RISC-V ARCHITECTURE TRAINING

@DEMO: PK (proxy kernel) and FESVR (front-end server)

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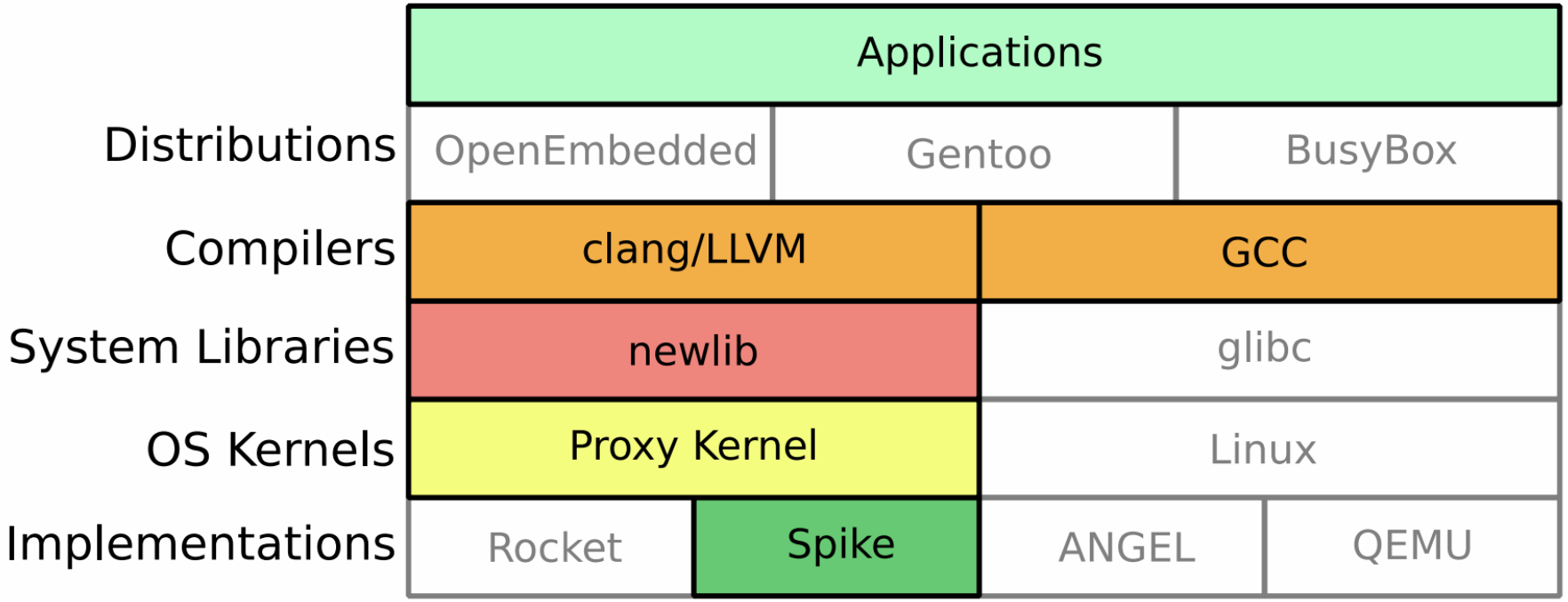
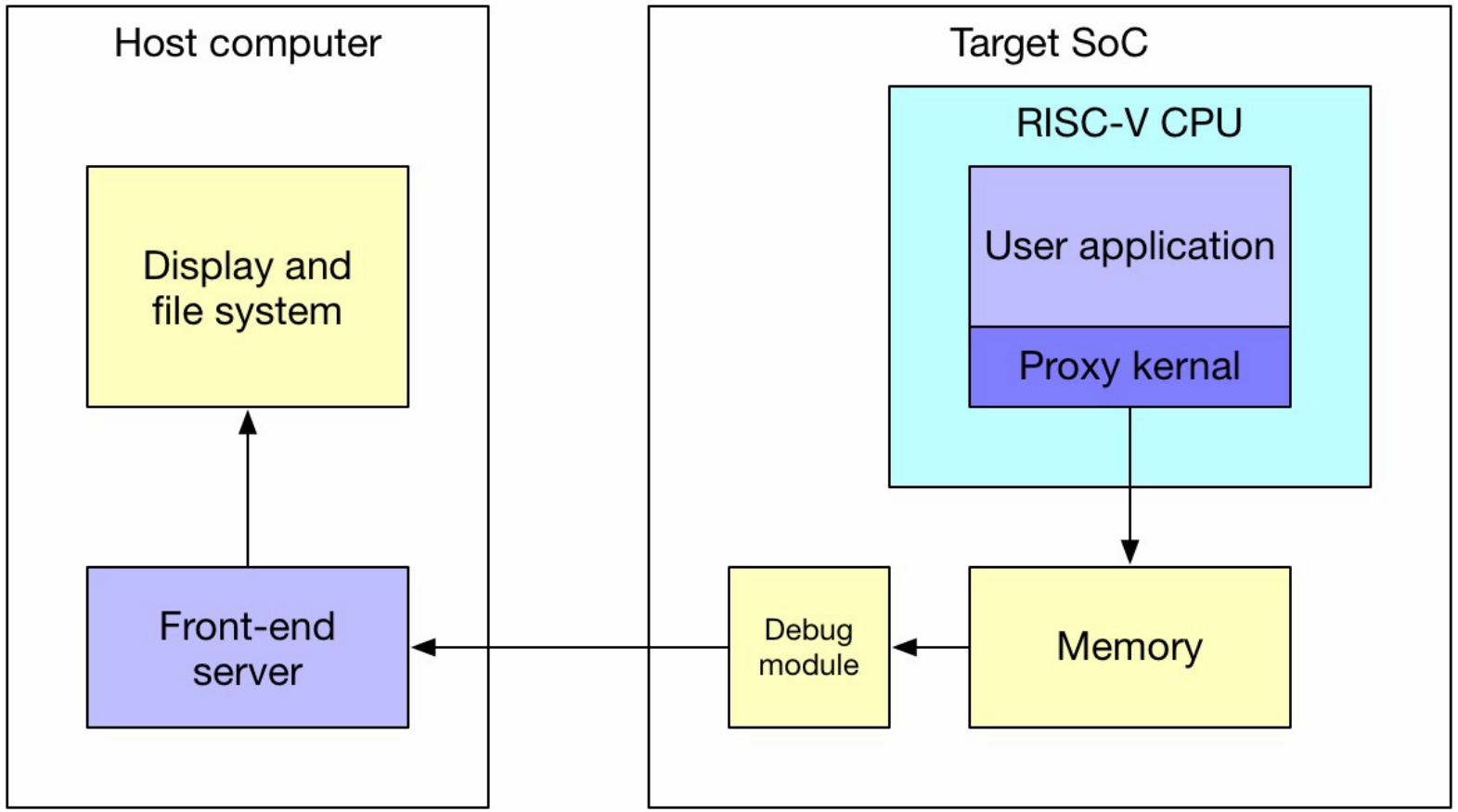
PK and FESVRV

