--------	--------	--------

Figure 5.3: Comparison of our parallel implementation (32 threads) of shortest edge collapse (scalability plot on the right, from 1 to 32 threads) of a model with 281,724 faces with the serial version in libigl. Both libigl and our output have 28,168 faces and comparable edge length (1.058 for libigl versus 1.061 for ours). Our serial method runs in 4.9s (5.84s on a single thread, 0.52s with 32 thread, leading to a speedup of 11 \times), while libigl runs in 2.74s.

the “decimate” implementation in libigl [Jacobson et al. 2016]. The serial libigl implementation is comparable when running the algorithms serially, and our parallel implementation is up to 9 times faster when using 32 threads (Figure 5.3).

QSLIM We use our framework to implement QSlim [Garland and Heckbert 1997]. QSlim collapse edges based on the planarity of the two adjacent faces measured with an error quadric. The algorithm continues to collapse until it reaches a target number of edges. We compare our implementation with the QSlim implementation in libigl [Jacobson et al. 2016]. The serial libigl implementation is 8 times faster than our implementation, due to their direct manipulations of elements in the queue with each collapse. But our parallel implementation is twice as fast when using 16 threads (Figure 5.4).

ISOTROPIC REMESHING We implemented the widely used algorithm for isotropic remeshing proposed in [Botsch and Kobbelt 2004]. This algorithm alternates edge collapse, edge flips, edge splits, and tangential smoothing to obtain a mesh that is isotropic (i.e. all elements have the same size) and where all triangles are close to equilateral. The process is guided by a user-provided target edge length L , and terminates when no local operation leads to either an improvement in the desired edge lengths or an improvement in vertex valence [Botsch and Kobbelt 2004].

In Figure 5.5, we compare our implementation of [Botsch and Kobbelt 2004] with the implementation in OpenFlipper [Möbius and Kobbelt 2010]. The OpenFlipper implementation is 2.5

wmtk-tex/figs/2d-qslim.pdf

Figure 5.4: Comparison of our parallel implementation of QSLim with the serial version in libigl for a model with 1,909,755 faces. Top right, the libigl output has 17,891 faces and takes 41.26s. Bottom left, our output has 17,906 faces and runs in 306.59s. Our implementation scales well: 347.59s with one thread and 13.88s with 32 (25× speedup).

wmtk-tex/figs/2d-remeshing-lucy.pdf

Figure 5.5: Example of uniform remeshing a model with 2,529,744 triangles (left) with the same target edge length. Middle, [Möbius and Kobbelt 2010] remeshes it to 78,322 faces in 31.96 seconds. On the right is our 32-thread implementation, which generates 71,640 triangles in 8.2 seconds (78.34s serial, 95.0s for a single thread, leading to a speedup of 11 \times). The difference in density of the meshes is due to differences in the detail of the implementation, which makes the two methods reach an average vertex valence of 5.999, and a similar target edge length (differ 0.01% of the bounding box diagonal length) with a different element budget.

times faster when running on a single thread, and our implementation becomes faster after 4 threads are used (Figure 5.5).

HARMONIC TRIANGULATIONS The *harmonic triangulations* algorithm has been introduced as an alternative to sliver exudation in the Delaunay tetrahedralization pipeline to efficiently reduce sliver tetrahedra. The original paper [Alexa 2019] proposes to use both flip and smoothing operations.

The code provided by the authors implements a reduced version of the algorithm proposed in the paper, restricting the optimization to 3-2 edge swap operations. We thus implemented both a reduced version for a fair comparison (Figure 5.6) and a complete version. Our more generic framework is twice as slower than the hand-optimized code written by the authors when running serially, and it is 2 times faster when running on 32 threads (Figure 5.6).

TETRAHEDRAL MESHING The TetWild algorithm is a tetrahedral meshing algorithm with minimal input requirements: given an input triangle soup, it can generate a tetrahedral mesh which approximates its volume. We take inspiration from the original algorithm introduced in [Hu et al. 2018, 2019d] with a few modifications: (1) we use the insertion operation [Hu et al. 2019d] (using rational coordinates) as a replacement for their BSP partitioning, as this simplifies the implementation, (2) we use the envelope proposed in [Wang et al. 2020] instead of sampling, and (3) we use 2-3 face

wmtk-tex/figs/3d-harmonic-gauss.pdf

Reference	Output	Timing
-----------	--------	--------

Figure 5.6: Example of *Harmonic Triangulations* starting with one million Gaussian distributed random points. Both our and the reference implementation reach a similar target number of tetrahedra (5.9 million for the reference and 6.1 million for ours, due to a difference in operation ordering) and a similar Mean Harmonic Index (0.547 for the reference and 0.554 for ours). Our method takes 3.82s with 32 threads (15.49s serial, 40.49s on a single thread, speedup of 11 \times), while the reference serial implementation takes 6.37s.

swapping, 3-2 and 4-4 edge swapping operations, to simplify the implementation. We show results on two models in Figure 5.7: the results are very similar to the original implementation, and our version is 2 times faster when using 8 threads. We experimentally observe that our framework scales well up to 8 threads, after that the algorithm becomes slower. This is because, as we increase the number of threads and partitions, the frequent conflict in tetrahedral mesh edge operations affects the parallel performance. We believe that this observation might be useful for the future design of high performance concurrent mesh generation algorithms.

5.4.1 PARALLELIZATION

Enabling the parallelization mechanism introduces a minor slowdown as visible in the difference between the pure serial and one thread timings on our applications, due to the additional cost of allocating mutexes and to acquire them. The algorithm scales well in all 5 applications (figures 5.3, 5.4, 5.5, 5.6, 5.7), obtaining a scaling speedup. We would like to remark that thanks to our specification and our runtime, the serial and parallel implementation of the 5 algorithms above is almost identical.

An inevitable drawback of parallelization is that the algorithms cannot efficiently preserve ordering. For instance, in shortest edge collapse, every thread will try collapse edges in its own partition independently from the others. If one of these collapses append on the partition's interface, the thread will need to acquire a lock. In case of failure the collapse is postponed to a later stage thus not respecting the order (Figure 5.8). This is a rare event that is more problematic for fast operations.

wmtk-tex/figs/3d-tetwild-dragon.pdf

Figure 5.7: Tetrahedralize a surface with 856, 294 faces. Original TetWild (top right) generates a mesh with 56, 761 tetrahedra in 287.58s; our reimplementation (bottom left) generates a mesh with 44, 866 tetrahedra in 153.33s with 8 threads (452.42s serial, 521.47s on a single thread, speedup of 3.4 \times). The difference in number of tetrahedra is likely due to the different order of scheduling of operations due to the partitioning.

5.4.2 ALGORITHM MODIFICATIONS

A major motivation to invent and develop this declarative language is enabling easy customization of meshing algorithms. As an example, we add an additional termination criteria to the shortest edge collapse and uniform surface refinement. Integrating the envelope check is straightforward with our approach, as it only requires adding the envelope check to the invariants. We use the open-source library proposed in [Wang et al. 2020], which allows to directly specify the maximal allowed surface deviation. The envelope adds a noticeable computational cost, which is ameliorated by our parallel implementation (figures 5.9 and 5.10).

5.4.3 LARGE-SCALE DATASET VALIDATION

To validate our framework we run our reimplementation of uniform remeshing and tetrahedral meshing on the Thingi10k dataset [Zhou and Jacobson 2016]. We run all experiments serially on

wmtk-tex/figs/limitation-of-para.pdf

Figure 5.8: Example of splitting the longest edge on the bottom of the triangle. On the left, the edge can be split to generate the dark blue edge. In the next iteration, the left edge is split (by the light blue edge) leading to a decreasing maximum edge-length. On the left, the bottom edge is locked illustrated by a dashed line. In this case the next iteration splits the left edge (dark blue edge) and so on. As long as the dashed edge is locked it will never be split preventing the maximum edge length to decrease.

wmtk-tex/figs/2d-sec-envelope.pdf

Input

Output

Timing

Figure 5.9: Shortest edge collapse with envelope containment of a model with 857, 976 faces. Our method successfully generates a mesh with 71, 298 faces in 37.49s with 32 threads (731.32s serial, 725.34s on a single thread, speedup of 20 \times).

wmtk-tex/figs/2d-remeshing-headport-env.pdf

Input	Output	Timing
-------	--------	--------

Figure 5.10: Uniform remeshing with envelope containment check of a model with 198,918 faces. Our method produces a mesh with 68,202 faces in 29.11s with 32 threads and 493.54s for a single thread (483.68s for the serial version) leading to a speedup of 16 \times .

wmtk-tex/figs/tri2d-stats.pdf

Figure 5.11: Timings, target edge length ratio, and valence for every model in the dataset. Most models finish within a minute. The target edge length ratio measure how well our algorithm simplifies the meshes to reach the desired edge length, with an optimal value of 1. Since uniform remeshing strives to generate regular meshes, for most model our algorithm is able to obtain the optimal valence of 6.

an individual node of an HPC cluster an Intel Xeon Platinum 8268 24C 205W 2.9GHz Processors limiting the runtime to 15 hours.

For uniform remeshing, Figure 5.11 shows the time, average edge length normalized by the target, and average valence of isotropic remeshing on the ten thousand models. Most of our models finish within 10 seconds with only a few requiring more than a minute. For almost all meshes, the algorithm succeeds at reaching the target edge length and valence of 6.

For TetWild we limit the number of iterations to 25 (Figure 5.12). We note that within the 15 hours limit only 2.5% models did not finish, and after 25 iterations 3% of the models still have some rational coordinates. Among the successful models, most finish within 20 minutes and succeed in achieving high-quality meshes (only 8 models have average AMIPS larger than 10).

wmtk-tex/figs/tet3d-stats.pdf

Figure 5.12: Timing, max and average AMIPS energy (capped at 20) for maximum 25 iterations of tetrahedral meshing. Most models finish within 20 minutes with only a few taking up to a day. Even by limiting the iterations to 25, most models reach an average AMIPS energy lower than 10, with optimal value at 3.

5.5 CONCLUDING REMARKS

This paper introduces a new declarative specification for mesh algorithms to allow an easier implementation, while at the same time obtaining competitive performances and exploiting parallel hardware.

Using this specification, we implement five popular mesh editing algorithms covering mesh generation and optimization on surfaces and volumes, which can be easily adapted for other use cases: we demonstrated that integrating an envelope check requires only a few lines of code.

The library we implemented supports shared memory parallelism, which leads to a good scaling on the machines we tested it on. We believe an exciting venue for future work would be the implementation of a library for our specification targeting MPI to distribute the computation over an HPC cluster. Having access to such a library would allow our five mesh editing applications to run on a distributed environment with minimal or no changes.

6 | CONCLUSION

6.1 OVERVIEW

This chapter relates to the outlook and conclusions.
Bijective Parameterization has received follow ups.

6.2 FUTURE DIRECTIONS

Geometry + Simulations
Geometry + Vision construction

A | SHELL SPECIFICS

A.1 PROOFS

A.1.1 THEOREM 3.2

Consider the unique (up to symmetry) piece-wise affine map T deforming Δ into a reference prism $\hat{\Delta} = \{u, v \geq 0 | u + v \leq 1\} \times [-1, 1]$, identifying the bottom surface of $\hat{\Delta}$ as $z = -1$, middle surface as $z = 0$, and top surface as $z = 1$. Let T_T be the affine map mapping tetrahedron T to the reference prism. Since the volume of all T is positive by assumption, the map T is bijective and orientation preserving. In particular, T transforms $\mathcal{V}(p)$ to the const vector $e_z = (0, 0, 1)$ since the edges of the prism are mapped to axis aligned edges of the reference prism. The projection operator $\mathcal{P}(p)$ can thus be equivalently defined as $\mathcal{P}(p) = T^{-1}(\mathcal{P}_z(T(p)))$, where \mathcal{P}_z is the projection over the z -axis in the reference prism. \mathcal{P}_z is bijective if the piecewise linear mesh intersecting Δ is composed of triangles with positive area (and the boundaries are mapped to the boundaries) [Lipman 2014], after being mapped to the reference Δ and projected by \mathcal{P}_z . In the reference domain, having positive area after projection (with a fixed boundary) is equivalent to that the dot product between the projection direction and the normal of every face is positive. In the top half of Δ , \mathcal{P}_z is bijective if for every point $T(p), p \in f T(n_p) \cdot [0, 0, 1] = T(n_p) \cdot T(\mathcal{V}(p)) > 0$ which holds since $n_p \cdot \mathcal{V}(p) > 0$ (from the definition of section) and the fact that T is orientation preserving and thus not changing the sign of the dot product.

