CSCB09 Software Tools and Systems Programming Process Communication – Sockets

Marcelo Ponce

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Department of Computer and Mathematical Sciences - UTSC

Today's class

Today we will discuss the following topics:

Sockets

Quick Recap

- Exit codes
- Signals

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- Signals
- Pipes

In order to use a pipe for interprocess communication what has to be true?

- 1. The process writing to the pipe must be the parent of the process reading from the pipe.
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- 3. The processes must be a parent/child combination but the direction of the pipe could go either way.
- 4. The processes must be on the same machine.

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Within the same machine

- Exit codes
- Signals
- Pipes

To communicate processes across different machines...

Within the same machine

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To communicate processes across different machines...

⇒ SOCKETS

- process → program running within a host
- within same host, two processes can communicate using *inter-process* communication (defined by OS)
- processes in different hosts communicate by exchanging messages

Client-Server Paradigm

Clients & Servers

client process: process that initiates communication

server process: process that waits to be contacted

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Server

- Running on end host
- Provides service
- E.g., Web server

Client

- Sometimes ON
- Initiates a request to the server when interested
- E.g., web browser
- Needs to know the server's address
- Shows up and initiates the connection

Server

- Always ON
- Serve services to many clients
- E.g., www.utsc.utoronto.ca
- Not initiate contact with the clients
- Needs a fixed address
- Always available at a known location

Sockets

Application Programming Interface – Sockets

- Socket Interface was originally provided by the Berkeley distribution of Unix (BSD)
- Now supported in virtually all operating systems
- Each protocol provides a certain set of services, and the API provides a syntax by which those services can be invoked in this particular OS

Application Programming Interface - Sockets

- Socket Interface was originally provided by the Berkeley distribution of Unix (BSD)
- Now supported in virtually all operating systems
- Each protocol provides a certain set of services, and the API provides a syntax by which those services can be invoked in this particular OS
- Once you set up the socket, the interface behaves like a file (or pipe)
 - entry in the file descriptor table
 - read/write/close
- Except sockets are full-duplex (2 way)

Across the network...

When you think about the task at hand of sending information across different machines and employing a complex and complicated underlying computer network (e.g. the internet), things may look even much more complicated...

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- Network Protocols
- layers
- clients & servers
- addresses, ports, TCP/IP
- packets
- routing

Because of the way in which computer networks have been designed, following a *layer approach* we can just focus in learning the basics of the "application" layer.

Across the network...

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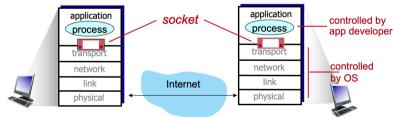
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Because of the way in which computer networks have been designed, following a layer approach we can just focus in learning the basics of the "application" layer. Having said that, if you are interested in learning the way in which computer networks work, consider taking

CSCD58

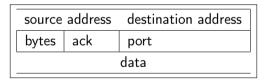
Sockets

- process sends/receives messages to/from its socket
- socket analogous to door
 - sending process shoves message out door
 - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process
- two sockets involved: one on each side



TCP/IP

- Transmission Control Protocol/Internet Protocol
- Tell us how to package up the data



