CS3230: Assignment for Week 11 Solutions

Due: Sunday, 17th Apr 2022, 11:59 pm SGT.

- 1. (a) Yes. This would be true if P = NP.
 - (b) Yes. This would be true if $P \neq NP$.
 - (c) No. If one NP-complete problem can be solved in polynomial time, then all of them can be solved in polynomial time, because they can be reduced to each other in polynomial time.
- 2. First, the problem 2022-LABEL belongs to NP: a certificate for a YES-instance is a labeling 1, 2, ..., 2022 of the graph such that no two adjacent vertices receive the same label. A verifier can check in polynomial time that each edge connects two vertices with different labels.

Next, to show that 2022-Label is NP-hard, we reduce from 2021-Label. Given an instance G of 2021-Label, as in the hint, we add an extra vertex v that is connected to all existing vertices to obtain an instance G' of 2022-Label. Clearly, this reduction can be done in polynomial time.

If G is a YES-instance of 2021-LABEL, then by labelling v with 2022, we get a valid labeling of G' with $1, 2, \ldots, 2022$. For the opposite direction, suppose G' is a YES-instance of 2022-LABEL. Assume without loss of generality that v is labelled 2022. Since v is connected to all remaining vertices, no other vertex can be labelled 2022. Hence, all other vertices are labelled $1, 2, \ldots, 2021$, which gives a valid labeling of G.

3. First, the problem 3-SATTWICE belongs to NP: a certificate for a YES-instance consists of two satisfying assignments, which the verifier can check in polynomial time.

Next, to show that 3-SATTWICE is NP-hard, we reduce from 3-SAT. Given an instance ϕ of 3-SAT, we add a clause $(x \lor x \lor \overline{x})$, where x is a new variable, to obtain an instance ϕ' of 3-SATTWICE. Clearly, this reduction can be done in polynomial time.

If ϕ is a YES-instance of 3-SAT, then by setting x to either True or False and using the same satisfying assignment of ϕ for the remaining variables, we get two satisfying assignments for ϕ' . For the opposite direction, a satisfying assignment for ϕ' immediately yields one for ϕ by ignoring x.