CARNEGIE MELLON UNIVERSITY
DEPARTMENT OF COMPUTER SCIENCE
15-445/645 – DATABASE SYSTEMS (FALL 2019)
PROF. ANDY PAVLO

Homework 2 (by Amadou Ngom)

Due: **Monday Sept 30, 2019 @ 11:59pm**

IMPORTANT:

- **Upload this PDF** with your answers to **Gradescope by 11:59pm on Monday Sept 30, 2019**.
- **Plagiarism**: Homework may be discussed with other students, but all homework is to be completed **individually**.

For your information:

- Graded out of **100** points; **4** questions total
- Rough time estimate: ≈1-4 hours (0.5-1 hours for each question)

Revision : 2019/09/25 15:06

Question	Points	Score
Cuckoo Hashing	20	
B+Tree	45	
Extendible Hashing	25	
Suffix Trees	10	
Total:	100	

Number of Days this Assignment is Late:

Number of Late Day You Have Left:

Question 1: Cuckoo Hashing . . . . . [20 points]

Consider the following cuckoo hashing schema:

1. Both tables have a size of 4.
2. The hashing function of the first table returns the lowest two bits: $h_1(\texttt{x}) = \texttt{x \& 0b11}$.
3. The hashing function of the second table returns the next two bits: $h_2(\texttt{x}) = \texttt{(x >> 2) \& 0b11}$
4. When replacement is necessary, first select an element in the second table.
5. The original content is shown in Figure 1.

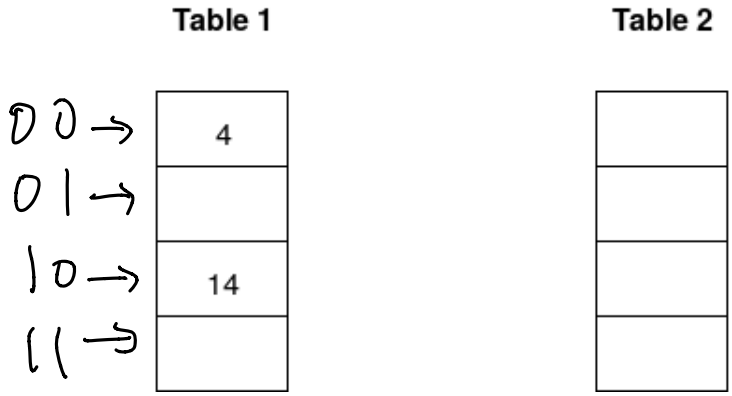
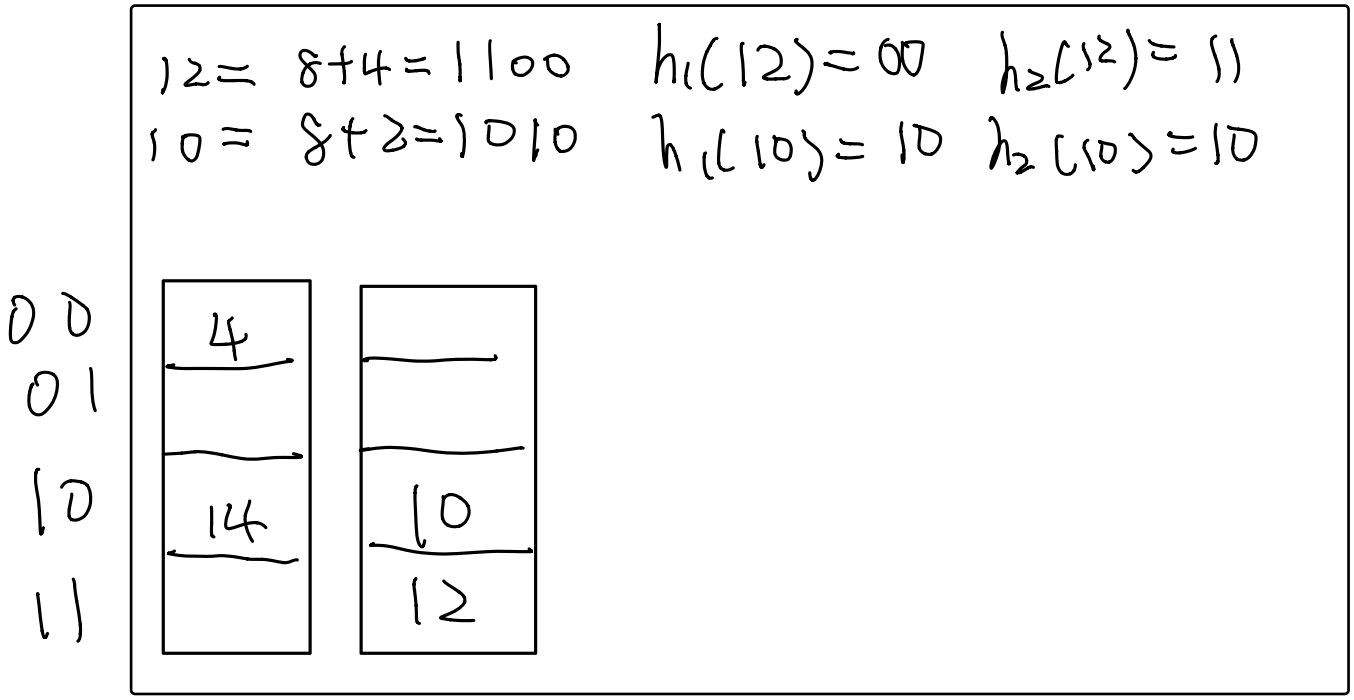


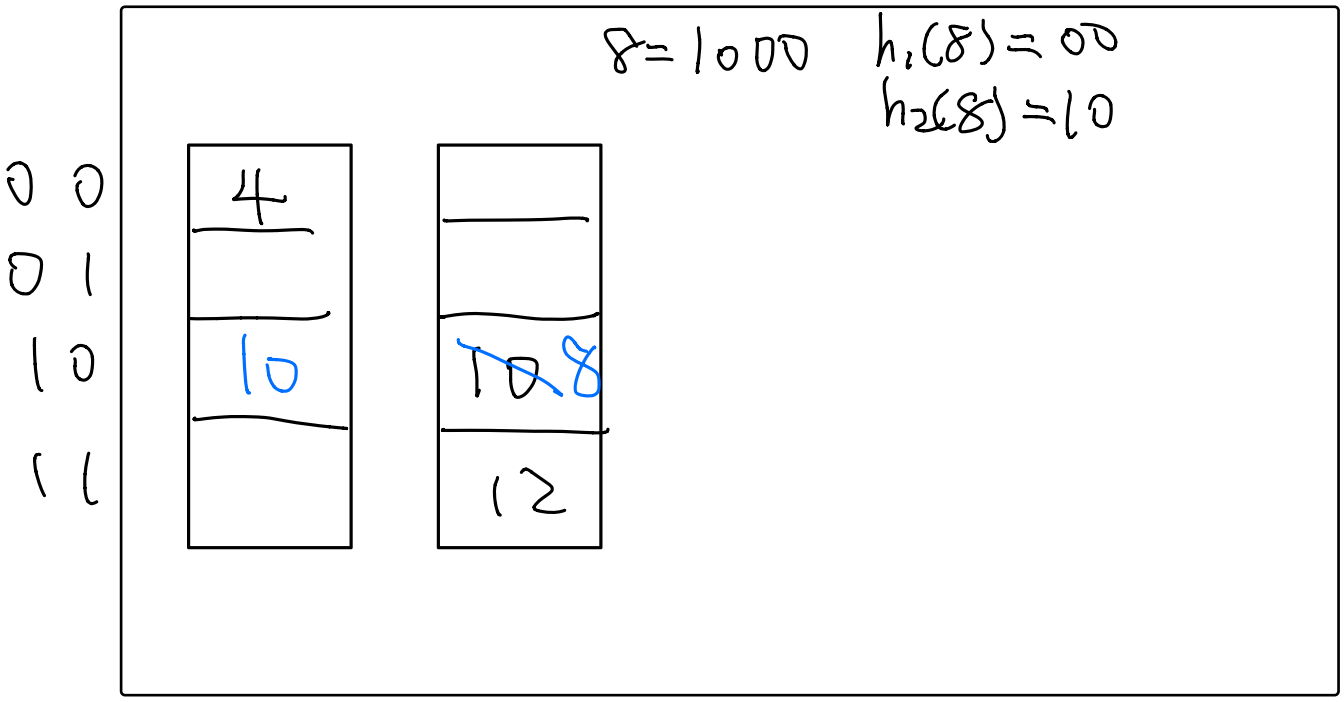
Figure 1: Initial contents of the hash tables.

Use the following template to answer the questions: `https://cmudb.io/fall2019-hw1`.

