

COE 115



Lecture 3

Indirect Addressing

Mov with indirect Addressing:

$\text{mov}\{\text{.b}\} [\text{Wso}], [\text{Wdo}] \quad ((\text{Wso})) \rightarrow (\text{Wdo})$

[] (brackets) indicate indirect addressing.

Source Effective Address (EAs) is the content of Wso, or (Wso). Destination Effective Address (EAd) is the content of Wdo, or (Wdo).

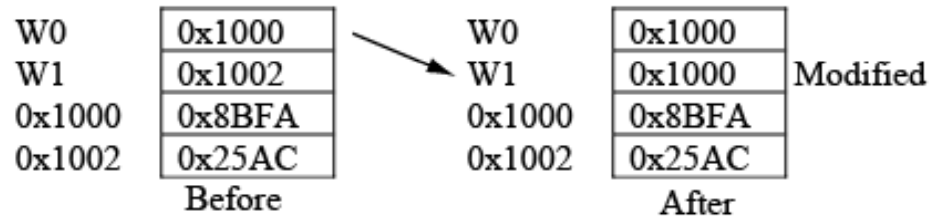
The MOV instruction copies the content of the Source Effective Address to the Destination Effect Address, or:

$(\text{EAs}) \rightarrow \text{EAd}$

which is:

$((\text{Wso})) \rightarrow (\text{Wdo})$

(a) Execute: `mov W0, W1` source, destination use register direct

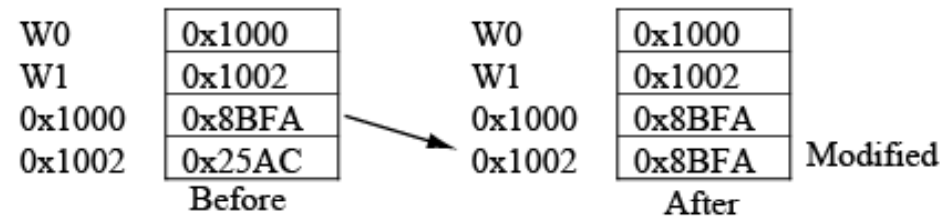


(b) Execute: `mov [W0], [W1]` source, destination use register indirect

Source Effective Address = (W0) = 0x1000

Destination Effective Address = (W1) = 0x1002

Operation is (0x1000) → 0x1002

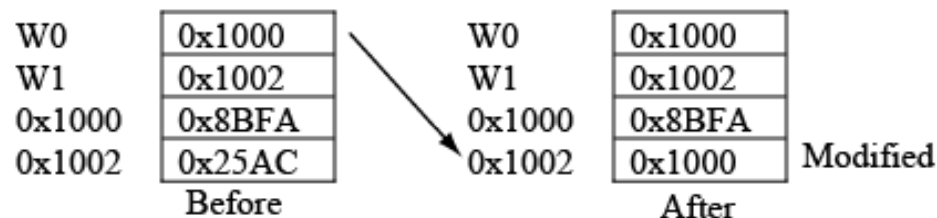


(c) Execute: `mov W0, [W1]` source uses register direct
destination uses register indirect

Source Effective Address = W0

Destination Effective Address = (W1) = 0x1002

Operation is (W0) → 0x1002



Indirect Addressing MOV Example

Why Indirect Addressing?

The instruction:

```
mov [W0], [W1]
```

Allows us to do a memory-memory copy with one instruction!

The following is illegal:

```
mov 0x1000, 0x1002
```

Instead, would have to do:

```
mov 0x1000, W0
```

```
mov W0, 0x1002
```

Indirect Addressing Coverage

- **There are six forms of indirect addressing**
- **The need for indirect addressing makes the most sense when covered in the context of C pointers**
- ***register indirect* the simplest form of indirect addressing, which is as shown on the previous slides.**
- **Most instructions that support register direct for an operand, also support indirect addressing as well for the same operand**
 - However, must check PIC24 datasheet and book to confirm.

