

Simplification

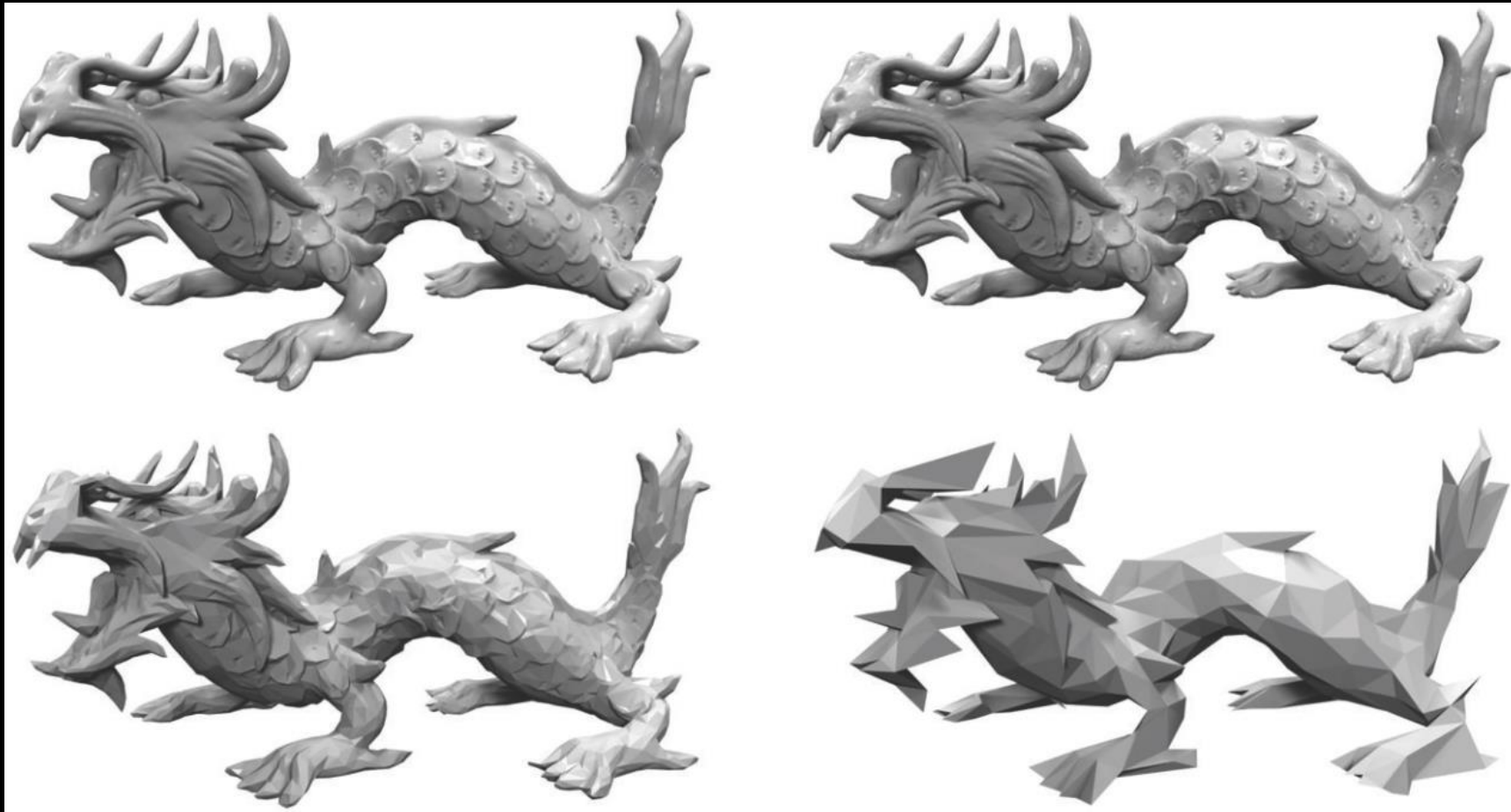
Xiao-Ming Fu

Outlines

- Definition
- Local operations
- Quadric error metric
- Variational shape approximation

Simplification and approximation

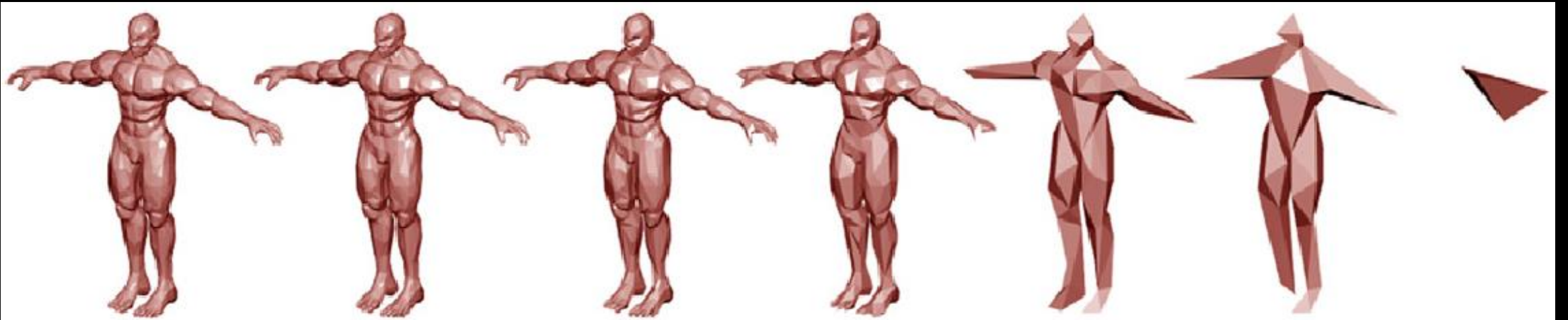
- Transform a given polygonal mesh into another mesh with **fewer** faces, edges, and vertices.



Simplification and approximation

- Transform a given polygonal mesh into another mesh with **fewer** faces, edges, and vertices.
- The simplification or approximation procedure is usually controlled by user-defined **quality criteria**.

Curvature-preserved criteria



2053

1500

1000

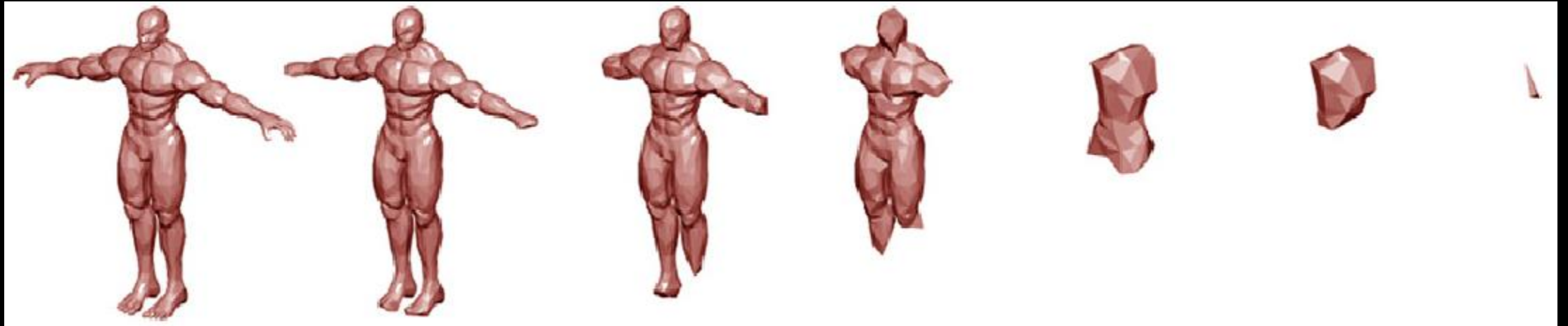
500

100

50

4

Curvature-removed criteria



2053

1500

1000

500

100

50

4

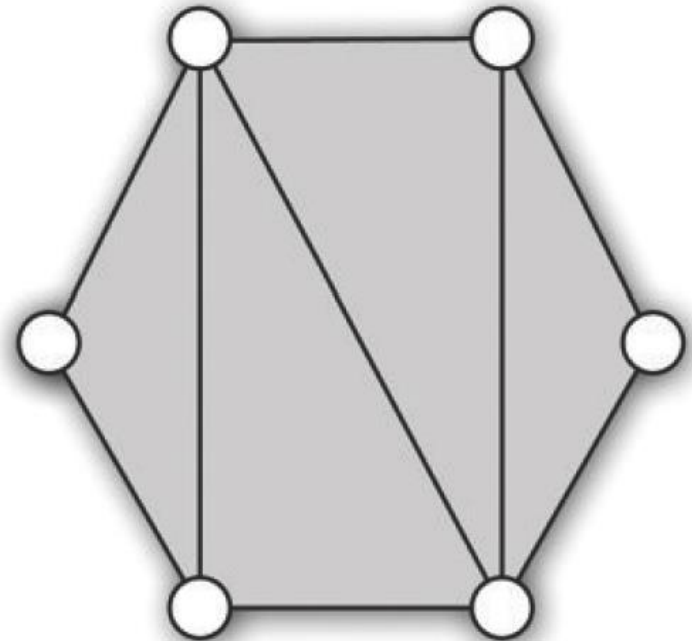
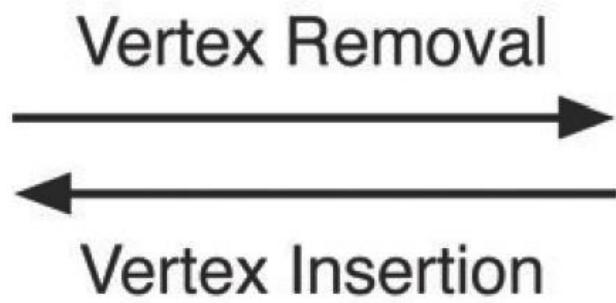
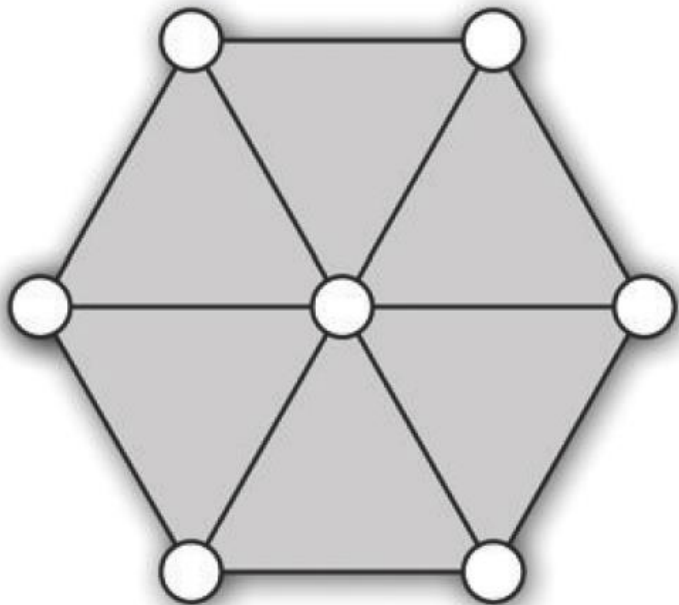
Simplification applications

