

CS2102: Database Systems

Lecture 4 — Entity Relationship Model (ER Model)

Errata

- Notation for renaming in RA expressions
 - Lecture

$$\rho_{\text{new}\leftarrow\text{old}}$$

■ Tutorial 2 (solutions)

$$\rho_{\text{old}\leftarrow\text{new}}$$

Quick Recap: SQL for Creating Databases

- Data Definition Language (DDL)
 - Create, modify and drop tables to implement a given DB schema
 - Specify integrity constraints (e.g., NOT NULL, PRIMARY KEY, FOREIGN KEY, CHECK)

- Data Manipulation Language (DML)
 - Insert, update and delete data from tables

Employees (id: integer, name: text, age: integer, role: text) **CREATE TABLE** Employees (INTEGER PRIMARY KEY. VARCHAR(50) NOT NULL, INTEGER. age VARCHAR(50) role **Employees** name age role **INSERT INTO Employees VALUES** (101, 'Sarah', 25, 'dev') (102, 'Judy', 35, 'sales');

Employees

id	name	age	role
101	Sarah	25	dev
102	Judy	35	sales

We Sneakily Skipped a Step

Open questions:

- Where does the database schema come from?
- What tables with which attributes do we need?
- What data integrity constraints are required?
- Table names, attribute names, data types, ...?

→ Database Design Process

Solution

- Storing the "dob" instead of "age" is arguably the preferred approach
- The value of "age" changes each year (no really a big deal)
- "dob" provides more detailed information compared the "age"

Which table is "better"?

```
CREATE TABLE Employees (

id INTEGER PRIMARY KEY,
name VARCHAR(50) NOT NULL,
age INTEGER,
role VARCHAR(50)
);
```

or

```
CREATE TABLE Employees (
id INTEGER PRIMARY KEY,
name VARCHAR(50) NOT NULL,
dob DATE,
role VARCHAR(100),
phone INTEGER
);
```

Database Design Process — 6 Common Steps

Requirement Analysis

- Identification and collection of user needs
- e.g., data /application / performance requirements

Conceptual DB Design

- Capturing requirements using a conceptual model
- RDBMS: Entity Relationship Model (ER Model)

Logical DB Design

- Mapping conceptual model to logical schema of DBMS
- RDBMS: Entity Relationship Model → Relational Schema

Schema Refinement

Checking schema / tables for redundancies and anomalies

Physical DB Design

- Implementing database based on final data schema
- Consideration of performance requirements

Security Design

 Identification users and user groups and their permissions to access which parts of the data

Overview

Entity Relationship Model

- Overview + ER diagrams
- Entity sets and attributes
- Relationship sets
- Cardinality & participation constraints

Relational Mapping

- From ER diagram to database tables
- Extended notations for ER diagrams
 - ISA hierarchies: generalization/specialization
 - Aggregation

Requirement Analysis: Online Airline Reservation System (OARS)

Users need to be able to make bookings from an origin to a destination airport which may comprise multiple connecting flights. Each flight has a flight number, the origin and destination airport, the distance in kilometers, the departure and arrival time, and the days of the week the flight is in operation.

A flight instance is the actual scheduled flight on a given day together with the assigned aircraft type. For example, flight SQ231 flies daily from Singapore to Sydney, typically with a Boeing 777-300ER (code: B77W).

For a valid booking, we need the user's name, sex, address, phone number(s), and the passport number. Users are only able to pay via credit card. When making a booking, the user can select the class, the seat number, as well as meal preferences (if available).

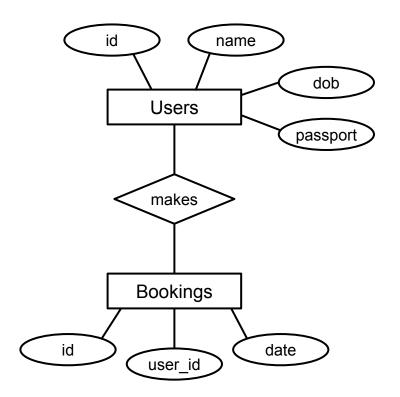
Entity Relationship Model

ER Model

- Most common model for conceptual database design
- Developed by Peter Chen (1976)
- Visualized using ER diagrams

Core concepts

- All data is described in terms of entities and their relationships
- Information about entities & relationships are described using attributes
- Certain data constraints can be described using additional annotations



Entities and Entity Sets

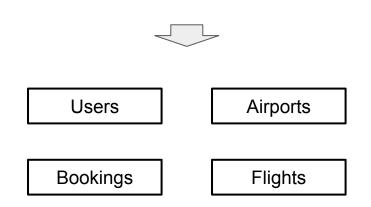
Entity

 Real-world things or objects that are distinguishable from other objects (e.g., an individual user, airport, flight, or booking)

Entity Set

- Collection of entities of the same type
- Represented by rectangles in ER diagrams
- Names are typically nouns

Users need to be able to make bookings from an origin to a destination airport which may comprise multiple connecting flights. Each flight has a flight number, [...]

