



GRIFFITH COLLEGE DUBLIN

Assignment Cover Sheet

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Signed: Yen Lung Chen

Date: 26/04/2018

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Consider the following schema:

Supplier(sid: integer, sname: string, address: string)

Part(pid: integer, pname: string, colour: string)

Catalog(sid: integer, pid: integer, cost: real)

The relation Supplier stores suppliers and the key of that relation is sid. The relation Part stores parts, and pid is the key of that relation. Finally, Catalogue stores which supplier supplies which part at which cost. The key is the combination of the two attributes sid and pid.

1. Write queries in relational algebra 3%

Supplier	Catalog	Part
<u>sid</u>	<u>sid</u>	<u>pid</u>
sname	pid	pname
address	cost	colour

Write the following queries in relational algebra.

1. Find the names of suppliers who supply some red part.

Answer:

$\Pi_{\text{sname}}(\sigma_{\text{colour}=\text{"red"}}(\text{Supplier} \bowtie \text{Part} \bowtie \text{Catalog}))$

Duplicate rows are automatically eliminated, as relation is a set.

Or the queries in the relational algebra could be

$\Pi_{\text{sname}}(\Pi_{\text{sid}}(((\sigma_{\text{colour}=\text{"red"}} \text{ Part}) \bowtie \text{Catalog}) \bowtie \text{Supplier}))$

2. Find the IDs of suppliers who supply some red or green part.

$(\sigma_{\text{colour}=\text{"red"} \vee \text{colour}=\text{"green"}}\text{Part})$
Pid
Pname
Colour

Catalog
<u>sid</u>
pid
cost

$((\sigma_{\text{colour}=\text{"red"} \vee \text{colour}=\text{"green"}}\text{Part}) \bowtie \text{Catalog})$
Pid
sid
Pname
Colour
Cost

Answer: $\Pi_{\text{sid}}((\sigma_{\text{colour}=\text{"red"} \vee \text{colour}=\text{"green"}}\text{Part}) \bowtie \text{Catalog})$

3. Find the IDs of suppliers who supply some red part or are based at 21 George Street.

$\Pi_{\text{pid}}(\sigma_{\text{colour}=\text{"red"}}\text{Part})$
Pid



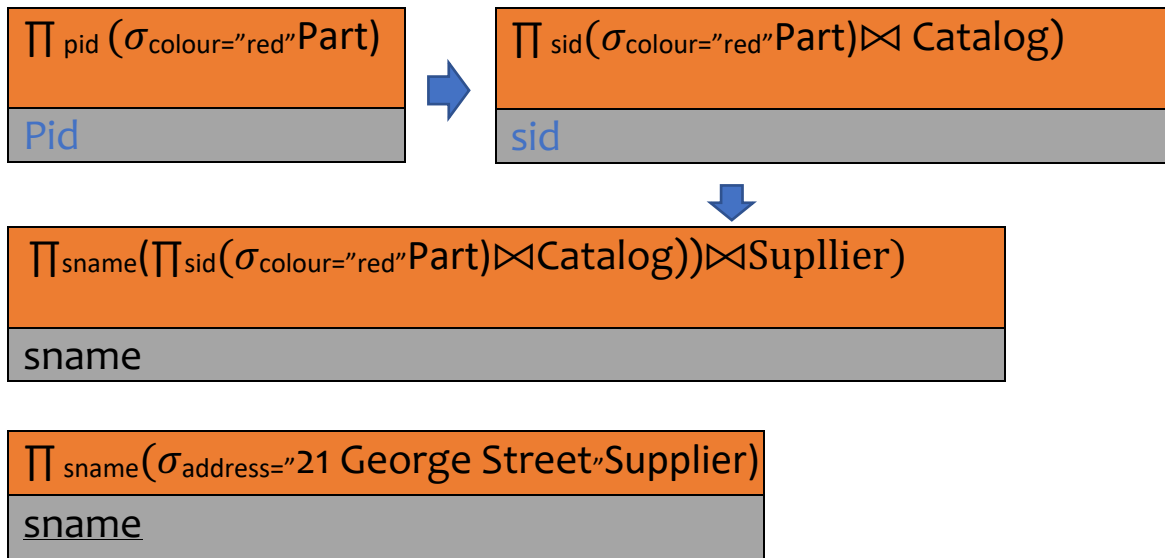
$\Pi_{\text{sid}}(\Pi_{\text{pid}}(\sigma_{\text{colour}=\text{"red"}}\text{Part}) \bowtie \text{Catalog})$
sid

$\Pi_{\text{sid}}(\sigma_{\text{address}=\text{"21 George Street"}}\text{Supplier})$
<u>sid</u>

