# Logical Agents: Knowledge Representation

CS3243: Introduction to Artificial Intelligence – Lecture 8

13 March 2023

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- 3. Wumpus World & Entailment
- 4. Inference Algorithms: Soundness & Completeness
- 5. Inference via Truth Table Enumeration

# Administrative Matters

#### Upcoming...

- Deadlines
  - TA6 (released last week)
    - Due in your Week 9 tutorial session
    - Submit the a physical copy (more instructions on the Tutorial Worksheet)
  - Prepare for the tutorial!
    - Participation marks = 5%
  - Project 2 (released Week 6)
    - Due this Sunday (19 March), 2359 hrs
  - Project 3 (released this week)
    - Due Week 12 Sunday (9 April), 2359 hrs

#### Midterm Appeals

- Appeals → to your tutor
- Deadline = Day of your Week 9 Tutorial

# A Problem with Problem-Solving Agents

#### **Problem-Solving Agents**

- Problem-solving agents try to find a solution via Search
- No real model of what the agent knows
  - Each state contains knowledge on state of entire environment
    - Knows actions and transition model
    - Implicit general facts about the environment
      - Route Finding Agent implicit knowledge that road lengths cannot be negative
      - 8-puzzle implicit knowledge that two number-tiles cannot occupy the same grid
  - Atomic representations limiting
    - Imagine a game of minesweeper where the environment is only partially observable;
       the agent would not know where all the mines actually were
    - A problem-solving agent would typically use a representation that includes all
      possible mine positions (with accompanying adjacent mine numbers) in an attempt
      to search for a viable solution from the current board



On to agents with generalised knowledge representations: Knowledge-Based Agents

## Knowledge-Based Agents: Logical Agents

#### **Knowledge-Based Agents**

- Represent agent domain knowledge using logical formulas
- General idea
  - Make inferences on existing information
    - Use existing knowledge to infer new information
  - States similar to CSPs
    - Represented as assignments of values to variables
- Agent structure



## Knowledge Base (KB)

- What is a knowledge base (KB)?
  - Set of sentences in a formal language
    - Sentences are expressive and parsable
  - Pre-populate with domain knowledge
    - Example: game rules, general rules/knowledge
- Declarative approach to problem-solving
  - TELL it what it needs to know
    - Update with percept/state/action information
  - ASK itself what to do
    - Make inferences that help determine what actions to take
      - Answers should follow from the KB

## **KB** Agent Function

What happened?

**function** KB-AGENT(percept) **returns** an action **persistent**: KB, a knowledge base t, a counter, initially 0, indicating time

Tell(KB, Make-Percept-Sentence(percept, t))

 $action \leftarrow Ask(KB, Make-Action-Query(t))$ 

What did I perceive at time *t*?

What did I do?

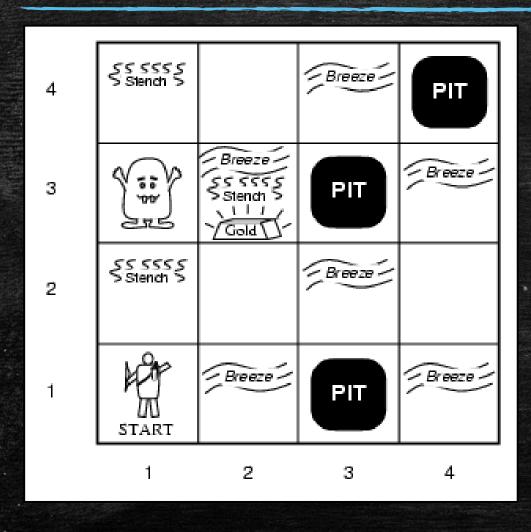
Tell(KB, Make-Action-Sentence(action, t))  $t \leftarrow t + 1$ return action

What is the best action at time *t*?

