

ECS 122A: Algorithm Design and Analysis

Week 3 Discussion

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A few words on logistics

- ▶ Homework 2 due tomorrow (Apr 15) at midnight
- ▶ Homework 3 is released, due next Tuesday (Apr 20) at midnight
- ▶ When submitting, please select page(s) to each question, preferably in PDFs
- ▶ Homework is graded by attempting (50%) + one selected problem (50%)
- ▶ Midterm 1 is next Thursday, will cover up to and including homework 3

Outline

- ▶ Divide and Conquer: Key idea
- ▶ Solve Divide and Conquer Recurrence: Master Theorem
- ▶ Variant of Maximum-Subarray Problem: Stock Investment

Divide and Conquer: Key idea

1. **Divide** the problem into a number of subproblems that are smaller instances of the **same** problem.
2. **Conquer** by solving the subproblems **recursively**.
3. **Combine** the solutions to the subproblems to produce the solution to the original problem.

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$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

So far, we demonstrated three problems that use divide and conquer paradigm:

- ▶ Merge Sort
- ▶ Maximum Subarray
- ▶ Matrix-matrix multiply (Strassen's algorithm)

Using Master Theorem

What is the asymptotic bound for the given recurrence?

$$T(n) = 3T\left(\frac{n}{4}\right) + n$$

$$a =$$

$$b =$$

$$f(n) =$$

Using Master Theorem

What is the asymptotic bound for the given recurrence? Does Master Theorem apply to this recurrence?

$$T(n) = 2T\left(\frac{n}{2}\right) + n \lg n$$

Stock Investment

Problem Statement:

We're doing a simulation in which we look at n consecutive days of a given stock, at some point in the past. Let's number the days $i = 1, 2, \dots, n$ and $p(i)$ is the price per share for the stock on that day. We want to know: When should we have bought and sold in order to have made as much money as possible? (If there was no way to make money during the n days, we should report this instead.)

For example, $n = 4$, $p(1) = 9$, $p(2) = 1$, $p(3) = 5$, $p(4) = 3$. Then we should answer "buy on day 2, sell on day 3".

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Rephrase the problem:

Input: array p of length n

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Rephrase the problem:

Input: array p of length n

Output: $\operatorname{argmax}\{p(j) - p(i)\}$ where $i \leq j$

Stock Investment

Revisit maximum subarray problem:

1. Divide $A[low \cdots high]$ into two subarrays of as equal size as possible by finding the midpoint mid
2. Conquer:
 - a. finding maximum subarrays of $A[low \cdots mid]$ and $A[mid + 1 \cdots high]$
 - b. finding a max-subarray that crosses the midpoint
3. Combine: returning the max of the three

Stock Investment

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Can we apply the same strategy on this problem?

Stock Investment

How does the conquer part work?

1. The optimal solution to $A[low \cdots mid]$
2. The optimal solution to $A[mid + 1 \cdots high]$
3. $\operatorname{argmax}\{p(j) - p(i)\}$ where $low \leq i \leq mid$ and $mid + 1 \leq j \leq high$

Stock Investment

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1. The optimal solution to $A[low \cdots mid]$
2. The optimal solution to $A[mid + 1 \cdots high]$
3. $\operatorname{argmax}\{p(j) - p(i)\}$ where $low \leq i \leq mid$ and $mid + 1 \leq j \leq high$ (equivalent to finding the index of min of $A[low \cdots mid]$ and that of max of $A[mid + 1 \cdots high]$)

Stock Investment

How does the conquer part work?

1. The optimal solution to $A[low \cdots mid]$
2. The optimal solution to $A[mid + 1 \cdots high]$
3. $\operatorname{argmax}\{p(j) - p(i)\}$ where $low \leq i \leq mid$ and $mid + 1 \leq j \leq high$ (**equivalent to** finding the index of \min of $A[low \cdots mid]$ and that of \max of $A[mid + 1 \cdots high]$)

How to describe the design of an algorithm in English

- ▶ First point out what strategy/method used, e.g. divide-and-conquer, binary search.
- ▶ Separate paragraphs if necessary, e.g. branches.
- ▶ Bullet-point format is also a good practice.

Stock Investment: Pseudocode ¹

STOCK-INVESTMENT(*A*, *low*, *high*)

```
1  // Base case: only one element
2  if low == high
3      return low, high, 0
4  else mid = low +  $\lfloor (\textit{high} - \textit{low}) / 2 \rfloor$ 
5      leftBuy, leftSell, leftGain = STOCK-INVESTMENT(A, low, mid)
6      rightBuy, rightSell, rightGain = STOCK-INVESTMENT(A, mid, high)
7      // Find the index of min(leftArray)
8      crossBuy = MIN-INDEX(A, LOW, MID)
9      // Find the index of max(rightArray)
10     crossSell = MAX-INDEX(A, MID, HIGH)
11     crossGain = A[crossSell] - A[crossBuy]
12     if max(leftGain, rightGain, crossGain) < 0
13         return "no gain"
14     elseif leftGain ≥ rightGain and leftGain ≥ crossGain
15         return leftBuy, leftSell, leftGain
16     elseif rightGain ≥ leftGain and rightGain ≥ crossGain
17         return rightBuy, rightSell, rightGain
18     else return crossBuy, crossSell, crossGain
```

¹See pp.[20-22] in the textbook for pseudocode conventions

Stock Investment: Time complexity

$$\begin{aligned}T(n) &= 2T\left(\frac{n}{2}\right) + \Theta(n) + \Theta(1) \\&= \Theta(n \log n)\end{aligned}$$

1. $2T(\frac{n}{2})$: the first two optimal solutions
2. $\Theta(n)$: the third solution is equivalent to finding the `min` and `max`, e.g. linear scan
3. $\Theta(1)$: compare the results of three solutions