Answer:

$\Pi_{\text{sid}}(\Pi_{\text{pid}}(\sigma_{\text{colour}=\text{"red"}}\text{Part}) \bowtie \text{Catalog}) \cup (\Pi_{\text{sid}}(\sigma_{\text{address}=\text{"21 George Street"}}\text{Supplier}))$

4. Find the names of suppliers who supply some red part or are based at 21 George Street.

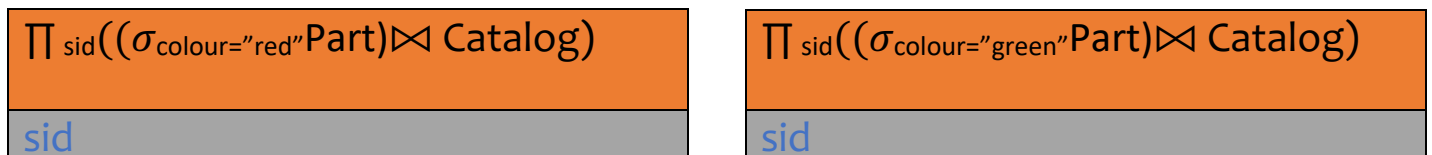


Answer:

$\Pi_{sname}(\Pi_{sid}(\sigma_{colour="red"}Part) \bowtie Catalog) \bowtie Supplier \cup \Pi_{sname}(\sigma_{address="21 George Street"}Supplier)$

5. Find the IDs of suppliers who supply some red part and some green part.

(Hint: use intersection of relations or join the same relation several times)



Answer:

$\Pi_{sid}((\sigma_{colour="red"}Part) \bowtie Catalog) \cap \Pi_{sid}((\sigma_{colour="green"}Part) \bowtie Catalog)$

6. Find pairs of IDs such that the supplier with the first ID charges more for some part than the supplier with the second ID. (Hint: you may want to create temporary relations to get two copies of Catalog)

- 1) Create two copies of Catalog and use Cartesian Product to compare each row of first table to each row of table 2
(Catalog 1 X Catalog 2)
- 2) Only compare the same product in each tuple and different sid
($\sigma_{\text{Catalog1.pid}=\text{Catalog2.pid} \wedge \text{Catalog1.sid} \neq \text{Catalog2.sid}}$ (Catalog 1 X Catalog 2))
- 3) Compare cost

($\sigma_{\text{Catalog1.pid}=\text{Catalog2.pid} \wedge \text{Catalog1.sid} \neq \text{Catalog2.sid} \wedge \text{Catalog1.cost} > \text{Catalog2.cost}}$ (Catalog 1 X Catalog 2))

4) Project the sid

5) Answer:

$\pi_{\text{Catalog1.sid, Catalog2.sid}}(\sigma_{\text{Catalog1.pid}=\text{Catalog2.pid} \wedge \text{Catalog1.sid} \neq \text{Catalog2.sid} \wedge \text{Catalog1.cost} > \text{Catalog2.cost}}(\text{Catalog 1 X Catalog 2}))$

7. Find the IDs of suppliers who supply only red parts. (Hint: A supplier supplies only red parts if it is not the case that the supplier offers a part that is not red. This question is a challenge!)

1) Find the sid supply part which is not red :

$\pi_{\text{sid}}((\sigma_{\text{colour} \neq \text{"red"}} \text{Part}) \bowtie \text{Catalog})$

2) A supplier supplies only red parts if it is not the case that the supplier offers a part that is not red.

3) Answer:

$\pi_{\text{sid}} \text{Supplier} - \pi_{\text{sid}}((\sigma_{\text{colour} \neq \text{"red"}} \text{Part}) \bowtie \text{Catalog})$

8. Find the IDs of suppliers who supply every part. (Hint: A supplier supplies every part if it is not the case that there is some part which they do not supply. Use set difference and cartesian product. This question is a challenge, too!)

1) Create a relation indicate each supplier ID supply all the part:

$$\pi_{sid, Supplier} X \pi_{pid} Part$$

2) Use division operator to find who actually supply every part:

$$\pi_{sid, pid}(Catalog) / \pi_{sid, Supplier} X \pi_{pid} Part$$

3) Find IDs of suppliers who supply every part:

Answer:

$$\pi_{sid}(\pi_{sid, pid}(Catalog) / \pi_{sid, pid}(Supplier X Part))$$

2. Queries in relational algebra, what do they mean? 3%

For each of the following relational algebra queries, say in English what they mean.

