

Algorithms and their applications

Assignment #1

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1. Definition of theta is described below.

Def of Θ

$f(n)$ is $\Theta(g(n))$ if

$$C_2 g(n) \leq f(n) \leq C_1 g(n) \text{ for every } n \geq n_0$$

Ans. (a)

(a) $(n^3+1)^6$ is $\Theta(n^6)$

(proof)

$$1 \cdot (n^3)^6 \leq (n^3+1)^6 \leq 2 \cdot (n^3)^6$$

for $n \geq n_0$

$\therefore (n^3+1)^6$ is $\Theta(n^6)$.

Ans. (b)

(b) $\sqrt{10n^2 + 7n + 3}$ is $\Theta(n)$

(proof)

$$\sqrt{10} n \leq \sqrt{10n^2 + 7n + 3} \leq \sqrt{11} n$$

for $n \geq n_0$

$\therefore \sqrt{10n^2 + 7n + 3}$ is $\Theta(n)$

Ans.(c)

$$(c) \ 2n \log(n+2)^2 + (n+2)^2 \log n \text{ is } \Theta(n^2 \log n)$$

$$\begin{aligned} \text{(proof)} \quad n^2 \log n &\leq \underbrace{2n \log(n+2)^2 + (n+2)^2 \log n}_{4n \log(n+2) + (n+2)^2 \log n} \leq 2n^2 \log n \end{aligned}$$

$$\therefore 2n \log(n+2)^2 + (n+2)^2 \log n \text{ is } \Theta(n^2 \log n)$$

Ans.(d)

$$(d) \ 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} \text{ is } \Theta(\ln(n))$$

$$\begin{aligned} \text{(proof)} \quad \ln(n) &= \int_1^n \frac{dx}{x} \leq 1 + \frac{1}{2} + \dots + \frac{1}{n} = \sum_{k=1}^n \frac{1}{k} \leq 1 + \int_2^n \frac{dx}{x} \\ &\Rightarrow 1 + \ln(n) \end{aligned}$$

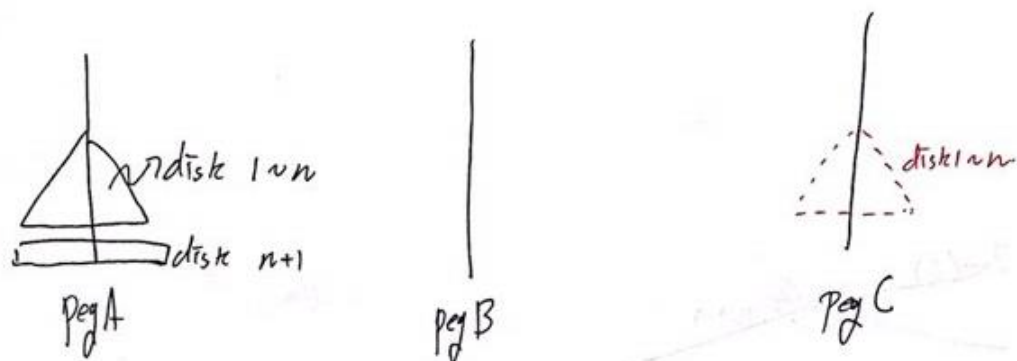
$$\text{Since } \Theta(1) < \Theta(\ln(n))$$

$$1 + \frac{1}{2} + \dots + \frac{1}{n} \text{ is } \Theta(\ln(n))$$

2. Restricted tower of Hanoi

Ans. (a)

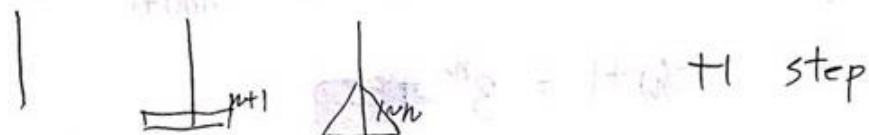
2 - (a)



Let $H(n)$ the total number of moving $\underbrace{\text{from A to C via B}}$.
And suppose there are $n+1$ disks to move.