- Adjust the complexity of a geometric data set
- Since many decimation schemes work iteratively, i.e., they decimate a mesh by removing one vertex at a time, they usually can be **inverted**.
 - Hierarchical method

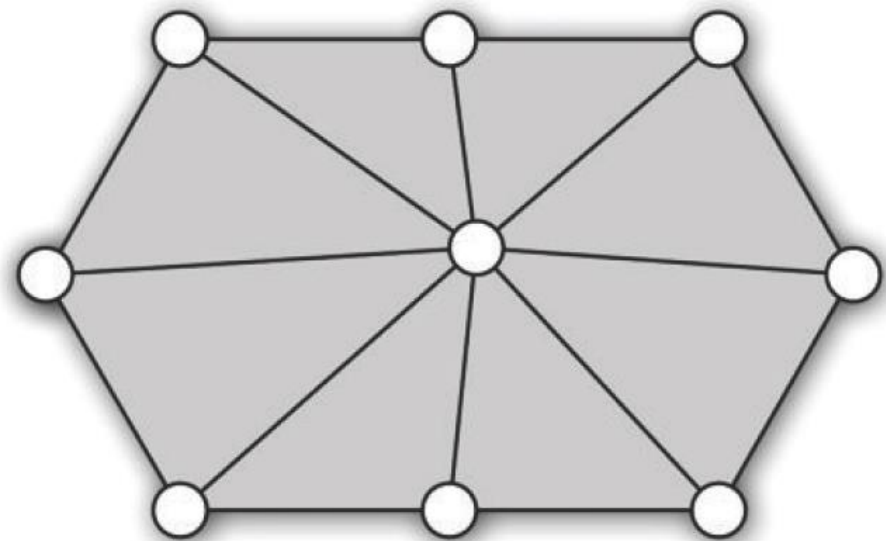
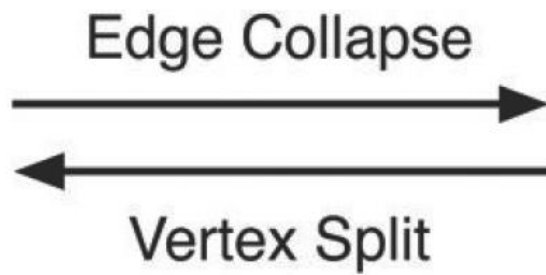
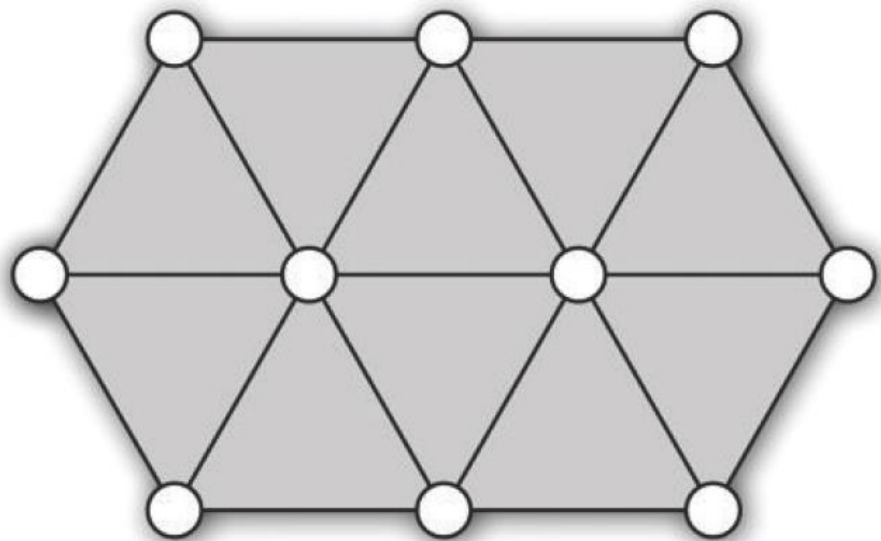
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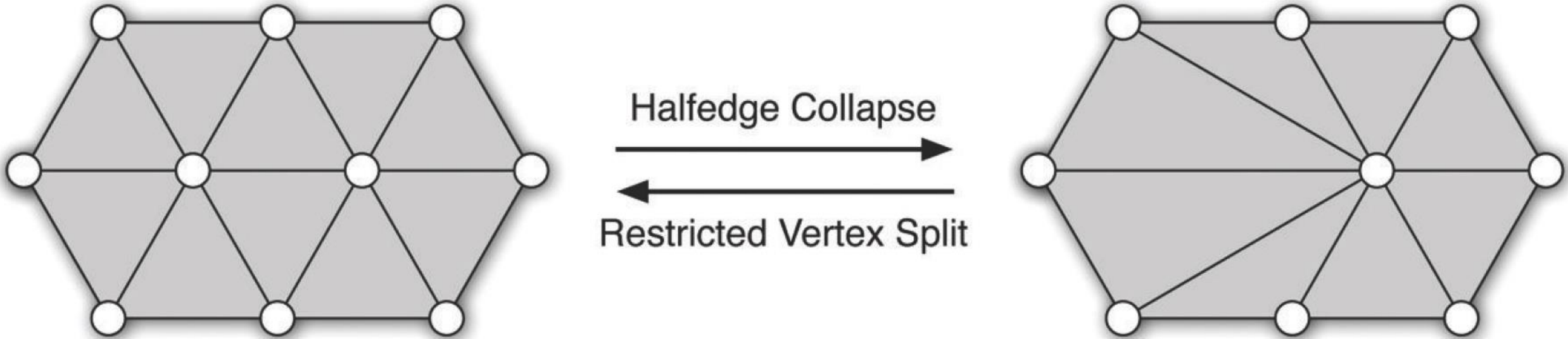
Vertex removal