1. $\pi_{\text{name}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier})$

Answer: Find the name of supplier who supply some red part less than cost 100

2. $\pi_{\text{name}}(\pi_{\text{sid}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog})) \times \text{Supplier})$

Answer: Find the name of supplier who supply some red part less than cost 100

3. $\pi_{\text{name}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier}) \cap \pi_{\text{name}}(\sigma_{\text{colour}=\text{"green"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier})$

$\pi_{\text{name}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier})$:
set of name in suppliers supplying some red part less than cost 100

$\pi_{\text{name}}(\sigma_{\text{colour}=\text{"green"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier})$:
set of name in suppliers supplying some green part less than cost 100

Answer: Find the name of supplier who supply some red part and the supplier who supply green part less than cost 100

Note: two distinct suppliers with the same name may be returned by this query even if the first doesn't supply a green part, and the second doesn't supply a red part.

4. $\pi_{\text{sid}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier}) \cap \pi_{\text{sid}}(\sigma_{\text{colour}=\text{"green"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier})$

Answer: Find the ID of supplier who supply some red part and green part both less than cost 100

5. $\pi_{\text{name}}(\pi_{\text{sid}, \text{name}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier}) \cap \pi_{\text{sid}, \text{name}}(\sigma_{\text{colour}=\text{"green"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier}))$

$\pi_{\text{name}}(\pi_{\text{sid}, \text{name}}(\sigma_{\text{colour}=\text{"red"}}(\text{Part}) \times \sigma_{\text{cost}<100}(\text{Catalog}) \times \text{Supplier}))$:
Project the name of suppliers from the table contains two attributes: sid and sname, indicate from which red part less than cost 100 supplied from.

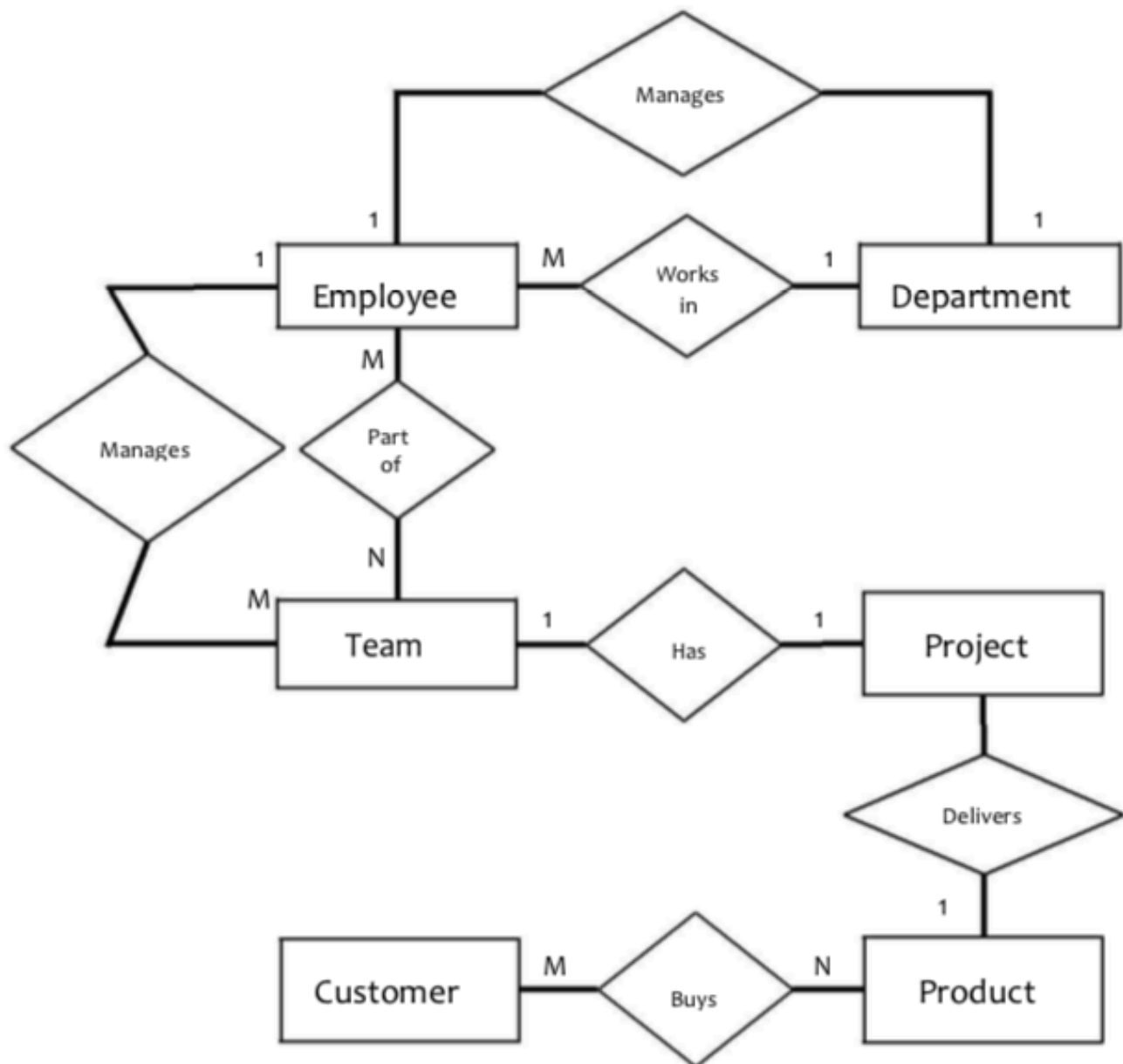
$\pi_{sid,sname}(\sigma_{colour="green"}(Part) \times \sigma_{cost < 100}(Catalog) \times Supplier)$: Project the name and ID of suppliers from the joined table indicates from which green part less than cost 100 supplied from.

