CS246 (Winter 2014) Mining Massive Data Sets

Introduction to proof techniques

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Disclaimer These notes may contain typos, mistakes or confusing points. Please contact the author so that we can improve them for next year.

1 Proof by induction

The proof by induction is usually illustrated by falling dominoes. Make the first domino fall (initialization or base case) and if your dominos are placed in such a way that each domino makes the next one fall (inheritance or induction), then all dominos will be down.

1.1 Principle

Say you want to prove that a property holds for all integers. To do a proof by induction, you will:

- 1. Base case: Prove that the property holds for n=0 (possibly another value, it really depends on the cases).
- 2. Inheritance: Assume that the result holds for $n \ge 0$ and prove that the property would hold for n+1.

1.2 Example

1. Let us start with a simple example. Prove that:

$$\forall n \ge 1, \ \sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}.$$

Solution We proceed in two steps:

- Base case (n=1): the left hand side is 1, as is the right hand side so the property holds for n=1.
- Now, we assume that the property holds for a given $n \ge 1$, and we want to prove that it holds for n+1.

¹There is a stronger variant of the proof by induction where you assume that the property holds for any integer $k \leq n$.

We compute:

$$\begin{split} \sum_{i=1}^{n+1} i^2 &= \sum_{i=1}^n i^2 + (n+1)^2 \\ &= \frac{n(n+1)(2n+1)}{6} + (n+1)^2 \qquad \text{(Induction hypothesis)} \\ &= \frac{n+1}{6} \left(n(2n+1) + 6(n+1) \right) \\ &= \frac{n+1}{6} \underbrace{\left(2n^2 + 7n + 6 \right)}_{(n+2)(2n+3)} \\ &= \frac{(n+1)(n+2)(2n+3)}{6} \\ &= \frac{(n+1)((n+1)+1)(2(n+1)+1)}{6}, \end{split}$$

which proves that the property holds for n+1 and concludes the proof by induction.

2. A more complicated example²: let $n \ge 1$, and two families of vector (e_1, \ldots, e_n) , and (f_1, \ldots, f_{n+1}) such that:

$$\forall 1 \le j \le n+1, \ f_j = \sum_{i=1}^n \lambda_i^{(j)} e_i.$$

Prove that (f_1, \ldots, f_{n+1}) is a dependent³ family.

Solution Let us denote by H(n) the property that we want to prove by induction:

H(n) =For any families $f = (f_1, \ldots, f_{n+1})$ and $e = (e_1, \ldots, e_n)$ such that:

$$\forall 1 \le j \le n+1, \ f_j = \sum_{i=1}^n \lambda_i^{(j)} e_i,$$

then f is dependent.

We proceed in two steps:

• Base case (n = 1): we have:

$$\exists \lambda_1, \lambda_2, \ f_1 = \lambda_1 e_1, \ f_2 = \lambda_2 e_1.$$

If both λ_1 and λ_2 are zero, then the family is clearly dependent.

Otherwise, we compute:

$$\lambda_1 e_2 - \lambda_2 e_1 = 0,$$

which proves that the family (e_1, e_2) is dependent.

• Inheritance: Assume that H(n) holds and let us try to prove that H(n+1) holds. We distinguish two cases:

$$\exists (\lambda_1, \dots, \lambda_n) \neq 0, \ \sum_{i=1}^n \lambda_i v_i = 0.$$

²Note that the following result basically says that any set of n + 1 vectors in a vector space of dimension n is always dependent.

³Recall that (v_1, \ldots, v_n) is a (linearly) dependent family if and only if:

- (a) Forall $1 \le j \le n+2$, $\lambda_{n+1}^{(j)} = 0$. Then, we can rewrite all the f_j 's in terms of (e_1, \ldots, e_n) and therefore using H(n) we conclude that (f_1, \ldots, f_{n+1}) is dependent and so will be (f_1, \ldots, f_{n+2}) .
- (b) One of the $\lambda_{n+1}^{(j)}$ is non-zero. We can assume (by changing the numbering if necessary) that $\lambda_{n+1}^{(n+2)} \neq 0$. We then define:

$$\forall 1 \le j \le n+1, \ \tilde{f}_j = f_j - \frac{\lambda_{n+1}^{(j)}}{\lambda_{n+1}^{(n+2)}} f_{n+2}.$$

Then, the coefficient on e_{n+1} is zero for all the \tilde{f}_j 's, and therefore, we can apply H(n) to $(\tilde{f}_1, \ldots, \tilde{f}_{n+1})$ and (e_1, \ldots, e_n) , which gives us that $(\tilde{f}_1, \ldots, \tilde{f}_{n+1})$ is dependent, that is:

$$\exists (\beta_1, \dots, \beta_{n+1}) \neq 0, \ \beta_1 \tilde{f}_1 + \dots + \beta_{n+1} \tilde{f}_{n+1} = 0,$$

which can be rewritten:

$$\exists (\beta_1, \dots, \beta_{n+1}) \neq 0, \ \beta_1 f_1 + \dots + \beta_{n+1} f_{n+1} - \sum_{j=1}^{n+1} \frac{\beta_j \lambda_{n+1}^{(j)}}{\lambda_{n+1}^{(n+2)}} f_{n+2} = 0,$$

which by definition means that (f_1, \ldots, f_{n+2}) is dependent and concludes the inheritance part of the proof.

