

## **LINEAR LOGIC Homework #2**

two separate Lemmas for Substitution

Jiawen LIU

UBID: jliu223      person #: 50245965

Types	$T ::= A B C$
Term	$t ::= x t  \lambda x.t$

Figure 1: syntax

$$\boxed{\Gamma \vdash t : A}$$

$$\begin{array}{c}
\frac{}{x : A \vdash x : A} \text{AXIOM} \qquad \frac{\Gamma, x : A \vdash t : B}{\Gamma \vdash \lambda x.t : A \rightarrow B} \text{ABS} \\
\\
\frac{\Gamma \cap \Delta = \emptyset \quad \Gamma \vdash t : A \rightarrow B \quad \Delta \vdash u : A}{\Gamma, \Delta \vdash t u : B} \text{APP} \qquad \frac{\Gamma, x : A, y : B \vdash u : C}{\Gamma, y : B, x : A \vdash u : C} \text{EXCHANGE} \\
\\
\frac{\Gamma \vdash t : B}{\Gamma, x : !A \vdash t : B} \text{WEAKING} \qquad \frac{\Gamma, x : A \vdash t : B}{\Gamma, x : !A \vdash t : B} \text{DERELICTION} \\
\\
\frac{\Gamma, x : !A, y : !A \vdash t : C}{\Gamma, x : !A \vdash t[x/y] : C} \text{CONTRACT} \qquad \frac{\Gamma \vdash t : A}{\Gamma \vdash t : !A} \text{PROMOTION}
\end{array}$$

Figure 2: Typing rules

**Lemma 1** (Preservation under Substitution).

1. If  $\Gamma \vdash t : A$  and  $\Delta, x : A \vdash u : B$  then  $\Gamma, \Delta \vdash u[t/x] : B$ .
2. If  $\Gamma \vdash t : !A$  and  $\Delta, x : !A \vdash u : B$  then  $\Gamma, \Delta \vdash u[t/x] : B$ .

*Proof.* Lemma 1.1 is proved in lecture.

Lemma 1.2 is proved by induction on the typing derivation of the second premise  $\Delta, x : !A \vdash u : B$ .

Assume we know:  $\Gamma \vdash t : !A$

**Case**

**SubCase1:**

$$\frac{\Delta \vdash u : B}{\Delta, x : !A \vdash u : B} \text{WEAKING}$$

TS:  $\Gamma, \Delta \vdash u[t/x] : B$

STS:  $\Gamma, \Delta \vdash u : B$ , because  $x$  doesn't show up in  $u$ .

By applying WEAKING rule on hypothesis, we get:  $\Delta, \Gamma \vdash u : B$ .

By applying WEAKING rule, we get:  $\Gamma, \Delta \vdash u : B$ .

This case is proved.

**SubCase2:**

$$\frac{\Delta, x : !A \vdash u : B}{\Delta, x : !A, \hat{x} : !A \vdash u : B} \text{WEAKING}$$

TS:  $\Gamma, \Delta, \hat{x} : !A \vdash u[t/x] : B$

By IH, we get:  $\Gamma, \Delta \vdash u[t/x] : B$

By apply WEAKING rule, we get:  $\Gamma, \Delta, \hat{x} : !A \vdash u[t/x] : B$

This case is proved.

**Case**

**SubCase1:**

$$\frac{\Delta, x : A \vdash u : B}{\Delta, x :!A \vdash u : B} \text{DERELICTION}$$

TS:  $!\Gamma, \Delta \vdash u[t/x] : B$ .

By IH of Lemma 1.1, we get:  $\Gamma, \Delta \vdash u[t/x] : B$ .

By applying the DERELICTION rule, we get:  $!\Gamma, \Delta \vdash u[t/x] : B$ .

This case is proved.

**SubCase2:**

$$\frac{\Delta, x :!A, \hat{x} : A \vdash u : B}{\Delta, x :!A, \hat{x} :!A \vdash u : B} \text{DERELICTION}$$

TS:  $!\Gamma, \Delta, \hat{x} :!A \vdash u[t/x] : B$ .

