

EXCHANGEABLE RANDOM GRAPHS, GRAPH LIMITS, AND GRAPHONS

Jiasen Yang

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Department of Statistics
Purdue University

Graph Data

Exchangeable Random Graphs

Graph Limits and Graphons

Graphon Estimation

GRAPH DATA

A **graph** $G = (V, E)$, where

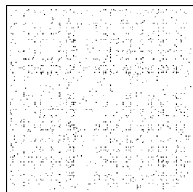
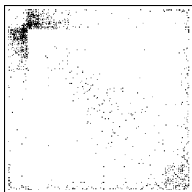
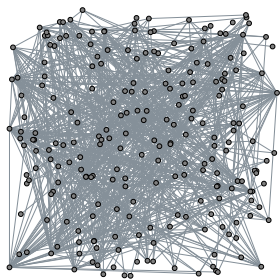
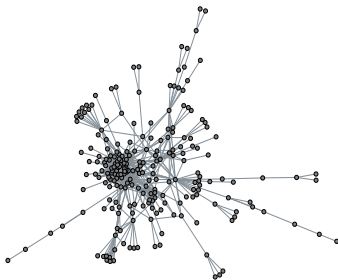
- $V = \{1, 2, \dots, n\}$ is the set of vertices;
- $E \subseteq V \times V$ is the set of edges,

can be represented by an **adjacency matrix** A , where

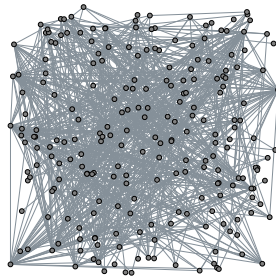
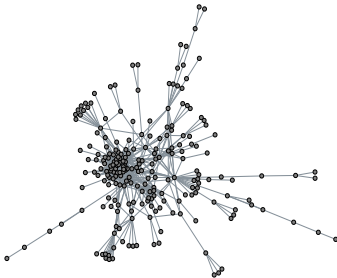
$$A_{ij} = \begin{cases} 1, & \text{if } (i, j) \in E; \\ 0, & \text{otherwise.} \end{cases}$$

Assume G is simple and undirected $\Rightarrow A$ is binary, symmetric, and $A_{ii} = 0, \forall i$.

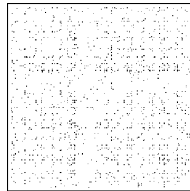
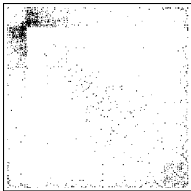
GRAPH DATA (CONT'D)



GRAPH DATA (CONT'D)



They are the same network!



EXCHANGEABLE RANDOM GRAPHS

Definition

An **exchangeable sequence** is an infinite sequence X_1, X_2, \dots of random variables whose joint distribution satisfies

$$\mathbb{P}(X_1 \in A_1, X_2 \in A_2, \dots) = \mathbb{P}(X_{\pi(1)} \in A_1, X_{\pi(2)} \in A_2, \dots)$$

for every permutation π of $\mathbb{N} := \{1, 2, \dots\}$ and collection A_1, A_2, \dots of (measurable) sets.

In words, order of observations does not matter.

i.i.d.* \Rightarrow exchangeable; exchangeable \nRightarrow *i.i.d.

Example (Polyá's Urn scheme)

Consider an urn with b black balls and w white balls. Draw a ball at random and note its color. Replace the ball together with a balls of the same color. Repeat the procedure *ad infinitum*. Let $X_i = \mathbb{I}\{\text{the } i\text{-th draw yields a black ball}\}$.

The sequence X_1, X_2, \dots is exchangeable but not *i.i.d.*/Markov.

$$\begin{aligned} \text{E.g., } P(X_1 = 1, X_2 = 1, X_3 = 0, X_4 = 1) \\ &= \frac{b}{b+w} \frac{b+a}{b+w+a} \frac{w+a}{b+w+2a} \frac{b+2a}{b+w+3a} \\ &= \frac{b}{b+w} \frac{w+a}{b+w+2a} \frac{b+a}{b+w+a} \frac{b+2a}{b+w+3a} \\ &= P(X_1 = 1, X_2 = 0, X_3 = 1, X_4 = 1). \end{aligned}$$

Theorem (de Finetti, 1931)

Let X_1, X_2, \dots be an infinite sequence of random variables with values in a space \mathcal{X} . The sequence X_1, X_2, \dots is *exchangeable* if and only if there is a *random probability measure* Θ on \mathcal{X} —i.e., a random variable with values in the set $\mathcal{M}(\mathcal{X})$ of probability distributions on \mathcal{X} —such that the X_i are *conditionally i.i.d.* given Θ and

$$\mathbb{P}(X_1 \in A_1, X_2 \in A_2, \dots) = \int_{\mathcal{M}(\mathcal{X})} \prod_{i=1}^{\infty} \theta(A_i) \nu(d\theta)$$

where ν is the distribution of Θ .

Adj. matrix of undirected **graph** = binary & symmetric **array**

⇒ de Finetti exchangeability too strong!

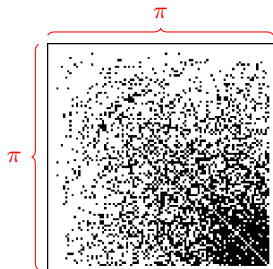
⇒ Need generalization of exchangeability to arrays.

EXCHANGEABLE RANDOM GRAPHS

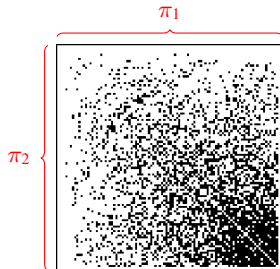
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⇒ Need generalization of exchangeability to arrays.



$$\sigma = \pi \otimes \pi$$



$$\sigma = \pi_1 \otimes \pi_2$$

Definition

A random array $(X_{ij})_{i,j \in \mathbb{N}}$ is called **separately exchangeable** if

$$(X_{ij}) \stackrel{d}{=} (X_{\pi(i)\pi'(j)})$$

holds for every pair of permutations π, π' of \mathbb{N} .

Definition

A random array $(X_{ij})_{i,j \in \mathbb{N}}$ is called **jointly exchangeable** if

$$(X_{ij}) \stackrel{d}{=} (X_{\pi(i)\pi(j)})$$

holds for every permutation π of \mathbb{N} .

ALDOUS-HOOVER THEOREM

Theorem (Aldous, 1981-Hoover, 1979)

A random array $(X_{ij})_{i,j \in \mathbb{N}}$ is

- *jointly exchangeable* iff. \exists a random measurable function $F : [0, 1]^3 \rightarrow \mathcal{X}$ s.t.

$$(X_{ij}) \stackrel{d}{=} (F(U_i, U_j, U_{\{\{i,j\}\}}))$$

- *separately exchangeable* iff. \exists a random measurable function $F : [0, 1]^3 \rightarrow \mathcal{X}$ s.t.

$$(X_{ij}) \stackrel{d}{=} (F(U_i, U'_j, U_{ij}))$$

where $U_i, U_{\{\{i,j\}\}}, U'_j, U_{ij} \stackrel{i.i.d.}{\sim} \text{Uniform}[0, 1]$.

Definition

An undirected *random* graph $G = (\mathbb{N}, E)$ with an infinite (countable) vertex set \mathbb{N} and a *random* edge set E is called an **exchangeable random graph** if its *random* adjacency matrix $A = (A_{ij})$ is a jointly exchangeable array.

Thus, G is exchangeable if its distribution is invariant under relabeling of the vertices.

Adjacency matrix $A = (A_{ij})$ symmetric & binary:

$$A_{ij} \stackrel{d}{=} F(U_i, U_j, U_{\{i,j\}}) \stackrel{d}{=} \mathbb{I}\{U_{\{i,j\}} < W(U_i, U_j)\}$$

EXCHANGEABLE RANDOM GRAPHS

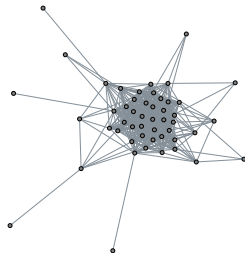
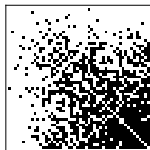
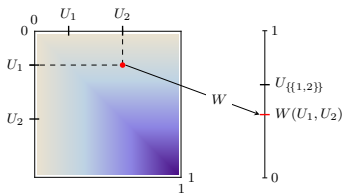
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Sample $W : [0, 1]^2 \rightarrow [0, 1]$ measurable & symmetric (**graphon**);

Sample $U_1, U_2, \dots \stackrel{i.i.d.}{\sim} \text{Uniform}[0, 1]$;

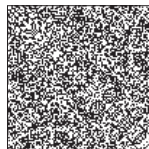
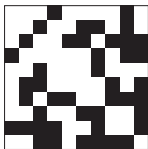
Sample $A_{ij} \sim \text{Bernoulli}(W(U_i, U_j))$ for $i < j$.



GRAPH LIMITS AND GRAPHONS

GRAPH LIMITS

Rado graph: $A_{ij} = A_{ji} \stackrel{i.i.d.}{\sim} \text{Bernoulli}(1/2), i \neq j \in \mathbb{N}$.



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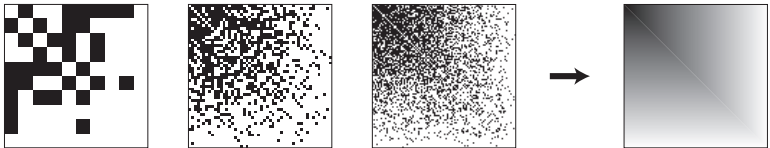


Limit graphon (graph function):

$$f : [0, 1]^2 \rightarrow [0, 1], (x, y) \mapsto \frac{1}{2}.$$

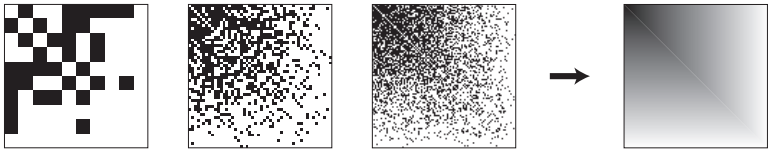
GRAPH LIMITS (CONT'D)

Growing uniform attachment: Let $G_1 = \bullet$. For $n \geq 2$, construct G_n from G_{n-1} by adding one new vertex, then, drawing an edge between each pair of non-adjacent vertices with probability $\frac{1}{n}$.



