

有 array 和 linked list big O 時間表

Data Structures

LINKED LISTS (CHAPTER 4)

Todo List

Active tasks (13)Design drafts in 3 different styles ASAP ⌚7

🚀 In Progress

Design

17 Apr 2024

Michael Martinez

Develop a wireframe ⌚3 ⌚1

🚀 In Progress

Operational

⌚

Sofia Brown

Write website copy ✓/1/3

🚀 In Progress

Operational

19 Apr 2024

Anastasia Novak

Write meta title & meta description for each page

📅 Scheduled

Operational

⌚

Anastasia Novak

Develop the website using the chosen CMS platform blocked ✓/0/4

📅 Scheduled

Operational

19 Apr 2024

David Thomas

Design the entire website in a chosen style

📅 Scheduled

Design

24 Apr 2024

Michael Martinez

Deploy the website to the development hosting server

✉️ New task

Operational

⌚

David Thomas

Prepare design files for web developer ✓/0/2

✉️ New task

Design

⌚

Michael Martinez

Implement responsive design

📅 Scheduled

Operational

⌚

David Thomas

Send new website link to the team

✉️ New task

Operational

⌚

David Thomas

Fix all the bugs reported by the team

✉️ New task

Operational

⌚

David Thomas

Review and comment on website design Feedback

✉️ New task

Important

⌚

Marry Williams

Deploy the website to the production environment

✉️ New task

Operational

⌚

David Thomas

+ Add task

Image credit: <https://bordio.com/to-do-list/>Completed tasks (3)

Research potential CMS platforms for website development

✅ Completed

Operational

⌚

David Thomas

Develop a structure for a new website ✓/4/4 ⌚2

✅ Completed

Operational

⌚

Sofia Brown

List

A number of connected items or names written or printed **consecutively**, typically one below the other.

For example:

- Check list
- Mailing list
- Todo list
- Wish list
- Memory management (OS)
- Data structures (stacks, queues, sets, hash tables, graphs)
- ...

Efficiency Way for Insertion/Deletion

Efficiency way: $O(1)$, $O(n)$, $O(n \log n)$?

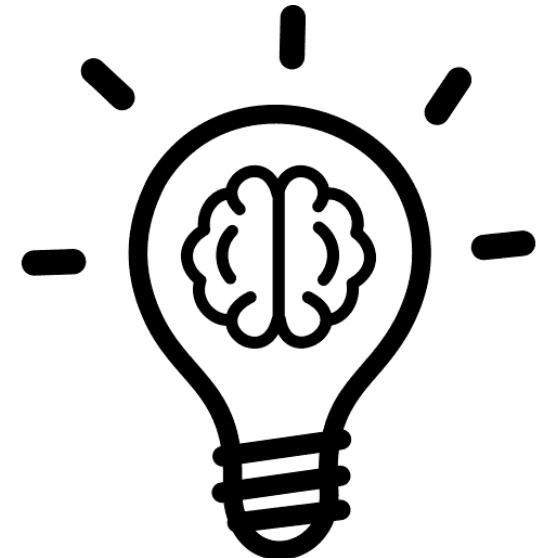
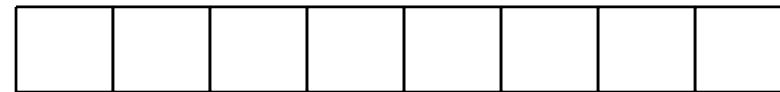


Image credit: <https://uxwing.com/idea-icon/>

Array

Array



Integer Array

64	34	25	12	22	11	90	8
----	----	----	----	----	----	----	---

Array

Operations

- Insert an element into an array
- Delete an element from an array
- Costly shifts ? improvement?

Efficiency Way for Insertion/Deletion

Traverse the entire data structures and insert/delete the target elements

Insertion

- Insert an elements before or after the target element

Deletion

- Delete the target element

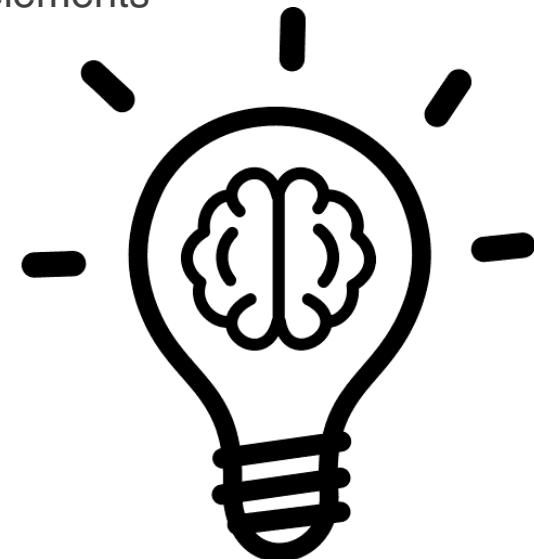


Image credit: <https://uxwing.com/idea-icon/>

What is This?