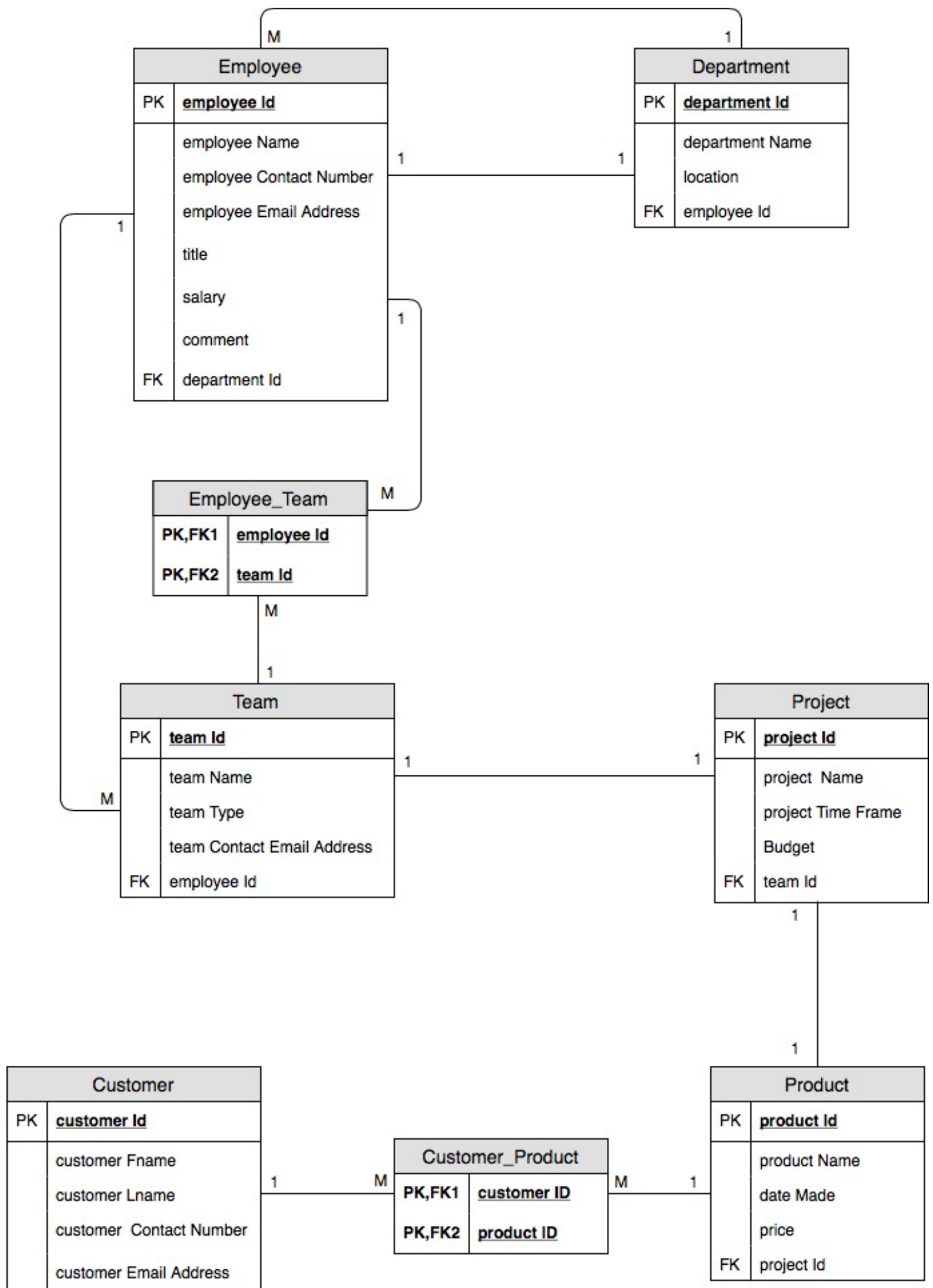
Answer:

Find the name of suppliers who supply some red part and some green part for less than 100.