A.1.2 PROPOSITION 3.4

For simplicity, we will show that a uniform lower bound δ for all vertices exists. By consider a triangle of T , extruded over the displacement field \mathcal{N} , we obtain a generalized half prism with six vertices $0, e_1, e_2, \delta n_0, e_1 + \delta n_1, e_2 + \delta n_2$, with δ a positive constant. The indices of the tetrahedra are:

$$(0, 2, 3, 4), (0, 3, 4, 5), (0, 1, 5, 3), (0, 1, 2, 3), (1, 2, 3, 4), (0, 1, 2, 4), (0, 1, 5, 4), (1, 3, 4, 5), (1, 2, 3, 5), (0, 2, 5, 4), (0, 1, 2, 5), (2, 3, 4, 5).$$

For a sufficiently small δ , the linear term will dominate the higher power of δ , which we omit. The dominant volume terms are listed in the following table: ($e_{ij} := e_j - e_i$)

(0 ,2, 3, 4)	vol1($0, e_2, \delta n_0, e_1 + \delta n_1$)	$\delta \langle e_2, n_0, e_1 \rangle$
(0 ,3, 4, 5)	vol2($0, \delta n_0, e_1 + \delta n_1, e_2 + \delta n_2$)	$\delta \langle n_0, e_1, e_2 \rangle$
(0 ,1, 5, 3)	vol3($0, e_1, e_2 + \delta n_2, \delta n_0$)	$\delta \langle e_1, e_2, n_0 \rangle$
(0 ,1, 2, 3)	vol4($0, e_1, e_2, \delta n_0$)	$\delta \langle e_1, e_2, n_0 \rangle$
(1 ,2, 3, 4)	vol5($e_1, e_2, \delta n_0, e_1 + \delta n_1$)	$\langle e_{12}, \delta n_0 - e_1, \delta n_1 \rangle$ $= \delta \langle e_2, -e_1, n_1 \rangle$
(0 ,1, 2, 4)	vol6($0, e_1, e_2, e_1 + \delta n_1$)	$\delta \langle e_1, e_2, n_1 \rangle$
(0 ,1, 5, 4)	vol7($0, e_1, e_2 + \delta n_2, e_1 + \delta n_1$)	$\langle e_1, e_2 + \delta n_2, \delta n_1 \rangle$ $= \delta \langle e_1, e_2, n_1 \rangle$
(1 ,3, 4, 5)	vol8($e_1, \delta n_0, e_1 + \delta n_1, e_2 + \delta n_2$)	$\delta \langle \delta n_0 - e_1, n_1, e_{12} + \delta n_2 \rangle$ $= \delta \langle -e_1, n_1, e_2 \rangle$
(1 ,2, 3, 5)	vol9($e_1, e_2, \delta n_0, e_2 + \delta n_2$)	$\langle e_{12}, \delta n_0 - e_1, e_{12} + \delta n_2 \rangle$ $= \delta \langle e_{12}, \delta n_0 - e_1, n_2 \rangle$ $= \delta \langle e_2, -e_1, n_2 \rangle$
(0 ,2, 5, 4)	vol10($0, e_2, e_2 + \delta n_2, e_1 + \delta n_1$)	$\delta \langle e_2, n_2, e_1 \rangle$
(0 ,1, 2, 5)	vol11($0, e_1, e_2, e_2 + \delta n_2$)	$\delta \langle e_1, e_2, n_2 \rangle$
(2, 3, 4, 5)	vol12($e_2, \delta n_0, e_1 + \delta n_1, e_2 + \delta n_2$)	$\langle \delta n_0 - e_2, \delta n_1 - e_{12}, \delta n_2 \rangle$ $= \delta \langle -e_2, e_1, n_2 \rangle$

For all the 12 tetrahedra the linear term is multiplied by a determinant containing the edges of the prism. We check directly that all these determinants are positive due to the assumption that $CN_i > 0$.

A.1.3 THEOREM 3.5

Consider a ball around a point p of the middle face M of the prism. If the radius $r(p) > 0$ is sufficiently small, it contains, at most, a single vertex of M , and parts of incident faces. As M is compact, there is a minimal value of r on M , and it is positive. The line along the normal, passing through v , intersects incident faces at v so it cannot intersect them at any other point. Thus, the top and bottom prism vertices can be obtained by displacing V by r in either direction along the normal. Consider the r -neighborhood of the M , i.e. the union of all balls of radius r centered at points of M . This is a convex set, containing only M and parts of vertex-adjacent faces. All top and bottom prism vertices are in this set. Thus tetrahedra connecting these vertices are also in the set, i.e., complete prisms. The fact that tetrahedra are nondegenerate is established in Proposition 3.4.

A.1.4 THEOREM 3.6

To show that \mathcal{T} is a section of \mathcal{S} we need to show that (1) $M_\Delta = \mathcal{T} \cap \Delta$ is a simply connected patch for any prism of \mathcal{S} and (2) for each prism Δ and every point $p \in M_\Delta$ the dot product between the face normal $n(p)$ and the pillars $\mathcal{V}_i, i = 1, 2, 3$ is strictly positive.

(1) is trivial since for every prism Δ of \mathcal{S} , Δ contains exactly one triangle of the topological bevel refinement face of \mathcal{T} .

To prove (2), consider the dot products of the normals of each of the triangles in the beveled region sharing at least a vertex with M_Δ , and the vectors along the pillars of M_Δ . We use colors to refer to the triangles of beveled prisms as shown in Figure A.1 left.

The prism corresponding to the pink triangle covers only the interior of the original triangle. It

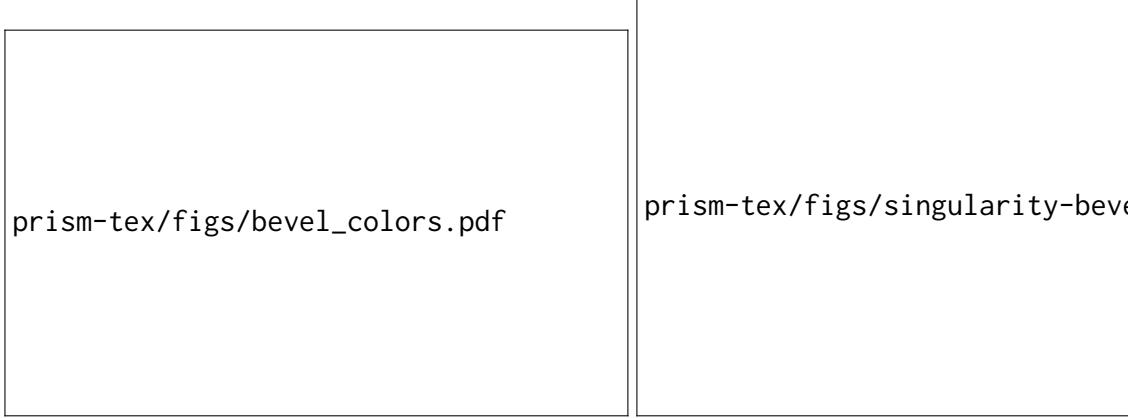


Figure A.1: Bevel patterns used in the proofs in Appendix A.1.4 and Appendix A.2

satisfies the dot product condition since its pillars are copied from the solution of Problem (3.3) at each vertex. A similar argument can be made for the green prisms: each prism gets its pillars from the edges incident at two of the vertices, which are compatible with the normal of the adjacent triangle (e.g., red and blue pillars have positive dot product with the normal of T_1). Finally, the prisms corresponding to orange triangles cover the same one ring that was used to compute the original pillars (e.g., red pillar is compatible with T_1 and T_2). Since the same pillar is used for each of the 3 vertices of the original prism, the dot product is unchanged.

A.1.5 THEOREM 3.7

We use A° to denote the interior of a set A . We first assert that the topological disk patch $\mathcal{T} \cap C$ coincides with $\mathcal{T} \cap C'$. To make the notation more concise, we define $\bar{\mathcal{T}}_C = \mathcal{T} \cap C$ and $\bar{\mathcal{T}}_{C'} = \mathcal{T} \cap C'$.

We show that $\bar{\mathcal{T}}_C \subset C'$ by contradiction: assume there is $p \in \bar{\mathcal{T}}_C$ but $p \notin C'$. Choose a point $q \in \partial\bar{\mathcal{T}}_C$, that belongs to single prisms in C and C' and denote the enclosing prism in C as Δ_q and the corresponding one in C' as Δ'_q . It follows from the positivity of the volumes of Δ_q and Δ'_q (Assumption 1), that there is $\varepsilon > 0$ such that the half-ball (for sufficiently small ε) $\text{Ball}(q, \varepsilon) \cap \Delta_q$ coincides with $\text{Ball}(q, \varepsilon) \cap \Delta'_q$. Furthermore, the validity of the initial shell guarantees that there exists an interior point $\bar{q} \in \text{Ball}(q, \varepsilon) \cap \Delta_q^\circ \cap \mathcal{T}$, and $\bar{q} \in \Delta_q'^\circ \cap \mathcal{T} \subset C'^\circ \cap \mathcal{T}$.

Since neither p or \bar{q} lies on the boundary side triangles of $\partial C'$ (since $p, \bar{q} \notin \partial\bar{\mathcal{T}}_C$), therefore there exists an interior path P (i.e., a path both in $\bar{\mathcal{T}}_C$ and in $\bar{\mathcal{T}}_{C'}$) that connects p and \bar{q} . By continuity the path P it will cross $\partial C'$ at some point $p_i \in \bar{\mathcal{T}}_C$. Since the $\partial\bar{\mathcal{T}}_{C'} = \partial\bar{\mathcal{T}}_C$ is a fixed boundary loop (by the definition of \mathbb{O}), and does not contain interior points, p_i can only cross on the top surface or bottom surface of C' which contradicts Assumption 2. The roles of C and C' can be inverted to complete the proof.

We then map M_C to a regular planar n -gon D , with vertices on ∂M_C mapped cyclically to the vertices of ∂D . It follows from [Floater 2003, Corollary 6.2] that such a convex combination mapping $\varphi: M_C \rightarrow D$ is bijective. Then, similarly to the construction in Theorem 3.2, we define $\phi: C \rightarrow D \times [-1, 1]$ through affine mapping induced by the tetrahedralization of the prisms Δ_i . Similarly, we define $\psi: C' \rightarrow D \times [-1, 1]$. The mapping ψ is also bijective because of Assumption 1 and the definition of \mathbb{O} it follows that the codomain of ψ is the same as ϕ . Note that in both

cases, the map transforms the vector field $\mathcal{V}(p)$ to the constant vector field $e_z = (0, 0, 2)$, and the projection is \mathcal{P}_z , as defined in Appendix A.1.1.

The dot product condition (Assumption 3) ensures that \mathcal{P}_z defines a bijection between D and $\psi(\mathcal{T}_C)$; and further, the image satisfies $\mathcal{P}_z(\psi(\Delta'_i \cap \mathcal{T})) = \psi(M'_i)$ where M'_i is the middle surface (a single triangle) of Δ'_i . Therefore, since ψ and \mathcal{P}_z are both continuous and bijective, they are homeomorphisms which preserve topology, thus $\Delta'_i \cap \mathcal{T}$ is a simply connected topological disk.

A.2 EXTENSION TO MESHES WITH SINGULARITIES

Most of the proofs and definition in Section 3.3 easily extends to meshes with singularity or pinched shells. For instance, both Theorem 3.2 and Definition 3.3 apply to pinched shells by just considering a prism made of 4 (2 for the top and 2 for the bottom slab) instead of 6. Propositions 3.4 and 3.5 both rely on per-vertex properties; by just excluding singular vertices (and neighboring prism) from the statements the proofs (and our algorithm) still holds.

Theorem 3.7 requires some minor additional considerations. As a consequence of beveling, no two singularities are adjacent. Therefore, we can always find a point $q \in \partial\mathcal{T}_C$ that belongs to a single prism (Appendix A.1.5) since every prism will have a positive volume because no singularities are adjacent.

Finally Proposition 3.6 requires a new proof since the beveling pattern used for singularities is different.

Proof. In Figure A.1 right, we highlight in red and orange the triangles not covered by the discussion in Appendix A.1.4. The prism corresponding to an orange middle triangle intersects the edge-adjacent triangle. And since the new pillars are computed as the average of the normals of the two adjacent triangles with dihedral angle strictly less than 360° (Section 3.3.5 inset), each dot product is positive. The prism for a pink middle triangle intersects only the interior of the original triangle, and the dot product between the pillars and the triangle normal are positive by construction. \square

A.3 REDUCTION TO A STANDARD QP

With slack variable t , Problem 3.2 can be rewritten as

$$\max_t t; d_v n_f \geq t, \|d_v\| = 1, t \geq \epsilon.$$

In the admissible region, substituting $t = 1/s$,

$$\min_s s; s d_v n_f \geq 1, \|d_v\| = 1, 0 < s \leq 1/\epsilon.$$

Substituting $x = s d_v$ (thus $\|x\| = s \|d\| = s$), we obtain the QP problem 3.3. Note that the lower bound on s is implied from $s d_v n_f \geq 1$, and the upper bound $\|x\| < 1/\epsilon$ is checked a posteriori.

A.4 BIJECTIVITY OF THE NONLINEAR PRISMATIC TRANSFORMATION

In this section, we show that the isoparametric transformation of prismatic finite element [Ciarlet 1991] is also bijective if I1 holds.

We follow the notation from Appendix A.1.2. The map is defined as (for the top slab)

$$f(u, v, \eta) = ue_1 + ve_2 + \eta n_0 + u\eta n_{01} + v\eta n_{02},$$

$$\det J_f = \langle (1 - u - v)n_0 + un_1 + vn_2, e_1 + \eta n_{01}, e_2 + \eta n_{02} \rangle,$$

where η is the variable corresponding to the thickness direction, and $n_{ij} = n_j - n_i$. Observe that $\det J_f(u, v, \eta_0)$ is linear in u, v for fixed η_0 , the extrema will only be achieved at the corner points [Knabner and Summ 2001]. Therefore, it is sufficient to check the signs at three edges $(0, 0, \eta)$, $(1, 0, \eta)$, $(0, 1, \eta)$ for $\eta \in [0, 1]$.

$$\begin{aligned} & \det J_f(0, 0, \eta) \\ &= \langle n_0, (1 - \eta)e_1 + \eta(e_1 + n_1 - n_0), (1 - \eta)e_2 + \eta(e_2 + n_2 - n_0) \rangle \\ &= (1 - \eta)^2 \langle n_0, e_1, e_2 \rangle + (1 - \eta)\eta \langle n_0, e_1, e_2 + n_2 \rangle \\ &\quad + (1 - \eta)\eta \langle n_0, e_1 + n_1, e_2 \rangle + \eta^2 \langle n_0, e_1 + n_1, e_2 + n_2 \rangle \\ &= (1 - \eta)^2 \text{vol}_4 + (1 - \eta)\eta(\text{vol}_3 + \text{vol}_1) + \eta^2 \text{vol}_2 > 0. \end{aligned}$$

By symmetry, it is easy to verify that the following are positive,

$$\det J_f(1, 0, \eta) = (1 - \eta)^2 \text{vol}_6 + (1 - \eta)\eta(\text{vol}_7 + \text{vol}_5) + \eta^2 \text{vol}_8,$$

$$\det J_f(0, 1, \eta) = (1 - \eta)^2 \text{vol}_{11} + (1 - \eta)\eta(\text{vol}_9 + \text{vol}_{10}) + \eta^2 \text{vol}_{12}.$$

A.5 DOUBLE SLAB

In Section 3.3.1, we explain that we want to ensure that our shell validity is independent from the enumeration of the faces; for this reason, we require that the conditions I1 and I2 are satisfied for all possible decompositions. Without using the double slab, different decompositions of the prism will result in different intersections with the middle surface. In other words, the projection operator will project the input mesh to a different set of triangles on the middle surface. The double slab makes the middle surface independent of the decomposition by construction. We note that Theorem 3.6 is valid only for double-slab prisms since the statement requires that one prism contains only one triangle.

A.6 UNPUBLISHED MATERIAL

(Disclaimer: this section is unpublished material. And serves as post publish notes for future references.)

