CS3243 : Introduction to Artificial Intelligence

Tutorial 3

NUS School of Computing

February 6, 2023

Admin

► Extra material for practice

Review

- Heuristics, more importantly on how do we come up with them
- ▶ Dominance, and how it helps (as we discussed last time)
- Relaxed problems, the recipe for obtaining heuristics

Review

Admissibility	h is admissible iff $\forall n: h(n) \leq h*(n)$ An admissible heuristic will never overestimate the cost to reach the goal
Consistency	h is consistent, iff for every node n and every successor n' of n generated by action a $h(n) \leq c(n,a,n') + h(n')$
Dominance	$h_1(n)$ dominates $h_2(n)$ iff $\forall n: h_1(n) \geq h_2(n)$

Review (8-puzzle)

Original Problem : A tile can move from square X to square Y if X is adjacent to Y and Y is blank

Review (8-puzzle)

- Original Problem : A tile can move from square X to square Y if X is adjacent to Y and Y is blank
- Relaxed Version 1 : A tile can move from square X to square Y if X is adjacent to Y and Y is blank

Review (8-puzzle)

- Original Problem : A tile can move from square X to square Y if X is adjacent to Y and Y is blank
- Relaxed Version 1 : A tile can move from square X to square Y if X is adjacent to Y and Y is blank
- Relaxed Version 2 : A tile can move from square X to square Y if X is adjacent to Y and Y is blank

- Design Admissible Heuristics for SameGame Puzzle
- Given information :
 - lnitial state : Grid with $n \times m$ tiles with c different colours
 - Neighbors: Directly adjacent tiles (Internal tiles: 4 neighbors, Edge tiles: 3 neighbors, Corner tiles: 2 neighbors)
 - Group : Set of ≥ 2 neighboring tiles of same color
 - Singleton : Tile not belonging to any group
 - ► Action/Move : Deleting a group (not singleton). Thereafter, vertical gravity and column shifting (right-to-left) applies (in transition model)
 - ► Goal state : Empty grid (assume solvable)
 - ▶ Transition cost : 1 (or ∞ if no groups exist)

► An example of the initial state :



▶ Top-left corner is (1,1), $1 \le ROW \le n$, and $1 \le COL \le m$

- ▶ Idea Admissible : $\forall n, \ h(n) \leq h^*(n)$
- ► We have two approaches

- ▶ Idea Admissible : $\forall n, \ h(n) \leq h^*(n)$
- ▶ We have two approaches
- Approach 1: By reasoning/inference (ie What information could I use at each state to get a estimate of the number of moves needed to reach the goal? Important: Must be an underestimate!)

- ▶ Idea Admissible : $\forall n, \ h(n) \leq h^*(n)$
- ▶ We have two approaches
- Approach 1: By reasoning/inference (ie What information could I use at each state to get a estimate of the number of moves needed to reach the goal? Important: Must be an underestimate!)
- ▶ Approach 2 (better approach): Relaxation of the game rules (ie. I now need less moves than what I originally need to reach the goal because the game has become "easier" = admissible. Important: We should be able to easily obtain the number of moves needed with the relaxed rules)

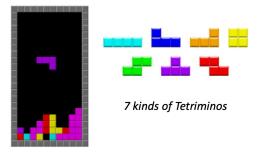
- ▶ Proof of admissibility : h(s) = Number of colours remaining.
- Each group contains exactly 1 colour
- Each remaining color has 1 or more groups
- ▶ Each move may reduce groups by more than 1 (example when groups of same color combine), but will never reduce remaining groups of same color below 1 (unless that color is removed by the move)
- ▶ Hence, each legal move will leave the number of colors remaining unchanged or reduce it by at most 1 if the color is removed by the move, ie. $h(s) \leq h^*(s)$

- ▶ Proof of admissibility : h(s) = Number of colours remaining.
- Each group contains exactly 1 colour
- Each remaining color has 1 or more groups
- ▶ Each move may reduce groups by more than 1 (example when groups of same color combine), but will never reduce remaining groups of same color below 1 (unless that color is removed by the move)
- ▶ Hence, each legal move will leave the number of colors remaining unchanged or reduce it by at most 1 if the color is removed by the move, ie. $h(s) \leq h^*(s)$
- Note: You may also prove admissibility by showing that heuristic was derived from a relaxed version of the game (I invite you to try this!)