Image credit: <https://www.britannica.com/technology/railroad-coupling>

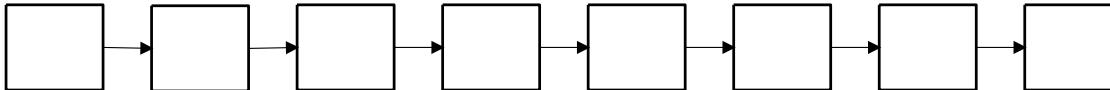
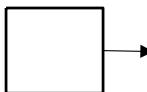
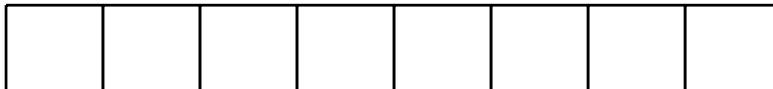
Linked List

Linked lists consist of a number of elements grouped, or linked, together in a specific order.

Array and Linked List

Something should be linked.

Add the linker will be easy to manipulate the order

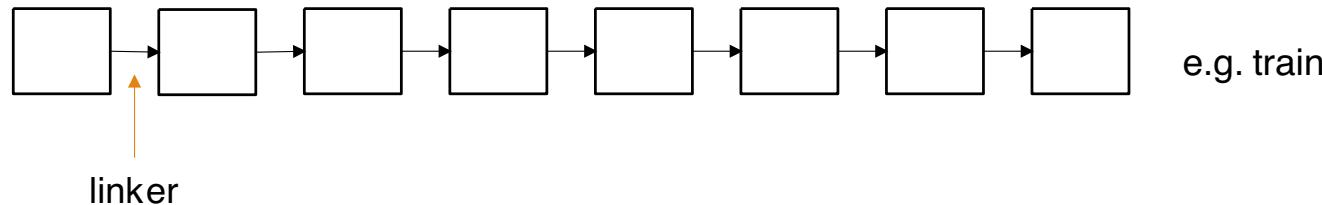


From Array to Linked List

1. Array



2. Linked list

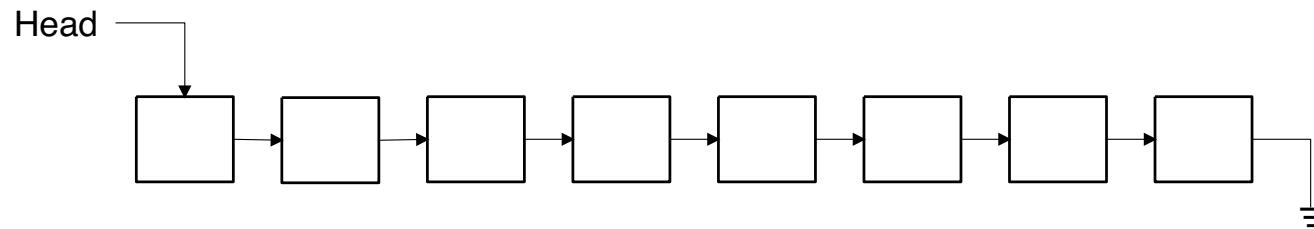


Something Missing?

1. Array



2. Linked list



Linked List

A linear data structure where elements (nodes) are stored in sequence, but not in contiguous memory locations.

Each node contains data and a reference (pointer) to the next node.

Unlike arrays, linked lists can grow or shrink during runtime, but accessing elements requires traversing from the head node.

ADT: Linked List

ADT LinkedList is

objects:

「鏈 序 列」 由 3 個

A finite sequence of nodes, each containing:

- an item (data value from set item)

- a reference (pointer) to the next node

指向下一節點

The sequence terminates with a null reference.

functions:

for all $L \in \text{LinkedList}$, $x \in \text{item}$, $p \in \text{position}$

$\text{LinkedList Create}()$::= return an empty linked list

$\text{Boolean IsEmpty}(L)$::= return true if L has no nodes
else return false

$\text{Integer Length}(L)$::= return the number of nodes in L

$\text{Item Retrieve}(L, p)$::= if position p is valid, return the item at position p
else return error

ADT: Linked List

LinkedList Insert(*L*, *p*, *x*) ::=
 

LinkedList Delete(*L*, *p*) ::=

LinkedList Update(*L*, *p*, *x*) ::=

Position Search(*L*, *x*) ::=

end *LinkedList*

if position *p* is valid, insert item *x* into list *L* at position *p* shift subsequent nodes if necessary
return updated list
else return error

if position *p* is valid, remove the node at position *p*
return updated list
else return error

if position *p* is valid, replace the item at position *p* with *x*
return updated list
else return error

return position of first occurrence of *x* in *L*
else return error

Linked List in C

```
struct list {  
    int integerValue;  
    struct list *nextPtr;  
};  
  
typedef struct list {  
    int integerValue;  
    struct list *nextPtr;  
} IntegerNode;
```

Linked List in C++

Node Class

```
class Node {  
public:  
    int data;  
    Node* next;  
  
    Node(int value) {  
        data = value;  
        next = nullptr;  
    }  
};
```

Linked List in C++

```
class LinkedList {  
private:  
    Node* head;  
  
public:  
    LinkedList() {  
        head = nullptr;  
    }
```

```
// Insert at end  
void insert(int value) {  
    Node* newNode = new Node(value);  
    if (head == nullptr) {  
        head = newNode;  
    } else {  
        Node* temp = head;  
        while (temp->next != nullptr) {  
            temp = temp->next;  
        }  
        temp->next = newNode;  
    }  
}
```