ADD {.B} Wb, Ws, Wd Instruction

- **Three operand addition, register-to-register form:**

ADD{.B} Wb, Ws, Wd $(Wb) + (Ws) \rightarrow Wd$

Wb, Ws, Wd are any of the 16 working registers *W0-W15*

ADD W0, W1, W2

$(W0) + (W1) \rightarrow W2$

ADD W2, W2, W2

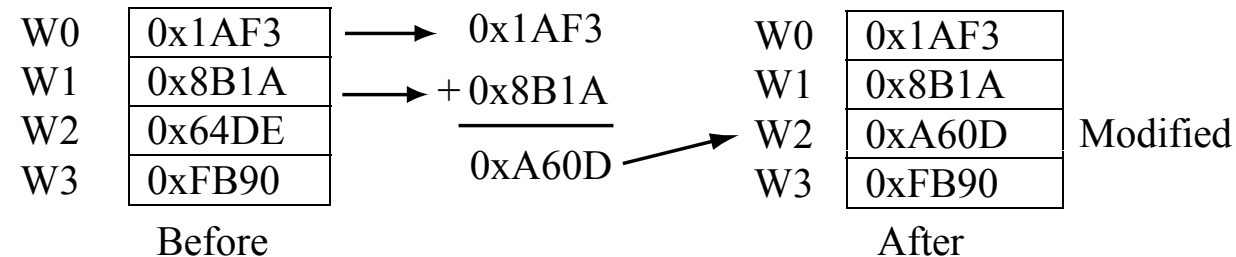
$W2 = W2 + W2 = W2 * 2$

ADD.B W0, W1, W2

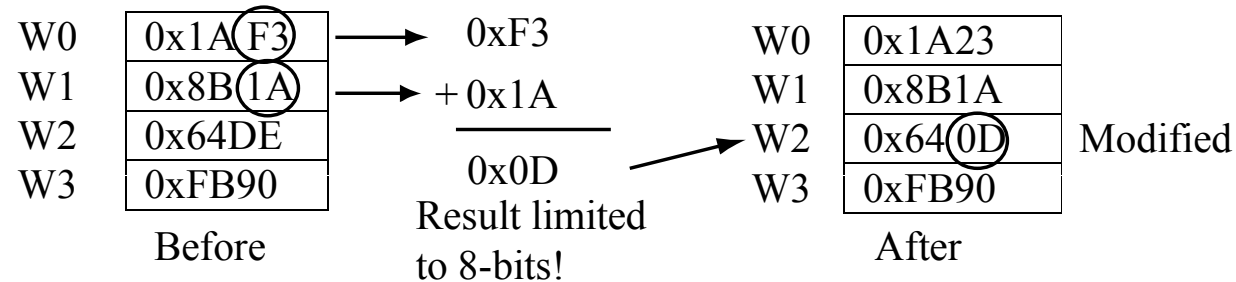
Lower 8 bits of *W0, W1* are added
and placed in the lower 8 bits of *W2*

ADD {.B} Wb, Ws, Wd Instruction Execution

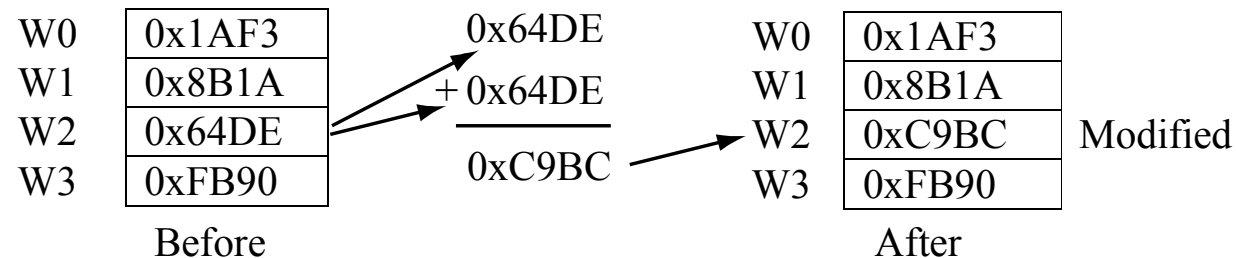
(a) Execute: `add W0,W1,W2`



(b) Execute: `add.b W0,W1,W2`



(c) Execute: `add W2,W2,W2`



SUB{.B} Wb, Ws, Wd Instruction

Three operand subtraction, register-to-register form:

SUB{.B} Wb, Ws, Wd (Wb) – (Ws) → Wd

Wb, Ws, Wd are any of the 16 working registers W0-W15.

Be careful:

while ADD Wx, Wy, Wz gives the same result as ADD Wy, Wx, Wz

The same is not true for

SUB Wx, Wy, Wz versus SUB Wy, Wx, Wz

SUB W0, W1, W2 (W0) – (W1) → W2

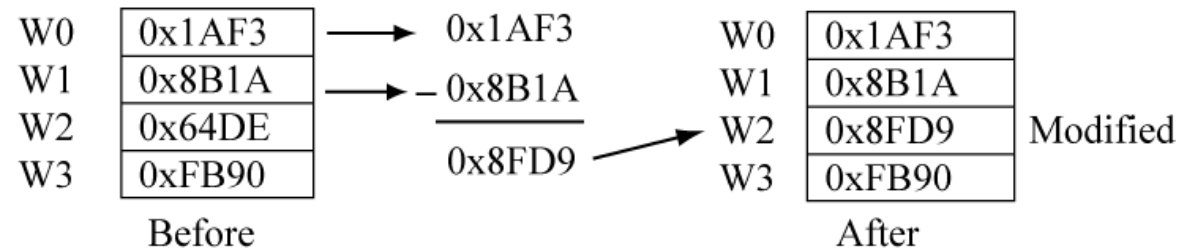
SUB W1, W0, W2 (W1) – (W0) → W2

SUB.B W0, W1, W2
lower 8-bits of W2

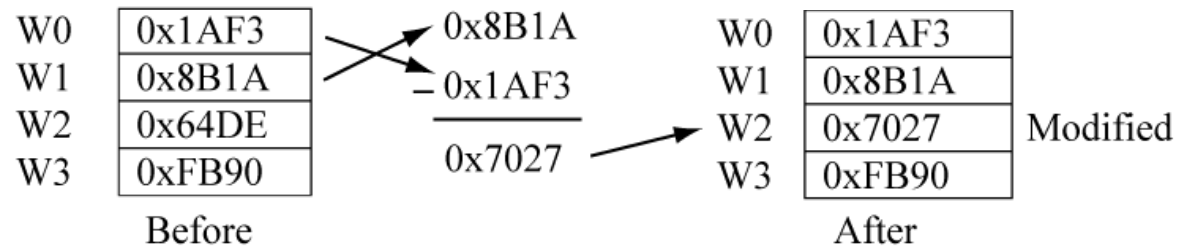
Lower 8 bits of W0, W1 are subtracted and placed in the