A.6.1 POINTLESS VARIANT

We extend the framework, by allowing the relaxation of definition of a section. Originally, any intersection between the triangle and the (different decomposition of the) prism is counted. Here, we explicitly forgive when the intersection is only one point on the vertex of mid surface and the triangle. This predicate can be easily implemented, making use of the bitwise equal coordinates. Therefore, the initial tracking section only follows edge-adjacency relationship, instead of the vertex adjacent one in the main text. With such adjusted condition for section, the optimization can be trivially accommodated. The only thing that is different (simpler) is the part for bevel, since we have less triangles to consider. The existing one is still valid, however, we can get away with simpler red-green intersection

A.6.2 DYNAMIC INTERSECTION CHECK

To make easier the guarantee of a self-intersection-free shell, substitute the AABB tree collision check with a dynamic hashgrid. We initialize two hashgrids at the beginning, one for top surface and another one for bottom surface. This is followed by additional shrinking to make sure they are clear. Each local operation is responsible to trial and keep the validity of the surfaces. Furthermore, the top/bottom surface should not interfere with the reference input surface, this can be accelerated through the maintenance of the tracking list, and is trial checked before a local operation is performed.

A.6.3 FEATURE SNAPPING SHELL

We can separate the definition of reference surface and input surface to make a feature snapping shell.

A.6.4 POSITIVE DETERMINANT FOR NATURAL PRISM MAP

This section follows and extends Appendix A.4. In fact, we observe that, for a quadratic Bezier curve piece $at^2 + bt(1-t) + c(1-t)^2$, $t \in [0, 1]$, the sufficient and necessary condition is $a > 0$, $c > 0$, $b > -2\sqrt{ac}$, and the last relation is equivalent to (but more clearly stated as) $b > 0 \mid b^2 - 4ac < 0$. The current text (all 12 positive tetra) impose an unnatural constraint on edge splits: in the case of an aggressively positive prism (all decompositions are all-positive), sub-prism may not be positive in the same sense. While this will not break anything, forbidding edge split will lead to unexpected lower quality in the optimization procedure.

B | CURVE MESHING RELATED

B.1 THE GEOMETRIC MAP IS BIJECTIVE

Consider a connected 3-dimensional compact manifold (curved) tetrahedral mesh $\mathcal{M} = \{\sigma_i\}, i = 1, \dots, n$ with $\sigma_i = g_i(\hat{\tau})$, $\hat{\tau}$ a regular unit tetrahedron, $\det(J_{g_i}) > 0$ at all points including the boundary, and σ_i and σ_j agree on a shared face. Let \mathcal{M}_D be the domain obtained by copies of $\hat{\tau}$ and identifying $\hat{\tau}_i$ along common faces. We then define the map

$$\sigma: M_D \rightarrow \mathbb{R}^3,$$

by setting $\sigma|_{\hat{\tau}_i} = \sigma_i$.

Proposition B.1. *Suppose $\sigma|_{\partial\mathcal{M}_D}$ is injective. Then σ is injective on the whole domain \mathcal{M}_D .*

Proof. Our argument closely follow from [Aigerman and Lipman 2013, Appendix B] and [Lipman 2014, Theorem 1].

We consider point $y \in \mathbb{R}^3$ in general position, but not in the faces, edges, vertices, or any plane spanned by a linear face of \mathcal{M} . For each tetrahedron τ , we construct the map $\hat{\Psi}$ as a composition of $\sigma|_{\partial\hat{\tau}}$ (restricted to the triangular faces of the regular tetrahedron) and the projection map χ to the unit sphere centered around y .

We parametrize the image of $\sigma|_{\partial\hat{\tau}}$ as $x(u, v)$ on the projective plane (note that since $\det J_{g_i} > 0$, the image is a non-degenerate surface homeomorphic to a sphere). Since $x(u, v)$, and its normal field $n(u, v) = \frac{\partial}{\partial u}x \times \frac{\partial}{\partial v}x$ are polynomial functions, Consider the parametric curve $n(u, v) \cdot (x(u, v) - y) = 0$, the algebraic curve partitions the surface into finite number of patches, and on each patch the orientation of $\hat{\Psi}$ is constant. We can further triangulate the partitions to create a (abstract simplicial) complex.

Similar to [Aigerman and Lipman 2013, Appendix B], we count the number of pre-images of y , which equals to the degree in general positions,

$$\deg(\sigma)(y) = \sum_{i=1}^n \deg(\hat{\Psi}|_{\tau_i}) = \deg(\hat{\Psi}|_{\partial\mathcal{M}}).$$

Since $\sigma|_{\partial\mathcal{M}_D}$ is injective, and furthermore

$$\deg(\hat{\Psi}|_{\partial\mathcal{M}}) = \deg\chi|_{\partial\mathcal{M}} = \begin{cases} 1, & y \in \mathcal{M} \\ 0, & y \notin \mathcal{M}. \end{cases}$$

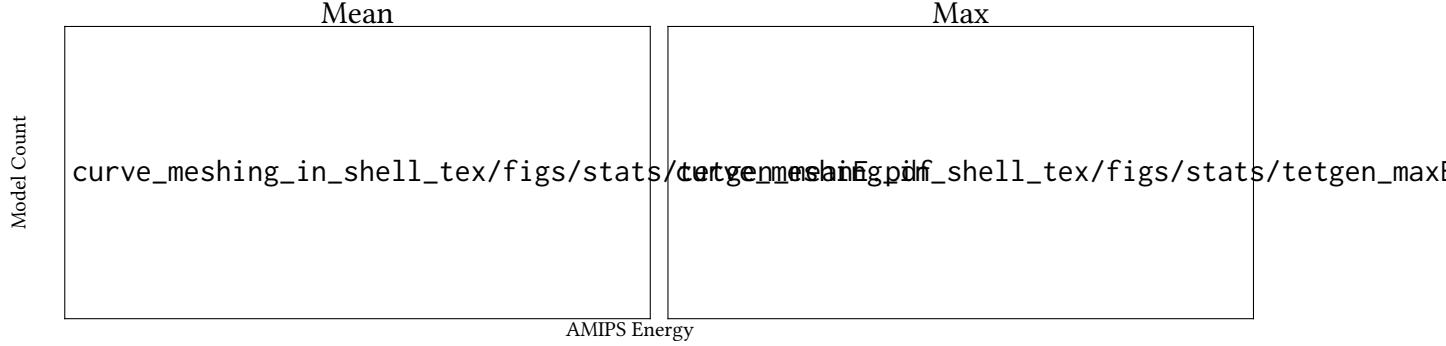


Figure B.1: Histogram of mean and maximum conformal AMIPS energy [Rabinovich et al. 2017] of the output of our method and TetGen.

Thus we have shown the map is injective for the general positions for y , and it remains to be shown that the map is an open map, which follows from the same argument of [Lipman 2014, Lemma 2].

□

B.2 LOCAL OPERATIONS

[Jiang et al. 2020, Fig. 11] introduce a set of *valid* local operations to modify the shell, including edge split, edge collapse, edge flip and vertex smoothing. The operations are an analogue of the triangle mesh edit operations [Dunyach et al. 2013], by simultaneously editing the shared connectivity of bottom, middle and top surface of the shell. [Jiang et al. 2020, Theorem 3.7] outlines invariant conditions, which maintains the shell projection to be bijective.

Our algorithm adopts and extends the local operations therein to the high order setting. In addition to the existing conditions, we also validate the curved volumetric mesh in the shell. The algorithm maintains the global intersection free bottom (top) surface with a dynamic hash grid [Teschner et al. 2003]. Then for each prism, we check the positivity (defined by the determinant of Jacobian of the geometric map) of the prismatic element (each decomposed into three tetrahedra).

In the presence of feature annotation and feature straightening (Section 4.5) more care is taken to maintain the valid correspondence between the grouped feature chains and the curved edges: edge flip is disabled on the edges annotated as features; collapse is only allowed when it does not degenerate the chain, and the two endpoints of the edge is on the same chain. Since we require a map from the original input edges, we insert additional degree of freedoms in the input mesh. When performing edge split, the insertion of the new vertex is queried among the pre-image of the current edge, as an existing vertex of the input. For vertex smoothing (more specifically pan), the target location is limited to the set of input vertices.

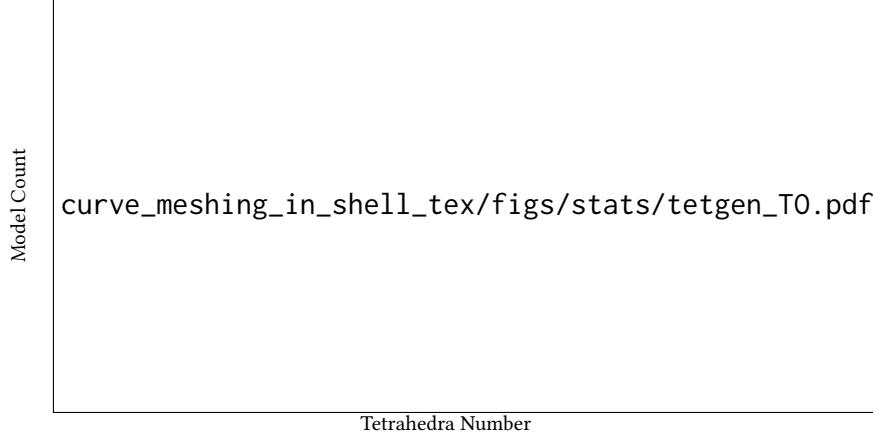


Figure B.2: Histogram of output tetrahedra number for TetGen and our method.

B.3 BOUNDARY PRESERVING TETGEN COMPARISON

We compared our conforming tetrahedral meshing algorithm (Section 4.4.3) with TetGen on the linear shells (triangle meshes) of 3522 models from the Thingi10k dataset, giving each model sufficient computing resources (2 hours maximum running time and 32GB memory usage). Inheriting the robustness from TetWild, our method successfully processed all the inputs while preserving the triangulation, while TetGen fails on 224 models (215 models are not conforming, and 9 models have no output).

In Figure B.1, we show the average and maximum element quality of the output of our method and TetGen. Our method has a better average and maximum output quality than TetGen. Note that in the quality plot, the “tail” of TetGen’s distribution is longer than ours. The maximum average energy of TetGen’s output and ours are 3×10^8 and 3×10^5 respectively. The largest maximum energy of TetGen’s output and ours are 3×10^{12} and 7×10^6 respectively. Our method generates denser output (Figure B.2), but our focus is on robustness instead of efficiency in this step.

B.4 TETWILD OVERVIEW

BIBLIOGRAPHY

- A. Freitag, L. and Ollivier-Gooch, C. (1998). Tetrahedral mesh improvement using swapping and smoothing. *International Journal for Numerical Methods in Engineering*, 40.
- Abgrall, R., Dobrzynski, C., and Froehly, A. (2012). A method for computing curved 2D and 3D meshes via the linear elasticity analogy: preliminary results. Research Report RR-8061, INRIA.
- Abgrall, R., Dobrzynski, C., and Froehly, A. (2014). A method for computing curved meshes via the linear elasticity analogy, application to fluid dynamics problems. *International Journal for Numerical Methods in Fluids*, 76(4):246–266.
- Aigerman, N., Kovalsky, S. Z., and Lipman, Y. (2017). Spherical orbifold tutte embeddings. *ACM Trans. Graph.*, 36(4).
- Aigerman, N. and Lipman, Y. (2013). Injective and bounded distortion mappings in 3d. *ACM Transactions on Graphics (TOG)*, 32(4):1–14.
- Aigerman, N. and Lipman, Y. (2015). Orbifold tutte embeddings. *ACM Trans. Graph.*, 34(6).
- Aigerman, N. and Lipman, Y. (2016). Hyperbolic orbifold tutte embeddings. *ACM Trans. Graph.*, 35(6).
- Aigerman, N., Poranne, R., and Lipman, Y. (2014). Lifted bijections for low distortion surface mappings. *ACM Trans. Graph.*, 33(4).
- Aigerman, N., Poranne, R., and Lipman, Y. (2015). Seamless surface mappings. *ACM Trans. Graph.*, 34(4).
- Ainsley, S., Vouga, E., Grinspun, E., and Tamstorf, R. (2012). Speculative parallel asynchronous contact mechanics. *ACM Trans. Graph.*, 31(6):151:1–151:8.
- Alauzet, F. and Marcum, D. (2014). A closed advancing-layer method with changing topology mesh movement for viscous mesh generation. In *Proceedings of the 22nd International Meshing Roundtable*, pages 241–261. Springer International Publishing, Cham.
- Alexa, M. (2019). Harmonic triangulations. *ACM Transactions on Graphics (Proceedings of SIGGRAPH)*, 38(4):54.

- Alexa, M. (2020). Conforming weighted delaunay triangulations. *ACM Transactions on Graphics (TOG)*, (6):1–16.
- Alliez, P., Cohen-Steiner, D., Yvinec, M., and Desbrun, M. (2005a). Variational tetrahedral meshing. In *ACM Transactions on Graphics (TOG)*, number 3. ACM.
- Alliez, P., Cohen-Steiner, D., Yvinec, M., and Desbrun, M. (2005b). Variational tetrahedral meshing. *ACM Transactions on Graphics*, (3):617.
- Alliez, P., De Verdire, E. C., Devillers, O., and Isenburg, M. (2003). Isotropic surface remeshing. In *2003 Shape Modeling International.*, pages 49–58. IEEE.
- Aluru, S. and Sevilgen, F. E. (1997). Parallel domain decomposition and load balancing using space-filling curves. In *Proceedings fourth international conference on high-performance computing*, pages 230–235. IEEE.
- Ando, R., Thürey, N., and Wojtan, C. (2013). Highly adaptive liquid simulations on tetrahedral meshes. *ACM Trans. Graph. (Proc. SIGGRAPH 2013)*.
- Aubry, R., Dey, S., Mestreau, E., and Karamete, B. (2017). Boundary layer mesh generation on arbitrary geometries. *International Journal for Numerical Methods in Engineering*, 112(2):157–173.
- Aubry, R. and Löhner, R. (2008). On the ‘most normal’normal. *Communications in Numerical Methods in Engineering*, 24(12):1641–1652.
- Aubry, R., Mestreau, E., Dey, S., Karamete, B., and Gayman, D. (2015). On the ‘most normal’normal—part 2. *Finite Elements in Analysis and Design*, 97:54–63.
- Babuska, I. and Guo, B. Q. (1988). The h-p version of the finite element method for domains with curved boundaries. *SIAM Journal on Numerical Analysis*, 25(4):837–861.
- Babuška, I. and Guo, B. (1992). The h, p and h-p version of the finite element method; basis theory and applications. *Advances in Engineering Software*, 15(3):159 – 174.
- Bajaj, C. L., Xu, G., Holt, R. J., and Netravali, A. N. (2002). Hierarchical multiresolution reconstruction of shell surfaces. *Computer Aided Geometric Design*, 19(2):89–112.
- Bank, R. E., Sherman, A. H., and Weiser, A. (1983). Some refinement algorithms and data structures for regular local mesh refinement. *Scientific Computing, Applications of Mathematics and Computing to the Physical Sciences*, 1:3–17.
- Bargteil, A. W. and Cohen, E. (2014). Animation of deformable bodies with quadratic bézier finite elements. *ACM Transactions on Graphics (TOG)*, 33(3):27.
- Bargteil, A. W., Wojtan, C., Hodgins, J. K., and Turk, G. (2007). A finite element method for animating large viscoplastic flow. In *ACM SIGGRAPH 2007 Papers*, SIGGRAPH ’07, page 16–es, New York, NY, USA. Association for Computing Machinery.