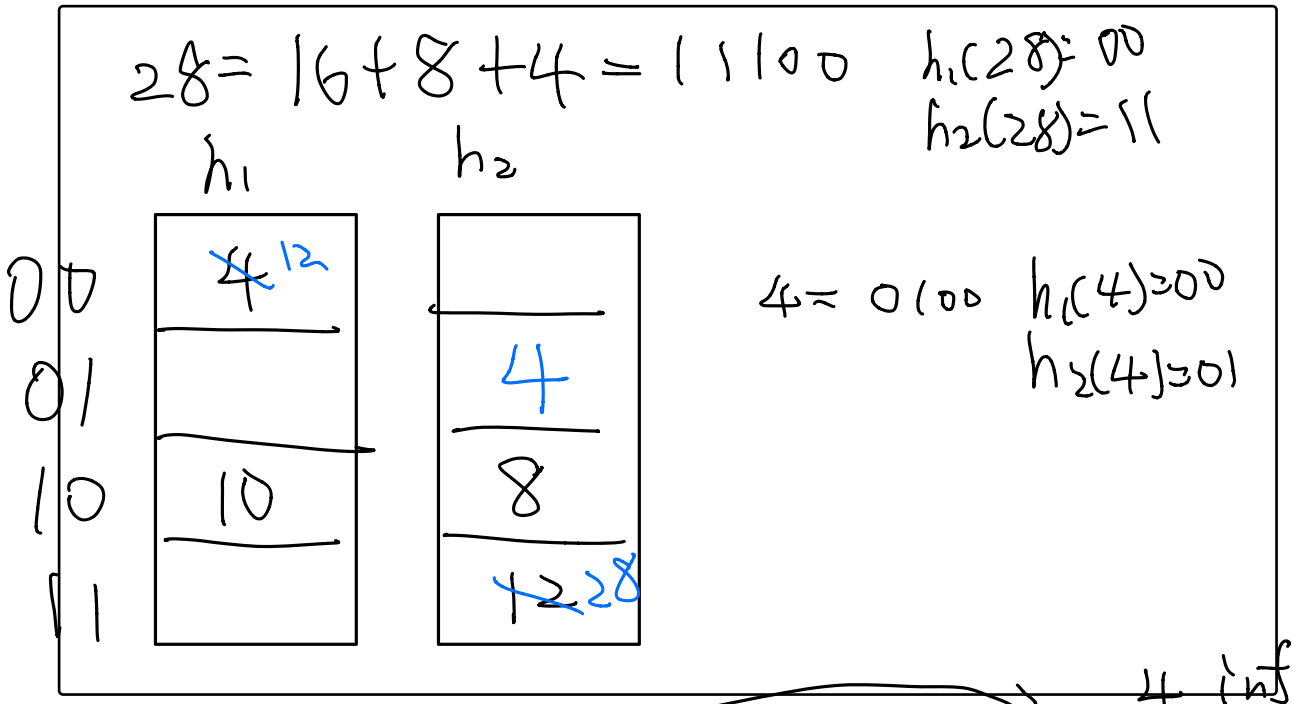
(a) **[4 points]** Insert keys 12 and 10. Draw the resulting two tables.



(b) **[4 points]** Then delete 14, and insert 8. Draw the resulting two tables.



(c) **[6 points]** Finally, insert 28. Draw the resulting two tables.

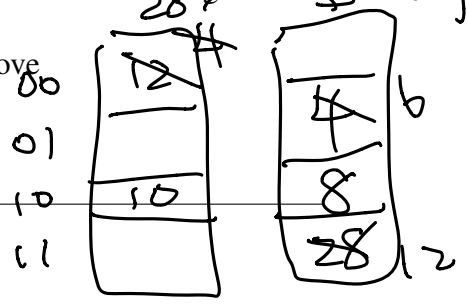


(d) **[6 points]** What is the smallest key that potentially causes an infinite loop given the tables in **(c)**

☐ 0 ☐ 2 ☐ 5 ☐ 6 ☐ 7 ☐ 9 ☐ None of the above

0010 0011 0110

4 inf loop



Question 2: B+Tree . . . . . [45 points]

Consider the following B+tree.

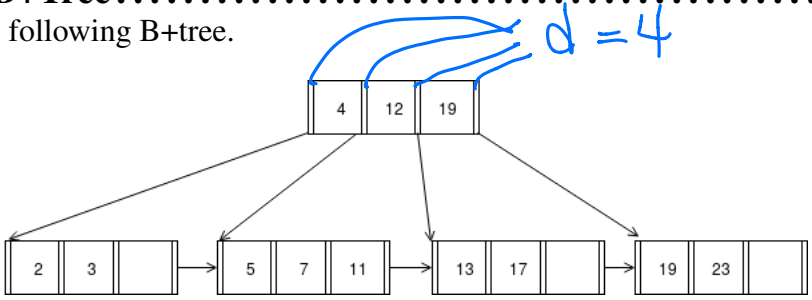


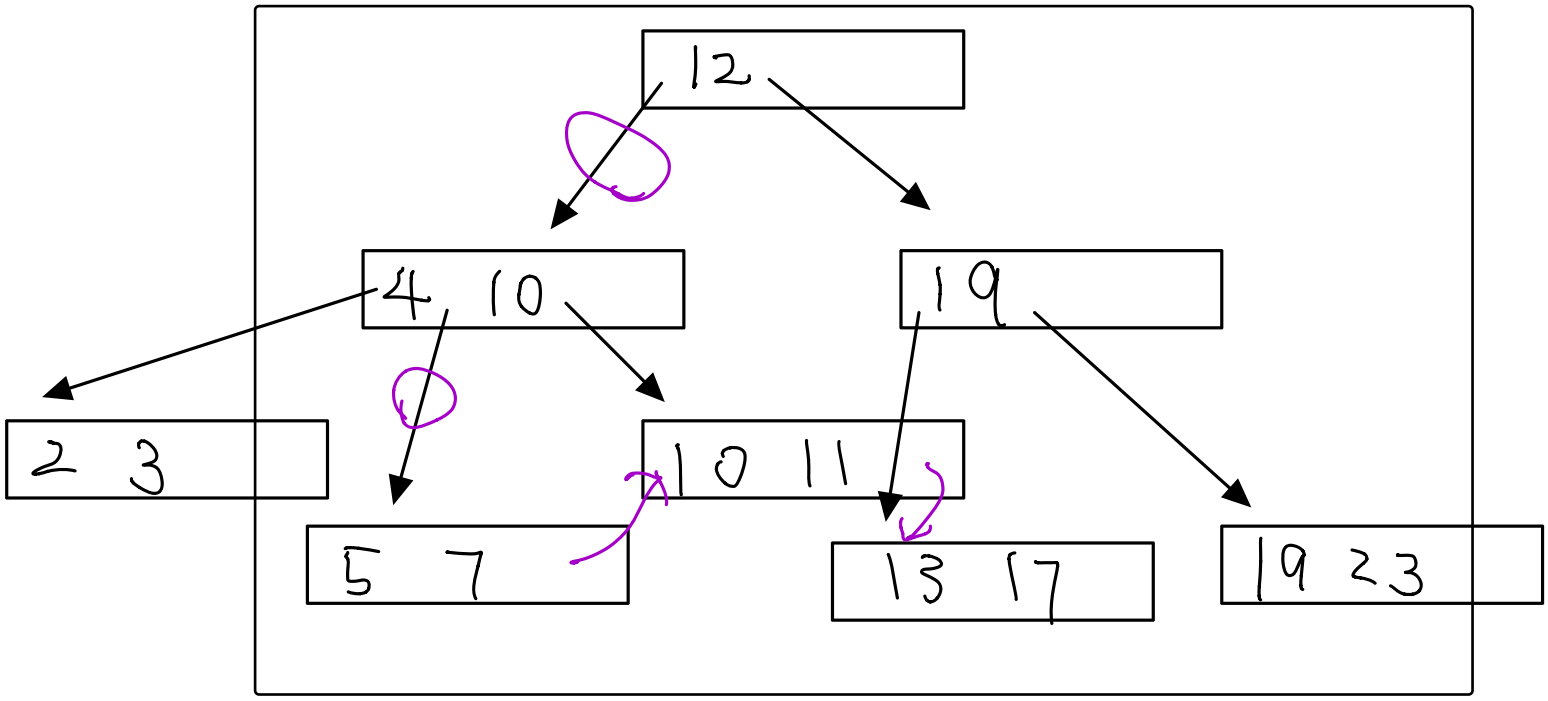
Figure 2: B+ Tree of order $d = 4$ and height $h = 2$.

When answering the following questions, be sure to follow the procedures described in class and in your textbook. You can make the following assumptions:

- A left pointer in an internal node guides towards keys $<$ than its corresponding key, while a right pointer guides towards keys $\geq$.
- A leaf node underflows when the number of **keys** goes bellow $\lceil \frac{d-1}{2} \rceil$.
- An internal node underflows when the number of **pointers** goes below $\lceil \frac{d}{2} \rceil$.

Use the following `draw.io` template for your answers:
`https://cmudb.io/fall2019-hw2`

(a) **[15 points]** Insert 10^* into the B+tree. Draw the resulting tree.