Edge collapse



Half-edge collapse



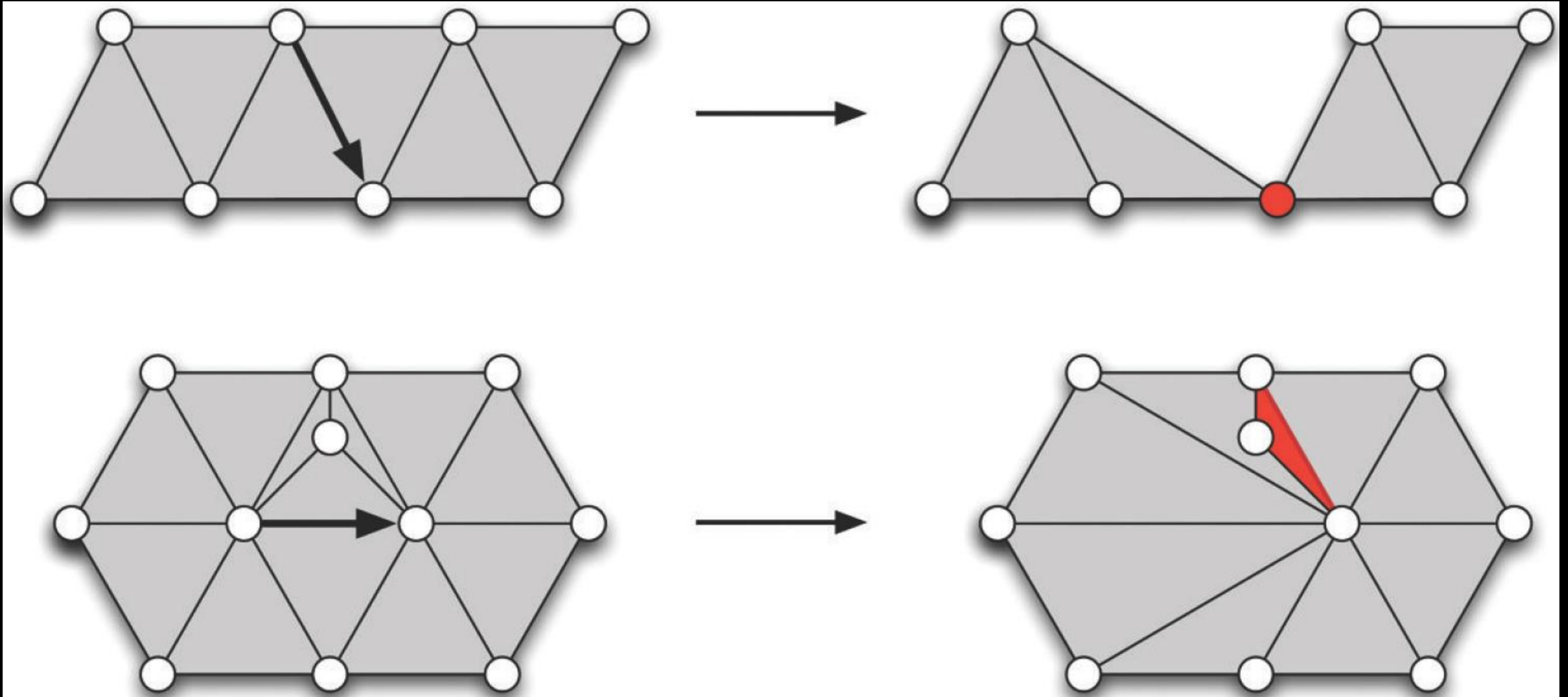
After collapse: $n(E) - 3$, $n(V) - 1$, $n(F) - 2$.