Given information :

- State : Partially filled tetris fields with a tetrimino that is about to be placed in the field next (but not placed yet).
- Initial state : Empty field with a starting tetrimino.
- Action : Sequence of lateral movement(s) and/or rotation(s) of a tetrimino (assume player has enough time to shift the tetrimino to an intended configuration before descent).
- Action cost: 1
- Transition model: Takes in a state, applies the sequence of actions on the tetrimino that enters the field, and outputs a state where the tetrimino of that specified configuration descended onto the field.
- Goal state: Completely filled board where there are no gaps (and every tetrimino fits perfectly). (assume solvable)

► An example of an intermediate state :



- Assume we start with a fixed number of tetriminos, N (comprising some of each kind), and all are required to be used to fill the board (i.e. there exists a way to place all these tetriminos such that the board is filled).
- Gap : An empty cell on the board.
- Blocked gap: If for the corresponding column where the gap belongs to, there exists an occupied cell somewhere above that gap.
- A row (or column) is complete if there are no gaps in that row (or column respectively).

▶ Let's determine the admissibility of different heuristics

- \blacktriangleright $h_1(n)$: number of unfielded tetriminos
- ► What do you think?

- \blacktriangleright $h_1(n)$: number of unfielded tetriminos
- ► What do you think?
- Admissible
- ▶ Justification : We need to correctly field each of the unfielded tetriminos (with 1 move each) to reach the goal.

- $ightharpoonup h_2(n)$: number of gaps
- ► What do you think?

- $ightharpoonup h_2(n)$: number of gaps
- ► What do you think?
- Inadmissible
- Recall that a gap is an empty cell. Can we think of a counter-example where fielding 1 tetrimino (with 1 move) could fill > 1 gap-cell?

- ▶ $h_3(n)$: number of incomplete rows
- ► What do you think?

- ▶ $h_3(n)$: number of incomplete rows
- ► What do you think?
- Inadmissible
- Recall an incomplete row = Row w/ any empty cell(s) Can we derive a counter-example where fielding 1 tetriminos (with 1 move) could fill > 1 incomplete row?

- $ightharpoonup h_4(n)$: number of blocked gaps
- ► What do you think?

- $ightharpoonup h_4(n)$: number of blocked gaps
- ► What do you think?
- Admissible
- ▶ Recall that a blocked gap = a gap w/ a cell above it occupied. Problem : Blocked gaps cannot be fielded! (assumption) Tetris fields with ≥ 1 blocked gap cannot be solved and hence h^* on these nodes is infinite while h_4 is finite, hence $h_4 < h^*$ on these nodes. For nodes on the optimal path, $h_4 = 0$ while h^* is ≥ 0 .

- $ightharpoonup max(h_1,h_2)$ is admissible
- ► What do you think?

- $ightharpoonup max(h_1,h_2)$ is admissible
- ► What do you think?
- ► False
- ▶ Recall that h_1 is admissible and h_2 is inadmissible. For nodes where $h_2 > h^*$, $max(h_1, h_2)$ will also be $> h^*$. Hence, $max(h_1, h_2)$ is inadmissible.

- $ightharpoonup min(h_2,h_3)$ is admissible
- ► What do you think?

- $ightharpoonup min(h_2,h_3)$ is admissible
- ► What do you think?
- ► False
- Recall that h_2 and h_3 are both inadmissible. But this does not mean that their min must be inadmissible as well! If the heuristics "complement" each other, ie. for nodes where h_2 overestimates $(h_2 > h^*)$, h_3 is guaranteed to not overestimate (and vice versa), then we are safe. But is it true in this case?

- $ightharpoonup min(h_2,h_3)$ is admissible
- ► What do you think?
- ► False
- Recall that h_2 and h_3 are both inadmissible. But this does not mean that their min must be inadmissible as well! If the heuristics "complement" each other, ie. for nodes where h_2 overestimates $(h_2 > h^*)$, h_3 is guaranteed to not overestimate (and vice versa), then we are safe. But is it true in this case?
- We need to look at the heuristics / previous counter-examples in detail.

- $ightharpoonup max(h_3,h_4)$ is inadmissible
- ► What do you think?

- $ightharpoonup max(h_3,h_4)$ is inadmissible
- ► What do you think?
- ► True
- Recall that h_3 is inadmissible while h_4 is admissible. For nodes where $h_3 > h^*$, $max(h_3, h_4)$ will also be $> h^*$.

- $ightharpoonup h_1$ dominates h_2
- ► What do you think?

- $ightharpoonup h_1$ dominates h_2
- ► What do you think?
- ► False
- Recall that h_1 is admissible while h_2 is inadmissible. There are nodes where $h_2 > h^* > h_1$. Hence h_1 does not dominate h_2 .