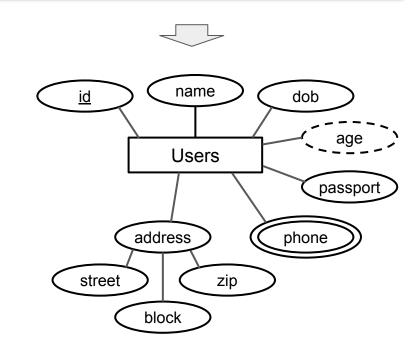


Attributes

- Attribute:
 - specific information describing an entity
 - represented by an oval in ER diagrams
- 4 subtypes of attributes
 - **Key attribute(s)**: uniquely identifies each entity (oval with the attribute name(s) underlines)
 - Composite attribute: composed of multiple other attributes (oval comprising of ovals)
 - Multivalued attribute: may consisting more than one value for a given entity (double-lined oval)
 - **Derived attribute**: derived from other attributes (dashed oval)

For a valid booking, we need the user's name, sex, address, phone number(s), and the passport number.

Users are only able to pay via credit card. [...]



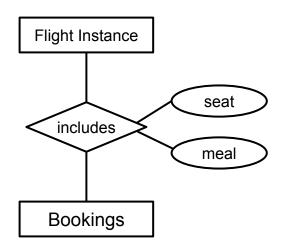
Relationships and Relationship Sets

Relationship

Association among two or more entities

Relationship Set

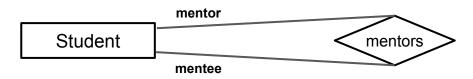
- Collection of relationships of the same type
- Represented by diamonds in ER diagrams
- Can have their own attributes that further describe the relationship
- Names are typically verbs



- Additional annotations to further specify relationships
 - Roles, degree, cardinalities, participation, dependencies

Relationship Roles

- Role
 - Descriptor of an entity set's participation in a relationship
 - Most of the time implicitly given by the name of the entity sets
 - Explicit role label only common in case of ambiguities (typically in case the same entity sets participates in the same relationship more than once)
- Example: Students can mentor other students

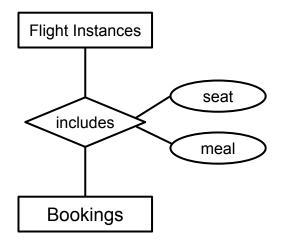


Degree of Relationship Sets

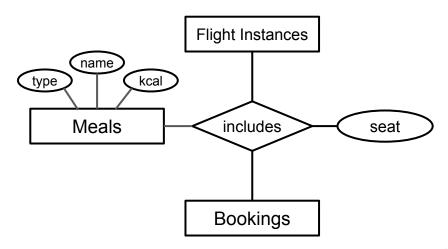
Degree

- In principle, no limitation how many entity roles participate in a relationship
- An *n*-ary relationship set involves *n* entity roles $\rightarrow n$ = degree of relationship set

 $n = 2 \rightarrow binary relationship set$

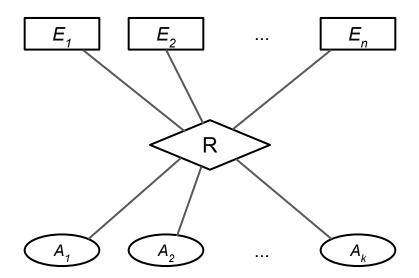


n = 3 → ternary relationship set



Degree of Relationship Sets

- General n-ary relationship set R
 - *n* participating entity sets $E_1, E_2, ..., E_n$
 - k relationship attributes $A_1, A_2, ..., A_k$



"In typical modeling, binary relationships are the most common and relationships with n>3 are very rare" - Peter Chen (2009)

Overview

Entity Relationship Model

- Overview + ER diagrams
- Entity sets and attributes
- Relationship sets
- Cardinality & participation constraints

Relational Mapping

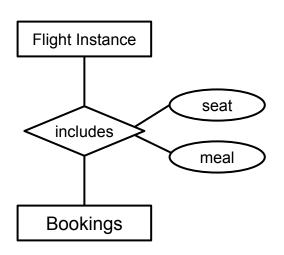
- From ER diagram to database tables
- Extended notations for ER diagrams
 - ISA hierarchies: generalization/specialization
 - Aggregation