Packet Details

make packets

00100010001111 10100010111 TCP

01100111001001

•Each TCP packet is given a header -sequence number -checksum

 put in an IP envelope with another header

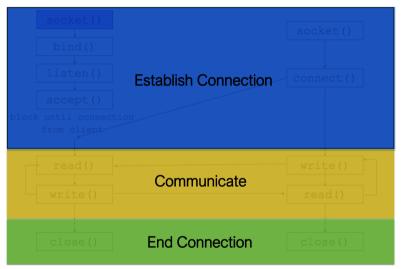




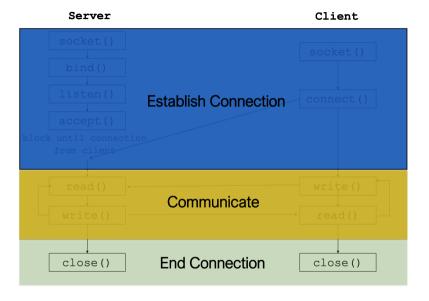


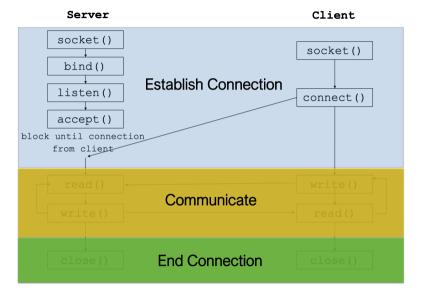


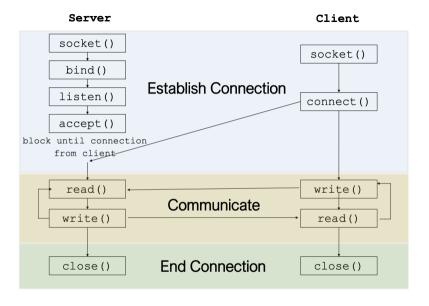
To 24.197.0.67 To 24.197.0.67 To 24.197.0.67 To 24.197.0.67 Server Client



Server Client **Establish Connection** read() write() Communicate write() read() **End Connection**







Connection-Oriented

Server

- Create a socket: socket()
- Assign an address to a socket: bind()
- Establish a queue for connections: listen()
- Get a connection from the queue: accept()

Client

- Create a socket: socket()
- Initiate a connection: connect()

Addressing Processes

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Addressing Processes

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- »» Q: does IP address of host on which process runs suffice for identifying the process?
 - » A: no, many processes can be running on same host

- identifier includes both IP address and port numbers associated with process on host.
- example port numbers:

• HTTP server: 80

mail server: 25

 to send HTTP message to gaia.cs.umass.edu web server:

• IP address: 128.119.245.12

• port number: 80

Addresses & Ports

- A socket pair is the two endpoints of the connection.
- An endpoint is identified by an IP address and a port.
- IPv4 addresses are 4 8-bit numbers:
 - 128.100.31.200 = werewolf
 - 128.100.31.201 = seawolf
 - 128.100.31.202 = skywolf
- Ports:
 - because multiple processes can communicate with a single machine we need another identifier.
- command-line tool to explore: nslookup

Ports

The *port* number is a 16-bit integer that ranges from 0 to 65535.

Types of port number space:

Pre-established ports – 0-1023

80: http 21: ftp

22: ssh 25: smtp (email)

23: telnet 194: irc

Registered port – 1024-49151

2709: supermon

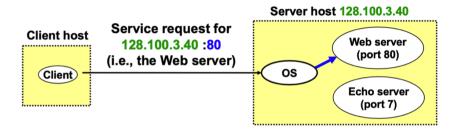
26000: quake

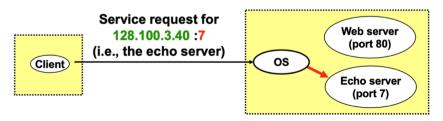
3724: world of warcraft

Dynamic port – 49152-65535
 pick ports in this range to avoid overlap

Number	Assignment
20	File Transfer Protocol (FTP) Data Transfer
21	File Transfer Protocol (FTP) Command Control
22	Secure Shell (SSH) Secure Login
23	Telnet remote login service, unencrypted text messages
25	Simple Mail Transfer Protocol (SMTP) email delivery
53	Domain Name System (DNS) service
67, 68	Dynamic Host Configuration Protocol (DHCP)
80	Hypertext Transfer Protocol (HTTP) used in the World Wide Web
110	Post Office Protocol (POP3)
119	Network News Transfer Protocol (NNTP)
123	Network Time Protocol (NTP)
143	Internet Message Access Protocol (IMAP) Management of digital mail
161	Simple Network Management Protocol (SNMP)
194	Internet Relay Chat (IRC)
443	HTTP Secure (HTTPS) HTTP over TLS/SSL
546, 547	DHCPv6 IPv6 version of DHCP

Using ports to identify services





Unix Socket API

Socket interface

- A collection of system calls to write a networking program at user-level.
- Originally provided in Berkeley UNIX
- Later adopted by all popular operating systems

System calls for sockets

- Client:
 - create, connect, write, read, close
- Server:
 - create, bind, listen, accept, read, write, close

In UNIX, everything is like a file

- All input is like reading a file
- All output is like writing a file
- File is represented by an integer file descriptor
- Data written into socket on one host can be read out of socket on other host

Socket:

- The point where a local application process attaches to the network
- An interface between an application and the network
- An application creates the socket

The interface defines operations for

- Creating a socket
- Attaching a socket to the network
- Sending and receiving messages through the socket
- Closing the socket

Socket: end point of communication

- Processes send messages to one another
- Messages traverse the underlying network

A process sends and receives through a "socket"

- Analogy: the doorway of the house
- Socket, as an API, supports the creation of network applications

Socket parameters

A socket connection has 5 general parameters:

 The protocol Example: TCP and UDP.