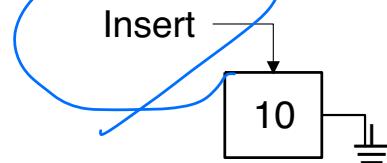
Linked List in C++

```
// Print the list
void print() {
    Node* temp = head;
    while (temp != nullptr) {
        std::cout << temp->data << " -> ";
        temp = temp->next;
    }
    std::cout << "NULL" << std::endl;
}
```

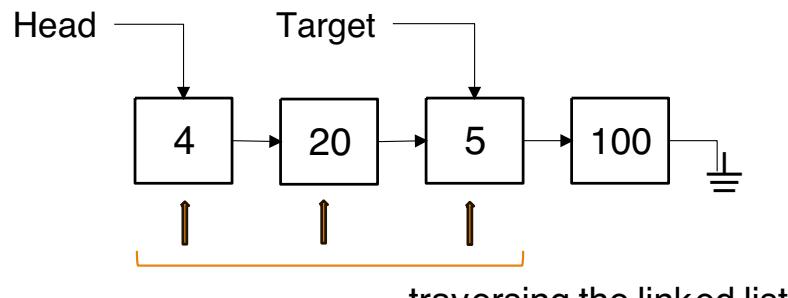
```
// Destructor to free memory
~LinkedList() {
    Node* current = head;
    while (current != nullptr) {
        Node* nextNode = current->next;
        delete current;
        current = nextNode;
    }
};
```

Operation: Insert Element (10) after the Target (5), integerLL

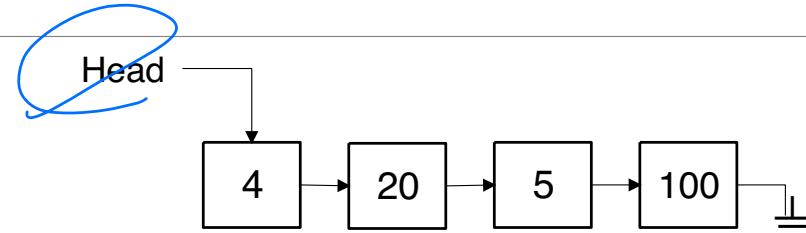
- 1 Initialize the node



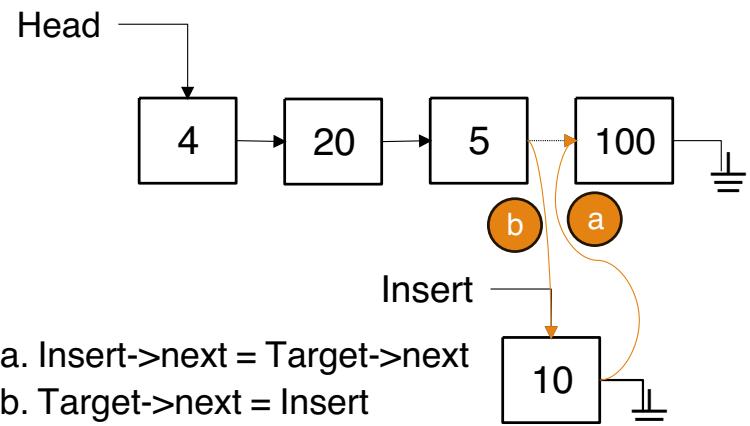
- 2 Traverse the linked list to find the target



Target = Head->next->next

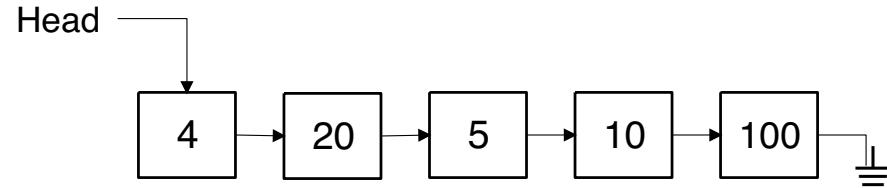


- 3 Insert the Element (10)



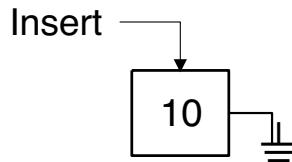
Operation: Insert Element (10) after the Target (5), integerLL

