Preliminaries

Definitions

• [Rate of Convergence] Let $(\alpha_n)_n$, $(\beta_n)_n$ be sequences that converges to α and 0 respectively. If $\exists K$ s.t. $|\alpha_n - \alpha| \le K |\beta_n|$ for large n, then $(\alpha_n)_n$ converges to α with rate of convergence $O(\beta_n)$. Write $\alpha_n = \alpha + O(\beta_n)$

- [Big 0] Let F, G be functions s.t. $\lim_{h\to 0} F(h) = L$ and $\lim_{h\to 0} G(h) = 0$. If $\exists K$ s.t. $|F(h) L| \le K|G(h)|$ for small enough h, then write F(h) = L + O(G(h))
- [Order of Convergence] Let $(p_n)_n$ be a sequence that converges to p, with $p_n \neq p \ \forall n$. If $\lim_{n\to\infty} \frac{|p_{n+1}-p|}{|p_n-p|^{\alpha}} = \lambda$, then $(p_n)_n$ converges to p of order α , with asymptotic error constant λ
 - \circ [Linearly Convergent] $\alpha = 1$ and $\lambda < 1$ implies linearly convergent
 - o [Quadratically Convergent] $\alpha = 2$ implies quadratically convergent
- [Superlinearly Convergent] Say sequence $(p_n)_n$ is superlinearly convergent to p if $\lim_{n\to\infty} \frac{|p_{n+1}-p|}{|p_n-p|} = 0$

Root Finding

Definitions

- [Horner Method] A method for evaluating polynomials via nesting
 - $P(x_0) = a_0 + x_0 (a_1 + x_0 (a_2 + \cdots))$
 - o $P'(x_0)$ and $P(x_0)$ can be evaluated in a single pass by considering P(x) = $(x-x_0)Q(x) + b_0$ with $Q(x) = b_1 + b_2x + \cdots$ and calculating $(b_i)_i$
- [Bisection Method] Finds solution to f(x) = 0 for continuous function f on interval [a, b]; requires that f(a)f(b) < 0
- [Fixed Point] Given function g, p is a fixed point for g if g(p) = p
- [Fixed Point Iteration] Start with p_0 and define $p_{n+1} = g(p_n)$ to generate $(p_n)_n$
- [Newton's Method] Root finding to f(x) = 0; $x_{n+1} = x_n \frac{f(x_n)}{f'(x_n)}$
 - View as a special fixed point iteration with $g(p) = p \frac{f(p)}{f'(p)}$
 - Fails if $f'(x_n) = 0$ for some x_n
 - o If f(p) = 0 and $f'(p) \neq 0$, then for starting values sufficiently close to p, Newton's method will converge at least quadratically
- [Secant's Method] Root finding; $x_{n+1} = x_n \frac{f(x_n)}{f(x_n) f(x_{n-1})}$
 - Doesn't need derivative evaluation (unlike Newton)
- [Simple Zero] Let $f \in C^1[a, b]$. Then f has a simple zero at $p \in (a, b)$ if f(p) = 0 but $f'(p) \neq 0$.
- [Zero of Multiplicity m] A solution p of f(x) = 0 is a zero of multiplicity m of f if for $x \neq p$, we can write $f(x) = (x - p)^m q(x)$ where $\lim_{x \to a} q(x) \neq 0$
- [Forward Difference] $\Delta p_n = p_{n+1} p_n$, $\Delta^k p_n = \Delta(\Delta^{k-1}p_n) = \Delta^{k-1}p_{n+1} \Delta^{k-1}p_n$ [Aitken's Δ^2 Method] Define $\hat{p}_n = p_n \frac{(p_{n+1}-p_n)^2}{p_{n+2}-2p_{n+1}+p_n} = p_n \frac{(\Delta p_n)^2}{\Delta^2 p_n} = \{\Delta^2\}(p_n)$. The assumption is that $(\hat{p}_n)_n$ converges more rapidly to p than $(p_n)_n$
- [Forward Difference] $\Delta p_n = p_{n+1} p_n$
- [Steffensen's Method]
 - $\circ \quad p_0^{(0)}; \ p_1^{(0)} = g(p_0^{(0)}); \ p_2^{(0)} = g(p_1^{(0)})$
 - $\circ \quad p_0^{(1)} = \{\Delta^2\} \Big(p_0^{(0)} \Big); \ p_1^{(1)} = g\Big(p_0^{(1)} \Big); \ p_2^{(1)} = g\Big(p_1^{(1)} \Big)$
 - $o p_0^{(2)} = \{\Delta^2\} (p_0^{(1)})$
- [Müller's Method] Finds a solution to f(x) = 0 given three initial points p_0, p_1, p_2 ; uses parabolas to interpolate and find closer root to p_2 ; can find real and complex roots.

$$o p_3 = p_2 - \frac{2c}{b + \operatorname{sign}(b)\sqrt{b^2 - 4ac}}$$

Theorems

- [2.1] Let $f \in C[a,b]$ with f(a)f(b) < 0. Then, bisection method generates $(p_n)_n$ with $|p_n-p| \leq \frac{b-a}{2^n}$ i.e. $p_n=p+O\left(\frac{1}{2^n}\right)$
- [2.3] Let $g \in C[a, b]$ and $g(x) \in [a, b]$. Then g has at least one fixed point in [a, b]. If in addition, g'(x) exists on (a,b) and $\exists K < 1$ s.t. $|g'(x)| \le K \ \forall x \in (a,b)$, then there is exactly one fixed point in [a, b].
- [2.4] Let $g \in C[a,b]$ be s.t. $g(x) \in [a,b] \ \forall x \in [a,b]$. Suppose g' exists on (a,b) and $\exists K \in [a,b]$ (0,1) s.t. $|g'(x)| \le K$ for $x \in (a,b)$. Then, for any number $p_0 \in [a,b]$, $\lim_{n \to \infty} p_n = p$.
 - o [2.5] $|p_n p| \le K^n \max(p_0 a, b p_0)$
 - \circ [2.8] If $g'(p) \neq 0$, then for any $p_0 \neq p$ in [a, b], the sequence $p_n = g(p_{n-1})$ converges linearly with asymptotic constant |g'(p)|

• [2.6] Let $f \in C^2[a,b]$. If $p \in (a,b)$ s.t. f(p)=0 and $f'(p)\neq 0$, then $\exists \delta>0$ s.t. Newton's method generates sequence $(p_n)_n$ converging to p for any initial $p_0 \in [p-\delta,p+\delta]$

- [2.9] Let p be a solution of g(x)=x. Suppose g'(p)=0 and g'' continuous with |g''(x)| < M on an open interval I containing p, then $\exists \delta > 0$ s.t. for $p_0 \in [p-\delta, p+\delta]$, the sequence $(p_n)_n$ converges at least quadratically to p. Moreover, for sufficiently large values of p, $|p_{n+1}-p| < \frac{M}{2}|p_n-p|^2$
- [2.11] Let $f \in C^1[a, b]$. Then f has a simple zero at $p \in (a, b)$ if and only if f(p) = 0 and $f'(p) \neq 0$.
- [2.14] Let $(p_n)_n$ converge linearly to limit p with $\lim_{n\to\infty}\frac{p_{n+1}-p}{p_n-p}<1$. Then Aitken's Δ^2 sequence $(\hat{p}_n)_n$ converges to p faster in the sense that $\lim_{n\to\infty}\frac{\hat{p}_n-p}{p_n-p}=0$
- [2.15] Suppose x = g(x) has solution p with $g'(p) \neq 1$. If there exists $\delta > 0$ s.t. $g \in C^3[p-\delta,p+\delta]$, then Steffensen's method gives quadratic convergence for any $p_0 \in [p-\delta,p+\delta]$

Interpolation

Definitions (General Interpolation)