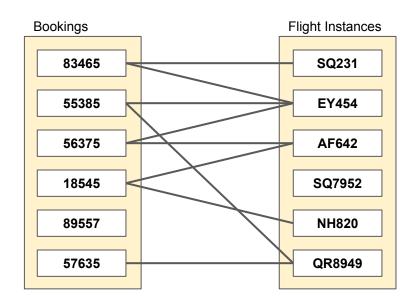
Cardinality Constraints

- Cardinalities of Relationship Sets
 - Describe how often an entity can participate in a relationship <u>at most</u>
- 3 basic cardinality constraints
 - Many-to-many (e.g., a flight can be performed by different aircrafts; an aircraft can perform different flights)
 - Many-to-one (e.g., a user can make many bookings, but each booking is done by one user)
 - One-to-one (e.g., a user is associated with one set of credit card details, and vice versa)
- Cardinality constraints can be specified using annotations in ER diagram
 - Note: different ways to specify cardinality constraints available

Many-to-Many

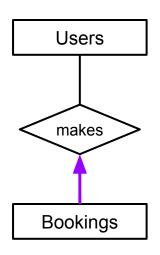
- Many-to-many relationship between bookings and flight instances
 - Each booking can include 0 or more flight instances (note that a booking with 0 flights might not meaningful; we will improve on that)
 - Each flight instance can be part of 0 or more bookings

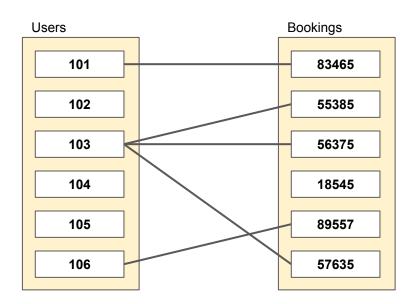




Many-to-One

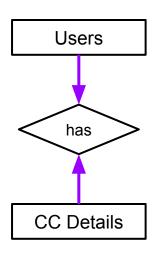
- Many-to-one relationship between users and bookings
 - Each user can make 0 or more bookings
 - Each booking is done by one 1 user at most (again, not perfect yet, and we will improve on that)

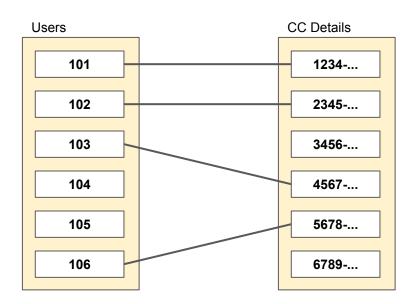




One-to-One

- One-to-one relationship between users and credit card details
 - Each user can provide only 1 set of credit card details at most
 - Each set of credit card details is associated with 1 user at most





Participation Constraints

- Limitation of (basic) cardinality constraints from previous examples
 - A booking can include 0 flights
 - A booking can be done by 0 users
 - A set of credit card details does not need to be associated with a user

an entity does not have to participate in a relation

→ Cardinality constraints (many-to-many, many-to-one, one-to-one) only specify some kind of "upper bound"

→ Participation constraints

- Is the participation of an entity in a relationship mandatory?
- Allow to specify a trivial lower bound

Participation Constraints

- Partial participation constraint (default)
 - Participation of an entity in a relationship is not mandatory
 - Example: A user made 0 or more bookings

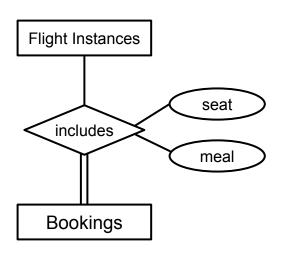


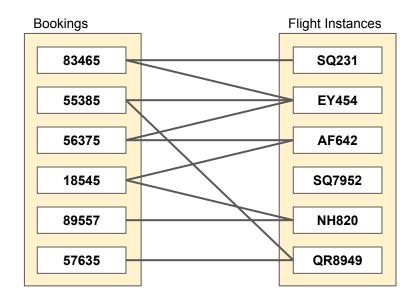
- Total participation constraint
 - Participation of an entity in a relationship is mandatory
 - Example: We only keep user that made at least one booking



Cardinality & Participation Constraints

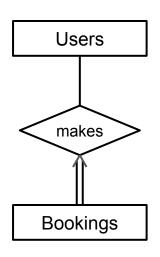
- Many-to-many relationship between bookings and flight instances
 - Each booking includes 1 or more flight instances
 - Each flight instance can be part of 0 or more bookings

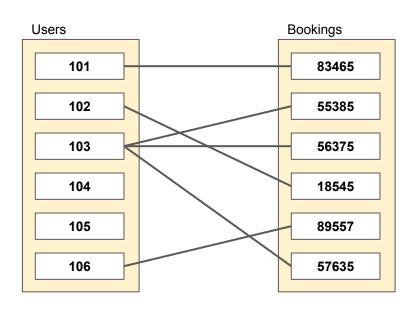




Cardinality & Participation Constraints

- Many-to-one relationship between users and bookings
 - Each user can make 0 or more bookings
 - Each booking is done by <u>exactly 1</u> user





