# HIGHT Algorithm Specification

2009.07



### 1. HIGHT

The HIGHT algorithm is a symmetric block cipher that can process data blocks of 64 bits, using a cipher key with length of 128 bits.

## 2. HIGHT encryption

The encryption operation is as shown in Figure 1. The transformation of a 64-bit block *P* into a 64-bit block *C* is defined as follows:

(1) 
$$P = P_7 /\!\!/ P_6 /\!\!/ P_5 /\!\!/ P_4 /\!\!/ P_3 /\!\!/ P_2 /\!\!/ P_1 /\!\!/ P_0$$
 (*P*<sub>i</sub> are plaintext bytes)

(2) 
$$X_{0,0} = P_0 \quad \boxplus \quad WK_0, \qquad X_{0,1} = P_1,$$

$$X_{0,2} = P_2 \oplus WK_1,$$
  $X_{0,3} = P_3,$ 

$$X_{0,4} = P_4 \boxplus WK_2, X_{0,5} = P_5,$$

$$X_{0,6} = P_6 \oplus WK_3,$$
  $X_{0,7} = P_{7.}$ 

(3) for i = 0 to 30:

$$X_{i+1,0} = X_{i,7} \oplus (F_0(X_{i,6}) \boxplus SK_{4i+3}),$$
  $X_{i+1,1} = X_{i,0},$   $X_{i+1,2} = X_{i,1} \boxplus (F_1(X_{i,0}) \oplus SK_{4i}),$   $X_{i+1,3} = X_{i,2},$   $X_{i+1,4} = X_{i,3} \oplus (F_0(X_{i,2}) \boxplus SK_{4i+1}),$   $X_{i+1,5} = X_{i,4},$   $X_{i+1,6} = X_{i,5} \boxplus (F_1(X_{i,4}) \oplus SK_{4i+2}),$   $X_{i+1,7} = X_{i,6}.$ 

for i = 31:

$$X_{i+1,0} = X_{i,0},$$
  $X_{i+1,1} = X_{i,1} \boxplus (F_1(X_{i,0}) \oplus SK_{124}),$   
 $X_{i+1,2} = X_{i,2},$   $X_{i+1,3} = X_{i,3} \oplus (F_0(X_{i,2}) \boxplus SK_{125}),$   
 $X_{i+1,4} = X_{i,4},$   $X_{i+1,5} = X_{i,5} \boxplus (F_1(X_{i,4}) \oplus SK_{126}),$   
 $X_{i+1,6} = X_{i,6},$   $X_{i+1,7} = X_{i,7} \oplus (F_0(X_{i,6}) \boxplus SK_{127}).$ 

(4) 
$$C_0 = X_{32,0} \oplus WK_4$$
,  $C_1 = X_{32,1}$ ,

$$C_2 = X_{32,2} \oplus WK_5,$$
  $C_3 = X_{32,3},$   $C_4 = X_{32,4} \boxplus WK_6,$   $C_5 = X_{32,5},$ 

$$C_6 = X_{32.6} \oplus WK_7$$
,  $C_7 = X_{32.7}$ 

(5) 
$$C = C_7 \parallel C_6 \parallel C_5 \parallel C_4 \parallel C_3 \parallel C_2 \parallel C_1 \parallel C_0$$
 ( $C_i$  are ciphertext bytes)