- Agent must be able to
  - Represent states, actions, etc.
  - Incorporate new percepts
  - Update internal representation of environment
  - Deduce hidden environment properties, and deduce actions

# An Example: The Wumpus Dungeon

### **About Wumpus World**



#### Performance Measure

- Optimise score
  - Obtain Gold: +1000Death: -1000
  - Each Action: -1
  - Fire Arrow: -10

#### Environment

- 4×4 grid of rooms
  - Agent
  - Wump
  - Gold

#### **Actuators**

- Turn left/right
- Move forward
- Fire arrow (kills Wumpus if facing it; uses up arrow)
- Grab gold
- Exit Wumpus dungeon (by climbing out at (1,1))

#### Sensors

- Rooms adjacent to Wumpus are SMELLY
- Rooms adjacent to Pit are BREEZY
- Gold glitters (can detect it if in same room)
- Bump into walls
- Hear scream if Wumpus killed

## **Properties of Wumpus World**

- Not fully observable
  - Only local perception
  - Don't know what is in unexplored rooms
- Deterministic
- Sequential
- Static
- Discrete
- Single Agent

1,4		2,4	3,4	4,4
1,3		2,3	3,3	4,3
1,2 <b>OK</b>		2,2 <b>P</b> ?	3,2	4,2
1,1 <b>OK</b>	A V	2,1 <b>A B OK</b>	3,1 <b>P</b> ?	4,1

**A** = Agent

 $\mathbf{B}$  = Breeze

**G** = Glitter, Gold

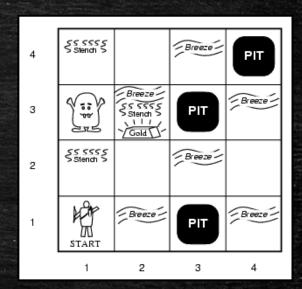
**OK** = Safe Square

 $\mathbf{P}$  = Pit

**S** = Stench

**V** = Visited

**W** = Wumpus



Start at (1,1); Infer that (1,2) and (2,1) are OK (i.e., safe)

Iteration 1: Move to (2,1)

Iteration 2: Move back to (1,1)

1		T	T
1,4	2,4	3,4	4,4
,	'	'	<b>,</b>
1.0		0.0	4.0
1,3	2,3	3,3	4,3
1,2 <b>A</b>	2.2	3,2	4,2
$\int_{0}^{\infty} S^{1}$		-,_	-,_
3	<b>P</b> ?		
OK			
1		0.4	1 1
	2,1	3,1	4,1
	2,1 <b>R</b>		4,
	В	<b>P?</b>	4,1
	<del> </del>	0.4	1 1

**A** = Agent

 $\mathbf{B}$  = Breeze

G = Glitter, Gold

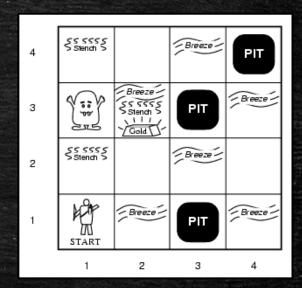
**OK** = Safe Square

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Start at (1,1); Infer that (1,2) and (2,1) are OK (i.e., safe)

Iteration 1: Move to (2,1)

Iteration 2: Move back to (1,1)

Iteration 3: Move to (1,2)

	1,4	2,4	3,4	4,4
7.00 2.00 1.00 1.00 1.00 1.00 1.00 1.00 1				
	1,3	2,3	3,3	4,3
No Breeze at (1,2)	<b>W</b> !			
	1,2 <b>A</b>	2,2	3,2	4,2
7 2.70 24 43.70 00 0	S	<b>P</b> ?		
	OK	OK		
No Stench	1,1	2,1	3,1	4,1
at (2,1)		<b>B</b>	<b>P</b> ?	
	OK V	OK V		

A = Agent
B = Breeze
G = Glitter, Gold

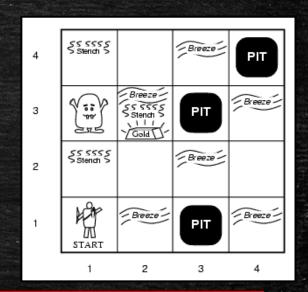
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Start at (1,1); Infer that (1,2) and (2,1) are OK (i.e., safe)

Iteration 1: Move to (2,1)

Iteration 2: Move back to (1,1)

Iteration 3: Move to (1,2)

Iteration 4: Move to (2,2)

1,4	2,4	3,4	4,4
1,3 <b>W</b> !	2,3 <b>OK</b>	3,3	4,3
1,2 S OK V	2,2 <b>A OK</b>	3,2 <b>OK</b>	4,2
1,1 OK V	2,1 B OK V	3,1 <b>P</b> !	4,1

