

10/1/22

Group

In Group Theory,  
 (e)  $\rightarrow$  Identity element  
 Date \_\_\_\_\_  
 Page \_\_\_\_\_

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Date \_\_\_\_\_  
 Page \_\_\_\_\_

Algebraic System :-

It is a set with a binary operation satisfying some number of properties.

$$G = \{1, -1, i, -i\}$$

# Abelian / Commutative group.

$$a \times b = b \times a$$

$$\textcircled{Q} \quad G = \{1, -1, i, -i\}$$

Show that it is a commutative group under the binary operation multiplication.

Closure

$$1 \times 1 = 1 \in G$$

$$(-1) \times (-1)$$

$$1 \times -1$$

Associativity

$$\begin{aligned} x &= 1, y = i, z = -i \\ x(y \cdot z) &= 1(i(-i)) \\ &= 1(1) \\ &= 1 \end{aligned}$$

$$\begin{aligned} (x \cdot y)z &= (1i)(-1) \\ &= i(-1) \\ &= -i^2 \\ &= -(-1) \\ &= 1 \end{aligned}$$

Identity

$$e = 1 \in G$$

$$\begin{aligned} 1 \cdot (1) &= 1; -1 \cdot 1 = -1; i \cdot 1 = i \\ -i \cdot 1 &= -i \end{aligned}$$

Inverse

$$\text{An } (1) = 1$$

$$1 \cdot 1 = 1$$

$$\text{An } (-1) = -1$$

$$-1 \times -1 = 1$$

$$\text{An } (i) = .$$

$\rightarrow (\mathbb{Z}, +)$  is a commutative group

$\rightarrow$

$\rightarrow$

\textcircled{Q}

Show that the operation add modulo 10 is a group over the set {0, 1, 2, 3, ..., 9}

~~\* = multiplication/addition modulo 10~~

Ary

$$a \oplus b = \begin{cases} a + b, & a+b \leq 10 \\ (a+b)-10, & a+b > 10 \end{cases}$$

Closure

$$0 \oplus 9 = 9 \in G$$

$$8 \oplus 9 = 7 \in G$$

Associative

$$x=2, y=5, z=9$$

$$\begin{aligned}(x+y)+z &= (2 \oplus 5) \oplus 9 \\ &= 7 \oplus 9 \\ &= 6\end{aligned}$$

$$\begin{aligned}x+(y+z) &= 2 \oplus (5 \oplus 9) \\ &= 2 \oplus 4 \\ &= 6\end{aligned}$$

Identity  
0

$$1 \oplus 0 = 1$$

$$2 \oplus 0 = 2$$

$$0 \oplus 9$$

Inverse

$$g_{nv} \cdot 0 = 10$$

$$g_{nv} \cdot 1 = 9 \quad \left| \begin{array}{l} 1 \oplus 9 = 1 \\ 2 \oplus 8 = 1 \end{array} \right.$$

$$g_n \cdot 2 = 8 \quad \left| \begin{array}{l} 2 \oplus 8 = 1 \\ 3 \oplus 7 = 1 \end{array} \right.$$

$$g_n \cdot 3 = 7 \quad \left| \begin{array}{l} 3 \oplus 7 = 1 \\ \dots \end{array} \right.$$

Commutative

$$a \oplus b = b \oplus a$$

- Q1 Prove that the Identity element of a group is unique.  
 Q2 Prove that the Inverse of an element of a group is unique.

Ans(1)

Let  $(G, *)$   
 $(e, f)$  are identity element

$$ae = ea = a \quad \text{--- (i)}$$

$$af = f \cdot e = a \quad \text{--- (ii)}$$

$$ae = af$$

~~$a^{-1}(ae) = a^{-1}f$~~

$$\Rightarrow e = f \quad (\text{Left cancellation Law})$$

So, it has only one identity element.