- Barnhill, R. E., Opitz, K., and Pottmann, H. (1992). Fat surfaces: a trivariate approach to triangle-based interpolation on surfaces. *Computer Aided Geometric Design*, 9(5):365–378.
- Bassi, F. and Rebay, S. (1997). High-order accurate discontinuous finite element solution of the 2d euler equations. *Journal of Computational Physics*, 138(2):251 – 285.
- Batista, V. H., Millman, D. L., Pion, S., and Singler, J. (2010). Parallel geometric algorithms for multi-core computers. *Computational Geometry*, 43(8):663–677.
- Baumgart, B. G. (1972). Winged edge polyhedron representation. Technical report, Stanford, CA, USA.
- Bender, J. and Deul, C. (2013). Adaptive cloth simulation using corotational finite elements. *Computers & Graphics*, 37(7):820–829.
- Bern, M., Eppstein, D., and Gilbert, J. (1994). Provably good mesh generation. *Journal of Computer and System Sciences*, 48(3):384 – 409.
- Bishop, C. J. (2016). Nonobtuse triangulations of pslgs. *Discrete & Computational Geometry*, 56(1):43–92.
- Blandford, D. K., Blelloch, G. E., and Kadow, C. (2006). Engineering a compact parallel delaunay algorithm in 3d. In *Proceedings of the twenty-second Annual Symposium on Computational Geometry*, pages 292–300.
- Blelloch, G. E., Miller, G. L., Hardwick, J. C., and Talmor, D. (1999). Design and implementation of a practical parallel delaunay algorithm. *Algorithmica*, 24(3):243–269.
- Boissonnat, J.-D., Devillers, O., Pion, S., Teillaud, M., and Yvinec, M. (2002). Triangulations in cgal. *Computational Geometry*, 22:5–19.
- Boissonnat, J.-D. and Oudot, S. (2005a). Provably good sampling and meshing of surfaces. *Graphical Models*, 67(5):405–451.
- Boissonnat, J.-D. and Oudot, S. (2005b). Provably good sampling and meshing of surfaces. *Graphical Models*, 67(5):405–451.
- Bommes, D., Campen, M., Ebke, H.-C., Alliez, P., and Kobbelt, L. (2013a). Integer-grid maps for reliable quad meshing. *ACM Trans. Graph.*, 32(4):98:1–98:12.
- Bommes, D., Lévy, B., Pietroni, N., Puppo, E., Silva, C., Tarini, M., and Zorin, D. (2013b). Quad-mesh generation and processing: A survey. In *Computer Graphics Forum*, volume 32, pages 51–76. Wiley Online Library.
- Bommes, D., Lévy, B., Pietroni, N., Puppo, E., a, C. S., Tarini, M., and Zorin, D. (2012). State of the art in quad meshing. In *Eurographics STARS*.

- Bommes, D., Zimmer, H., and Kobbelt, L. (2009). Mixed-integer quadrangulation. *ACM Trans. Graph.*, 28(3):77:1–77:10.
- Borrell, R., Cañas, J. C., Mira, D., Taha, A., Koric, S., Vázquez, M., and Houzeaux, G. (2018). Parallel mesh partitioning based on space filling curves. *Computers & Fluids*, 173:264–272.
- Botsch, M. and Kobbelt, L. (2003). Multiresolution surface representation based on displacement volumes. In *Computer Graphics Forum*, volume 22, pages 483–491. Wiley Online Library.
- Botsch, M. and Kobbelt, L. (2004). A remeshing approach to multiresolution modeling. In *Proceedings of the 2004 Eurographics/ACM SIGGRAPH symposium on Geometry processing*, pages 185–192.
- Botsch, M., Pauly, M., Gross, M. H., and Kobbelt, L. (2006). Primo: coupled prisms for intuitive surface modeling. In *Symposium on Geometry Processing*, number CONF, pages 11–20.
- Botsch, M., Steinberg, S., Bischoff, S., and Kobbelt, L. (2002). Openmesh-a generic and efficient polygon mesh data structure.
- Boubekeur, T. and Alexa, M. (2008). Phong tessellation. In *ACM Transactions on Graphics (TOG)*, volume 27, page 141. ACM.
- Bowers, J., Wang, R., Wei, L.-Y., and Maletz, D. (2010). Parallel poisson disk sampling with spectrum analysis on surfaces. *ACM Transactions on Graphics (TOG)*, 29(6):1–10.
- Bridson, R. and Doran, C. (2014). Quartet: A tetrahedral mesh generator that does isosurface stuffing with an acute tetrahedral tile. <https://github.com/crawforddoran/quartet>.
- Brisson, E. (1989). Representing geometric structures in d dimensions: Topology and order. In *Proceedings of the Fifth Annual Symposium on Computational Geometry*, SCG ’89, page 218–227, New York, NY, USA. Association for Computing Machinery.
- Brochu, T., Edwards, E., and Bridson, R. (2012). Efficient geometrically exact continuous collision detection. *ACM Trans. Graph.*, 31(4).
- Bronson, J. R., Levine, J. A., and Whitaker, R. T. (2013). Lattice cleaving: Conforming tetrahedral meshes of multimaterial domains with bounded quality. In *Proceedings of the 21st International Meshing Roundtable*, pages 191–209. Springer Berlin Heidelberg, Berlin, Heidelberg.
- Busaryev, O., Dey, T. K., and Levine, J. A. (2009). Repairing and meshing imperfect shapes with delaunay refinement. In *2009 SIAM/ACM Joint Conference on Geometric and Physical Modeling*, SPM ’09, pages 25–33. ACM.
- Busaryev, O., Dey, T. K., and Wang, H. (2013). Adaptive fracture simulation of multi-layered thin plates. *ACM Trans. Graph.*, 32(4).
- Calderon, S. and Boubekeur, T. (2017). Bounding proxies for shape approximation. *ACM Trans. Graph.*, 36(4).

- Campen, M., Capouillez, R., Shen, H., Zhu, L., Panozzo, D., and Zorin, D. (2021). Efficient and robust discrete conformal equivalence with boundary. *ACM Trans. Graph.*, 40(6).
- Campen, M., Shen, H., Zhou, J., and Zorin, D. (2019). Seamless parametrization with arbitrary cones for arbitrary genus. *ACM Trans. Graph.*, 39(1).
- Campen, M., Silva, C. T., and Zorin, D. (2016). Bijective maps from simplicial foliations. *ACM Trans. Graph.*, 35(4):74:1–74:15.
- Campen, M. and Zorin, D. (2017). Similarity maps and field-guided t-splines: a perfect couple. *ACM Trans. Graph.*, 36(4).
- Canann, S. A., Muthukrishnan, S. N., and Phillips, R. K. (1996). Topological refinement procedures for triangular finite element meshes. *Engineering with Computers*, 12(3):243–255.
- Canann, S. A., Stephenson, M. B., and Blacker, T. (1993). Optismoothing: An optimization-driven approach to mesh smoothing. *Finite Elements in Analysis and Design*, 13(2):185 – 190.
- Cardoze, D., Cunha, A., Miller, G. L., Phillips, T., and Walkington, N. (2004). A b  zier-based approach to unstructured moving meshes. In *Proceedings of the Twentieth Annual Symposium on Computational Geometry*, SCG ’04, pages 310–319, New York, NY, USA. ACM.
- Chazal, F. and Cohen-Steiner, D. (2005). A condition for isotopic approximation. *Graphical Models*, 67(5):390–404.
- Chazal, F., Lieutier, A., Rossignac, J., and Whited, B. (2010). Ball-map: Homeomorphism between compatible surfaces. *International Journal of Computational Geometry & Applications*, 20(03):285–306.
- Chen, L. and Xu, J.-c. (2004). Optimal delaunay triangulations. *Journal of Computational Mathematics*, 22(2):299–308.
- Chen, M.-B. (2010). The merge phase of parallel divide-and-conquer scheme for 3d delaunay triangulation. In *International Symposium on Parallel and Distributed Processing with Applications*, pages 224–230. IEEE.
- Chen, Y., Tong, X., Wang, J., Wang, J., Lin, S., Guo, B., Shum, H.-Y., and Shum, H.-Y. (2004). Shell texture functions. In *ACM Transactions on Graphics (TOG)*, volume 23, pages 343–353. ACM.
- Cheng, S.-W., Dey, T. K., Edelsbrunner, H., Facello, M. A., and Teng, S.-H. (2000). Silver exudation. *Journal of the ACM (JACM)*, 47(5):883–904.
- Cheng, S.-W., Dey, T. K., and Levine, J. A. (2008a). A practical delaunay meshing algorithm for a large class of domains. In *Proceedings of the 16th International Meshing Roundtable*, pages 477–494. Springer.

- Cheng, S.-W., Dey, T. K., and Levine, J. A. (2008b). A practical delaunay meshing algorithm for a large class of domains. In *Proceedings of the 16th International Meshing Roundtable*, pages 477–494, Berlin, Heidelberg. Springer Berlin Heidelberg.
- Cheng, S.-W., Dey, T. K., and Shewchuk, J. (2012). *Delaunay Mesh Generation*. Chapman and Hall/CRC, Boca Raton, Florida.
- Cheng, X.-X., Fu, X.-M., Zhang, C., and Chai, S. (2019). Practical error-bounded remeshing by adaptive refinement. *Computers & Graphics*, 82:163 – 173.
- Chentanez, N., Feldman, B. E., Labelle, F., O’Brien, J. F., and Shewchuk, J. R. (2007). Liquid simulation on lattice-based tetrahedral meshes. In *Proceedings of the 2007 ACM SIGGRAPH/Eurographics Symposium on Computer Animation*, SCA ’07, page 219–228, Goslar, DEU. Eurographics Association.
- Cheshmi, K., Kaufman, D. M., Kamil, S., and Dehnavi, M. M. (2020). Nasoq: Numerically accurate sparsity-oriented qp solver. *ACM Transactions on Graphics*, 39(4).
- Chew, L. P. (1989a). Constrained delaunay triangulations. *Algorithmica*, 4(1-4):97–108.
- Chew, L. P. (1989b). Constrained delaunay triangulations. *Algorithmica*, 4(1):97–108.
- Chew, L. P. (1993). Guaranteed-quality mesh generation for curved surfaces. In *Proceedings of the ninth annual symposium on Computational geometry*, pages 274–280. ACM.
- Chew, L. P., Chrisochoides, N., and Sukup, F. (1997). Parallel constrained delaunay meshing. *ASME APPLIED MECHANICS DIVISION-PUBLICATIONS-AMD*, 220:89–96.
- Chrisochoides, N. (2006). Parallel mesh generation. In *Numerical solution of partial differential equations on parallel computers*, pages 237–264. Springer.
- Chrisochoides, N. and Nave, D. (2003). Parallel delaunay mesh generation kernel. *International Journal for Numerical Methods in Engineering*, 58(2):161–176.
- Ciarlet, P. G. (1991). Basic error estimates for elliptic problems. *Finite Element Methods (Part 1)*.
- Cignoni, P., Ganovelli, F., et al. (2021). Vcg library.
- Cignoni, P., Montani, C., Perego, R., and Scopigno, R. (1993). Parallel 3d delaunay triangulation. *Computer Graphics Forum*, 12(3):129–142.
- Clausen, P., Wicke, M., Shewchuk, J. R., and O’Brien, J. F. (2013). Simulating liquids and solid-liquid interactions with lagrangian meshes. *ACM Trans. Graph.*, 32(2).
- Cohen, J., Manocha, D., and Olano, M. (1997). Simplifying polygonal models using successive mappings. In *Proceedings. Visualization’97 (Cat. No. 97CB36155)*, pages 395–402. IEEE.

- Cohen, J., Varshney, A., Manocha, D., Turk, G., Weber, H., Agarwal, P., Brooks, F., and Wright, W. (1996). Simplification envelopes. In *Siggraph*, volume 96, pages 119–128.
- Cohen-Steiner, D., De Verdiere, E. C., and Yvinec, M. (2002a). Conforming delaunay triangulations in 3d. In *Proceedings of the eighteenth annual symposium on Computational geometry*, pages 199–208. ACM.
- Cohen-Steiner, D., de Verdière, E. C., and Yvinec, M. (2002b). Conforming delaunay triangulations in 3D. In *Proceedings of the eighteenth annual symposium on Computational geometry - SCG '02*, number 2, pages 217 – 233. ACM Press.
- Collins, G. and Hilton, A. (2002). Mesh decimation for displacement mapping. In *Eurographics (Short Papers)*.
- Community, B. O. (2018). *Blender - a 3D modelling and rendering package*. Blender Foundation, Stichting Blender Foundation, Amsterdam.
- Conway, J. H., Burgiel, H., and Goodman-Strauss, C. (2016). *The symmetries of things*. CRC Press.
- Coreform (2020). Cubit. <https://coreform.com/products/coreform-cubit/>.
- Coxeter, H. S. M. (1973). *Regular polytopes*. Courier Corporation.
- Crane, K. (2013). *Conformal geometry processing*. California Institute of Technology.
- Crane, K., Weischedel, C., and Wardetzky, M. (2013). Geodesics in heat: A new approach to computing distance based on heat flow. *ACM Transactions on Graphics (TOG)*, 32(5).
- Cuilliere, J.-C., Francois, V., and Drouet, J.-M. (2013). Automatic 3D mesh generation of multiple domains for topology optimization methods. In *Proceedings of the 21st International Meshing Roundtable*, pages 243–259. Springer Berlin Heidelberg, Berlin, Heidelberg.
- Degener, P., Meseth, J., and Klein, R. (2003). An adaptable surface parameterization method. In *Proceedings of the 12th International Meshing Roundtable*, pages 201–213.
- Dey, S., O’bara, R. M., and Shephard, M. S. (1999a). Curvilinear mesh generation in 3d. In *IMR*, pages 407–417. IMR.
- Dey, T. K., Edelsbrunner, H., Guha, S., and Nekhayev, D. V. (1999b). Topology preserving edge contraction. *Publ. Inst. Math.(Beograd)(NS)*, 66(80):23–45.
- Dey, T. K. and Levine, J. A. (2008). Delpsc: A delaunay mesher for piecewise smooth complexes. In *Proceedings of the twenty-fourth annual symposium on Computational geometry - SCG '08*, pages 220–221, New York, NY, USA. ACM Press.
- Dey, T. K. and Ray, T. (2010). Polygonal surface remeshing with delaunay refinement. *Engineering with computers*, 26(3):289–301.

