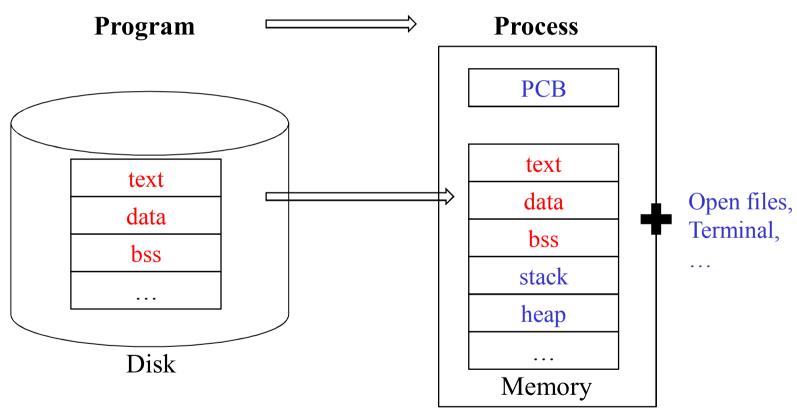
Process control

Process

- a program in execution
- program image + environment
 - text(code) /data / stack / heap segment
 - kernel data structure, address space, open files, ...

```
/* test_program.c */
#include <stdio.h>
int a,b;
int glob_var = 3;
void test_func(void) {
   int local_var, *buf;
   buf = (int *) malloc(10,000 * sizeof(int));
int main(int argc, char *argv[]) {
   int i = 1;
   int local_var;
   a = i + 1;
                                          $ gcc -o test_program test_program.c
   printf("value of a = \%d\n", a);
   test_func();
```

\$./test_program



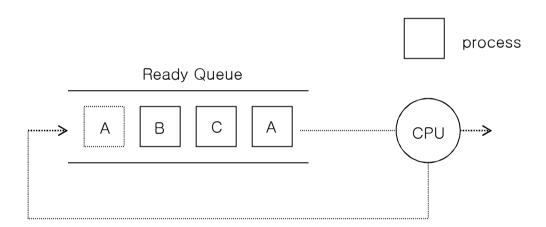
* bss: Block Started by Symbol

PCB: Process Control Block

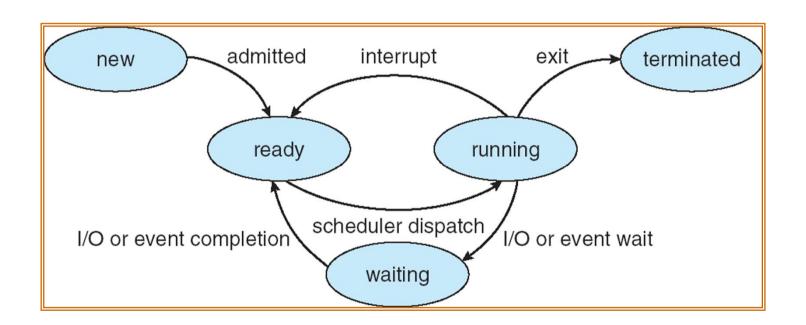
UNIX is a multitasking system.

```
$ ps –е
PID TTY
              TIME CMD
1?
       00:00:00 init
2?
       00:00:00 migration/0
       00:00:00 ksoftirqd/0
       00:00:00 watchdog/0
4?
5?
       00:00:00 migration/1
       00:00:00 ksoftirqd/1
6?
       00:00:00 watchdog/1
7?
8?
       00:00:00 migration/2
29122 pts/9 00:00:00 bash
29151?
           00:00:00 sshd
29170?
           00:00:00 sshd
29171 pts/10 00:00:00 bash
30465?
           00:00:00 httpd
```

UNIX is a time sharing system.



State of a process



Process identifiers

process id

- Every process has a unique process ID.
- As processes terminate, IDs become candidates for reuse.

```
$ ps -е
PID TTY
              TIME CMD
       00:00:00 init
       00:00:00 migration/0
       00:00:00 ksoftirqd/0
       00:00:00 watchdog/0
5?
       00:00:00 migration/1
29122 pts/9 00:00:00 bash
29151 ?
           00:00:00 sshd
29170?
           00:00:00 sshd
29171 pts/10 00:00:00 bash
30465?
           00:00:00 httpd
```

Process identifiers

- Process 0 (scheduler process or swapper)
 - Part of the kernel.
 - System process.
- Process 1 (init process)
 - /sbin/init
 - Invoked at the end of the bootstrap procedure.
 - Initialize the UNIX system with /etc/rc* or /etc/inittab.
 - Never dies.
 - User process with superuser privilege.