According to Euler's formula: $2 - 2m = n(V) + n(F) - n(E)$.

Half-edge collapsing would not change the genus of a mesh.

OpenMesh: `collapse()`, `is_collapse_ok()`.

Topologically illegal (half-)edge collapses



Outlines

- Definition
- Local operations
- **Quadratic error metric**
- Variational shape approximation

Incremental algorithms

- Removing one vertex at a time
- The iterative decimation procedure can take **arbitrary user-defined criteria** into account, according to **which the next removal operation is chosen**.

Quadric error metric (QEM)

- The squared distance of a point x from the plane P_i :

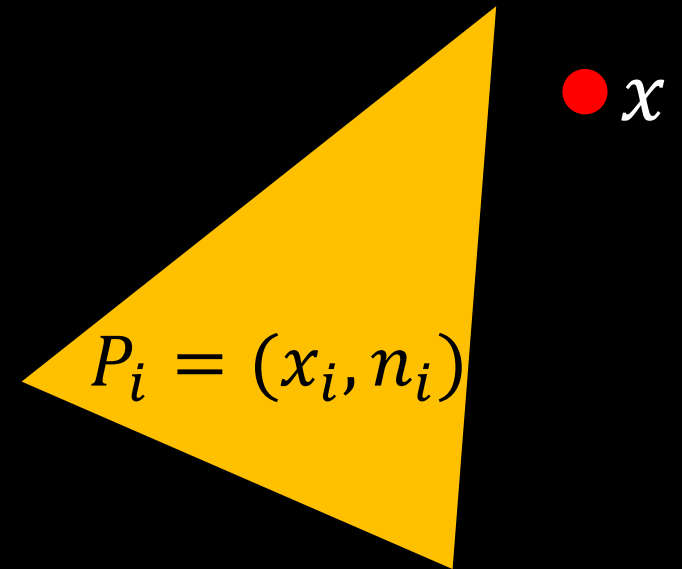
$$d(x, P_i) = (n_i^T x - d_i)^2$$
$$d_i = n_i^T x_i$$

Denote $\bar{x} = (x, 1)$ and $\bar{n}_i = (n_i, -d_i)$.

Then:

$$d(x, P_i) = (\bar{n}_i^T \bar{x})^2 = \bar{x}^T \bar{n}_i \bar{n}_i^T \bar{x} =: \bar{x}^T Q_i \bar{x}$$

