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https://algo.monster/problems/amazon_online_assessment_questions

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两道题蒟蒻都可以找到，蒟蒻1481 1120

1. remove the distinct product id

$O(n)$ to calculate the frequent. PQ $O(k \log k)$ kic the

1481. Least Number of Unique Integers after K Removals number of unique element.

```
public int findLeastNumOfUniqueInts(int[] arr, int k) {
    if(k == arr.length){
        return 0;
    }
    Map<Integer, Integer> map = new HashMap<>();
    Integer[] nums = new Integer[arr.length];
    int index = 0;
    for(int n : arr){
        nums[index++] = n;
        map.put(n, map.getOrDefault(n, 0) + 1);
    }
    if(k == 0){
        return map.size();
    }
    Arrays.sort(nums, new Comparator<Integer>(){
        public int compare(Integer a, Integer b){
            int n1 = map.get(a), n2 = map.get(b);
            return (n1 != n2) ? (n1 - n2) : (a - b);
        }
    });
    Set<Integer> hs = new HashSet<>();
    for(int i = k; i < nums.length; i++){
        hs.add(nums[i]);
    }
    return hs.size();
}
```

2. find the root component

1120. Maximum Average Subtree

Time $O(n)$ space $O(n)$ n is the number of node

```
public class Solution {

    private double res;

    public double maximumAverageSubtree(TreeNode root) {
        dfs(root);
        return res;
    }

    private int[] dfs(TreeNode cur) {
        // 如果是空树, 则节点数和总和都是0
        if (cur == null) {
            return new int[]{0, 0};
        }

        int[] left = dfs(cur.left), right = dfs(cur.right);
        int count = left[0] + right[0] + 1, sum = left[1] + right[1] + cur.val;

        // 枚举cur为树根的情形, 更新答案
        res = Math.max(res, (double) sum / count);
        return new int[]{count, sum};
    }
}

class TreeNode {
    int val;
    TreeNode left, right;

    public TreeNode(int val) {
        this.val = val;
    }
}
```

1041. Robot Bounded In Circle

```
public boolean isRobotBounded(String instructions) {
    int x = 0;
    int y = 0;
    //Keep track of directions
    String dir = "North";
    //    N
    // W<----->E
    //    S
```

```

//Caluculate position
for(char c : instructions.toCharArray()){
    if(c == 'G'){
        if(dir.equals("North")){y++;}
        else if(dir.equals("East")){x++;}
        else if(dir.equals("South")){y--;}
        else {x--;}
    }
    else if(c == 'L'){
        if(dir.equals("North")){dir = "West";}
        else if(dir.equals("East")){dir = "North";}
        else if(dir.equals("South")){dir = "East";}
        else {dir = "South";}
    }
    else{
        if(dir.equals("North")){dir = "East";}
        else if(dir.equals("East")){dir = "South";}
        else if(dir.equals("South")){dir = "West";}
        else {dir = "North";}
    }
}
//Check if calutucated position is starting position
if(x == 0 && y == 0){
    return true;
}
//check if the final faced direction is not North(Strarting
DIRECTION)
if(dir.equals("North")){
    return false;
}
return true;
}

```

Optimize Box Weight

Given a list of integers, partition it into two subsets S1 and S2 such that the sum of S1 is greater than that of S2. And also the number of elements in S1 should be minimal.

Return S1 in increasing order.

Notice if more than one subset A exists, return the one with the maximal sum.

Examples:

Input:

```
nums = [4, 5, 2, 3, 1, 2]
```

Output:

```
[4, 5]
```

Explanation:

We can divide the numbers into two subsets A = [4, 5] and B = [1, 2, 2, 3]. The sum of A is 9 which is greater than the sum of B which is 8. There are other ways to divide but A = [4, 5] is of minimal size of 2.

Input:

```
nums = [10, 5, 3, 1, 20]
```

Output:

```
[20]
```

Input:

```
nums = [1, 2, 3, 5, 8]
```

Output:

```
[5, 8]
```

// "static void main" must be defined in a public class.

```
public class OptimizeBox {  
    public static void main(String[] args) {  
        int[] ans1 = optimize(new int[]{1, 1, 2, 1});  
        print("ans1", ans1);  
        int[] ans2 = optimize(new int[]{3, 7, 6, 2});  
        print("ans2", ans2);  
    }  
    private static int[] optimize(int[] arr) {  
        long arraySum = 0;  
        for (int num : arr) {  
            arraySum += num;  
        }  
        Arrays.sort(arr);  
        long max = 0;  
        int idx = 0;  
  
        while (idx < arr.length && max * 2 < arraySum) {
```

```

        max += arr[idx];
        idx++;
    }

    idx--; // idx now is the first element in box A

    int[] res = new int[arr.length - idx];
    for (int i = 0; i < arr.length - idx; i++) {
        res[i] = arr[idx + i];
    }
    return res;
}

```

storage optimization

Create a function that takes in a matrix, and an array of rows and an array of columns to be removed from the matrix, and return the size of the biggest cubic space after removing the shelves and columns.

biggestStorage(matrix, rows, columns)

```
return biggestCubicSize; // i.e. 3x4 = 12
```

$S: O(1)$

$T: O(N \cdot \log(N) + M \cdot \log(M))$

$N \neq h \dots$

$M \neq \text{vertical cuts}$

solution leetcode: 切蛋糕 cut cake

```

public int maxArea(int h, int w, int[] horizontalCuts, int[] verticalCuts) {
    Arrays.sort(horizontalCuts);
    Arrays.sort(verticalCuts);
    int maxH = horizontalCuts[0];
    int maxW = verticalCuts[0];
    for (int i = 1; i < horizontalCuts.length; i++) {
        maxH = Math.max(maxH, horizontalCuts[i] - horizontalCuts[i - 1]);
    }
    for (int i = 1; i < verticalCuts.length; i++) {
        maxW = Math.max(maxW, verticalCuts[i] - verticalCuts[i - 1]);
    }
    maxH = Math.max(maxH, h - horizontalCuts[horizontalCuts.length - 1]);
    maxW = Math.max(maxW, w - verticalCuts[verticalCuts.length - 1]);
    return (int)((long)maxH * maxW % (1000000007));
}

```

CloudFront Caching

A company owns N warehouses, identified as `warehouse[0 to N-1]`. The owner would like to measure the maintenance cost. A warehouse will be able to share the stock with other connected warehouses, making it less costly to restock.