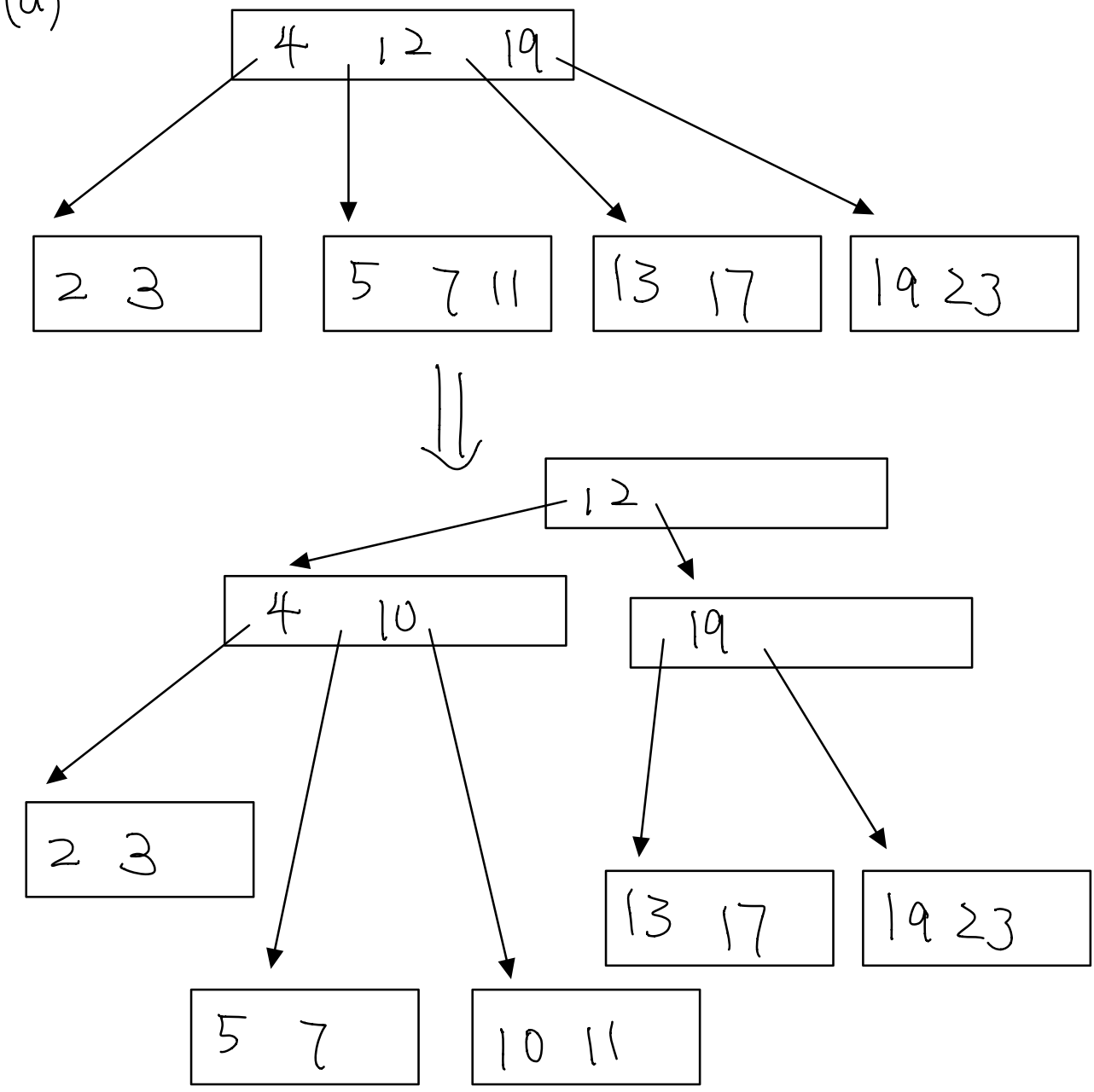


(b) **[5 points]** How many pointers (parent-to-child and sibling-to-sibling) do you chase to find all keys between 5 and 15?

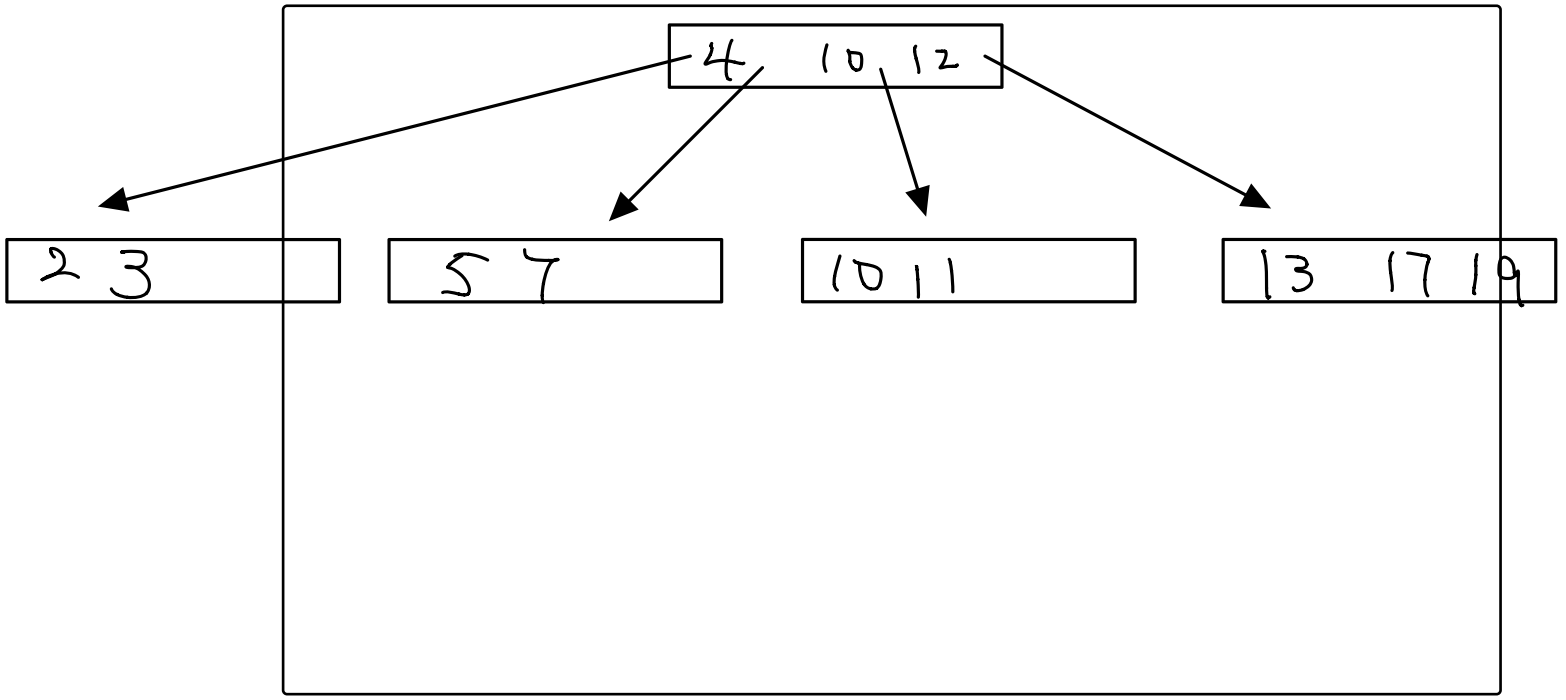
☐ 2 ☐ 3 ☑ 4 ☐ 5 ☐ 6 ☐ 7

Find 5, then keep going right

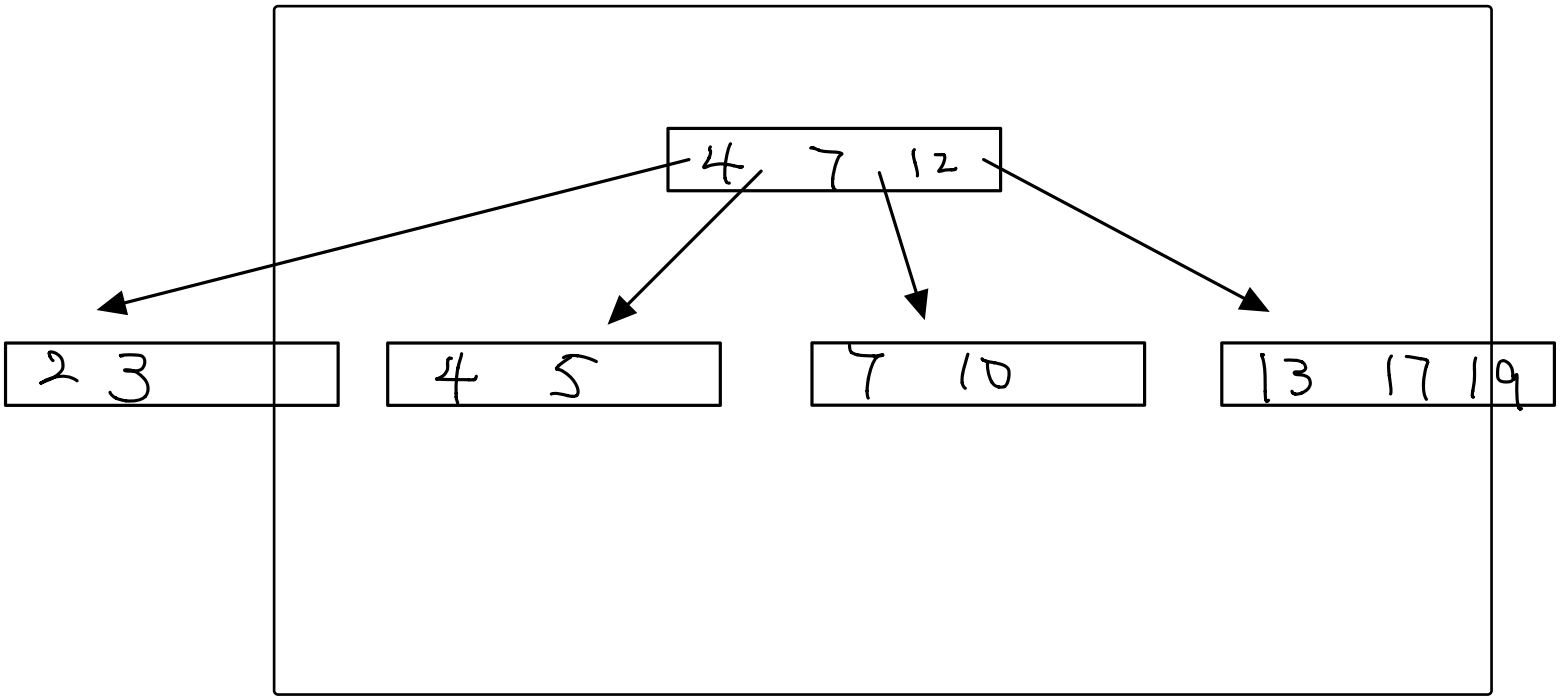
(a)