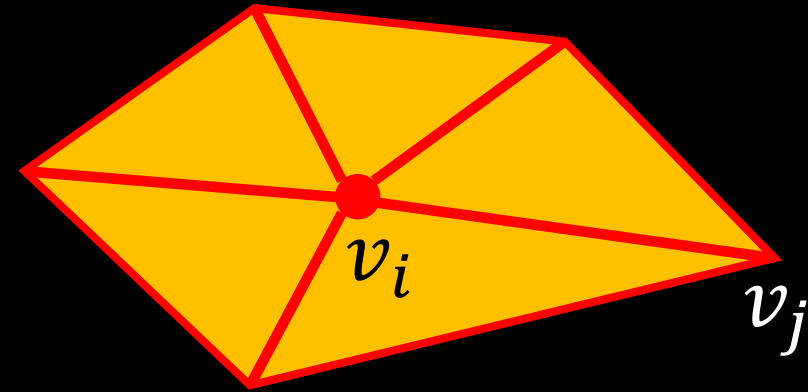
Quadratic error Matrix



Quadratic error Matrix Q

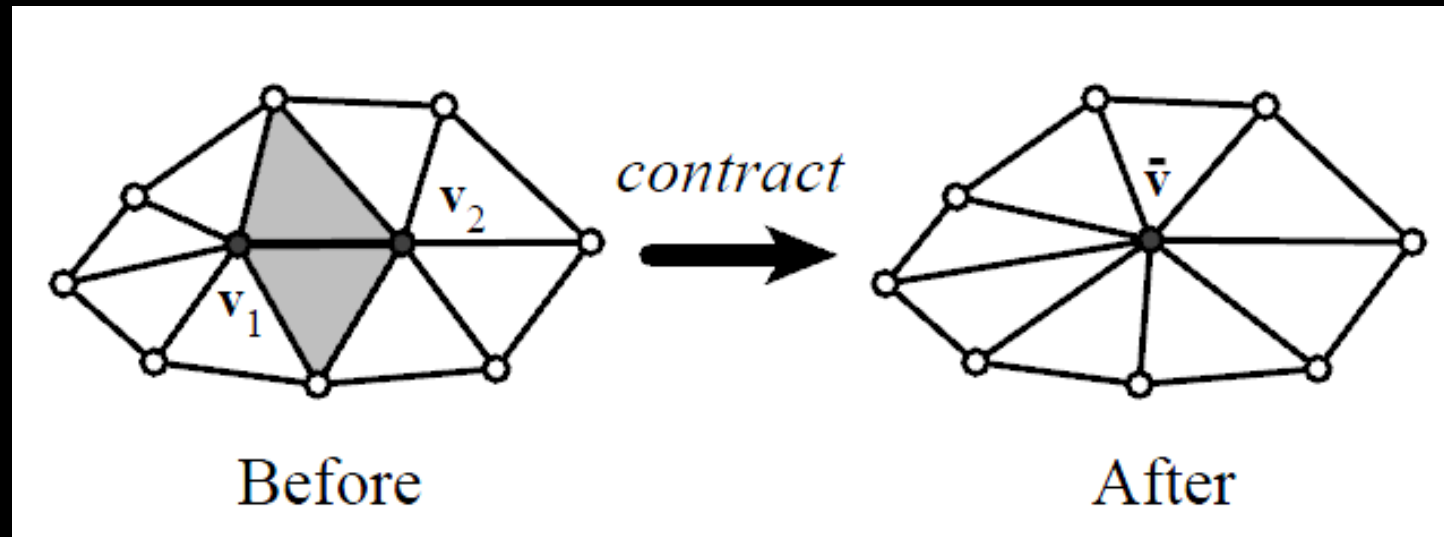
- On vertices

$$Q_i^v = \sum_{j \in \Omega(i)} Q_j$$



- On edge

$$Q^e = Q_1^v + Q_2^v$$



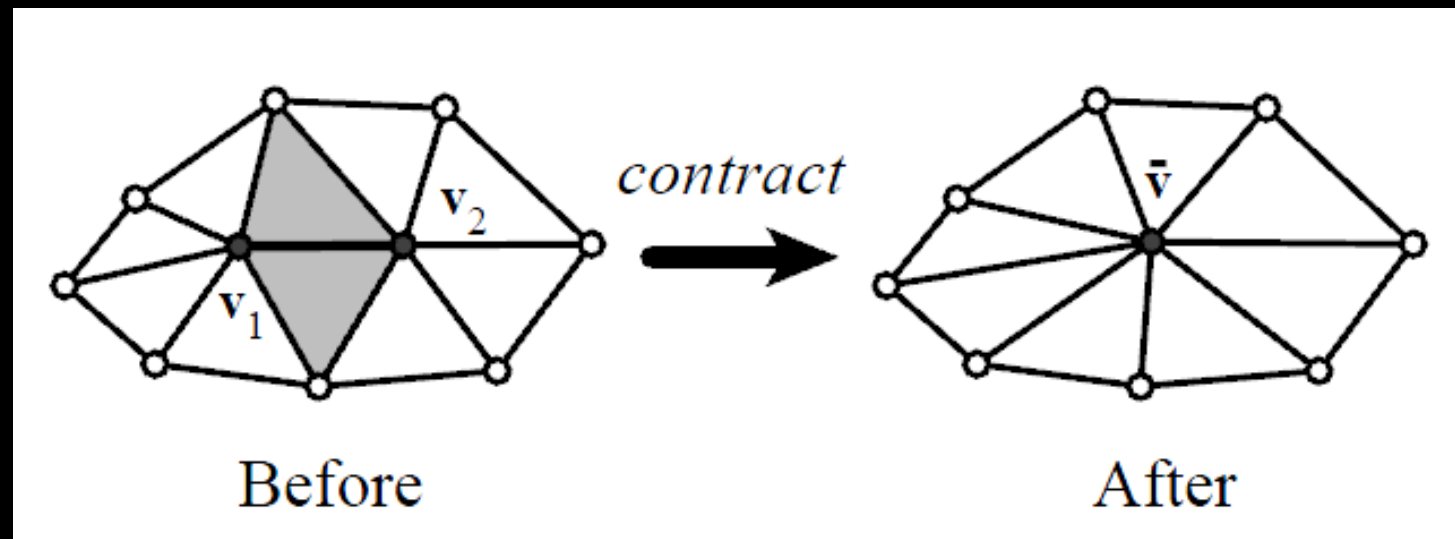
QEM error

- QEM error On edge:

$$\bar{v} = \arg \min_v v^T Q^e v$$

Note: Q^e may not a full rank matrix

- Q on \bar{v} is just Q^e .



QEM Algorithm

Input: a mesh

Output: a simplified mesh

Initialization:

Compute the Q^e matrices for all the edges.

Compute the optimal contraction target \bar{v} for each edge.