By IH, we get:  $!\Gamma, \Delta, \hat{x} : A \vdash u[t/x] : B$ .

By applying DERELICTION rule, we get:  $!\Gamma, \Delta, \hat{x} :!A \vdash u[t/x] : B$ .

This case is proved.

**Case**

**SubCase1:**

$$\frac{\Delta, x :!A, y :!A \vdash v : B}{\Delta, x :!A \vdash v[x/y] : B} \text{CONTRACT}$$

TS:  $!\Gamma, \Delta \vdash u[t/x] : B$

STS:  $!\Gamma, \Delta \vdash v[x/y][t/x] : B$

STS:  $!\Gamma, \Delta \vdash v[t/y][t/x] : B$

By IH on  $\Delta, x :!A, y :!A \vdash v : B$ , we get:

$!\Gamma, \Delta, y :!A \vdash v[t/x] : B$

**[[ Rename all the variables in premise 1:  $!\Gamma \vdash t :!A$ , we get:**

$!\Gamma' \vdash t[\Gamma'/\Gamma] :!A$

**Apply Induction Hypothesis again on  $!\Gamma, \Delta, y :!A \vdash v[t/x] : B$  with premise  $!\Gamma' \vdash t[\Gamma'/\Gamma] :!A$ , we get:**

$!\Gamma', !\Gamma, \Delta \vdash v[t/x][t[\Gamma'/\Gamma]/y] : B$

**Apply CONTRACT rule on every pair of same variables in  $\Gamma', \Gamma$  with different name, we get:**

$!\Gamma, \Delta \vdash v[t/x][t[\Gamma'/\Gamma]/y][\Gamma/\Gamma'] : B$ .

**Because  $v[t/x][t[\Gamma'/\Gamma]/y][\Gamma/\Gamma'] = v[t/x][t/y] = v[t/y][t/x]$ . ]]**

this case is proved.

**SubCase2:**

$$\frac{\Delta, x :!A, \hat{x} :!A, y :!A \vdash v : B}{\Delta, x :!A, \hat{x} :!A \vdash v[\hat{x}/y] : B} \text{CONTRACT}$$

where  $v[\hat{x}/y] = u$ .

TS:  $!\Gamma, \Delta, \hat{x} :!A \vdash u[t/x] : B$

STS:  $!\Gamma, \Delta, \hat{x} :!A \vdash v[\hat{x}/y][t/x] : B$

We know  $\Delta, \hat{x} :!A, y :!A, x :!A \vdash v : B$  by applying EXCHANGE rule on  $\Delta, x :!A, \hat{x} :!A, y :!A \vdash v : B$ .

By IH on  $\Delta, \hat{x} :!A, y :!A, x :!A \vdash v : B$ , we get:

$!\Gamma, \Delta, \hat{x} :!A, y :!A \vdash v[t/x] : B$

By applying CONTRACT rule on  $!\Gamma, \Delta, \hat{x} :!A, y :!A \vdash v[t/x] : B$ , we get:

$!\Gamma, \Delta, \hat{x} :!A \vdash v[t/x][\hat{x}/y] : B$ .

Because  $v[t/x][\hat{x}/y] \equiv v[\hat{x}/y][t/x]$ , this case is proved.

**Case**

$$\frac{!\Delta, x :!A \vdash u : C}{!\Delta, x :!A \vdash u :!C} \text{PROMOTION}$$

where  $B =!C$ .

TS:  $!\Gamma, !\Delta \vdash u[t/x] : B$ .

STS:  $!\Gamma, !\Delta \vdash u[t/x] :!C$

By IH, we get:  $!\Gamma, !\Delta \vdash u[t/x] : C$

By applying the PROMOTION rule, we get:  $!\Gamma, !\Delta \vdash u[t/x] : !C$ .

This case is proved. □

**Theorem 0.1** (Preservation). *If  $\Gamma \vdash t : A$  and  $t \rightarrow t'$  then  $\vdash t' : A$*

*Proof.* of Theorem 0.1 by induction on the derivation of first premise:  $\Gamma \vdash t : A$ : □