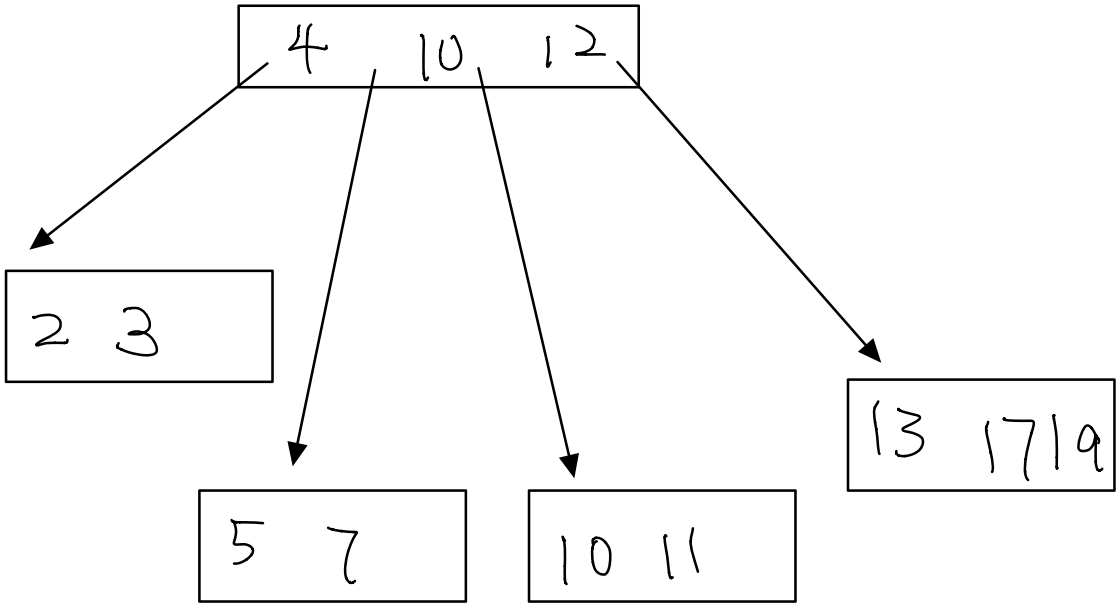
(c) **[15 points]** Then delete 23^*. Draw the resulting tree.



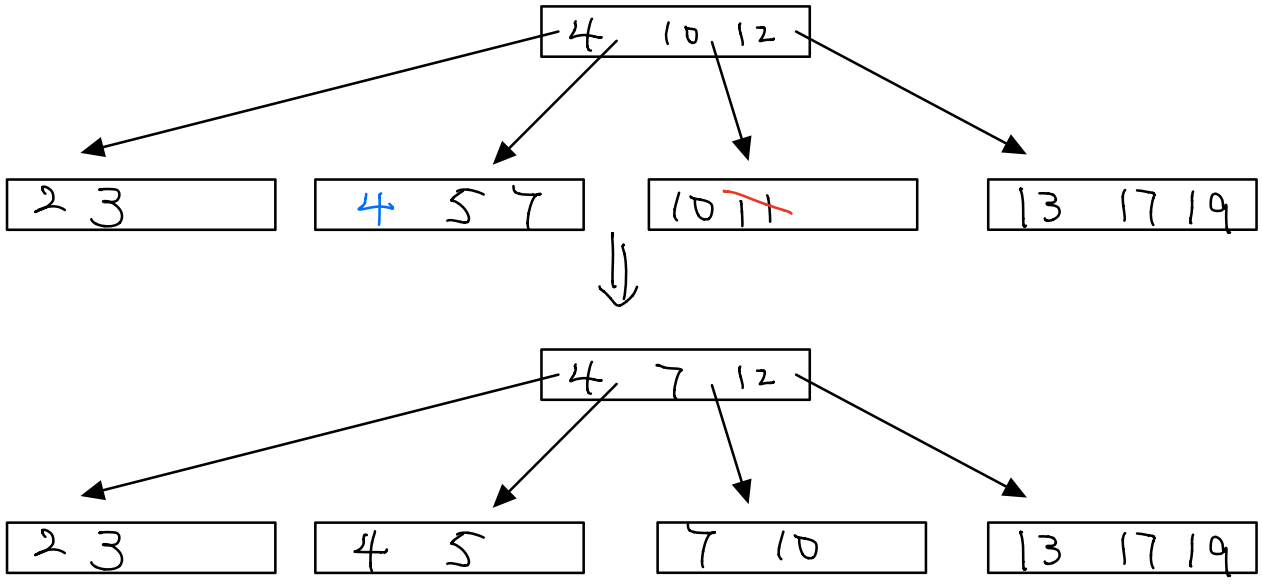
(d) **[10 points]** Finally insert 4^* and delete 11^*. Draw the resulting tree.



(b)


4 10 12
2 3
5 7
10 11
13 17 19

(c)


4 10 12
2 3
4 5 7
10 11
13 17 19
4 7 12
2 3
4 5
7 10
13 17 19

Question 3: Extendible Hashing . . . . . [25 points]

Consider an extendible hashing structure such that:

- Each bucket can hold up to two records.
- The hashing function uses the lowest g bits, where g is the global depth.

(a) Starting from an empty table, insert keys 15, 3, 7, 14.

i. **[3 points]** What is the global depth of the resulting table?
☐ 0 ☐ 1 ☐ 2 ☑ 3 ☐ 4 ☐ None of the above

ii. **[3 points]** What is the local depth the bucket containing 14?
☐ 0 ☑ 1 ☐ 2 ☐ 3 ☐ 4 ☐ None of the above

iii. **[3 points]** What is the local depth of the bucket containing 3?
☐ 0 ☐ 1 ☐ 2 ☑ 3 ☐ 4 ☐ None of the above

(b) Starting from the result in **(a)**, you insert keys 1, 9, 23, 11, 17.

i. **[4 points]** Which key will first cause a split (without doubling the size of the table)?
☐ 1 ☐ 9 ☐ 23 ☐ 11 ☑ 17 ☐ None of the above

ii. **[4 points]** Which key will first make the table double in size?
☐ 1 ☐ 9 ☑ 23 ☐ 11 ☐ 17 ☐ None of the above

(c) Now consider the table below, along with the following deletion rules:

1. If two buckets have the same local depth d, and share the first $d-1$ bits of their indexes (e.g. `010` and `110` share the first 2 bits), then they can be merged if the total capacity fits in a single bucket. The resulting local depth is $d-1$.
2. If the global depth g becomes strictly greater than all local depths, then the table can be halved in size. The resulting global depth is $g-1$.

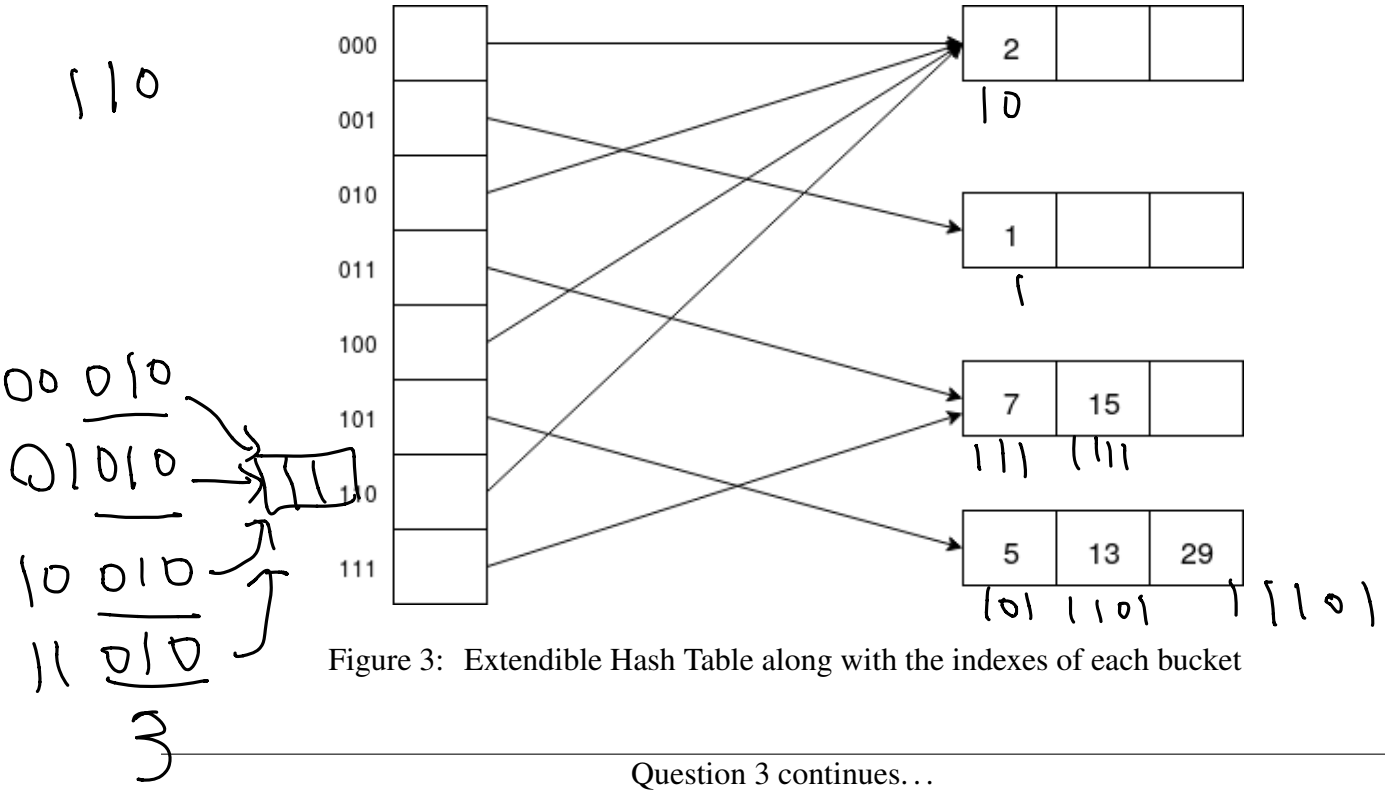


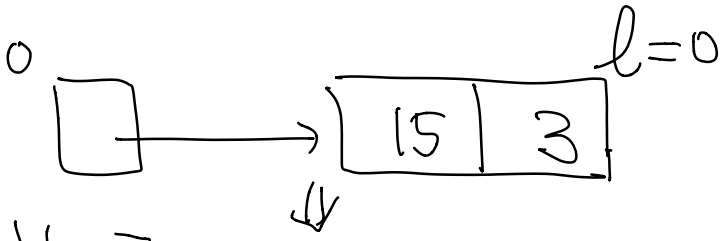
Figure 3: Extendible Hash Table along with the indexes of each bucket