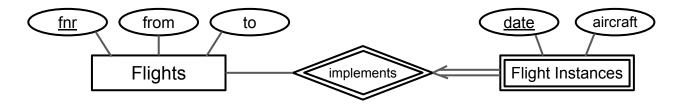
Dependency Constraints

Weak entity sets

- Entity set that does not have its own key
- A weak entity can only be uniquely identify by considering the primary key of the owner entity
- A weak entity's existence depends on the existence of its owner entity
- Weak entity set and identifying relation set are represented via double-lined rectangles / diamonds

Example

■ A flight instance is the actual scheduled flight (with a unique flight number) on a given day



Dependency Constraints

Requirements

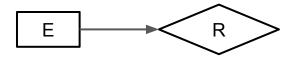
- Many-to-one relationship (identifying relationship) from weak entity set to owner entity set
- Weak entity set must have total participation in identifying relationship

Partial key

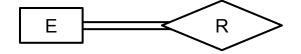
- Set of attributes of weak entity set that uniquely identifies a weak entity for a given owner entity
- Example: Given a flight (e.g. SQ231), the date identifies the exact instance of that flight



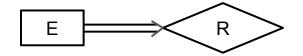
Summary of Participation Constraints



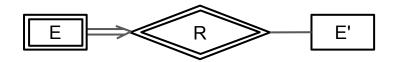
Each instance of E participates in at most one instance of R.



Each instance of E participates in at least one instance of R.



Each instance of E participates in exactly one instance of R.



E is a weak entity set with identifying owner E' and identifying relationship set R.

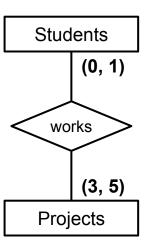
Alternative Representations (Cardinality Constraints)

Many-to-One One-to-One Many-to-Many Flight Instance Flight Instance Users Users Users Users m m includes includes has has makes makes n **Bookings Bookings CC Details CC** Details **Bookings** Bookings

Alternative Representations (Cardinality Constraints)

Min/Max notation

Specification of precise lower and upper bounds



A student works on exactly 1 project, or no project at all.

A project is assigned to teams comprising 3 to 5 students.

Quick Quiz: Why is this more precise notation in practice often not that useful?

Solution

- With what we've learned so far, we could not enforce these constraints
- At least not the (3, 5) constraint

Overview

Entity Relationship Model

- Overview + ER diagrams
- Entity sets and attributes
- Relationship sets
- Cardinality & participation constraints

Relational Mapping

- From ER diagram to database tables
- Extended notations for ER diagrams
 - ISA hierarchies: generalization/specialization
 - Aggregation

Database Design Process — 6 Common Steps

Requirement Analysis

Identification and collection of user needs

• e.g., data /application / performance requirements

Conceptual DB Design

- Capturing requirements using a conceptual model
- RDBMS: Entity Relationship Model (ER Model)

Logical DB Design

- Mapping conceptual model to logical schema of DBMS
- RDBMS: Entity Relationship Model → Relational Schema

Schema Refinement

• Checking schema / tables for redundancies and anomalies

Physical DB Design

Implementing database based on final data schema

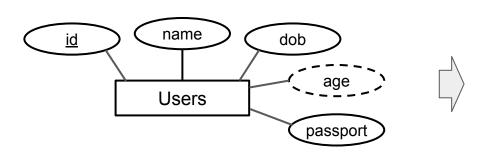
Consideration of performance requirements

Security Design

 Identification users and user groups and their permissions to access which parts of the data

Entity Sets

- Straightforward mapping from entity sets to tables (except composite & multivalued attributes)
 - Name of entity set → name of table
 - Attributes of entity set → attributes of table
 - Key attributes of entity set → primary key of table



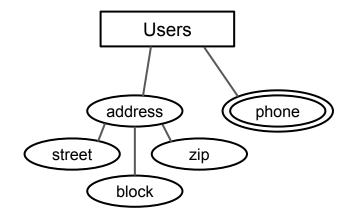
```
id INTEGER,
name VARCHAR(100),
dob DATE,
age INTEGER,
passport VARCHAR(20),
PRIMARY KEY (id)
);
```

Note: The ER diagram does not specify UNIQUE or NOT NULL constraints that are potentially meaningful when creating a table.

Composite & Multivalued Attributes

 Problem: Tables can only hold atomic values (ignoring complex data types support by some DBMS)

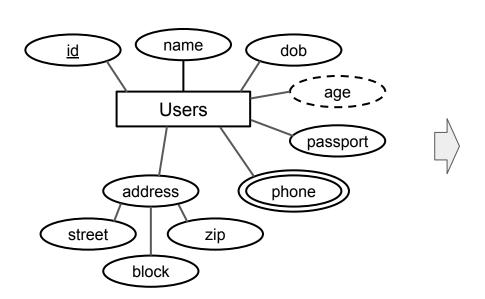
- 2 principle solutions + 1 alternative
 - Convert composite & multivalued attributes into a set of single-valued attributes
 - Create additional tables with a foreign key constraint referencing table of original entity set (typically only meaningful for multivalued attributes)
 - Convert composite & multivalued attributes to one single-valued attribute (if meaningful)



Note: One can design the ER diagram without composite and multivalued attributes using additional entity and relationship set which yield the same result as the proposed solutions.