A = Agent

 $\mathbf{B}$  = Breeze

**G** = Glitter, Gold

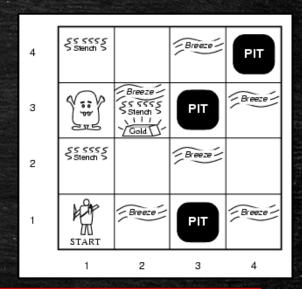
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Start at (1,1); Infer that (1,2) and (2,1) are OK (i.e., safe)

Iteration 1: Move to (2,1)

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Iteration 3: Move to (1,2)

Iteration 4: Move to (2,2)

1,4	2,4	3,4	4,4
1,3 <b>W</b> !	2,3 <b>A SBG OK</b>	3,3	4,3
1,2 <b>S</b>		3,2	4,2
OK V	OK V	OK	
1,1	2,1	3,1	4,1
	В	<b>P</b> !	
OK V	OK V		

A = Agent

 $\mathbf{B}$  = Breeze

**G** = Glitter, Gold

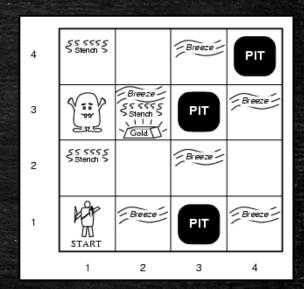
**OK** = Safe Square

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Start at (1,1); Infer that (1,2) and (2,1) are OK (i.e., safe)

Iteration 1: Move to (2,1)

Iteration 2: Move back to (1,1)

Iteration 3: Move to (1,2)

Iteration 4: Move to (2,2)

Iteration 5: Move to (2,3)

# Logic

#### Review of Logic

#### Logic

- Formal language for knowledge representation (KR)
- Allows the inference of conclusions about environment

#### Syntax

- Defines sentences in the language
- Semantics
  - Defines meaning of sentences
- Truth value
  - Statement result given observed values
    - Defines validity of a sentence within the environment
      - i.e., given value assignments that hold in the environment

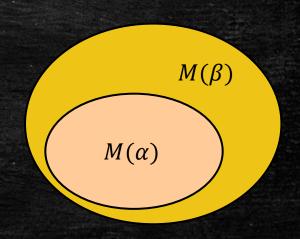
#### Review of Logic

- Example of KR: language of arithmetic
  - Syntax
    - $x + 2 \ge y$  is a sentence
    - x2y +> is not a sentence
  - Truth values
    - $x + 2 \ge y$  is true in a world where x = 7, y = 1
    - $x + 2 \ge y$  is false in a world where x = 0, y = 6

# Entailment

#### Entailment

- Modelling
  - v models  $\alpha$  if  $\alpha$  is true under v
    - v corresponds to one set of value assignments (applied to sentences  $\alpha$ )
    - v corresponds to one instance of the environment (known part of a state)
  - For example
    - $\alpha = (q \in \mathbb{Z}_+) \land (\forall n, m \in \mathbb{Z}_+ : q = nm \Rightarrow n \lor m = 1)$
    - For which values of q will  $\alpha$  be true?
- Let  $M(\alpha)$  be the set of all models for  $\alpha$
- Entailment (⊨) means that one thing follows from the another
  - $-\alpha \models \beta$  or equivalently  $M(\alpha) \subseteq M(\beta)$
  - Example:
    - $[\alpha = (q \text{ is prime})] \models [\beta = (q \text{ is odd}) \lor (q = 2)]$

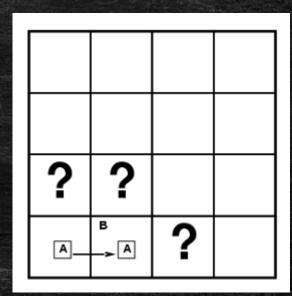


#### Situation:

- Detected Nothing at (1,1)
- Moved Right to (2,1)
- Detected Breeze at (2,1)

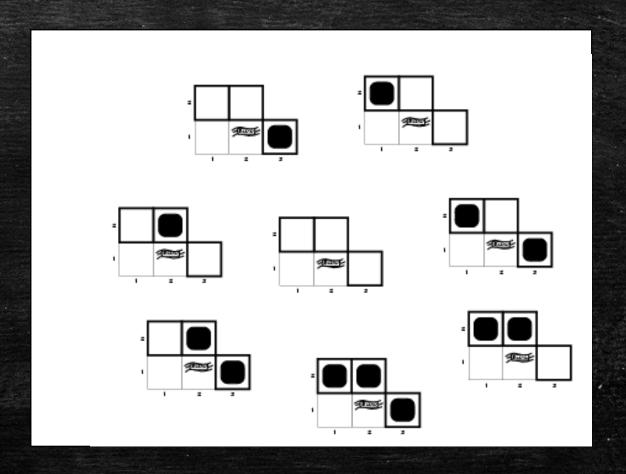
#### Consider possible models for KB with pits