SUB{.B} *Wb, Ws, Wd* Instruction Execution

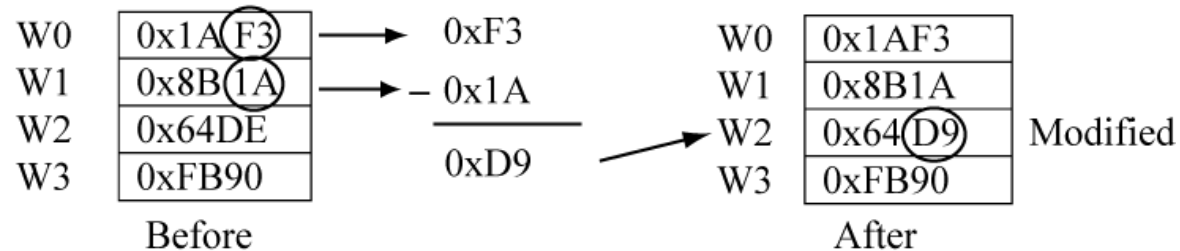
(a) Execute: `sub W0,W1,W2`



(b) Execute: `sub W1,W0,W2`



(c) Execute: `sub.b W0,W1,W2`



Subtraction/Addition with Literals

- **Three operand addition/subtraction with literals:**

ADD{.B} *Wb*, #*lit5*, *Wd*

$(Wb) + \#lit5 \rightarrow Wd$

SUB{.B} *Wb*, #*lit5*, *Wd*

$(Wb) - \#lit5 \rightarrow Wd$

#lit5 is a 5-bit unsigned literal; the range 0-31. Provides a convenient method of adding/subtracting a small constant using a single instruction

Examples

ADD W0,#4,W2

$(W0) + 4 \rightarrow W2$

SUB.B W1,#8, W3

$(W1) - 8 \rightarrow W3$

ADD W0, #60, W1

illegal, 60 is greater than 31!

ADD {.B} *f* {,WREG} Instruction

Two operand addition form:

<i>ADD</i> {.B} <i>f</i>	$(f) + (WREG) \rightarrow f$
<i>ADD</i> {.B} <i>f</i> , WREG	$(f) + (WREG) \rightarrow WREG$

WREG is W0, *f* is limited to first 8192 bytes of memory.

One of the operands, either *f* or WREG is always destroyed!

<i>ADD</i> 0x1000	$(0x1000) + (WREG) \rightarrow 0x1000$
<i>ADD</i> 0x1000, WREG	$(0x1000) + (WREG) \rightarrow WREG$
<i>ADD.B</i> 0x1001, WREG	$(0x1001) + (WREG.lsb) \rightarrow WREG.lsb$

Assembly Language Efficiency

- **The effects of the following instruction:**

ADD 0x1000 $(0x1000) + (WREG) \rightarrow 0x1000$

Can also be accomplished by:

MOV 0x1000 , W1	$(0x1000) \rightarrow W1$
ADD W0, W1, W1	$(W0) + (W1) \rightarrow W1$
MOV W1, 0x1000	$(W1) \rightarrow 0x1000$

This takes three instructions and an extra register. However, in this class we are only concerned with the **correctness** of your assembly language, and not the **efficiency**. Use whatever approach you best understand!!!!

ADD {.B}f {,WREG} Instruction Execution

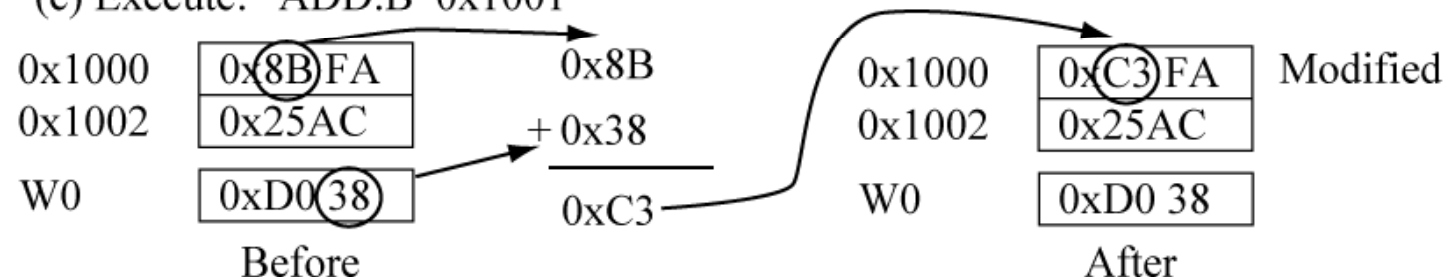
(a) Execute: ADD 0x1000



(b) Execute: ADD 0x1000, WREG



(c) Execute: ADD.B 0x1001



SUB{.B}f{,WREG} Instruction

- **Two operand subtraction form:**

$SUB\{.B\}f$	$(f) - (WREG) \rightarrow f$
$SUB\{.B\}f, WREG$	$(f) - (WREG) \rightarrow WREG$

WREG is W0, f is limited to first 8192 bytes of memory.

One of the operands, either f or WREG is always destroyed!