4 Final

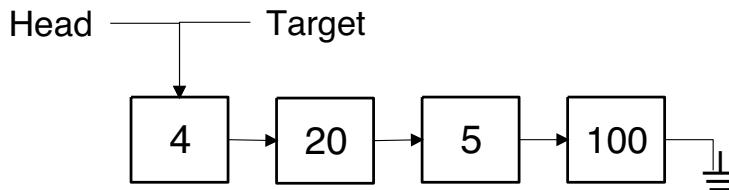


Operation: Insert Element (10) in the Beginning of the Linked List

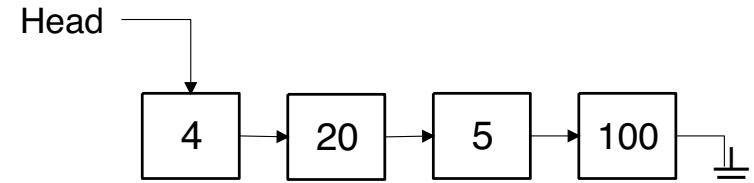
- 1 Initialize the node



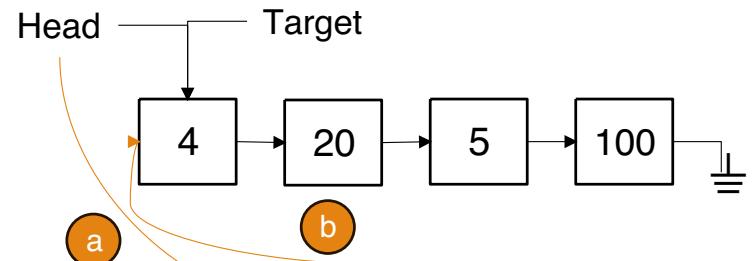
- 2 No need to traverse the linked list



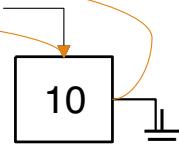
Target = Head



- 3 Insert the Element (10)

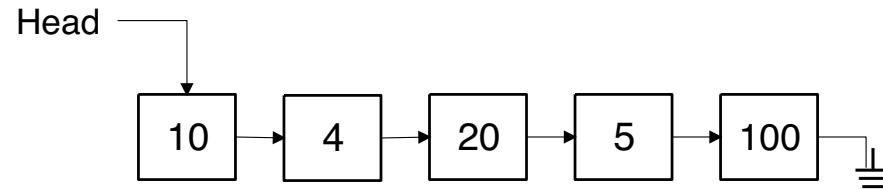


- a. Head= Insert
b. Insert->next = Target



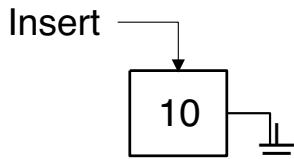
Operation: Insert Element (10) in the Beginning of the Linked List

4 Final

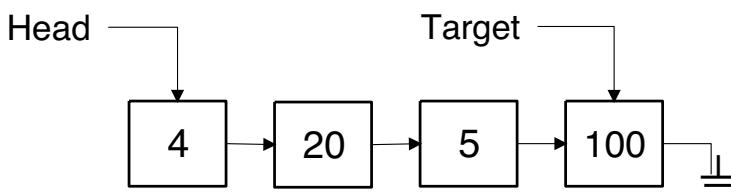


Operation: Insert Element (10) in the End of the Linked List

- 1 Initialize the node



- 2 Traverse the linked list to find the end

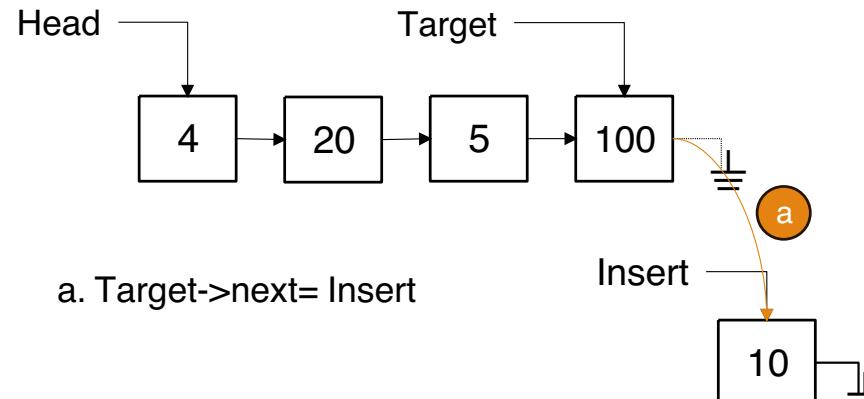


```
Target = Head  
While (Target->next != NULL){  
    Target = Target->next;  
}
```

Head

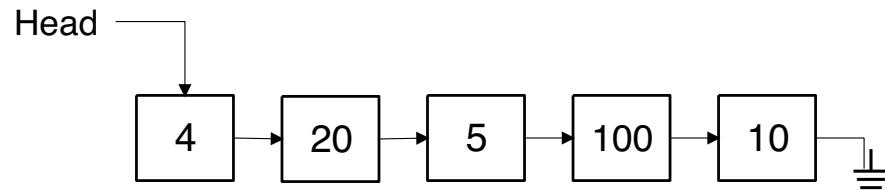


- 3 Insert the Element (10)



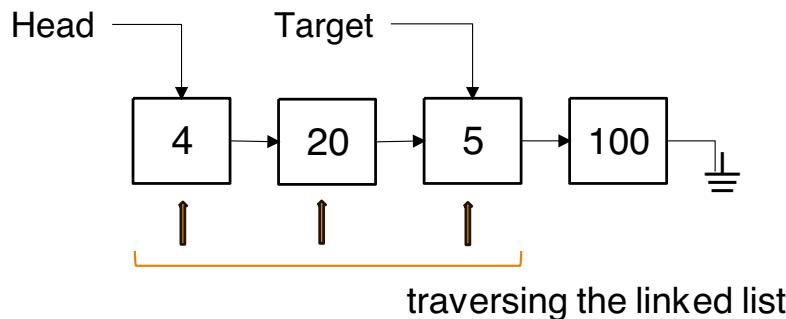
Operation: Insert Element (10) in the End of the Linked List