3. Mapping & Normalisation 9%

a) Mapping the following Entity Relationship Diagram into relational schema. Choose appropriate attributes for each entity type. 5%





b) Database normalization – consider the following relation: 5%

Employee (emp-id, frst-name, last-name, address, phone, role-id, role-name)

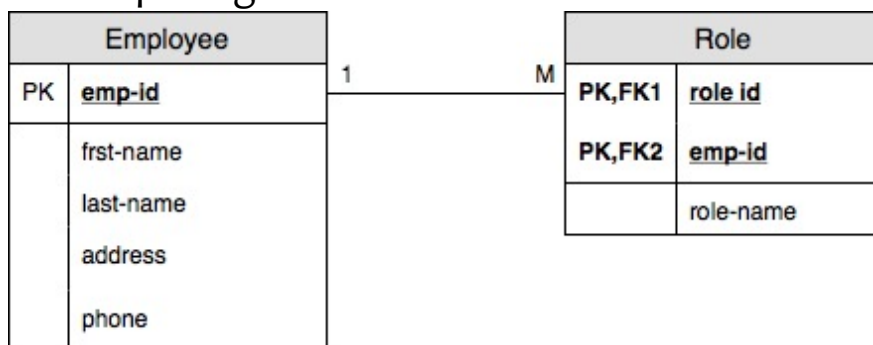
Let's assume each employee has one address and one phone number but each employee within the company is given more than one roles.

For example an employee may take responsibility of providing telephone support and face to face support.

- State whether the given table structure is in 1NF. If not, make appropriate changes so that it is in 1NF.

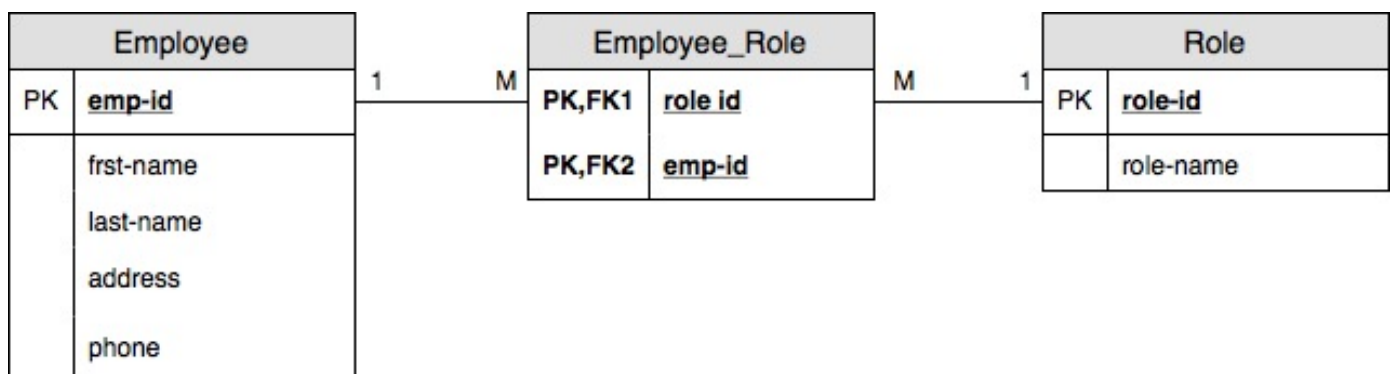
No, the given table structure is not in 1NF. A relation is in 1NF if, and only if, it contains no repeating attributes or groups of attributes. In other words each of the column in a table should contain one value and just one value.

Decomposing:



- State whether the given table structure is in 2NF. If not, make appropriate changes so that it is in 2NF.