- [Lagrange Interpolating Polynomial] $P(x) = \sum_{i=0}^{n} f(x_i) \prod_{j \neq i} \frac{(x x_i)}{(x_i x_j)} = \sum_{i=0}^{n} f(x_i) L_i(x)$
 - $\text{o} \quad \text{Then } f(x) = \sum_{k=0}^{n} f(x_k) L_k(x) + \frac{(x x_0) \dots (x x_n)}{(n+1)!} f^{(n+1)} \left(\xi(x) \right) \text{ for some } \xi(x) \in I$
- [Neville's Method] Let Q(x) interpolate f(x) at $x_0, ..., x_k$ and $\hat{Q}(x)$ interpolate f(x) at $x_1, ..., x_{k+1}$, then $P(x) = \frac{(x-x_{k+1})Q(x)-(x-x_0)\hat{Q}(x)}{x_0-x_{k+1}}$ interpolates all points.
 - $\circ Q_i(x) = f(x_i)$
 - $Q_{i,i+1,\dots,j,j+1}(x) = \frac{(x-x_{j+1})Q_{i,\dots,j}(x) (x-x_i)Q_{i+1,\dots,j+1}(x)}{x_i x_{j+1}}$
 - o Can do on-the-fly memoisation

Neville	Divided Difference
$ \begin{array}{ccc} f(x_0) & \searrow \\ f(x_1) & \searrow & Q_{0,1}(x) & \searrow \end{array} $	$ \begin{array}{ccc} f[x_0] & \searrow \\ f[x_1] & \searrow & f[x_0, x_1] & \searrow \end{array} $
$f(x_2)$ $\stackrel{\longrightarrow}{\searrow}$ $Q_{1,2}(x)$ $\stackrel{\longrightarrow}{\searrow}$ $Q_{0,1,2}(x)$	$f[x_2] \stackrel{\longrightarrow}{\searrow} f[x_1, x_2] \stackrel{\longrightarrow}{\searrow} f[x_0, x_1, x_2]$
E S	1 N
$f(x_{n-2}) \stackrel{\longrightarrow}{\searrow} Q_{n-3,n-2}(x) \stackrel{\longrightarrow}{\searrow} Q_{n-4,n-3,n-2}(x) \stackrel{\longrightarrow}{\searrow} \cdots$	$f[x_{n-2}] \stackrel{\longrightarrow}{\searrow} f[x_{n-3}, x_{n-2}] \stackrel{\longrightarrow}{\searrow} f[x_{n-4}, x_{n-3}, x_{n-2}] \stackrel{\longrightarrow}{\searrow} \cdots$
$f(x_{n-1}) \stackrel{\longrightarrow}{\searrow} Q_{n-2,n-1}(x) \stackrel{\longrightarrow}{\searrow} Q_{n-3,n-2,n-1}(x) \stackrel{\longrightarrow}{\searrow} \cdots \stackrel{\searrow}{\searrow}$	$f[x_{n-1}] \stackrel{\longrightarrow}{\searrow} f[x_{n-2}, x_{n-1}] \stackrel{\longrightarrow}{\searrow} f[x_{n-3}, x_{n-2}, x_{n-1}] \stackrel{\longrightarrow}{\searrow} \cdots \searrow$
$f(x_n) \longrightarrow Q_{n-1,n}(x) \longrightarrow Q_{n-2,n-1,n}(x) \longrightarrow \cdots \longrightarrow Q_{0,1,\dots,n}(x)$	$f[x_n] \longrightarrow f[x_{n-1}, x_n] \longrightarrow f[x_{n-2}, x_{n-1}, x_n] \longrightarrow \cdots \longrightarrow f[x_0, \cdots, x_n]$

- [Divided Difference] Given $(x_0, f(x_0)), \dots, (x_n, f(x_n))$, produces the coefficients of the interpolating polynomial of the specific form $P_n(x) = a_0 + a_1(x x_0) + \dots + a_n(x x_0) \dots (x x_{n-1})$
 - \circ [Zeroth Divided Difference] $f[x_i] = f(x_i)$
 - $\circ f[x_i, x_i] = f'(x_i)$
 - $\circ f[x_i, ..., x_j] = \frac{f[x_{i+1}, ..., x_j] f[x_i, ..., x_{j-1}]}{x_i x_i}$
 - $\circ \quad a_n = f[x_0, \dots, x_{n-1}]$
 - $o P_n(x) = f[x_0] + \sum_{i=1}^n f[x_0, \dots, x_i](x x_0) \dots (x x_{i-1})$
 - o Can also do on-the-fly memoisation
- [Osculating Polynomial] Let $x_0, x_1, ..., x_n$ be n+1 distinct numbers in [a, b] and m_i be a nonnegative integer. Suppose that $f \in C^m[a, b]$ where $m = \max m_i$. Then, the osculating polynomial approximating f is P(x) of the least degree s.t. $\frac{\mathrm{d}^k}{\mathrm{d}x^k}P(x_i) = \frac{\mathrm{d}^k f(x_i)}{\mathrm{d}x^k}$ for $i \in \{0, ..., n\}$ and $k \in \{0, ..., m\}$
- [Hermite Polynomial] A Hermite polynomial is an osculating polynomial for $m_i = 1$.
 - Let $f \in C^1[a,b]$ with $x_0, ..., x_n \in [a,b]$ distinct. Then the unique polynomial P(x) of least degree agreeing with f and f' at $x_0, ..., x_n$ is the Hermite polynomial of degree at most 2n + 1 satisfying $H(x_i) = f(x_i)$, $H'(x_i) = f'(x_i)$
 - o If $f \in C^{2n+2}[a,b]$, then $f(x) = H_{2n+1}(x) + \frac{(x-x_0)^2 \dots (x-x_n)^2}{(2n+2)!} f^{(2n+2)}(\xi(x))$ for $\xi(x) \in (a,b)$

Theorems

- [Weierstrass Approximation Theorem] Let $f \in C([a,b])$. For every $\epsilon > 0$, $\exists P(x)$ s.t. $|f(x) P(x)| < \epsilon \ \forall x \in [a,b]$
- [Lagrange Interpolation] Let $(x_i, y_i)_{i=1}^n$ be points. Then $P(x) = \sum_{i=1}^n y_i \prod_{j \neq i} \frac{x x_j}{x_i x_j}$ interpolates the points.
- [3.3] Let $x_0, ..., x_n$ be distinct points in [a, b] and $f \in C^{n+1}[a, b]$. Then $\forall x \in [a, b], \exists \xi_x$ s.t. $f(x) = P(x) + \frac{f^{(n+1)}(\xi_x)}{(n+1)!}(x x_0) ... (x x_n)$

$$0 g(t) = f(t) - P(t) - (f(x) - P(x)) \prod_{i=0}^{n} \frac{t - x_i}{x - x_i}$$

• [3.6] Suppose $f \in C^n([a,b])$ and x_0, \dots, x_n distinct in [a,b]. Then $\exists \xi \in (a,b)$ with $f[x_0, \dots, x_n] = \frac{f^{(n)}(\xi)}{n!}$

- [Forward Difference] $P_n(x) = P_n(x_0 + sh) = f(x_0 + sh) + \sum_{k=1}^n {s \choose k} \Delta^k f(x_0)$
 - o $h = x_{i+1} x_i$ (equal difference)
 - $\circ x = x_0 + sh$

Definitions (Splines)

- [Cubic Spline Interpolant] Let f be defined on [a, b]
 - $S_i(x)$: cubic polynomial on $[x_i, x_{i+1}]$ for $j \in \{0, ..., n-1\}$
 - $\circ S_i(x_i) = f(x_i)$
 - $\circ S_i(x_{i+1}) = f(x_{i+1})$
 - o $S'_{i+1}(x_{i+1}) = S'_i(x_{i+1})$
 - $\circ S_{i+1}''(x_{i+1}) = S_i''(x_{i+1})$
 - o [Natural] $S''(x_0) = S''(x_n) = 0$
 - o [Clamped] $S'(x_0) = f'(x_0)$ and $S'(x_n) = f'(x_n)$
- [Guidepoint] Helps to specify the derivative at an endpoint
 - For (x_0, y_0) , the guidepoint is $(x_0 + \alpha_0, y_0 + \beta_0)$
 - For (x_1, y_1) , the guidepoint is $(x_1 \alpha_1, y_1 \beta_1)$
 - Cubic Hermite polynomial satisfies $x'(0) = \alpha_0$, $x'(1) = \alpha_1$, $y'(0) = \beta_0$, $y'(1) = \beta_1$
 - $(x_1) = [2(x_0 x_1) + (\alpha_0 + \alpha_1)]t^3 + [3(x_1 x_0) (\alpha_1 + 2\alpha_0)]t^2 + \alpha_0 t + x_0$
 - $(y(t) = [2(y_0 y_1) + (\beta_0 + \beta_1)]t^3 + [3(y_1 y_0) (\beta_1 + 2\beta_0)]t^2 + \beta_0 t + y_0$
- [Cubic Bezier Polynomial] Given n+1 data points $(x(t_0), y(t_0)), ..., (x(t_n), y(t_n))$ and $\frac{\mathrm{d}y}{\mathrm{d}x}|_{t_i}$, find 2n cubic Hermite polynomials satisfying: $x_i(t) = x(t), y_i(t) = y(t)$.

 - $y(t) = [2(y_0 y_1) + 3(\beta_0 + \beta_1)]t^3 + [3(y_1 y_0) 3(\beta_1 + 2\beta_0)]t^2 + 3\beta_0 t + y_0$
 - \circ $t \in [0,1]$
- [Bezier Curve Algorithm]
 - o Input: $n, (x_0, y_0), ..., (x_n, y_n), (x_0^+, y_0^+), ..., (x_{n-1}^+, y_{n-1}^+), (x_1^-, y_1^-), ..., (x_n^-, y_n^-)$
 - $\qquad \qquad \circ \quad \text{Output:} \left\{ a_0^{(i)}, a_1^{(i)}, a_2^{(i)}, a_3^{(i)}, b_0^{(i)}, b_1^{(i)}, b_2^{(i)}, b_3^{(i)} \right\} \\$

Differentiation

Forward Difference Formula

- [Forward Difference Formula] $f'(x_0) = \frac{f(x_0 + h) f(x_0)}{h}$ where h > 0
- [Backward Difference Formula] $f'(x_0) = \frac{f(x_0 + h) f(x_0)}{h}$ where h < 0
- [Error Formula] $\left| \left(f(x_0 + h) \frac{f(x_0)}{h} \right) f'(x_0) \right| = \left| \frac{h}{2} f''(\xi) \right|$ for $\xi \in (x_0, x_0 + h)$
- [Three Point Formula] $O(h^2)$

$$L_0(x) = \frac{(x-x_1)(x-x_2)}{(x_0-x_1)(x_0-x_2)}, L'_0(x) = \frac{2x-x_1-x_2}{(x_0-x_1)(x_0-x_2)}$$

- o [Three-Point Endpoint Formula]
 - $f'(x_0) = \frac{1}{2h} \left[-3f(x_0) + 4f(x_0 + h) f(x_0 + 2h) \right] + \frac{h^2}{3} f^{(3)}(\xi_0)$
- $\xi_0 \in (x_0, x_0 + 2h)$ [Three-Point Midpoint Formula]

•
$$f'(x_0) = \frac{1}{2h} [-f(x_0 - h) + f(x_0 + h)] - \frac{h^2}{6} f^{(3)}(\xi_1)$$

- $\xi_1 \in (x_0 h, x_0 + h)$
- [Five Point Formula] $O(h^4)$
 - [Five-Point Midpoint Formula]

$$f'(x_0) = \frac{1}{12h} \left(f(x_0 - 2h) - 8f(x_0 - h) + 8f(x_0 + h) - f(x_0 + 2h) \right) + \frac{h^4}{30} f^{(5)}(\xi)$$