$SUB\ 0x1000$	$(0x1000) - (WREG) \rightarrow 0x1000$
$SUB\ 0x1000, WREG$	$(0x1000) - (WREG) \rightarrow WREG$
$SUB.B\ 0x1001, WREG$	$(0x1001) - (WREG.lsb) \rightarrow WREG.lsb$

Increment

Increment operation, register-to-register form:

$\text{INC}\{\text{.B}\} Ws, Wd \quad (Ws) + 1 \rightarrow Wd$

Increment operation, memory to memory/WREG form:

$\text{INC}\{\text{.B}\} f \quad (f) + 1 \rightarrow f$

$\text{INC}\{\text{.B}\} f, \text{WREG} \quad (f) + 1 \rightarrow \text{WREG}$

(f must be in first 8192 locations of data memory)

Examples:

$\text{INC } W2, W4 \quad (W2) + 1 \rightarrow W4$

$\text{INC.B } W3, W3 \quad (W3.\text{lsb}) + 1 \rightarrow W3.\text{lsb}$

$\text{INC } 0x1000 \quad (0x1000) + 1 \rightarrow 0x1000$

$\text{INC.B } 0x1001, \text{WREG} \quad (0x1001) + 1 \rightarrow \text{WREG.lsb}$

Decrement

Decrement operation, register-to-register form:

$\text{DEC}\{.B\} Ws, Wd \quad (Ws) - 1 \rightarrow Wd$

Increment operation, memory to memory/WREG form:

$\text{DEC}\{.B\} f \quad (f) - 1 \rightarrow f$

$\text{DEC}\{.B\} f, \text{WREG} \quad (f) - 1 \rightarrow \text{WREG}$

(f must be in first 8192 locations of data memory)

Examples:

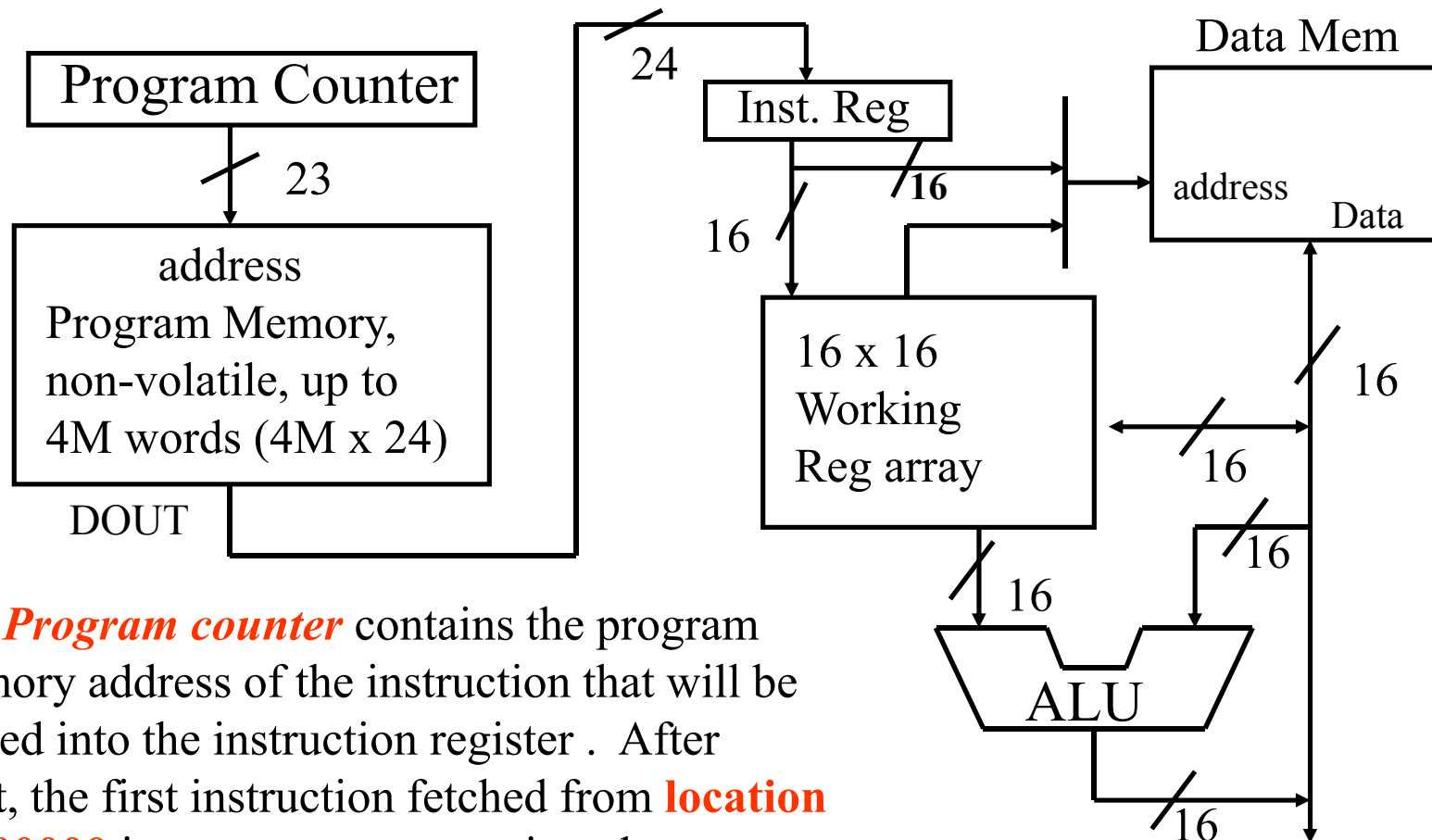
$\text{DEC } W2, W4$

$\text{DEC.B } W3, W3$

$\text{DEC } 0x1000$

$\text{DEC.B } 0x1001, \text{WREG}$

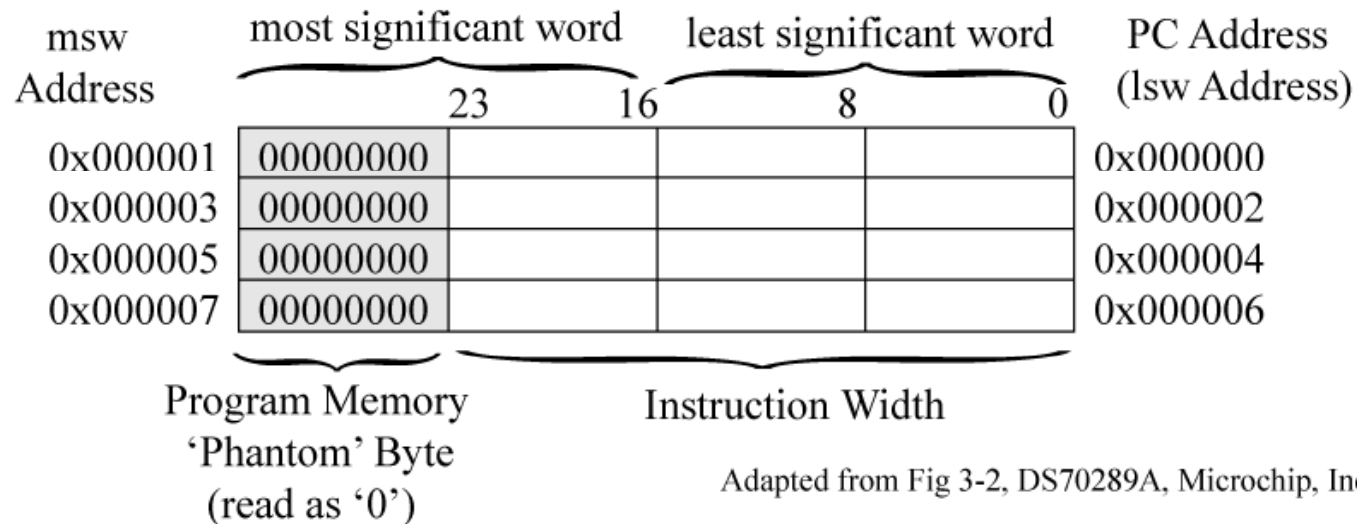
How is the instruction loaded?



The **Program counter** contains the program memory address of the instruction that will be loaded into the instruction register. After reset, the first instruction fetched from **location 0x000000** in program memory, i.e., the program counter is reset to **0x000000**.

17 x 17 Multiplier
not shown

Program Memory Organization



PC is 23-bits wide, but instructions start on even word boundaries, so the PC can address 4M instructions ($M = 2^{20}$).

An instruction is 24 bits (3 bytes). Program memory should be viewed as words (16-bit addressable), with the upper byte of the upper word of an instruction always reading as '0'. Instructions must start on even-word boundaries. Instructions are addressed by the Program counter (PC).

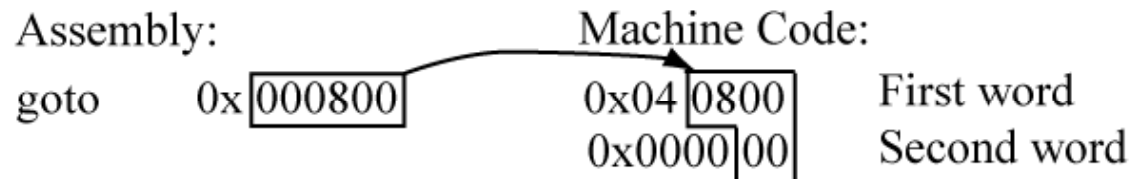
Goto location (goto)

How can the program counter be changed?

`goto Expr lit23 → PC`

Expr is a label or expression that is resolved by the linker to a 23-bit program memory address known as the *target address* (this must be an even address).

The GOTO instruction requires two instruction words:



```
BBBB BBBB BBBB BBBB BBBB BBBB
2222 1111 1111 1100 0000 0000
3210 9876 5432 1098 7654 3210
```

```
0000 0100 nnnn nnnn nnnn nnn0
```

```
0000 0000 0000 0000 0nnn nnnn
```

nn..nn0 = 23-bit value that is loaded into the PC

A GOTO instruction is an unconditional jump.

Valid addressing modes

- What are valid addressing modes for instructions**

Complete information can be found in table 19-2 of the PIC24H32GP202 datasheet.

TABLE 19-2: INSTRUCTION SET OVERVIEW (CONTINUED)

Base Instr #	Assembly Mnemonic	Assembly Syntax	Description	# of Words	# of Cycles	Status Flags Affected
40	MOV	MOV <i>f</i> , <i>Wn</i>	Move <i>f</i> to <i>Wn</i>	1	1	None
		MOV <i>f</i>	Move <i>f</i> to <i>f</i>	1	1	N,Z
		MOV <i>f</i> , <i>WREG</i>	Move <i>f</i> to <i>WREG</i>	1	1	N,Z
		MOV #lit16, <i>Wn</i>	Move 16-bit literal to <i>Wn</i>	1	1	None
		MOV.b #lit8, <i>Wn</i>	Move 8-bit literal to <i>Wn</i>	1	1	None
		MOV <i>Wn</i> , <i>f</i>	Move <i>Wn</i> to <i>f</i>	1	1	None
		MOV <i>Wso</i> , <i>Wdo</i>	Move <i>Ws</i> to <i>Wd</i>	1	1	None
		MOV <i>WREG</i> , <i>f</i>	Move <i>WREG</i> to <i>f</i>	1	1	N,Z
		MOV.D <i>Wns</i> , <i>Wd</i>	Move Double from <i>W(ns):W(ns + 1)</i> to <i>Wd</i>	1	2	None
		MOV.D <i>Ws</i> , <i>Wnd</i>	Move Double from <i>Ws</i> to <i>W(nd + 1):W(nd)</i>	1	2	None

Wso, Wsd, Wn

- MOV Wso, Wdo**

Symbols used in opcode descriptions

Field	Description
Wnd	One of 16 destination working registers $\in \{W0..W15\}$
Wns	One of 16 source working registers $\in \{W0..W15\}$
WREG	W0 (working register used in file register instructions)
Ws	Source W register $\in \{Ws, [Ws], [Ws++] , [Ws--], [++Ws], [--Ws] \}$
Wso	Source W register $\in \{Wns, [Wns], [Wns++] , [Wns--], [++Wns], [--Wns], [Wns+Wb] \}$
Wd	Destination W register $\in \{Wd, [Wd], [Wd++] , [Wd--], [++Wd], [--Wd] \}$
Wdo	Destination W register $\in \{Wnd, [Wnd], [Wnd++] , [Wnd--], [++Wnd], [--Wnd],$
Wn	One of 16 working registers $\in \{W0..W15\}$
Wb	Base W register $\in \{W0..W15\}$

ADD forms

- **ADD Wb, Ws, Wd**

Field	Description
Ws	Source W register $\in \{ Ws, [Ws], [Ws++] , [Ws--], [++Ws], [--Ws] \}$
Wd	Destination W register $\in \{ Wd, [Wd], [Wd++] , [Wd--], [++Wd], [--Wd] \}$
Wb	Base W register $\in \{ W0..W15 \}$

- **Legal:**

ADD W0, W1, W2

ADD W0, [W1], [W4]

- **Illegal**

ADD [W0],W1,W2 ;first operand illegal!


Simple Program (example)

- **Sample programs written in C, translated (compiled) to PIC 24uC assembly language**

C Program equivalent

```
#define avalue 100  
uint8 i,j,k;
```

A uint8 variable is
8 bits (1 byte)



```
i = avalue;    // i = 100  
i = i + 1;     // i++, i = 101  
j = i;         // j is 101  
j = j - 1;     // j--, j is 100  
k = j + i;     // k = 201
```

Where are variables stored?

When writing assembly language, can use any free data memory location to store values, it your choice.

A logical place to begin storing data in the first free location in data memory, which is 0x0800 (Recall that 0x0000-0x07FF is reserved for SFRs).

Assign i to 0x0800, j to 0x0801, and k to 0x0802. Other choices could be made.

C to PIC24 Assembly

$i = 100;$	<div><pre>mov.b #100,W0 ;WREG = 100 = 0x64 mov.b WREG,0x0800 ;i = WREG</pre></div>
$i = i + 1;$	<div><pre>inc.b 0x0800 ;i = i + 1</pre></div>
$j = i;$	<div><pre>mov.b 0x0800,WREG ;WREG = i mov.b WREG,0x0801 ;j = WREG</pre></div>
$j = j - 1;$	<div><pre>dec.b 0x0801 ;j = j - 1</pre></div>
$k = j + i;$	<div><pre>mov.b 0x0800,WREG ;WREG = i add.b 0x0801,WREG ;WREG = j + WREG mov.b WREG,0x0802 ;k = WREG</pre></div>

i is location 0x0800, j is location 0x0801, k is location 0x0802

Comments: The assembly language program operation is not very clear. Also, multiple assembly language statements are needed for one C language statement. Assembly language is more *primitive* (operations less powerful) than C.

PIC24 Assembly to PIC24 Machine Code

- **Could perform this step manually by determining the instruction format for each instruction from the data sheet.**
- **Much easier to let a program called an *assembler* do this step automatically**
- **The MPLAB Integrated Design Environment (IDE) is used to assemble PIC24 programs and simulate them**
 - Simulate means to execute the program without actually loading it into a PIC24 microcontroller

```

.include "p24Hxxxx.inc"
.global __reset
.bss    ;reserve space for variables
i:      .space 1
j:      .space 1
k:      .space 1
.text                    ;Start of Code section
__reset: ; first instruction located at __reset label
    mov #__SP_init, W15    ;;initialize stack pointer
    mov #__SPLIM_init,W0
    mov W0,SPLIM           ;;initialize Stack limit reg.
    avalue = 100
; i = 100;
    mov.b #avalue, W0      ; W0 = 100
    mov.b WREG,i           ; i = 100
; i = i + 1;
    inc.b i                ; i = i + 1
; j = i
    mov.b i,WREG           ; W0 = i
    mov.b WREG,j           ; j = W0
; j = j - 1;
    dec.b j                ; j= j - 1
; k = j + i
    mov.b i,WREG           ; W0 = i
    add.b j,WREG           ; W0 = W0+j (WREG is W0)
    mov.b WREG,k           ; k = W0
done:
    goto done             ;loop forever

```

mptst_byte.s

This file can be assembled
by the MPLAB™
assembler into PIC24
machine code and
simulated.

Labels used for memory
locations 0x0800 (i),
0x0801(j), 0x0802(k) to
increase code clarity

mptst_byte.s (cont.)

```
.include "p24Hxxxx.inc"
```

Include file that defines various labels for a particular processor. '*.include*' is an assembler directive.

```
.global __reset
```

Declare the `__reset` label as global – it is needed by linker for defining program start

```
.bss      ;reserve space for variables  
i:        .space 1  
j:        .space 1  
k:        .space 1
```

The *.bss* assembler directive indicates the following should be placed in data memory. By default, variables are placed beginning at the first free location, 0x800. The *.space* assembler directive reserves space in bytes for the named variables. `i`, `j`, `k` are labels, and labels are case-sensitive and must be followed by a ':' (colon).