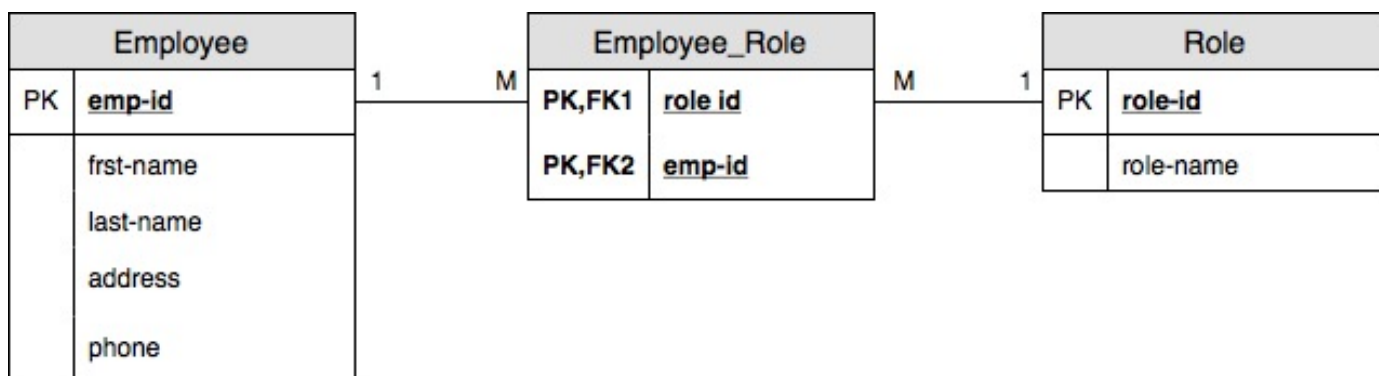
No, the given table structure is not in 2NF. No non-key fields may depend on a portion of the primary key in 2NF. Check all of the non-key attributes against each part of the key attributes to ensure they are functionally dependent on it.



- State whether the given table structure is in 3NF. If not, make appropriate changes so that it is in 3NF.

No, the given table structure is not in 3NF.

Third Normal Form (3NF): No fields may depend on other non-key fields. In other words, each field in a record should contain information about the entity that is defined by the primary key.



4. SQL 5%

Run the SQL – DDL and DML statements provided in assignment2__DDL_DML.sql file (on Moodle). This will create 3 tables and populate them with data.

As the file is in .sql format, to save time you can copy and paste when executing these statements.

a) Read the statements when you execute them so you understand the tables, columns, and their relationships. Once you have run all the statements successfully, now create and execute the following queries:

```

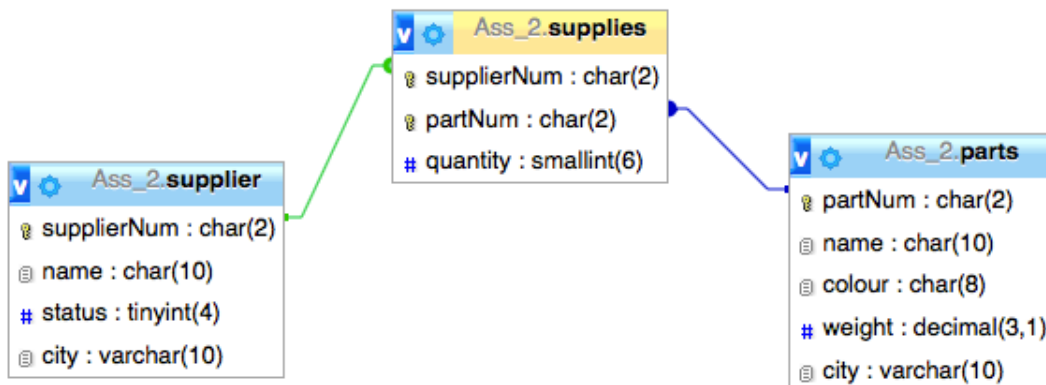
mysql> use Ass_2;
Reading table information for completion of table and column names
You can turn off this feature to get a quicker startup with -A

Database changed
mysql> show databases;
+-----+
| Database |
+-----+
| information_schema |
| Ass_2 |
| animals |
| c9 |
| mysql |
| performance_schema |
| phpmyadmin |
| registration-schema |
| sakila |
+-----+
9 rows in set (0.00 sec)

mysql> use Ass_2;
Database changed
mysql>

```

b) How many relationships exist between these tables? Specify their type & cardinality.



```
mysql> describe parts;
```

Field	Type	Null	Key	Default	Extra
partNum	char(2)	NO	PRI	NULL	
name	char(10)	NO		NULL	
colour	char(8)	NO		NULL	
weight	decimal(3,1)	NO		NULL	
city	varchar(10)	NO		NULL	