Starting from the table above, delete keys 2, 7, 13, 15, 29.

i. **[4 points]** Which deletion first causes a reduction in a local depth.
☐ 2 ☐ 7 ☑ 13 ☐ 15 ☐ 29 ☐ None of the above

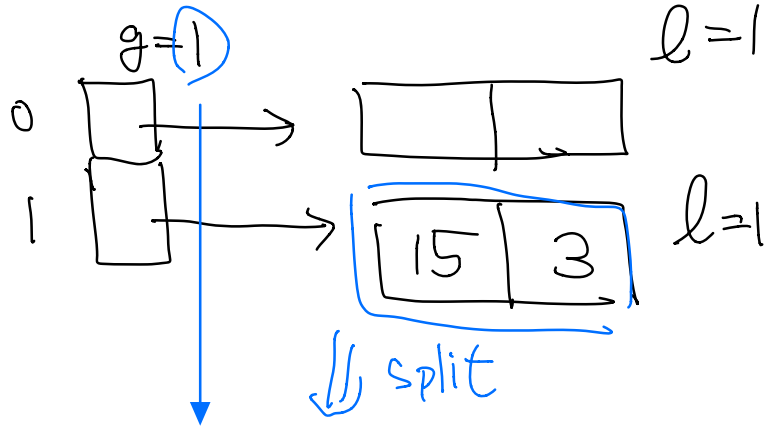
ii. **[4 points]** Which deletion first causes a reduction in global depth.
☐ 2 ☐ 7 ☑ 13 ☐ 15 ☐ 29 ☐ None of the above

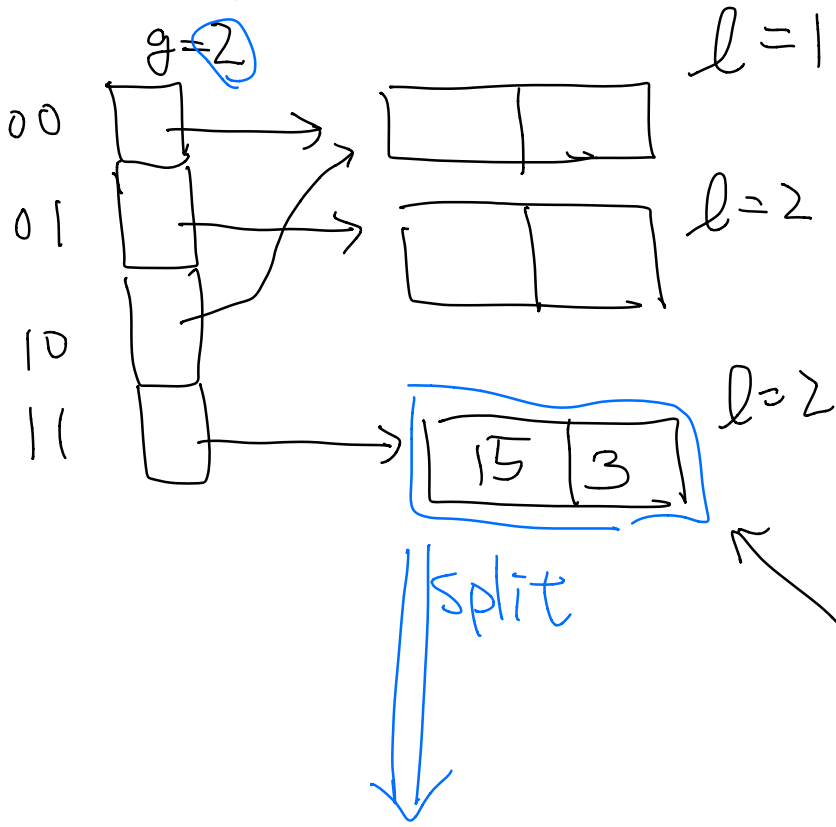
(a) $g=0$



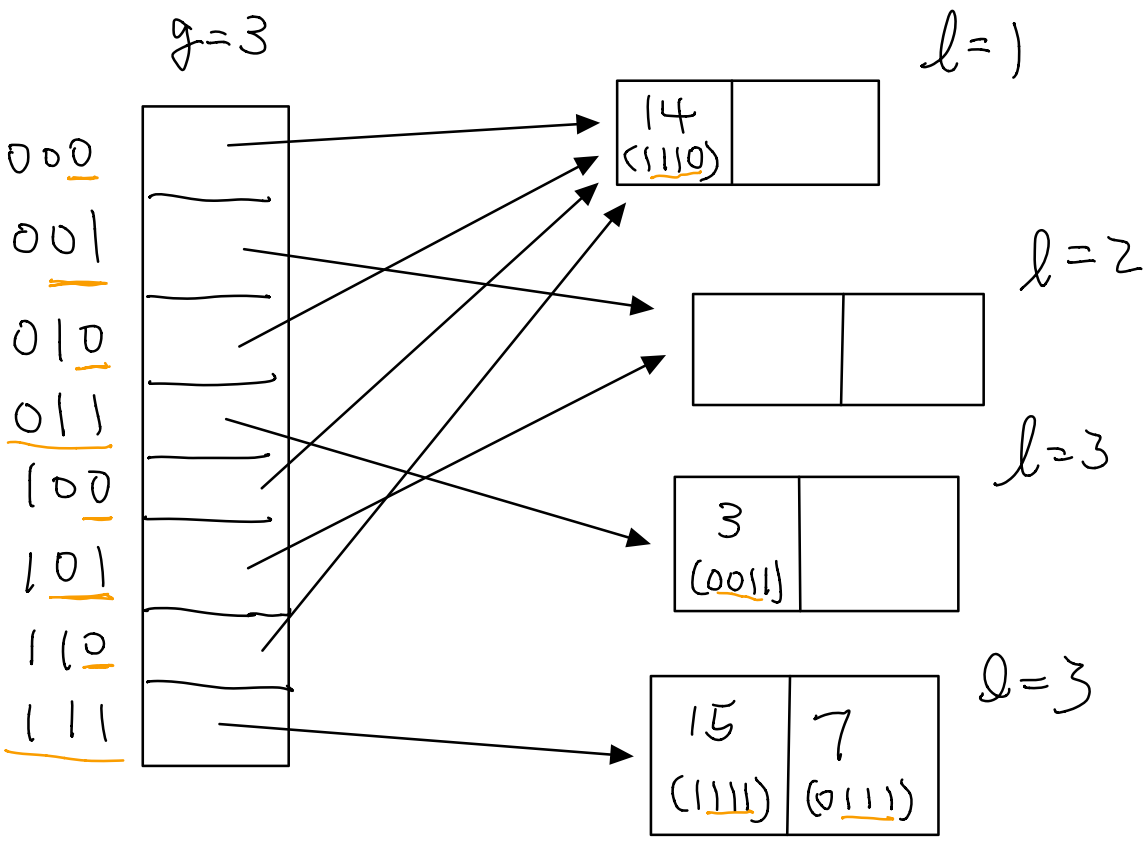
15 = 8+4+2+1
= 1111
3 = 0011
7 = 0111
14 = 1110

add 7

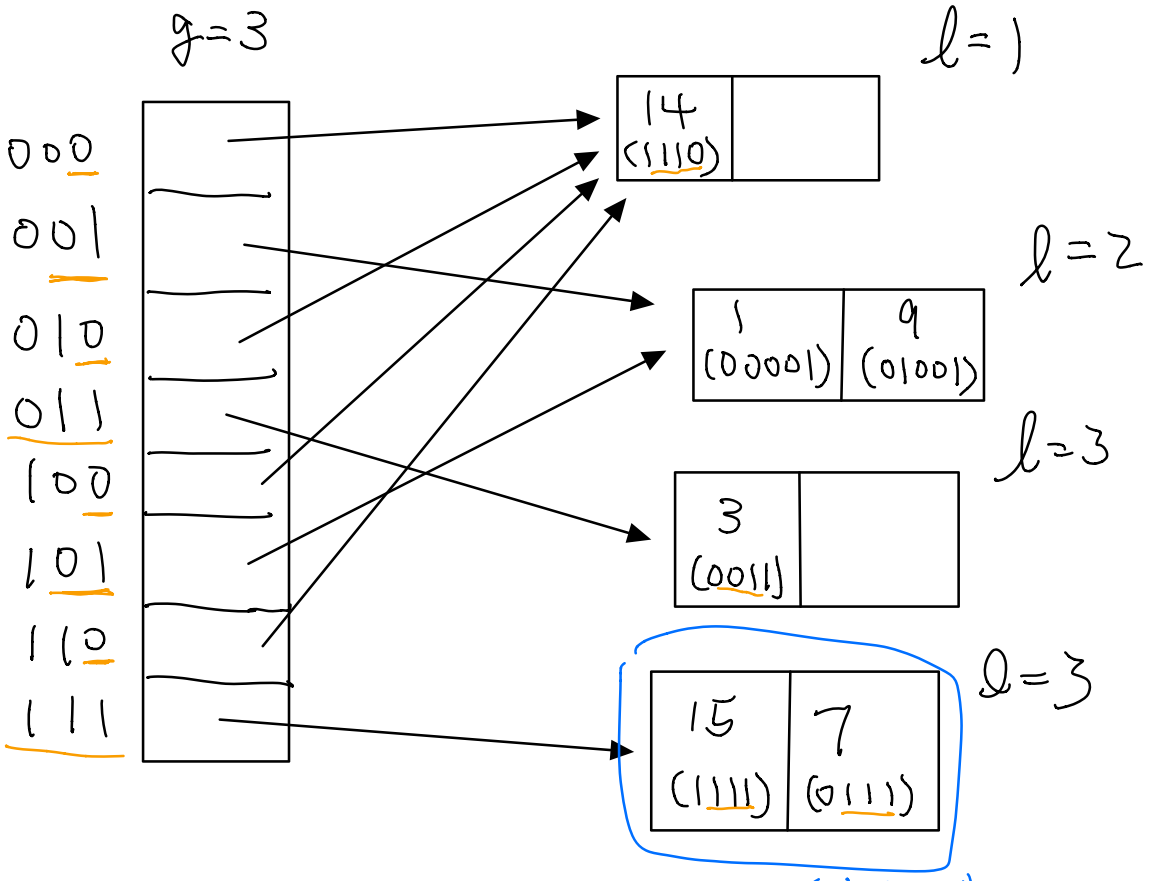




7 = 0111
⇒ still full
⇒ split


g=3
l=1
000
001
010
011
100
101
110
111
14
(1110)
l=2
l=3
3
(0011)
l=3
15
(1111)
7
(0111)

(b)



Split when 23 is inserted

1, 9, 23, 11, 17

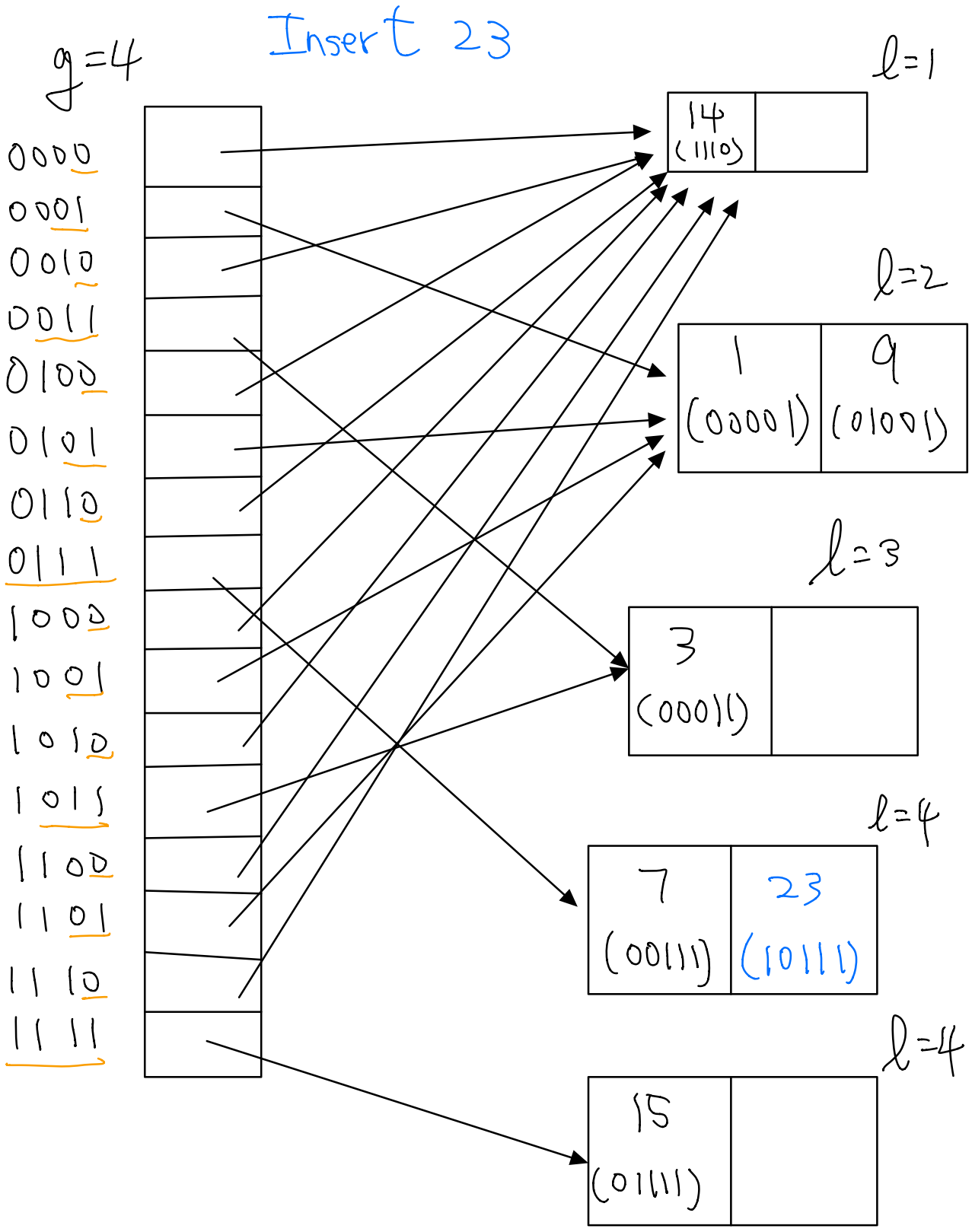
1 = 00001

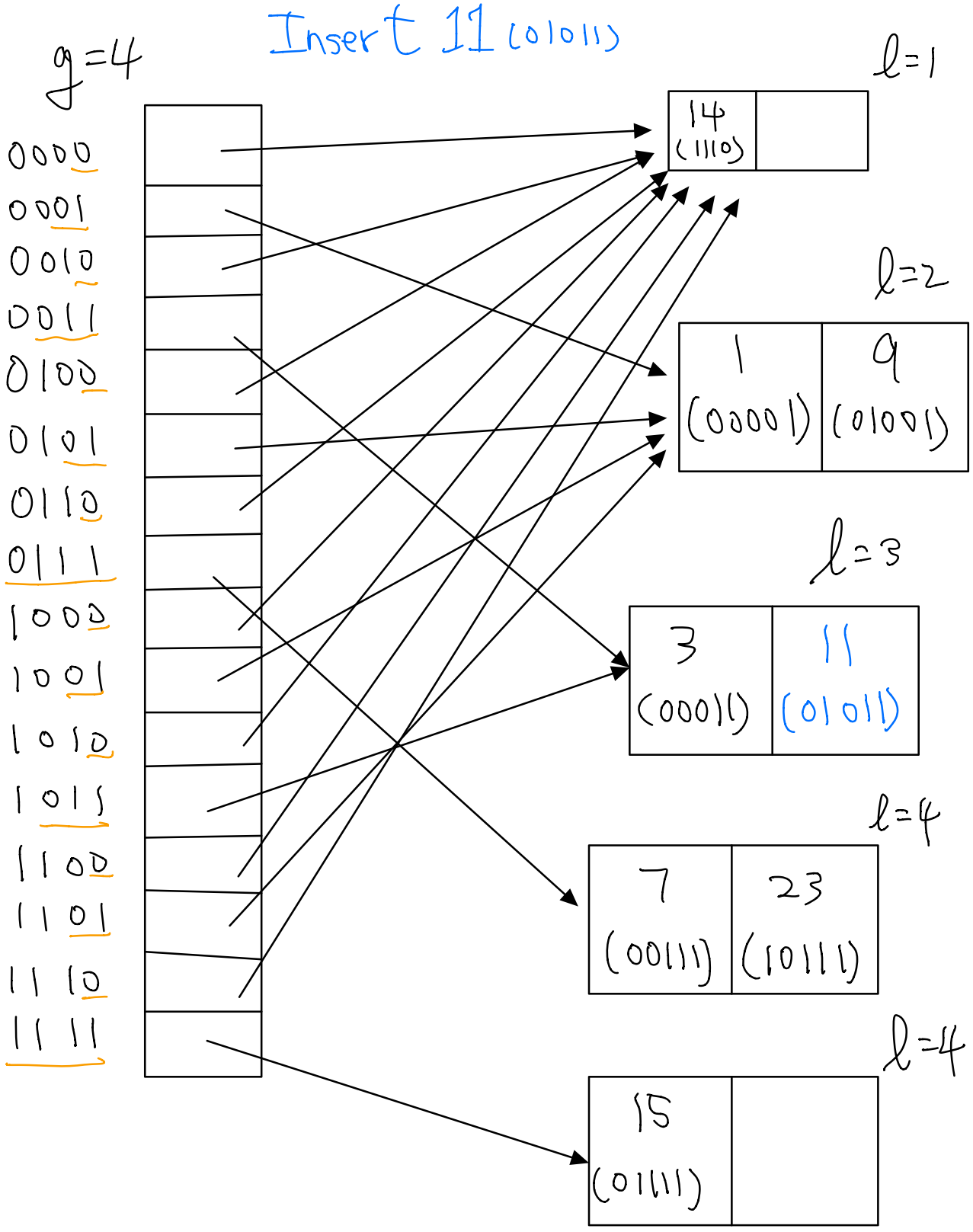
9 = 01001

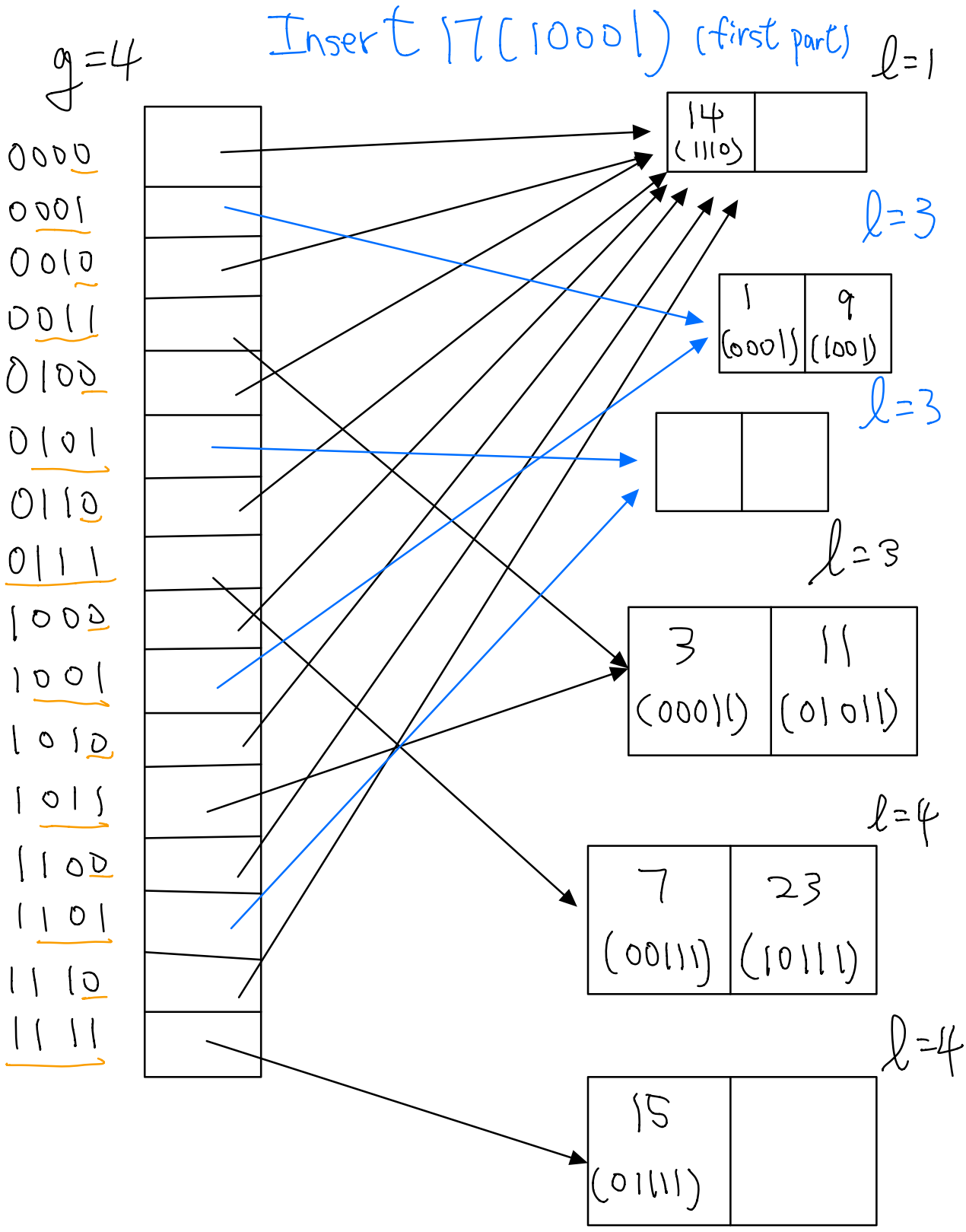
23 = 16+4+2+1 = 10111

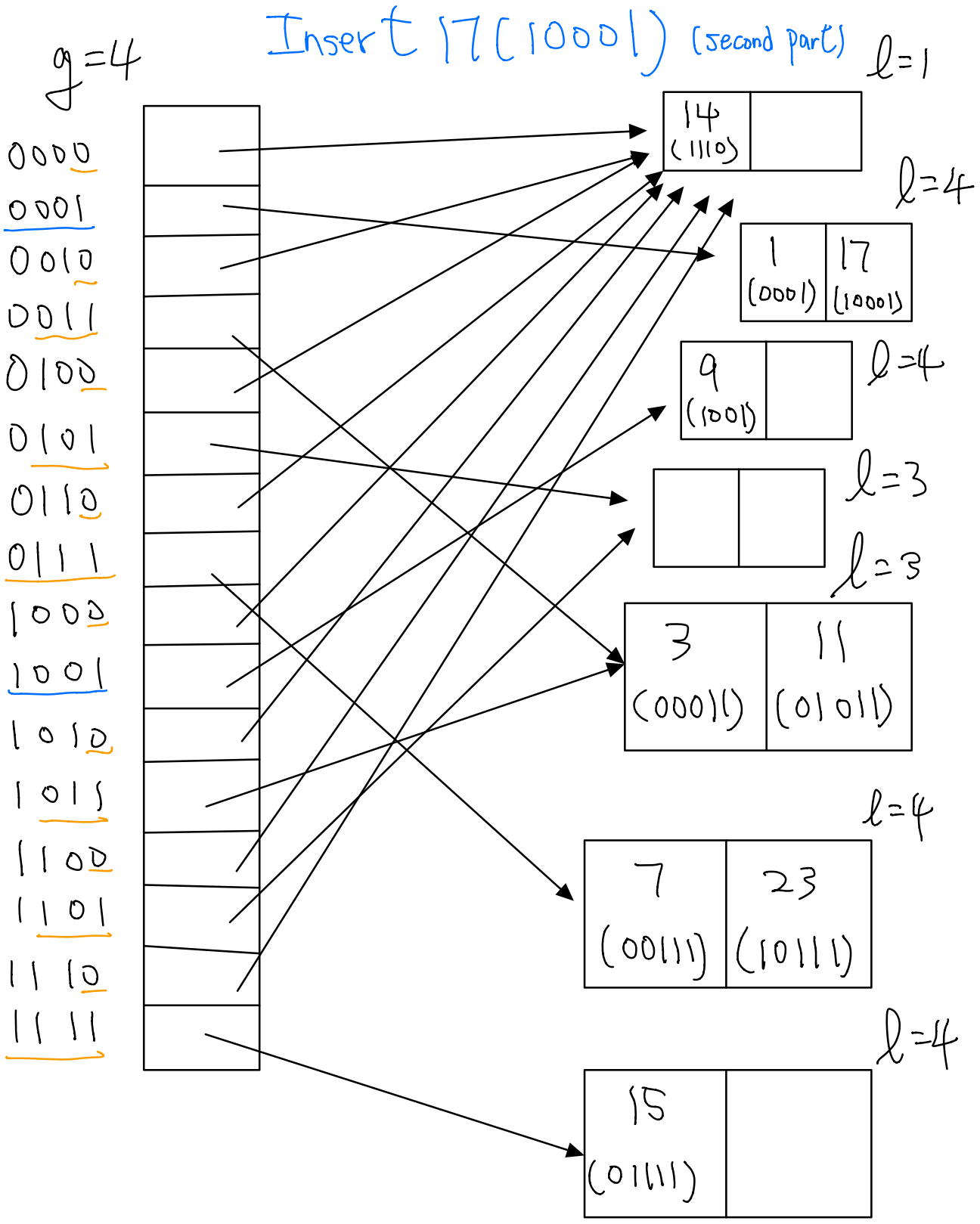