An *assembler directive* is not a PIC24 instruction, but an instruction to the assembler program. Assembler directives have a leading '.' period, and are not case sensitive.

mptst_byte.s (cont.)

```
.text ←  
__reset: mov #__SP_init, W15  
        mov #__SPLIM_init, W0  
        mov W0, SPLIM
```

‘.text’ is an assembler directive that says what follows is code. Our first instruction must be labeled as ‘__reset’.

These move instructions initialize the stack pointer and stack limit registers – this will be discussed in a later chapter.

```
avalue = 100
```

The equal sign is an assembler directive that equates a label to a value.

mptst_byte.s (cont.)

```

; i = 100;
mov.b #avalue, W0    ; W0 = 100
mov.b WREG,i         ; i = 100

; i = i + 1;
inc.b i              ; i = i + 1
; j = i

mov.b i,WREG         ; W0 = i
mov.b WREG,j         ; j = W0
; j = j - 1;
dec.b j              ; j = j - 1
; k = j + i
mov.b i,WREG         ; W0 = i
add.b j,WREG         ; W0 = W0+j
mov.b WREG,k         ; k = W0

```

The use of labels and comments greatly improves the clarity of the program.

It is hard to over-comment an assembly language program if you want to be able to understand it later.

Strive for at least a comment every other line; refer to lines

(WREG is W0)

mptst_byte.s (cont.)

```
done: ←  
goto  done ;loop forever
```

A label that is the target of a *goto* instruction. Labels are **case sensitive** (instruction mnemonics and assembler directives are not case sensitive).

```
.end
```

A comment

An assembler directive specifying the end of the program in this file.

An Alternate Solution

C Program equivalent

```
#define avalue 100
uint8 i,j,k;

i = avalue;    // i = 100
i = i + 1;     // i++, i = 101
j = i;         // j is 101
j = j - 1;     // j--, j is 100
k = j + i;     // k = 201
```

Previous approach took 9 instructions, this one took 11 instructions. Use whatever approach that you best understand.

```
;Assign variables to registers
;Move variables into registers.
;use register-to-register operations for
computations;
```

```
;write variables back to memory
;assign i to W1, j to W2, k to W3
```

```
mov #100,W1      ; W1 (i) = 100
inc.b W1,W1      ; W1 (i) = W1 (i) + 1
mov.b W1,W2      ; W2 (j) = W1 (i)
dec.b W2,W2      ; W2 (j) = W2 (j) -1
add.b W1,W2,W3   ; W3 (k) = W1 (i) + W2 (j)
;;write variables to memory
mov.b W1,W0      ; W0 = i
mov.b WREG,i     ; 0x800 (i) = W0
mov.b W2,W0      ; W0 = j
mov.b WREG,j     ; 0x801 (j) = W0
mov.b W3,W0      ; W3 = k
mov.b WREG,k     ; 0x802 (k) = W0
```


Clock Cycles vs. Instruction Cycles

The clock signal used by a PIC24 μ C to control instruction execution can be generated by an off-chip oscillator or crystal/capacitor network, or by using the internal RC oscillator within the PIC24 μ C.

For the PIC24H family, the maximum clock frequency is 80 MHz.

An **instruction cycle (FCY)** is **two clock (FOSC)** cycles.

A PIC24 instruction takes 1 or 2 **instruction (FCY)** cycles, depending on the instruction (see Table 19-2, PIC24HJ32GP202 data sheet). If an instruction causes the program counter to change (i.e, GOTO), that instruction takes 2 instruction cycles.

An add instruction takes 1 instruction cycle.

How much time is this if the clock frequency (FOSC) is 80 MHz (1 MHz = 1.0×10^6 = 1,000,000 Hz)?

$1 / \text{frequency} = \text{period}$

$1 / 80 \text{ MHz} = 12.5 \text{ ns}$ (1 ns = 1.0×10^{-9} s)

1 Add instruction @ 80 MHz takes 2 clocks * 12.5 ns = 25 ns (or 0.025 μ s).

By comparison, an Intel Pentium add instruction @ 3 GHz takes 0.33 ns (330 ps). An Intel Pentium could emulate a PIC24HJ32GP202 faster than a PIC24HJ32GP202 can execute! But you can't put a Pentium in a toaster, or buy one from Digi-key for \$5.00.

How long does mpstst_bytes.s take to execute

- Beginning at the `__reset` label, and ignoring the *goto* at the end, takes 12 instruction cycles, which is 24 clock cycles.

	Instruction Cycles
<code>mov #__SP_init, W15</code>	1
<code>mov #__SPLIM_init, W0</code>	1
<code>mov W0, SPLIM</code>	1
<code>mov.b #avalue, W0</code>	1
<code>mov.b WREG, i</code>	1
<code>inc.b i</code>	1
<code>mov.b i, WREG</code>	1
<code>mov.b WREG, j</code>	1
<code>dec.b j</code>	1
<code>mov.b i, WREG</code>	1
<code>add.b j, WREG</code>	1
<code>mov.b WREG, k</code>	1
V 0.2 Total	12


What if we used 16-bit variable instead of 8-bit variables?