- 3 Boolean choices ⇒ 8 possible models
  - Pit or No Pit at: (1,2), (2,2), (3,1)
  - All  $(2^3 = 8)$  permutations for the above (each a possible model)

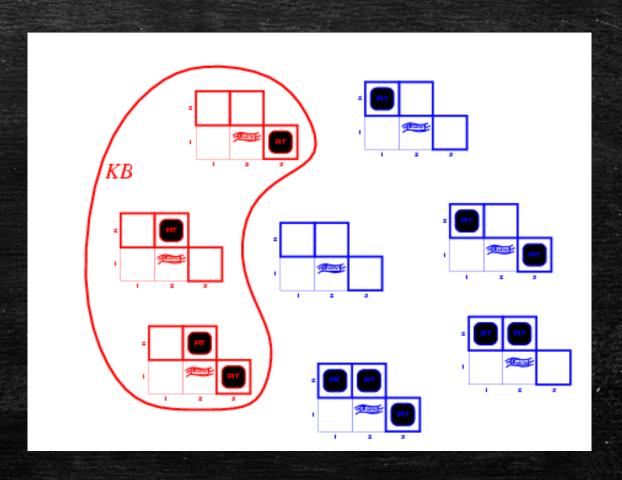


Within this example, we will only deal with pits only; we ignore the Wumpus.

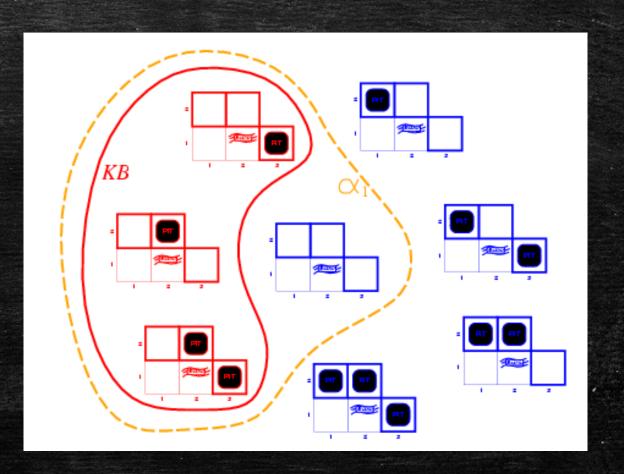
Possible 8 models



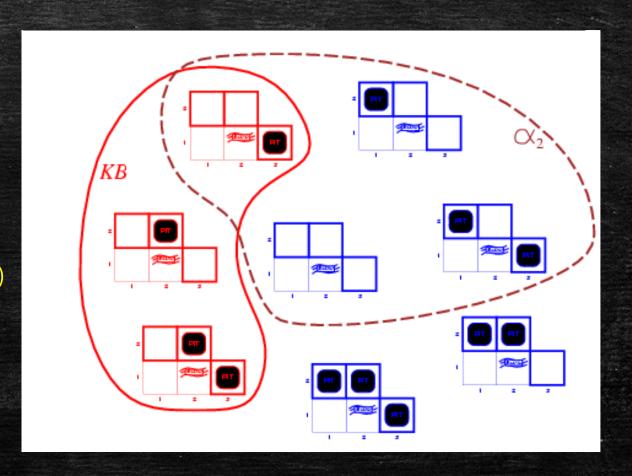
- KB = rules + percepts
- Percepts
  - No Breeze at (1,1)
  - Breeze at (2,1)
- Relevant rules
  - No Pit at (1,1)
  - No Pit at (1,2)
  - No Pit at (2,1)
  - Breeze ⇒Pit in Adjacent Room



- KB = rules + percepts
- Let  $\alpha_1 = (1,2)$  is Safe
- We observe that:  $M(KB) \subseteq M(\alpha_1)$
- Or rather:  $KB \models \alpha_1$
- We may thus infer that it is safe for the agent to move to (1,2)



- KB = rules + percepts
- Let  $\alpha_2 = (2,2)$  is Safe
- We observe that: KB  $\not\models \alpha_2$
- Since:  $M(KB) \nsubseteq M(\alpha_1)$
- May NOT infer that it is safe for the to move to (2,2)
  - Exist some models where KB is True but  $\alpha_2$  is False
  - For entailment, we want all  $\alpha_2$  True when KB True
- Also, cannot infer unsafe!