Process identifiers

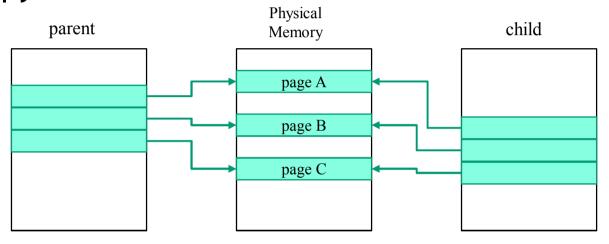
#include <unistd.h>

pid t fork(void);

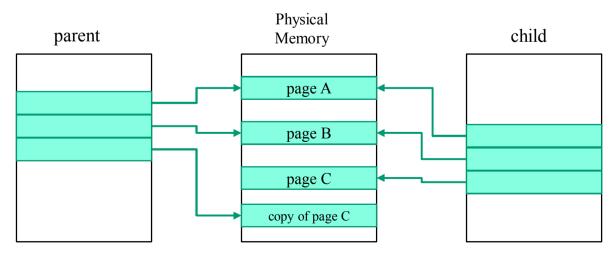
Returns: 0 in child, process ID of child in parent, -1 on error

- Create a child process.
 - It is called once but returns twice.
 - Return value of child: 0
 - Return value of parent: PID of child
 - Both child and parent continue executing with the instruction that follows the call to fork().
 - The child is a copy of parent.
 - Child gets a copy of the parent's data, heap, and stack.
 - Parent and child often share the text segment.

Copy on write



• After parent process write a page C.



Example

```
#include "apue.h"
int glob = 6; /* external variable in initialized data */
char buf[] = "a write to stdout\n";
int
main(void)
                /* automatic variable on the stack */
  int
         var;
  pid t pid;
  var = 88;
  if (write(STDOUT_FILENO, buf, sizeof(buf)-1) != sizeof(buf)-1)
    err_sys("write error");
  printf("before fork\n"); /* we don't flush stdout */
```

Example(cont.)

```
#include "apue.h"
int glob = 6;
char buf[] = "a write to stdout\n";
int main(void)
  int
         var;
  pid t pid;
  var = 88:
  if (write(STDOUT FILENO, buf, sizeof(buf)-1) != sizeof(buf)-1)
    err sys("write error");
  printf("before fork\n"); /* we don't flush stdout */
  if ((pid = fork()) < 0) {
                                 /* pid is child's pid(non-0). */
    err sys("fork error");
                                 /* child */
  else if (pid == 0) 
                                 /* modify variables */
    glob++;
    var++;
  } else {
                                  /* parent */
    sleep(2);
  printf("pid = \%d, glob = \%d, var = \%d\n", getpid(), glob, var);
  exit(0);
```

```
#include "apue.h"
int glob = 6;
char buf[] = "a write to stdout\n";
int main(void)
  int
         var;
  pid_t pid;
  var = 88:
  if (write(STDOUT FILENO, buf, sizeof(buf)-1) != sizeof(buf)-1)
     err sys("write error");
  printf("before fork\n"); /* we don't flush stdout */
  if ((pid = fork()) < 0) {
                                  /* pid is 0. */
     err sys("fork error");
                                  /* child */
  } else if (pid == 0) {
                                  /* modify variables */
     glob++;
     var++;
  } else {
                                  /* parent */
     sleep(2);
  printf("pid = \%d, glob = \%d, var = \%d\n", getpid(), glob, var);
  exit(0);
```

Parent process

Child process



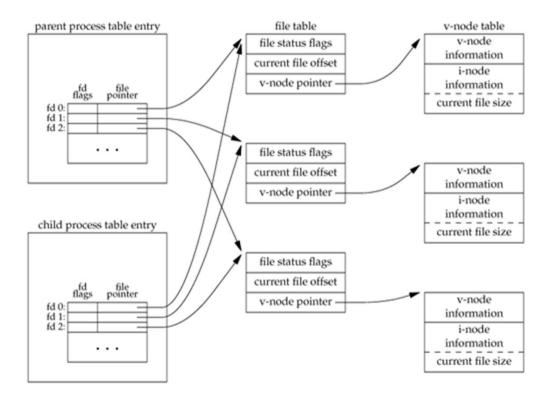
◉ 실행

```
$ ./a.out
a write to stdout
before fork
pid = 430, glob = 7, var = 89
pid = 429, glob = 6, var = 88
$ ./a.out > temp.out
$ cat temp.out
a write to stdout
before fork
pid = 432, glob = 7, var = 89
before fork
pid = 431, glob = 6, var = 88
```

We never know if the child starts executing before the parent or vice versa. - It depends on the process scheduling algorithm.