5 rows in set (0.00 sec)

```
mysql> describe supplier;
```

Field	Type	Null	Key	Default	Extra
supplierNum	char(2)	NO	PRI	NULL	
name	char(10)	NO		NULL	
status	tinyint(4)	NO		NULL	
city	varchar(10)	NO		NULL	

4 rows in set (0.00 sec)

```
mysql> describe supplies;
```

Field	Type	Null	Key	Default	Extra
supplierNum	char(2)	NO	PRI	NULL	
partNum	char(2)	NO	PRI	NULL	
quantity	smallint(6)	NO		NULL	

3 rows in set (0.00 sec)

c) List all the records in supplier, parts, & supplies tables. One table at a time.

```
mysql> select* from parts;
```

partNum	name	colour	weight	city
P1	Nut	Red	12.0	London
P2	Bolt	Green	17.0	Paris
P3	Screw	Blue	17.0	Oslo
P4	Screw	Red	14.0	London
P5	Cam	Blue	12.0	Paris
P6	Cog	Red	19.0	London

6 rows in set (0.00 sec)

```
mysql> select* from supplier;
```

supplierNum	name	status	city
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	20	Paris
S4	Clark	20	London
S5	Adams	30	Athens

5 rows in set (0.00 sec)

```
mysql> select *from supplies;
```

supplierNum	partNum	quantity
S1	P1	300
S1	P2	200
S1	P3	400
S1	P4	200
S1	P5	100
S1	P6	100
S2	P1	300
S2	P2	400
S3	P2	200
S4	P2	200
S4	P4	300
S4	P5	400

12 rows in set (0.01 sec)

d) Update the supplier table to reflect a change in supplier's status to the value 20 for all suppliers from Paris.

```
mysql> update supplier set status=20 where city="Paris";
Query OK, 1 row affected (0.01 sec)
Rows matched: 2  Changed: 1  Warnings: 0
```

```
mysql> select * from supplier;
```

supplierNum	name	status	city
S1	Smith	20	London
S2	Jones	20	Paris
S3	Blake	20	Paris
S4	Clark	20	London
S5	Adams	30	Athens

```
5 rows in set (0.01 sec)
```

e) Show the number of suppliers in each city ordered from highest to lowest.

```
mysql> select city, count(city) from supplier group by city order by count(city) desc;
```

city	count(city)
Paris	2
London	2
Athens	1

```
3 rows in set (0.00 sec)
```

f) List only the name of all the parts except the Red parts.

```
mysql> select name from parts
-> where colour != "Red";
```

name
Bolt
Screw
Cam

```
3 rows in set (0.00 sec)
```

g) Show all entries from the supplies table with their corresponding part names and supplier names. Rename the columns to appropriate ones.

```
mysql> select supplier.name Sname, supplies.supplierNum, supplies.partNum, parts.name Pname, quantity from supplier join supplies on supplier.supplierNum=supplies.supplierNum join parts on supplies.partNum=parts.partNum;
```

Sname	supplierNum	partNum	Pname	quantity
Smith	S1	P1	Nut	300
Smith	S1	P2	Bolt	200
Smith	S1	P3	Screw	400
Smith	S1	P4	Screw	200
Smith	S1	P5	Cam	100
Smith	S1	P6	Cog	100
Jones	S2	P1	Nut	300
Jones	S2	P2	Bolt	400
Blake	S3	P2	Bolt	200
Clark	S4	P2	Bolt	200
Clark	S4	P4	Screw	300
Clark	S4	P5	Cam	400

```
12 rows in set (0.00 sec)
```

h) Show the names of all suppliers that appear more than once in the supplies table.


```
mysql> select name from supplier join supplies on supplier.supplierNum=supplies.supplierNum group by supplies.supplierNum having count(supplies.supplierNum)>1;
```

name
Smith
Jones
Clark

```
3 rows in set (0.00 sec)
```

i) Supplier with supplierNum = S4 has closed down his business. Delete all the records related to this supplier from all relevant tables.

Note: MySQL provides a more effective way called `ON DELETE CASCADE` referential action for a [foreign key](#) that allows you to delete data from child tables automatically when you delete the data from the parent table. For this assignment the original database file is altered in .txt and imported to myphpadmin. In addition, the action of this query will affect the following two queries due to deletion of supplierNum = S4.