- DiCarlo, A., Paoluzzi, A., and Shapiro, V. (2014). Linear algebraic representation for topological structures. *Computer-Aided Design*, 46:269–274. 2013 SIAM Conference on Geometric and Physical Modeling.
- Dobrev, V., Knupp, P., Kolev, T., Mittal, K., and Tomov, V. (2019). The target-matrix optimization paradigm for high-order meshes. *SIAM Journal on Scientific Computing*, 41(1):B50–B68.
- Dobrzynski, C. and El Jannoun, G. (2017). High order mesh untangling for complex curved geometries. Research Report RR-9120, INRIA Bordeaux, équipe CARDAMOM.
- Dompierre, J., Labb  , P., Vallet, M.-G., and Camarero, R. (1999). How to subdivide pyramids, prisms, and hexahedra into tetrahedra. *IMR*, 99:195.
- Doran, C., Chang, A., and Bridson, R. (2013). Isosurface stuffing improved: Acute lattices and feature matching. In *ACM SIGGRAPH 2013 Talks on - SIGGRAPH ’13*, pages 38:1–38:1, New York, NY, USA. ACM Press.
- Dougherty, R., Faber, V., and Murphy, M. (2004). Unflippable tetrahedral complexes. *Discrete & Computational Geometry*, 32(3):309–315.
- Du, Q. and Wang, D. (2003a). Tetrahedral mesh generation and optimization based on centroidal voronoi tessellations. *International journal for numerical methods in engineering*, 56(9):1355–1373.
- Du, Q. and Wang, D. (2003b). Tetrahedral mesh generation and optimization based on centroidal voronoi tessellations. *International journal for numerical methods in engineering*, 56(9):1355–1373.
- Dufourd, J.-F. (1991). An obj3 functional specification for boundary representation. In *Proceedings of the first ACM symposium on Solid modeling foundations and CAD/CAM applications*, pages 61–72.
- Dunyach, M., Vanderhaeghe, D., Barthe, L., and Botsch, M. (2013). Adaptive remeshing for real-time mesh deformation.
- Dyer, R., Zhang, H., and M  ller, T. (2007). Delaunay mesh construction.
- Ebke, H.-C., Campen, M., Bommes, D., and Kobbelt, L. (2014). Level-of-detail quad meshing. *ACM Transactions on Graphics (TOG)*, 33(6):184.
- Engvall, L. and Evans, J. A. (2016). Isogeometric triangular bernstein–b  zier discretizations: Automatic mesh generation and geometrically exact finite element analysis. *Computer Methods in Applied Mechanics and Engineering*, 304:378 – 407.
- Engvall, L. and Evans, J. A. (2017). Isogeometric unstructured tetrahedral and mixed-element bernstein–b  zier discretizations. *Computer Methods in Applied Mechanics and Engineering*, 319:83 – 123.
- Engvall, L. and Evans, J. A. (2018). Mesh quality metrics for isogeometric bernstein–b  zier discretizations. *arXiv:1810.06975*.

- Engwirda, D. (2016). Conforming restricted delaunay mesh generation for piecewise smooth complexes. *CoRR*.
- Ezuz, D., Solomon, J., and Ben-Chen, M. (2019). Reversible harmonic maps between discrete surfaces. *ACM Trans. Graph.*, 38(2).
- Feng, L., Alliez, P., Busé, L., Delingette, H., and Desbrun, M. (2018a). Curved optimal delaunay triangulation. *ACM Transactions on Graphics (TOG)*.
- Feng, L., Alliez, P., Busé, L., Delingette, H., and Desbrun, M. (2018b). Curved optimal delaunay triangulation. *ACM Trans. Graph.*, 37(4):61:1–61:16.
- Floater, M. (2003). One-to-one piecewise linear mappings over triangulations. *Mathematics of Computation*, 72(242):685–696.
- Floater, M. S. (1997). Parametrization and smooth approximation of surface triangulations. *Computer Aided Geometric Design*, 14:231–250.
- Floater, M. S. and Hormann, K. (2005). Surface parameterization: a tutorial and survey. In *In Advances in Multiresolution for Geometric Modelling, Mathematics and Visualization*, pages 157–186. Springer Verlag.
- Fortunato, M. and Persson, P.-O. (2016). High-order unstructured curved mesh generation using the winslow equations. *Journal of Computational Physics*, 307:1 – 14.
- Foteinos, P. and Chrisochoides, N. (2011). Dynamic parallel 3d delaunay triangulation. In *Proceedings of the 20th International Meshing Roundtable*, pages 3–20. Springer.
- Fu, X.-M. and Liu, Y. (2016). Computing inversion-free mappings by simplex assembly. *ACM Trans. Graph.*, 35(6):216:1–216:12.
- Fu, X.-M., Liu, Y., and Guo, B. (2015a). Computing locally injective mappings by advanced mips. *ACM Trans. Graph.*, 34(4):71:1–71:12.
- Fu, X.-M., Liu, Y., and Guo, B. (2015b). Computing locally injective mappings by advanced mips. *ACM Transactions on Graphics (TOG)*, 34(4):71.
- Funke, D. and Sanders, P. (2017). Parallel d-d delaunay triangulations in shared and distributed memory. In *2017 Proceedings of the Nineteenth Workshop on Algorithm Engineering and Experiments (ALENEX)*, pages 207–217. SIAM.
- Galtier, J. and George, P. L. (1996). Prepartitioning as a way to mesh subdomains in parallel. In *in 5th International Meshing Roundtable*. Citeseer.
- Gao, X., Shen, H., and Panozzo, D. (2019). Feature preserving octree-based hexahedral meshing. *Computer Graphics Forum*, 38(5):135–149.

- Gargallo-Peiró, A., Roca, X., Peraire, J., and Sarrate, J. (2015). Optimization of a regularized distortion measure to generate curved high-order unstructured tetrahedral meshes. *International Journal for Numerical Methods in Engineering*, 103(5):342–363.
- Gargallo Peiró, A., Roca Navarro, F. J., Peraire Guitart, J., and Sarrate Ramos, J. (2013). High-order mesh generation on cad geometries. In *Adaptive Modeling and Simulation 2013*, pages 301–312. Centre Internacional de Mètodes Numèrics en Enginyeria (CIMNE).
- Garimella, R. V. and Shephard, M. S. (2000). Boundary layer mesh generation for viscous flow simulations. *International Journal for Numerical Methods in Engineering*, 49(1-2):193–218.
- Garland, M. and Heckbert, P. (1999). *Quadric-Based Polygonal Surface Simplification*. PhD thesis, USA. AAI9950005.
- Garland, M. and Heckbert, P. S. (1997). Surface simplification using quadric error metrics. In *Proceedings of the 24th annual conference on Computer graphics and interactive techniques*, pages 209–216. ACM Press/Addison-Wesley Publishing Co.
- Garland, M. and Heckbert, P. S. (1998). Simplifying surfaces with color and texture using quadric error metrics. In *Proceedings Visualization'98 (Cat. No. 98CB36276)*, pages 263–269. IEEE.
- Gärtner, B. and Schönherr, S. (2000). An efficient, exact, and generic quadratic programming solver for geometric optimization. In *Proceedings of the sixteenth annual symposium on Computational geometry*, pages 110–118.
- George, P. and Borouchaki, H. (2012). Construction of tetrahedral meshes of degree two. *International Journal for Numerical Methods in Engineering*, 90(9):1156–1182.
- Geuzaine, C., Johnen, A., Lambrechts, J., Remacle, J.-F., and Toulorge, T. (2015). *The Generation of Valid Curvilinear Meshes*, pages 15–39. Springer International Publishing, Cham.
- Geuzaine, C. and Remacle, J.-F. (2009). Gmsh: A 3-d finite element mesh generator with built-in pre- and post-processing facilities. *International Journal for Numerical Methods in Engineering*, 79(11):1309–1331.
- Ghasemi, A., Taylor, L. K., and Newman, III, J. C. (2016). Massively parallel curved spectral/finite element mesh generation of industrial cad geometries in two and three dimensions. *Fluids Engineering Division Summer Meeting*, (50299).
- Gillespie, M., Springborn, B., and Crane, K. (2021). Discrete conformal equivalence of polyhedral surfaces. *ACM Trans. Graph.*, 40(4).
- Gotsman, C. and Surazhsky, V. (2001). Guaranteed intersection-free polygon morphing. *Computers & Graphics*, 25(1):67–75.
- Grinspun, E., Krysl, P., and Schröder, P. (2002). Charms: A simple framework for adaptive simulation. *ACM Trans. Graph.*, 21(3):281–290.

- Gu, X., Guo, R., Luo, F., Sun, J., and Wu, T. (2018a). A discrete uniformization theorem for polyhedral surfaces ii. *Journal of Differential Geometry*, 109(3):431–466.
- Gu, X., Luo, F., Sun, J., and Wu, T. (2018b). A discrete uniformization theorem for polyhedral surfaces. *Journal of Differential Geometry*, 109(2):223–256.
- Guennebaud, G., Jacob, B., et al. (2010). Eigen v3.
- Guéziec, A. (1996). *Surface simplification inside a tolerance volume*. IBM TJ Watson Research Center.
- Guibas, L. and Stolfi, J. (1985). Primitives for the manipulation of general subdivisions and the computation of voronoi. *ACM Trans. Graph.*, 4(2):74–123.
- Guigue, P. and Devillers, O. (2003). Fast and robust triangle-triangle overlap test using orientation predicates. *Journal of graphics tools*, 8(1):25–32.
- Gumhold, S., Borodin, P., and Klein, R. (2003). Intersection free simplification. *International Journal of Shape Modeling*, 9(02):155–176.
- Guo, H.-X., Liu, X., Yan, D.-M., and Yang, L. (2020). Cut-enhanced polycube-maps for feature-aware all-hex meshing. *ACM Transactions on Graphics (TOG)*, 39(4):106:1–106:14.
- Hahmann, S. and Bonneau, G. . (2003). Polynomial surfaces interpolating arbitrary triangulations. *IEEE Transactions on Visualization and Computer Graphics*, 9(1):99–109.
- Haines, R. (2014). Moss: Multiple orthogonal strand system. In *Proceedings of the 22nd International Meshing Roundtable*, pages 75–91. Springer International Publishing, Cham.
- Harmon, D. (2010). *Robust, efficient, and accurate contact algorithms*. PhD thesis, Columbia University.
- Harmon, D., Panozzo, D., Sorkine, O., and Zorin, D. (2011). Interference-aware geometric modeling. *ACM Trans. Graph.*, 30(6):137:1–137:10.
- Harmon, D., Vouga, E., Smith, B., Tamstorf, R., and Grinspun, E. (2009). Asynchronous contact mechanics. *ACM Trans. Graph.*, 28(3):87:1–87:12.
- Hart, G. W. (2018). Conway notation for polyhedra. URL: http://www.ergehart.co/virtual-polyhedra/conway_notation.html.
- Hass, J. and Trnkova, M. (2020). Approximating isosurfaces by guaranteed-quality triangular meshes. In *Computer Graphics Forum*, volume 39, pages 29–40. Wiley Online Library.
- Hoppe, H. (1996). Progressive meshes. In *Proceedings of the 23rd annual conference on Computer graphics and interactive techniques*, pages 99–108. ACM.

- Hoppe, H., DeRose, T., Duchamp, T., Halstead, M., Jin, H., McDonald, J., Schweitzer, J., and Stuetzle, W. (1994). Piecewise smooth surface reconstruction. In *Proceedings of the 21st annual conference on Computer graphics and interactive techniques*, pages 295–302.
- Hormann, K. and Greiner, G. (2000a). MIPS: An efficient global parametrization method. In *Curve and Surface Design: Saint-Malo 1999*, pages 153–162.
- Hormann, K. and Greiner, G. (2000b). Mips: An efficient global parametrization method. Technical report, ERLANGEN-NUERNBERG UNIV (GERMANY) COMPUTER GRAPHICS GROUP.
- Hormann, K., Lévy, B., and Sheffer, A. (2007). Mesh parameterization: Theory and practice. In *ACM SIGGRAPH 2007 Courses*, SIGGRAPH '07, New York, NY, USA. ACM.
- Hormann, K. and Sukumar, N., editors (2017). *Generalized Barycentric Coordinates in Computer Graphics and Computational Mechanics*. CRC Press, Boca Raton, FL.
- Hu, K., Yan, D., Bommes, D., Alliez, P., and Benes, B. (2017). Error-bounded and feature preserving surface remeshing with minimal angle improvement. *IEEE Transactions on Visualization and Computer Graphics*, 23(12):2560–2573.
- Hu, K., Yan, D.-M., Bommes, D., Alliez, P., and Benes, B. (2016). Error-bounded and feature preserving surface remeshing with minimal angle improvement. *IEEE transactions on visualization and computer graphics*, 23(12):2560–2573.
- Hu, Y., Li, T.-M., Anderson, L., Ragan-Kelley, J., and Durand, F. (2019a). Taichi: a language for high-performance computation on spatially sparse data structures. *ACM Transactions on Graphics (TOG)*, 38(6):1–16.
- Hu, Y., Schneider, T., Gao, X., Zhou, Q., Jacobson, A., Zorin, D., and Panozzo, D. (2019b). Triwild: robust triangulation with curve constraints. *ACM Transactions on Graphics (TOG)*, 38(4):52.
- Hu, Y., Schneider, T., Gao, X., Zhou, Q., Jacobson, A., Zorin, D., and Panozzo, D. (2019c). Triwild: Robust triangulation with curve constraints. *ACM Trans. Graph.*, 38(4):52:1–52:15.
- Hu, Y., Schneider, T., Wang, B., Zorin, D., and Panozzo, D. (2019d). Fast tetrahedral meshing in the wild. *arXiv preprint arXiv:1908.03581*.
- Hu, Y., Schneider, T., Wang, B., Zorin, D., and Panozzo, D. (2020). Fast tetrahedral meshing in the wild. *ACM Trans. Graph.*, 39(4).
- Hu, Y., Zhou, Q., Gao, X., Jacobson, A., Zorin, D., and Panozzo, D. (2018). Tetrahedral meshing in the wild. *ACM Trans. Graph.*, 37(4):60–1.
- Huang, J., Liu, X., Jiang, H., Wang, Q., and Bao, H. (2007). Gradient-based shell generation and deformation. *Computer Animation and Virtual Worlds*, 18(4-5):301–309.