• The local and remote address Example: 128.100.3.40

- The local and remote port number
 - Some ports are reserved (e.g., 80 for HTTP)
 - Root access require to listen on port numbers below 1024

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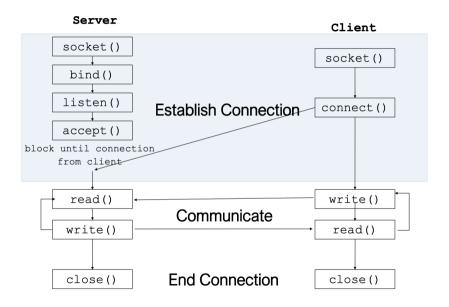
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Socket Family

- PF_INET denotes the Internet family
- PF_UNIX denotes the Unix pipe facility
- PF_PACKET denotes direct access to the network interface (i.e., it bypasses the TCP/IP protocol stack)

Socket Type

- SOCK_STREAM is used to denote a byte stream
- SOCK_DGRAM is an alternative that denotes a message oriented service, such as that provided by UDP



Typical Client Program

Prepare to communicate

- Create a socket
- Determine server address and port number
- Initiate the connection to the server

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Exchange data with the server

- Write data to the socket
- Read data from the socket
- Do something with the data (eg. render a webpage)

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- Create a socket
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Exchange data with the server

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socket() create the socket descriptor

connect() connect to the remote server

read(),write() communicate
 with the server

close() end communication
by closing socket
descriptor

Close the socket

Prepare to communicate

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Exchange data with the client over new socket

- Receive data from the socket
- Send data to the socket
- Close the socket

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- Receive data from the socket.
- Send data to the socket.
- Close the socket

Repeat with the next connection request

CSCB09|Week 09: Process Comm. - Sockets

bind() associate the local address

listen() wait for incoming connections from clients

accept() accept incoming connection

read(), write() communicate with the client

> close() end communication by closing socket descriptor

Server side

```
int socket(int family, int type, int protocol);
```

- family specifies protocol family:
 - AF_INET: IPv4
 - AF_LOCAL: UNIX domain
- typeSOCK_STREAM, SOCK_DGRAM, SOCK_RAW
- protocol set to 0 except for RAW sockets
- returns a socket descriptor

Creating a Socket

```
int sockfd = socket(address_family, type, protocol);
```

- The socket number returned is the socket descriptor for the newly created socket
- TCP SOCK_STREAM: Connection oriented
 int sockfd = socket (PF_INET, SOCK_STREAM, 0);
- UDP SOCK_DGRAM: Connectionless
 int sockfd = socket (PF_INET, SOCK_DGRAM, 0);

Bind to a name

```
int bind(int sockfd, const struct sockaddr *servaddr, socklen_t addrlen);
```

sockfd – returned by socket()

- Binds the newly created socket to the specified address i.e. the network address of the local participant (the server)
- Address is a data structure which combines IP and port
- sin_addr.s_addr can be set to INADDR_ANY to communicate on any network interface
- addrlen is the size of struct sockaddr_in

Set up queue in kernel

```
int listen(int sockfd, int backlog);
```

- after calling listen(), a socket is ready to accept connections
 - prepares a queue in the kernel where partially completed connections wait to be accepted
- defines how many connections can be pending on the specified socket
- backlog is the maximum number of partially completed connections that the kernel should queue

Complete the connection

```
int accept(int sockfd, struct sockaddr *cliaddr, socklen_t *addrlen);
```

- carries out the passive open
- blocks waiting for a connection (from the queue)
- returns a new descriptor which refers to the TCP connection with the client
- sockfd is the listening socket
- cliaddr stores the address of the client on return
- reads and writes on the connection will use the socket returned by accept

Client side

- socket() same as server, to say "how" we are going to talk
 int connect(int sockfd, const struct sockaddr *servaddr, socklen_t
 addrlen);
- the kernel will choose a dynamic port and source IP address.
- returns 0 on success and -1 on failure setting errno.

Sending and Receiving Data

Sending Data

```
write(int sockfd, void *buf, size_t len);
Arguments: socket descriptor, pointer to buffer of data, and length of the buffer
Returns the number of characters written, and -1 on error
```

Receiving data

```
read(int sockfd, void *buf, size_t len);
Arguments: socket descriptor, pointer to buffer to place the data, size of the buffer
Returns the number of characters read (where 0 implies "end of file"), and -1 on
error
```

Important Network Details

Byte Order - aka "Endianess"

Big-endian

The most significant byte (the "big end") of the data is placed at the byte with the lowest address

Little-endian

The least significant byte (the "little end") of the data is placed at the byte with the lowest address

$$\mathbf{91}, \mathbf{329_{10}} = \boxed{\texttt{C1}} \boxed{\texttt{64}} \boxed{\texttt{01}} \boxed{\texttt{00}}$$

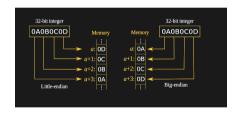
Machines differ in how they store data
 E.g. Intel is little-endian, Sparc is big-endian, ...

Network Byte Order

- To communicate between machines with unknown or different "endian-ness" we convert numbers to network byte order (big-endian) before we send them.
- There are functions provided to do this:

```
unsigned long htonl(unsigned long)
unsigned short htons(unsigned short)
unsigned long ntohl(unsigned long)
unsigned short ntohs(unsigned short)
```

 Use htons() and htonl() to convert to network byte order
 Use ntohs() and ntohl() to convert to host order



Network new-line

"\r\n" rather than just "\n"

Application: Client-Server Model

with TCP

Client-Server Model with TCP

Server

- Passive open
 Prepares to accept connection, does not actually establish a connection
- Hearing from multiple clients
 Allow a backlog of waiting clients
 ... in case several try to start a connection at once
- Create a socket for each client
 Upon accepting a new client
 ... create a new socket for the
 communication

Client

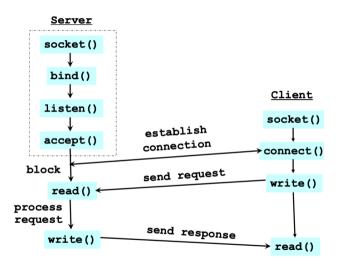
- Application performs active open
- It say swho it wants to communicate with
- Client invokes
 int connect (int socket, struct
 sockaddr *address, int addr_len)
- Connect
 Does not return until TCP has successfully established a connection at which application is free to begin sending data
 Address contains remote machine's address

Client-Server Model with TCP

In practice

- The client usually specifies only remote participant's address and let's the system fill in the local information
- Whereas a server usually listens for messages on a well-known port
- A client does not care which port it uses for itself, the OS simply selects an unused one

Once a connection is established, the application process invokes two operation int send (int socket, char *msg, int msg_len, int flags) int recv (int socket, char *buff, int buff_len, int flags)



```
gethostbyname() // get address for given host name (e.g. 128.100.3.40 for name
     "cs.toronto.edu"):
getservbyname() // get port and protocol for a given service e.g. ftp, http (e
    .g. "http" is port 80, TCP)
getsockname() // get local address and local port of a socket
getpeername() // get remote address and remote port of a socket
struct sockaddr { // Generic address, "connect(), bind(), accept()"
  u short sa family: // <svs/socket.h>
  char sa_data[14];
};
struct sockaddr in { // Client & server addresses TCP/UDP address (inc.port
    #)
  u_short sa_family; // <netinet/in.h>
  u_short sa_family;
  u_short sin_port;
  struct in_addr sin_addr:
  char sin zero[8]:
}:
struct in addr { // IP address <netinet/in.h>
  u long s addr:
};
```

```
// Convert between system's representation of IP addresses and readable
    strings (e.g. "128.100.3.40")
unsigned long inet_addr(char* str);
char * inet_ntoa(struct in_addr inaddr);

// Relevant header files:
<sys/types.h>
<sys/socket.h>
<netinet/in.h>
<arpa/inet.h>

// man pages
socket, accept, bind , listen
```