- PK (proxy kernel) & FESVR (front-end server)
  - For debug and system bring up
  - PK is an abstraction of kernel that provides system services through FESVR running on host



PK and FESVR (cont'd)

- PK is running on target CPU, while FESVR is running on host computer



PK and FESVR / code example

`printf` in modified version of `spike`

```
# Original spike
> cd ~/riscv-training/lab/22-lab.system-call
> spike ../pk hello.elf
bbl loader
Hello world!

# Modified spike (print out all system call)
> cd ~/riscv-training/lab/22-lab.system-call
> ~/riscv-git/riscv-tools/riscv-isa-sim/build/spike ../pk hello.elf
bbl loader
fesvr::sys_getmainvars (0x 8000d860 200 0 0 0 0 0)
fesvr::sys_openat (0x ffffffffffffff9c 8000d8b5 a 0 0 0 0)
fesvr::sys_pread (0x 3 8000d600 40 0 0 0 0)
...
fesvr::sys_pread (0x 3 80830000 1000 a000 0 0 0)
fesvr::sys_fstat (0x 1 80024df0 0 0 0 0 0)
fesvr::sys_pread (0x 3 8082c000 1000 6000 0 0 0)
fesvr::sys_write (0x 1 80834230 d 0 0 0 0)
Hello world!
fesvr::sys_exit (0x 0 0 0 0 0 0 0)
```

PK and FESVR / system call

`tohost` and `fromhost`

- Memory location: shared knowledge between PK and FESVR
- Both are 32-bit size, that can be read/write with single access
- Follow producer-consumer model
  - `tohost` is written by PK, cleared by FESVR
  - `fromhost` is written by FESVR, cleared by PK

`magicmem`

- `tohost` and `fromhost` are too small to communicate real data structure. They only store the address of `magicmem`
- Syscall type, arguments and return values are stored in `magicmem`

Syscall entry point in PK

`~/riscv-git/riscv-tools/riscv-pk/pk/syscall.c`

Syscall handler in FESVR

`~/riscv-git/riscv-tools/riscv-isa-sim/fesvr/syscall.cc`

PK and FESVR / system call workflow

> Please remember, this is for debugging hardware and system bring-up, sometimes co-processor running enviroment. And it's good for understanding application/kernel interaction.

Target side (PK on RISC-V)	Host side (FESVR on x86)
User-level code: `ecall` and trap into PK	
Write syscall arguments into **magicmem**	
Write address of **magicmem** into **tohost**	Looping: read **tohost** until it's non-0
	Read **magicmem**
Looping: read **tohost** until it's 0	Write 0 to **tohost**
	Deal with syscall. Write return values into **magicmem**
Looping: read **fromhost** until it's non-0	Write address of **magicmem** into **fromhost**
Read **magicmem**	
Write 0 to **fromhost**	Looping: read **fromhost** until it's 0

PK and FESVR / verification exit

Another very useful scenario

- To pass exit code in verification. It's embedded inside the `riscv-tests` verification suite.
- `RVTEST_PASS` and `RVTEST_FAIL` in `~/riscv-git/riscv-tools/riscv-tests/env/p/riscv_test.h`

```
#define RVTEST_PASS                                                     \
        fence;                                                          \
        li TESTNUM, 1;                                                  \
        ecall

#define TESTNUM gp
#define RVTEST_FAIL                                                     \
        fence;                                                          \
1:      beqz TESTNUM, 1b;                                               \
        sll TESTNUM, TESTNUM, 1;                                        \
        or TESTNUM, TESTNUM, 1;                                         \
        ecall
```


감사합니다
Natick
Grazie
Danke
Ευχαριστίες
Dalu
Obrigado
Thank You
Köszönöm
Tack
Спасибо
Dank
Gracias
谢谢
Merci
Seé
ありがとう

@LAB: new system call

Add new system call to synchronize system time between target and host

Read CSR of `mtime` to get system time on RISC-V CPU

Need to use assembly code to get `mtime`