#### Questions about the Lecture?

- Was anything unclear?
- Do you need to clarify anything?

- Ask on Archipelago
  - Specify a question
  - Upvote someone else's question



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# Propositional Logic

## Review of Propositional Logic: Syntax

- A simple language for logic illustrates basic ideas
- Defines allowable sentences
- Sentences are represented by symbols e.g., s<sub>1</sub>, s<sub>2</sub>
  - Formed over basic variables
- Logical connectives for constructing complex sentences from simpler ones
  - If s is a sentence, ¬s is a sentence (negation)
  - If s<sub>1</sub> and s<sub>2</sub> are sentences:
    - s₁ ∧ s₂ is a sentence (conjunction)
    - s<sub>1</sub> v s<sub>2</sub> is a sentence (disjunction)
    - $s_1 \Rightarrow s_2$  is a sentence (implication)
    - $s_1 \Leftrightarrow s_2$  is a sentence (biconditional *iff*.)

### Review of Propositional Logic: Semantics

- A model
  - Truth assignment to the given basic variables
  - Given n variables, 2<sup>n</sup> truth assignments
- All other sentences' truth value are derived according to logical rules
  - Example
    - Given  $x_1$  = True;  $x_2$  = False;  $x_3$  = True
    - What is the truth value for  $(x_1 \land \neg x_2) \Rightarrow \neg (x_3 \lor (\neg x_1 \land x_2))$ ?
      - Recall that  $X \Rightarrow Y$  is true if X false, or X true and Y true

#### Wumpus World KB

#### Notation

- $P_{ij}$  = True  $\Leftrightarrow$  Pit at (i, j)
- $B_{ij}$  = True  $\Leftrightarrow$  Breeze at (i, j)

#### - Given

- $R_1: \neg P_{1,1}$
- $R_2: \neg B_{1,1}$
- $R_3: B_{2,1}$
- Rules: "Pits cause a breeze in adjacent squares"
  - $-R_4: B_{1,1} \Leftrightarrow \overline{(P_{1,2} \vee P_{2,1})} \text{i.e.}, \neg B_{1,1} \Leftrightarrow \overline{\neg (P_{1,2} \vee P_{2,1})}$
  - $R_5: B_{2,1} \Leftrightarrow \left(P_{1,1} \vee P_{2,2} \vee P_{3,1}\right)$

Within this example, again, we will only deal with pits only; we ignore the Wumpus.

KB is true iff  $\Lambda_{k=1,...,5} R_k$  is true

## Inference: Objectives & Application

- Given a KB, infer something non-obvious about the environment
- Mimic logical human reasoning
- After exploring 3 squares, we have some understanding of the Wumpus World
- Inference ⇒ Deriving knowledge out of percepts

Given KB and  $\alpha$ , we want to know if KB  $\models \alpha$ 

What  $\alpha$ ?

Based on domain: e.g., is (1,2) safe?

# Properties of Inference Algorithms

#### Soundness & Completeness

- KB ⊢<sub>A</sub> α
  - Means: "sentence  $\alpha$  is derived (i.e., inferred) from KB by inference algorithm  $\mathcal{A}$ "
- Soundness
  - $\mathcal{A}$  is sound if KB  $\vdash_{\mathcal{A}} \alpha$  implies KB  $\vDash \alpha$
  - This means that A will not infer nonsense
    - For all sentences inferred from the KB by A, S
    - The KB will entail each  $\alpha$  in S
- Completeness
  - $\mathcal{A}$  is complete if KB  $\models \alpha$  implies KB  $\vdash_{\mathcal{A}} \alpha$
- This means that A can infer any sentence that the KB entails
  - If KB entails a sentence (any sentence describing a superset of the KB)
  - A can infer that sentence