11 = 8+2+1 = 01011

17 = 16+1 = 10001

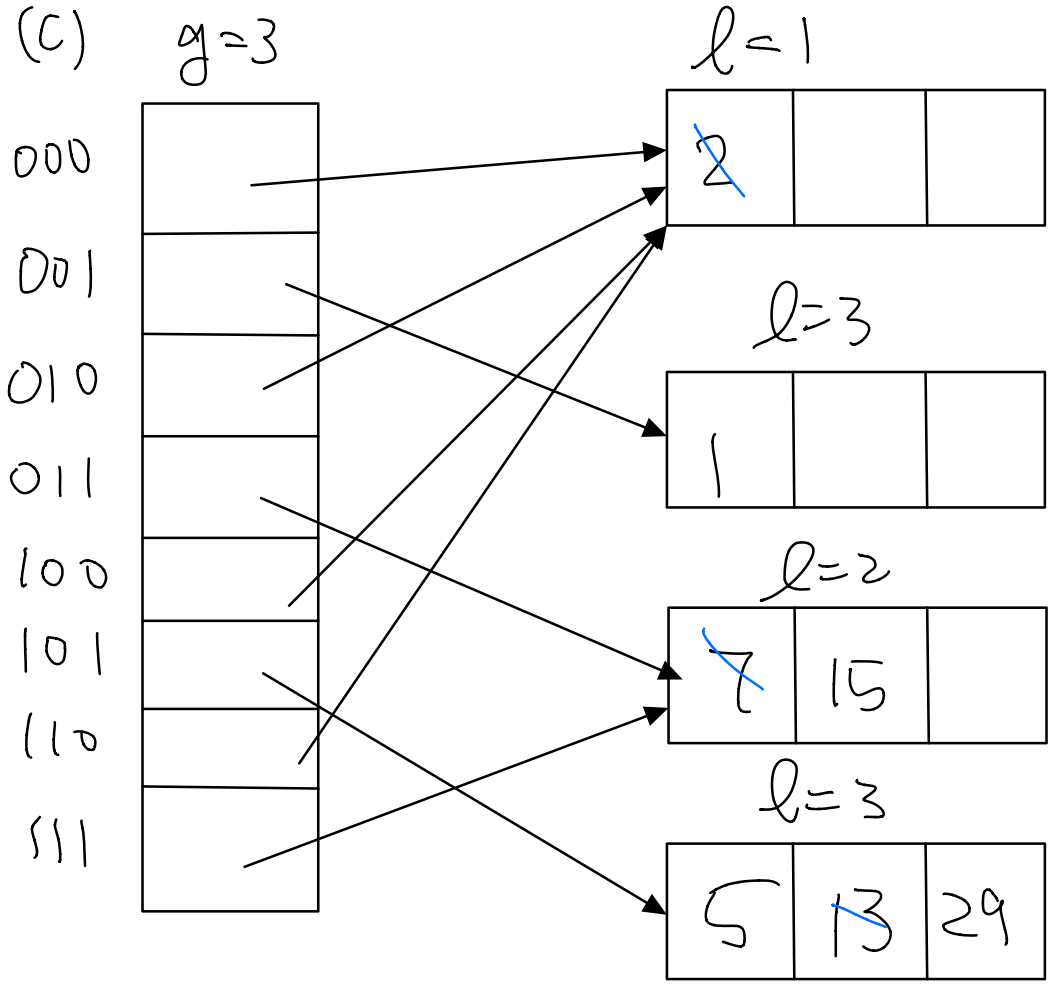

Insert 23
g=4
0000
0001
0010
0011
0100
0101
0110
0111
1000
1001
1010
1011
1100
1101
1110
1111
l=1
14
(1110)
l=2
1
(0001)
9
(1001)
l=3
3
(0011)
l=4
7
(0111)
23
(1011)
l=4
15
(1111)