Ans(2)

Let  $(G, *)$   
 $a \in G$

$$g(a) = b \quad c$$

$$a \cdot b = ba = e \quad \text{--- (i)}$$

$$ac = ca = e \quad \text{--- (ii)}$$

$$ab = ac$$

$$\boxed{b = c}$$

NOTE

Let  $(G, *)$  be a group and  $H$  is a subset of  $G$  ( $H \subseteq G$ ) and  $H$  is a finite set. Then  $(H, *)$  is a subgroup of  $(G, *)$  if  $H$  satisfies closure property on  $*$ .  
 (Important Theorem)

Ans

$(H, *) \longrightarrow$

|                                |
|--------------------------------|
| $\text{① } H \text{ C.G}$      |
| $\text{② } (H, *) \rightarrow$ |
| $\text{③ closure}$             |
| $\text{④ associative}$         |
| $\text{⑤ Identity}$            |
| $\text{⑥ inverse}$             |

11/22

$H_1, H_2$  are sub-groups of  $G$   
 then  $H_1 \cap H_2$  is also a subgroup of  $G$

$G \rightarrow$  group  
 $H_1, H_2$  are subgroup of  $G$

$H_1 \cap H_2$  is a subgroup of  $G$

$H_1 CG, H_2 CG$

$H_1 \cap H_2 \subset H_1, CG$

$H_1 \cap H_2 \subset CG$

$H_1 \cap H_2$

$a, b \in H_1 \cap H_2 \Rightarrow a, b \in H_1, H_2$

$a \in H_1, H_2 \Rightarrow a \in H_1, a \in H_2$

$b \in H_1, H_2 \Rightarrow b \in H_1, b \in H_2$

$a, b \in H_1 \Rightarrow ab \in H_1$

$a b \in H_2$

$a b \in H_1, H_2$

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 Page \_\_\_\_\_

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Date \_\_\_\_\_  
 Page \_\_\_\_\_

# Cosets and Lagrange's theorem

Cosets

Let  $G$  be any group,  $H$  is a subgroup of  $G$

$$H = \{h_1, h_2, h_3, \dots, h_m\}$$

$$a \in G, a \notin H$$

$$Ha$$

→ Left coset

→ Right coset

$$aH = \{ah_1, ah_2, ah_3, ah_4, \dots, ah_m\}$$

$$Ha = \{h_1a, h_2a, h_3a, h_4a, \dots, h_ma\}$$

order ( $H$ ) =  $k$  then order ( $aH$ ) =  $k$

Theorem

# Let  $G$  be any group,  $H$  is a subgroup of  $G$ . Let  $a, b$  be any 2 elements of  $G$ , then  $aH$  &  $bH$  are 2 cosets of  $H$  f  $G$ . Either  $aH = bH$  or  $aH \cap bH = \emptyset$

Contradiction to proof it

→ Let us assume that  $aH \neq bH$  have some common elements.

Let

$$H = \{h_1, h_2, h_3, \dots, h_m\}$$

$$aH = \{ah_1, ah_2, ah_3, \dots, ah_m\}$$

$$bH = \{bh_1, bh_2, bh_3, \dots, bh_m\}$$

Let  $ah_1 = bh_2$

$$\Rightarrow a = bh_2^{-1}h_1 \quad \text{--- (1)}$$

Let  $x \in aH$

$$\Rightarrow x = ah_3 \quad (h_3 \in H)$$
$$\Rightarrow x = bh_2h_1^{-1}h_3 \quad \text{--- (from eq ①)}$$

$$h_1 \in H$$
$$h_1^{-1} \in H$$
$$\Rightarrow x = b h_1 h_1^{-1} h_3$$

$$\text{So, } h_2 h_1^{-1} h_3 = h_4 \in H$$

$$x = bh_4$$

$$bh_4 \in bH$$

$$x \in bH$$

$$\forall x \in aH \Rightarrow x \in bH$$

$$[aH \subset bH] \quad \text{--- ①}$$

$$x \in bH$$

$$x = bh_3$$

$$x = ah_1 h_2^{-1} h_3$$

$$x = ah_1 \in aH$$

$$\forall x \in bH \Rightarrow x \in aH$$

$$[bH \subset aH] \quad \text{--- ②}$$

from ① + ②

$$aH = bH$$

### Lagrange's Theorem

The order of any finite group divides the order of any finite group.