- Hutchinson, D., Preston, M., and Hewitt, T. (1996). Adaptive refinement for mass/spring simulations. In *Proceedings of the Eurographics Workshop on Computer Animation and Simulation '96*, page 31–45, Berlin, Heidelberg. Springer-Verlag.
- Jacobson, A., Kavan, L., and Sorkine-Hornung, O. (2013). Robust inside-outside segmentation using generalized winding numbers. *ACM Trans. Graph.*, 32(4).
- Jacobson, A., Panozzo, D., et al. (2014). libigl: A simple C++ geometry processing library. <http://igl.ethz.ch/projects/libigl/>.
- Jacobson, A., Panozzo, D., Schüller, C., Diamanti, O., Zhou, Q., Pietroni, N., et al. (2016). libigl: A simple c++ geometry processing library, 2016.
- Jameson, A., Alonso, J., and McMullen, M. (2002). Application of a non-linear frequency domain solver to the euler and navier-stokes equations. In *40th AIAA Aerospace Sciences Meeting & Exhibit*.
- Jamin, C., Alliez, P., Yvinec, M., and Boissonnat, J.-D. (2015a). Cgalmesh: a generic framework for delaunay mesh generation. *ACM Transactions on Mathematical Software (TOMS)*, 41(4):23.
- Jamin, C., Alliez, P., Yvinec, M., and Boissonnat, J.-D. (2015b). CGALmesh: A generic framework for delaunay mesh generation. *ACM Transactions on Mathematical Software*, 41(4):1–24.
- Jiang, Z., Schaefer, S., and Panozzo, D. (2017). Simplicial complex augmentation framework for bijective maps. *ACM Transactions on Graphics*, 36(6).
- Jiang, Z., Schneider, T., Zorin, D., and Panozzo, D. (2020). Bijective projection in a shell. *ACM Transactions on Graphics (TOG)*, 39(6):1–18.
- Jiao, X. and Heath, M. T. (2004). Overlaying surface meshes, part i: Algorithms. *International Journal of Computational Geometry & Applications*, 14(06):379–402.
- Jiao, X. and Wang, D. (2012). Reconstructing high-order surfaces for meshing. *Engineering with Computers*, 28(4):361–373.
- Jimenez, P., Thomas, F., and Torras, C. (2001). 3d collision detection: a survey. *Computers & Graphics*, 25(2):269 – 285.
- Jin, M., Kim, J., Luo, F., and Gu, X. (2008). Discrete surface ricci flow. *IEEE Transactions on Visualization and Computer Graphics*, 14(5):1030–1043.
- Jin, Y., Song, D., Wang, T., Huang, J., Song, Y., and He, L. (2019). A shell space constrained approach for curve design on surface meshes. *Computer-Aided Design*, 113:24–34.
- Johnen, A., Remacle, J.-F., and Geuzaine, C. (2013). Geometrical validity of curvilinear finite elements. *Journal of Computational Physics*, 233:359–372.

- Karras, T. (2012). Maximizing parallelism in the construction of bvhs, octrees, and k-d trees. In *Proceedings of the Fourth ACM SIGGRAPH/Eurographics conference on High-Performance Graphics*, pages 33–37.
- Karypis, G. and Kumar, V. (1998). A fast and high quality multilevel scheme for partitioning irregular graphs. *SIAM Journal on scientific Computing*, 20(1):359–392.
- Kazhdan, M., Solomon, J., and Ben-Chen, M. (2012). Can mean-curvature flow be modified to be non-singular? *Comput. Graph. Forum*, 31(5):1745–1754.
- Kharevych, L., Springborn, B., and Schröder, P. (2006). Discrete conformal mappings via circle patterns. *ACM Trans. Graph.*, 25(2):412–438.
- Kjolstad, F., Chou, S., Lugato, D., Kamil, S., and Amarasinghe, S. (2017). Taco: A tool to generate tensor algebra kernels. In *2017 32nd IEEE/ACM International Conference on Automated Software Engineering (ASE)*, pages 943–948. IEEE.
- Kjolstad, F., Kamil, S., Ragan-Kelley, J., Levin, D. I., Sueda, S., Chen, D., Vouga, E., Kaufman, D. M., Kanwar, G., Matusik, W., et al. (2016). Simit: A language for physical simulation. *ACM Transactions on Graphics (TOG)*, 35(2):1–21.
- Klingner, B. and Shewchuk, J. (2007). Aggressive tetrahedral mesh improvement. pages 3–23.
- Klingner, B. M. (2009). *Tetrahedral mesh improvement*. PhD thesis, University of California at Berkeley.
- Klingner, B. M., Feldman, B. E., Chentanez, N., and O’Brien, J. F. (2006). Fluid animation with dynamic meshes. In *ACM SIGGRAPH 2006 Papers*, SIGGRAPH ’06, page 820–825, New York, NY, USA. Association for Computing Machinery.
- Knabner, P. and Summ, G. (2001). The invertibility of the isoparametric mapping for pyramidal and prismatic finite elements. *Numerische Mathematik*, 88(4):661–681.
- Knupp, P. M. (2000). Achieving finite element mesh quality via optimization of the jacobian matrix norm and associated quantities. part i, a framework for surface mesh optimization. *International Journal for Numerical Methods in Engineering*, 48(3):401–420.
- Knupp, P. M. (2002). Achieving finite element mesh quality via optimization of the jacobian matrix norm and associated quantities. part ii, a framework for volume mesh optimization and the condition number of the jacobian matrix. *International Journal for Numerical Methods in Engineering*, 48(8):1165–1185.
- Kobbelt, L., Campagna, S., Vorsatz, J., and Seidel, H.-P. (1998). Interactive multi-resolution modeling on arbitrary meshes. In *Siggraph*, volume 98, pages 105–114.
- Koch, S., Matveev, A., Jiang, Z., Williams, F., Artemov, A., Burnaev, E., Alexa, M., Zorin, D., and Panizzo, D. (2019). Abc: A big cad model dataset for geometric deep learning. In *The IEEE Conference on Computer Vision and Pattern Recognition (CVPR)*.

- Kovalsky, S. Z., Galun, M., and Lipman, Y. (2016). Accelerated quadratic proxy for geometric optimization. *ACM Trans. Graph.*, 35(4):134:1–134:11.
- Kraemer, P., Untereiner, L., Jund, T., Thery, S., and Cazier, D. (2014). Cgogn: N-dimensional meshes with combinatorial maps. In *Proceedings of the 22nd International Meshing Roundtable*, pages 485–503. Springer.
- Kraevoy, V. and Sheffer, A. (2004). Cross-parameterization and compatible remeshing of 3d models. *ACM Transactions on Graphics (TOG)*, 23(3):861–869.
- Kremer, M., Bommes, D., and Kobbelt, L. (2013). Openvolumemesh—a versatile index-based data structure for 3d polytopal complexes. In *Proceedings of the 21st International Meshing Roundtable*, pages 531–548. Springer.
- Krishnamurthy, V. and Levoy, M. (1996). Fitting smooth surfaces to dense polygon meshes. In *Proceedings of the 23rd annual conference on Computer graphics and interactive techniques*, pages 313–324.
- Labelle, F. and Shewchuk, J. R. (2007). Isosurface stuffing: Fast tetrahedral meshes with good dihedral angles. In *ACM SIGGRAPH 2007 papers on - SIGGRAPH '07*, page 57, New York, NY, USA. ACM Press.
- Lachat, C., Dobrzynski, C., and Pellegrini, F. (2014). Parallel mesh adaptation using parallel graph partitioning. In *5th European conference on computational mechanics (ECCM V)*, volume 3, pages 2612–2623. CIMNE-International Center for Numerical Methods in Engineering.
- Lee, A., Moreton, H., and Hoppe, H. (2000). Displaced subdivision surfaces. In *Proceedings of the 27th annual conference on Computer graphics and interactive techniques*, pages 85–94.
- Lee, A. W., Sweldens, W., Schröder, P., Cowsar, L. C., and Dobkin, D. P. (1998). Maps: Multiresolution adaptive parameterization of surfaces. In *Siggraph*, volume 98, pages 95–104.
- Lengyel, J., Praun, E., Finkelstein, A., and Hoppe, H. (2001). Real-time fur over arbitrary surfaces. In *Proceedings of the 2001 symposium on Interactive 3D graphics*, pages 227–232. ACM.
- Lévy, B. (2015). Geogram.
- Lévy, B., Petitjean, S., Ray, N., and Maillot, J. (2002). Least squares conformal maps for automatic texture atlas generation. *ACM Trans. Graph.*, 21(3):362–371.
- Li, L. and Volkov, V. (2005). Cloth animation with adaptively refined meshes. In *Proceedings of the Twenty-Eighth Australasian Conference on Computer Science - Volume 38*, ACSC '05, page 107–113, AUS. Australian Computer Society, Inc.
- Li, M., Ferguson, Z., Schneider, T., Langlois, T., Zorin, D., Panizzo, D., Jiang, C., and Kaufman, D. M. (2020). Incremental potential contact: Intersection- and inversion-free large deformation dynamics. *ACM Trans. Graph. (SIGGRAPH)*, 39(4).

- Li, M., Kaufman, D. M., Kim, V. G., Solomon, J., and Sheffer, A. (2018). Optcuts: joint optimization of surface cuts and parameterization. *ACM Transactions on Graphics (TOG)*, 37(6):1–13.
- Li, Y., Liu, Y., Xu, W., Wang, W., and Guo, B. (2012). All-hex meshing using singularity-restricted field. *ACM Trans. Graph.*, 31(6).
- Lienhardt, P. (1994). N-dimensional generalized combinatorial maps and cellular quasi-manifolds. *International Journal of Computational Geometry & Applications*, 4(03):275–324.
- Lin, H., Chen, W., and Bao, H. (2007). Adaptive patch-based mesh fitting for reverse engineering. *Computer-Aided Design*, 39(12):1134 – 1142.
- Linardakis, L. and Chrisochoides, N. (2006). Delaunay decoupling method for parallel guaranteed quality planar mesh refinement. *SIAM Journal on Scientific Computing*, 27(4):1394–1423.
- Lipman, Y. (2012). Bounded distortion mapping spaces for triangular meshes. *ACM Trans. Graph.*, 31(4):108:1–108:13.
- Lipman, Y. (2013). Construction of injective mappings of meshes. *CoRR*, abs/1310.0955.
- Lipman, Y. (2014). Bijective mappings of meshes with boundary and the degree in mesh processing. *SIAM Journal on Imaging Sciences*, 7(2):1263–1283.
- Litke, N., Droske, M., Rumpf, M., and Schröder, P. (2005). An image processing approach to surface matching. In *Symposium on Geometry Processing*, volume 255.
- Liu, L., Zhang, L., Xu, Y., Gotsman, C., and Gortler, S. J. (2008). A local/global approach to mesh parameterization. In *Proceedings of the Symposium on Geometry Processing*, SGP ’08, pages 1495–1504, Aire-la-Ville, Switzerland, Switzerland. Eurographics Association.
- Liu, S., Ferguson, Z., Jacobson, A., and Gingold, Y. I. (2017). Seamless: seam erasure and seam-aware decoupling of shape from mesh resolution. *ACM Trans. Graph.*, 36(6).
- Liu, Y.-J., Xu, C.-X., Fan, D., and He, Y. (2015). Efficient construction and simplification of delaunay meshes. *ACM Transactions on Graphics (TOG)*, 34(6).
- Lo, S. (2012). Parallel delaunay triangulation in three dimensions. *Computer Methods in Applied Mechanics and Engineering*, 237:88–106.
- Loriot, S., Rouxel-Labbé, M., Tournois, J., and Yaz, I. O. (2020). Polygon mesh processing. In *CGAL User and Reference Manual*. CGAL Editorial Board, 5.0.3 edition.
- Loseille, A., Alauzet, F., and Menier, V. (2017). Unique cavity-based operator and hierarchical domain partitioning for fast parallel generation of anisotropic meshes. *Computer-Aided Design*, 85:53–67.