The method to evaluate the maintenance cost is

- I. If a warehouse is not connected to any others, its maintenance cost is 1.
- II. If multiple warehouses are connected, the total maintenance cost of the group of connected warehouses will be the ceiling of the square root of K , where K is the number of warehouses in the group.
- III. A warehouse connected to any of the warehouses in a group will be able to share stock with all in the group. For example, if warehouse 0 and warehouse 1 are connected, and warehouse 1 and warehouse 2 are connected, consider 0, 1 and 2 in the same group.
- IV. The total maintenance cost is the sum of all costs.

Given the number of warehouses N and a 2d array, where every subarray is a pair of connected warehouses, build a function that return the total maintenance cost.

Function

`int costEvaluation(int n, int[][] connections)`

Examples

Input

$n = 4$

`connections = [[0, 2], [1, 2]]`

Output

3

Explanation

Warehouses 0, 1, 2 are in the same group, the cost is $\text{ceiling}(\text{sqrt}(3)) = 2$.

Warehouse 3 costs 1.

Total cost is $2 + 1 = 3$.

Input

$n = 10$

`connections = [[2, 6], [3, 5], [0, 1], [2, 9], [5, 6]]`

Output

8

Explanation

Warehouses 0, 1 are in the same group, the cost is $\text{ceiling}(\text{sqrt}(2)) = 2$.

Warehouses 2, 3, 5, 6, 9 are in the same group, the cost is $\text{ceiling}(\text{sqrt}(5)) = 3$.

Warehouse 4 costs 1.

Warehouse 7 costs 1.

solution:

```
public static int cost(int n, int[][] arr) {
    if (n < 2) return n;
    // build graph
    HashMap<Integer, List<Integer>> graph = new HashMap<>();
    for (int i = 0; i < n; i++) {
        graph.put(i, new ArrayList<>());
    }
    for (int[] a : arr) {
        int x = a[0];
        int y = a[1];
        graph.get(x).add(y);
        graph.get(y).add(x);
    }

    int island = 0;
    int res = 0;
    HashSet<Integer> visited = new HashSet<>();
    for (Map.Entry<Integer, List<Integer>> e : graph.entrySet()) {
        if (e.getValue().size() == 0) {
            island++;
        } else {
            int curK = e.getKey();
            if (!visited.contains(curK)) {
                HashSet<Integer> curGroup = new HashSet<>();
                Queue<Integer> q = new LinkedList<>();
                q.add(curK);
                curGroup.add(curK);
                while (!q.isEmpty()) {
                    int cur = q.poll();
                    List<Integer> curV = graph.get(cur);
                    for (int i = 0; i < curV.size(); i++) {
                        int iCurV = curV.get(i);
                        if (!curGroup.contains(iCurV) && !visited.contains(iCurV)) {
                            curGroup.add(iCurV);
                            visited.add(iCurV);
                            q.add(iCurV);
                        }
                    }
                }
                res += (int) Math.ceil(Math.sqrt(curGroup.size()));
            }
        }
    }
    return res + island;
}
```

union find solution:

```

public static int test(int n, int[][] arr) {
    UnionFind uf = new UnionFind(n);
    for (int i = 0; i < arr.length; i++) {
        uf.union(arr[i][0], arr[i][1]);
    }
    int res = 0;
    for (int i = 0; i < n; i++) {
        if (uf.parent[i] == i) {
            res += Math.ceil(Math.sqrt(uf.rank[i]));
        }
    }
    return res;
}

static class UnionFind {
    int[] parent;
    int[] rank;
    int size;
    public UnionFind(int n) {
        parent = new int[n];
        rank = new int[n];
        size = n;
        for (int i = 0; i < n; i++) {
            parent[i] = i;
            rank[i] = 1;
        }
    }

    int find(int x) {
        int root = x;
        if (parent[x] != x) {
            root = find(parent[x]);
        }
        return root;
    }

    void union(int a, int b) {
        int rootA = find(a);
        int rootB = find(b);
        if (rootA == rootB) return;
        if (rank[rootA] > rank[rootB]) {
            parent[rootB] = rootA;
            rank[rootA] += rank[rootB];
        } else {
            parent[rootA] = rootB;
            rank[rootB] += rank[rootA];
        }
    }
}

```

```

1 import java.util.*;
2 import java.io.*;
3 import java.lang.*;
4
5 public class Solution {
6     public static int costEvaluation(int n, int[][] connections) {
7         if(connections == null || connections.length == 0) return 0;
8         int res = 0;
9         boolean[] visited = new boolean[n];
10        Map<Integer, List<Integer>> graph = new HashMap<>();
11        for (int i = 0 ; i < n ; i++) {
12            graph.put(i, new ArrayList<>());
13        }
14        for (int[] c : connections) {
15            int from = c[0];
16            int to = c[1];
17            graph.get(from).add(to);
18            graph.get(to).add(from);
19        }
20        for(int i = 0 ; i < n ; i++) {
21            if(!visited[i]) {
22                int num = dfs(i, graph, visited);
23                res += Math.ceil(Math.sqrt(num));
24            }
25        }
26        return res;
27    }
28
29    private static int dfs(int cur, Map<Integer, List<Integer>> graph, boolean[] visited) {
30        if(visited[cur]) return 0;
31        int res = 1;
32        visited[cur] = true;
33        for(int nei : graph.get(cur)) {
34            res += dfs(nei, graph, visited);
35        }
36        return res;
37    }
38 }
39

```

$O(n)$ build graph.

