# **Chosen Ciphertext Attack On RSA**

#### Aim:

To demonstrate a Chosen Ciphertext Attack (CCA) on an RSA-based encryption scheme by manipulating an encrypted message and using a decryption oracle to retrieve the original plaintext.

# **Description:**

This program simulates RSA encryption and decryption and showcases how an attacker can exploit a decryption oracle to recover the original message. It performs the following steps:

- 1. Generates an RSA key pair (public and private keys).
- 2. Encrypts a plaintext message using the public key.
- 3. Simulates a Chosen Ciphertext Attack by modifying the ciphertext with a random multiplier and querying the decryption oracle.
- 4. Uses the decrypted modified message to recover the original plaintext.

#### Code:

```
import java.math.BigInteger;
import java.security.SecureRandom;
public class CCAttack {
  // Generate a large prime number
  private static BigInteger generatePrime(int bitLength) {
     SecureRandom rand = new SecureRandom();
     return BigInteger.probablePrime(bitLength, rand);
  }
  // Generate RSA keys
  public static BigInteger[] generateKeyPair(int bitLength) {
     BigInteger p = generatePrime(bitLength / 2);
     BigInteger q = generatePrime(bitLength / 2);
     BigInteger n = p.multiply(q);
     BigInteger phi = (p.subtract(BigInteger.ONE)).multiply(q.subtract(BigInteger.ONE));
    // Choose public exponent e (commonly 65537)
    BigInteger e = BigInteger.valueOf(65537);
    // Ensure e and phi(n) are coprime (\gcd(e, phi(n)) = 1)
     while (e.gcd(phi).compareTo(BigInteger.ONE) != 0) {
       e = e.add(BigInteger.TWO); // Try next odd number
     }
    // Calculate the private exponent d such that e * d \equiv 1 \pmod{\varphi(n)}
```

```
BigInteger d = e.modInverse(phi);
  // Public key is (e, n), Private key is (d, n)
  return new BigInteger[] { e, n, d };
}
// Encrypt a message using the public key (e, n)
public static BigInteger encrypt(BigInteger message, BigInteger e, BigInteger n) {
  return message.modPow(e, n);
}
// Decrypt a message using the private key (d, n)
public static BigInteger decrypt(BigInteger cipherText, BigInteger d, BigInteger n) {
  return cipherText.modPow(d, n);
}
public static void main(String[] args) {
  int bitLength = 512; // Key size in bits
  // Step 1: Generate RSA keys
  BigInteger[] keys = generateKeyPair(bitLength);
  BigInteger e = keys[0]; // Public exponent
  BigInteger n = keys[1]; // Modulus
  BigInteger d = keys[2]; // Private exponent
  System.out.println("Public Key (e, n): (" + e + ", " + n + ")");
  System.out.println("Private Key (d, n): (" + d + ", " + n + ")");
  // Step 2: Encrypt a secret message
  String message = "hello world!";
  BigInteger messageBigInt = new BigInteger(1, message.getBytes());
  System.out.println("Original Message: " + message);
  BigInteger cipherText = encrypt(messageBigInt, e, n);
  System.out.println("Encrypted Message: " + cipherText);
  // Step 3: Simulating a Chosen Ciphertext Attack
  System.out.println("\n=== Chosen Ciphertext Attack (CCA) Simulation ===");
  // The attacker chooses a random number r (mod n)
  BigInteger r = new BigInteger(bitLength / 2, new SecureRandom()).mod(n);
  System.out.println("Attacker chosen r: " + r);
  // Attacker creates a modified ciphertext: C' = C * r^e mod n
  BigInteger modifiedCipherText = cipherText.multiply(r.modPow(e, n)).mod(n);
  System.out.println("Modified Ciphertext (C'): " + modifiedCipherText);
```

```
// The attacker queries the decryption oracle to get the decrypted C'
BigInteger decryptedModifiedMessage = decrypt(modifiedCipherText, d, n);
System.out.println("Decryption Oracle Output (M'): " + decryptedModifiedMessage);

// The attacker recovers the original message using: M = M' / r mod n
BigInteger recoveredMessage

decryptedModifiedMessage.multiply(r.modInverse(n)).mod(n);
String recoveredText = new String(recoveredMessage.toByteArray());
System.out.println("\nRecovered Message by Attacker: " + recoveredText);
}
```

# **Output:**

```
Public Key (e, n): (65537, 9315505000465798780100184572299921255152332751612956029377255887624257454991491582266914488755807905719511772485512656245195643284242886232285807840170569)
Private Key (d, n): (9010612219066153108444251345143446911006251091600904963826281417035826000378540271873513624550491982420426704937929792754844973065339969165056935440676993, 93155
050006457987801001845722999212255152332751612956029377255887624257454991491582266914488755807905719511772485512656245195643284242886232285807840170569)
Original Message: hello world!
Encrypted Message: 3962560270177850592499119125028567194097752129062890659218668041571501839805861714627170476884235211604898046083468235900040440313483093597441311680943041

=== Chosen Ciphertext Attack (CCA) Simulation ===

Attacker chosen r: 102526967538661770521239333779796428193845186161762361697474690447778571926591

Modified Ciphertext (C'): 177702038628703520037766724950641582744823353845149723029493681816490205407888126106800239103620444762033812024784666263385903485764230529560528681371035

Decryption Oracle Output (M'): 3312549386543861656559196559576759021664018368057955716021529525833136743810939061346351383169065446596639

Recovered Message by Attacker: hello world!
```

# Code analysis:

#### 1. Key Generation:

- Two large prime numbers p and q are generated to compute n = p \* q.
- Euler's totient function phi(n) = (p-1) \* (q-1) is calculated.
- A common public exponent e = 65537 is chosen, ensuring it's coprime with phi(n).
- The private exponent d is computed as  $d = e^{(-1)} \mod phi(n)$ .

# 2. Encryption:

- Converts the plaintext message to a BigInteger.
- Encrypts the message using cipherText = message^e mod n.

# 3. **Decryption:**

Decrypts ciphertext using message = cipherText^d mod n.

### 4. Chosen Ciphertext Attack (CCA) Simulation:

- Attacker selects a random integer r.
- Modifies the original ciphertext  $C' = C * r \le mod n$ .
- Queries the decryption oracle for  $M' = C' \wedge d \mod n$ .
- Recovers the original message using  $M = M' * r^{(-1)} \mod n$ .

# Time complexity:

- **Key Generation:** O(n^2 log n log log n) (Probable prime generation and modular inverse calculations)
- **Encryption:** O(log e \* log n) (Modular exponentiation using square-and-multiply algorithm)
- **Decryption:** O(log d \* log n) (Modular exponentiation)
- **CCA Attack:** O(log e \* log n) + O(log n) (Ciphertext modification, modular exponentiation, and modular inverse)

# **Space complexity:**

- **Key Storage:** O(log n) for p, q, n, e, d
- **Encryption/Decryption:** O(log n) for storing messages and ciphertexts
- CCA Attack: O(log n) for storing intermediate values like r, C', and M'