While $N_v > n$ and $Cost_{min} < t$

 The error $\bar{v}^T Q^e \bar{v}$ becomes the **cost** of the edge.

 Place all the edges in a priority queue keyed on cost with minimum **cost** edge at the top.

 Remove the edge of the least cost from the heap , collapse this edge, and update the costs of all edges involving.

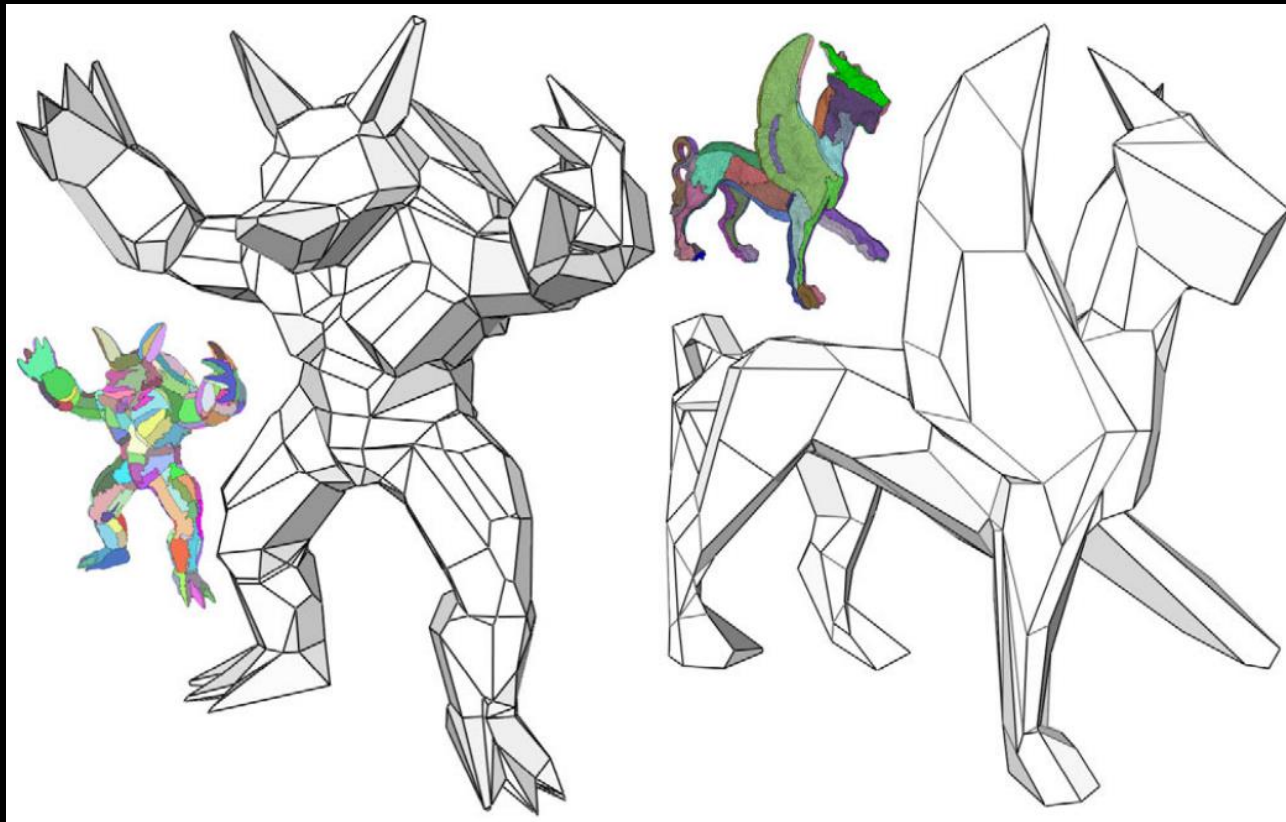
End;