- $\xi \in (x_0 2h, x_0 + 2h)$ [Five-Point Endpoint Formula]

$$f'(x_0) = \frac{1}{12h} \left(-25f(x_0) + 48f(x_0 + h) - 36f(x_0 + 2h) + 16f(x_0 + 3h) - 3f(x_0 + 4h) \right) + \frac{h^4}{5} f^{(5)}(\xi)$$

- $\bullet \quad \xi \in (x_0, x_0 + 4h)$
- [Second Derivative Midpoint Formula] $O(h^2)$ if $f^{(4)}$ bounded

$$o f''(x_0) = \frac{1}{h^2} (f(x_0 - h) - 2f(x_0) + f(x_0 + h)) - \frac{h^2}{12} f^{(4)}(\xi)$$

- $\circ \quad \xi \in (x_0 h, x_0 + h)$
- [Round off Error] $\frac{\xi}{h} + \frac{h^2}{\epsilon M}$ where M is a bound for |f|

Extrapolation

- Combine inaccurate O(h) approximations to get formulas with higher order error
- M: estimand, $N_1(h)$: estimator

o
$$M = N_1(h) + K_1h + K_2h^2 + K_3h^3 + \cdots$$
 where K_i are constants

$$0 \quad M = N_1 \left(\frac{h}{2}\right) + K_1 \frac{h}{2} + K_2 \left(\frac{h}{2}\right)^2 + K_3 \left(\frac{h}{2}\right)^3 + \cdots$$

[Normal Extrapolation]

$$O(N_2(h) = 2N_1(\frac{h}{2}) - N_1(h)$$

$$\circ N_{j+1}(h) = \frac{2^{j}}{2^{j}-1} N_{j}\left(\frac{h}{2}\right) - \frac{1}{2^{j}-1} N_{j}(h), M - N_{j}(h) = O(h^{j})$$

• [Even Power Extrapolation] $N_{j+1}(h) = N_j \left(\frac{h}{2}\right) + \frac{N_j \left(\frac{h}{2}\right) - N_j(h)}{4^{j-1}}$ where $M - N_{j+1}(h) = O\left(h^{2(j+1)}\right)$

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	O(h)	$O(h^2)$	$O(h^3)$	$O(h^4)$	
	$N_1(h) \searrow$				
	$N_1(\frac{h}{2})$	$N_2(h)$			
	$N_1(\frac{h}{4})\stackrel{ ightarrow}{\searrow}$	$N_2(\frac{h}{2}) \stackrel{\rightarrow}{\searrow}$	$N_3(h)$		
[Richardson Extrapolation	$N_1(rac{h}{8}) ightarrow$	$N_2(\frac{\overline{h}}{4}) \rightarrow$	$N_3(\frac{h}{2}) \rightarrow$	$N_4(h)$	
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Integration

Definition

[Degree of Precision] The degree of precision of a quadrature formula is the largest positive integer n s.t. the formula is exact for x^k for $k \in \{0,1,...,n\}$

Method of Quadrature

- [Quadrature] Select set of distinct nodes $x_0 < \cdots < x_n$ from interval [a, b] and integrate the
 - Lagrange interpolating polynomial $P_n(x) = \sum_{i=0}^n f(x_i) L_i(x)$ $\int_a^b f(x) dx = \sum_{i=0}^n a_i f(x_i) + \frac{1}{(n+1)!} \int_a^b \prod_{i=0}^n (x-x_i) f^{(n+1)} (\xi(x)) dx \text{ where } a_i = 0$ $\int_a^b L_i(x) dx$
- [Trapezoidal Rule] $x_0 = a, x_1 = b, h = b a$

$$\circ \int_a^b f(x) \, \mathrm{d}x = \frac{h}{2} \Big(f(x_0) + f(x_1) \Big) - \frac{h^3}{12} f''(\xi)$$

- Degree of precision: 1
- [Simpson Rule] $x_0 = a, x_1 = \frac{a+b}{2}, x_2 = b, h = \frac{b-a}{2}$

$$\circ \int_a^b f(x) \, \mathrm{d}x = \frac{h}{3} \Big(f(x_0) + 4f(x_1) + f(x_2) \Big) - \frac{h^5}{90} f^{(4)}(\xi)$$

- Degree of precision: 3
- [Closed Newton-Cotes Formula] $x_0 = a$, $x_n = b$, $h = \frac{b-a}{n}$

o
$$n \text{ even: } \int_a^b f(x) \, \mathrm{d}x = \sum_{i=0}^n a_i f(x_i) + \frac{h^{n+3} f^{(n+2)}(\xi)}{(n+2)!} \int_0^n t^2 (t-1) \dots (t-n) \, \mathrm{d}t$$

o
$$n \text{ odd: } \int_a^b f(x) dx = \sum_{i=0}^n a_i f(x_i) + \frac{h^{n+2} f^{(n+1)}(\xi)}{(n+1)!} \int_0^n t(t-1) \dots (t-n) dt$$

[Open Newton-Cotes Formula]

Composite Methods

[Composite Simpson's Rule] Let $f \in C^4[a,b]$, n even, $h = \frac{b-a}{n}$ and $x_j = a + jh$. Then $\exists \xi \in C^4[a,b]$ (a,b) s.t.

$$\circ \int_{a}^{b} f(x) \, \mathrm{d}x = \frac{h}{3} \left(f(a) + 2 \sum_{i=1}^{\frac{n}{2}-1} f(x_{2i}) + 4 \sum_{i=1}^{\frac{n}{2}} f(x_{2i-1}) + f(b) \right) - \frac{b-a}{180} h^4 f^{(4)}(\xi)$$

[Composite Trapezoidal Rule]

$$\int_{a}^{b} f(x) dx = \frac{h}{2} \left(f(a) + 2 \sum_{i=1}^{n-1} f(x_i) + f(b) \right) - \frac{b-a}{12} h^2 f''(\xi)$$

[Composite Midpoint Rule]

$$\int_{a}^{b} f(x) dx = 2h \sum_{i=0}^{\frac{n}{2}} f(x_{2i}) + \frac{b-a}{6} h^{2} f''(\xi)$$

[Round Off Error Stability] Round off error does not depend on number of calculations performed i.e. independent of composite integration techniques and n

Romberg Integration

- $R_{n,k}$: n+1 number of points subdividing the interval, k
- $R_{k,j} = R_{k,j-1} + \frac{1}{4^{j-1}-1} (R_{k,j-1} R_{k-1,j-1})$ for k = j, j+1, ...
- [Algorithm]

$$R_{1,1} = \frac{h}{2} (f(a) + f(b))$$

o
$$R_{n,1} = \frac{1}{2} (R_{n-1,1} + 2h_n \sum_{j=1}^{2^{n-2}} f(a + (k-0.5)h))$$
 where $h_n = \frac{b-a}{2^{n-1}}$

- $R_{n,1}$ is just dividing the interval into 2^{n-1} pieces and use composite method
- o $R_{n,i} = R_{n,i-1} + \frac{R_{n,i-1} R_{1,i-1}}{4^{i-1} 1}$ Other Romberg terms come for free

 - For Simpson's, $R_{n,i} = I + O(h_i^{2i+2})$

$$\begin{array}{c|ccccc} O(h_k^4) & O(h_k^6) & O(h_k^8) & O(h_k^{10}) \\ \hline R_{1,1} & & & & \\ R_{2,1} & & R_{2,2} & & & \\ \hline R_{3,1} & & R_{3,2} & & R_{3,3} & & \\ R_{4,1} \rightarrow & R_{4,2} \rightarrow & R_{4,3} \rightarrow & R_{4,4} \end{array}$$