- Lu, Q., Shephard, M. S., Tendulkar, S., and Beall, M. W. (2013). Parallel curved mesh adaptation for large scale high-order finite element simulations. In Jiao, X. and Weill, J.-C., editors, *Proceedings of the 21st International Meshing Roundtable*, pages 419–436, Berlin, Heidelberg. Springer Berlin Heidelberg.
- Lu, Q., Shephard, M. S., Tendulkar, S., and Beall, M. W. (2014). Parallel mesh adaptation for high-order finite element methods with curved element geometry. *Engineering with Computers*, 30(2):271–286.
- Luo, F. (2004). Combinatorial Yamabe flow on surfaces. *Communications in Contemporary Mathematics*, 6(05):765–780.
- Luo, X., Shephard, M. S., Lee, L.-Q., Ng, C., and Ge, L. (2008). Tracking adaptive moving mesh refinements in 3d curved domains for large-scale higher order finite element simulations. In *Proceedings of the 17th International Meshing Roundtable*, pages 585–601. Springer.
- Luo, X., Shephard, M. S., and Remacle, J.-F. (2001). The influence of geometric approximation on the accuracy of high order methods. *Rensselaer SCOREC report*, 1.
- Luo, X., Shephard, M. S., Remacle, J.-F., O’Bara, R. M., Beall, M. W., Szabó, B. A., and Actis, R. (2002). p-version mesh generation issues. In *IMR*.
- Lyche, T. and Muntingh, G. (2015). Simplex spline bases on the Powell-Sabin 12-split: Part i. *arXiv preprint arXiv:1505.01798*.
- Mahmoud, A. H., Porumbescu, S. D., and Owens, J. D. (2021). Rxmesh: A gpu mesh data structure. *ACM Trans. Graph.*, 40(4).
- Mandad, M., Cohen-Steiner, D., and Alliez, P. (2015a). Isotopic approximation within a tolerance volume. *ACM Transactions on Graphics (TOG)*, 34(4):64.
- Mandad, M., Cohen-Steiner, D., and Alliez, P. (2015b). Isotopic approximation within a tolerance volume. *ACM Trans. Graph.*, 34(4):64:1–64:12.
- Mäntylä, M. (1987). *An introduction to solid modeling*. Computer Science Press, Inc.
- Marcon, J., Peiró, J., Moxey, D., Bergemann, N., Bucklow, H., and Gammon, M. R. (2019). A semi-structured approach to curvilinear mesh generation around streamlined bodies. In *AIAA Scitech 2019 Forum*, page 1725.
- Maron, H., Galun, M., Aigerman, N., Trope, M., Dym, N., Yumer, E., Kim, V. G., and Lipman, Y. (2017). Convolutional neural networks on surfaces via seamless toric covers. *ACM Trans. Graph.*, 36(4):71–1.
- Marot, C., Pellerin, J., and Remacle, J.-F. (2019). One machine, one minute, three billion tetrahedra. *International Journal for Numerical Methods in Engineering*, 117(9):967–990.

- Marot, C. and Remacle, J.-F. (2020). Quality tetrahedral mesh generation with hxt. *arXiv preprint arXiv:2008.08508*.
- Marschner, Z., Palmer, D., Zhang, P., and Solomon, J. (2020). Hexahedral mesh repair via sum-of-squares relaxation. In *Computer Graphics Forum*, volume 39, pages 133–147. Wiley Online Library.
- Mezger, J., Thomaszewski, B., Pabst, S., and Straßer, W. (2007). A finite element method for interactive physically based shape modelling with quadratic tetrahedra.
- Mezger, J., Thomaszewski, B., Pabst, S., and Straßer, W. (2009). Interactive physically-based shape editing. *Computer Aided Geometric Design*, 26(6):680 – 694. Solid and Physical Modeling 2008.
- Misztal, M. K. and Bærentzen, J. A. (2012). Topology-adaptive interface tracking using the deformable simplicial complex. *ACM Trans. Graph.*, 31(3).
- Mitchell, J. S., Mount, D. M., and Papadimitriou, C. H. (1987). The discrete geodesic problem. *SIAM Journal on Computing*, 16(4):647–668.
- Möbius, J. and Kobbelt, L. (2010). Openflipper: An open source geometry processing and rendering framework. In *International Conference on Curves and Surfaces*, pages 488–500. Springer.
- Molino, N., Bridson, R., and Fedkiw, R. (2003). Tetrahedral mesh generation for deformable bodies. In *Proc. Symposium on Computer Animation*.
- Moxey, D., Ekelschot, D., Keskin, U., Sherwin, S., and Peiro, J. (2016). High-order curvilinear meshing using a thermo-elastic analogy. *Computer-Aided Design*, 72:130 – 139. 23rd International Meshing Roundtable Special Issue: Advances in Mesh Generation.
- Moxey, D., Green, M., Sherwin, S., and Peiró, J. (2015). An isoparametric approach to high-order curvilinear boundary-layer meshing. *Computer Methods in Applied Mechanics and Engineering*, 283:636 – 650.
- Moxey, D., Turner, M., Marcon, J., and Peiro, J. (2018). Nekmesh: An open-source high-order mesh generator.
- Müller, M., Chentanez, N., Kim, T.-Y., and Macklin, M. (2015). Air meshes for robust collision handling. *ACM Trans. Graph.*, 34(4).
- Murphy, M., Mount, D. M., and Gable, C. W. (2001a). A point-placement strategy for conforming delaunay tetrahedralization. *International Journal of Computational Geometry & Applications*, 11(06):669–682.
- Murphy, M., Mount, D. M., and Gable, C. W. (2001b). A point-placement strategy for conforming delaunay tetrahedralization. *International Journal of Computational Geometry & Applications*, 11(06):669–682.

- Myles, A., Pietroni, N., and Zorin, D. (2014). Robust field-aligned global parametrization. *ACM Trans. Graph.*, 33(4):135:1–135:14.
- Narain, R., Pfaff, T., and O'Brien, J. F. (2013). Folding and crumpling adaptive sheets. *ACM Trans. Graph.*, 32(4).
- Narain, R., Samii, A., and O'Brien, J. F. (2012). Adaptive anisotropic remeshing for cloth simulation. *ACM Trans. Graph.*, 31(6).
- Nguyen, H. (2007). *Gpu gems 3*. Addison-Wesley Professional.
- Oden, J. (1994). Optimal h-p finite element methods. *Computer Methods in Applied Mechanics and Engineering*, 112(1):309 – 331.
- Okusanya, T. and Peraire, J. (1996). Parallel unstructured mesh generation.
- Okusanya, T. and Peraire, J. (1997). 3-d parallel unstructured mesh generation. In *Proc. Joint ASME/ASCE/SES Summer Meeting*. Citeseer.
- Ovsjanikov, M., Ben-Chen, M., Solomon, J., Butscher, A., and Guibas, L. (2012). Functional maps: A flexible representation of maps between shapes. *ACM Trans. Graph.*, 31(4).
- Ovsjanikov, M., Corman, E., Bronstein, M., Rodolà, E., Ben-Chen, M., Guibas, L., Chazal, F., and Bronstein, A. (2017). Computing and processing correspondences with functional maps. In *ACM SIGGRAPH 2017 Courses*, SIGGRAPH '17.
- Pal, K., Schüller, C., Panozzo, D., Sorkine-Hornung, O., and Weyrich, T. (2014). Content-aware surface parameterization for interactive restoration of historical documents. *Computer Graphics Forum (proceedings of EUROGRAPHICS issue)*, 33(2).
- Palmer, D., Bommes, D., and Solomon, J. (2020). Algebraic representations for volumetric frame fields. *ACM Trans. Graph.*, 39(2).
- Panozzo, D., Baran, I., Diamanti, O., and Sorkine-Hornung, O. (2013). Weighted averages on surfaces. *ACM Transactions on Graphics (TOG)*, 32(4):60.
- Peiró, A. G., Gironés, E. R., Navarro, F. J., and Ramos, J. S. (2015). On curving high-order hexahedral meshes.
- Peiró, A. G., Roca, X., Peraire, J., and Sarrate, J. (2014). Defining quality measures for validation and generation of high-order tetrahedral meshes. In Sarrate, J. and Staten, M., editors, *Proceedings of the 22nd International Meshing Roundtable*, pages 109–126, Cham. Springer International Publishing.
- Peiró, J., Sherwin, S. J., and Giordana, S. (2008). Automatic reconstruction of a patient-specific high-order surface representation and its application to mesh generation for cfd calculations. *Medical & Biological Engineering & Computing*, 46(11):1069–1083.

- Peng, J., Kristjansson, D., and Zorin, D. (2004). Interactive modeling of topologically complex geometric detail. In *ACM Transactions on Graphics (TOG)*, volume 23, pages 635–643. ACM.
- Persson, P.-O. and Peraire, J. (2009). Curved mesh generation and mesh refinement using lagrangian solid mechanics. In *47th AIAA Aerospace Sciences Meeting including The New Horizons Forum and Aerospace Exposition*.
- Pfaff, T., Narain, R., de Joya, J. M., and O'Brien, J. F. (2014). Adaptive tearing and cracking of thin sheets. *ACM Trans. Graph.*, 33(4).
- Phong, B. T. (1975). Illumination for computer generated pictures. *Communications of the ACM*, 18(6):311–317.
- Pietroni, N., Tarini, M., and Cignoni, P. (2010). Almost isometric mesh parameterization through abstract domains. *IEEE Transactions on Visualization and Computer Graphics*, 16(4):621–635.
- Pointwise (2018). High order mesh generation at pointwise. <https://www.pointwise.com/theconnector/2016-Q3/High-Order-Mesh-Generation-at-Pointwise.html>. Accessed: 2018-11-14.
- Popović, J. and Hoppe, H. (1997). Progressive simplicial complexes. In *Proceedings of the 24th Annual Conference on Computer Graphics and Interactive Techniques*, SIGGRAPH '97, page 217–224, USA. ACM Press/Addison-Wesley Publishing Co.
- Poranne, R. and Lipman, Y. (2014). Provably good planar mappings. *ACM Trans. Graph.*, 33(4):76:1–76:11.
- Porumbescu, S. D., Budge, B., Feng, L., and Joy, K. I. (2005). Shell maps. In *ACM Transactions on Graphics (TOG)*, volume 24, pages 626–633. ACM.
- Poya, R., Sevilla, R., and Gil, A. J. (2016). A unified approach for a posteriori high-order curved mesh generation using solid mechanics. *Computational Mechanics*, 58(3):457–490.
- Praun, E., Sweldens, W., and Schröder, P. (2001). Consistent mesh parameterizations. In *Proceedings of the 28th annual conference on Computer graphics and interactive techniques*, pages 179–184.
- Rabinovich, M., Poranne, R., Panozzo, D., and Sorkine-Hornung, O. (2017). Scalable locally injective mappings. *ACM Trans. Graph.*, 36(2):16:1–16:16.
- Ragan-Kelley, J., Barnes, C., Adams, A., Paris, S., Durand, F., and Amarasinghe, S. (2013). Halide: a language and compiler for optimizing parallelism, locality, and recomputation in image processing pipelines. *Acm Sigplan Notices*, 48(6):519–530.
- Remacle, J.-F. (2017). A two-level multithreaded delaunay kernel. *Computer-Aided Design*, 85:2–9.
- Remacle, J.-F., Toulorge, T., and Lambrechts, J. (2013). Robust untangling of curvilinear meshes. In *Proceedings of the 21st International meshing roundtable*, pages 71–83. Springer.

- Requicha, A. G. (1980). Representations for rigid solids: Theory, methods, and systems. *ACM Comput. Surv.*, 12(4):437–464.
- Roca, X., Gargallo-Peiró, A., and Sarrate, J. (2012). Defining quality measures for high-order planar triangles and curved mesh generation. In Quadros, W. R., editor, *Proceedings of the 20th International Meshing Roundtable*, pages 365–383, Berlin, Heidelberg. Springer Berlin Heidelberg.
- Ruiz-Gironés, E., Gargallo-Peiró, A., Sarrate, J., and Roca, X. (2017). An augmented lagrangian formulation to impose boundary conditions for distortion based mesh moving and curving. *Procedia Engineering*, 203:362 – 374. 26th International Meshing Roundtable, IMR26, 18-21 September 2017, Barcelona, Spain.
- Ruiz-Gironés, E., Roca, X., and Sarrate, J. (2016a). High-order mesh curving by distortion minimization with boundary nodes free to slide on a 3d cad representation. *Computer-Aided Design*, 72:52 – 64. 23rd International Meshing Roundtable Special Issue: Advances in Mesh Generation.
- Ruiz-Gironés, E., Sarrate, J., and Roca, X. (2016b). Generation of curved high-order meshes with optimal quality and geometric accuracy. *Procedia Engineering*, 163:315 – 327. 25th International Meshing Roundtable.
- Ruppert, J. (1995a). A delaunay refinement algorithm for quality 2-dimensional mesh generation. *Journal of algorithms*, 18(3):548–585.
- Ruppert, J. (1995b). A delaunay refinement algorithm for quality 2-dimensional mesh generation. *Journal of Algorithms*, 18(3):548–585.
- Sacht, L., Jacobson, A., Panozzo, D., Schüller, C., and Sorkine-Hornung, O. (2013). Consistent volumetric discretizations inside self-intersecting surfaces. *Computer Graphics Forum (proceedings of EUROGRAPHICS/ACM SIGGRAPH Symposium on Geometry Processing)*, 32(5):147–156.
- Sacht, L., Vouga, E., and Jacobson, A. (2015). Nested cages. *ACM Transactions on Graphics (TOG)*, 34(6):170.
- Sadek, E. A. (1980). A scheme for the automatic generation of triangular finite elements. *International Journal for Numerical Methods in Engineering*, 15(12):1813–1822.
- Salari, K. and Knupp, P. (2000). Code verification by the method of manufactured solutions. Technical report.
- Sander, P. V., Snyder, J., Gortler, S. J., and Hoppe, H. (2001). Texture mapping progressive meshes. In *ACM SIGGRAPH*, pages 409–416.
- Schlömer, N. (2020). nschloe/meshio: Input/output for many mesh formats. *Zenodo. doi*, 10.
- Schmidt, P., Born, J., Campen, M., and Kobbelt, L. (2019). Distortion-minimizing injective maps between surfaces. *ACM Transactions on Graphics (TOG)*, 38.