Shows the CREATE TABLE statement that creates the named table to make sure `ON DELETE CASCADE` clause at the end of the foreign key constraint definition.

```
CREATE TABLE `supplies` (
  `supplierNum` char(2) NOT NULL,
  `partNum` char(2) NOT NULL,
  `quantity` smallint(6) NOT NULL,
  PRIMARY KEY (`supplierNum`,`partNum`),
  KEY `partNum` (`partNum`),
  CONSTRAINT `supplies_ibfk_1` FOREIGN KEY (`supplierNum`) REFERENCES `supplier` (`supplierNum`) ON DELETE CASCADE,
  CONSTRAINT `supplies_ibfk_2` FOREIGN KEY (`partNum`) REFERENCES `parts` (`partNum`) ON DELETE CASCADE
) ENGINE=InnoDB DEFAULT CHARSET=latin1
```

Delete:

```
mysql> delete from supplier where supplierNum="S4";
Query OK, 1 row affected (0.00 sec)
```

```
mysql> select* from supplier;
```

supplierNum	name	status	city
S1	Smith	20	London
S2	Jones	10	Paris
S3	Blake	20	Paris
S5	Adams	30	Athens

```
4 rows in set (0.00 sec)
```

Check

```
mysql> select* from supplies;
```

supplierNum	partNum	quantity
S1	P1	300
S1	P2	200
S1	P3	400
S1	P4	200
S1	P5	100
S1	P6	100
S2	P1	300
S2	P2	400
S3	P2	200

```
9 rows in set (0.00 sec)
```

j) List all the parts except those with a quantity of 200.

Note: assuming exclude the total quantity of 200 of each parts

1) Join two tables

```
mysql> select * from supplies join parts on supplies.partNum=parts.partNum;
```

supplierNum	partNum	quantity	partNum	name	colour	weight	city
S1	P1	300	P1	Nut	Red	12.0	London
S2	P1	300	P1	Nut	Red	12.0	London
S1	P2	200	P2	Bolt	Green	17.0	Paris
S2	P2	400	P2	Bolt	Green	17.0	Paris
S3	P2	200	P2	Bolt	Green	17.0	Paris
S1	P3	400	P3	Screw	Blue	17.0	Oslo
S1	P4	200	P4	Screw	Red	14.0	London
S1	P5	100	P5	Cam	Blue	12.0	Paris
S1	P6	100	P6	Cog	Red	19.0	London

9 rows in set (0.00 sec)

2) Sum and group

```
mysql> select parts.name,supplies.partNum, sum(quantity) from supplies join parts on supplies.partNum=parts.partNum group by supplies.partNum;
```

name	partNum	sum(quantity)
Nut	P1	600
Bolt	P2	800
Screw	P3	400
Screw	P4	200
Cam	P5	100
Cog	P6	100

6 rows in set (0.00 sec)

3) List all the parts except those with a quantity of 200.

```
mysql> select parts.name,supplies.partNum, sum(quantity) from supplies join parts on supplies.partNum=parts.partNum group by supplies.partNum having sum(quantity)!="200";
```

name	partNum	sum(quantity)
Nut	P1	600
Bolt	P2	800
Screw	P3	400
Cam	P5	100
Cog	P6	100

5 rows in set (0.00 sec)

k) List part names, their colour, and the supplier(s) who supply them.

1) Join three tables required

```
mysql> select * from supplier join supplies on supplier.supplierNum=supplies.supplierNum join parts on supplies.partNum=parts.partNum;
```

supplierNum	name	status	city	supplierNum	partNum	quantity	partNum	name	colour	weight	city
S1	Smith	20	London	S1	P1	300	P1	Nut	Red	12.0	London
S1	Smith	20	London	S1	P2	200	P2	Bolt	Green	17.0	Paris
S1	Smith	20	London	S1	P3	400	P3	Screw	Blue	17.0	Oslo
S1	Smith	20	London	S1	P4	200	P4	Screw	Red	14.0	London
S1	Smith	20	London	S1	P5	100	P5	Cam	Blue	12.0	Paris
S1	Smith	20	London	S1	P6	100	P6	Cog	Red	19.0	London
S2	Jones	10	Paris	S2	P1	300	P1	Nut	Red	12.0	London
S2	Jones	10	Paris	S2	P2	400	P2	Bolt	Green	17.0	Paris
S3	Blake	20	Paris	S3	P2	200	P2	Bolt	Green	17.0	Paris

9 rows in set (0.00 sec)

2) List columns

```
mysql> select supplier.name, parts.name, parts.colour from supplier join supplies on supplier.supplierNum=supplies.supplierNum join parts on supplies.partNum=parts.partNum;
```

name	name	colour
Smith	Nut	Red
Smith	Bolt	Green
Smith	Screw	Blue
Smith	Screw	Red
Smith	Cam	Blue
Smith	Cog	Red
Jones	Nut	Red
Jones	Bolt	Green
Blake	Bolt	Green

9 rows in set (0.00 sec)