1. It will take $H(n)$ to move n disk from A to C.
+ $H(n)$ step

2. then move disk_{n+1} from A to B



3. move backward $\text{disk}_{1 \text{ to } n}$ from C to A
↳ move disk_{n+1} to C



4. then move $\text{disk}_{1 \text{ to } n}$ from A to C again
+ $H(n)$ step

$$\text{Sum (1 to 4)} : H(n+1) = 3H(n) + 2$$

Ans.(b)

```
# 2-(b). Hanoi problem.
```

```
# Calculate summation number of moving disks
```

```
# n is natural number here.
```

```
def Hanoi_sum(n):
```

```
    if n == 1:    # Return end value for the recursive func which condition is that the disk is only one.
        return 2
```

```
    else:
        return 3*Hanoi_sum(n-1)+2    # Return recursive function that justified in 2-(a)
```

```
In [12]: Hanoi_sum(5)
```

(example output) **Out[12]: 242**

Ans.(c)

$$2-(c) \quad H_{(n+1)} = 3H_{(n)} + 2$$

$$H_{(n+1)} + 1 = 3(H_{(n)} + 1) \Rightarrow \frac{H_{(n+1)} + 1}{H_{(n)} + 1} = 3$$

$$H_{(n)} + 1 = 3^n$$

$$H_{(n)} = 3^n - 1 \quad (H_{(1)} = 2)$$

3. Selection Algorithm

```
# 3. Selection Algorithm
```

```
# Calculate entire number of case - choosing 'k' in 'n' candidates.
```

```
# n is natural number // k is integer which is >= 0
```

```
def Combination(n, k):
```

```
    if k == 0 or n == k:    #Return end(k=0) / exceptional(n=k) value for the recursive function.
        return 1
```

```
    else:
        return Combination(n-1, k-1) + Combination(n-1, k)
```

```
    """
```

```
# Description.
```

```
    Pulling k out of n equals the sum of the cases in which, after selecting particular one,
    the number of cases in which k is pulled out of it (n-1), and the number in which k-1 is selected including
    it.
```

```
    """
```

```
)
```

```
In [13]: Combination(5,2)
```

(example output) **Out[13]: 10**

4. Newton's method

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```
# Input ans as your starting rough guessing number
# tol is abbreviation of tolerance, the parameter that you think that the gap is good enough
def MySqrt(num, ans, tol):
    if abs(ans*ans - num) <= tol:
        return ans # if the gap btw guessing num and real one is close enough, than return the newly guessing number.
    else:
        return MySqrt(num, (ans*ans + num)/(2*ans), tol) # Newton method that defined as average btw (ans & num/ans)
```

In [16]: MySqrt(2, 1.5, 0.0001)

(example output) Out[16]: 1.4142156862745099

5. Infix to prefix using manual process

5- (a)

<infix>

$A * B + C * D$

→ $(A * B) + (C * D)$

<prefix>

$AB * CD * +$

5- (b)

<infix>

$(A - 2 * (B + C) - D * E) * F$

→ $((((A - 2 * (B + C)) - (D * E)) * F))$

<prefix>

$A B C + 2 * - D E * - F *$

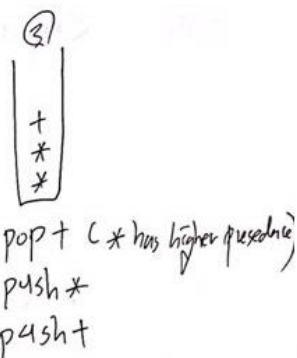
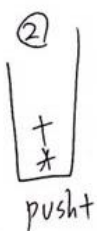
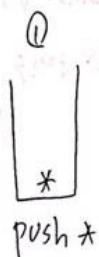
6. Prefix to Infix using *manual process*

6-(a) $\langle \text{postfix} \rangle$ $AB* C - D +$ $\langle \text{infix} \rangle$ $A*B - C + D$
 $\rightarrow (A*B) - C + D$

6-(b) $\langle \text{postfix} \rangle$ $AB C + * D -$ $\langle \text{infix} \rangle$ $A * (B + C) - D$
 $\rightarrow (A * (B + C) - D)$

7. Infix to postfix using *stack*

7-(a) $A * B + C * D$
 sol) $A^{\textcircled{1}} * B^{\textcircled{2}} + C^{\textcircled{3}} * D^{\textcircled{4}}$
 $\xrightarrow{\text{Scan operation}}$



\therefore postfix: $AB*CD*+$

7-(b)

① ② ③ ④ ⑤ ⑥

$(A - 2 * (B + C) - D * E) * F$

scan operator

①

-

push -

②

*

push *

③

+

*

push +

④

-

+

*

push -

⑤

-

*

+

*

pop -

push *

⑥

*

+

+

*

push *

postfix: A B C + 2 * - D E * - F *