- Schneider, T., Hu, Y., Gao, X., Dumas, J., Zorin, D., and Panozzo, D. (2019). A large scale comparison of tetrahedral and hexahedral elements for finite element analysis.
- Schreiner, J., Asirvatham, A., Praun, E., and Hoppe, H. (2004). Inter-surface mapping. *ACM Trans. Graph.*, 23(3):870–877.
- Schüller, C., Kavan, L., Panozzo, D., and Sorkine-Hornung, O. (2013). Locally injective mappings. In *Symposium on Geometry Processing*, pages 125–135.
- Sharp, N. and Crane, K. (2020). A Laplacian for Nonmanifold Triangle Meshes. *Computer Graphics Forum (SGP)*, 39(5).
- Sharp, N., Soliman, Y., and Crane, K. (2019). Navigating intrinsic triangulations. *ACM Transactions on Graphics (TOG)*, 38(4):55.
- Sheffer, A., Praun, E., and Rose, K. (2006). Mesh parameterization methods and their applications. *Found. Trends. Comput. Graph. Vis.*, 2(2):105–171.
- Shephard, M. S., Flaherty, J. E., Jansen, K. E., Li, X., Luo, X., Chevaugeon, N., Remacle, J.-F., Beall, M. W., and O’Bara, R. M. (2005). Adaptive mesh generation for curved domains. *Applied Numerical Mathematics*, 52(2):251 – 271. ADAPT ’03: Conference on Adaptive Methods for Partial Differential Equations and Large-Scale Computation.
- Sherwin, S. and Peiró, J. (2002). Mesh generation in curvilinear domains using high-order elements. *International Journal for Numerical Methods in Engineering*, 53(1):207–223.
- Shewchuk, J. (1996a). Triangle: Engineering a 2d quality mesh generator and delaunay triangulator. *Applied computational geometry towards geometric engineering*, pages 203–222.
- Shewchuk, J. R. (1996b). Triangle: Engineering a 2d quality mesh generator and delaunay triangulator. In Lin, M. C. and Manocha, D., editors, *Applied Computational Geometry Towards Geometric Engineering*, pages 203–222, Berlin, Heidelberg. Springer Berlin Heidelberg.
- Shewchuk, J. R. (1997). Adaptive precision floating-point arithmetic and fast robust geometric predicates. *Discrete & Computational Geometry*, 18(3).
- Shewchuk, J. R. (1998a). Tetrahedral mesh generation by delaunay refinement. In *Proceedings of the fourteenth annual symposium on Computational geometry*, pages 86–95. ACM.
- Shewchuk, J. R. (1998b). Tetrahedral mesh generation by delaunay refinement. In *Proceedings of the fourteenth annual symposium on Computational geometry - SCG ’98*, pages 86–95, New York, NY, USA. ACM Press.
- Shewchuk, J. R. (2002a). Constrained delaunay tetrahedralizations and provably good boundary recovery. In *IMR*, pages 193–204.

- Shewchuk, J. R. (2002b). Constrained delaunay tetrahedralizations and provably good boundary recovery. In *Eleventh International Meshing Roundtable*, pages 193–204. Sandia National Laboratories.
- Shoham, M., Vaxman, A., and Ben-Chen, M. (2019). Hierarchical functional maps between subdivision surfaces. *Computer Graphics Forum*, 38(5):55–73.
- Si, H. (2015a). Tetgen, a delaunay-based quality tetrahedral mesh generator. *ACM Transactions on Mathematical Software (TOMS)*, 41(2):1–36.
- Si, H. (2015b). Tetgen, a delaunay-based quality tetrahedral mesh generator. *ACM Trans. Math. Softw.*, 41(2):11:1–11:36.
- Si, H. and Gärtner, K. (2005). Meshing piecewise linear complexes by constrained delaunay tetrahedralizations. In *IMR*, pages 147–163. Springer.
- Si, H. and Gartner, K. (2005). Meshing piecewise linear complexes by constrained delaunay tetrahedralizations. In *Proceedings of the 14th international meshing roundtable*, pages 147–163, Berlin, Heidelberg. Springer Berlin Heidelberg.
- Si, H. and Shewchuk, J. R. (2014a). Incrementally constructing and updating constrained delaunay tetrahedralizations with finite-precision coordinates. *Eng. Comput. (Lond.)*, 30(2):253–269.
- Si, H. and Shewchuk, J. R. (2014b). Incrementally constructing and updating constrained delaunay tetrahedralizations with finite-precision coordinates. *Engineering with Computers*, 30(2):253–269.
- Sieger, D. and Botsch, M. (2019). The polygon mesh processing library. <http://www.pmp-library.org>.
- Simnett, T. J. R., Laycock, S. D., and Day, A. M. (2009). An Edge-based Approach to Adaptively Refining a Mesh for Cloth Deformation. In Tang, W. and Collomosse, J., editors, *Theory and Practice of Computer Graphics*. The Eurographics Association.
- Smith, J. and Schaefer, S. (2015). Bijective parameterization with free boundaries. *ACM Trans. Graph.*, 34(4):70:1–70:9.
- Solomon, J. (2015). *Numerical Algorithms: Methods for Computer Vision, Machine Learning, and Graphics*. A. K. Peters, Ltd., Natick, MA, USA.
- Sorger, C., Frischmann, F., Kollmannsberger, S., and Rank, E. (2014). Tum.geoframe: automated high-order hexahedral mesh generation for shell-like structures. *Eng. Comput.*, 30(1):41–56.
- Sorkine, O. and Alexa, M. (2007). As-rigid-as-possible surface modeling. In *Proceedings of the Fifth Eurographics Symposium on Geometry Processing*, SGP ’07, pages 109–116, Aire-la-Ville, Switzerland. Eurographics Association.
- Sorkine, O., Cohen-Or, D., Goldenthal, R., and Lischinski, D. (2002). Bounded-distortion piecewise mesh parameterization. In *Proceedings of the Conference on Visualization*, pages 355–362.

- Springborn, B. (2020). Ideal hyperbolic polyhedra and discrete uniformization. *Discrete & Computational Geometry*, 64(1):63–108.
- Springborn, B., Schröder, P., and Pinkall, U. (2008). Conformal equivalence of triangle meshes. *ACM Trans. Graph.*, 27(3):1–11.
- Stees, M. and Shontz, S. M. (2017). A high-order log barrier-based mesh generation and warping method. *Procedia Engineering*, 203:180 – 192. 26th International Meshing Roundtable, IMR26, 18-21 September 2017, Barcelona, Spain.
- Stellato, B., Banjac, G., Goulart, P., Bemporad, A., and Boyd, S. (2017). OSQP: An operator splitting solver for quadratic programs. *ArXiv e-prints*.
- Stephenson, K. (2005). *Introduction to circle packing: The theory of discrete analytic functions*. Cambridge University Press.
- Steve L. Karman, J T. Erwin, R. S. G. and Stefanski, D. (2016). High-order mesh curving using wcn mesh optimization. In *46th AIAA Fluid Dynamics Conference, AIAA AVIATION Forum*.
- Sun, J., Wu, T., Gu, X., and Luo, F. (2015). Discrete conformal deformation: algorithm and experiments. *SIAM Journal on Imaging Sciences*, 8(3):1421–1456.
- Surazhsky, V. and Gotsman, C. (2001). Morphing stick figures using optimized compatible triangulations. In *Computer Graphics and Applications, 2001. Proceedings. Ninth Pacific Conference on*, pages 40–49. IEEE.
- Suwelack, S., Lukarski, D., Heuveline, V., Dillmann, R., and Speidel, S. (2013). Accurate surface embedding for higher order finite elements. In *Proceedings of the 12th ACM SIGGRAPH/Eurographics Symposium on Computer Animation*, SCA’13, pages 187–192, New York, NY, USA. ACM.
- Teschner, M., Heidelberger, B., Müller, M., Pomerantes, D., and Gross, M. H. (2003). Optimized spatial hashing for collision detection of deformable objects. In *Vmv*, volume 3, pages 47–54.
- The CGAL Project (2020). *CGAL User and Reference Manual*. CGAL Editorial Board, 5.0.3 edition.
- Thiery, J.-M., Tierny, J., and Boubekeur, T. (2012). Cager: Cage-based reverse engineering of animated 3d shapes. *Comput. Graph. Forum*, 31(8):2303–2316.
- Tong, W.-h. and Kim, T.-w. (2009). High-order approximation of implicit surfaces by g1 triangular spline surfaces. *Computer-Aided Design*, 41(6):441–455.
- Toulorge, T., Geuzaine, C., Remacle, J.-F., and Lambrechts, J. (2013). Robust untangling of curvilinear meshes. *Journal of Computational Physics*, 254:8 – 26.
- Toulorge, T., Lambrechts, J., and Remacle, J.-F. (2016). Optimizing the geometrical accuracy of curvilinear meshes. *Journal of Computational Physics*, 310:361 – 380.

- Tournois, J., Wormser, C., Alliez, P., and Desbrun, M. (2009a). Interleaving delaunay refinement and optimization for practical isotropic tetrahedron mesh generation. *ACM Transactions on Graphics*, 28(3):Art–No.
- Tournois, J., Wormser, C., Alliez, P., and Desbrun, M. (2009b). Interleaving delaunay refinement and optimization for practical isotropic tetrahedron mesh generation. *ACM Transactions on Graphics*, 28(3):1.
- Tozoni, D. C., Dumas, J., Jiang, Z., Panetta, J., Panozzo, D., and Zorin, D. (2020). A low-parametric rhombic microstructure family for irregular lattices. *ACM Transactions on Graphics (TOG)*, 39(4):101–1.
- Turner, M. (2017). *High-order mesh generation for CFD solvers*. PhD thesis, Imperial College London.
- Turner, M., Peiró, J., and Moxey, D. (2016). A variational framework for high-order mesh generation. *Procedia Engineering*, 163:340 – 352. 25th International Meshing Roundtable.
- Tutte, W. T. (1963). How to draw a graph. *Proceedings of the London Mathematical Society*, 13(3):743–768.
- Villard, J. and Borouchaki, H. (2002). Adaptive meshing for cloth animation. In *In Proceedings of the 11th International Meshing Roundtable (IMR 2002)*, pages 243–252.
- Wang, B., Ferguson, Z., Schneider, T., Jiang, X., Attene, M., and Panozzo, D. (2021). A large-scale benchmark and an inclusion-based algorithm for continuous collision detection. *ACM Trans. Graph.*, 40(5).
- Wang, B., Schneider, T., Hu, Y., Attene, M., and Panozzo, D. (2020). Exact and efficient polyhedral envelope containment check. *ACM Trans. Graph.*, 39(4).
- Wang, L., Wang, X., Tong, X., Lin, S., Hu, S., Guo, B., and Shum, H.-Y. (2003). View-dependent displacement mapping. In *ACM Transactions on graphics (TOG)*, volume 22, pages 334–339. ACM.
- Wang, R., Liu, L., Yang, Z., Wang, K., Shan, W., Deng, J., and Chen, F. (2016). Construction of manifolds via compatible sparse representations. *ACM Transactions on Graphics (TOG)*, 35(2):1–10.
- Wang, X., Tong, X., Lin, S., Hu, S., Guo, B., and Shum, H.-Y. (2004). Generalized displacement maps. In *Proceedings of the Fifteenth Eurographics conference on Rendering Techniques*, pages 227–233. Eurographics Association.
- Weber, O. and Zorin, D. (2014). Locally injective parametrization with arbitrary fixed boundaries. *ACM Trans. Graph.*, 33(4):75:1–75:12.

- Wicke, M., Ritchie, D., Klingner, B. M., Burke, S., Shewchuk, J. R., and O'Brien, J. F. (2010). Dynamic local remeshing for elastoplastic simulation. *ACM Trans. Graph.*, 29(4).
- Wojtan, C. and Turk, G. (2008). Fast viscoelastic behavior with thin features. *ACM Trans. Graph.*, 27(3):1–8.
- Xia, S. and Qian, X. (2017). Isogeometric analysis with bézier tetrahedra. *Computer Methods in Applied Mechanics and Engineering*, 316:782 – 816. Special Issue on Isogeometric Analysis: Progress and Challenges.
- Xie, Z. Q., Sevilla, R., Hassan, O., and Morgan, K. (2013). The generation of arbitrary order curved meshes for 3d finite element analysis. *Computational Mechanics*, 51(3):361–374.
- Ye, K., Ni, W., Krieger, M., Ma'ayan, D., Wise, J., Aldrich, J., Sunshine, J., and Crane, K. (2020). Penrose: from mathematical notation to beautiful diagrams. *ACM Transactions on Graphics (TOG)*, 39(4):144–1.
- Yerry, M. A. and Shephard, M. S. (1983). A modified quadtree approach to finite element mesh generation. *IEEE Computer Graphics and Applications*, 3(1):39–46.
- Yu, Y., Wei, X., Li, A., Liu, J. G., He, J., and Zhang, Y. J. (2020). Hexgen and hex2spline: Polycube-based hexahedral mesh generation and spline modeling for isogeometric analysis applications in ls-dyna.
- Yvert, A., Hahmann, S., and Bonneau, G.-P. (2005a). Hierarchical triangular splines. *ACM Trans. Graph.*, 24(4):1374–1391.
- Yvert, A., Hahmann, S., and Bonneau, G.-P. (2005b). Smooth adaptive fitting of 3d models using hierarchical triangular splines. In *International Conference on Shape Modeling and Applications 2005 (SMI'05)*, pages 13–22. IEEE.
- Zayer, R., Steinberger, M., and Seidel, H.-P. (2017). A gpu-adapted structure for unstructured grids. *Comput. Graph. Forum*, 36(2):495–507.
- Zhang, E., Mischaikow, K., and Turk, G. (2005). Feature-based surface parameterization and texture mapping. *ACM Trans. Graph.*, 24(1):1–27.
- Zhang, P., Vekhter, J., Chien, E., Bommes, D., Vouga, E., and Solomon, J. (2020). Octahedral frames for feature-aligned cross fields. *ACM Trans. Graph.*, 39(3).
- Zhang, S., Li, Z., Zhang, H., and Yong, J. (2011). Multi-resolution mesh fitting by b-spline surfaces for reverse engineering. In *2011 12th International Conference on Computer-Aided Design and Computer Graphics*, pages 251–257.
- Zhou, K., Synder, J., Guo, B., and Shum, H.-Y. (2004). Iso-charts: Stretch-driven mesh parameterization using spectral analysis. In *Proceedings of the 2004 Eurographics/ACM SIGGRAPH Symposium on Geometry Processing*, SGP '04, pages 45–54, New York, NY, USA. ACM.

- Zhou, M., Xie, T., Seol, S., Shephard, M. S., Sahni, O., and Jansen, K. E. (2012). Tools to support mesh adaptation on massively parallel computers. *Engineering with Computers*, 28(3):287–301.
- Zhou, Q., Grinspun, E., Zorin, D., and Jacobson, A. (2016). Mesh arrangements for solid geometry. *ACM Transactions on Graphics (TOG)*, 35(4):1–15.
- Zhou, Q. and Jacobson, A. (2016). Thingi10k: A dataset of 10,000 3d-printing models. *arXiv preprint arXiv:1605.04797*.
- Ziel, V., Bériot, H., Atak, O., and Gabard, G. (2017). Comparison of 2d boundary curving methods with modal shape functions and a piecewise linear target mesh. *Procedia Engineering*, 203:91 – 101. 26th International Meshing Roundtable, IMR26, 18-21 September 2017, Barcelona, Spain.
- Zomorodian, A. and Edelsbrunner, H. (2000). Fast software for box intersections. In *Proceedings of the sixteenth annual symposium on Computational geometry*, pages 129–138.
- Zorin, D. (2000). Subdivision for modeling and animation. *SIGGRAPH2000 course notes*.