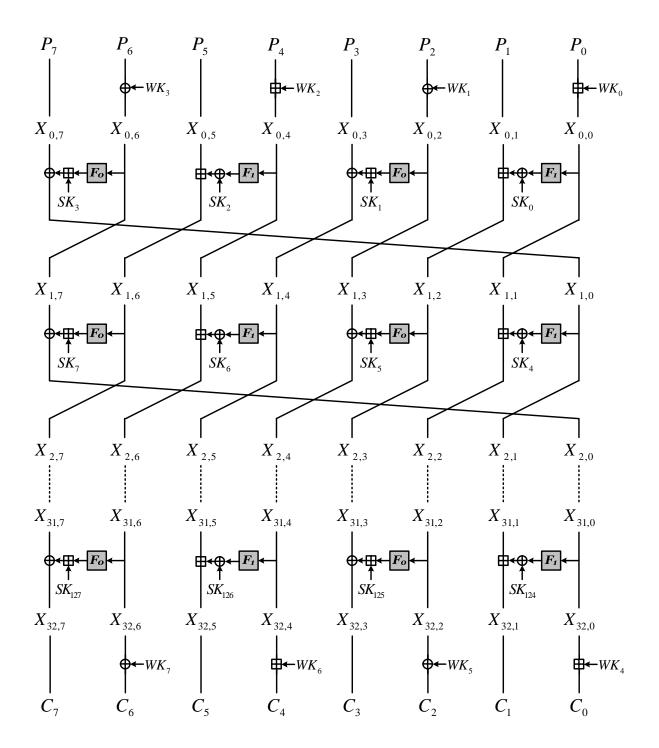


Figure 1. Encryption procedure of HIGHT

# 3. HIGHT decryption

The decryption operation is identical in operation to encryption apart from the following two modifications.

- (1) All  $ext{ } ext{ } ext$
- (2) The order in which the keys  $WK_i$  and  $SK_i$  are applied is reversed.

### 4. HIGHT functions

### 4.1 The functions $F_0$ and $F_1$

The HIGHT algorithm uses two functions, namely,  $F_0$  and  $F_1$  which are now defined.

### 4.2 Function F<sub>0</sub>

The  $F_0$  function is used for encryption and decryption. The function  $F_0$  is defined as follows:

$$F_0(x) = (x <<<_1) \oplus (x <<<_2) \oplus (x <<<_7)$$

### 4.3 Function F<sub>1</sub>

The  $F_1$  function is used for encryption and decryption. The function  $F_1$  is defined as follows:

$$F_1(x) = (x <<<_3) \oplus (x <<<_4) \oplus (x <<<_6)$$

# 5. HIGHT key schedule

The key scheduling part accepts a 128-bit master key  $K=K_{15} \parallel K_{14} \parallel \cdots \parallel K_0$  and yields 8 whitening key bytes  $WK_i$  and 128 subkey bytes  $SK_i$ , as shown below.

The generation of whitening keys is defined as follows.

for 
$$i = 0, 1, 2, 3$$
:  
 $WK_i = K_{i+12}$   
for  $i = 4, 5, 6, 7$ :  
 $WK_i = K_{i-4}$ 

The 128 subkeys are used for encryption and decryption, 4 subkeys per round. The generation of subkeys is defined as follows.

(1) 
$$s_0 = 0$$
,  $s_1 = 1$ ,  $s_2 = 0$ ,  $s_3 = 1$ ,  $s_4 = 1$ ,  $s_5 = 0$ ,  $s_6 = 1$ 

$$\delta_0 = s_6 \parallel s_5 \parallel s_4 \parallel s_3 \parallel s_2 \parallel s_1 \parallel s_0$$
(2) for  $i = 1$  to 127:

$$S_{i+6} = S_{i+2} \oplus S_{i-1}$$

$$\delta_i = \varsigma_{i+6} \parallel \varsigma_{i+5} \parallel \varsigma_{i+4} \parallel \varsigma_{i+3} \parallel \varsigma_{i+2} \parallel \varsigma_{i+1} \parallel \varsigma_i$$

(3) for 
$$i = 0$$
 to 7:

for 
$$j = 0$$
 to 7:

$$SK_{16 \cdot i + j} = K_{j \cdot i \mod 8} \quad \boxplus \quad \delta_{16 \cdot i + j}$$

for 
$$j = 0$$
 to 7:

$$SK_{16 \cdot i + j + 8} = K_{(j - i \mod 8) + 8} \boxplus \delta_{16 \cdot i + j + 8}$$