Determine if an inference algorithm is complete and sound

#### Soundness & Completeness

#### More on completeness

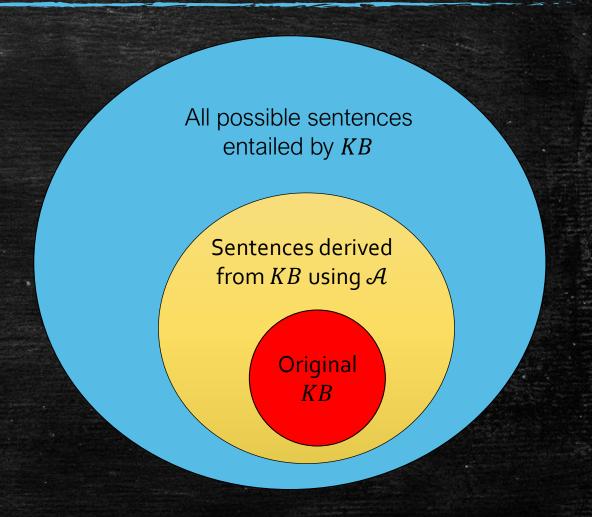
- If A is incomplete
- A cannot reach all possible conclusions

#### Given

- Y = all possible sentences entailed by KB
- X = all sentences derived from KB using A

#### Then

- X = Y: sound, complete
- X ⊂ Y: sound, not complete
- Y ⊂ X: not sound, complete
- X ⊄ Y, Y ⊄ X, X ≠ Y: not sound, not complete



## **Truth Table Enumeration**

## Truth Table Enumeration Example: Wumpus World

Can we infer that (1,2) is safe from pits?

$$\alpha_1 = \neg P_{1,2}$$

$B_{1,1}$	$B_{2,1}$	$P_{1,1}$	$P_{1,2}$	$P_{2,1}$	$P_{2,2}$	$P_{3,1}$	KB	$\alpha_1$
false	true							
false	false	false	false	false	false	true	false	true
:	:	:	:	:	:	:	:	:
false	true	false	false	false	false	false	false	true
false	true	false	false	false	false	true	$\underline{true}$	$\underline{true}$
false	true	false	false	false	true	false	$\underline{true}$	$\underline{true}$
false	true	false	false	false	true	true	$\underline{true}$	$\underline{true}$
false	true	false	false	true	false	false	false	true
:	:	:	:	:	:	:	:	:
true	false	false						

$$R_1: \neg P_{1,1}$$
 $R_2: \neg B_{1,1}$ 
 $R_3: B_{2,1}$ 
 $R_4: \neg B_{1,1} \Leftrightarrow \neg (P_{1,2} \lor P_{2,1})$ 
 $R_5: B_{2,1} \Leftrightarrow (P_{1,1} \lor P_{2,2} \lor P_{3,1})$ 

#### **KB** true

KB is true iff  $\Lambda_{k=1,...,5} R_k$  is true

Recall that a truth table contains every possible truth assignment (2<sup>7</sup> models in this example)

Does  $\overline{KB}$  entail  $\alpha_1$ ? (Whenever KB true,  $\alpha_1$  true?)

#### Truth Table Enumeration

**function** TT-ENTAILS?(KB,  $\alpha$ ) **returns** true or false

```
symbols \leftarrow a list of the proposition symbols in KB and \alpha
  return TT-CHECK-ALL(KB, \alpha, symbols, \{\})
function TT-CHECK-ALL(KB, \alpha, symbols, model) returns true or false
  if EMPTY?(symbols) then
      if PL-True?(KB, model) then return PL-True?(\alpha, model)
      else return true
                             // when KB is false, always return true
  else
      P \leftarrow FIRST(symbols)
      rest \leftarrow REST(symbols)
      return (TT-CHECK-ALL(KB, \alpha, rest, model \cup \{P = true\})
              and
              TT-CHECK-ALL(KB, \alpha, rest, model \cup \{P = false \})
```

**inputs**: KB, the knowledge base, a sentence in propositional logic

 $\alpha$ , the query, a sentence in propositional logic

Checks all  $2^n$  truth assignments to verify KB entails  $\alpha$ 

Depth-first enumeration

Recursive step generates the 2<sup>n</sup> possible assignments to the n symbols

O(2<sup>n</sup>) time complexity O(n) space complexity

Implements definition of entailment directly

(guarantees soundness)

Finite models to check (guarantees completeness)

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