

# *Design and Analysis of Algorithms*



*Algorithms*

CS3230

GR3330

## **Tutorial**

Week 3

# Question 1



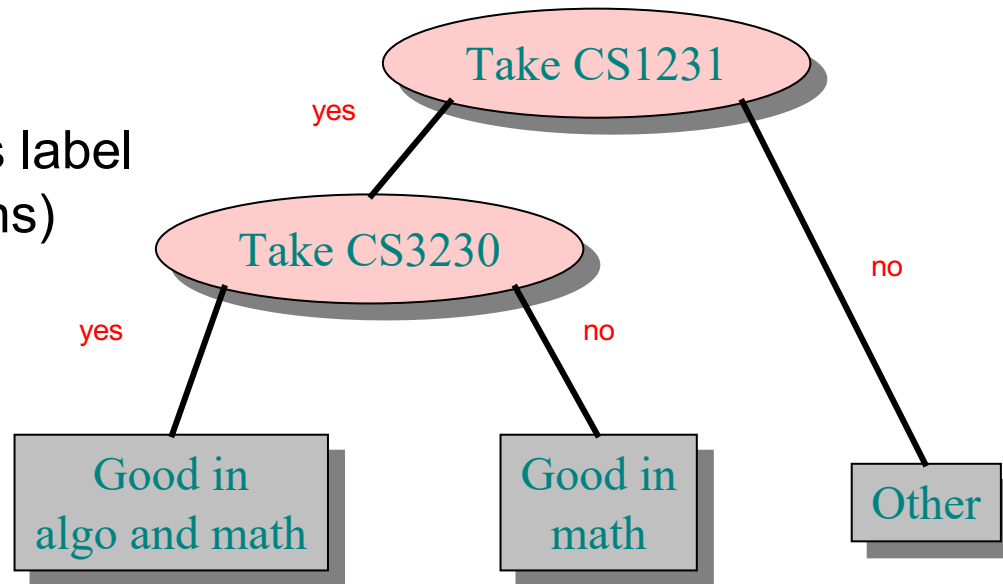
- Given an **unsorted array** of  $n$  real numbers  $A[1..n]$  and a query number  $x$ . You need to develop a function **search**( $x, A$ ) which returns an integer  $i$  if  $A[i]=x$ ; and returns  $-1$  otherwise.
- We have two assumptions:
  - Assume comparison model
  - Assume each comparison returns  $<$ , or  $>$ , or  $=$  between  $x$  and an element of  $A$ .
- What is the lower bound of the number of comparisons?

A.  $n$       B.  $\lfloor \lg n \rfloor + 1$       C.  $n \lfloor \lg n \rfloor$

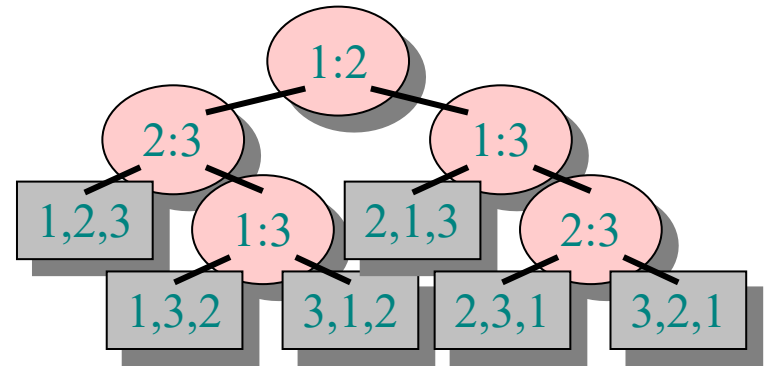
# What is a decision tree?

- A **decision tree** is a tree-like model.
  - Every node is a comparison.
  - Every branch represents the outcome of the comparison
  - Every leaf represents a class label (decision after all comparisons)
- We can use decision tree to model any comparison-based algorithm.

```
if (take CS1231) then
  if (take CS3230) then
    return "Good in Algo &
    Math"
  else
    return "Good in Math"
else
  return "other"
```



# Decision-tree model



A **decision tree** *can model the execution of any comparison sort*:

- One tree for each input size  $n$ .
- View the algorithm as splitting whenever it compares two elements.
- The tree contains the comparisons along all possible instruction traces.
- The running time of the algorithm = the length of the path taken.
- **Worst-case running time = height of tree.**

# Lower bound for decision-tree sorting



- **Theorem:** Any decision tree that can sort  $n$  elements must have height  $\Omega(n \lg n)$ .
- **Proof:** The tree must contain  $\geq n!$  leaves, since there are  $n!$  possible permutations. A height- $h$  binary tree has  $\leq 2^h$  leaves. Thus,  $n! \leq 2^h$ .

$$\begin{aligned} \therefore h &\geq \lg(n!) \quad (\lg \text{ is monotonically increasing}) \\ &\geq \lg((n/e)^n) \quad (\text{Stirling's formula}) \\ &= n \lg n - n \lg e \\ &= \Omega(n \lg n). \quad \square \end{aligned}$$

# Question 2



- Given a **sorted array** of  $n$  real numbers  $A[1..n]$  and a query number  $x$ . You need to develop a function  $\text{search}(x, A)$  which returns an integer  $i$  if  $A[i]=x$ ; and returns  $-1$  otherwise.
- We have two assumptions:
  - Assume comparison model
  - Assume each comparison returns  $<$ , or  $>$ , or  $=$ .
- What is the lower bound of the number of comparisons?

A.  $n$       B.  $\lfloor \lg n \rfloor + 1$       C.  $n \lfloor \lg n \rfloor$

# Question 3



Which of the following is true?

- A.  $f(n) = o(g(n))$
- B.  $f(n) = \Theta(g(n))$
- C.  $f(n) = \omega(g(n))$

when  $f(n) = \ln(n)$  and  $g(n) = \log_{10}(n)$ .

# Question 4



Which of the following is true?

- A.  $f(n) = o(g(n))$
- B.  $f(n) = \Theta(g(n))$
- C.  $f(n) = \omega(g(n))$

when  $f(n) = n^{2.5}$  and  $g(n) = n^2 \log^4 n$ .



# Question 5



Which of the following is true?

- A.  $f(n) = o(g(n))$
- B.  $f(n) = \Theta(g(n))$
- C.  $f(n) = \omega(g(n))$

when  $f(n) = 3^n$  and  $g(n) = 2^n$ .

# Question 6 (If time allows)



- Ali has 81 coconuts, all of which have the same weight, except for one which is heavier. He does not know which is the heavier coconut. Ali's friend has a balance scale, but will charge Ali one dollar for each use of the scale.
- What is the maximum amount of money that Ali has to pay to guarantee that he can find the heaviest coconut, assuming that Ali uses an optimal algorithm?

A. 3      B. 4      C. 5      D. 6