C Program equivalent

```
#define avalue 2047
```

```
uint16 i,j,k;
```

A uint16 variable is
16 bits (1 byte)



```
i = avalue;    // i = 2047
```

```
i = i + 1;     // i++, i = 2048
```

```
j = i;         // j is 2048
```

```
j = j - 1;     // j--, j is 2047
```

```
k = j + i;     // k = 4095
```

```

.include "p24Hxxxx.inc"
.global __reset
.bss ;reserve space for variables
i:   .space 2
j:   .space 2
k:   .space 2

.text ;Start of Code section
__reset: ; first instruction located at __reset label
    mov #__SP_init, w15 ;initialize stack pointer
    mov #__SPLIM_init, W0
    mov W0, SPLIM ;initialize stack limit reg
    avalue = 2048

; i = 2048;
    mov #avalue, W0 ; W0 = 2048
    mov WREG, i ; i = 2048

; i = i + 1;
    inc i ; i = i + 1

; j = i
    mov i, WREG ; W0 = i
    mov WREG, j ; j = W0

; j = j - 1;
    dec j ; j= j - 1

; k = j + i
    mov i, WREG ; W0 = i
    add j, WREG ; W0 = W0+j (WREG is W0)
    mov WREG, k ; k = W0

done:
    goto done ;loop forever

```

Reserve 2 bytes for each variable. Variables are now stored at 0x0800, 0x0802, 0x0804

Instructions now perform WORD (16-bit) operations (the .b qualifier is removed).

An alternate Solution (16-bit variables)

C Program equivalent

```
#define avalue 2047
uint16 i,j,k;

i = avalue;    // i = 2047
i = i + 1;     // i++, i = 2048
j = i;         // j is 2048
j = j - 1;     // j--, j is 2047
k = j + i;     // k = 4095
```

Previous approach took 9 instructions, this one took 8 instructions. In this case, this approach is more efficient!

```
;Assign variables to registers
;Move variables into registers.
;use register-to-register operations for
computations;
```

```
;write variables back to memory
;assign i to W1, j to W2, k to W3
```

```
mov #2047,W1    ; W1 (i) = 2047
inc W1,W1       ; W1 (i) = W1 (i) + 1
mov W1,W2       ; W2 (j) = W1 (i)
dec W2,W2       ; W2 (j) = W2 (j) -1
add W1,W2,W3    ; W3 (k) = W1 (i) + W2 (j)
;;write variables to memory
mov W1,i        ; 0x800 (i) = W1
mov W2,j        ; 0x802 (j) = W2
mov W3,k        ; 0x804 (k) = W3
```

How long does `mptst_word.s` take to execute?

Ignoring the *goto* at the end, takes 12 instruction cycles, which is 24 clock cycles.

	Instruction Cycles
<code>mov #__SP_init, W15</code>	1
<code>mov #__SPLIM_init, W0</code>	1
<code>mov W0, SPLIM</code>	1
<code>mov #avalue, W0</code>	1
<code>mov WREG, i</code>	1
<code>inc i</code>	1
<code>mov i, WREG</code>	1
<code>mov WREG, j</code>	1
<code>dec j</code>	1
<code>mov i, WREG</code>	1
<code>add j, WREG</code>	1
<code>mov WREG, k</code>	1
Total	12

16 bit operations vs. 8 bit operations

The 16-bit version of the *mptst* program requires the same number of instruction bytes and the same number of instruction cycles as the 8-bit version.

This is because the PIC24 family is a 16-bit microcontroller; its natural operation size is 16 bits, so 16-bit operations are handled as efficiently as 8-bits operations

On an 8-bit processor, like the PIC18 family, the 16-bit version would take roughly double the number of instructions and clock cycles as the 8-bit version.

On the PIC24, a 32-bit version of the *mptst* program will take approximately twice the number of instructions and clock cycles as the 16-bit version. We will look at 32-bit operations later in the semester.

Review: Units

In this class, units are always used for physical quantity:

Time	Frequency
milliseconds (ms = 10^{-3} s)	kilohertz (kHz = 10^3 Hz)
microseconds (μ s = 10^{-6} s)	megahertz (MHz = 10^6 Hz)
nanoseconds (ns = 10^{-9} s)	gigahertz (GHz = 10^9 Hz)

For a frequency of 1.25 kHz, what is the period in μ s?

period = $1/f = 1/(1.25 \text{ e}3) = 8.0 \text{ e} -4$ seconds

Unit conversion= $8.0\text{e-}4 \text{ (s)} * (1\text{e}6 \text{ }\mu\text{s})/1.0 \text{ (s)} = 8.0\text{e}2 \text{ }\mu\text{s} = 800 \text{ }\mu\text{s}$

PIC24H Family

- **Microchip has an extensive line of PICmicro[®] microcontrollers, with the PIC24 family introduced in 2005.**
- **The PIC16 and PIC18 are older versions of the PICmicro[®] family, have been several previous generations.**
- **Do not assume that because something is done one way in the PIC24, that it is the most efficient method for accomplishing that action.**
- **The datasheet for the PIC24 is found on the class UVLE site.**

PICmicro Survey

	PIC16F87x	PIC18F242	PIC24H
Instruction width	14 bits	16 bits	24 bits
Program Memory	8K instruction	8K instructions	~10K instructions
Data Memory	386 bytes	1537 bytes	2048 bytes
Clock Speed	Max 20 MHz 4 clks = 1 instruction	Max 40 MHz 4 clks = 1 instruction	Max 80 MHz 2 clks = 1 instruction

The PIC24H can execute about 6x faster than the PIC18F242

Summary

- **Understand the PIC24 basic architecture (program and data memory organization)**
- **Understand the operation of *mov, add, sub, inc, dec, goto* instructions and their various addressing mode forms**
- **Be able to convert simple C instruction sequences to PIC24 assembly language**
 - Be able to assemble/simulate a PIC24 μ C assembly language program in the MPLAB IDE
- **Understand the relationship between instruction cycles and machine cycles**