$O(G) = H$  is a subgroup of  $G$

$$O(H) \mid O(G)$$

$O \rightarrow \text{order}$

$$O(G) = n$$

$$G = \{g_1, g_2, g_3, \dots, g_n\}$$

$$O(H) = m$$

$$H = \{h_1, h_2, \dots, h_m\}$$

Left coset of  $H$  in  $G$

$$g_1 H, g_2 H, g_3 H, g_4 H, \dots, g_n H$$

Let the number of distinct cosets be  $k$

$$\begin{array}{|c|} \hline G \\ \hline aH & bH & cH & \dots & g_n H \\ \hline \end{array}$$

$$D = \{g_1 H, g_2 H, \dots, g_n H\}$$

$$O(H) = m$$

$$O(mH) = m$$

$$O(D) = O(g_1 H) + O(g_2 H) + \dots + O(g_n H)$$

$$g_1 H \cup g_2 H \cup g_3 H \cup \dots \cup g_n H = G$$

$$O(G) = O(g_1 H) + O(g_2 H) + \dots + O(g_n H)$$

$$n = m + m + \dots + m \quad (\text{k times})$$

$$n = mk$$

$$n = mk$$

$$m/n$$

$$\mathcal{O}(H) \mid \mathcal{O}(G)$$

12/01/22

# Normal ~~sub-group~~  $\rightarrow$  (NSG) $G$  be any group $N$  is a subgroup of  $G$ 

$$\forall g \in G \quad gNg^{-1} \subseteq N \quad n \in N$$

# Theorem 1

Let  $G$  be any group $N$  is a subgroup of  $G$  $N$  is a  $N$ -subgroup of  $G$  iff

$$xNx^{-1} = N \quad \forall x \in G$$

Proof

$$N = \{n_1, n_2, \dots, n_k\}$$

$$xNx^{-1} = \{xn_1x^{-1}, xn_2x^{-1}, xn_3x^{-1}, \dots\}$$

① Let  $N$  is a  $N$ -subgroup of  $G$ 

$$xNx^{-1} = N$$

$$xNx^{-1} \subseteq N \quad \& \quad N \subseteq xNx^{-1}$$

$$y \in xNx^{-1}$$

$$\Rightarrow y = xnx^{-1} \quad n \in N$$

 $N$  is ~~NSG~~ of  $G$ 

$$xnx^{-1} \in N$$

$$y \in N$$

$$[xNx^{-1} \subseteq N]$$

$$\begin{matrix} x \\ \times \\ n \\ \times \\ x^{-1} \\ \hline \end{matrix} \in N$$

$$n^{-1}n(n^{-1})^{-1}$$

$$n(n^{-1}nn)$$

$$\cancel{n} = nn^{-1}n$$

$$(nx)n^{-1} \in xNx^{-1}$$

$$n \in xNx^{-1}$$

$$[N \subseteq xNx^{-1}]$$

Theorem 2

Let  $G$  be any group $N$  is a normal sub-group of  $G$  $N$  is a NSG iff  $xH = Hx$ Proof $N$  is a NSG of  $G$ 

$$nN = Nx$$

 $N$  is a NSG iff  $xNx^{-1} = N$ 

$$xNn^{-1} = N$$

$$(xNn^{-1})x = Nx$$

$$xNx^{-1}x = Nx$$

$$xNe = Nx$$

$$\underline{xN = Nx}$$

 $(e = \text{Identity})$ Let,  $Nx = xN$  $N$  is a NSG of  $G$ 

$$xN = Nx$$

$$xNx^{-1} = (Nx)x^{-1}$$

$$xNx^{-1} = Nx^{-1}x$$

$$xNx^{-1} = Ne$$

$$\underline{xNx^{-1} = N}$$

 $N$  is a NSG of  $G$ .