File sharing

Sharing of open files between parent and child after fork()



- Information shared by child and parent.
 - real-uid(gid), effective-uid(gid)
 - controlling terminal
 - current working directory, root directory
 - signal handlers
 - environment
 - resource limits
- differences between child and parent.
 - the return value from fork.
 - PID and PPID
 - child's resource utilizations are set to 0.
 - pending signals

- Two main reasons for fork to fail
 - if there are already too many processes in the system.
 - if the total number of processes for this real user ID exceeds the system's limit.
- Two uses for fork
 - A process wants to duplicate itself.
 - Parent and child can each execute different sections of code at the same time.
 - Common for network servers.
 - A process wants to execute a different program.
 - Common for shells.
 - Child does exec() right after it returns from fork().

#include <unistd.h>

pid t vfork(void);

Returns: 0 in child, process ID of child in parent, -1 on error

- same calling sequence and same return values as fork().
- intended to create a new process when the purpose of the new process is to exec() a new program.
 - Does not copy the address space of parent into the child.
 - The child calls exec() or exit() right after the vfork().
 - The child runs in the address space of the parent.
 - Provides an efficiency.
- vfork() guarantees that the child runs first.

Example

```
#include "apue.h"
int glob = 6; /* external variable in initialized data */
int
main(void)
               /* automatic variable on the stack */
  int var;
  pid_t pid;
  var = 88;
  printf("before vfork\n"); /* we don't flush stdio */
  if ((pid = vfork()) < 0) {
    err_sys("vfork error");
  } else if (pid == 0) { /* child */
    glob++; /* modify parent's variables */
    var++;
             /* child terminates */
    exit(0);
```

Example(cont.)

```
/*
 * Parent continues here.
 */
printf("pid = %d, glob = %d, var = %d\n", getpid(), glob, var);
exit(0);
}
```

◉ 실행

```
$ ./a.out
before vfork
pid = 29039, glob = 7, var = 89
$
```

exit()

```
#include <stdlib.h>
void exit(int status);
void _Exit(int status);

#include <unistd.h>
void _exit(int status);
```

- causes normal program termination and the value of status is returned to the parent.
- _exit and _Exit
 - return to the kernel immediately.
- exit
 - performs cleanup processing and then returns to the kernel.
 - all open streams are flushed and closed.

Process termination

- Two types of process termination
 - normal termination
 - return from main
 - calling exit
 - calling _exit or _Exit
 - abnormal termination
 - calling abort
 - Receipt of a signal
- Regardless of how a process terminate, the same code is eventually executed.
 - Close open descriptors, release the memory, ...

Process termination

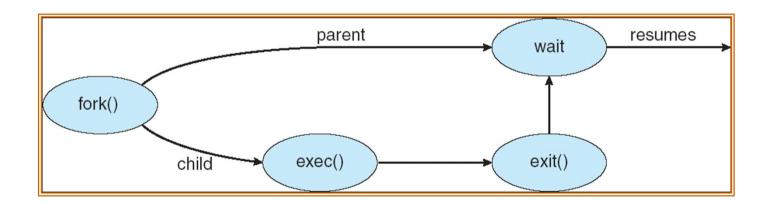
Termination status

- Terminating process notify its parent how it terminated.
- normal termination
 - pass an exit status as argument to exit() or _exit().
 - return value from main()
 - exit status is converted into a termination status.
- abnormal termination
 - Kernel generates a termination status to indicate the reason for the abnormal termination.
- The parent of the terminated process can obtain the termination status from wait() or waitpid().

Process termination

- zombie state
 - Kernel has to keep some information
 - PID, termination status, and CPU usage time
 - This information is available when parent calls wait().
 - the process that has terminated, but whose parent has not yet waited for it, is called a zombie.
- If the parent terminates before the child?
 - init process becomes the parent of the orphaned process.
- If an orphaned process is terminated?
 - init calls wait() to fetch the termination status.