$O(n)$ visited all nodes

排序, $O(n \log n)$

第一题 利口舅其散 套了个算餐馆距离壳子 restaurant

973. K Closest Points to Origin

```
public int[][] kClosest(int[][] points, int k) {
    PriorityQueue<Distance> pq = new PriorityQueue<>((a, b)->{
        if (a.distance < b.distance) {
            return -1;
        } else if (a.distance > b.distance) {
            return 1;
        } else {
            return 0;
        }
    });

    for (int[] p: points) {
        double dis = (double)(p[0] * p[0] + p[1] * p[1]);
        pq.add(new Distance(p, dis));
    }

    int[][] res = new int[k][2];
    while (k > 0) {
        res[k-1] = pq.poll().p;
        k--;
    }
    return res;
}

private class Distance {
    int[] p;
    double distance;
    Distance(int[] _p, double _distance) {
        p = _p;
        distance = _distance;
    }
}
```

第二题 Prime Air Route 换了个壳子, sort后 two pointer, 处理好重复的地方。

prime air route

Examples

Example 1:

Input:

maxTravelDist = 7000

forwardRouteList = [[1,2000],[2,4000],[3,6000]]

returnRouteList = [[1,2000]]

Output:

[[2,1]]

Explanation:

There are only three combinations, [1,1], [2,1], and [3,1], which have a total of 4000, 6000, and 8000 miles, respectively. Since 6000 is the largest use that does not exceed 7000, [2,1] is the only optimal pair.

Example 2:

Input:

maxTravelDist = 10000

forwardRouteList = [[1, 3000], [2, 5000], [3, 7000], [4, 10000]]

returnRouteList = [[1, 2000], [2, 3000], [3, 4000], [4, 5000]]

Output:

[[2, 4], [3, 2]]

Explanation:

There are two pairs of forward and return shipping routes possible that optimally utilizes the given aircraft.

10

11

12

13

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21

22

23

24

Tes

solution:

```

public static List<int[]> getPair(int[][] a, int[][] b, int target) {
    Arrays.sort(a, Comparator.comparingDouble(o -> o[1]));
    Arrays.sort(b, Comparator.comparingDouble(o -> o[1]));
    List<int[]> res = new ArrayList<>();
    int max = -1;
    int aInd = 0;
    int bInd = b.length - 1;

    while (aInd < a.length && bInd >= 0) {
        //System.out.println("aInd: " + aInd);
        //System.out.println("bInd: " + bInd);
        int aVal = a[aInd][1];
        int bVal = b[bInd][1];
        int sum = aVal + bVal;
        //System.out.println("sum: " + sum);

        if (sum > target) {
            bInd--;
        } else {
            if (sum >= max) {
                if (sum > max) {
                    res.clear();
                    max = sum;
                }
                res.add(new int[]{a[aInd][0], b[bInd][0]});
                int index = bInd - 1;
                while (index >= 0 && b[index][1] == b[index + 1][1]) {
                    res.add(new int[]{a[aInd][0], b[index][0]});
                    index--;
                }
            }
            aInd++;
        }
    }

    return res;
}

```

solution2:

```

public List<int[]> getPair(final int[][] A, final int[][] B, final int target) {
    int sum = -1;
    final List<int[]> result = new ArrayList<>();
    for (int[] a : A) {
        for (int[] b : B) {
            final int s = a[1] + b[1];
            if (s > target || s < sum) {
                continue;
            }
            if (s > sum) {
                result.clear();
                sum = s;
            }
            result.add(new int[]{a[0], b[0]});
        }
    }
    return result;
}

```

solution3:

```

List<int[]> getOptimalRoute(int maxTravelDist, int[][] forwardRoute, int[][] returnRoute) { // 按旅程从大到小排列
    Arrays.sort(forwardRoute, (a,b)->b[1]-a[1]);
    TreeMap<Integer, Stack<Integer>> map=new TreeMap<>();
    for(var rr:returnRoute) {
        // 按旅程分组
        var stack=map.getOrDefault(rr[1],new LinkedList<>());
        stack.push(rr[0]);
        map.put(rr[1],stack);
    }
    var dest=new ArrayList<int[]>();
    for(var fr:forwardRoute) {
        // 贪心, 永远取和最接近最大里程的组
        var re=map.floorEntry(maxTravelDist-fr[1]);
        // 如果没有合适的, 说明最小值和这个值的和都大于里程数了
        if(re==null) continue;
        // 随便取一个
        var stack=re.getValue();
        dest.add(new int[] {fr[0],stack.pop()});
        // 如果空了删除主键
        if (stack.isEmpty()) map.remove(re.getKey());
    }
    return dest;
}

```

shopping options 类似lc454

T: $O(n^2)$ S: $O(n^2)$

Disappeared shopping options question, retrieved from google cache

A customer wants to buy a pair of jeans, a pair of shoes, a skirt, and a top but has a limited budget in dollars. Given different pricing options for each product, determine how many options our customer has to buy 1 of each product. You cannot spend more money than the budgeted amount.

Example

priceOfJeans = [2, 3]

priceOfShoes = [4]

priceOfSkirts = [2, 3]

priceOfTops = [1, 2]

budgeted = 10

The customer must buy shoes for 4 dollars since there is only one option. This leaves 6 dollars to spend on the other 3 items. Combinations of prices paid for jeans, skirts, and tops respectively that add up to 6 dollars or less are [2, 2, 2], [2, 2, 1], [3, 2, 1], [2, 3, 1].

There are 4 ways the customer can purchase all 4 items.

Function Description

Complete the `getNumberOfOptions` function in the editor below. The function must return an integer which represents the number of options present to buy the four items.

getNumberOfOptions has 5 parameters:

int[] priceOfJeans: An integer array, which contains the prices of the pairs of jeans available.

int[] priceOfShoes: An integer array, which contains the prices of the pairs of shoes available.

int[] priceOfSkirts: An integer array, which contains the prices of the skirts available.

int[] priceOfTops: An integer array, which contains the prices of the tops available.

int dollars: the total number of dollars available to shop with.

Constraints

$1 \leq a, b, c, d \leq 10^3$

$1 \leq \text{dollars} \leq 10^9$

$1 \leq \text{price of each item} \leq 10^9$

Note: a, b, c and d are the sizes of the four price arrays

solution:

```

public static int getNumberOfOptions(
    List<Integer> priceOfJeans,
    List<Integer> priceOfShoes,
    List<Integer> priceOfSkirts,
    List<Integer> priceOfTops,
    int dollars
) {
    // combine 2 lists together and sort
    List<Integer> a = combine(priceOfJeans, priceOfShoes, dollars);
    List<Integer> b = combine(priceOfSkirts, priceOfTops, dollars);

    // search
    int left = 0;
    int right = b.size() - 1;
    int counts = 0;

    while (left < a.size() && right >= 0) {
        //System.out.println("a.get(left): " + a.get(left) + " b.get(right): " + b.get(right));
        int sum = a.get(left) + b.get(right);
        if (sum <= dollars) {
            counts += right + 1;
            left++;
        } else {
            right--;
        }
    }
    return counts;
}

private static List<Integer> combine(List<Integer> a, List<Integer> b, int target) {
    List<Integer> sumList = new ArrayList<Integer>();
    for (int i : a) {
        for (int j : b) {
            int sum = i + j;
            if (sum <= target) {
                sumList.add(sum);
            }
        }
    }
    sumList.sort((x, y) -> x - y);
    return sumList;
}

```

Music Runtime

Amazon music is working on a new feature to pair songs together to play while on bus. The goal of this feature is to select two songs from the list that will end exactly 30

seconds before the listener reaches their stop. You are tasked with writing the method that selects the songs from a list. Each song is assigned a unique id, numbered from 0 to N-1.

Assumptions:

1. You will pick exactly 2 songs.
2. you cannot pick the same song twice.
3. if you have multiple pairs with the same total time, select the pair with the longest song.
4. there are atleast 2 songs available.

write an algorithm to return ID's of the 2 songs whose combined runtime will finish exactly 30 seconds before the bus arrives, keeping the original order. If no such pair is possible, return a pair with <-1,-1>.

input:

rideDuration = 90

songDuration = {1,10,25,35,60}

output:

[2,3]

solution:

You should add a condition that `map.get(complement) != i` (or it may return the same element twice.)

```
1  class Solution {
2      public int numPairsDivisibleBy60(int[] time) {
3          Map<Integer, Integer> map = new HashMap<>();
4          int count = 0;
5          for(int t : time) {
6              if(t % 60 == 0) {
7                  if(map.containsKey(0)) {
8                      count += map.get(0);
9                  }
10                 map.put(0, map.getOrDefault(0, 0) + 1);
11             } else {
12                 int diff = 60 - (t % 60);
13                 if(map.containsKey(diff)) {
14                     count += map.get(diff);
15                 }
16                 map.put(t % 60, map.getOrDefault(t % 60, 0) + 1);
17             }
18         }
19         return count;
20     }
21 }
```

$$T: O(n) \quad S: O(n)$$

```
private static int[] test(int rideDuration, int[] songDuration){
    int target = rideDuration - 30;
    int max = Integer.MIN_VALUE;
    int[] res = new int[]{-1, -1};
    if (rideDuration <= 30) return res;
    Map<Integer, Integer> map = new HashMap<>();
    for (int i = 0; i < songDuration.length; i++){
        if (map.containsKey(target - songDuration[i])){
            if (songDuration[i] > max || songDuration[map.get(target - songDuration[i])] > max) {
                res[0] = map.get(target - songDuration[i]);
                res[1] = i;
                max = Math.max(songDuration[i], songDuration[map.get(target - songDuration[i])]);
            }
        }
        map.put(songDuration[i], i);
    }
    return res;
}
```

Treasure Island

You have a map that marks the location of a treasure island. Some of the map area has jagged rocks and dangerous reefs. Other areas are safe to sail in. There are other explorers trying to find the treasure. So you must figure out a shortest route to the treasure island.

Assume the map area is a two dimensional grid, represented by a matrix of characters.

You must start from the top-left corner of the map and can move one block up, down, left or right at a time. The treasure island is marked as **X** in a block of the matrix. **X** will not be at the top-left corner. Any block with dangerous rocks or reefs will be marked as **D**. You must not enter dangerous blocks. You cannot leave the map area. Other areas **O** are safe to sail in. The top-left corner is always safe. Output the minimum number of steps to get to the treasure.

Example:

Input:

```
[[ 'O', 'O', 'O', 'O'],
 [ 'D', 'O', 'D', 'O'],
 [ 'O', 'O', 'O', 'O'],
 [ 'X', 'D', 'D', 'O']]
```

Output: 5

Explanation: Route is (0, 0), (0, 1), (1, 1), (2, 1), (2, 0), (3, 0) The minimum route takes 5 steps.

Solution:

BFS

```
private static final int[][] DIRS = {{1, 0}, {0, 1}, {-1, 0}, {0, -1}};

public static int minSteps(char[][] grid) {
    Queue<Point> q = new ArrayDeque<>();
    q.add(new Point(0, 0));
    grid[0][0] = 'D'; // mark as visited
    for (int steps = 1; !q.isEmpty(); steps++) {
        for (int sz = q.size(); sz > 0; sz--) {
            Point p = q.poll();

            for (int[] dir : DIRS) {
                int r = p.r + dir[0];
                int c = p.c + dir[1];

                if (isSafe(grid, r, c)) {
                    if (grid[r][c] == 'X') return steps;
                    grid[r][c] = 'D';
                    q.add(new Point(r, c));
                }
            }
        }
    }
    return -1;
}

private static boolean isSafe(char[][] grid, int r, int c) {
    return r >= 0 && r < grid.length && c >= 0 && c < grid[0].length && grid[r][c] != 'D';
}

private static class Point {
    int r, c;
    Point(int r, int c) {
        this.r = r;
        this.c = c;
    }
}
```

treasure island 2

You have a map that marks the locations of treasure islands. Some of the map area has jagged rocks and dangerous reefs. Other areas are safe to sail in. There are other

explorers trying to find the treasure. So you must figure out a shortest route to one of the treasure islands.