1.3 Common mistakes

• Forgetting the base case. Usually, the base case is easy compared to the induction step, but you still have to clearly mention that you looked at it.

2 Proof by contrapositive

2.1 Principle

Say you want to prove something of the form:

$$A \Rightarrow B$$
.

The usual way to prove this is to assume that A is true and then try to get to B. Sometimes it is easier to prove that:

$$\neg B \Rightarrow \neg A$$
.

2.2 Example

1. Prove that for any $a, b \in \mathbb{R}$:

a + b is irrational $\Rightarrow a$ or b is irrational.

Solution $\neg B$ in our case will be:

a and b are rational.

Now we want to prove $\neg A$, that is:

a + b is rational.

This is now straightforward: indeed, we can write:

$$\exists p, q, u, v \in \mathbb{N}, \ a = \frac{p}{q} \text{ and } b = \frac{u}{v}.$$

We then compute:

$$a+b = \frac{pv + qu}{qv},$$

which proves that a + b is rational and concludes the proof by contrapositive.

2. Prove that for $n \in \mathbb{N}$:

8 does not divide $n^2 - 1 \Rightarrow n$ is even.

Hint: Notice that any odd number n can be written as n = 4k + r with $k \in \mathbb{N}$ and $r \in \{1, 3\}$.

Solution The contrapositive is:

$$n \text{ is odd} \Rightarrow 8 \text{ divides } n^2 - 1.$$

Now, let us do the proof: let us take n an odd integer. By using the hint, we write n as 4k + r with $r \in \{1, 3\}$.

We compute:

$$n^2 - 1 = 16k^2 + 8rk + r^2 - 1.$$

Now by noticing that 8 divides $16k^2$, 8rk and r^2-1 (because this quantity is either 0 or 8), we conclude that 8 divides n^2-1 which is what we wanted.

2.3 Common mistakes

• Computing $\neg A$ and $\neg B$ may be tricky in some cases. You should spend some time practicing if you are not confident on this point.

For instance, try to find the contrapositive of:

$$\lim_{x \to a} f(x) = b \Rightarrow (\forall \epsilon > 0, \ \exists \alpha > 0, \ ||x - a|| \le \alpha \Rightarrow ||f(x) - b|| \le \epsilon)$$

Solution The contrapositive is:

$$(\exists \epsilon > 0, \ \forall \alpha > 0, \ ||x - a|| \le \alpha \Rightarrow ||f(x) - b|| \le \epsilon) \Rightarrow \lim_{x \to a} f(x) \ne b.$$

3 Proof by contradiction

3.1 Principle

A proof by contradiction works as follows: assume that the property you are trying to prove is wrong and find a contradiction.

3.2 Example

1. Prove that $\sqrt{2}$ is irrational.

Solution Assume by contradiction that $\sqrt{2}$ is rational, that is⁴:

$$\exists p, q \text{ such that } \gcd(p,q) = 1 \text{ and } \sqrt{2} = \frac{p}{q}.$$

Now, our goal is to find a contradiction.

We compute:

$$p^2 = 2q^2 \Rightarrow p^2$$
 is even $\Rightarrow p$ is even.

Therefore, we can write p = 2k and:

$$2 = \frac{(2k)^2}{q^2} \Leftrightarrow q^2 = 2k^2 \Rightarrow q \text{ is even.}$$

Because both p and q are even, we have $\gcd(p,q) \ge 2$, which contradicts our assumption that $\gcd(p,q) = 1$.

Therefore $\sqrt{2}$ is irrational.

2. Show that any function k-Lipschitz with k < 1 has at most one fixed point⁵.

Solution Let f be a k-Lipschitz function.

Assume by contradiction that f has 2 distinct fixed points that we call u and v with $v \neq u$. Then:

$$\underbrace{||f(u) - f(v)||}_{=||u-v||} \le k||u-v|| \Leftrightarrow (1-k)||u-v|| \le 0.$$

But 1 - k > 0 and ||u - v|| > 0 because u and v are distinct. This is a contradiction and therefore our initial assumption is wrong, that is f has at most one fixed point.

$$\forall x, y \in \mathbb{R}, \ ||f(x) - f(y)|| \le k||x - y||.$$

⁴The fact that we can assume gcd(p,q) = 1 comes from the fact that you can reduce a fraction to its minimal form.

⁵Recall that f is k-Lipschitz on \mathbb{R} if and only if: