

Exact diagonalization studies

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Studies of small spin chains

Using basis states incorporating conservation laws (symmetries)

- magnetization conservation, momentum states, parity, spin inversion
- discussion without group theory
 - only basic quantum mechanics and common sense needed

Key elements of an actual program implementation

- pseudocode segments; actual program available for practice

How to characterize different kinds of ground states

- critical ground state of the Heisenberg chain
- quantum phase transition to a valence-bond solid in a J_1-J_2 chain

Complete diagonalization of the hamiltonian

To find the ground state (maybe excitations, T>0 properties) of the Heisenberg S=1/2 chain

$$\begin{aligned} H &= J \sum_{i=1}^N \mathbf{S}_i \cdot \mathbf{S}_{i+1} = J \sum_{i=1}^N [S_i^x S_{i+1}^x + S_i^y S_{i+1}^y + S_i^z S_{i+1}^z], \\ &= J \sum_{i=1}^N [S_i^z S_{i+1}^z + \frac{1}{2}(S_i^+ S_{i+1}^- + S_i^- S_{i+1}^+)] \end{aligned}$$

Simplest way; enumerate the states

- construct the hamiltonian matrix using bits

$$|0\rangle = |\downarrow, \downarrow, \downarrow, \dots, \downarrow\rangle \quad (= 0\dots000)$$

$$|1\rangle = |\uparrow, \downarrow, \downarrow, \dots, \downarrow\rangle \quad (= 0\dots001)$$

$$|2\rangle = |\downarrow, \uparrow, \downarrow, \dots, \downarrow\rangle \quad (= 0\dots010)$$

$$|3\rangle = |\uparrow, \uparrow, \downarrow, \dots, \downarrow\rangle \quad (= 0\dots011)$$

$$\begin{aligned} H_{ij} &= \langle i | H | j \rangle \\ i, j &= 0, \dots, 2^N - 1 \end{aligned}$$

bit representation perfect for S=1/2 systems

- use >1 bit/spin for S>1/2, or integer vector
- construct H by examining/flipping bits

Diagonalizing the matrix

- on the computer
- gives the eigenvalues and eigenvectors

If \mathbf{U} is the matrix whose columns are the eigenvectors of \mathbf{H} , then

$$\langle n | A | n \rangle = [U^{T*} A U]_{nn}$$

is the expectation value of some operator A in the n:th eigenstate

Use some “canned” diagonalization subroutine

- Useful subroutines available at <http://gams.nist.gov>

Problem: Matrix size $M=2^N$ becomes too large quickly

- maximum $N \approx 20$
- M^2 matrix elements to store, time to diagonalize $\propto M^3$

Pseudocode; construction of the hamiltonian matrix

Let $a[i]$ refer to the i :th bit of an integer a

Define a function $\text{flip}(a,i,j)$

- “flips” ($0 \leftrightarrow 1$) bits i and j of the integer a
- In F90 the bit-level function $\text{ieor}(a,2^{**}i)$ can be used to flip bit i of a
- bits i and j can be flipped using $\text{ieor}(a,2^{**}i+2^{**}j)$

		j	i					
a		0	1	0	1	0	0	1
$2^i + 2^j$		0	0	0	1	1	0	0
$\text{ieor}(a, 2^i + 2^j)$		0	1	0	0	1	0	1

The S=1/2 Heisenberg chain hamiltonian
can be constructed according to

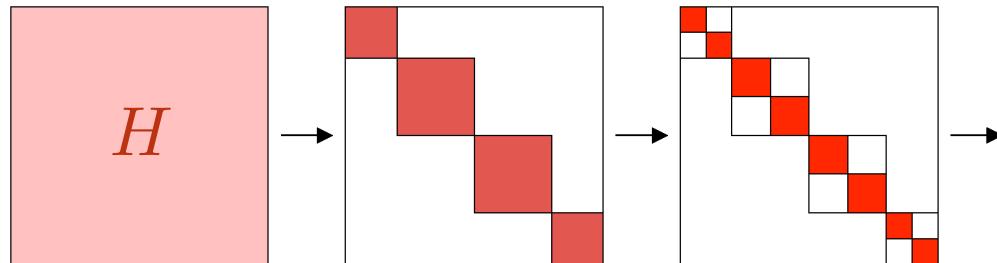
j is the “right” nearest-neighbor of i

- periodic boundary conditions

```
do a = 0, 2N - 1
    do i = 0, N - 1
        j = mod(i + 1, N)
        if (a[i] = a[j]) then
            H(a, a) = H(a, a) + 1/4
        else
            H(a, a) = H(a, a) - 1/4
            b = flip(a, i, j); H(a, b) = 1/2
        endif
    enddo
enddo
```

Using conservation laws (symmetries) for block-diagonalization

We can choose the basis in such a way that the H becomes block-diagonal



- the blocks can be diagonalized individually
- we can reach larger N (but not much larger, $N \approx 40$ is max)

Simplest example: magnetization conservation

$$m_z = \sum_{i=1}^N S_i^z$$

- blocks correspond to fixed values of m_z
- no H matrix elements between states of different m_z
- block-diagonalization just amounts to including states with only given m_z

Number of states in the largest block ($m_z = 0$):

$$\frac{N!}{(N/2)!(N/2)!}$$

Other symmetries (conserved quantum numbers)

- can be used to further split the blocks
- but more complicated
 - basis states have to be constructed to obey symmetries

Pseudocode: using magnetization conservation

Constructing the basis in the block of n_\uparrow spins \uparrow

Store state-integers in ordered list s_a , $a=1,\dots,M$

```
do  $s = 0, 2^N - 1$ 
    if ( $\sum_i s[i] = n_\uparrow$ ) then  $a = a + 1$ ;  $s_a = s$  endif
enddo
 $M = a$ 
```

Example; $N=4$, $n_\uparrow=2$

$s_1=3$	(0011)
$s_2=5$	(0101)
$s_3=6$	(0110)
$s_4=9$	(1001)
$s_5=10$	(1010)
$s_6=12$	(1100)

```
subroutine findstate( $s, b$ )
do
     $b_{\min} = 1$ ;  $b_{\max} = M$ 
     $b = b_{\min} + (b_{\max} - b_{\min})/2$ 
    if ( $s < s_b$ ) then
         $b_{\max} = b - 1$ 
    elseif ( $s > s_b$ ) then
         $b_{\min} = b + 1$ 
    else
        exit
    endif
enddo
```

Finding the location b
of a state-integer s in the list
• using bisection in the ordered list

Hamiltonian construction

```
do  $a = 1, M$ 
    do  $i = 0, N - 1$ 
         $j = \text{mod}(i + 1, N)$ 
        if ( $s_a[i] = s_a[j]$ ) then
             $H(a, a) = H(a, a) + \frac{1}{4}$ 
        else
             $H(a, a) = H(a, a) - \frac{1}{4}$ 
             $s = \text{flip}(s_a, i, j)$ 
            call findstate( $s, b$ )
             $H(a, b) = H(a, b) + \frac{1}{2}$ 
        endif
    enddo
enddo
```

Momentum states (translationally invariant systems)

A periodic chain (ring) is translationally invariant

- the eigenstates have a momentum (crystal momentum)
- quantum number k

$$T|n\rangle = e^{ik}|n\rangle \quad k = m\frac{2\pi}{N}, \quad m = 0, \dots, N-1,$$

The operator T translates the state by one lattice spacing

- for a spin basis state

$$T|S_1^z, S_2^z, \dots, S_N^z\rangle = |S_N^z, S_1^z, \dots, S_{N-1}^z\rangle$$

$[T, H] = 0 \rightarrow$ momentum blocks of H

- can use eigenstates of T with given k as basis

A momentum state can be constructed from any **representative** state

$$|a(k)\rangle = \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{-ikr} T^r |a\rangle, \quad |a\rangle = |S_1^z, \dots, S_N^z\rangle$$

Construct ordered list of representatives
If $|a\rangle$ and $|b\rangle$ are representatives, then

$$T^r |a\rangle \neq |b\rangle \quad r \in \{1, \dots, N-1\}$$

4-site examples

$$(0011) \rightarrow (0110), (1100), (1001)$$

$$(0101) \rightarrow (1010)$$

The sum can contain several copies of the same state

If $T^R |a\rangle = |a\rangle$ for some R

- the total weight for this component is

$$1 + e^{-ikR} + e^{-i2kR} + \dots + e^{-ik(N-R)}$$

- vanishes (state incompatible with k) unless $kR=n2\pi$
- the total weight of the representative is then N/R

$$kR = n2\pi \rightarrow \frac{mR}{N} = n \rightarrow m = n \frac{N}{R} \rightarrow \text{mod}(m, N/R) = 0$$

$$|a(k)\rangle = \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{-ikr} T^r |a\rangle$$

Normalization of a state $|a(k)\rangle$ with periodicity R_a

$$\langle a(k) | a(k) \rangle = \frac{1}{N_a} \times R_a \times \left(\frac{N}{R_a} \right)^2 = 1 \rightarrow N_a = \frac{N^2}{R_a}$$

Pseudocode; basis construction

```

do  $s = 0, 2^N - 1$ 
    call checkstate( $s, R$ )
    if  $R \geq 0$  then  $a = a + 1; s_a = s; R_a = R$  endif
enddo
 $M = a$ 

```

M = size of
the H-block

Uses a subroutine **checkstate**(s, R)

- R = periodicity if state-integer s is a new representative
- store in list $R_a, a=1, \dots, M$
- $R = -1$ if some translation of $|s\rangle$ gives a smaller integer

Translations of the representative; cyclic permutation

Define function **cyclebits**(t,N)

- cyclic permutations of first N bits of integer t
- F90 function `ishiftc(t,-1,N)`

Pseudocode; **checkstate()** subroutine

```
subroutine checkstate(s, R)
R = -1
if ( $\sum_i s[i] \neq n_{\uparrow}$ ) return
t = s
do i = 1, N
    t = cyclebits(t, N)
    if (t < s) then
        return
    elseif (t = s) then
        if ( $\text{mod}(k, N/i) \neq 0$ ) return
        R = i; return
    endif
enddo
```

The representative is the lowest integer among all translations

The Hamiltonian matrix. Write $S = 1/2$ chain hamiltonian as

$$H_0 = \sum_{j=1}^N S_j^z S_{j+1}^z, \quad H_j = \frac{1}{2}(S_j^+ S_{j+1}^- + S_j^- S_{j+1}^+), \quad j = 1, \dots, N$$

Act with H on a momentum state

$$H|a(k)\rangle = \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{-ikr} T^r H |a\rangle = \frac{1}{\sqrt{N_a}} \sum_{j=0}^N \sum_{r=0}^{N-1} e^{-ikr} T^r H_j |a\rangle,$$

$H_j|a\rangle$ is related to some representative: $H_j|a\rangle = h_a^j T^{-l_j} |b_j\rangle$

$$H|a(k)\rangle = \sum_{j=0}^N \frac{h_a^j}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{-ikr} T^{(r-l_j)} |b_j\rangle$$

Shift summation index r and use definition of momentum state

$$H|a(k)\rangle = \sum_{j=0}^N h_a^j e^{-ikl_j} \sqrt{\frac{N_{b_j}}{N_a}} |b_j(k)\rangle \quad \rightarrow \text{matrix elements}$$

$$\langle a(k)|H_0|a(k)\rangle = \sum_{j=1}^N S_j^z S_j^z,$$

$$\langle b_j(k)|H_{j>0}|a(k)\rangle = e^{-ikl_j} \frac{1}{2} \sqrt{\frac{R_a}{R_{b_j}}}, \quad |b_j\rangle \propto T^{-l_j} H_j |a\rangle,$$

Pseudocode; hamiltonian construction

First, some elements needed; recall

$$H_j |a\rangle = h_a^j T^{-l_j} |b_j\rangle$$

Finding the representative r of a state-integer s

```
subroutine representative(s, r, l)
  r = s; t = s; l = 0
  do i = 1, N - 1
    t = cyclebits(t, N)
    if (t < r) then r = t; l = i endif
  enddo
```

$$|r\rangle = T^l |s\rangle$$

```
subroutine findstate(s, b)
  do
    b_min = 1; b_max = M
    b = b_min + (b_max - b_min)/2
    if (s < s_b) then
      b_max = b - 1
    elseif (s > s_b) then
      b_min = b + 1
    else
      exit
    endif
    if (b_min > b_max) then
      b = -1; exit
    endif
  enddo
```

Finding the location of the representative in the state list

- may not be there, if the new state is incompatible with k
- **b=-1** for not found in list

Construct all the matrix elements

```
do a = 1, M
    do i = 0, N - 1
        j = mod(i + 1, N)
        if (s_a[i] = s_a[j]) then
            H(a, a) = H(a, a) + 1/4
        else
            H(a, a) = H(a, a) - 1/4
            s = flip(s_a, i, j)
            call representative(s, r, l)
            call findstate(r, b)
            if (b ≥ 0) then
                H(a, b) = H(a, b) + 1/2 √(R_a/R_b) e^{i2πkl/N}
            endif
        endif
    enddo
enddo
```

Reflection symmetry (parity) Define a reflection (parity) operator

$$P|S_1^z, S_2^z, \dots, S_N^z\rangle = |S_N^z, \dots, S_2^z, S_1^z\rangle$$

Consider a hamiltonian for which $[H, P] = 0$ and $[H, T] = 0$; but note that $[P, T] \neq 0$

Can we still exploit both P and T at the same time? Consider the state

$$|a(k, p)\rangle = \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{-ikr} T^r (1 + pP) |a\rangle, \quad p = \pm 1$$

This state has momentum k, but does it have parity p? Act with P

$$\begin{aligned} P|a(k, p)\rangle &= \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{-ikr} T^{-r} (P + p) |a\rangle \\ &= p \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} e^{ikr} T^r (1 + pP) |a\rangle = p|a(k, p)\rangle \text{ if } k = 0 \text{ or } k = \pi \end{aligned}$$

k=0,π momentum blocks are split into p=+1 and p=-1 sub-blocks

- $[T, P] = 0$ in the $k=0, \pi$ blocks
- physically clear because $-k=k$ on the lattice for $k=0, \pi$
- we can exploit parity in a different way for other $k \rightarrow$
- **semi-momentum states**

Semi-momentum states

Mix momenta $+k$ and $-k$ for $k \neq 0, \pi$. Introduce function

$$C_k^\sigma(r) = \begin{cases} \cos(kr), & \sigma = +1 \\ \sin(kr), & \sigma = -1. \end{cases}$$

Useful trigonometric relationships

$$\begin{aligned} C_k^\pm(-r) &= \pm C_k^\pm(r), \\ C_k^\pm(r+d) &= C_k^\pm(r)C_k^+(d) \mp C_k^\mp(r)C_k^-(d). \end{aligned}$$

Semi-momentum state

$$|a^\sigma(k)\rangle = \frac{1}{\sqrt{N_a}} \sum_{r=0}^{N-1} C_k^\sigma(r) T^r |a\rangle$$

$$k = m \frac{2\pi}{N}, \quad m = 1, \dots, N/2 - 1, \quad \sigma = \pm 1$$

States with same k , different σ are orthogonal

$$\langle a^{-\sigma}(k) | a^\sigma(k) \rangle = \frac{1}{N_a} \sum_{r=1}^{R_a} \sin(kr) \cos(kr) = 0,$$

Normalization of semi-momentum states

$$N_a = \left(\frac{N}{R_a} \right)^2 \sum_{r=1}^{R_a} [C_k^\sigma(r)]^2 = \frac{N^2}{2R_a}$$

Hamiltonian: ac with H

$$H|a^\pm(k)\rangle = \sum_{j=0}^N h_a^j \sqrt{\frac{R_a}{R_{b_j}}} \left(C_k^+(l_j) |b_j^\pm(k)\rangle \mp C_k^-(l_j) |b_j^\mp(k)\rangle \right),$$

The matrix elements are

$$\langle b^\tau(k) | H_j | a^\sigma(k) \rangle = \tau^{(\sigma-\tau)/2} h_a^j \sqrt{\frac{N_{b_j}}{N_a}} C_k^{\sigma\tau}(l_j)$$

σ is not a conserved quantum number

- H and T mix $\sigma=+1$ and $\sigma=-1$ states
- the H matrix is twice as large as for momentum states

Why are the semi-momentum states useful then?

Because we can construct a real-valued basis:

Semi-momentum states with parity

This state has definite parity with $p=+1$ or $p=-1$

$$|a^\sigma(k, p)\rangle = \frac{1}{\sqrt{N_a^\sigma}} \sum_{r=0}^{N-1} C_k^\sigma(r)(1 + pP)T^r|a\rangle.$$

- $(k, -1)$ and $(k, +1)$ blocks
- roughly of the same size as original k blocks
- but these states are real, not complex!
- For $k \neq 0, \pi$, the $p=-1$ and $p=+1$ states are degenerate

r	T^r	$T^r P$
0	27 [0 0 0 1 1 0 1 1]	216 [1 1 0 1 1 0 0 0]
1	54 [0 0 1 1 0 1 1 0]	177 [1 0 1 1 0 0 0 1]
2	108 [0 1 1 0 1 1 0 0]	99 [0 1 1 0 0 0 1 1]
3	216 [1 1 0 1 1 0 0 0]	198 [1 1 0 0 0 1 1 0]
4	177 [1 0 1 1 0 0 0 1]	141 [1 0 0 0 1 1 0 1]
5	99 [0 1 1 0 0 0 1 1]	27 [0 0 0 1 1 0 1 1]
6	198 [1 1 0 0 0 1 1 0]	54 [0 0 1 1 0 1 1 0]
7	141 [1 0 0 0 1 1 0 1]	108 [0 1 1 0 1 1 0 0]

P,T transformations

example: $N=8$; note that

- $T^5 P |a\rangle = |a\rangle$

such P,T relationships
will affect normalization
and H-elements

Normalization: We have to check whether or not

$$T^m P|a\rangle = |a\rangle \text{ for some } m \in \{1, \dots, N-1\}$$

Simple algebra gives

$$N_a^\sigma = \frac{N^2}{R_a} \times \begin{cases} 1, & T^m P|a\rangle \neq |a\rangle \\ 1 + \sigma p \cos(km), & T^m P|a\rangle = |a\rangle \end{cases}$$

In the latter case the $\sigma=-1$ and $\sigma=+1$ states are not orthogonal

- calculation of the overlap gives

$$\langle a^\mp(k, p) | a^\pm(k, p) \rangle = -p, \quad (T^m P|a\rangle = |a\rangle \text{ for some } m)$$

Then only one of them should be included in the basis

- convention: **use $\sigma=+1$ if $1+\sigma p \cos(km) \neq 0$, else $\sigma=-1$**

If both $\sigma=+1$ and $\sigma=-1$ are present:

- **we store 2 copies of the same representative**
- we will store the σ value along with the periodicity of the representative

Pseudocode: semi-momentum, parity basis construction

```
do  $s = 0, 2^N - 1$ 
    call checkstate( $s, R, m$ )
    do  $\sigma = \pm 1$  (do only  $\sigma = +1$  if  $k = 0$  or  $k = N/2$ )
        if ( $m \neq -1$ ) then
            if ( $1 + \sigma p \cos(ikm2\pi/N) = 0$ )  $R = -1$ 
            if ( $\sigma = -1$  and  $1 - \sigma p \cos(ikm2\pi/N) \neq 0$ );  $R = -1$ 
        endif
        if  $R > 0$  then  $a = a + 1$ ;  $s_a = s$ ;  $R_a = \sigma R$ ;  $m_a = m$  endif
    enddo
enddo
```

In the subroutine **checkstate()**, we now find whether

$$T^m P |a\rangle = |a\rangle \text{ for some } m \in \{1, \dots, N-1\}$$

$m=-1$ if there is no such transformation

if $\sigma=-1$ we check whether the $\sigma=+1$ state is
orthogonal to the $\sigma=-1$ state or not

- if not orthogonal use $\sigma=+1$
- $R=-1$ for not including

the subroutine **checkstate()**
is modified to gives us:
 • periodicity R (R=-1 if incompatible)
 • m>0 if $T^m \text{Pls} >= \text{ls}$
 • m=-1 if no such relationship

check all translations of ls

construct reflected state Pls

check all translations of Pls

```

subroutine checkstate( $s, R, m$ )
 $R = -1$ 
if ( $\sum_i s[i] \neq n_{\uparrow}$ ) return
 $t = s$ 
do  $i = 1, N$ 
     $t = \text{cyclebits}(t, N)$ 
    if ( $t < s$ ) then
        return
    elseif ( $t = s$ ) then
        if ( $\text{mod}(k, N/i) \neq 0$ ) return
         $R = i$ ; exit
    endif
enddo
 $t = \text{reflectbits}(s, N); m = -1$ 
do  $i = 0, R - 1$ 
    if ( $t < s$ ) then
         $R = -1$ ; return
    elseif ( $t = s$ ) then
         $m = i$ ; return
    endif
     $t = \text{cyclebits}(t, N)$ 
enddo

```

Hamiltonian : Act with an operator H_j on a representative state:

$$H_j |a\rangle = h_a^j P^{q_j} T^{-l_j} |b_j\rangle$$

We can write H acting on a basis state as

$$H|a^\sigma(k, p)\rangle = \sum_{j=0}^N \frac{h_a^j (\sigma p)^{q_j}}{\sqrt{N_a^\sigma}} \sum_{r=0}^{N-1} C_k^\sigma(r + l_j)(1 + pP)T^r |b_j\rangle$$

Using the properties (trigonometry) of the C-functions:

$$\begin{aligned} H|a^\sigma(k, p)\rangle &= \sum_{j=0}^N h_a^j (\sigma p)^{q_j} \sqrt{\frac{N_{b_j}^\sigma}{N_a^\sigma}} \times \\ &\quad \left(\cos(kl_j) |b_j^\sigma(k, p)\rangle - \sigma \sqrt{\frac{N_{b_j}^{-\sigma}}{N_{b_j}^\sigma}} \sin(kl_j) |b_j^{-\sigma}(k, p)\rangle \right) \end{aligned}$$

If, for some m , $T^m P |b_j\rangle = |b_j\rangle$ then

$$\sqrt{\frac{N_{b_j}^{-\sigma}}{N_{b_j}^\sigma}} = \sqrt{\frac{1 - \sigma p \cos(km)}{1 + \sigma p \cos(km)}} = \frac{|\sin(km)|}{1 + \sigma p \cos(km)}$$

$$\langle b_j^\mp(k, p) | b_j^\pm(k, p) \rangle = -p$$

else the ratio is one and the + and – states are orthogonal

The matrix elements are

diagonal in σ

$$\langle b_j^\sigma(k, p) | H_j | a^\sigma(k, p) \rangle = h_a^j (\sigma p)^{q_j} \sqrt{\frac{N_{b_j}^\sigma}{N_a^\sigma}} \times$$

$$\begin{cases} \cos(kl_j), & P|b_j\rangle \neq T^m|b_j\rangle \\ \frac{\cos(kl_j) + \sigma p \cos(k[l_j - m])}{1 + \sigma p \cos(km)}, & P|b_j\rangle = T^m|b_j\rangle \end{cases}$$

off-diagonal in σ

$$\langle b_j^{-\sigma}(k, p) | H_j | a^\sigma(k, p) \rangle = h_a^j (\sigma p)^{q_j} \sqrt{\frac{N_{b_j}^{-\sigma}}{N_a^\sigma}} \times$$

$$\begin{cases} -\sigma \sin(kl_j), & P|b_j\rangle \neq T^m|b_j\rangle, \\ \frac{-\sigma \sin(kl_j) + p \sin(k[l_j - m])}{1 - \sigma p \cos(km)}, & P|b_j\rangle = T^m|b_j\rangle, \end{cases}$$

Pseudocode: semi-momentum, parity hamiltonian

If 2 copies of the same representative, $\sigma=-1$ and $\sigma=+1$:

- do both in the same loop iteration
- examine the previous and next element
- carry out the loop iteration only if representative found for the first time

```
do a = 1, M
    if (a > 1 and sa = sa-1) then
        cycle
    elseif (a < M and sa = sa+1) then
        n = 2
    else
        n = 1
    endif
    ...
enddo
```

n is the number of copies
of the representative

```
do i = a, a + n - 1
    H(a, a) = H(a, a) + Ez
enddo
```

diagonal matrix elements

- **E_z** = diagonal energy

```

 $s = \text{flip}(s_a, i, j)$ 
call representative( $s, r, l, q$ )
call findstate( $r, b$ )
if ( $b \geq 0$ ) then
    if ( $b > 1$  and  $s_b = s_{b-1}$ ) then
         $m = 2; b = b - 1$ 
    elseif ( $b < M$  and  $s_b = s_{b+1}$ ) then
         $m = 2$ 
    else
         $m = 1$ 
    endif
    do  $j = b, b + m - 1$ 
    do  $i = a, a + n - 1$ 
         $H(i, j) = H(i, j) + \text{helement}(i, j, l, q)$ 
    enddo
    enddo
endif

```

construct
off-diagonal
matrix elements

```

subroutine representative( $s, r, l, q$ )
...
 $t = \text{reflectbits}(s, N); q = 0$ 
do  $i = 1, N - 1$ 
     $t = \text{cyclebits}(t, N)$ 
    if ( $t < r$ ) then  $r = t; l = i; q = 1$  endif
enddo

```

helement()
computes the
values based on

- stored info
- and l, q

find the
representative r of s

- translation and
reflection numbers l, q

Using spin-inversion symmetry

Spin inversion operator: $Z|S_1^z, S_2^z, \dots, S_N^z\rangle = | -S_1^z, -S_2^z, \dots, -S_N^z\rangle$

In the magnetization block $m^z=0$ we can use eigenstates of Z

$$|a^\sigma(k, p, z)\rangle = \frac{1}{\sqrt{N_a^\sigma}} \sum_{r=0}^{N-1} C_k^\sigma(r)(1 + pP)(1 + zZ)T^r|a\rangle,$$

$$Z|a^\sigma(k, p, z)\rangle = z|a^\sigma(k, p, z)\rangle, \quad z = \pm 1$$

Normalization: must check how a representative transforms under Z,P,T

- | | | | |
|----|-----------------------------------|---------------------------------|---|
| 1) | $T^m P a\rangle \neq a\rangle$, | $T^m Z a\rangle \neq a\rangle$ | $T^m PZ a\rangle \neq a\rangle$ |
| 2) | $T^m P a\rangle = a\rangle$, | $T^m Z a\rangle \neq a\rangle$ | $T^m PZ a\rangle \neq a\rangle$ |
| 3) | $T^m P a\rangle \neq a\rangle$, | $T^m Z a\rangle = a\rangle$ | $T^m PZ a\rangle \neq a\rangle$ |
| 4) | $T^m P a\rangle \neq a\rangle$, | $T^m Z a\rangle \neq a\rangle$ | $T^m PZ a\rangle = a\rangle$ |
| 5) | $T^m P a\rangle = a\rangle$, | $T^n Z a\rangle = a\rangle$ | $\Rightarrow T^{m+n} PZ a\rangle = a\rangle$ |

For cases 2,4,5 only $\sigma=+1$ or $\sigma=-1$ included

$$N_a^\sigma = \frac{2N^2}{R_a} \times \begin{cases} 1, & 1) \\ 1 + \sigma p \cos(km), & 2) \\ 1 + z \cos(km), & 3) \\ 1 + \sigma pz \cos(km), & 4) \\ [1 + \sigma p \cos(km)][1 + z \cos(kn)], & 5) \end{cases}$$

Hamiltonian: acting on a state gives a transformed representative

$$H_j |a\rangle = h_a^j P^{q_j} Z^{g_j} T^{-l_j} |b_j\rangle$$

$$q_j \in \{0, 1\}, \quad g_j \in \{0, 1\}, \quad l_j = \{0, 1, \dots, N-1\}$$

After some algebra we can obtain the matrix elements

diagonal in σ

$$\langle b_j^\sigma(k, p) | H_j | a^\sigma(k, p) \rangle = h_a^j (\sigma p)^{q_j} z^{g_j} \sqrt{\frac{N_{b_j}^\tau}{N_a^\sigma}} \times$$

$$\begin{cases} \cos(kl_j), & 1), 3) \\ \frac{\cos(kl_j) + \sigma p \cos(k[l_j - m])}{1 + \sigma p \cos(km)}, & 2), 5) \\ \frac{\cos(kl_j) + \sigma pz \cos(k[l_j - m])}{1 + \sigma pz \cos(km)}, & 4) \end{cases}$$

off-diagonal in σ

$$\langle b_j^{-\sigma}(k, p) | H_j | a^\sigma(k, p) \rangle = h_a^j (\sigma p)^{q_j} z^{g_j} \sqrt{\frac{N_{b_j}^\tau}{N_a^\sigma}} \times$$

$$\begin{cases} -\sigma \sin(kl_j), & 1), 3) \\ \frac{-\sigma \sin(kl_j) + p \sin(k[l_j - m])}{1 - \sigma p \cos(km)}, & 2), 5) \\ \frac{-\sigma \sin(kl_j) + pz \sin(k[l_j - m])}{1 - \sigma pz \cos(km)}, & 4) \end{cases}$$

Example: block sizes

$k=0, m_z=0$ (largest block)

N	$(p = \pm 1, z = \pm 1)$			
	$(+1, +1)$	$(+1, -1)$	$(-1, +1)$	$(-1, -1)$
8	7	1	0	2
12	35	15	9	21
16	257	183	158	212
20	2518	2234	2136	2364
24	28968	27854	27482	28416
28	361270	356876	355458	359256
32	4707969	4690551	4685150	4700500

Total spin S conservation

- difficult to exploit
- complicated basis states
- calculate S using $\mathbf{S}^2 = \mathbf{S}(\mathbf{S}+1)$

$$\begin{aligned}\mathbf{S}^2 &= \sum_{i=1}^N \sum_{j=1}^N \mathbf{S}_i \cdot \mathbf{S}_j \\ &= 2 \sum_{i < j} \mathbf{S}_i \cdot \mathbf{S}_j + \frac{3}{4} N\end{aligned}$$

Full diagonalization; expectation values

shorthand block label: $j=(m_z, k, p)$ or $j=(m_z=0, k, p, z)$

$$D_j^{-1} H_j D_j = E_j, \quad \langle n_j | A | n_j \rangle = [D_j^{-1} A D_j]_{nn}$$

$T > 0$: sum over all blocks j and states in block $n=0, M_j-1$

$$\langle A \rangle = \frac{1}{Z} \sum_j \sum_{n=0}^{M_j-1} e^{-\beta E_{j,n}} [D_j^{-1} A_j U_j]_{nn}, \quad Z = \sum_j \sum_{n=0}^{M_j-1} e^{-\beta E_{j,n}}$$

E_j = diagonal (energy) matrix, $E_{j,n}$ = energies, $n=0, \dots, M_j-1$

Full diagonalization limited to small N ; $N=20-24$

Example: Thermodynamics

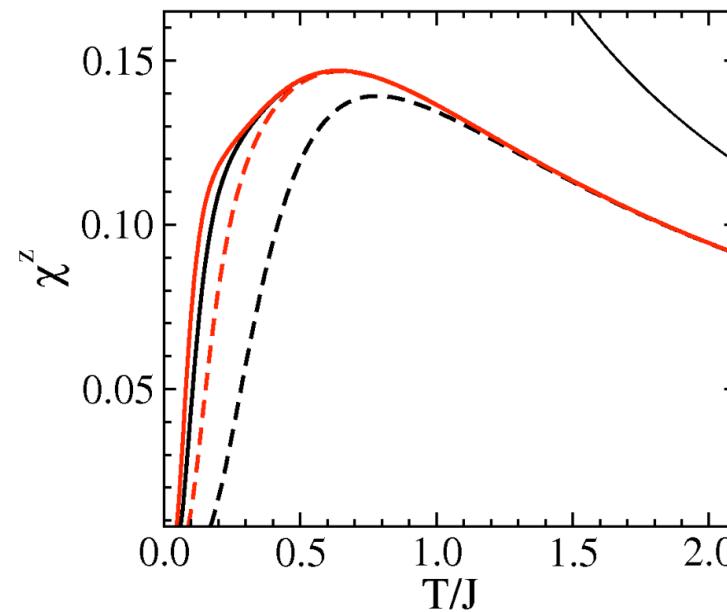
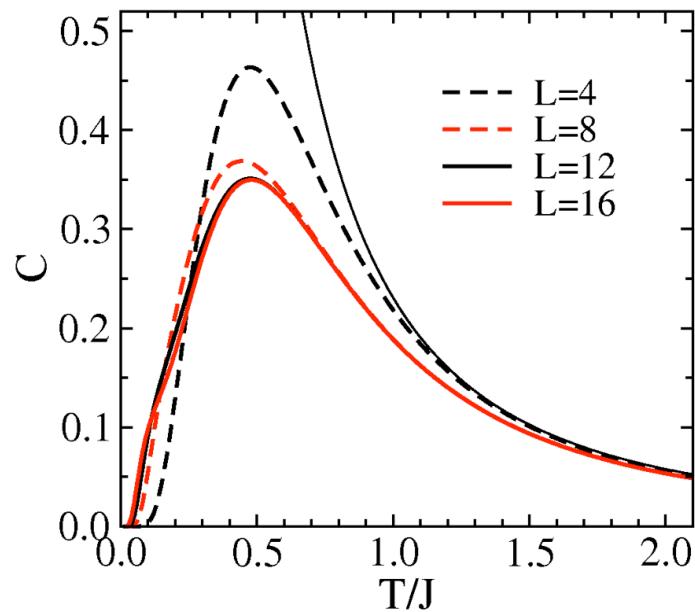
some quantities can be computed using only the magnetization $m_z=0$ sector

- spin-inversion symmetry can be used, smallest blocks
- spin-S state is **(2S+1)**-fold degenerate (no magnetic field) → weight factor
- possible spin dependence of expectation value → average over **$m_z=-S, \dots, S$**

$$C = \frac{d\langle H \rangle}{dt} = \frac{1}{T^2} (\langle H^2 \rangle - \langle H \rangle^2)$$

$$\chi^z = \frac{d\langle m_z \rangle}{dh_z} = \frac{1}{T} (\langle m_z^2 \rangle - \langle m_z \rangle^2)$$

$$\langle m_z \rangle = 0, \quad \langle m_z^2 \rangle = \frac{\langle m_x^2 + m_y^2 + m_z^2 \rangle}{3} = \frac{\langle S^2 \rangle}{3} = \frac{S(S+1)}{3}$$



