

⇒ CMU: Claire Le Goues

interprocedural control-flow analysis.

⇒ example: flow, in function call & return:

$$\text{flow}(x \leftarrow g(y)) \sigma = [x \rightarrow L_r] \sigma, \text{ where } \sigma(y) \in L_a.$$

$$\text{flow}(\text{return}) \sigma = \sigma.$$

$$\text{where } \sigma(x) \in L_r.$$

if  $L_a = L_r = T$ :

fun:  $g(x) : \text{int}$

$y = 10/x.$

return  $y.$

$z = 5.$

$w = g(z).$

$$\sigma: [w \rightarrow L_r] \sigma_0 \text{ where } \sigma_0(z) \in L_a.$$

$$\Rightarrow \Rightarrow \sigma_0 = \{z \rightarrow L_a\} \sigma_1^*$$

$$[\text{return } y] \sigma_{z1} = \sigma_1 : \sigma_1(y) \in L_r.$$

$$\Rightarrow \sigma_1 = \{y \rightarrow L_r\}.$$

$$\Rightarrow \sigma = \{w \rightarrow L_r, z \rightarrow L_a, y \rightarrow L_r, \quad z \leftarrow x, x \leftarrow y.\}$$

$$y \leftarrow w.$$

⇒ Harvard: Stephen Chong:

⇒ call graph: ⇒ 1 big CFG with call graph.

⇒ interprocedural CFG. ⇒ treat argument, return value as assignments.

⇒ problem: all different calls of a function, its flow paths are merged

⇒ inline: copy a new function's CFG every time it is called.

⇒ problem: no recursive call

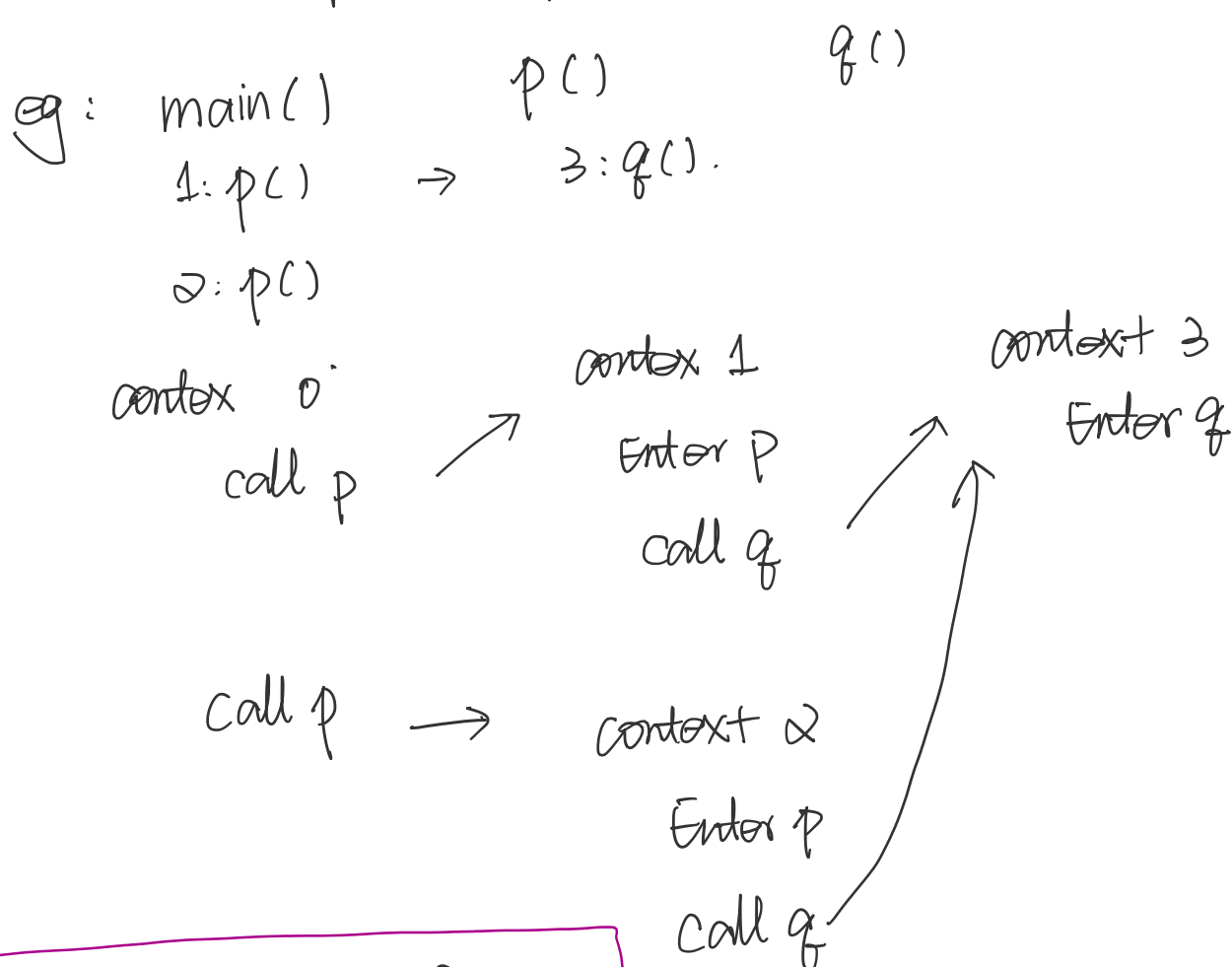
size increase exponentially.

⇒ context Sensitive analysis.

Syntactic Context:

⇒ only produce 1 copy if the nested call

or multiple call of a function isn't explicitly inlining in program.



Call-site Stack Context Sensitive

context 0 → context 0:1 → context 0:1:3

→ context 0:2 → context 0:2:3.

similar for nested call.

others: caller stack. less precise.