- $ightharpoonup h_2$ dominates h_4
- ▶ What do you think?

- $ightharpoonup h_2$ dominates h_4
- ► What do you think?
- ▶ Recall that h_2 is inadmissible while h_4 is admissible. But this does not mean that h_2 dominates h_4 ! Even while some nodes have $h_2 > h^* > h_4$, there could be other nodes where $h_2 < h_4$. Hence, we need to look at the heuristics in detail.

- $ightharpoonup h_2$ dominates h_4
- ► What do you think?
- ▶ Recall that h_2 is inadmissible while h_4 is admissible. But this does not mean that h_2 dominates h_4 ! Even while some nodes have $h_2 > h^* > h_4$, there could be other nodes where $h_2 < h_4$. Hence, we need to look at the heuristics in detail.
- True
- ▶ Recall that a blocked gap is an empty cell that has an additional requirement that its top is blocked by an occupied cell above it. Since each blocked gap is also a gap but not all gaps are blocked gaps, blocked gaps has to be a subset of gaps. Hence, $h_2 > h_4$ at all nodes and h_2 dominates h_4 .

- ▶ h_3 does not dominate h_2
- ► What do you think?

- $ightharpoonup h_3$ does not dominate h_2
- What do you think?
- Recall that both h_3 and h_2 are inadmissible hence we cannot conclude anything from their properties. Let's look at the heuristic in detail.

- \blacktriangleright h_3 does not dominate h_2
- What do you think?
- Recall that both h_3 and h_2 are inadmissible hence we cannot conclude anything from their properties. Let's look at the heuristic in detail.
- ► True
- Since each incomplete row has to have ≥ 1 gap, h_2 has to be $\geq h_3$ and hence h_3 does not dominate h_2 .

- \blacktriangleright h_3 does not dominate h_2
- What do you think?
- Recall that both h_3 and h_2 are inadmissible hence we cannot conclude anything from their properties. Let's look at the heuristic in detail.
- ► True
- Since each incomplete row has to have ≥ 1 gap, h_2 has to be $\geq h_3$ and hence h_3 does not dominate h_2 .
- OR give a counter-example.

- ▶ h_4 does not dominate $h_2/2$
- ► What do you think?

- \blacktriangleright h_4 does not dominate $h_2/2$
- ► What do you think?
- ► True
- ▶ In the initial state, $h_2/2 = \#$ of gaps / 2 = no. of cells in grid / 2. But $h_4 = \#$ of blocked gaps = 0. Since $h_4 < h_2/2$, h_4 does not dominate h_2 .

▶ Dominance of Heuristics in Question 3

- ▶ Dominance of Heuristics in Question 3
- Sort of a Giveaway

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$
- \blacktriangleright h_1 and h_3 ?

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$
- \blacktriangleright h_1 and h_3 ?
- $\blacktriangleright h_1$ and h_3 do not have a dominance relationship

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$
- \blacktriangleright h_1 and h_3 ?
- $ightharpoonup h_1$ and h_3 do not have a dominance relationship
- $ightharpoonup h_1$ and h_4 ?

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$
- \blacktriangleright h_1 and h_3 ?
- $\blacktriangleright h_1$ and h_3 do not have a dominance relationship
- $ightharpoonup h_1$ and h_4 ?
- $ightharpoonup h_1$ and h_4 do not have a dominance relationship

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$
- \blacktriangleright h_1 and h_3 ?
- $ightharpoonup h_1$ and h_3 do not have a dominance relationship
- $ightharpoonup h_1$ and h_4 ?
- $ightharpoonup h_1$ and h_4 do not have a dominance relationship
- $ightharpoonup h_3$ and h_4 ?

- Dominance of Heuristics in Question 3
- Sort of a Giveaway
- Note from Prof : You should only discuss the dominance part, justifications are to be done by students.
- Any thoughts?
- $\blacktriangleright h_1, h_3, \text{ and } h_4 \text{ are admissible heuristics}$
- \blacktriangleright h_1 and h_3 ?
- $lacktriangledown h_1$ and h_3 do not have a dominance relationship
- $ightharpoonup h_1$ and h_4 ?
- lacktriangle h_1 and h_4 do not have a dominance relationship
- $ightharpoonup h_3$ and h_4 ?
- $ightharpoonup h_3$ dominates h_4

Thank you!

If you have any questions, please don't hesitate. Feel free to ask! We are here to learn together!