


Book: Introduction to Algorithms 3e

Chap: 17.1 Aggregate Analysis

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17.1-1 理想情况下没有上限, 不为 $O(1)$, 而是 $O(k)$

17.1-2  不停抖动为 $\Theta(nk)$
附加一段推导:

$$\begin{aligned} T_n &= \sum_{i=0}^{k-1} \lfloor \frac{n}{2^i} \rfloor \\ &= \sum_{i=0}^{k-1} \left(\frac{n}{2^i} - \left\{ \frac{n}{2^i} \right\} \right) \\ &= n \sum_{i=0}^{k-1} \frac{1}{2^i} \Delta x - \sum_{i=0}^{k-1} \left\{ \frac{n}{2^i} \right\} \\ &= n \cdot \left(2 - \frac{1}{2^{k-1}} \right) - \underbrace{\sum_{i=0}^{k-1} \left\{ \frac{n}{2^i} \right\}}_{\in [0, k)} \end{aligned}$$

$$\frac{T_n}{n} < 2 - \frac{1}{2^{k-1}} = O(1)$$

17.1-3 $T_n = \sum_{\substack{1 \leq i \leq n \\ m}} \lfloor i = 2^m \rfloor (2^m - 1) + n$ (使用艾弗森约定)

$$= \sum_m \{ 0 \leq m \leq \lceil \lg n \rceil \} (2^m - 1) + n$$

$$= \sum_0^{\lceil \lg n \rceil} (2^x - 1) \Delta x + n \text{ (使用《Concrete Mathematics》的定和式技巧)}$$

$$= (2^x - x) \Big|_0^{\lceil \lg n \rceil} + n$$

$$= 2^{\lceil \lg n \rceil} - \lceil \lg n \rceil - 1 + n$$

$$T_n/n = (2^{\lceil \lg n \rceil} - \lceil \lg n \rceil - 1 + n)/n$$

$$\leq (2n - \lceil \lg n \rceil)/n = \Theta(1)$$