Assume the map area is a two dimensional grid, represented by a matrix of characters.

You must start from one of the starting point (marked as S) of the map and can move one block up, down, left or right at a time. The treasure island is marked as X. Any block with dangerous rocks or reefs will be marked as D. You must not enter dangerous blocks. You cannot leave the map area. Other areas O are safe to sail in. Output the minimum number of steps to get to any of the treasure islands.

Example:

Input:

```
[['S', 'O', 'O', 'S', 'S'],  
 ['D', 'O', 'D', 'O', 'D'],  
 ['O', 'O', 'O', 'O', 'X'],  
 ['X', 'D', 'D', 'O', 'O'],  
 ['X', 'D', 'D', 'D', 'O']]
```

Output: 3

Explanation:

You can start from (0,0), (0, 3) or (0, 4). The treasure locations are (2, 4) (3, 0) and (4, 0). Here the shortest route is (0, 3), (1, 3), (2, 3), (2, 4).

solution

```

final int[][] dirs = {{0, 1}, {0, -1}, {1, 0}, {-1, 0}};
public int shortestPath(char[][] islands){
    if(islands.length == 0 || islands[0].length == 0) return -1;
    int R = islands.length, C = islands[0].length;
    Queue<int[]> queue = new LinkedList<>();
    int steps = 1;
    // add all sources to queue and set 'visited'.
    for(int i = 0; i < R; ++i){
        for(int j = 0; j < C; ++j){
            if(islands[i][j] == 'S'){
                queue.add(new int[]{i, j}); islands[i][j] = 'D';
            }
        }
    }
    while(!queue.isEmpty()){
        int size = queue.size();
        for(int i = 0; i < size; ++i){
            int[] pos = queue.poll();
            for(int[] dir: dirs){
                int x = pos[0] + dir[0], y = pos[1] + dir[1];
                if(x < 0 || x >= R || y < 0 || y >= C || islands[x][y] == 'D') continue;
                if(islands[x][y] == 'E') return steps;
                queue.add(new int[]{x, y});
                islands[x][y] = 'D';
            }
        }
        ++steps;
    }
    return -1;
}

```

items in container

Given a string *s* consisting of items as "*" and closed compartments as an open and close "|", an array of starting indices *startIndices*, and an array of ending indices *endIndices*, determine the number of items in closed compartments within the substring between the two indices, inclusive.

- An item is represented as an asterisk *
- A compartment is represented as a pair of pipes | that may or may not have items between them.

Example:

$T: O(n)$ $\xi: O(n)$

`s = "|**|*|*|"`

`startIndices = [1, 1]`

`endIndices = [5, 6]`

```
public static List<Integer> numberOfItems(String s, List<Integer> startIndices, List<Integer> endIndices) {
    int[] mem = new int[s.length()];
    int count = 0;
    for (int i = 0; i < s.length(); ++i) {
        if (s.charAt(i) == '|') {
            mem[i] = count;
        } else {
            ++count;
        }
    }
    List<Integer> ans = new ArrayList<>();
    for (int i = 0; i < startIndices.size(); ++i) {
        int start = startIndices.get(i) - 1;
        int end = endIndices.get(i) - 1;

        while (start < end && s.charAt(start) != '|') ++start;
        while (start < end && s.charAt(end) != '|') --end;

        ans.add(mem[end] - mem[start]);
    }
    return ans;
}
```

The String has a total 2 closed compartments, one with 2 items and one with 1 item. For the first pair of indices, (1,5), the substring is '|**|*|'. There are 2 items in a compartment.

For the second pair of indices, (1,6), the substring is '|**|*|' and there 2+1=3 items in compartments.

Both of the answers are returned in an array. [2,3].

solution:

turnstile

9. Turnstile

A university has exactly one turnstile. It can be used either as an exit or an entrance. Unfortunately, sometimes many people want to pass through the turnstile and their directions can be different. The i^{th} person comes to the turnstile at `time[i]` and wants to either exit the university if `direction[i] = 1` or enter the university if `direction[i] = 0`. People form 2 queues, one to exit and one to enter. They are ordered by the time when they came to the turnstile and, if the times are equal, by their indices.

If some person wants to enter the university and another person wants to leave the university at the same moment, there are three cases:

- If in the previous second the turnstile was not used (maybe it was used before, but not at the previous second), then the person who wants to leave goes first.
- If in the previous second the turnstile was used as an exit, then the person who wants to leave goes first.
- If in the previous second the turnstile was used as an entrance, then the person who wants to enter goes first.

Passing through the turnstile takes 1 second.

For each person, find the time when they will pass through the turnstile.

Function Description

Complete the function `getTimes` in the editor below.

`getTimes` has the following parameters:

`int time[n]`: an array of n integers where the value at index i is the time in seconds when the i^{th} person will come

Function Description

Complete the function `getTimes` in the editor below.