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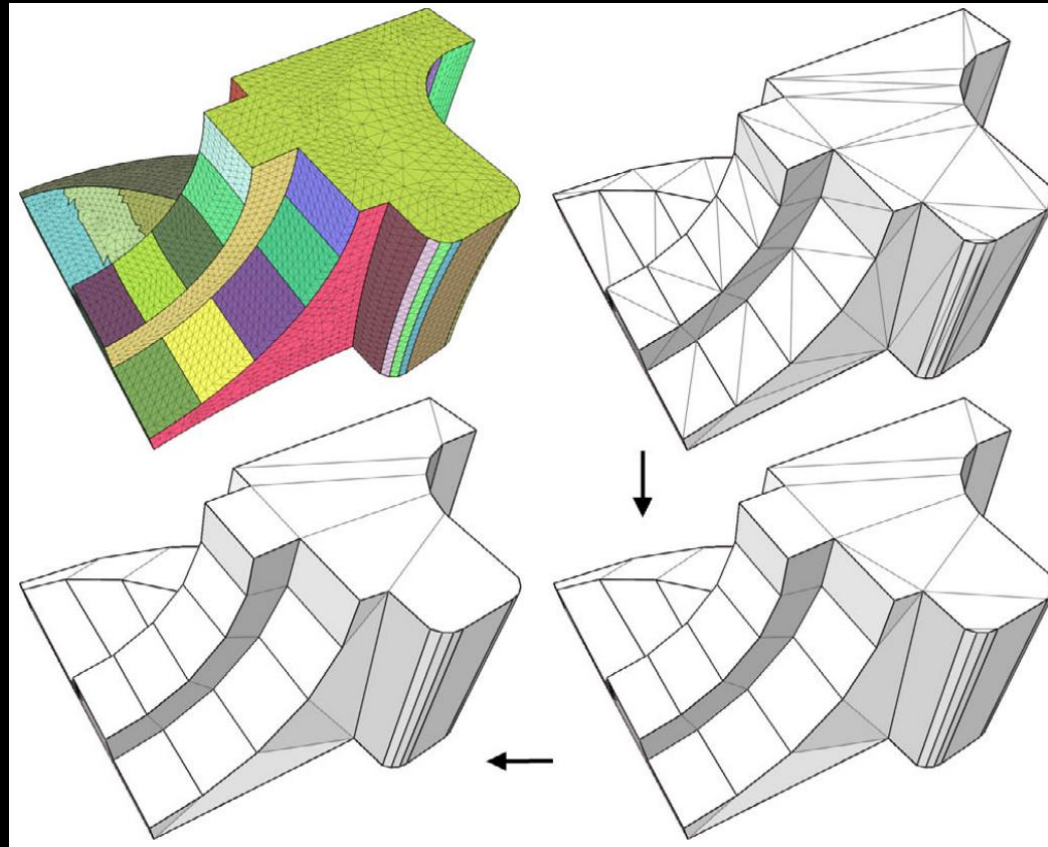
Variational shape approximation (VSA)

- VSA is highly sensitive to features and symmetries and produces anisotropic meshes of high approximation quality.



Variational shape approximation (VSA)

- The input shape is approximated by a set of proxies.
 - A plane in space through the point x_i with normal direction n_i .



Region representation

$$R_1 \cup \dots \cup R_k = M$$

M : a triangle mesh

$R = \{R_1, \dots, R_k\}$: a partition of M into k regions.

Proxies: $P = \{P_1, \dots, P_k\}$, $P_i = (x_i, n_i)$

Distance metrics between R_i and P_i

- The squared orthogonal distance of x from the plane P_i .

$$L^2(R_i, P_i) = \int_{x \in R_i} (n_i^T x - n_i x_i)^2 dA$$

- A measure of the normal field:

$$L^{2,1}(R_i, P_i) = \int_{x \in R_i} \|n(x) - n_i\|^2 dA$$

Goal of VSA

- Given a number k and an error metric E (L^2 or $L^{2,1}$), find a set $R = \{R_1, \dots, R_k\}$ of regions and a set $P = \{P_1, \dots, P_k\}$ of proxies such that the global distortion

$$E(R, P) = \sum_{i=1}^k E(R_i, P_i)$$

is minimized.

Lloyd's clustering algorithm

- The algorithm iteratively alternates between a **geometry partitioning** phase and a **proxy fitting** phase.
- Geometry partitioning phase
 - a set of regions that best fit a **given** set of proxies
- Proxy fitting phase
 - the partitioning is kept **fixed** and the proxies are adjusted
- Initialization
 - randomly picking k triangles as R
 - The planes of k triangles are used to initialize P

Geometry partitioning phase

- **Modifies** the set R of regions to achieve a **lower approximation error** while keeping the proxies P **fixed**.

```
partition( $\mathcal{R} = \{\mathcal{R}_1, \dots, \mathcal{R}_k\}$ ,  $\mathcal{P} = \{P_1, \dots, P_k\}$ )
```

```
// find the seed triangles and initialize the priority queue
```

```
queue =  $\emptyset$ 
```

```
for  $i = 1$  to  $k$  do
```

```
    select the triangle  $t \in \mathcal{R}_i$  that minimizes  $E(t, P_i)$ 
```

```
     $\mathcal{R}_i = \{t\}$ 
```

```
    set  $t$  to conquered
```

```
    for all neighbors  $r$  of  $t$  do
```

```
        insert  $(r, P_i)$  into queue
```

```
// grow the regions
```

```
while the queue is not empty do
```

```
    get  $(t, P_i)$  from the queue that minimizes  $E(t, P_i)$ 
```

```
    if  $t$  is not conquered then
```

```
        set  $t$  to conquered
```

```
         $\mathcal{R}_i = \mathcal{R}_i \cup \{t\}$ 
```

```
        for all neighbors  $r$  of  $t$  do
```

```
            if  $r$  is not conquered then
```

```
                insert  $(r, P_i)$  into queue
```

Proxy fitting phase

- The partition R is kept fixed, the proxies P_i are adjusted in order to minimize approximation error.
- L^2 metric
 - The best proxy is the least-squares fitting plane.
- $L^{2,1}$ metric
 - The proxy normal n_i is just the area-weighted average of the triangle normals.