Compared
with leading
high-T forms
 $\chi = (1/4)/T$
 $C = (3/13)/T^2$

The Lanczos method

If we need only the ground state and a small number of excitations

- can use “Krylov space” methods, which work for much larger matrices
- basis states with 10^7 states or more can be easily handled (30-40 spins)

The Krylov space and “projecting out” the ground state

Start with an arbitrary state $|\Psi\rangle$

- it has an expansion in eigenstates of H ; act with a high power Λ of H

$$H^\Lambda |\Psi\rangle = \sum_n c_n E_n^\Lambda |n\rangle = E_0^\Lambda \left(c_0 |0\rangle + c_1 \left(\frac{E_1}{E_0} \right)^\Lambda |1\rangle + \dots \right)$$

For large Λ , if the state with largest $|E_n|$ dominates the sum

- one may have to subtract a constant, $H - C$, to ensure ground state
- even better to use linear combination of states generated for different Λ

$$|\psi_a\rangle = \sum_{m=0}^{\Lambda} \psi_a(m) H^m |\Psi\rangle, \quad a = 0, \dots, \Lambda$$

- diagonalize H in this basis

In the **Lanczos basis**, H is tridiagonal, convenient to generate and use

- Normally $M=50-200$ basis states is enough; easy to diagonalize H

Constructing the Lanczos basis

First: construct **orthogonal but not normalized basis $\{f_m\}$** . Define

$$N_m = \langle f_m | f_m \rangle, \quad H_{mm} = \langle f_m | H | f_m \rangle$$

The first state $|f_0\rangle$ is arbitrary, e.g., random. The next one is

$$|f_1\rangle = H|f_0\rangle - a_0|f_0\rangle$$

Demand orthogonality

$$\langle f_1 | f_0 \rangle = \langle f_0 | H | f_0 \rangle - a_0 \langle f_0 | f_0 \rangle = H_{00} - a_0 N_0 \rightarrow a_0 = H_{00}/N_0$$

The next state and its overlaps with the previous states

$$|f_2\rangle = H|f_1\rangle - a_1|f_1\rangle - b_0|f_0\rangle$$

$$\langle f_2 | f_1 \rangle = H_{11} - a_1 N_1, \quad \langle f_2 | f_0 \rangle = N_1 - b_0 N_0$$

For orthogonal states

$$a_1 = H_{11}/N_1, \quad b_0 = N_1/N_0$$

All subsequent states are constructed according to

$$|f_{m+1}\rangle = H|f_m\rangle - a_m|f_m\rangle - b_{m-1}|f_{m-1}\rangle$$

$$a_m = H_{mm}/N_m, \quad b_{m-1} = N_m/N_{m-1}$$

Easy to prove orthogonality of all these states ($\langle f_{m+1} | f_m \rangle = 0$ is enough)

The hamiltonian in the Lanczos basis

Rewrite the state generation formula

$$H|f_m\rangle = |f_{m+1}\rangle + a_m|f_m\rangle + b_{m-1}|f_{m-1}\rangle$$

Because of the orthogonality, the only non-0 matrix elements are

$$\langle f_{m-1}|H|f_m\rangle = b_{m-1}N_{m-1} = N_m$$

$$\langle f_m|H|f_m\rangle = a_mN_m$$

$$\langle f_{m+1}|H|f_m\rangle = N_{m+1}$$

But the f-states or not normalized. The normalized states are:

$$|\phi_m\rangle = \frac{1}{\sqrt{N_m}}|f_m\rangle$$

In this basis the H-matrix is

$$\langle \phi_{m-1}|H|\phi_m\rangle = \sqrt{b_{m-1}}$$

$$\langle \phi_m|H|\phi_m\rangle = a_m$$

$$\langle \phi_{m+1}|H|\phi_m\rangle = \sqrt{b_m}$$

Potential problem:

The normalization constants N_m can become very large (think of E_0^\wedge)

Solution:

generate the normalized basis directly

- start with $|\Phi_0\rangle$ arbitrary, normalized, and then

$$|\phi_1\rangle = \frac{1}{N_1} (H|\phi_0\rangle - a_0|\phi_0\rangle).$$

$$|\phi_{m+1}\rangle = \frac{1}{N_{m+1}} (H|\phi_m\rangle - a_m|\phi_m\rangle - N_m|\phi_{m-1}\rangle) = \frac{|\gamma_m\rangle}{N_{m+1}}$$

The definition of N_m is different, and no b_m :

$$a_m = \langle \phi_m | H | \phi_m \rangle$$

$$N_m = \langle \gamma_m | \gamma_m \rangle^{-1/2}$$

Generate $|\gamma_m\rangle$ first, normalize to get N_{m+1}

The H-matrix is

$$\langle \phi_{m-1} | H | \phi_m \rangle = N_m$$

$$\langle \phi_m | H | \phi_m \rangle = a_m$$

$$\langle \phi_{m+1} | H | \phi_m \rangle = N_{m+1}$$

Lanczos basis generation in practice

Here: generate the orthogonal basis $\{\Phi_m\}$ directly

$$|\phi_m\rangle = \sum_{a=1}^M \phi_m(a)|a\rangle, \quad m = 0, \dots, \Lambda$$

in a given symmetry block of size M

The coefficients $\Phi_m(a)$ are stored as $\Lambda+1$ vectors of size M

- may store only the vectors Φ_{m-1} and Φ_m to generate Φ_{m+1}
 - but basis has to be re-generated when computing expectation values
 - stabilization by “re-orthogonalization” (later) requires storage of all Φ_m

The main computational effort is in acting with the hamiltonian; $H|\Phi_m\rangle$

- implement as a subroutine **hoperation**(ϕ, γ), where $|\gamma\rangle = H|\phi\rangle$
- state normalization implemented as **normalize**(ϕ, n)
 - ϕ = vector to normalize, $n = \langle \phi | \phi \rangle$ before normalization

Pseudocode; Lanczos basis generation

Initial random state

```
do i = 1, M
     $\phi_0(i)$  = randomfloat(0, 1)
enddo
call normalize( $\phi_0$ ,  $n_0$ )
```

second state

```
call hoperation( $\phi_0$ ,  $\phi_1$ )
 $a_0 = \langle \phi_0 | \phi_1 \rangle$ ;  $\phi_1 = \phi_0 - a_0 | \phi_1 \rangle$ 
call normalize( $\phi_1$ ,  $n_1$ )
```

Generate the rest of the states

```
do m = 1,  $\Lambda - 1$ 
    call hoperation( $\phi_m$ ,  $\phi_{m+1}$ )
     $a_m = \langle \phi_m | \phi_{m+1} \rangle$ 
     $\phi_{m+1} = \phi_{m+1} - a_m \phi_m - n_m \phi_{m-1}$ 
    call normalize( $\phi_{m+1}$ ,  $n_{m+1}$ )
enddo
```

Note: the H-matrix can be constructed and diagonalized after each step

- follow evolution of energy versus Λ
- stop based on some convergence criterion on E_0 (or higher energy)
- expectation values converge slower than energies

The subroutine **hoperation**(ϕ, γ) implements

$$H|\phi\rangle = |\gamma\rangle = \sum_{a=1}^M \sum_{b=1}^M \phi(a) \langle b|H|a\rangle |b\rangle$$

in a given symmetry block (M = block size)

We do not want to store H as an $M \times M$ matrix (too big). Two options:

- carry out the operations on the fly; only the vectors are stored
- store H in a compact form; only non-0 elements (sparse matrix)

Storing H speeds up the Lanczos iterations

- but may require a lot of memory

Compact storage of H : For each $a=1, M$

- e_a is the number of non-0 elements $\langle b|H|a\rangle$
- labels $i=s_a+1, s_a+e_a$ will refer to these matrix elements; $s_a = \sum_{c=1}^{a-1} e_a$
- $H(i)$ contains the values of the matrix elements $\langle b|H|a\rangle$
- $B(i)$ contains the corresponding “target” state index b
- The hamiltonian is symmetric
 - store only elements with $b \leq a$ (divide diagonal elements by 2)

Pseudocode; hamiltonian operation with compact storage

```
subroutine hoperation( $\phi, \gamma$ )
 $\gamma = 0; i = 0$ 
do  $a = 1, M$ 
    do  $j = 1, e_a$ 
         $i = i + 1$ 
         $\gamma(B(i)) = \gamma(B(i)) + H(i)\phi(a)$ 
         $\gamma(a) = \gamma(a) + H(i)\phi(B(i))$ 
    enddo
enddo
```

$$H|\phi\rangle = |\gamma\rangle = \sum_{a=1}^M \sum_{b=1}^M \phi(a) \langle b | H | a \rangle |b\rangle$$