`getTimes` has the following parameters:

`int time[n]`: an array of n integers where the value at index i is the time in seconds when the i^{th} person will come to the turnstile

`int direction[n]`: an array of n integers where the value at index i is the direction of the i^{th} person

Returns:

`int[n]`: an array of n integers where the value at index i is the time when the i^{th} person will pass the turnstile

Constraints

- $1 \leq n \leq 10^5$
- $0 \leq \text{time}[i] \leq 10^9$ for $0 \leq i \leq n - 1$
- $\text{time}[i] \leq \text{time}[i + 1]$ for $0 \leq i \leq n - 2$
- $0 \leq \text{direction}[i] \leq 1$ for $0 \leq i \leq n - 1$

► Input Format For Custom Testing

▼ Sample Case 0

Sample Input 0

▼ Sample Case 0

Sample Input 0

STDIN	Function
-----	-----
4	→ time[] size n = 4
0	→ time = [0, 0, 1, 5]
0	
1	
5	
4	→ direction[] size n = 4
0	→ direction = [0, 1, 1, 0]
1	
1	
0	

Sample Output 0

```
2
0
1
5
```

solution:

```
public static int[] getTime(int[] time, int[] direction) {
    Queue<Integer> exitQ = new LinkedList<>();
    Queue<Integer> entryQ = new LinkedList<>();
    int[] res = new int[time.length];
    int lastTime = -2;
    for (int i = 0; i < direction.length; i++) {
        if (direction[i] == 1) exitQ.offer(i);
        else entryQ.offer(i);
    }
    Queue<Integer> lastQ = exitQ;
    while (!entryQ.isEmpty() && !exitQ.isEmpty()) {
        int currentTime = lastTime + 1;
        int peekEntryTime = time[entryQ.peek()];
        int peekExitTime = time[exitQ.peek()];
        Queue<Integer> q;
        if (currentTime < peekExitTime && currentTime < peekEntryTime) {
            q = (peekEntryTime < peekExitTime) ? entryQ : exitQ;
            int personIdx = q.poll();
            res[personIdx] = time[personIdx];
            lastTime = time[personIdx]; // time
            lastQ = q;
        } else {
            if (currentTime >= peekExitTime && currentTime >= peekEntryTime) {
                q = lastQ;
            } else {
                q = currentTime >= peekEntryTime ? entryQ : exitQ;
            }
            int personId = q.poll();
            res[personId] = currentTime;
            lastTime = currentTime;
            lastQ = q;
        }
    }
    lastQ = entryQ.size() > 0 ? entryQ : exitQ;
    while (!lastQ.isEmpty()) {
        int currentTime = lastTime + 1;
        int personId = lastQ.poll();
        if (currentTime < time[personId]) {
            currentTime = time[personId];
        }
        res[personId] = currentTime;
        lastTime = currentTime;
    }
    return res;
}

enum Dir {
    EXIT,
    ENTER;
}

static class Person {
    int idx;
    int time;

    public Person(int i, int time) {
        this.idx = i;
        this.time = time;
    }
}
```

$O(n)$

$O(n)$

```

public class Turnstile {
    public static int[] getTimes(int[] time, int[] direction) {
        Queue<Integer> enters = new LinkedList<Integer>();
        Queue<Integer> exits = new LinkedList<Integer>();
        int n = time.length;
        for(int i = 0; i < n; i++) {
            Queue<Integer> q = direction[i] == 1 ? exits : enters;
            q.offer(i);
        }

        int[] result = new int[n];
        int lastTime = -2;
        Queue<Integer> lastQ = exits;
        while(enters.size() > 0 && exits.size() > 0) {
            int currentTime = lastTime + 1;
            int peekEnterTime = time[enters.peek()];
            int peekExitTime = time[exits.peek()];
            Queue<Integer> q;
            if (currentTime < peekEnterTime && currentTime < peekExitTime) {
                // The turnstile was not used
                // Take whoever has the earliest time or
                // if enter == exit, take exit
                q = (peekEnterTime < peekExitTime) ? enters : exits;
                int personIdx = q.poll();
                result[personIdx] = time[personIdx];
                lastTime = time[personIdx]; // time
                lastQ = q;
            } else {
                // Turnstile was used last second
                if (currentTime >= peekEnterTime && currentTime >= peekExitTime) {
                    // Have people waiting at both ends
                    // Prioritize last direction
                    q = lastQ;
                } else {
                    // current >= enters.peek() || current >= exits.peek()
                    q = currentTime >= peekEnterTime ? enters : exits; // take whatever that's queuing
                }
                int personIdx = q.poll();
                result[personIdx] = currentTime;
                lastTime = currentTime; // time
                lastQ = q;
            }
        }

        Queue<Integer> q = enters.size() > 0 ? enters : exits;
        while(q.size() > 0) {
            int currentTime = lastTime + 1;
            int personIdx = q.poll();
            if (currentTime < time[personIdx]) {
                // The turnstile was not used
                currentTime = time[personIdx];
            }

            result[personIdx] = currentTime;
            lastTime = currentTime; // time
        }

        return result;
    }
}

```

Optimal Utilization

Given 2 lists `a` and `b`. Each element is a pair of integers where the first integer represents the unique id and the second integer represents a value. Your task is to find an element from `a` and an element from `b` such that the sum of their values is less or equal to `target` and as close to `target` as possible. Return a list of ids of selected elements. If no pair is possible, return an empty list.

Example 1:

Input:

```
a = [[1, 2], [2, 4], [3, 6]]
```

```
b = [[1, 2]]
```

```
target = 7
```

Output: `[[2, 1]]`

Explanation:

There are only three combinations `[1, 1]`, `[2, 1]`, and `[3, 1]`, which have a total sum of 4, 6 and 8, respectively.

Since 6 is the largest sum that does not exceed 7, `[2, 1]` is the optimal pair.

Example 2:

Input:

```
a = [[1, 3], [2, 5], [3, 7], [4, 10]]
```

```
b = [[1, 2], [2, 3], [3, 4], [4, 5]]
```

```
target = 10
```

Output: `[[2, 4], [3, 2]]`

Explanation:

There are two pairs possible. Element with id = 2 from the list ``a`` has a value 5, and element with id = 4 from the list ``b`` also has a value 5.

Combined, they add up to 10. Similarly, element with id = 3 from ``a`` has a value 7, and element with id = 2 from ``b`` has a value 3.

These also add up to 10. Therefore, the optimal pairs are `[2, 4]` and `[3, 2]`.

solution:

$O(1)$

```

private List<int[]> getPairs(List<int[]> a, List<int[]> b, int target) {
    Collections.sort(a, (i,j) -> i[1] - j[1]);
    Collections.sort(b, (i,j) -> i[1] - j[1]);
    List<int[]> result = new ArrayList<>();
    int max = Integer.MIN_VALUE;
    int m = a.size();
    int n = b.size();
    int i = 0;
    int j = n-1;
    while(i<m && j >= 0) {
        int sum = a.get(i)[1] + b.get(j)[1];
        if(sum > target) {
            --j;
        } else {
            if(max <= sum) {
                if(max < sum) {
                    max = sum;
                    result.clear();
                }
                result.add(new int[]{a.get(i)[0], b.get(j)[0]});
                int index = j-1;
                while(index >= 0 && b.get(index)[1] == b.get(index+1)[1]) {
                    result.add(new int[]{a.get(i)[0], b.get(index--)[0]});
                }
            }
            ++i;
        }
    }
    return result;
}

```

$O(n \log n)$

$O(n)$

描述是说把一个array 变为sorted 需要多少次swap 实质和315是一样的

用的就睡mergesort, testcase 全过

Given an array and a sorting algorithm, the sorting algorithm will do a selection swap. Find

the number of swaps to sort the array.

Let N be the length of nums.

T: $O(N \log(N))$ Merge Sort

$S O(N)$

Example 1:

Input: [5, 4, 1, 2]

Output: 5

Explanation:

Swap index 0 with 1 to form the sorted array [4, 5, 1, 2].

Swap index 0 with 2 to form the sorted array [1, 5, 4, 2].

Swap index 1 with 2 to form the sorted array [1, 4, 5, 2].

Swap index 1 with 3 to form the sorted array [1, 2, 5, 4].

Swap index 2 with 3 to form the sorted array [1, 2, 4, 5].

solution :

<https://www.geeksforgeeks.org/number-swaps-sort-adjacent-swapping-allowed/>

<https://www.geeksforgeeks.org/counting-inversions/>

- Letter-logs: All words (except the identifier) consist of lowercase English letters.
- Digit-logs: All words (except the identifier) consist of digits.

Reorder these logs so that:

1. The letter-logs come before all digit-logs.
2. The letter-logs are sorted lexicographically by their contents. If their contents are the same, then sort them lexicographically by their identifiers.
3. The digit-logs maintain their relative ordering.

Return *the final order of the logs*.

Example 1:

Input: logs = ["dig1 8 1 5 1","let1 art can","dig2 3 6","let2 own kit dig","let3 art zero"]

Output: ["let1 art can","let3 art zero","let2 own kit dig","dig1 8 1 5 1","dig2 3 6"]

Explanation:

The letter-log contents are all different, so their ordering is "art can", "art zero", "own kit dig".

The digit-logs have a relative order of "dig1 8 1 5 1", "dig2 3 6".

Example 2:

Input: logs = ["a1 9 2 3 1","g1 act car","zo4 4 7","ab1 off key dog","a8 act zoo"]

Output: ["g1 act car","a8 act zoo","ab1 off key dog","a1 9 2 3 1","zo4 4 7"]

Complexity Analysis

solution:

Let N be the number of logs in the list and M be the maximum length of a single log.

- Time Complexity: $\mathcal{O}(M \cdot N \cdot \log N)$
 - First of all, the time complexity of the `Arrays.sort()` is $\mathcal{O}(N \cdot \log N)$, as stated in the [API specification](#), which is to say that the `compare()` function would be invoked $\mathcal{O}(N \cdot \log N)$ times.
 - For each invocation of the `compare()` function, it could take up to $\mathcal{O}(M)$ time, since we compare the contents of the logs.
 - Therefore, the overall time complexity of the algorithm is $\mathcal{O}(M \cdot N \cdot \log N)$.
- Space Complexity: $\mathcal{O}(M \cdot \log N)$
 - For each invocation of the `compare()` function, we would need up to $\mathcal{O}(M)$ space to hold the parsed logs.
 - In addition, since the implementation of `Arrays.sort()` is based on quicksort algorithm whose space complexity is $\mathcal{O}(\log n)$, assuming that the space for each element is $\mathcal{O}(1)$. Since each log could be of $\mathcal{O}(M)$ space, we would need $\mathcal{O}(M \cdot \log N)$ space to hold the intermediate values for sorting.
 - In total, the overall space complexity of the algorithm is $\mathcal{O}(M + M \cdot \log N) = \mathcal{O}(M \cdot \log N)$.

```

public String[] reorderLogFiles(String[] logs) {
    Comparator<String> myComp = new Comparator<String>() {
        @Override
        public int compare(String log1, String log2) {
            // split each log into two parts: <identifier, content>
            String[] split1 = log1.split(" ", 2);
            String[] split2 = log2.split(" ", 2);

            boolean isDigit1 = Character.isDigit(split1[1].charAt(0));
            boolean isDigit2 = Character.isDigit(split2[1].charAt(0));

            // case 1). both logs are letter-logs
            if (!isDigit1 && !isDigit2) {
                // first compare the content
                int cmp = split1[1].compareTo(split2[1]);
                if (cmp != 0)
                    return cmp;
                // logs of same content, compare the identifiers
                return split1[0].compareTo(split2[0]);
            }

            // case 2). one of logs is digit-log
            if (!isDigit1 && isDigit2)
                // the letter-log comes before digit-logs
                return -1;
            else if (isDigit1 && !isDigit2)
                return 1;
            else
                // case 3). both logs are digit-log
                return 0;
        }
    };

    Arrays.sort(logs, myComp);
    return logs;
}

```

Transaction logs

A Company parses logs of online store user transactions/activity to flag fraudulent activity.

The log file is represented as an Array of arrays. The arrays consist of the following data:

[<# of transactions>]

For example:

[345366 89921 45]

Note: the data is space delimited

So, the log data would look like:

[

[345366 89921 45],

[029323 38239 23]

...

