

Software Security Project - Report

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1 Introduction

Elf file is something really important and can be found almost everywhere so we will try to do this project to learn more about them and manipulate it using the C language. For that the goal is to inject a small piece of assembly code, modify the current section header and to finally be able to hijack the got entry allowing us to execute our code on a binary elf file.

Before starting the long way to inject our code we need to know the structure to the elf file to be able to know where we have to inject and what we have to modify.

The structure of the Elf file from the [elf man](#) is representing like that is the next one.

Name	Information
Executable Header	general information about the binary
Program Header (optional)	Provides the execution view, contains information about the file.
Sections	The code and data of programs
Section Header (optional)	Each denoting the property of its related section

2 Initialize ELF FILE for reading

I used the C function getopt to create an arguments parser so that the program must take 5 argument

- -r nameOfElfFile
- -b nameOfBinaryToBeInjected
- -c newSectionName
- -a AddressOfInjectedCode
- -e

The option -e will decide whether we want to overwrite the entry point or not.

Then we will use for the rest of the project the libelf to be able to get the information of an elf file. With the function getElf i will be able to get the elf file of the argument passed in the command line.

An elf file contains an executable header that can be obtained from the libelf function elf64_getehdr and if the result of e_ident[4] is 2 then that means that it is a 64bits executable.

3 Find the PT_NOTE segment header

The libelf function [elf_getphdrnum](#) will give us the number of program headers that can also be found in the executable header.

Then the program header can be obtained through the function [elf64_getphdr](#) and because the program header is an array of program headers we can iterate over this array and inspect the p_type to check if the name is PT_NOTE which is for auxilray information. then we will store the index when we found the first PT_NOTE.

4 Code injection

Now after having the PT_NOTE segment we will append our injected code to the elf binary for that as I previously said, we will create an assembly code and we will compile it using nasm which is an assembler for x86 CPU architecture.

- The code must be compiled with nasm assembler
- It must save the context and restore it before quitting since the call will break the current one
- We can't use any .data section since the date binary already has this section.
- the next time we will manipulate it will be in the Execute injected code section.

And now we have to append our code, this task is really simple, we just have to open a file descriptor for thoses 2file and use the function fseek to position ourself inside the specific file so we will fseek at the end of the elf file and do an fwrite to write the injected binary at the end of the elf file.

Now we have injected our binary but the Elf specification requires that the offset and address are congruent modulo 4096. To ensure correct alignment we will modify the address so that the difference with the file offset becomes zero modulo 4096. To do that we have to create a delta integer.

$$\sigma = 4096 - ((elfSize - argumentAddress) \mod 4096) \quad (1)$$

and replace our base address by adding the difference.
argumentAddress += 4096 - σ

Then our address is align with 4096 so that

$$((argumentAddress - elfSize) \mod 4096) \iff 0 \quad (2)$$

5 Overwriting the concerned section header

Now we have to get the number of section header in the .shstrtab section that can be get using the libelf function `elf_getshdrstrndx` then we will loop over all sections headers to find the .note.ABI-tag to do that the libelf function `elf_nextscn` can be use to get the section header descriptor to loop until we find NULL then we will use `elf64_getshdr` to get the section header depending of the scn with the section header we can get the name using `elf_strpt` and if we find the section .note.ABI-tag we will modify it.

- sh_type = SHT_PROGBITS Because our code is a program data
- sh_flags = SHF_EXECINSTR | SHF_ALLOC Because our code must be run
- sh_addr = BaseAddress

- `sh_offset = ElfSize` Where the code has been injected
- `sh_size = BinaryInjectedSize`
- `sh_addralign = 16`

and we can just write back using the same `fwrite` and `fseek` method as the `codeInjection`. We will write back the section depending on the index of `.note.ABI-tag` inside the section header.

6 section headers calibration

6.1 Reorder sections

We will now reorder the sections header by section address for that we can create an array that will contains every section header during the loop in the previous task then we will pass it through the function `reorder` and it will check wheter the new address of `.note.ABI-tag` have to be swapped left or right. with a simple loop over the left or right section of `.note.ABI-tag` and a simple comparison between their address. To swap the section header inside the array we will just create a temporary memory that will get a section header then to a simple swap.

When everything is done we will loop over the array and write back each section header to the elf binary.

6.2 Set section names

Now we have to set the name that was passed through the command line argument

We can just get the `.shstrtab` which is the string table in ELF. The string table is used for all other references so each symbol from an ELF object has a member called `st_name` which is an index into this string table. then the field `sh_name` is an index inside the `.shstrtab` so we can use it to create an offset.

$$offsetShstrtabAbiTag = shstrtab -> shOffset + shdrAbitag -> shName \quad (3)$$

When we have the offset of the section `abiTag` in our string table we can just overwrite it with a simple `fwrite`.

Be aware that before changing the name we have to check that the new name length is lower than `.note.ABI-tag` else we will overwrite another section name.

When everything is done we can check the result with the command line `redelf -S ./date`.

1							
2	[Nr]	Name	Type	Address		Offset	
3		Size	EntSize	Flags	Link	Info	Align
4	[0]		NULL	0000000000000000			00000000
5		0000000000000000	0000000000000000			0	0
6	[1]	.interp	PROGBITS	0000000000400238			00000238

```

7      0000000000000001c 0000000000000000 A      0      0      1
8      [ 2] .note.gnu.build-id NOTE      0000000000400274 00000274
9      0000000000000024 0000000000000000 A      0      0      4
10     ...
11     ...
12     [17] .eh_frame      PROGBITS      000000000040e860 0000e860
13     000000000000143c 0000000000000000 A      0      0      8
14     [18] injected      PROGBITS      0000000000602b70 00010b70
15     0000000000000041 0000000000000000 AX      0      0      16
16     [19] .init_array    INIT_ARRAY    000000000060fe10 0000fe10
17     0000000000000008 0000000000000000 WA      0      0      8
18     [20] .fini_array    FINI_ARRAY    000000000060fe18 0000fe18
19     0000000000000008 0000000000000000 WA      0      0      8
20     ...
21

```

Our injected has been change with the right informations and placed at the good place depending of its address.

7 Overwriting the PT_NOTE program header

Before trying to execute our injected code, we now have to modify the program header PT_NOTE index

- `phdr[index_PT_NOTE].p_type = PT_LOAD` Allow the segment to be loadable
- `phdr[index_PT_NOTE].p_offset = elfSize`
- `phdr[index_PT_NOTE].p_vaddr = baseAddress`
- `phdr[index_PT_NOTE].p_paddr = baseAddress`
- `phdr[index_PT_NOTE].p_filesz = injectedSize`
- `phdr[index_PT_NOTE].p_memsz = injectedSize`
- `phdr[index_PT_NOTE].p_flags = PF_X | PF_R` for readable and executable
- `phdr[index_PT_NOTE].p_align = 4096`

And as simple as it is we just have to write it back to the elf binary.

8 Execute the injected code

Now that we have done all this. We can now try to execute our injected code to print a string on the output. 2 methods.

8.1 Entry Point Modification

If we pass the option `-e` so we will modify the entry point for that we just modify the executable header `e_entry` by changing it to our `baseAddress`. then we also have to modify the assembly code to jump back to the old entry point. to jump back we can just do that

```
1
2      mov rax, 0x4022e0      ;0x4022e0 is the old entry point
3      jmp rax
4
```

at the end of our assembly code that means we will jump at the address content of the `rax` register.

8.2 Hijacking GOT Entries

But if we don't have the `-e` option we have to overwrite the got entry.

PLT is the Procedure Linkage Table which is used to call external functions whose address isn't known when linking and is done at run time by the dynamic linker.
GOT is the Global Offsets Table that will be used to resolve addresses.

Now we have to find a function that will be used to hijack the program. For that the command `ltrace`.

`Ltrace` intercepts and records the dynamic library calls which are called by the executed process and the signals which are received by that process. It can also intercept and print the system calls executed by the program. cf `ltrace` manual page

so we can `ltrace` our `date` program and see a function `getenv` that will be used to overwrite for that we can just `objdump` the `.plt` section and search the `getenv` address we check the jump address and we have using `objdump -d -j .plt ./date`

```
1
2      0000000000401700 <getenv@plt>:
3      401700: ff 25 22 e9 20 00      jmpq    *0x20e922(%rip)      # 610028 <
      gmon_start@plt+0x20e4c8>
4      401706: 68 02 00 00 00      pushq   $0x2
5      40170b: e9 c0 ff ff ff      jmpq    4016d0 <__ctype_toupper_loc@plt-0x10>
6
```

We can see in the `.plt` section the entry for `getenv` call at the address `0x610028` and we can see the relation between the `.got.plt` which one start his address at `0x600000` using `objdump -d -j .got.plt ./date` We can now modify the address of our program in the `got.plt` section by modifying the 29th bytes of this section by our `injectedCode` address so that when the program run `getenv`

function the jump performed point now to our program instead of the libc function `getenv`.

So we just have to search over the section to get the `.got.plt` section then we overwrite it at offset `+ 0x28` and run the program then we will have our string with the date print on the output.

We have been able to hijack the GOT entries successfully.

9 Complementary Information and Complexity

To run this project I choose to use 2assembly code depending on whether we want to modify the entry point or to the GOT entries so the program launch test is different.

Moreover because the binary is always modified when we run each program we are forced to keep a backup of our date binary.

To launch the program please refer to the Makefile on the git associated to this report.

Take care, even if ELF file are numerous there isn't a good documentation about it.

In conclusion we have learned a lot of things on the elf binary and have been able to hijack it to inject our code and execute it. We learned the way to use the C language to manipulate every elf content. But there's still more to learn with dynamic binary !