## # Theorem 3

 $G$  is any group $N$  is a NSG of  $G$ Then  $N$  is a NSG iff  $NxNy = Nx y$ Proof $N$  is NSG of  $G$ 

$$NxNy = Nx y$$

$$\stackrel{LHS}{=} (Nx)(Ny)$$

$$= N(xN)y$$

$$= N(Nx)y \quad (Nx = xN)$$

$$= NNxy.$$

$$= N^2xy$$

$$= Nx y$$

$$\text{Let, } NxNy = Nx y$$

 $\# N$  is a NSG of  $G$ 

$$\Rightarrow NxN x^{-1} = Nx x^{-1}$$

$$\Rightarrow NxN x^{-1} = Ne$$

$$\Rightarrow NxN x^{-1} = N$$

$$\Rightarrow n_1 n_2 x^{-1} \in N$$

$$\Rightarrow n_1^{-1}(n_1 n_2 x^{-1}) \in n_1^{-1}N$$

$$\Rightarrow n_1^{-1}n_1 n_2 x^{-1} \in n_1^{-1}N$$

$$\Rightarrow x n_2 x^{-1} \in n_1^{-1}N$$

$$\Rightarrow x n_2 x^{-1} \in N$$

$$\Rightarrow n_2 x^{-1} \in N$$

 $N$  is a NSG of  $G$ .

## # Theorem 4

Every subgroup of abelian group is normal SG

 $G \rightarrow$  abelian group $N \rightarrow$  SG of  $G$  $N$  is a NSG of  $G$ 

$$n \in N, x n x^{-1} \in N$$

$$= n(nx^{-1})$$

$$= x(n^{-1})n$$

$$= (nx^{-1})n \quad (\text{Ab.})$$

$$= xn$$

## # Theorem 5

 $G$  is any grp. $N_1, N_2$  are 2 NSG of  $G$  $N_1 \cap N_2$  is also a NSG of  $G$ . $N_1 \cap N_2$  is also a SG of  $G$ Proof

$$\forall n \in N_1 \cap N_2, \exists n \in G$$

$$\Rightarrow nx^{-1} \in N_1 \cap N_2$$

$$n \in N_1 \cap N_2$$

$$n \in N_1 \text{ & } n \in N_2$$

$$nnx^{-1} \in N_1, \quad nnx^{-1} \in N_2$$

$$nnx^{-1} \in N_1 \cap N_2$$

We have to prove it

(1)  $N_1 \cap N_2$  is a SG of  $G$ (2)  $N_1 \cap N_2$  satisfies Normal property

## # Theorem ①

Let  $G$  be any grp

$H$  is a SG of  $G$

$N$  is a NSG of  $G$

$H \cap N$  is a NSG of  $H$

Proof

To prove:  $\vdash$

①  $H \cap N$  is a SG of  $H$

②  $H \cap N$  satisfies Normal property

$H \cap N$  is a SG of  $G$

$H \cap N$  is itself a group

$H \cap N \subset H$

$H \cap N$  is a SG of  $H$

$\forall \alpha \in H \cap N, \exists h \in H$   
 $h \alpha h^{-1} \in H \cap N$

$\forall \alpha \in H \cap N, \exists h \in H$

$h \alpha h^{-1} \in H$

$\forall \alpha \in H \cap N, \exists h \in H$

$h \alpha h^{-1} \in H$

$h \alpha h^{-1} \in H \cap N$

## # Homomorphism

$(G, *)$        $(G', \cdot)$

$\phi: G \rightarrow G'$

~~$\phi(a * b) = \phi(a) * \phi(b)$~~

$$\phi(a * b) = \phi(a) \cdot \phi(b)$$

(1)  $\forall \alpha \in G \cap N, \exists h \in H$

$h \alpha h^{-1} \in H \cap N$

$\alpha \in H \cap N$

$h \alpha h^{-1} \in N$  (By normal property)

$h \in H, h^{-1} \in H, \alpha \in H$

$h \alpha h^{-1} \in H$

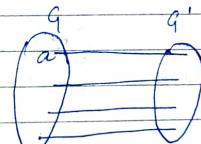
$h \alpha h^{-1} \in H \cap N$

## # Homomorphism

Let  $(G, *)$        $(G', \cdot)$

$\phi: G \rightarrow G'$

$$\phi(a * b) = \phi(a) \cdot \phi(b)$$



In Let  $(G, *)$  &  $(G', \cdot)$  be two groups over respective operations

$\phi: G \rightarrow G'$

$$\forall x \in G \Rightarrow \phi(x) = 1$$

$$\phi(x * y) = \phi(x) \cdot \phi(y)$$

$$\begin{aligned} \phi(x) &= \log x & \phi(y) &= \log y \\ \phi(x \cdot y) &= \log(x \cdot y) & & \\ &= \log x + \log y & & \\ &= \phi(x) + \phi(y) & & \end{aligned}$$

Isomorphism:

$$\textcircled{Q} \quad \phi: G \rightarrow G'$$

$e$  is the identity element of  $G$ .

s.t.

$$(i) \phi(e) = e'$$

$$(ii) \forall a \in G \Rightarrow \phi(a^{-1}) = (\phi(a))^{-1}$$

$$\rightarrow (i) ae = a$$

$$\Rightarrow \phi(ae) = \phi(a)$$

$$\Rightarrow \phi(a)\phi(e) = \phi(a) \quad (\because \phi \text{ is a homomor.})$$

$$\Rightarrow \phi(a)\phi(e) = \phi(a)e' \quad (\phi(a)e' = \phi(a))$$

$$\Rightarrow \phi(e) = e'$$

$$(ii) e' = \phi(e)$$

$$\Rightarrow e' = \phi(a \cdot a^{-1}) \quad (\because a \cdot a^{-1} = e)$$

$$\Rightarrow e' = \phi(a)\phi(a^{-1}) \quad (\because \phi \text{ is a homomor.})$$

$$\Rightarrow \text{Inv. of } \phi(a) \text{ is } \phi(a^{-1})$$

$$\Rightarrow (\phi(a))^{-1} = \phi(a^{-1})$$

## # Kernel of a homomorphism

$$\phi: G \rightarrow G'$$

If we take these elements as a set ( $K\phi$ )

$$n \in K\phi \Rightarrow \phi(n) = e'$$

If  $\phi$  is a homomorphism from  $G \rightarrow G'$   
 $K$  is a NSG of  $G$ .

- (i)  $K$  is a SG of  $G$
- (ii)  $K$  satisfies Normal Property  
i.e.  $gKg^{-1} \subseteq K$

(a)  $K \subseteq G$ , closure

$$\forall a, b \in K \Rightarrow ab \in K$$

$$a \in K \Rightarrow \phi(a) = e'$$

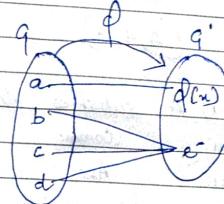
$$b \in K \Rightarrow \phi(b) = e'$$

$$\begin{aligned} \phi(ab) &= \phi(a)\phi(b) \quad (\phi \text{ is a mono. mor.}) \\ &= e'e' \\ &= e' \end{aligned}$$

$$\Rightarrow ab \in K$$

$$\begin{aligned} (iii) \phi(gkg^{-1}) &= \phi(g)\phi(k)\phi(g^{-1}) = \phi(g)e'\phi(g^{-1}) \\ &= \phi(g)\phi(g^{-1}) = \phi(gg^{-1}) \quad (\text{mono. mor.}) \\ &= \phi(e) \end{aligned}$$

$$\begin{aligned} \phi(gkg^{-1}) &= e' \\ gkg^{-1} &\in K \end{aligned}$$



Ring:

$(A, +, \cdot)$

(i)  $A$  satisfies closure +

comm.

(ii) Ass.

(iii) Add. Identity

(iv) Add. Inv.

(v) Closure \*

(vi) Associative \*

(vii)  $a.(b+c) = ab+ac$

Q  $(z, +, \cdot)$        $(R, +, \cdot)$   
Ring                  Ring

$(N, +, \cdot)$   
Doesn't satisfy  
Identity & Inverse  
not satisfied.