- A process that calls wait() or waitpid()
 - Block, if all of its children are still running.
 - Return immediately if a child has terminated.
 - Return immediately with an error, if it doesn't have any child processes.



```
#include <sys/wait.h>

pid_t wait(int *statloc);

Both return: process ID if OK, 0 (see later), or -1 on error
```

wait()

- If a child has already terminated and is a zombie, wait returns immediately with that child's status.
- Otherwise, it blocks the caller until a child terminates.
- If the caller blocks and has multiple children, wait returns when one terminates.

statloc argument

Store termination status in the location pointed to by staloc.

- Macro to examine the termination status
 - WIFEXITED(status)
 - 정상적으로 종료한 경우에 참 값을 리턴
 - WEXITEDSTATUS(status)
 - exit 함수의 인자에서 하위 8 비트 값을 리턴
 - WIFSIGNALED(status)
 - 시그널에 의해 종료한 경우에 참 값을 리턴
 - WTERMSIG(status)
 - 시그널 번호를 리턴
 - WIFSTOPPED(status)
 - 실행이 일시 중단된 경우에 참 값을 리턴
 - WSTOPSIG(status)
 - 실행을 일시 중단시킨 시그널 번호를 리턴

Example

```
#include <sys/types.h>
#include <sys/wait.h>
void pr_exit(int status)
    if (WIFEXITED(status))
         printf("normal termination, exit status = %d\n",
              WEXITSTATUS(status));
    else if (WIFSIGNALED(status))
         printf("abnormal termination, signal number = %d\n",
              WTERMSIG(status));
    else if (WIFSTOPPED(status))
         printf("child stopped, signal number = %d\n",
              WSTOPSIG(status));
```

Example(cont.)

```
#include <sys/types.h>
#include <sys/wait.h>
int main(void)
     pid_t pid;
           status;
     int
                                 F
     if ((pid = fork()) < 0)
          printf("fork error");
     else if (pid == 0)
                                              /* child */
          exit(7);
     if (wait(&status) != pid)
                                              /* wait for child */
          printf("wait error");
                                              /* and print its status */
     pr_exit(status);
```

Example(cont.)

```
if (\text{pid} = \text{fork}()) < 0)
      printf("fork error");
else if (pid == 0)
                               /* child */
                               /* generates SIGABRT */
      abort();
if (wait(&status) != pid)
                               /* wait for child */
      printf("wait error");
pr_exit(status);
                                /* and print its status */
if (\text{pid} = \text{fork}()) < 0)
      printf("fork error");
else if (pid == 0)
                               /* child */
      status = 0;
                               /* divide by 0 generates SIGFPE */
                               /* wait for child */
if (wait(&status) != pid)
      printf("wait error");
pr_exit(status);
                                /* and print its status */
```



◉ 실행

```
$ ./a.out
normal termination, exit status = 7
abnormal termination, signal number = 6
abnormal termination, signal number = 8
```

Each signal number is defined in <signal.h>

SIGABRT: 6 SIGFPE: 8

- Why waitpid() is required?
 - If we have more than one child, wait() returns on termination of any of the children.
 - What if we want to wait for a specific process to terminate?
- waitpid()
 - Waits for a <u>specific</u> process with *pid* argument.
 - Provides some controls with options argument.

pid argument

- pid == -1
 - Waits for any child process.
 - In this case, waitpid() is equivalent to wait().
- pid > 0
 - Waits for the child whose process ID equals pid.
- pid == 0
 - Waits for any child whose process group ID equals that of the calling process.
- pid < -1</p>
 - Waits for any child whose process group ID equals the absolute value of pid.

options argument

- WNOHANG
 - Not block(return immediately) if a child specified by pid is not terminated.
- WUNTRACED
 - Not block if a child specified by pid has stopped.

wait() and waitpid()

Example

```
#include "apue.h"
#include <sys/wait.h>
int main(void)
  pid_t pid;
  if ((pid = fork()) < 0) {
     err_sys("fork error");
  } else if (pid == 0) {
                                /* first child */
    if ((pid = fork()) < 0)
       err_sys("fork error");
     else if (pid > 0)
                                  /* parent from second fork == first child */
       exit(0);
    /* We're the second child; our parent becomes init. */
     sleep(2);
     printf("second child, parent pid = %d\n", getppid());
     exit(0);
```

wait() and waitpid()

Example

```
if (waitpid(pid, NULL, 0) != pid) /* wait for first child */
    err_sys("waitpid error");

/*
    * We're the parent (the original process); we continue executing,
    * knowing that we're not the parent of the second child.
    */
    exit(0);
}
```

連 실행

```
$ ./a.out
$ second child, parent pid = 1
```

Execute a program.