Further storage compactification possible

- small number of different elements
- use mapping $\langle b | H | a \rangle \rightarrow \text{integer}$
- many operations on $|a\rangle$ give same $|b\rangle$
 - add up all contributions before storing

Operator expectation values

Diagonalizing the tri-diagonal matrix → eigenstates in the Lanczos basis

- eigenvectors \mathbf{v}_n , energies E_n
- only some number of low-energy states ($<< \Lambda$) are correct eigenstates of H

To compute expectation values we normally go back to the original basis

$$\psi_n(a) = \sum_{m=0}^{\Lambda} v_n(m) \phi_m(a), \quad a = 1, \dots, M$$

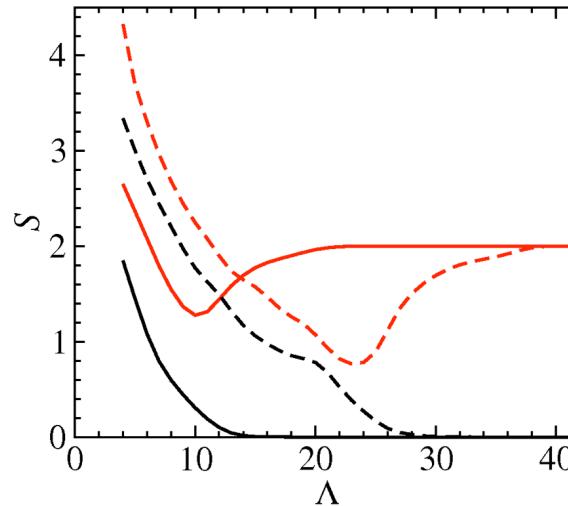
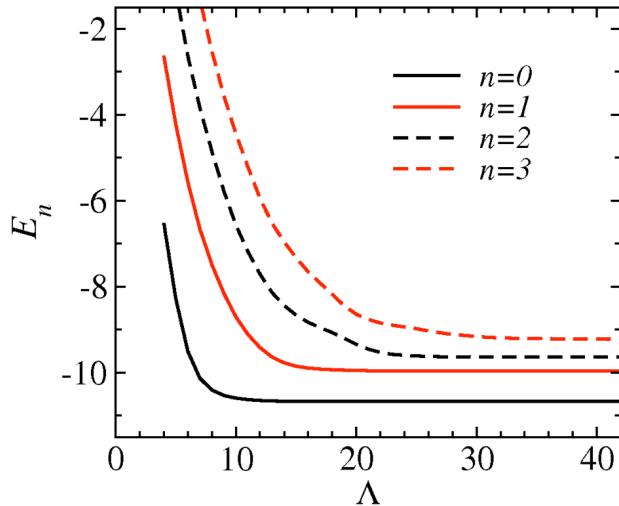
To compute $\langle \psi_n | O | \psi_n \rangle$ first construct

$$\begin{aligned} O|\psi_n\rangle &= |\psi_n^O\rangle = \sum_{a=1}^M \psi_n(a) O|a\rangle \\ &= \sum_{a=1}^M \sum_{b=1}^M \psi_n(a) |b\rangle \langle b| O|a\rangle \\ &= \sum_{b=1}^M \psi_n^O(b) |b\rangle \quad \psi_n^O(b) = \sum_{a=1}^M \psi_n(a) \langle b| O|a\rangle \end{aligned}$$

Then evaluate the scalar product

$$\langle \psi_n | O | \psi_n \rangle = \langle \psi_n | \psi_n^O \rangle = \sum_{a=1}^M \psi_n(a) \psi_n^O(a)$$

Convergence properties of the Lanczos method

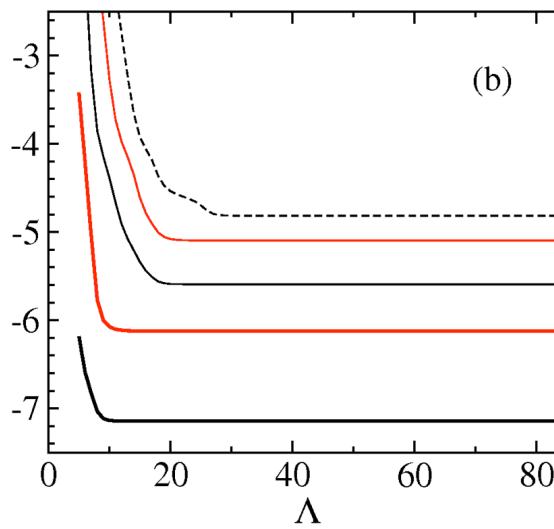
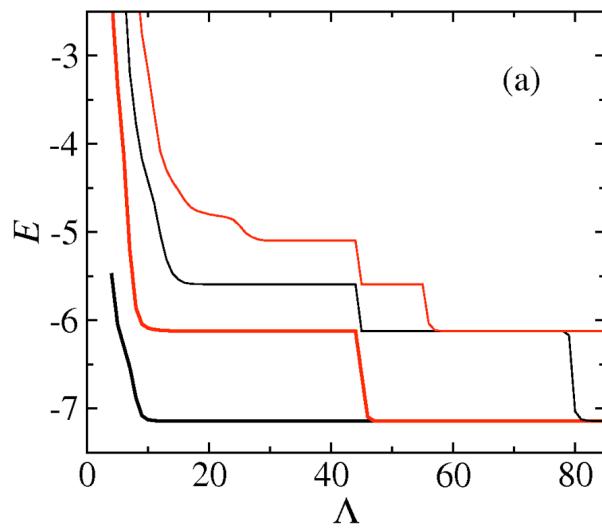


Example: 24-site chain
 $m_z = 0, k = 0, p = 1, z = 1$
block size $M = 28416$

Ground state converges first, then successively excited states

Loss of orthogonality: accumulation of numerical error \rightarrow basis becomes non-orthogonal

- higher states collapse down onto lower ones
- can be cured with re-orthogonalization



Example: 16-site chain
 $m_z = 0, k = 0, p = 1, z = 1$
block size $M = 212$

- (a) non-orthogonality
- (b) re-orthogonalized

Re-orthogonalization procedure

For each state generated, remove all components of prior states, $i=1,\dots,m$

- easy if we work with the normalized basis and all states are stored

$$|\phi_m\rangle \rightarrow \frac{|\phi_m\rangle - q|\phi_i\rangle}{1 - q^2}, \quad q = \langle\phi_i|\phi_m\rangle$$

Pseudocode: modify state generation

```
do m = 1, Λ - 1
    call hoperation(φm, φm+1)
    am = ⟨φm|φm+1⟩; φm+1 = φm+1 - amφm - nmφm-1
    call normalize(φm+1, nm+1)
    do i = 1, m
        q = ⟨φm+1|φi⟩; φm+1 = (φm+1 - qφi)/(1 - q2)
    enddo
enddo
```

Note: the Lanczos method can only generate a single state of a multiplet

- some random linear combination of degenerate states

Example: 2 degenerate states i, j:

$$H^\Lambda |\Psi\rangle = \sum_{m \neq i, j} c_m E_m^\Lambda |\psi_m\rangle + E_{i,j}^m (c_i |\psi_i\rangle + c_j |\psi_j\rangle)$$

The mixing of the duplet is determined by c_i, c_j of the initial state

Spin correlations in the Heisenberg chain

Let's look at the (staggered) spin correlation function

$$C(r) = \langle \mathbf{S}_i \cdot \mathbf{S}_{i+r} \rangle (-1)^r$$

versus the distance r and at $r=N/2$ versus system size N

Theory (bosonization conformal field theory) predicts (for large r, N)

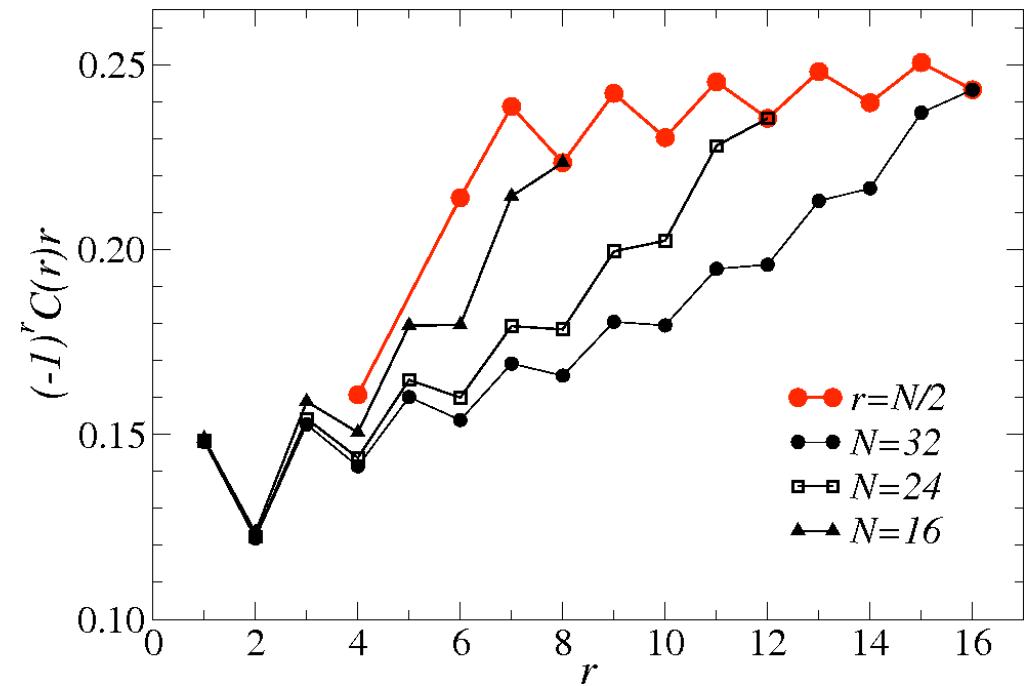
$$C(r) \propto \frac{\ln^{1/2}(r/r_0)}{r}$$

Plausible based on N up to 32

- other methods for larger N

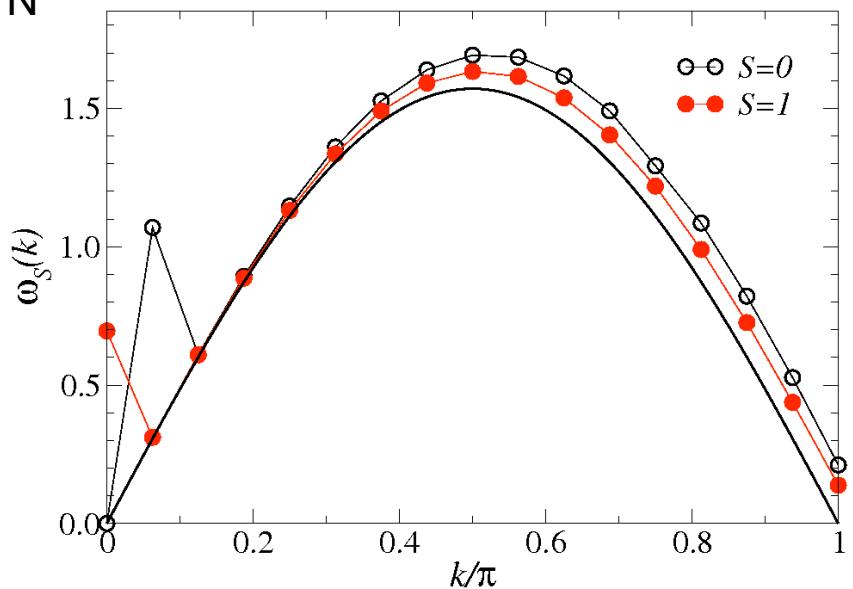
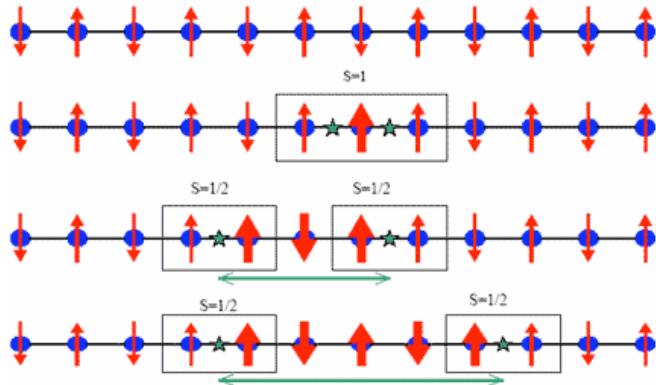
Power-law correlations are a sign of a “critical” state;
at the boundary between

- ordered (antiferromagnetic)
- disordered (spin liquid)



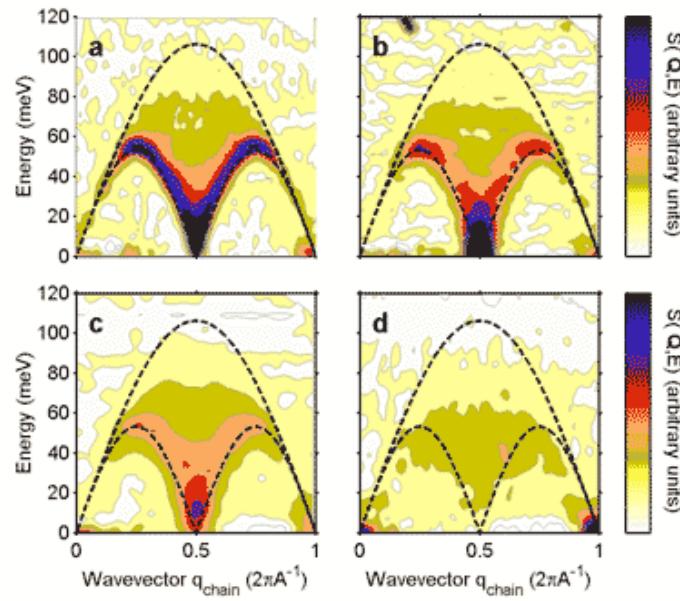
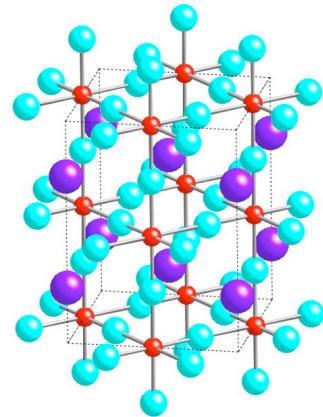
Excitations of the Heisenberg chain

- the ground state is a singlet ($S=0$) for even N
- the first excited state is a triplet ($S=1$)
- can be understood as pair of “spinons”



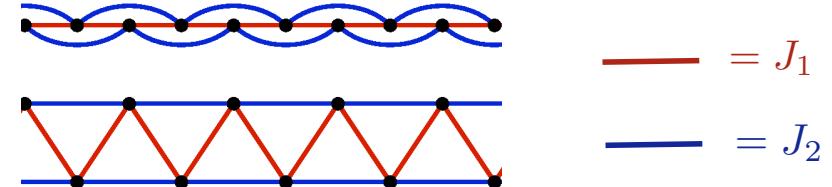
Neutron scattering experiments

- quasi-one-dimensional KCuF_3
- B. Lake et al., Nature Materials 4 329-334 (2005)



Heisenberg chain with frustrated interactions

$$H = \sum_{i=1}^N [J_1 \mathbf{S}_i \cdot \mathbf{S}_{i+1} + J_2 \mathbf{S}_i \cdot \mathbf{S}_{i+2}]$$

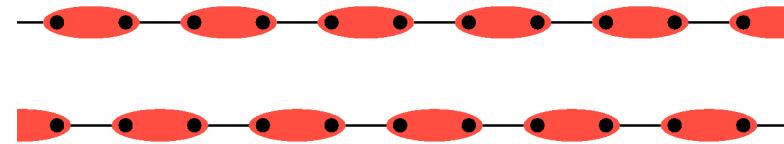


For the special point $J_2/J_1=0.5$, this model has an exact solution

Singlet-product states

$$|\Psi_A\rangle = |(1, 2)(3, 4)(5, 6) \dots\rangle$$

$$|\Psi_B\rangle = |(1, N)(3, 2)(5, 4) \dots\rangle$$



It is not hard to show that these are eigenstates of H (we will do later)

$$(a, b) = (\uparrow_a \downarrow_b - \downarrow_a \uparrow_b)/\sqrt{2}$$

The system has this kind of order (with fluctuations, no exact solution) for all $J_2/J_1 > 0.2411\dots$. This is a **quantum phase transition** between