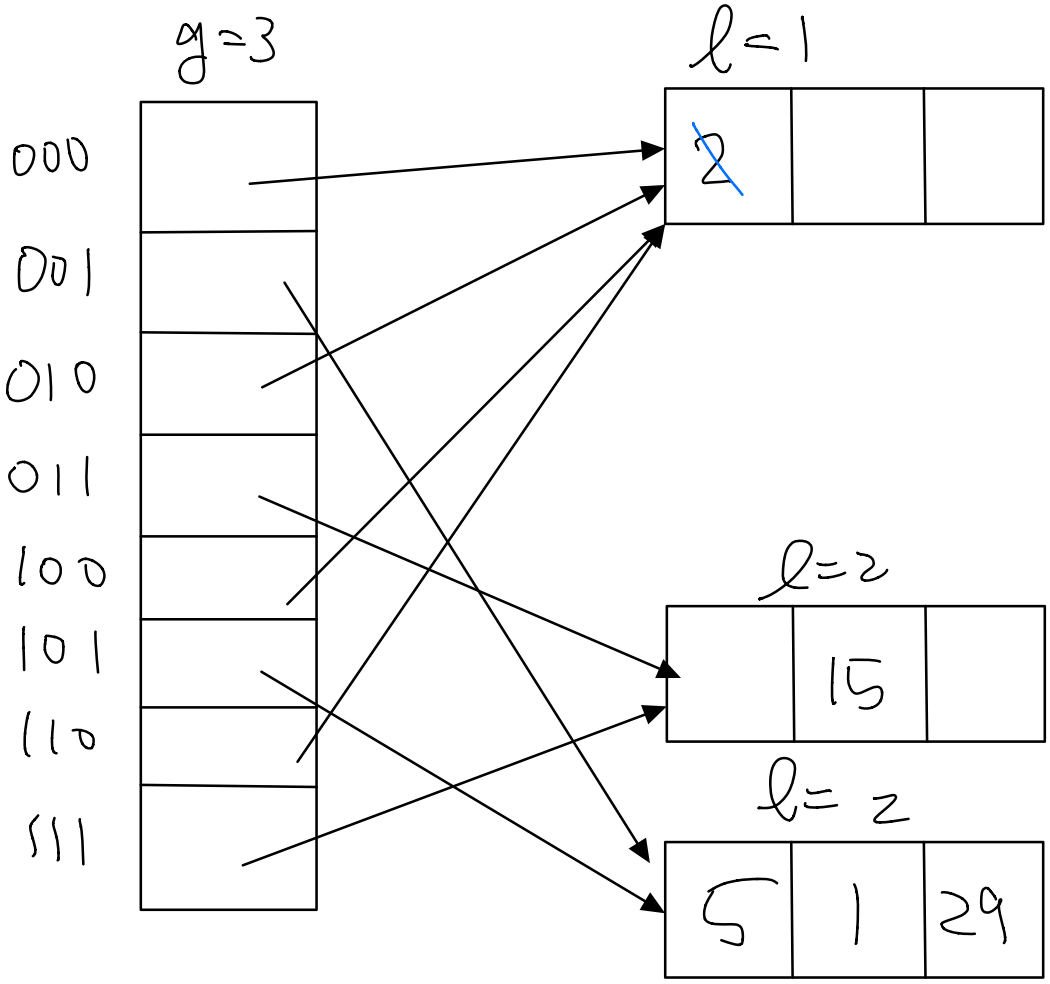

Insert 11 (01011)
g=4
0000
0001
0010
0011
0100
0101
0110
0111
1000
1001
1010
1011
1100
1101
1110
1111
l=1
14
(1110)
l=2
1
(0001)
9
(01001)
l=3
3
(00011)
11
(01011)
l=4
7
(00111)
23
(10111)
l=4
15
(01111)


Insert 17 (10001) (first part)
g=4
0000
0001
0010
0011
0100
0101
0110
0111
1000
1001
1010
1011
1100
1101
1110
1111
14 (1110)
l=1
l=3
1 (0001)
9 (1001)
l=3
l=3
3 (00011)
11 (01011)
l=4
7 (00111)
23 (10111)
l=4
15 (01111)

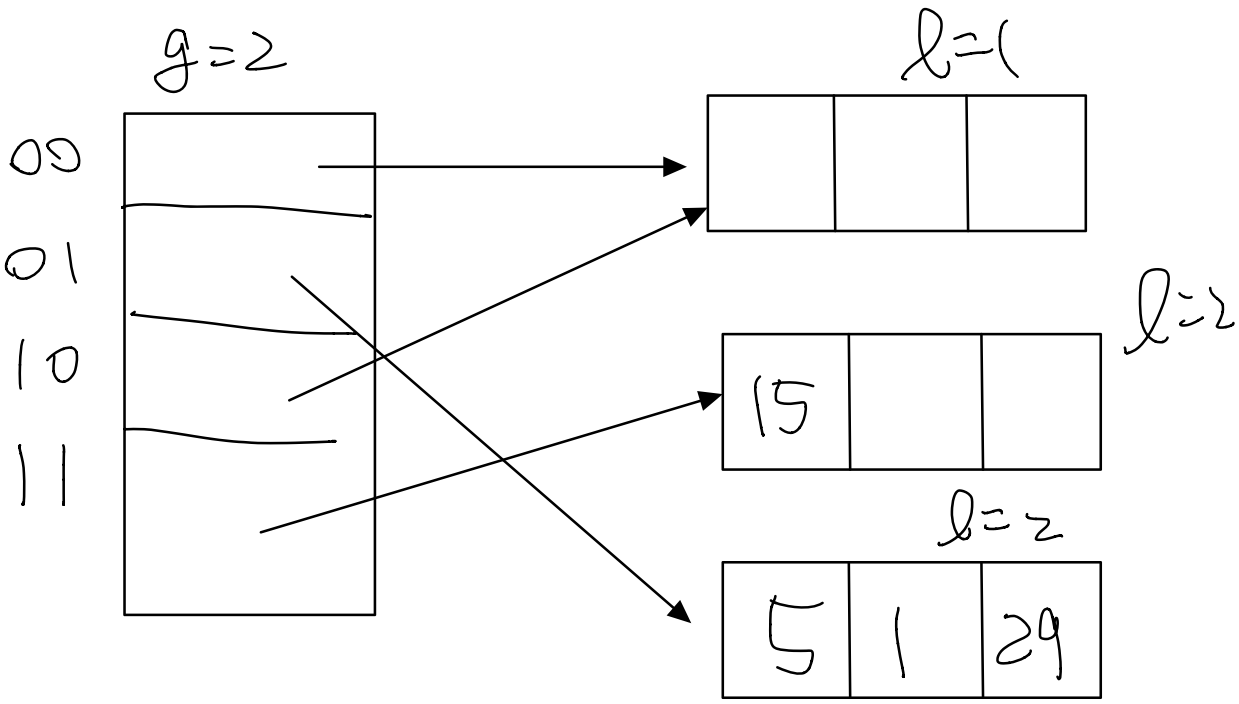

Insert 17 (10001) (second part)
g=4
0000
0001
0010
0011
0100
0101
0110
0111
1000
1001
1010
1011
1100
1101
1110
1111
14 (1110)
l=1
1 (0001)
17 (10001)
l=4
9 (1001)
l=4
l=3
3 (00011)
11 (01011)
l=3
7 (00111)
23 (10111)
l=4
15 (01111)
l=4

(c)





⇒ halve the size


g=2
00
01
10
11
l=1
l=2
15
l=2
5
1
29

Question 4: Suffix Trees ..... [10 points]

Consider the following suffix tree for **unsigned 32-bit integers**.

one hex = 4 bits

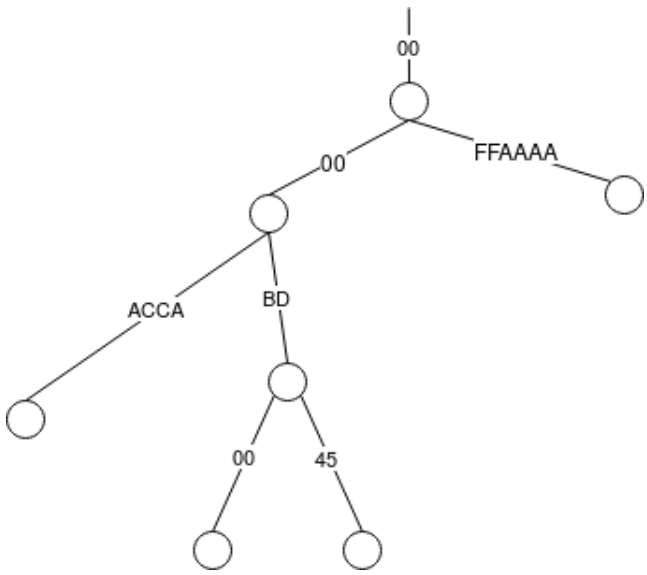


Figure 4: Suffix Tree

(a) **[3 points]** Which of the following elements belong to the suffix tree. Select all that apply.

☐ `0x45BD0000` ☐ `0x0000CAAC` ☐ `0xFFAAAA00` ☐ `0xACCA0000` ☐ `0xBD000000`
☑ None of the above

(b) **[7 points]** Insert the key `0x00FFAABB`. Draw the resulting tree using this template: `https://cmudb.io/fall2019-hw4`.