- When a process calls one of the exec functions, the process is completely replaced by the new program.
- New program starts executing at its main() function.
- exec() merely replaces the current process(text, data) with a brand new program from disk.

- pathname(filename) argument
 - pathname(filename) of a file to be executed.
 - execl, execv, execle, execve take a pathname argument.
 - execlp, execvp take a filename argument.
 - If filename contains a '/', it is taken as a pathname.
 - Otherwise, the executable file is searched for in the directories specified by PATH.
 - E.g. PATH=/bin:/usr/bin:/usr/local/bin/:.

Argument list

- arg0, arg1, ..., argn in the execl, execlp, and execle.
- The list of arguments is terminated by a NULL pointer.
 - execl("/bin/ls", "ls", "-al", 0);

Argument vector

- *argv[] in execv, execve and execvp.
- The array of pointers is terminated by a NULL pointer.
 - char *argv[3] = {"ls", "-al", 0};
 - execv("/bin/ls", argv);

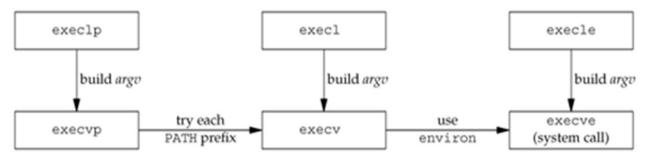
Environment variable

- *envp[] in execve and execle
- A pointer to an array of pointers to the environment strings.
- The other functions use the environ variable.
 - char *env[2] = {"TERM=vt100", 0};
 - execle("/bin/ls", "ls", 0, env);

Differences among the six exec functions

| function | pathname | filename | arg list | argv[] | environ | envp[] |
|------------------|----------|----------|----------|--------|---------|--------|
| execl | 0 | | 0 | | 0 | |
| execlp | | 0 | 0 | | 0 | |
| execle | 0 | | 0 | | | 0 |
| execv | 0 | | | 0 | 0 | |
| execvp | | 0 | | 0 | 0 | |
| execve | 0 | | | 0 | | 0 |
| (letter in name) | | p | l | ٧ | | e |

- execve is a system call.
 - execl, execv, execle, execlp, execvp are library functions.
 - Relationship of the six exec functions.



Example

```
#include "apue.h"
#include <sys/wait.h>
char
      *env_init[] = { "USER=unknown", "PATH=/tmp", NULL };
int
main(void)
  pid_t pid;
  if ((pid = fork()) < 0) {
    err_sys("fork error");
  } else if (pid == 0) { /* specify pathname, specify environment */
    if (execle("/home/sar/bin/echoall", "echoall", "myarg1",
         "MY ARG2", (char *)0, env_init) < 0)
       err_sys("execle error");
```

Example(cont.)

```
if (waitpid(pid, NULL, 0) < 0)
    err_sys("wait error");

if ((pid = fork()) < 0) {
    err_sys("fork error");
} else if (pid == 0) { /* specify filename, inherit environment */
    if (execlp("echoall", "echoall", "only 1 arg", (char *)0) < 0)
        err_sys("execlp error");
}

exit(0);
}</pre>
```

echoall

```
#include "apue.h"
int
main(int argc, char *argv[])
  int
           **ptr;
  char
  extern char **environ;
                                            /* echo all command-line args */
  for (i = 0; i < argc; i++)
    printf("argv[%d]: %s\n", i, argv[i]);
  for (ptr = environ; *ptr != 0; ptr++) /* and all env strings */
    printf("%s\n", *ptr);
  exit(0);
```

■ 실행

```
$./a.out
argv[0]: echoall
argv[1]: myarg1
argv[2]: MY ARG2
USER=unknown
PATH=/tmp
$ argv[0]: echoall
argv[1]: only 1 arg
USER=sar
LOGNAME=sar
SHELL=/bin/bash
HOME=/home/sar
```