GRAPH LIMITS (CONT'D)

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Limit graphon:

$$f : [0, 1]^2 \rightarrow [0, 1], (x, y) \mapsto 1 - \max(x, y).$$

Definition

A **labeled graphon** is a symmetric, Lebesgue-measurable a.e. function $W : [0, 1]^2 \rightarrow [0, 1]$. A labeled graphon determines the equivalence class of graphons

$$[W] = \{W^\phi : (x, y) \mapsto W(\phi(x), \phi(y))\},$$

where ϕ is an invertible, measure-preserving transformation of $[0, 1]$. Such equivalence classes are called **unlabeled graphons**.

Definition

The **cut distance** between two labeled graphons W and U is

$$\delta_{\square}(W, U) = \inf_{\substack{\phi, \psi \text{ m.p.t.} \\ \text{of } [0,1]}} \sup_{\substack{S, T \subseteq [0,1] \\ \text{measurable}}} \left| \int_{S \times T} W^{\phi}(x, y) - U^{\psi}(x, y) \right|.$$

Theorem

Every graphon is the δ_{\square} -limit of a sequence of finite graphs.

Theorem (Lovász-Szegedy)

*Let \mathcal{G} be the space of unlabeled graphons (modulo weak isometry). The metric space $(\mathcal{G}, \delta_{\square})$ is **compact**.*

GRAPHON ESTIMATION

EXCHANGEABLE RANDOM GRAPHS

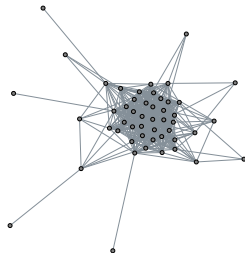
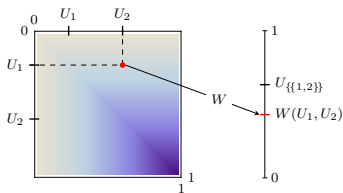
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Sample $W : [0, 1]^2 \rightarrow [0, 1]$ measurable & symmetric (graphon);

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Sample $A_{ij} \sim \text{Bernoulli}(W(U_i, U_j))$ for $i < j$.



Problem: Estimate graphon W from a set of observed networks.

Nonparametric regression: estimate W from $\{(U_i, U_j), A_{ij}\}_{i,j \in \mathbb{N}}$.

Challenge: design points $\{(U_i, U_j)\}$ are latent.

Proposed estimators:

- Stochastic blockmodel approximation [Airoldi *et.al.*, 2013], [Gao *et.al.*, 2015]
- Histogram estimator (sorting-and-smoothing) [Chan and Airoldi, 2014]
- Gaussian process model [Lloyd *et.al.*, 2012], [Orbanz and Roy, 2015]

E. M. Airoldi, T. B. Costa, and S. H. Chan. (2013) Stochastic Blockmodel Approximation of a Graphon: Theory and Consistent Estimation. *NIPS '13*.

Gao, C., Lu, Y. and Zhou, H. H. (2015). Rate-Optimal Graphon Estimation. *Ann. Statist.* 43(6):2624-2652.

S.H. Chan and E.M. Airoldi (2014). A Consistent Histogram Estimator for Exchangeable Graph Models. *ICML '14*.

J.R. Lloyd, P. Orbanz, Z. Ghahramani, and D.M. Roy (2012). Random Function Priors for Exchangeable Arrays with Applications to Graphs and Relational Data. *NIPS '12*.

P. Orbanz and D. M. Roy (2015). Bayesian Models of Graphs, Arrays and Other Exchangeable Random Structures. *IEEE Trans. Pattern Anal. Mach. Intell.* 37(2):437-461.

- Community detection
- Link prediction:

$$\mathbb{P}((i, j) \in E) = \mathbb{P}(A_{ij} = 1) = W(U_i, U_j)$$

- Network comparison & hypothesis testing

Real-world graphs are **sparse** (finite # of edges per vertex).

- Power-law degree distribution;
- Small-world phenomena.

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Theorem (Misspecification)

If a random graph is **exchangeable**, it is either **dense** or empty.

Proof.

For a random graph G_n with n vertices, the expected proportion of present edges is

$$p := \int_{[0,1]^2} W(x,y) dx dy.$$

If $p = 0$, G_n is empty; if $p > 0$, G_n has $p \cdot \binom{n}{2} = \Theta(n^2)$ edges in expectation. □

The limit object of a convergent sequence of **sparse** graphs in the cut metric is always the **empty** graphon.

Representation theorems for **sparse** random graphs?

THANK YOU!