]

Write a function to parse the log data to find distinct users that meet or cross a certain threshold.

The function will take in 2 inputs:

logData: Log data in form an array of arrays

threshold: threshold as an integer

Output:

It should be an array of userids that are sorted.

If same userid appears in the transaction as userid1 and userid2, it should count as one occurrence, not two.

Example:

Input:

logData:

[

[345366 89921 45],

[029323 38239 23],

[38239 345366 15],

[029323 38239 77],

[345366 38239 23],

[029323 345366 13],

[38239 38239 23]

...

]

threshold: 3

Output: [345366 , 38239, 029323]

Explanation:

Given the following counts of userids, there are only 3 userids that meet or exceed the threshold of 3.

345366 -4 , 38239 -5, 029323-3, 89921-1

$O(m \times n)$

```

static List<String> processLogs(List<String> logs, int threshold) {
    Map<String, Integer> map = new HashMap<>();
    for (String logLine : logs) {
        String[] log = logLine.split(" ");
        map.put(log[0], map.getDefault(log[0], 0) + 1);
        if (log[0] != log[1]) {
            map.put(log[1], map.getDefault(log[1], 0) + 1);
        }
    }

    List<String> userIds = new ArrayList<>();
    for (Map.Entry<String, Integer> entry : map.entrySet()) {
        if (entry.getValue() >= threshold) {
            userIds.add(entry.getKey());
        }
    }

    Collections.sort(userIds, new Comparator<String>() {
        @Override
        public int compare(String s1, String s2) {
            return Integer.parseInt(s1) - Integer.parseInt(s2);
        }
    });

    return userIds;
}

```

AMAZON MUSIC PAIRS 60min

Amazon Music is working on a new “community radio station” feature which will allow users to iteratively populate

the playlist with their desired songs. Considering this radio station will also have other scheduled shows to be

aired, and to make sure the community soundtrack will not run over other scheduled shows, the user-submitted songs

will be organized in full-minute blocks. Users can choose any songs they want to add to the community list, but

only in pairs of songs with durations that add up to a multiple of 60 seconds (e.g. 60, 120, 180).

As an attempt to insist on the highest standards and avoid this additional burden on users, Amazon will let them

submit any number of songs they want, with any duration, and will handle this 60-second matching internally. Once

the user submits their list, given a list of song durations, calculate the total number of distinct song pairs that

can be chosen by Amazon Music.

A better way can reach $O(N)$.

We just need to get the remainders of each element, for example, [60, 60, 60] will become [0, 0, 0]. Then we know there are 3 elements that can be divided by 60. So we just need to pick 2 of them to form a pair, so there will be $n*(n-1)/2$ pairs.

In other cases, if we have x with count of a and (60 - x) with count of b, then we can form a*b pairs.

So in summary:

- 1): get the remainders of each element;
- 2): count the number of each remainder;
- 3): analyze the counts to find the valid pair

The time complexity is $O(N)$, and space complexity is $O(N)$.

See the code below:

```
int findPair(vector<int>& nums) {  
    int res = 0;  
    vector<int> cts(60, 0);  
    for(auto &a : nums) ++cts[a%60];  
    for(int i=1; i<30; ++i) res += cts[i]*cts[60-i];  
    res += cts[0]*(cts[0]-1)/2 + cts[30]*(cts[30]-1)/2;  
    return res;  
}
```

same thing with java:

```
1  class Solution {  
2      public int numPairsDivisibleBy60(int[] time) {  
3          Map<Integer, Integer> map = new HashMap<>();  
4          int count = 0;  
5          for(int t : time) {  
6              if(t % 60 == 0) {  
7                  if(map.containsKey(0)) {  
8                      count += map.get(0);  
9                  }  
10                 map.put(0, map.getOrDefault(0, 0) + 1);  
11             } else {  
12                 int diff = 60 - (t % 60);  
13                 if(map.containsKey(diff)) {  
14                     count += map.get(diff);  
15                 }  
16                 map.put(t % 60, map.getOrDefault(t % 60, 0) + 1);  
17             }  
18         }  
19         return count;  
20     }  
21 }
```

```

public int numPairsDivisibleBy60(int[] time) {
    int res = 0;
    int[] rem = new int[60];
    for (int t : time) rem[t % 60]++;
    for (int i = 1; i < 30; i++) {
        res += rem[i] * rem[60 - i];
    }
    res += rem[0] * (rem[0] - 1) / 2 + rem[30] * (rem[30] - 1) / 2;
    return res;
}

```

space is o1 not on

subfiles into a single file

Write an Algorithm to output the minimum possible time to merge the given N subfiles into a single file

Input: The input to the function/method consists of two arguments:

numOfSubFiles: an integer representing the number of subfiles;

files: a list of integers representing the size of the compressed subfiles

Output: Return an integer representating the minimum time required to merge all the subfiles

Constraints::

$2 \leq \text{numOfSubFiles} \leq 10^6$

$1 \leq \text{files} \leq 10^6$

Example:

input:

numOfSubFiles = 4

files = [4,8,6,12]

Output: 58

Explanation:

The optimal way to merge subfiles is as follows:

Step 1: Merge the files of size 4 and 6 (time required is 10). Size of subfiles after merging.

[8,10,12]

Step 2: Merge the files of size 8 and 10 (time required is 18). Size of subfiles after merging. [18,12]

Step 3: Merge the files of size 18 and 12 (time required is 30)

Total time required to merge the file is $10 + 18 + 30 = 58$.

solution:

```
public static int cost(int n, int[] arr) {  
    if (n < 2) return 0;  
    Arrays.sort(arr);  
    int res = 0;  
    int prevSum = arr[0];  
    for (int i = 1; i < arr.length; i++) {  
        prevSum += arr[i];  
        res += prevSum;  
    }  
    return res;  
}
```

pq solution: