

# A Dynamical Point of View on Quantum Optimization

**Jiaqi (Jimmy) Leng**

Applied Mathematics & Statistics, and Scientific Computing (AMSC) Program  
Joint Center for Quantum Information & Computer Science  
University of Maryland, College Park

“Quantum Hamiltonian Descent”, arXiv:2303.01471

“A quantum-classical performance separation in nonconvex optimization”, manuscript in preparation



JOINT CENTER FOR  
QUANTUM INFORMATION  
AND COMPUTER SCIENCE

# Quantum computing 101

## Classical (digital) Computer

- Unit of information: bit ( $b \in \{0,1\}$ )
- Classical information  $\rightarrow$  bitstrings (e.g., 10001001)
- Computation = manipulation of finite-length bitstrings
- Readout/Return: bitstrings

## Quantum Computer

- Unit of information: quantum bit, or **qubit**  
 $|\psi\rangle = a|0\rangle + b|1\rangle, a, b \in \mathbb{C}, |a|^2 + |b|^2 = 1$
- Quantum information  $\rightarrow$  quantum state (i.e., superposition of  $|bitstring\rangle$ )

$$|GHZ\rangle = \frac{1}{\sqrt{2}} (|000\rangle + |111\rangle)$$

- Computation = unitary (i.e., preserving 2-norm) operations on quantum states
- Readout: quantum measurement  $\rightarrow$  a sample of bitstrings

Measure GHZ state:  
000: 50%, 111: 50%, others: 0%

### Digital Quantum Computers

Use quantum gates & quantum circuits, not ready in 5-10 years.  
Must do error correction.

### Analog Quantum Computers

Solving problems by emulating a real quantum system. Ready & easy to scale. No error correction.

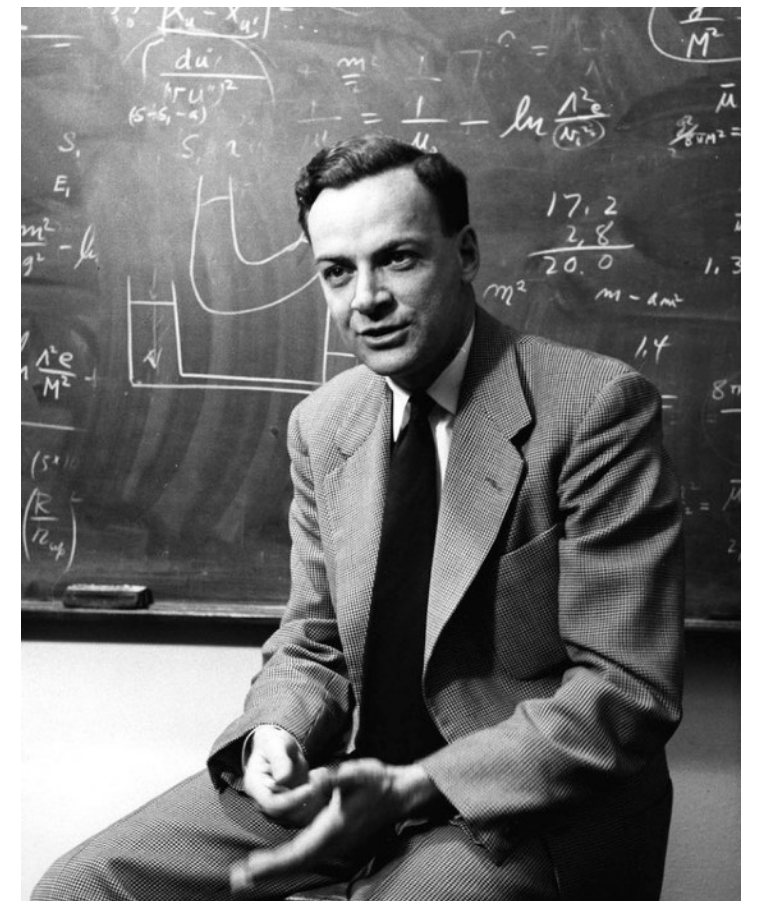
# Hamiltonian simulation

- **Quantum simulation** (or **Hamiltonian simulation**): a prominent application of quantum computers. **Exponential** quantum speedup expected.

$$i\frac{\partial}{\partial t}\psi = \hat{H}(t)\psi$$

$\hat{H}(t)$ : quantum Hamiltonian (Hermitian/self-adjoint)

- **Feynman** (“*Simulating Physics with Computers*”, 1982): to simulate quantum systems, we would need to build quantum computers.
- **Applications**: quantum chemistry, quantum field theory, condensed matter physics, numerical optimization, etc.



# Overview

- Part I: Quantum Hamiltonian Descent (QHD)
- Part II: A Quantum-Classical Performance Separation
- Part III: Implementation on Quantum Computers

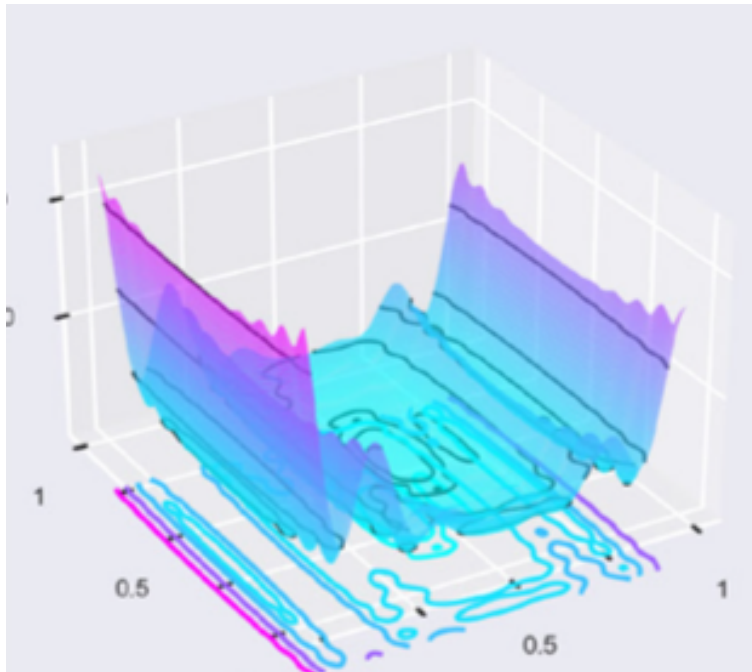
# Part I:

## Quantum Hamiltonian Descent

# Problem formulation

## Continuous optimization

$$\min_{x \in \mathbb{R}^d} f(x)$$



- Important in practice: machine learning, operations research, scientific computing, etc.
- A **challenging problem** for quantum: different nature in (quantum) algorithm design, requires new mathematical tools to prove convergence.
- **Opportunity:** new *primitives* of quantum speedups!

## Classical algorithms: Gradient Descent (GD)

- Standard GD:  $x_{k+1} = x_k - s \nabla f(x_k)$ .
- Nesterov's accelerated GD:  $x_{k+1} = y_k - s \nabla f(y_k)$ ,  $y_{k+1} = x_{k+1} + \frac{k}{k+3} (x_{k+1} - x_k)$ .

# A Lagrangian formulation of accelerated methods

Su, Boyd, & Candes  
NeurIPS 2014

$$\begin{aligned}x_k &= y_{k-1} - s \nabla f(y_{k-1}), \\ y_k &= x_k + \frac{k-1}{k+2}(x_k - x_{k-1}).\end{aligned}$$

Nesterov's accelerated GD



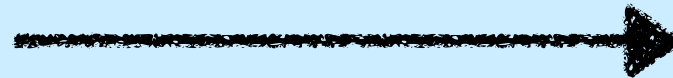
$$\ddot{X} + \frac{3}{t}\dot{X} + \nabla f(X) = 0,$$

ODE model

Wibisono, Wilson, & Jordan  
PNAS Nov 2016, 113 (47)

$$\begin{aligned}x_k &= y_{k-1} - s \nabla f(y_{k-1}), \\ y_k &= x_k + \frac{k-1}{k+2}(x_k - x_{k-1}).\end{aligned}$$

Nesterov's accelerated GD



$$\mathcal{L}(t, \dot{X}, X) = t^3 \left( \frac{1}{2} |\dot{X}|^2 - f(X) \right)$$

Lagrangian formulation

$$H(t, X, P) = \frac{1}{2t^3} |P|^2 + t^3 f(X)$$

Hamiltonian formulation

- Accelerated GD can be modeled by [accelerated gradient flows](#).
- GD algorithms can be generated by discretizing continuous-time dynamics.

# Accelerated Hamiltonian flows: classical v.s. quantum

We can make the *classical* dynamics *quantum*!!

**Classical  
Hamiltonian  
Systems**

*Path integral formulation of  
Quantum Mechanics*

Due to Feynman

**Quantum  
Hamiltonian  
Evolution**

$$H(t, X, P) = \frac{1}{2t^3} |P|^2 + t^3 f(X)$$



$$\hat{H}(t) = \frac{1}{t^3} \left( -\frac{1}{2} \Delta \right) + t^3 f(x)$$

**Bregman Lagrangian**

$$\mathcal{L}(X, \dot{X}, t) = e^{-\varphi_t} \left( \frac{1}{2} |\dot{X}|^2 \right) - e^{\chi_t} f(X)$$

**Quantum propagator**

$$K(b, t_b; a, t_a) \propto \sum_{\text{paths from a to b}} e^{\frac{i}{\hbar} S[X(t)]}$$

**Schrodinger equation (QHD)**

$$i\hbar \frac{\partial \Psi}{\partial t} = \left[ e^{\varphi_t} \left( -\frac{\hbar^2}{2} \nabla^2 \right) + e^{\chi_t} f(x) \right] \Psi$$

**Infinitesimal expansion**

$$\Psi(x, t + \epsilon) \propto \int \exp \left\{ \frac{i}{\hbar} \epsilon \mathcal{L} \left( \frac{x+y}{2}, \frac{x-y}{\epsilon}, t \right) \right\} \Psi(y, t) dy$$



# Quantum Hamiltonian Descent

$$H(t, X, P) = e^{\varphi_t} \left( \frac{1}{2} |P|^2 \right) + e^{\chi_t} f(X)$$

Accelerated gradient flows

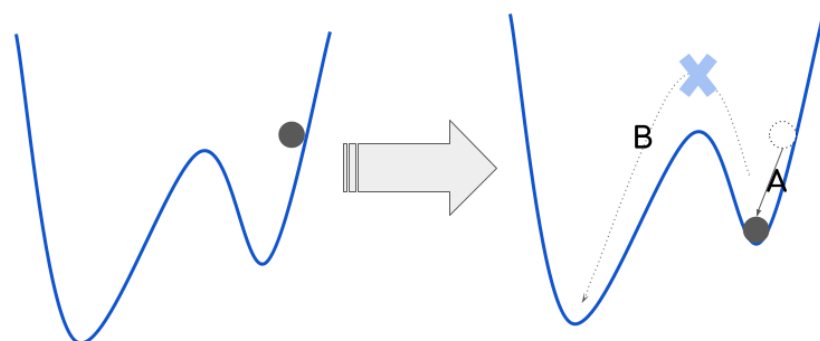


$$\hat{H}(t) = e^{\varphi_t} \left( -\frac{1}{2} \nabla^2 \right) + e^{\chi_t} f(x)$$

Quantum Hamiltonian Descent (QHD)

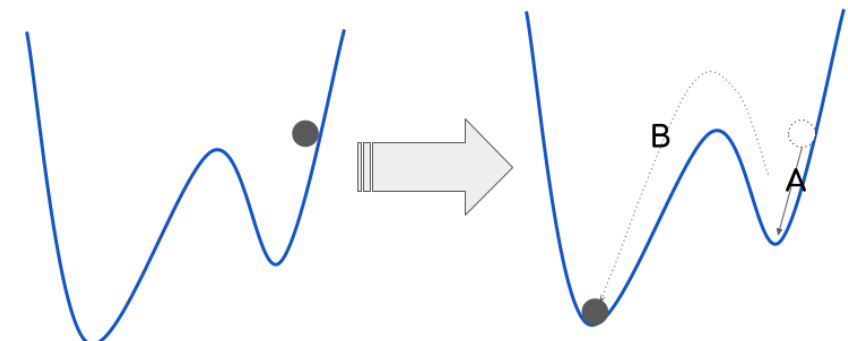
- This is a quantum dynamical system with *damping* ( $e^{\varphi_t - \chi_t} \rightarrow 0$ ).
- The energy damping allows the quantum system to *converge* to a low-energy configuration (i.e., minimizing the objective  $f$ ).
- **Intuition:** “Path integral” of classical GD  $\rightarrow$  does it help **nonconvex** optimization?

Classical Gradient methods



Optimization paths = A  
B is prohibited by classical mechanics.

Quantum Hamiltonian Descent (QHD)



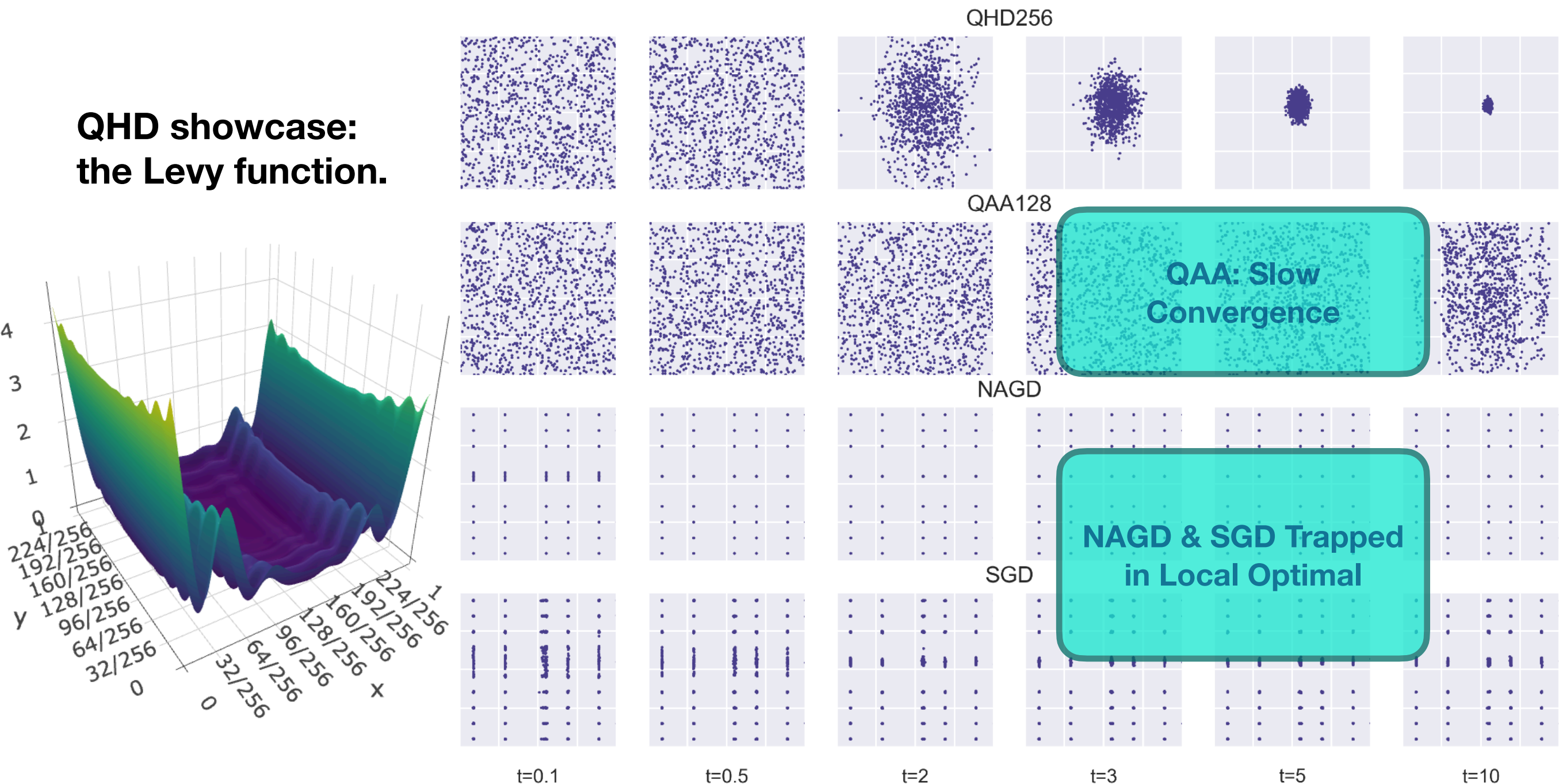
Optimization paths =  $\int_{\text{all possible paths}}$   
A, B both contribute to quantum optimization paths.

# Quantum Hamiltonian Descent - numerical example

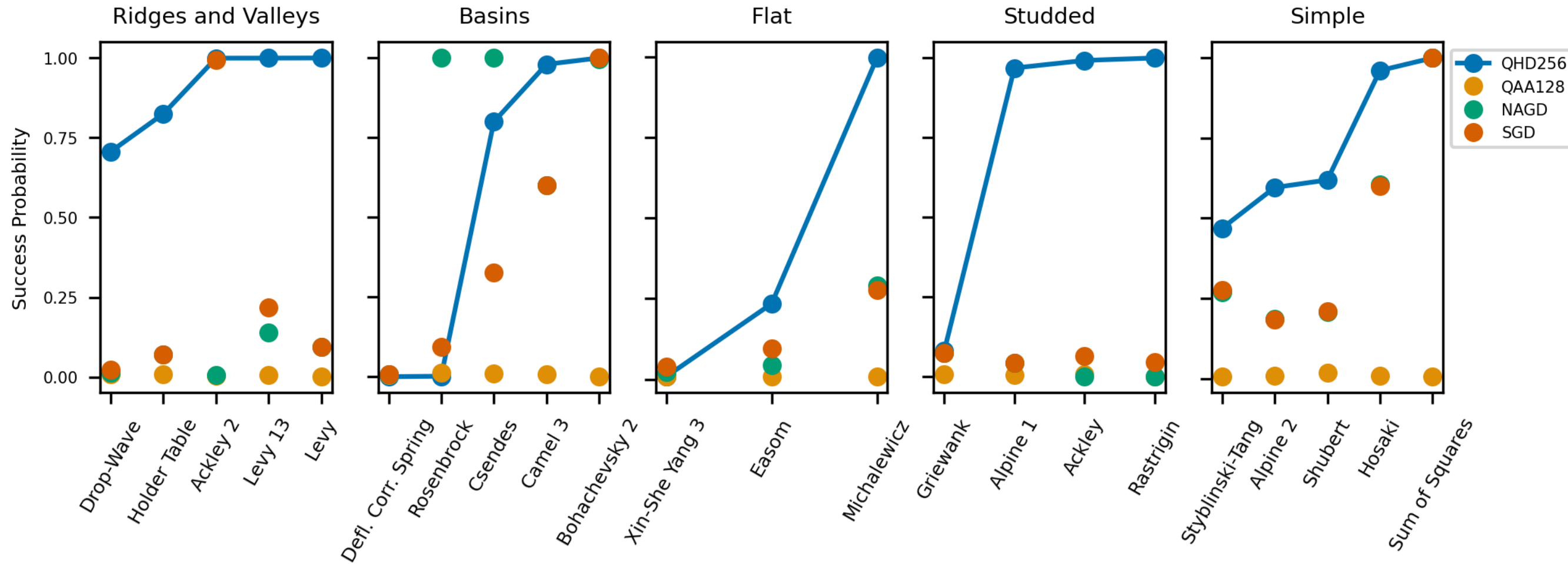
**QAA:** quantum adiabatic algorithm. **NAGD:** Nesterov's accelerated gradient descent.

**SGD:** stochastic gradient descent.

**QHD showcase:  
the Levy function.**



# Behavior of QHD on 2D benchmark



## Measure in Success Probability for some given time:

the probability of generating a solution within radius  $r$  of the optimal solution from random initial points.

**QAA:** use standard *linear interpolation* in the adiabatic method (i.e.,  $H_T(t) = (1 - t/T)H_0 + (t/T)H_1$ ) and simulate  $H_T(t)$  for  $T=10$  over a **128-by-128** grid on  $[0,1] \times [0,1]$ .

**QHD:** simulate QHD Hamiltonian with Nesterov's AGD parameters (i.e.,  $H(t) = -\frac{1}{2t^3}\nabla^2 + t^3f(x)$ ) from  $t = 0$  to 10 over a **256-by-256** grid on  $[0,1] \times [0,1]$ .

# Convergence guarantee: convex optimization

$$H(t, P, X) = e^{\alpha_t} \left( -\frac{e^{-\gamma_t}}{2} P^2 + e^{\gamma_t + \beta_t} f(X) \right) \longrightarrow \frac{d}{dt} (e^{-\alpha_t} \dot{X}_t + X_t) = e^{\alpha_t + \beta_t} \nabla f(X_t)$$

**Theorem** [Wibisono, Wilson, Jordan (2016)]

Assume  $f(x)$  is continuously differentiable and convex, and  $\dot{\beta}_t \leq e^{\alpha_t}$ ,  $\dot{\gamma}_t = e^{\alpha_t}$  (aka, *ideal scaling condition*). Then, the solution  $X(t)$  satisfies

$$f(X_t) - f(x^*) \leq O(e^{-\beta_t}).$$

$$\hat{H}(t) = e^{\alpha_t} \left( -\frac{e^{-\gamma_t}}{2} \Delta + e^{\gamma_t + \beta_t} f(x) \right) \longrightarrow \frac{\partial}{\partial t} \Psi(t, x) = \hat{H}(t) \Psi(t, x)$$

**Theorem** [Leng, Hickman, Li, Wu (2023)]

Assume  $f(x)$  is continuously differentiable and convex, and  $\dot{\beta}_t \leq e^{\alpha_t}$ ,  $\dot{\gamma}_t = e^{\alpha_t}$  (aka, *ideal scaling condition*). Define  $\mathbb{E}[f]_{\sim \Psi(t)} = \langle \Psi(t) | f | \Psi(t) \rangle = \int f | \Psi(t) |^2 dx$ , then we have

$$\mathbb{E}[f]_{\sim \Psi(t)} - f(x^*) \leq O(e^{-\beta_t}).$$

# A Lyapunov function approach

## Theorem [Leng, Hickman, Li, Wu (2023)]

Assume  $f(x)$  is continuously differentiable and convex, and  $\dot{\beta}_t \leq e^{\alpha_t}$ ,  $\dot{\gamma}_t = e^{\alpha_t}$  (aka, ideal scaling condition). Define  $\mathbb{E}[f]_{\sim \Psi(t)} = \langle \Psi(t) | f | \Psi(t) \rangle = \int f | \Psi(t) |^2 dx$ , then we have

$$\mathbb{E}[f]_{\sim \Psi(t)} - f(x^*) \leq O(e^{-\beta_t}).$$

- We construct a Lyapunov function ( $\langle O \rangle_t = \langle \Psi_t | O | \Psi_t \rangle$ ,  $\hat{p} = -i \nabla$ ):

$$\mathcal{W}(t) = \langle \hat{J}^2 / 2 \rangle_t + e^{\beta_t} \langle f \rangle_t$$

$$\hat{J} := e^{-\gamma_t} \hat{p} + \hat{x}.$$

- We can prove this Lyapunov function is **non-increasing in t**.
- Therefore, we have  $e^{\beta_t} \mathbb{E}[f]_{\sim \Psi_t} \leq \mathcal{W}(t) \leq \mathcal{W}(0)$ ,  $\mathbb{E}[f]_{\sim \Psi_t} \leq \mathcal{W}(0) e^{-\beta_t} \leq O(e^{-\beta_t})$ .

- A sanity check: we have the same converge rate as classical [WWJ16].
- Nesterov's GD:  $e^{\alpha_t} = 2/t$ ,  $e^{\beta_t} = e^{\gamma_t} = t^2$ , implies the convergence rate  $O(t^{-2})$ .
- QHD is more stable & robust with larger time discretization steps.
- Question: can we achieve  $O(e^{-\sqrt{\beta}t})$  convergence rate for  $\beta$ -strongly convex  $f$ ?

# Convergence guarantee: nonconvex optimization

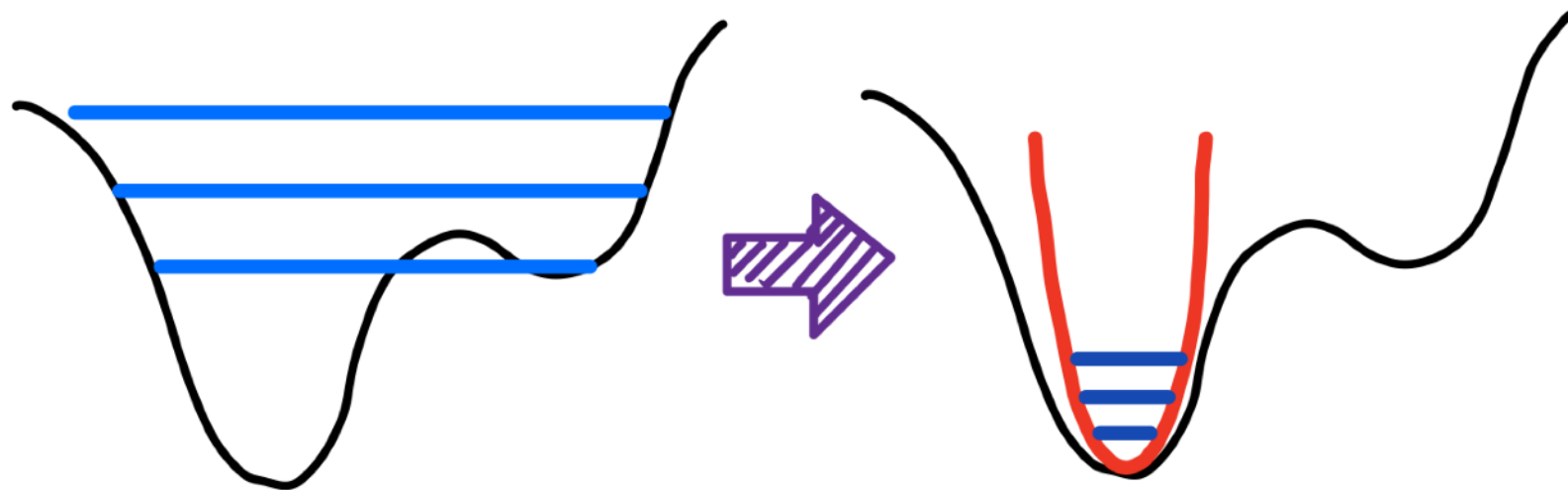
$$\hat{H}(t) = \left( -\frac{e^{\varphi_t}}{2} \Delta + e^{\chi_t} f(x) \right) \longrightarrow \frac{\partial}{\partial t} \Psi(t, x) = \hat{H}(t) \Psi(t, x)$$

## Theorem (informal) [Leng, Hickman, Li, Wu (2023)]

Suppose  $f(x)$  be smooth, unbounded at infinity, and has a unique non-degenerate global minimum  $x^*$ . Let the initial wave  $\Psi_0$  be in the low-energy subspace of  $H(0)$  and  $e^{\varphi_t - \chi_t} \rightarrow 0$ ,  $|\dot{\varphi}_t|, |\dot{\chi}_t| \ll 1$  (i.e.,  $H(t)$  is slow-varying), then

$$\lim_{t \rightarrow \infty} \mathbb{E}[f]_{\sim \Psi_t} = f(x^*).$$

- Slow-varying  $H(t)$   $\rightarrow$  the wave function  $\Psi(t)$  stays in the low-energy subspace for all  $t$ .
- The low-energy subspace of  $H(t)$  will *migrate* to the global minimum of  $f$  as  $t \rightarrow \infty$ .





# Quantum Hamiltonian Descent: complexity analysis

1. Prepare an initial state  $|\psi_0\rangle$ .
2. Simulating the Schrodinger equation:  $i\partial_t\psi = \hat{H}(t)\psi$ ,  $\psi(0) = \psi_0$ .
3. At time  $t = T$ , *measure* the final state  $|\Psi_T\rangle$  (i.e., to sample from the corresponding distribution  $|\Psi(T)|^2$ ).
4. Ideally, the measurement results will cluster around the global minimizer of  $f$ .

$$\hat{H}(t) = e^{\varphi_t} \left( -\frac{1}{2} \nabla^2 \right) + e^{\chi_t} f(x)$$

## Runtime of QHD

**QHD is comparable to GD in terms of **simplicity** and **resource cost**:**

- Running QHD = time-dependent Hamiltonian simulation [Childs, Leng, Li, Liu & Zhang, 2022]
- The runtime of QHD is  $\tilde{O}(dT)$  ( $d$ : dimension of  $f$ ,  $T$ : total evolution time).
- Resource cost (on quantum computer) is **comparable** to that of classical GD.
- Resource cost (on classical computer) is **exponential** in  $d$ .

Therefore, *QHD is a quantum-upgraded version of GD*. We will demonstrate how powerful QHD could be by itself. One could also build on top of QHD like what we've done with GD.

# QHD as accelerated Wasserstein gradient flow

Let's consider a simple QHD evolution:

$$i \frac{\partial}{\partial t} \psi(t, x) = \left[ e^{-\alpha t} \left( -\frac{1}{2} \Delta \right) + e^{\alpha t} f \right] \psi(t, x)$$

We apply the Madelung transform:  $\psi(t, x) = \sqrt{\rho(t, x)} e^{iS(t, x)}$

$$\begin{cases} \partial_t \rho_t + e^{-\alpha t} \nabla \cdot (\rho_t \nabla S_t) = 0, \\ \partial_t S_t + e^{-\alpha t} \left[ \frac{1}{2} (\nabla S_t)^2 + \frac{1}{8} \frac{\delta \mathcal{F}(\rho_t)}{\delta \rho_t} \right] + e^{\alpha t} f(x) = 0 \end{cases}$$

$\mathcal{F}(\rho) = \int |\nabla \log \rho|^2 \rho$  is the *Fisher information functional*

The above system of PDEs is the Euler-Lagrange equation generated by:

$$\mathcal{L}(\rho, \sigma, t) = e^{\alpha t} \left( \frac{1}{2} g_W(\sigma, \sigma) - \int f \rho \, dx \right) - e^{-\alpha t} \left( \frac{1}{8} \mathcal{F}(\rho) \right)$$

- QHD = accelerated Wasserstein gradient flow with *Fisher regularization*
- This explains why QHD works well for **nonconvex** problems (Fisher information is convex, see [Li, Lu, Wang (2020)]).
- Question: can we use this interpretation to prove faster convergence rate?

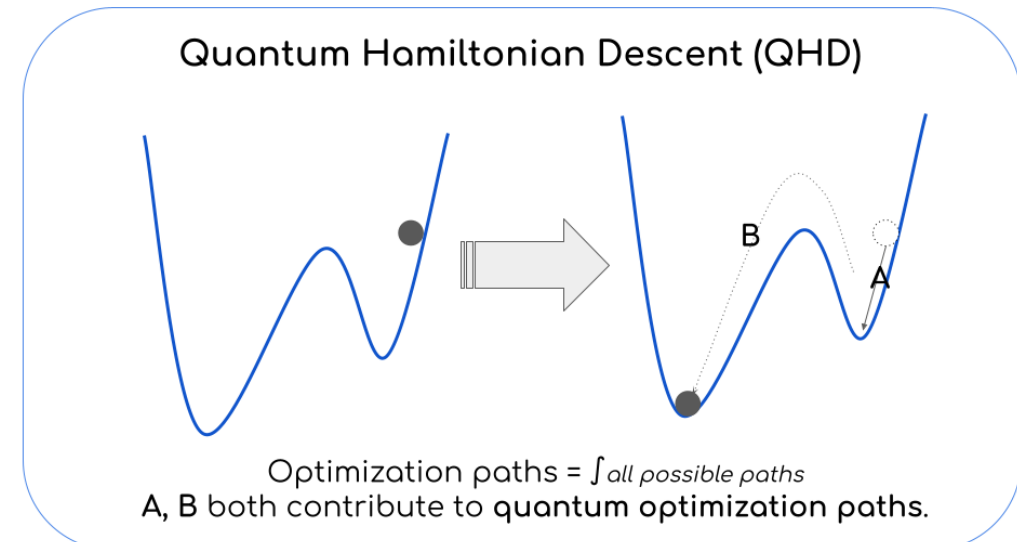
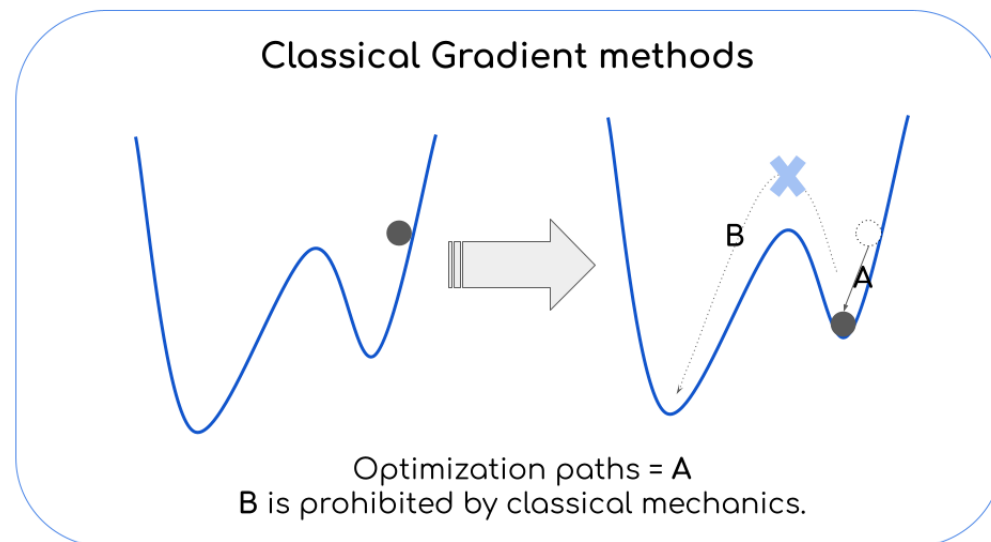


# Part II:

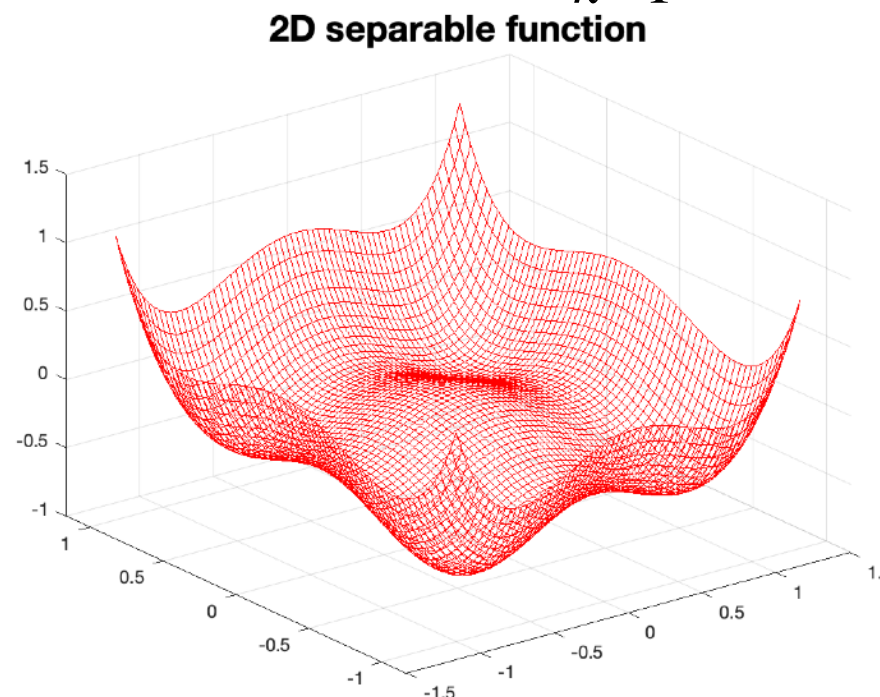
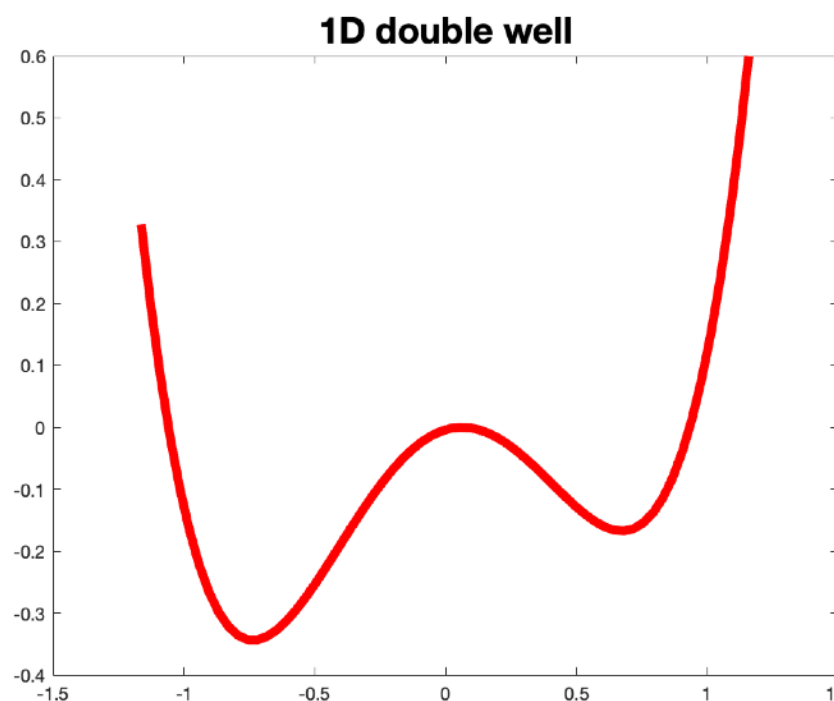
## A Quantum-Classical Performance Separation

# Quantum Hamiltonian Descent — practical setting?

1-dim nonconvex model problem: double-well potential —  $f(x)$



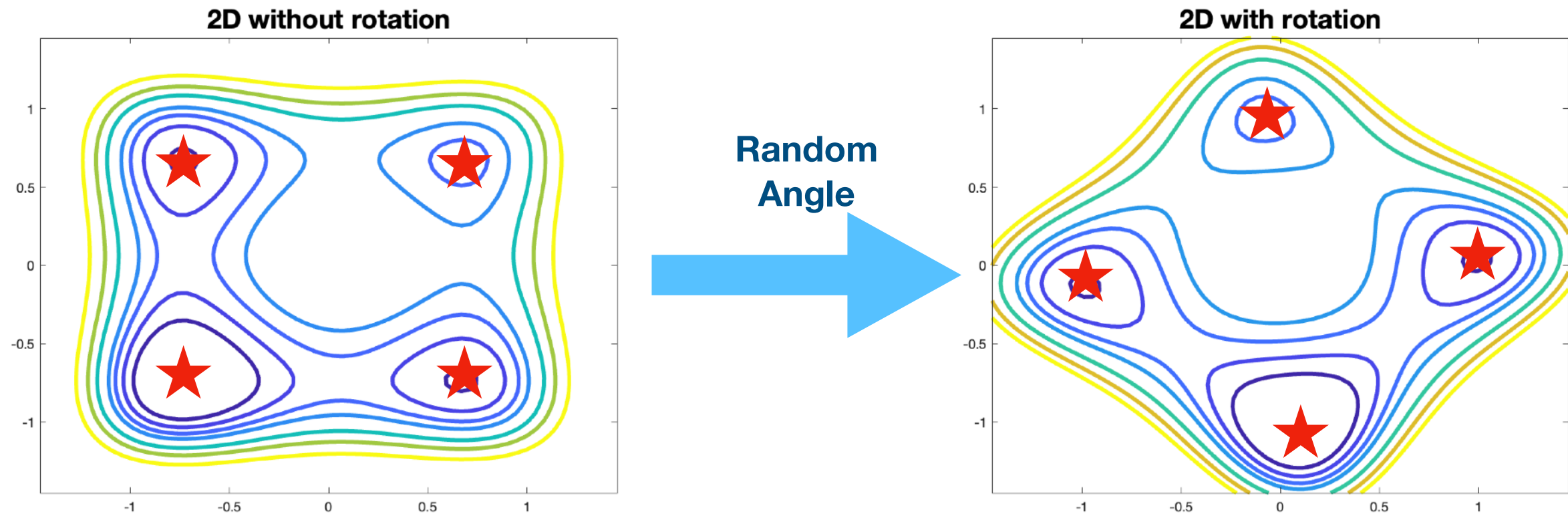
d-dim objective function:  $F(x_1, \dots, x_d) = \sum_{k=1}^d f(x_k),$



- $F(x)$  has  $2^d$  local minima! (Only **one unique** global minimum).
- $F(x)$  is **separable**: not difficult for classical algorithms if the closed-form formula is given.

# Construction of the optimization instances

Our instances = d-dim separable functions + random rotation



$$F(x_1, \dots, x_d) = \sum_{k=1}^d f(x_k) \mapsto F_U(x) = F(Ux)$$

- $F_U(x)$  still has  $2^d$  local minima, with a **unique** global minimum.
- $F(x)$  is **non-separable**: difficult to recover the rotation even with the closed-form formula!

# QHD: a polynomial-time quantum algorithm

Given an optimization problem  $f(x): \mathbb{R}^d \rightarrow \mathbb{R}$  with a unique global minimizer  $x^*$ .

$x$  is a  **$\delta$ -approximate** solution if  $\|x - x^*\| < \delta$ .

## **Theorem (Informal) [Leng, Zheng, Wu (2023)]**

Let  $f(x): \mathbb{R} \rightarrow \mathbb{R}$  be a double-well potential function. Define  $F_U(x) = F(Ux)$  where

$F(x) = \sum_{k=1}^d f(x_k)$  and  $U$  is an arbitrary orthogonal matrix. For any small  $\delta > 0$ , QHD can

produce a  $\delta$ -approximate solution with probability at least  $2/3$  using

- $\tilde{\mathcal{O}}(d^3/\delta^2)$  quantum queries to  $F_U$ , and
- $\tilde{\mathcal{O}}(d^4/\delta^2)$  additional 1- and 2-qubit gates.

Manuscript in preparation.

- The QHD Hamiltonian (more precisely, the Laplacian operator) is rotationally invariant.
- The ground state of the QHD Hamiltonian is the *vehicle* of quantum optimization.
- We use an adiabatic theorem for unbounded Hamiltonian.

# A quantum-classical performance separation

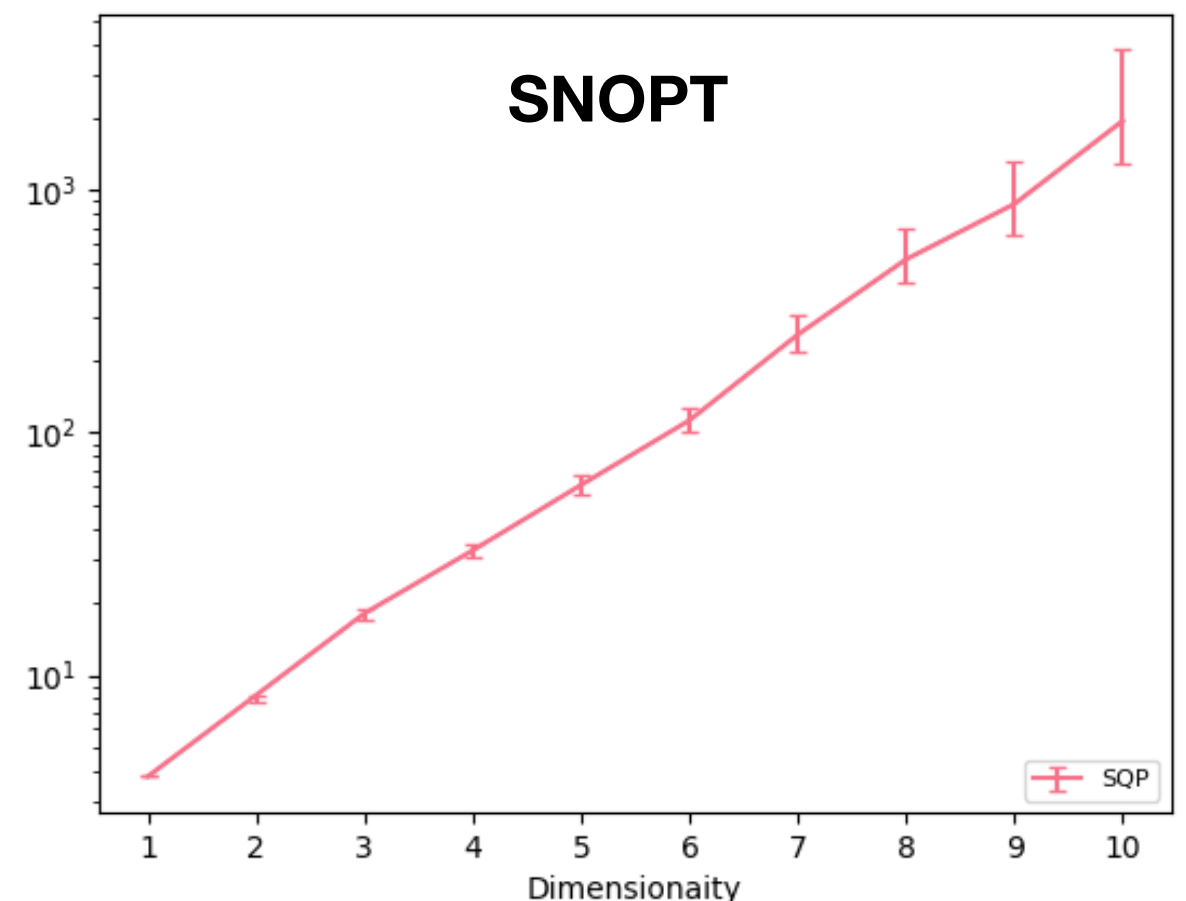
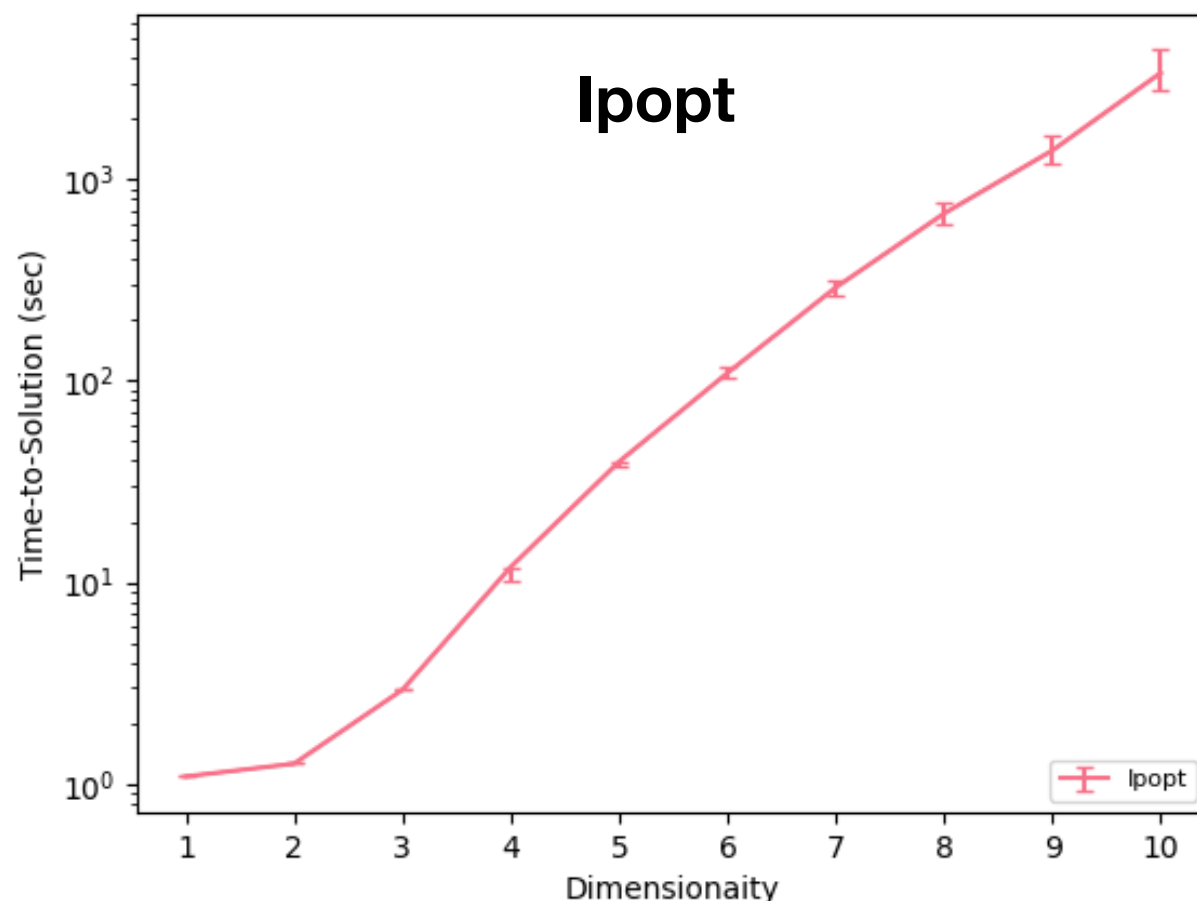
## Time-To-Solution (TTS)

$$\text{TTS} = t_f \left\lceil \frac{\ln(1 - 0.99)}{\ln(1 - P_g)} \right\rceil$$

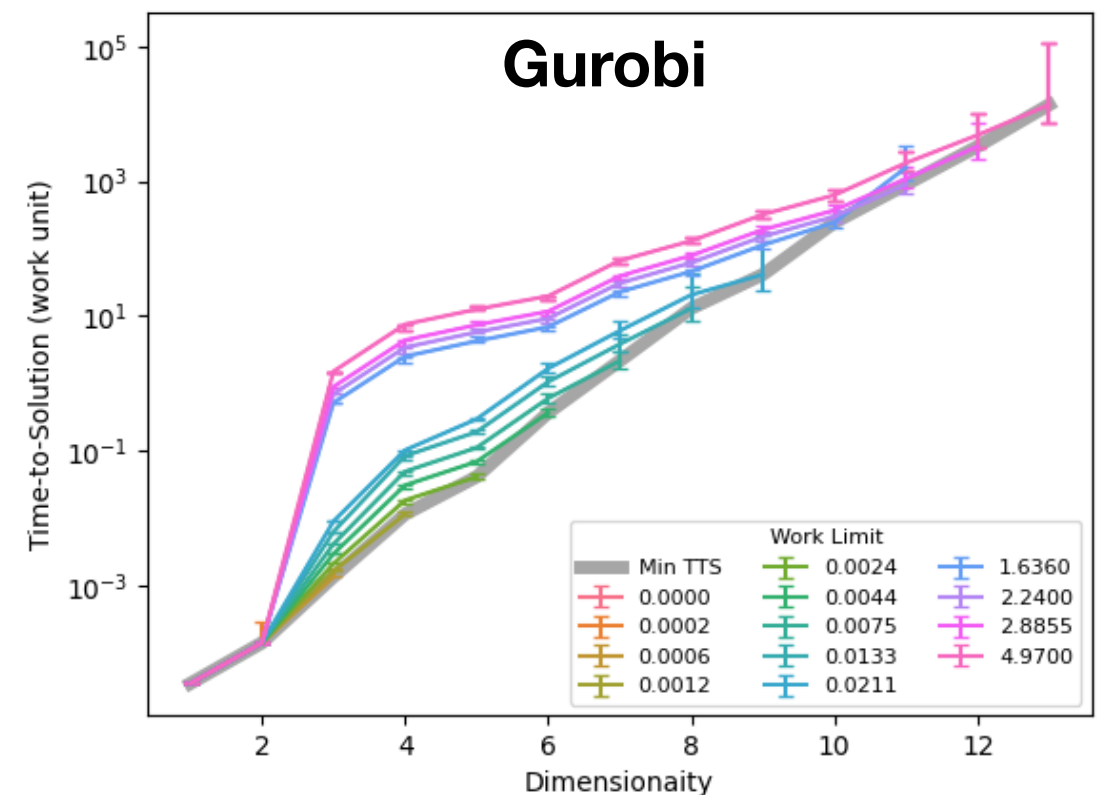
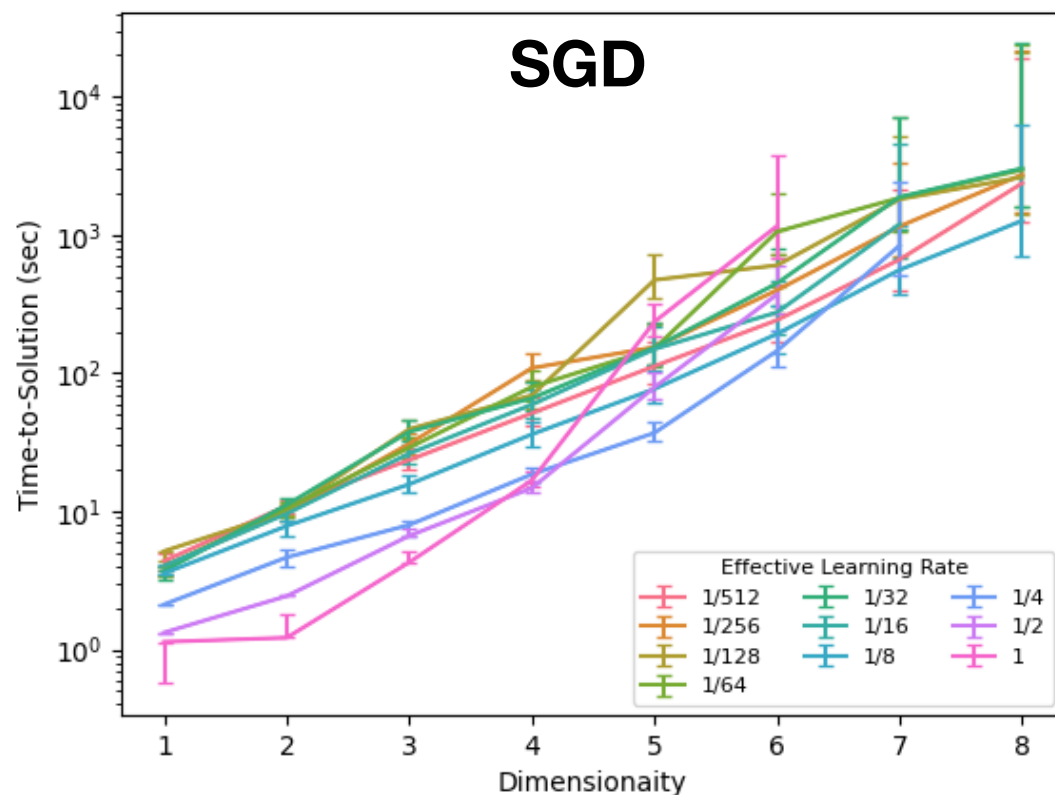
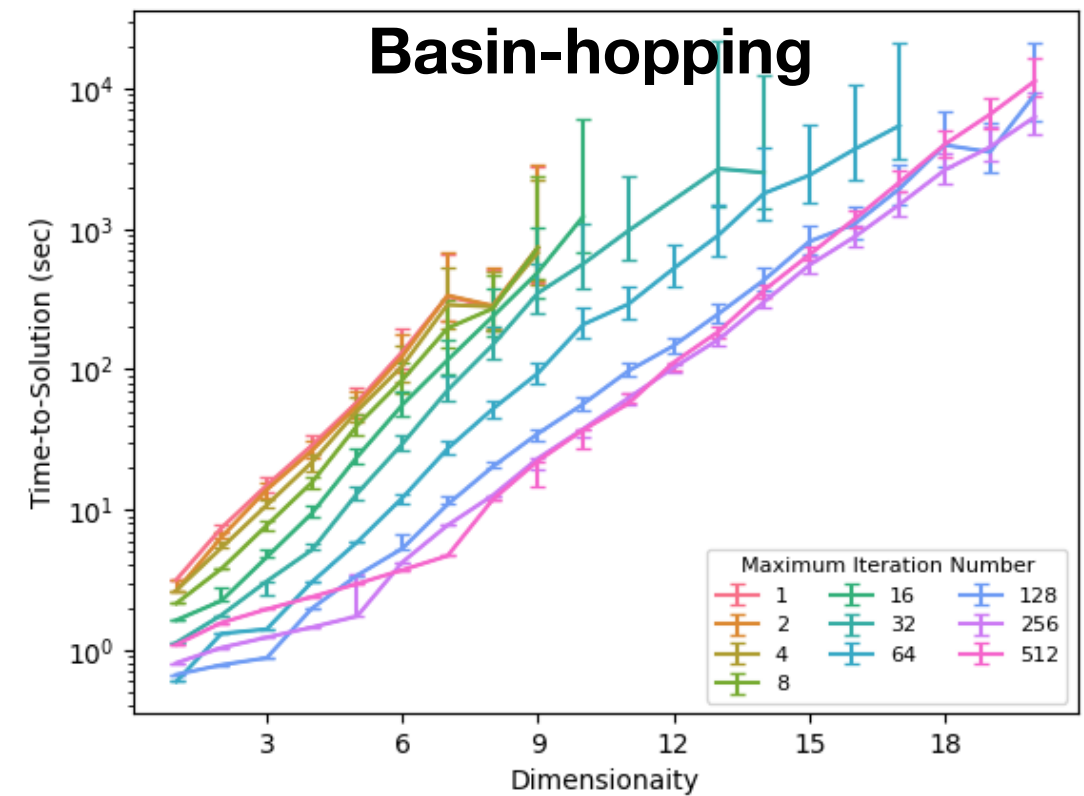
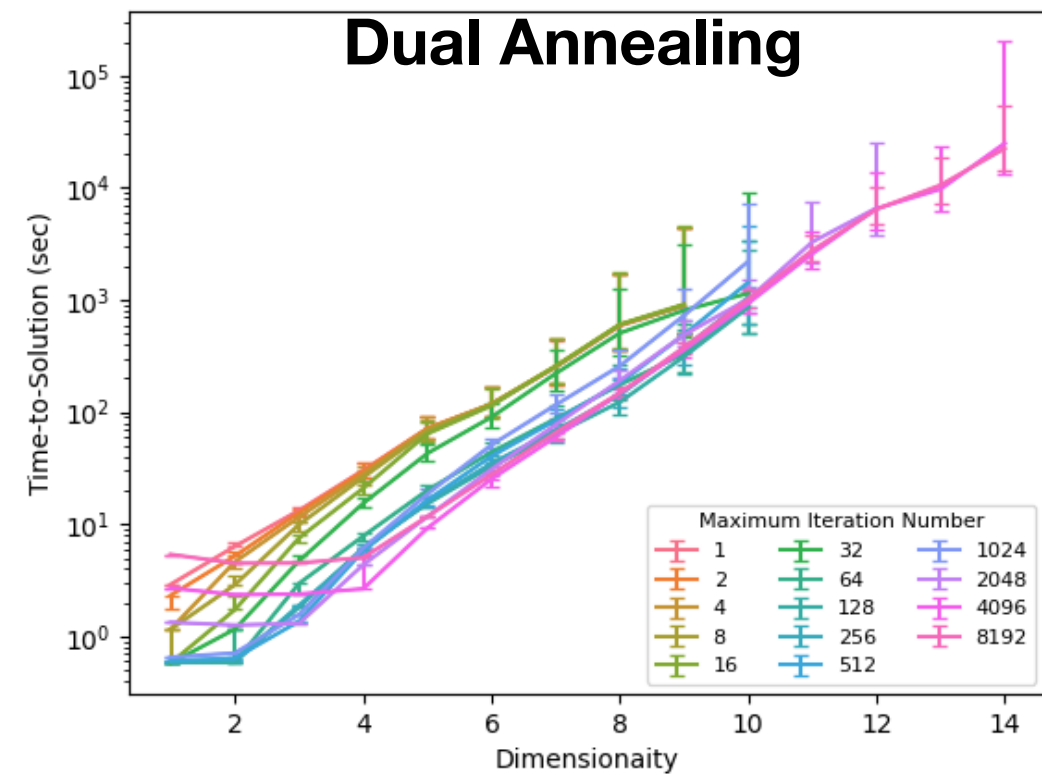
- $t_f$  - algorithm/solver running time (wall clock time)
- $p_g$  - success probability per run

**QHD TTS** -  $\mathcal{O}(d^4)$  (given fixed  $\delta$ )

**Classical TTS** - numerical results suggest **super-polynomial** scaling in  $d$



# A quantum-classical performance separation



# Part III:

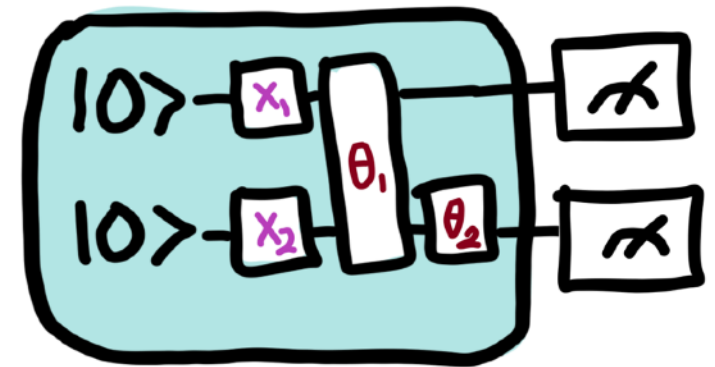
Implementation on Quantum Computers

# QHD with digital quantum computers

$$H(t) = e^{\varphi(t)}\left(-\frac{1}{2}\nabla^2\right) + e^{\chi(t)}f(x), t \geq 0$$

## Digital Quantum Implementation

- The algorithm is effectively a **time-dependent Hamiltonian simulation** in the real space (we know poly-time algorithms).
- Requires **hundreds of millions of gates!**
- **Conclusion:** the digital implementation is **far from being feasible** in near term!



Corresponding **T-gate** Count with Digital Quantum Computing (before fault-tolerance)

Dimensions	3-qubit format	16-qubit format	32-qubit format
50	5.49e+8	7.8386e+9	2.672e+10
60	6.588e+8	9.4063e+9	3.2064e+10
75	8.235e+8	1.1758e+10	4.008e+10



# QHD with analog quantum computers

## Analog Quantum Implementation

- **Analog simulation:** problem solving by *emulating real quantum systems*.
- Abstraction: **Quantum Ising Machine (QIM)**, e.g., D-Wave, QuEra, etc.)

$$H(t) = -\frac{A(t)}{2} \left( \sum_j \sigma_x^{(j)} \right) + \frac{B(t)}{2} \left( \sum_j h_j \sigma_z^{(j)} + \sum_{j>k} J_{j,k} \sigma_z^{(j)} \sigma_z^{(k)} \right)$$

- **Programmability:** coefficients  $h_j, J_{j,k}$  and functions  $A(t), B(t)$ .



- QHD is formulated as a quantum Hamiltonian evolution → suitable for analog implementation!
- Patter mismatch:  $\text{QHD } H(t) = e^{\varphi_t(-\Delta/2)} + e^{\chi_t f}$
- We develop a new technique named **Hamiltonian embedding**: mapping our target Hamiltonian (QHD) to a “diagonal block” of the machine Hamiltonian (QIM).
- If  $H = H_0 \oplus H_1$ , we have  $e^{-iHt} = e^{-iH_0t} \oplus e^{-iH_1t}$ .
- We implement QHD on today’s analog quantum computers (D-Wave’s **advantage\_system6.1**).
- **Advantage:** resource-efficiency → first large-scale empirical study for nonlinear optimization using quantum computers.

# Hamiltonian embedding of QHD

$$i \frac{\partial}{\partial t} \Psi(t, x) = \left[ \underbrace{-\frac{e^{\varphi_t}}{2} \nabla^2}_{\text{Kinetic part}} + \underbrace{e^{\chi_t} f(x)}_{\text{Potential part}} \right] \Psi(t, x),$$

Finite difference 

$$i \frac{d}{dt} |\phi_t\rangle = \left[ \underbrace{-\frac{e^{\varphi_t}}{2} \hat{A}}_{\text{Kinetic part}} + \underbrace{e^{\chi_t} \hat{F}}_{\text{Potential part}} \right] |\phi_t\rangle$$

- **Kinetic operator:**  $\hat{L} = \frac{1}{(1/r)^2} \begin{bmatrix} -2 & 1 & & \\ 1 & -2 & 1 & \\ \dots & \dots & \dots & \dots \\ & 1 & -2 & 1 \\ & & 1 & -2 \end{bmatrix} \rightarrow \hat{A} = \frac{1}{(1/r)^2} \begin{bmatrix} 0 & 1 & & \\ 1 & 0 & 1 & \\ \dots & \dots & \dots & \dots \\ & 1 & 0 & 1 \\ & & 1 & 0 \end{bmatrix}$   
Tri-diagonal
- **Potential operator:**  $\hat{F} = \begin{bmatrix} f(a_0) & & & \\ & f(a_1) & & \\ \dots & \dots & \dots & \dots \\ & & f(a_{r-1}) & \\ & & & f(a_r) \end{bmatrix}$   
Diagonal

**Hamming states — an orthonormal basis**

$$|H_j\rangle = \frac{1}{\sqrt{C_j}} \sum_{|b|=j} |b\rangle \quad \text{where } C_j = \binom{N}{j} \text{ and } 0 \leq j \leq N$$

$$|H_0\rangle = |0000\rangle, |H_1\rangle = \frac{1}{2}(|1000\rangle + |0100\rangle + |0010\rangle + |0001\rangle), \text{ etc.}$$

**Lemma (informal version).** [Leng, Hickman, Li, Wu (2023)]

Given  $n$  qubits, the subspace  $\mathcal{S}$  spanned by all  $(n + 1)$  Hamming states is an invariant subspace of the QIM Hamiltonian. The projection of the QIM Hamiltonian into the subspace  $\mathcal{S}$  approximates the discretized QHD Hamiltonian.

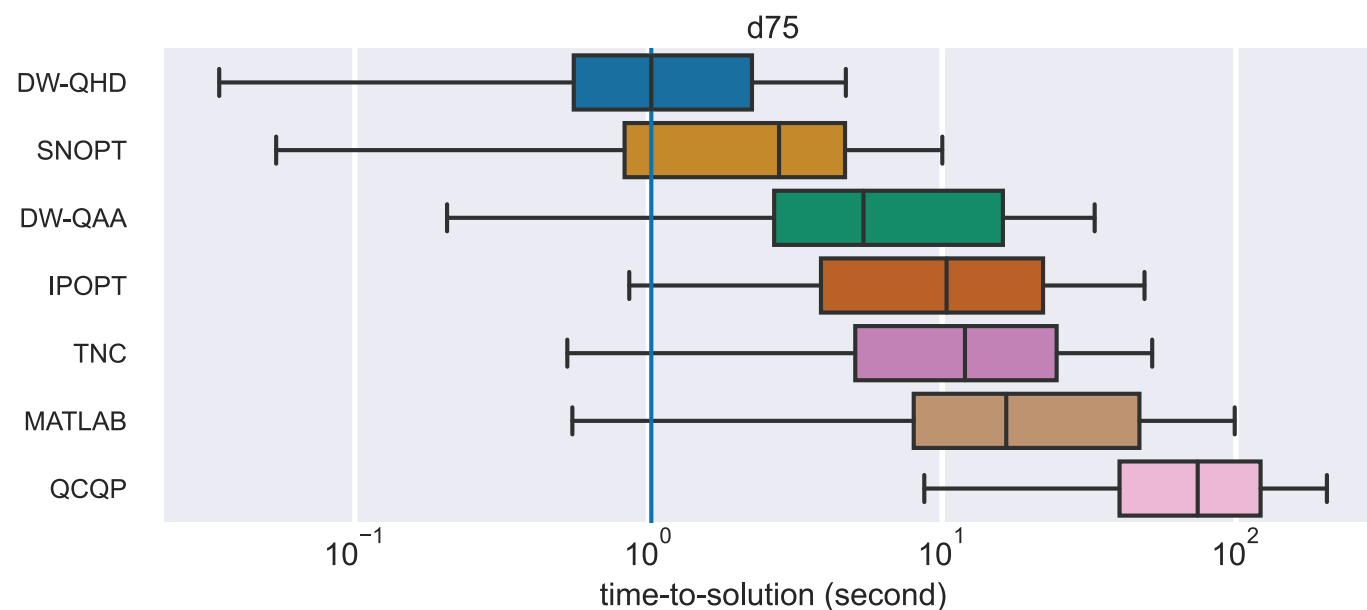
# Large-scale empirical study on real quantum computers

We identify a class of *non-trivial* and *self-interesting* optimization problems that can be mapped to QIMs — **Quadratic Programming (QP) with box constraints**.

$$\begin{array}{ll} \text{minimize} & f(\mathbf{x}) = \frac{1}{2} \mathbf{x}^\top \mathbf{Q} \mathbf{x} + \mathbf{b}^\top \mathbf{x} \\ \text{subject to} & \mathbf{0} \preceq \mathbf{x} \preceq \mathbf{1}, \end{array}$$

**QP** — NP-hard with indefinite Q

(For implementation details, see our paper [arXiv:2303.01471](https://arxiv.org/abs/2303.01471) Appendix F)



**Time-To-Solution (TTS)**

The lower, the better!!

$$\text{TTS} = t_f \left[ \frac{\ln(1 - 0.99)}{\ln(1 - P_g)} \right]$$

$t_f$  - quantum anneal time + post-processing  
or classical runtime (**wall-clock time**)

$P_g$  - success probability per run

- **DW-QHD** is better than the rest, including **DW-QAA**, and **classical GD, interior points**, and **some local search heuristic**.
- Assuming DW-QHD is no worse than the ideal QHD, this provides **a very strong empirical evidence supporting QHD**.
- QHD does not beat SOTA classical solvers (e.g., Gurobi, CPLEX) in D75. However, such branch-and-bound solvers are **not scalable!**

# Summary & Future Work

- **QHD** is an upgraded version of **classical GD and variants**.
- QHD leverages the continuous structure of the problem and converges faster than QAA.
- QHD has different solution path compared to classical GD.
- **QHD** can be used as *subroutines* for more complicated algorithms like branch-and-bound.

## All data & codes are available online!

- Source code (Github): <https://github.com/jiaqileng/quantum-hamiltonian-descent>
- Raw data (Box): <https://umd.app.box.com/s/vq747fvjnt8qrkbpexhoh44n0q9m0i>
- Website: <https://jiaqileng.github.io/quantum-hamiltonian-descent/>



arXiv:2303.01471

# Thank You!



QHD Website