# Commutative Ring / Ring with unit element

$(R, +, \cdot)$   
 $a, b \in R \Rightarrow ab = ba$

} Commutative  
Ring

$(R, +, \cdot)$   
 $\forall n \in R \Rightarrow ne = n$   
(multiplicative ring)

} Ring with  
unit elements.

## Zero Divisors of a Ring

Let  $(A, +, \cdot)$  be a Ring

$x \neq 0 \text{ & } A \text{ f } y \in A, y \neq 0$ , then  $y$  is 0 divisor

of  $x$  if  $xy = 0$  or  $yx = 0$ .

$$p) \quad n = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \quad y = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$n \neq 0$

$y \neq 0$

$$ny = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

$$= \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix} = 0$$

## Field

$(A, +, \cdot)$

|        |                               |
|--------|-------------------------------|
| $+$    | cl., Comm., Ass., Iden., Inv. |
| $*$    |                               |
| $+, *$ | Distributive                  |

## Integral Domain

Suppose  $(A, +, \cdot)$  is a ring. This ring is called integral domain if it satisfies:

- i)  $ab = ba$  (commutative prop. for mult.).
- ii)  $1 \in A$  such that  $x \cdot 1 = 1 \cdot x = x$  (Iden. mul.)
- iii)  $A$  has no zero divisor.

OR)

A commutative ring with unit element is called an integral domain if it has no zero divisor.

Q.  $(\mathbb{Z}, +, \cdot) \rightarrow$  Ring

- ↳ comm. prop. for mult.
- ↳ multiplicative identity prop.
- ↳ has no zero divisor.

$\therefore (\mathbb{Z}, +, \cdot)$  is an integral domain.

Thm 1:

Every field is an integral domain.

Ex:  $(D, +, \cdot)$  is a field.  $\begin{matrix} \nearrow + \text{ prop.} \\ \searrow \times \text{ "} \\ \text{Distributive} \end{matrix}$

Let  $a, b \in D$ .

$$ab = 0 \quad \text{--- (1).}$$

(i) Suppose,  $a \neq 0 \neq ab = 0$ .

$a \in D \Rightarrow a^{-1} \text{ exists } \neq a^{-1} \neq 0$ .

$$ab = 0$$

$$\Rightarrow a^{-1}(ab) = a^{-1}0$$

$$\Rightarrow (a^{-1}a)b = a^{-1}0$$

$$\Rightarrow 1b = a^{-1}0 \quad (\because aa^{-1} = I)$$

$$\Rightarrow b = 0 \quad (\text{Identity prop.})$$

(ii) — ,  $b \neq 0 \neq ab = 0$

$b \in D, \Rightarrow b^{-1} \text{ exists } \neq b^{-1} \neq 0$

$$ab = 0$$

$$\Rightarrow (ab)b^{-1} = 0b^{-1}$$

$$\Rightarrow a(bb^{-1}) = 0b^{-1}$$

$$\Rightarrow a1 = 0b^{-1}$$

$$a = 0$$

i.e.  $\begin{cases} \text{for } ab = 0 \\ \text{any one out of } (a \text{ or } b) \end{cases}$  must be 0.

$\therefore D$  has no zero divisor.

$\Rightarrow D$  is an integral domain.

Thm 2:

Every finite integral domain is a field.

Let  $D$  is a finite integral domain.

Show that:  $D$  is a field.

Suppose:  $D = \{x_1, x_2, \dots, x_n\}$ .

Let  $x$  is any element of  $D$ .

$$xD = \{xx_1, xx_2, \dots, xx_n\} \in D$$

(closure prop)

If possible, let  $xx_i = xx_j$

$$\Rightarrow x(x_i - x_j) = 0$$

either  $x=0$  or  $x_i - x_j = 0$

$\therefore D$  has no zero divisor

$\therefore$  Product of two non-zero elements  $\neq 0$ .

$$\therefore x_i = x_j$$

Let  $xx_i = x$  ( $\because$  every element of  $xD =$  any 1 ele of  $D$ ?)  
 $\nexists x_i$  is identity element.

$$x_i = 1$$

Since,  $x_i = 1$

so, out of  $x_1, x_2, \dots, x_n$  one element must be 1.

$$\text{Let } xx_j = 1$$

$$\Rightarrow x_j = x^{-1}$$

so,  $\forall x \in D \quad \exists x^{-1} \in D$

$$\text{where } xx^{-1} = x^{-1}x = 1$$

Hence,  $D$  is a field.

Q. Let  $(R, +, \cdot)$  be a Ring Field  
 $a, b \in R$ .

$$(i) a \cdot 0 = 0 \cdot a = 0$$

$$(ii) a(-b) = (-a)b = -ab$$

→ (i) We know that, since  $R$  is a ring  
it must satisfy identity prop.

$$\therefore a + 0 = a \quad (\text{Ident. add.})$$

$$\Rightarrow (a + 0) \cdot 0 = a \cdot 0 \quad (\text{mult. Ident.})$$

$$\Rightarrow a \cdot 0 + 0 \cdot 0 = a \cdot 0 \quad (\text{Dist.})$$

$$\Rightarrow a \cdot 0 + 0 = a \cdot 0$$

$$\Rightarrow a \cdot 0 = a \cdot 0$$

$$x + 0 = x \quad (\text{Add. Iden.})$$

$$\text{Prf. } x = 0$$

$$\Rightarrow 0 + 0 = 0$$

$$\Rightarrow a(0+0) = a \cdot 0 \quad (\text{Factor.})$$

$$\Rightarrow a \cdot 0 + a \cdot 0 = a \cdot 0 \quad (\text{Dist.})$$

$$\Rightarrow a \cdot 0 + a \cdot 0 = a \cdot 0 + 0 \quad (\text{Ident.})$$

$$\Rightarrow a \cdot 0 = 0 \quad \text{---(i)} \quad (\text{L.C. law})$$

$$x + 0 = x$$

$$\Rightarrow 0 + 0 = 0 \quad (x = 0)$$

$$\Rightarrow (0+0)a = 0 \cdot a \quad (\text{Factor.})$$

$$\Rightarrow 0a + 0a = 0a + 0 \quad (\text{Dist.})$$

$$\Rightarrow 0a + 0a = 0a + 0 \quad (\text{Add. ident.})$$

$$\Rightarrow 0a = 0 \quad \text{---(ii)} \quad (\text{L.C. law})$$

from (i) & (ii):

$$a \cdot 0 = 0 \cdot a = 0$$

$$(ii) \quad x + (-x) = 0 \quad - (i) \quad (\text{Add. Inv.})$$

$$\Rightarrow b + (-b) = 0 \quad (x = -b)$$

$$\Rightarrow a(b + (-b)) = a \cdot 0 \quad (\text{clscer})$$

$$\Rightarrow ab + a(-b) = a \cdot 0 \quad (\text{Distr.})$$

$$\Rightarrow ab + a(-b) = 0 \quad (\because a \cdot 0 = 0)$$

$\Rightarrow ab$  is additive inv. of  $a(-b)$

$$\therefore a(-b) = -ab. \quad (ii)$$

putting  $n=0$  in eq (i).

$$\Rightarrow a + (-a) = 0 \quad (\text{Add. Inv.})$$

$$\Rightarrow (a + (-a))b = 0b \quad (\text{clscer})$$

$$\Rightarrow ab + (-a)b = 0b \quad (\text{Distr.})$$

$$\Rightarrow ab + (-a)b = 0 \quad (0 \cdot a = 0).$$

$\Rightarrow ab$  is add. inv. of  $(-a)b$ .

$$\therefore (-a)b = -ab. \quad — (iii)$$

from (ii) & (iii) :

$$a(-b) = (-a)b = -ab$$

Codes & Group Codes

- Word is a sequence of letters
- Code is a collection of words used to represent diff. messages.
- Any word in a code is called a codeword.
- Block code is a code consisting of equal length words.
- Binary alphabet : Generally in a code only these are used i.e. 0 & 1.

Binary Operation ( $\oplus$ ).

$$\begin{aligned}x &= 101100 \\y &= \underline{101001} \\x \oplus y &= \underline{\underline{000101}}\end{aligned}$$

| x | y | $x \oplus y$ |
|---|---|--------------|
| 1 | 1 | 0            |
| 1 | 0 | 1            |
| 0 | 1 | 1            |
| 0 | 0 | 0            |

$A = \{x, y, z, \dots\}$  all the binary seq<sup>n</sup> of length  $n$ .  
 $(A, \oplus)$  is a grp if :

- (i)  $x \oplus y \in A$  closer
- (ii)  $(x \oplus y) \oplus z = x \oplus (y \oplus z)$  Ass.
- (iii)  $x \oplus e = x$  ( $e = 00000$ ). Idem.
- (iv)  $x \oplus (-x) = 0$  Inv.

— weight of a word : no. of  $\perp$  present in  
( $w$ ) the word.

$$x = 10101 \quad wt = 3$$

$$y = 01111 \quad wt = 4.$$

→ Distance ( $d(x, y)$ ):  
weight.

$$d(x, y) = w(x \oplus y)$$

$$= w(11010)$$

$$= 3.$$

In words: Distance b/w  $x$  &  $y$  is no. of pos's where  $x \neq y$  differs.

Properties:

$$(i) \quad d(x, y) = d(y, x)$$

$$(ii) \quad d(x, y) \leq d(x, z) + d(z, y).$$

$$\begin{aligned} (iii) \quad \text{LHS: } d(x, y) &= w(x \oplus y) \\ &= w(x \oplus 0 \oplus y) \\ &= w(x \oplus z \oplus z \oplus y) \\ &\leq w(x \oplus z) + w(z \oplus y) \\ &\leq d(x, z) + d(z, y) \end{aligned}$$

## Distance of a block code

Let  $G$  be any block code. The distance of  $G$  is the min. distance b/w any pair of distinct code words of  $G$ .

## Max<sup>m</sup> likelihood decoding criterion

Let  $x_1, x_2, x_3, \dots, x_n$  denotes the code words in a block code  $G$  &  $y$  is the received word.

Let us find out  $P(x_i|y)$   $i=1, 2, \dots, n$  where  $P(x_i|y)$  is the probab. that  $x_i$  is the transmitted word under the cond' that  $y$  is the received word.

$$P(x_i|y) = \frac{P(x_i, y)}{P(y)}$$

If  $P(x_k|y)$  is the largest conditional probab. among  $P(x_1|y), P(x_2|y), \dots, P(x_n|y)$  then we conclude that  $x_k$  was the received word.

## Min<sup>m</sup> distance decoding criterion

Similarly, if  $d(x_k, y)$  is min<sup>m</sup> then  $x_k$  is considered as the transmitted word. This criterion is called min<sup>m</sup> d— criterion.

NOTE: When min<sup>m</sup> distance decoding criterion is followed then a code of distance  $2t+1$  can correct  $t$  or fewer transmission errors.

Pf: Let  $x$  is a codeword which is transmitted &  $y$  is the receiving word of the error  $\leq t$ . So  $d(x, y) \leq t$

Let  $x_1$  be another codeword

But,  $d(x, x_1) \geq 2t+1$  ( $\because$  Given).

Now,

$$d(x, x_1) \leq d(x, y) + d(y, x_1)$$

$$\Rightarrow 2t+1 \leq t + d(y, x_1)$$

$$\Rightarrow d(y, x_1) \geq t+1$$

Since, the min<sup>m</sup> dist. occurs in case of  $x$ , so  $x$  is the transmitted word.

## Group code :

We know that  $(A, \oplus)$  is a group.  
 $(G, \oplus)$  is called a group code if  
 $(G, \oplus)$  is a subgrp of  $(A, \oplus)$  where  
A is a set of all binary sequences  
of length n.