Composite & Multivalued Attributes

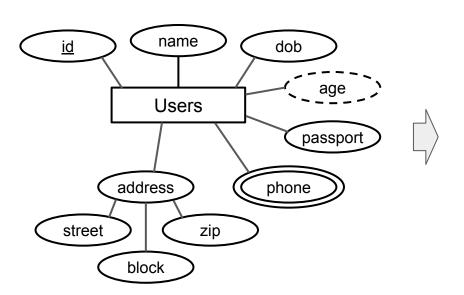
- Conversion to single-valued attributes
 - Requires an upper limit in case of multivalued attributes



```
CREATE TABLE Users (
     id
                INTEGER.
                VARCHAR(100),
     name
     dob
                date.
                INTEGER,
     age
                VARCHAR(20),
     passport
     street
                VARCHAR(50),
                VARCHAR(6),
     block
     zip
                INTEGER.
     phone1
                INTEGER,
     phone2
                INTEGER,
     phone3
                INTEGER.
     PRIMARY KEY (id)
```

(Composite &) Multivalued Attributes

Additional table with foreign key constraint

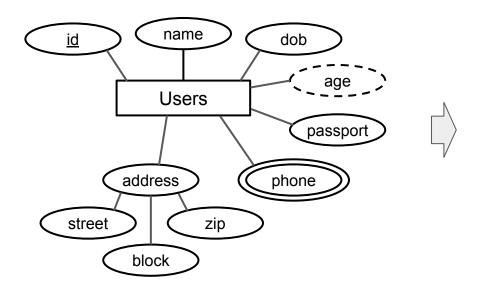


```
CREATE TABLE Users (
     id
                INTEGER,
                 VARCHAR(100),
     name
     PRIMARY KEY (id)
);
CREATE TABLE PhoneNumbers (
                 INTEGER.
     user id
     phone
                 INTEGER,
     FOREIGN KEY (user id) REFERENCES Users (id)
```

Composite & Multivalued Attributes

Quick Quiz: What are the problems with this approach and when is it meaningful (enough)?

Convert to single-valued attribute



```
CREATE TABLE Users (
     id
                INTEGER,
                VARCHAR(100),
     name
     dob
                date.
                INTEGER,
     age
                VARCHAR(20),
     passport
                VARCHAR(200),
     address
                VARCHAR(200),
     phone
     PRIMARY KEY (id)
```

id	name	dob	age	passport	address	phone
101	Alice	15-02-2000	21	KEJR4A90	15 Computing Drive, Singapore 117418	65-1111-2222, 65-2222-3333, 65-3333-4444

Quick Quiz

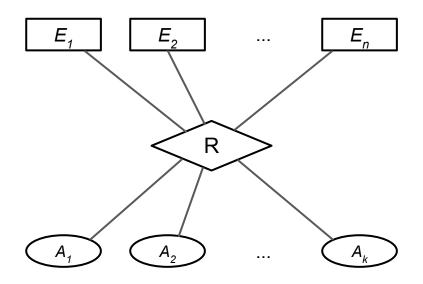
Solution

- Modeling "address" and "phone" as a single-values string might be OK-ish
 if we never use these attributes to select rows
- If we only need to get the address or all phone numbers for a given user the this solution might be good enough
- However, queries using "address" or "phone" to filter rows will become unnecessarily complicated or even impossible
- A query such as "Return all users with addresses with the ZIP code 123456" are possible since SQL supports string pattern matching and even regular expression. The performance would degrade, though.
- More intricate queries might still be formulated but the complexity of the SQL query would quickly blow up

Relationship Sets

General n-ary relationship set R

- *n* participating entity sets $E_1, E_2, ..., E_n$
- k relationship attributes $A_1, A_2, ..., A_k$
- Let $Key(E_i)$ be the attributes of the selected key of entity set E_i

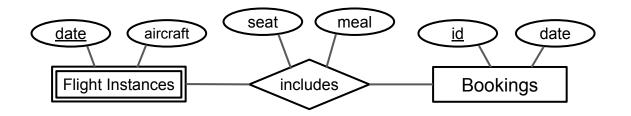


→ Attributes of relationship set R

- $Key(E_1)$, $Key(E_2)$, ..., $Key(E_n)$ key attributes of all participating entity sets E_i
- $A_1, A_2, ..., A_k$ all relationship attributes of R

Cardinality Constraints: Many-to-Many

Quick Quiz: Where does "flight_nr" come from?

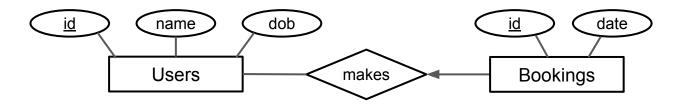


Quick Quiz

Solution

- "Flight Instances" is weak entity set with "Flights" being the owner entity set
- Thus, "Flight Instances" is identified by the key of "Flights" (i.e., "fnr") and its own partial key "date"

Cardinality Constraints: Many-to-One

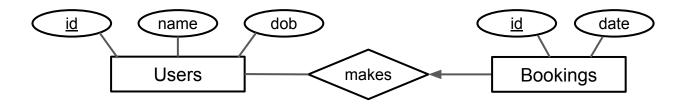


- Approach 1: Represent "makes" with a separate table
 - Similar to Many-to-Many but with different primary key!

```
CREATE TABLE Makes (
    user_id INTEGER,
    booking_id INTEGER,
    PRIMARY KEY (booking_id),
    FOREIGN KEY (user_id) REFERENCES Users (id),
    FOREIGN KEY (booking_id) REFERENCES Bookings (id)
);
```

Cardinality Constraints: Many-to-One

Quick Quiz: Which is generally the preferred approach?



- Approach 2: Combine "makes" and "Bookings" into one table
 - Possible because given a booking, we can uniquely identify the user who made it

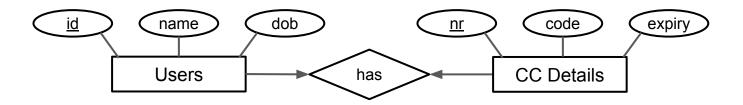
```
CREATE TABLE Bookings (
    id INTEGER,
    date DATE,
    user_id INTEGER,
    PRIMARY KEY (id),
    FOREIGN KEY (user_id) REFERENCES Users (id)
);
```

Quick Quiz

Solution

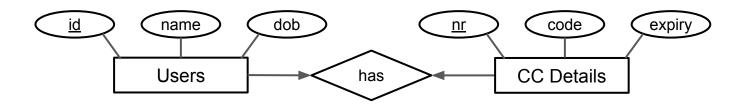
- Approach 2 is generally the preferred approach as it leads to a smaller number of table
- Less tables also means that queries might need less join operations (which are typically the more expensive operations).

Cardinality Constraints: One-to-One



- Approach 1: Represent "has" with a separate table
 - Similar to Many-to-One but primary key can be chosen

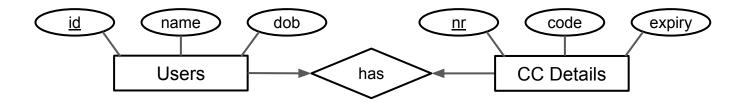
Cardinality Constraints: One-to-One



Approach 2: Combine "has" and "Users" or "has" and "CC Details"

```
id INTEGER,
name VARCHAR(100),
dob DATE,
cc_nr CHAR(16),
PRIMARY KEY (id),
FOREIGN KEY (cc_nr) REFERENCES CCDetails (nr)
);
```

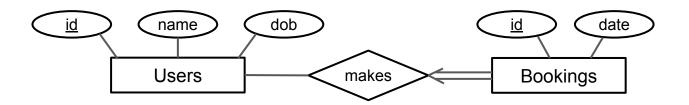
Cardinality Constraints: One-to-One



Approach 3: Combine "has", "Users", and "CC Details"

```
id INTEGER,
name VARCHAR(100),
dob DATE,
cc_nr CHAR(16) UNIQUE,
cc_code CHAR(3),
cc_expiry DATE,
PRIMARY KEY (id)
);
```

Cardinality & Participation Constraints

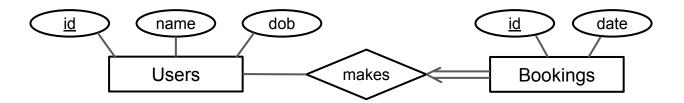


• Approach 1: Represent "makes" with a separate table

```
CREATE TABLE Makes (
    user_id INTEGER NOT NULL,
    booking_id INTEGER,
    PRIMARY KEY (booking_id),
    FOREIGN KEY (user_id) REFERENCES Users (id),
    FOREIGN KEY (booking_id) REFERENCES Bookings (id)
);
```

- Schema does <u>not</u> enforce total participation of "Bookings" w.r.t. "Makes"
- e.g.: "Makes" can be empty while both "Users" and "Bookings" are non-empty

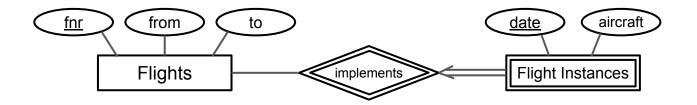
Cardinality & Participation Constraints



- Approach 2: Combine "makes" and "Bookings" into one table
 - Enforces total participation via NOT NULL constraint

```
CREATE TABLE Bookings (
    id INTEGER,
    date DATE,
    user_id INTEGER NOT NULL,
    PRIMARY KEY (id),
    FOREIGN KEY (user_id) REFERENCES Users (id)
);
```

Weak Entity Sets



```
fnr VARCHAR(10),
from VARCHAR(10),
to VARCHAR(10),
PRIMARY KEY (fnr)
);
```

```
CREATE TABLE FlightInstances (
fnr VARCHAR(10),
date DATE,
aircraft VARCHAR(10),
PRIMARY KEY (fnr, date),
FOREIGN KEY (fnr) REFERENCES Flights (fnr)
ON DELETE CASCADE
);
```

ER Design & Relational Mapping — Basic Guidelines

- Guidelines for ER design
 - An ER diagram should capture as many of the constraints as possible
 - An ER diagram must not impose any constraints that are not required
- Guidelines for relational mapping

(i.e., from ER diagram to relational database schema)

- The relational schema should enforce as many if the constraints as possible using column and/or table constraints
- The relational schema should not impose and constraints that are not required

Overview

Entity Relationship Model

- Overview + ER diagrams
- Entity sets and attributes
- Relationship sets
- Cardinality & participation constraints

Relational Mapping

■ From ER diagram to database tables

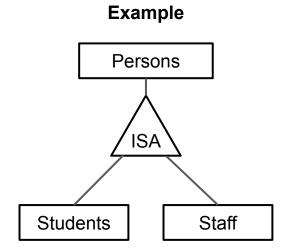
Extended notations for ER diagrams

- ISA hierarchies: generalization/specialization
- Aggregation

Extended Concepts — ISA Hierarchies

- ISA hierarchies
 - Special type of relationship: "is a"
 - Used to model generalization/specialization of entity sets

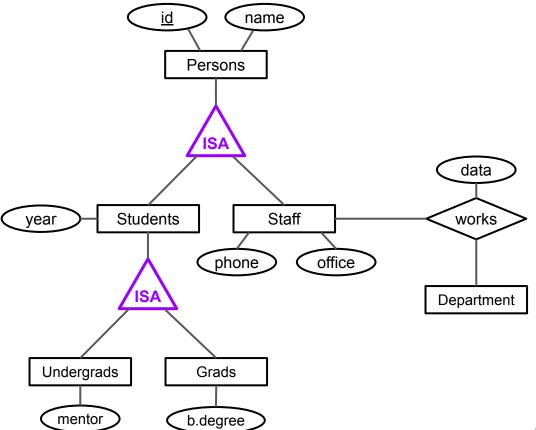
Subclass₁ Subclass₁ ... Subclass₁



ISA Hierarchies

Interpretation

- Every entity in a subclass is an entity in its superclass
- Each subclass has specific attributes and/or relationships



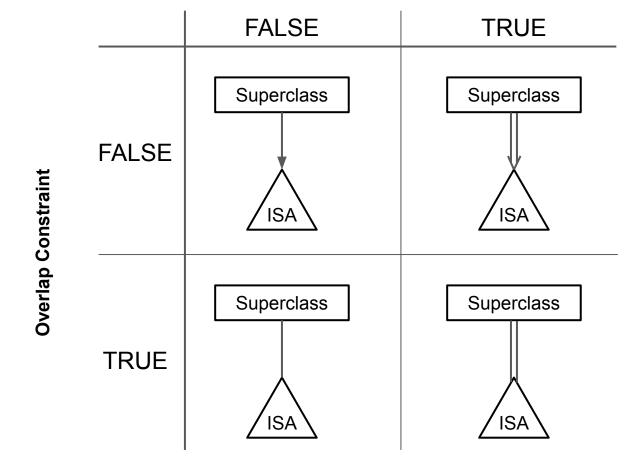
ISA Hierarchies — Constraints

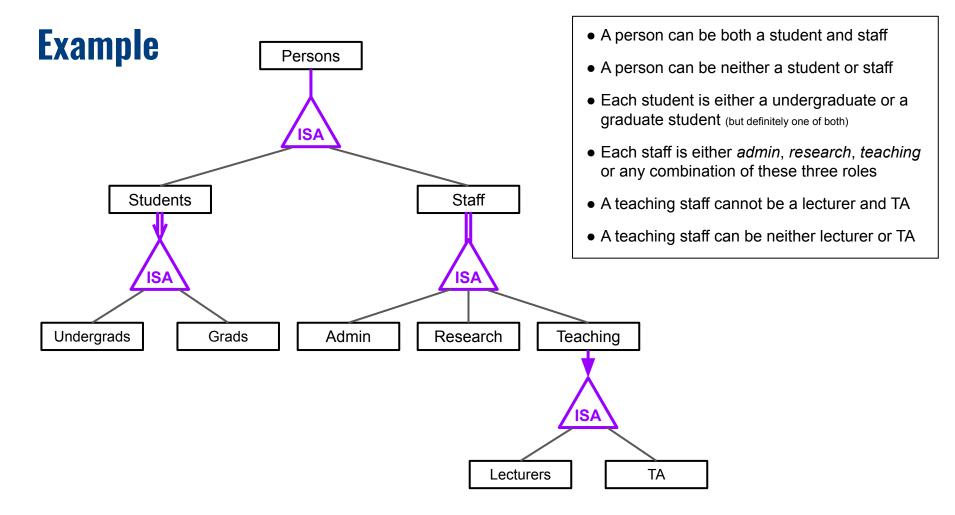
- Overlap constraint: Can a superclass entity belong to multiple subclasses?
 - TRUE → a superclass entity can belong to multiple subclasses (e.g., a person can be both student and staff)
 - FALSE → otherwise (e.g., a student is either a graduate or undergraduate)
- Covering constraint: Does a superclass entity have to belong to a subclass?
 - TRUE → every superclass entity has to belong to a subclass (e.g., there is no student that is neither a graduate or undergraduate)
 - FALSE → otherwise (e.g., not every person is a student or staff)

ISA Hierarchies

Covering Constraint

Notation in ER Diagram

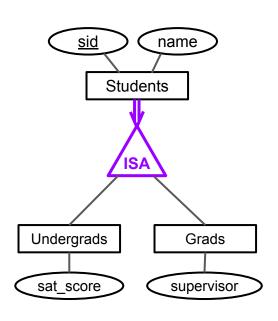




ISA Hierarchies: Relational Mapping

Basic approach: One relation per subclass and superclass

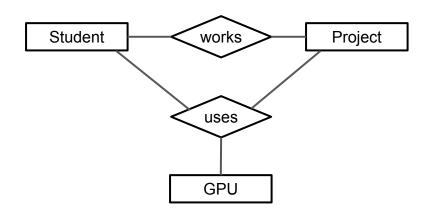
```
CREATE TABLE Students (
     sid
                CHAR(20) PRIMARY KEY,
                VARCHAR(50)
     name
CREATE TABLE Undergrads (
     sid
                CHAR(20) PRIMARY KEY,
                NUMERIC,
     sat score
     FOREIGN KEY (sid) REFERENCES Students (sid) ON DELETE CASCADE
);
CREATE TABLE Grads (
     sid
                CHAR(20) PRIMARY KEY,
                CHAR(8),
     supervisor
     FOREIGN KEY (sid) REFERENCES Students (sid) ON DELETE CASCADE,
     FOREIGN KEY (supervisor) REFERENCES Staff (id) ON DELETE SET NULL
```



Extended Concepts — Aggregation

- Concepts of ER diagrams so far
 - Only relationships between entity sets
 - No relationships between entity sets and relationship sets

Motivational example



Limitations:

- Relationship between "works" and "uses" not explicitly captured
- "works" and "uses" are kind of redundant relationships
- **→** Aggregation

Extended Concepts — Aggregation

- Aggregation basic idea
 - Abstraction that treats relationships as higher-level entities
 - Example: treat Student-works-Project as an entity set
- Notation in ER diagram (2 equivalent alternatives)

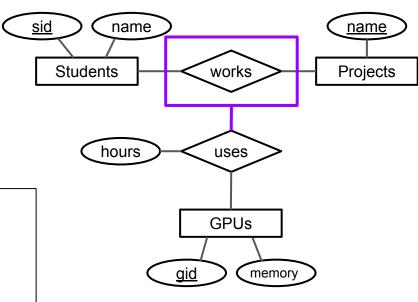


Aggregation — **Relational Mapping**

Schema definition of "uses"

- Primary key of aggregation relationship → (sid, pname)
- Primary key of associated entity set "GPUs" → gid
- Descriptive attributes of "uses" → hours

```
CREATE TABLE Uses (
gid INTEGER,
sid CHAR(20),
pname VARCHAR(50),
hours NUMERIC,
PRIMARY KEY (gid, sid, pname),
FOREIGN KEY (gid) REFERENCES GPUs (gid),
FOREIGN KEY (sid, pname) REFERENCES works (sid, pname)
);
```



Summary

- Entity-Relationship (ER) model
 - Basic concepts: entity sets, relationship sets, attributes
 - Cardinality constraints and participation constraints
 - Extended concepts: ISA hierarchies, aggregation

Visualized using ER diagrams

- Relational Mapping
 - Mapping ER diagram to database schema
 - Not all constraints of ER diagram may be captured
- Outlook for next lecture
 - SQL for querying a database (recommendation: study RA)