- a critical state
- a valence-bond-solid (VBS) state

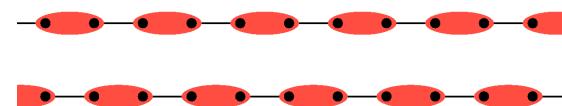
The symmetry is not broken for finite N

- the ground state is a superposition of the two ordered states

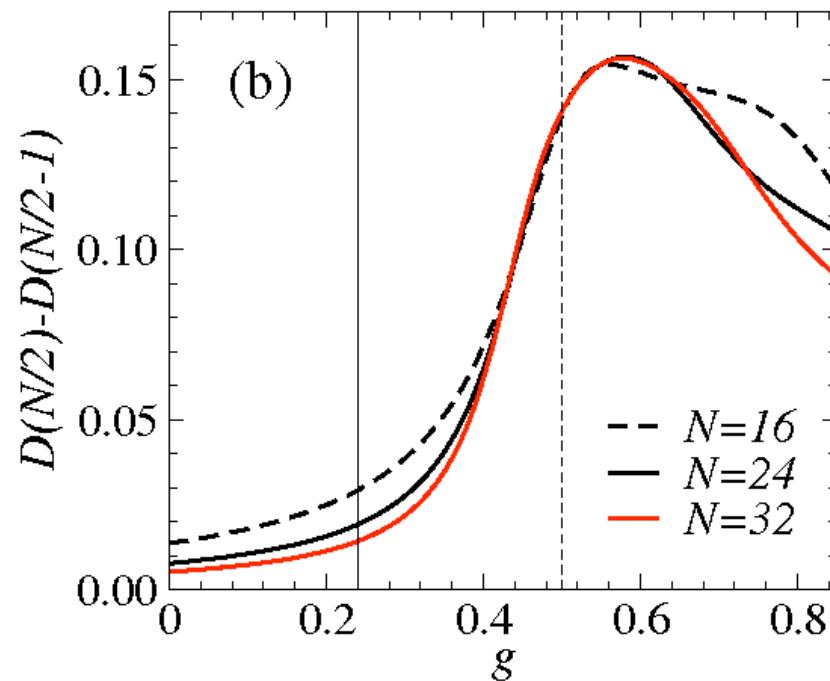
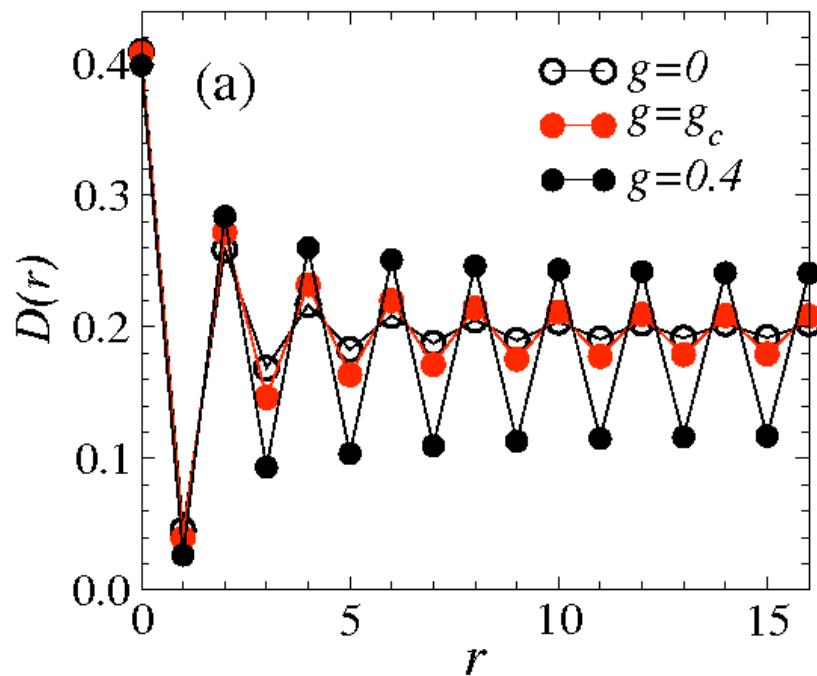
$$|\Psi_0\rangle \sim |\Psi_A\rangle + |\Psi_B\rangle, \quad |\Psi_1\rangle \sim |\Psi_A\rangle - |\Psi_B\rangle$$

The VBS state can be detected in finite systems using “dimer” correlations

$$D(r) = \langle B_i B_{i+r} \rangle = \langle (\mathbf{S}_i \cdot \mathbf{S}_{i+1}) (\mathbf{S}_{i+r} \cdot \mathbf{S}_{i+1+r}) \rangle$$



Results from Lanczos diagonalization; different coupling ratios $g=J_2/J_1$



It is not easy to detect the transition this way

- much larger systems are needed for observing a sharp transition
- other properties can be used to accurately determine the critical point g_c

Determining the transition point using level crossings

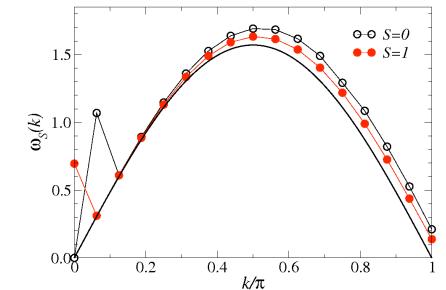
Lowest excitation for the $g=0$ Heisenberg chain is a triplet

- this can be expected for all $g < g_c$

The VBS state is 2-fold degenerate for infinite N

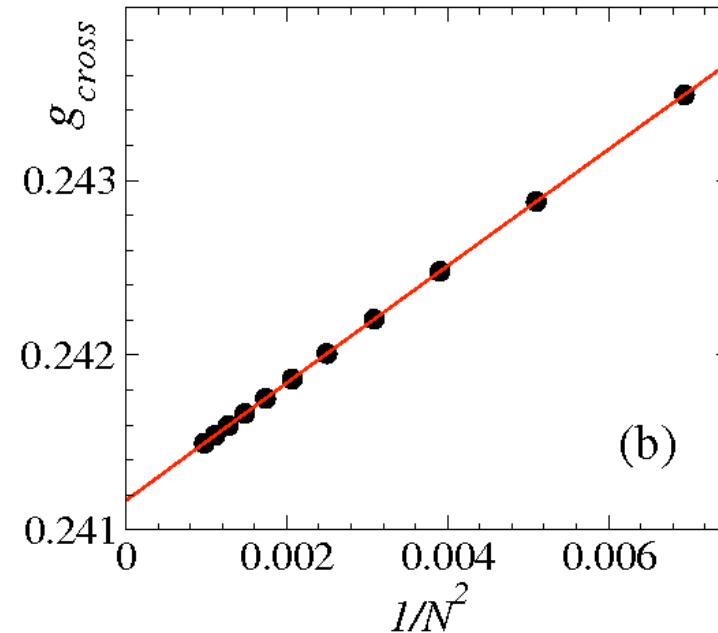
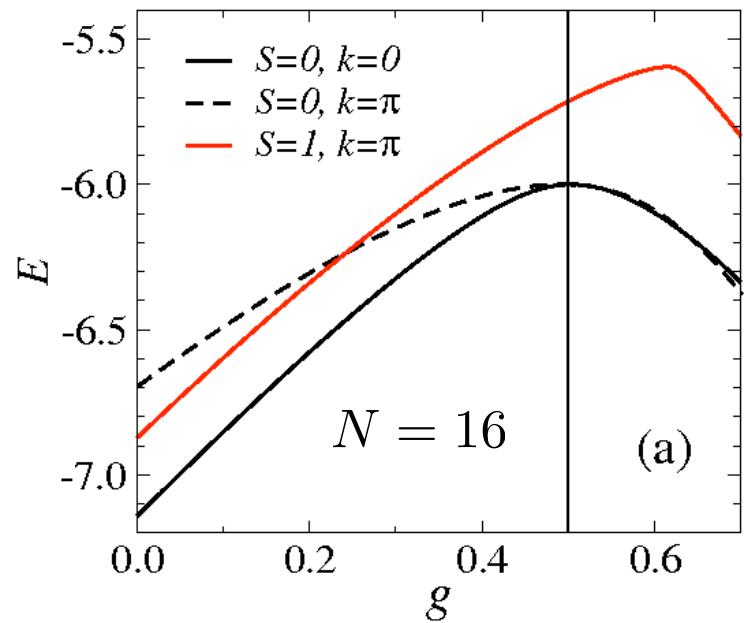
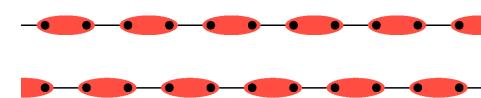
- and for any N at $g=1/2$
- these two states are singlets
- gap between them closes exponentially as $N \rightarrow \infty$
- the lowest excitation is the second singlet

The two lowest excited state should cross at g_c



$$|\Psi_0\rangle \sim |\Psi_A\rangle + |\Psi_B\rangle$$

$$|\Psi_1\rangle \sim |\Psi_A\rangle - |\Psi_B\rangle$$



Extrapolating point for different N up to 32 gives $g_c = 0.2411674(2)$