4 Final



Operation: Delete the Target (5)

- 1 Traverse the linked list to find the target

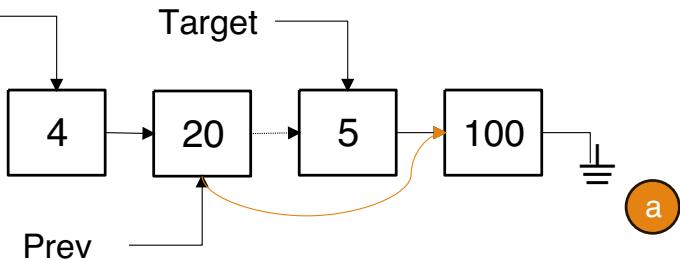
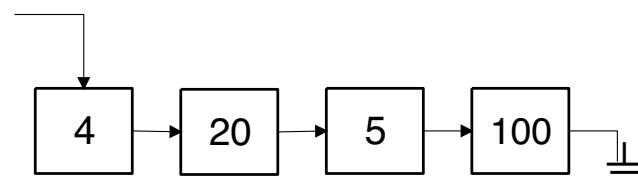


Target = Head

Prev = Head

```
While (Target->value !== Value){  
    Prev = Target  
    Target = Target->next  
}
```

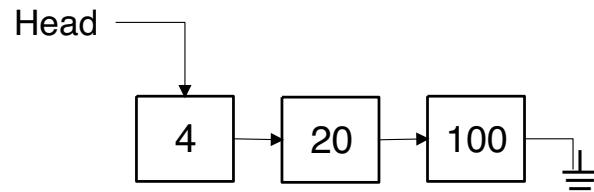
- 2 Delete the Element (5)



- a. Prev->next = Target->next
- b. free(Target)

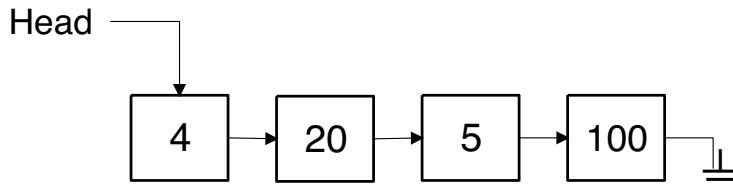
Operation: Delete the Target (5)

4 Final



Operation: Delete the Target (First)

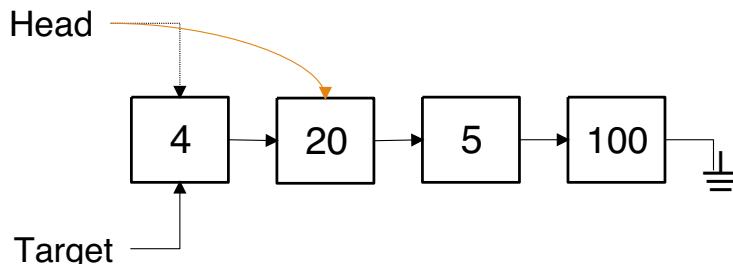
- 1 No need to traverse the linked list



Target = Head

Prev = Head

- 2 Delete the first element (4)

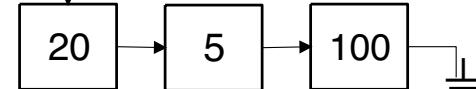


Head

4
20
5
100

3 Final

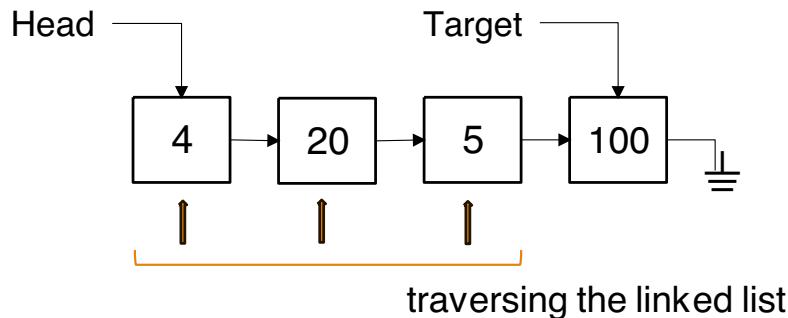
Head



Head = Target->next
free(target)

Operation: Delete the Target (Last)

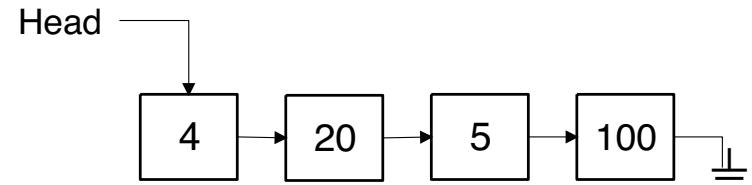
- 1 Traverse the linked list to fine the end



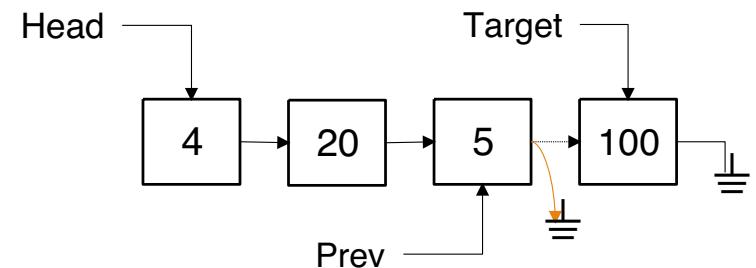
Target = Head

Prev = Head

```
While (Target->next != NULL){  
    Prev = Target  
    Target = Target->next  
}
```



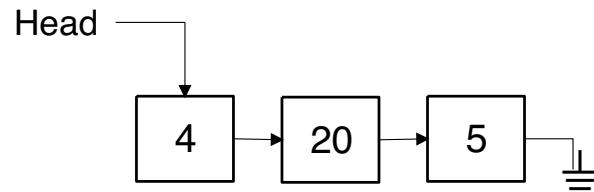
- 2 Delete the first element (4)



Prev->next = NULL
free(target)

Operation: Delete the Target (Last)

③ Final



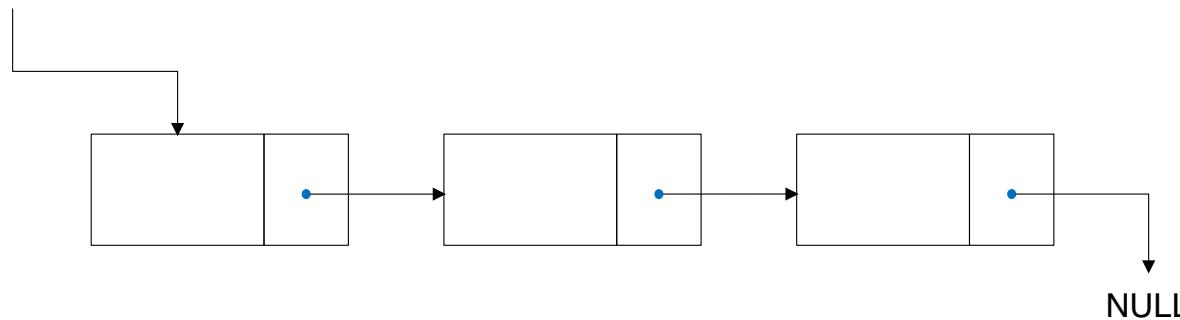
Type of Linked List (Variation)

1. Singly linked list (basic form)
2. Doubly linked list (two-way navigation)
3. Circular linked list (wrap-around structure)
 - Singly linked List
 - Doubly linked list

Singly Linked List



Head



Singly Linked List in C

```
struct SNode {  
    int data;  
    struct SNode *next;  
};
```

Singly Linked List in C++

```
class SNode {  
public:  
    int data;  
    SNode* next;  
    SNode(int val) : data(val), next(nullptr) {}  
};
```

Practice Problems: Traversing the Linked List and Observe the Memory Address

```
#include <stdio.h>
#include <stdlib.h>

// Define the node structure
typedef struct Node {
    int data;
    struct Node* next;
} Node;
```

Practice Problems: Traversing the Linked List and Observe the Memory Address

```
// Create a new node with given value
Node* createNode(int value) {
    Node* newNode = (Node*)malloc(sizeof(Node));

    if (newNode == NULL) {
        printf("Memory allocation failed.\n");
        exit(1);
    }
    newNode->data = value;
    newNode->next = NULL;
    return newNode;
}
```

Practice Problems: Traversing the Linked List and Observe the Memory Address

```
// Traverse and print memory information
void traverseList(Node* head) {
    Node* current = head;
    int index = 0;

    while (current != NULL) {
        printf("Node %d: Value = %d, Address = %p, Next = %p\n",
               index, current->data, (void*)current, (void*)current->next);
        current = current->next;
        index++;
    }
}
```

Practice Problems: Traversing the Linked List and Observe the Memory Address

```
int main() {  
    // Create linked list: 10 -> 20 -> 30 -> NULL  
    Node* head = createNode(10);  
    head->next = createNode(20);  
    head->next->next = createNode(30);  
  
    // Traverse and print  
    printf("Traversing the linked list:\n");  
    traverseList(head);  
}  
  
// Free memory  
Node* current = head;  
while (current != NULL) {  
    Node* temp = current;  
    current = current->next;  
    free(temp);  
}  
  
return 0;
```

Doubly Linked List

next(指後面)

prev(指前面)

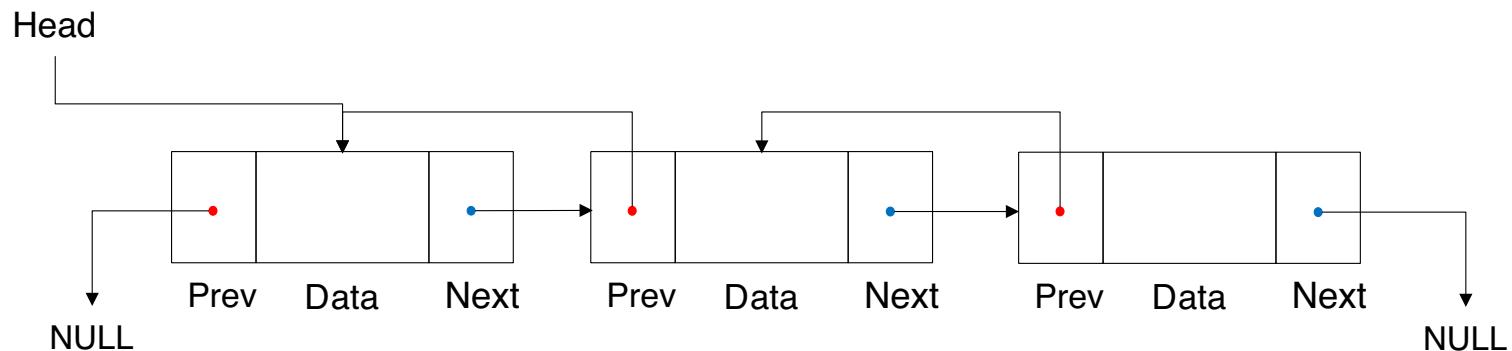
Doubly-linked lists, as their name implies, are composed of elements linked by two pointers.

Each element of a doubly-linked list consists of three parts: in addition to the data and the next pointer, each elements includes a pointer to the previous elements, called the prev pointer.

Improvement:

提供双向導航

Doubly Linked List



Doubly Linked List in C

```
struct DNode {  
    int data;  
    struct DNode *prev;  
    struct DNode *next;  
};
```

Doubly Linked List in C++

```
class DNode {  
public:  
    int data;  
    DNode* prev;  
    DNode* next;  
    DNode(int val) : data(val), prev(nullptr), next(nullptr) {}  
};
```

(

Circular Linked List

The circular linked list is another form of linked list that provides additional flexibility in traversing elements.

A circular linked list may be singly-linked or doubly-linked, but its distinguishing feature is that it has no tail.

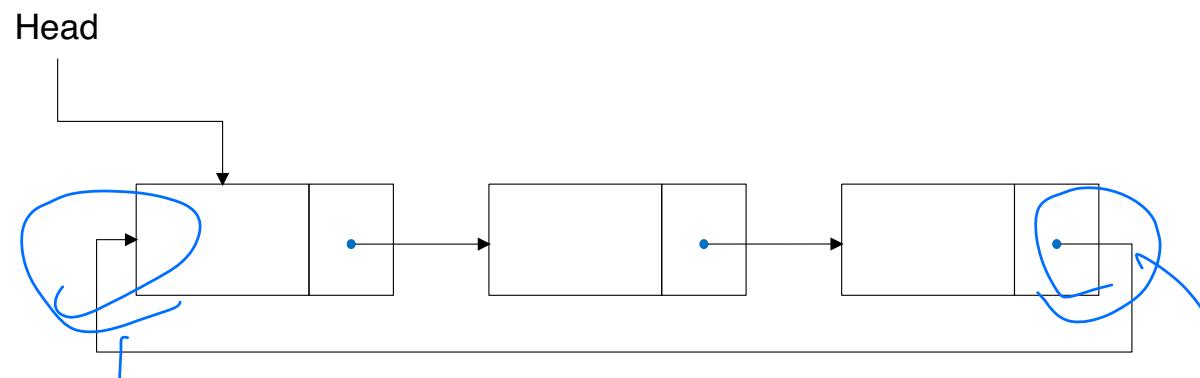
Improvement:

可是單向 (僅 next)

双向 (next, prev)

但 結尾 指向 head

Circular Linked List



Circular Linked List in C

```
struct CNode {  
    int data;  
    struct CNode *next;  
};
```

Circular Linked List in C++

```
class CNode {  
public:  
    int data;  
    CNode* next;  
    CNode(int val) : data(val), next(nullptr) {}  
};
```

Trade-offs: Variations of Linked List

1. traversal direction
2. insert/delete flexibility
3. memory overhead

Comparison: Array vs. Linked List

Aspect	Array	Linked List
Memory Allocation	Contiguous block of memory	Non-contiguous, each node allocated dynamically
Size 存取 索引 Access (Indexing)	Fixed (static) – must be defined at declaration (in C) Direct access with index ($O(1)$)	Dynamic – can grow/shrink at runtime Sequential access only ($O(n)$)
插入 删除 Insertion/Deletion	Costly ($O(n)$) due to shifting elements	Efficient ($O(1)$) if pointer to node is known → 地址 pointer 已知
Memory Usage 快取 Cache Performance	No extra overhead	Requires extra memory for pointers 分散
Implementation Simplicity	Easier to implement and use	More complex due to pointer handling
Use Cases	When size is known and frequent random access is needed	When frequent insertion/deletion is required and size is unpredictable

Comparison: Array vs. Linked List

Operation	Array (Dynamic)	Linked List
Access (by index)	$O(1)$	$O(n)$ (must traverse)
Search	$O(n)$	$O(n)$
Insert at front	$O(n)$	$O(1)$
Insert at middle	$O(n)$	$O(n)$
Insert at end	$O(1)$ amortized	$O(n)^*$ or $O(1)^{**}$
Delete	$O(n)$	$O(n)$ or $O(1)^{***}$

- * $O(n)$ unless tail pointer is used
- ** With tail pointer and singly linked list
- *** If node pointer is known, deletion is $O(1)$

std::list

Element access

front access the first element
(public member function)

back access the last element
(public member function)

Iterators

begin
cbegin (C++11) returns an iterator to the beginning
(public member function)

end
cend (C++11) returns an iterator to the end
(public member function)

rbegin
crbegin (C++11) returns a reverse iterator to the beginning
(public member function)

rend
crend (C++11) returns a reverse iterator to the end
(public member function)

Capacity

empty checks whether the container is empty
(public member function)

size returns the number of elements
(public member function)

max_size returns the maximum possible number of elements
(public member function)

Modifiers

clear clears the contents
(public member function)

insert inserts elements
(public member function)

insert_range (C++23) inserts a range of elements
(public member function)

emplace (C++11) constructs element in-place
(public member function)

erase erases elements
(public member function)

push_back adds an element to the end
(public member function)

emplace_back (C++11) constructs an element in-place at the end
(public member function)

append_range (C++23) adds a range of elements to the end
(public member function)

pop_back removes the last element
(public member function)

push_front inserts an element to the beginning
(public member function)

emplace_front (C++11) constructs an element in-place at the beginning
(public member function)

prepend_range (C++23) adds a range of elements to the beginning
(public member function)

pop_front removes the first element
(public member function)

resize changes the number of elements stored
(public member function)

swap swaps the contents
(public member function)

Operations

merge merges two sorted lists
(public member function)

splice transfers elements from another list
(public member function)

remove
remove_if removes elements satisfying specific criteria
(public member function)

reverse reverses the order of the elements
(public member function)

unique removes consecutive duplicate elements
(public member function)

sort sorts the elements
(public member function)

Key Observation of Linked List

```
struct Node {  
    int data;  
    struct Node* next;  
};
```

```
struct Node* temp = head;  
int i = 0;  
while (temp != NULL) {  
    printf("Node %d -> Value: %d, Address: %p, Next: %p\n", i, temp->data, (void*)temp,  
    (void*)temp->next);  
    temp = temp->next;  
    i++;  
}
```

Key Observation of Linked List

Node Position

- head (start of the list)
- head->next
- head->next

Value of Each Node

- head->data → 0
- head->next->data → 1
- head->next->next->data → 2

Node Memory Location (Node Address)

- (void *) head
- (void *) head->next
- (void *) head->next->next

Next Pointer Address (Link to Next Node)

- (void *) head->next
- (void *) head->next->next

Summary

1. Create: always dynamic

- Memory is allocated node-by-node at runtime.
- Each node contains data + pointer to next node.

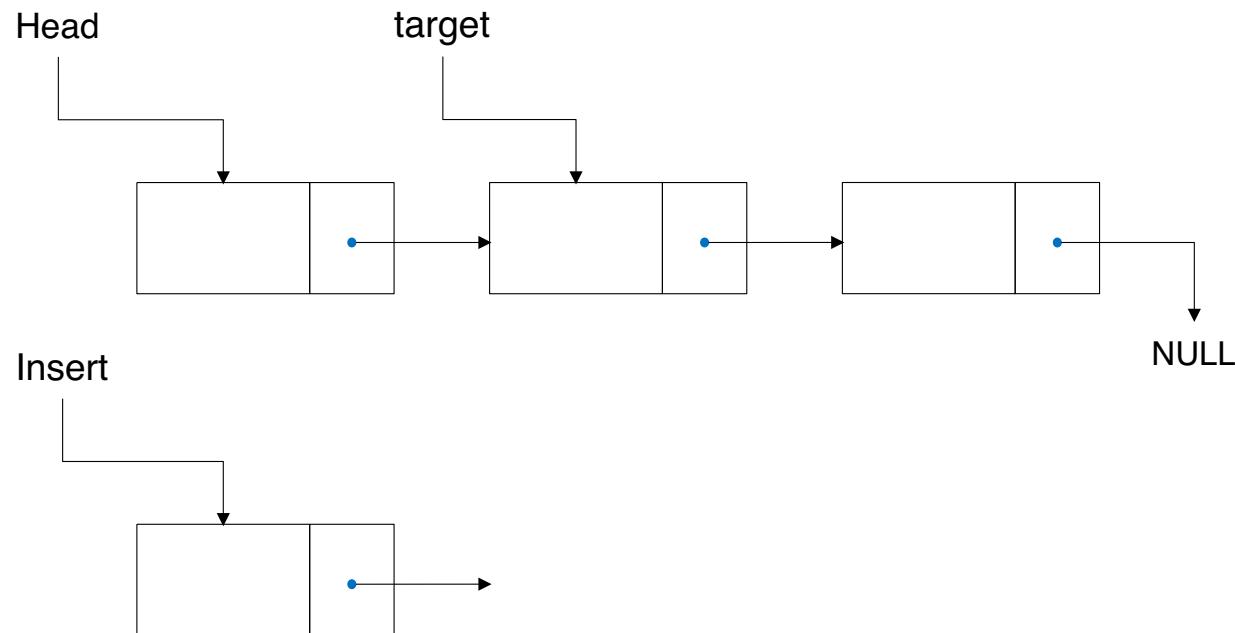
2. Retrieve: 只有在循環存取可以檢索

- Sequential Access Only
- Search the target
 - Unsorted
 - Sorted

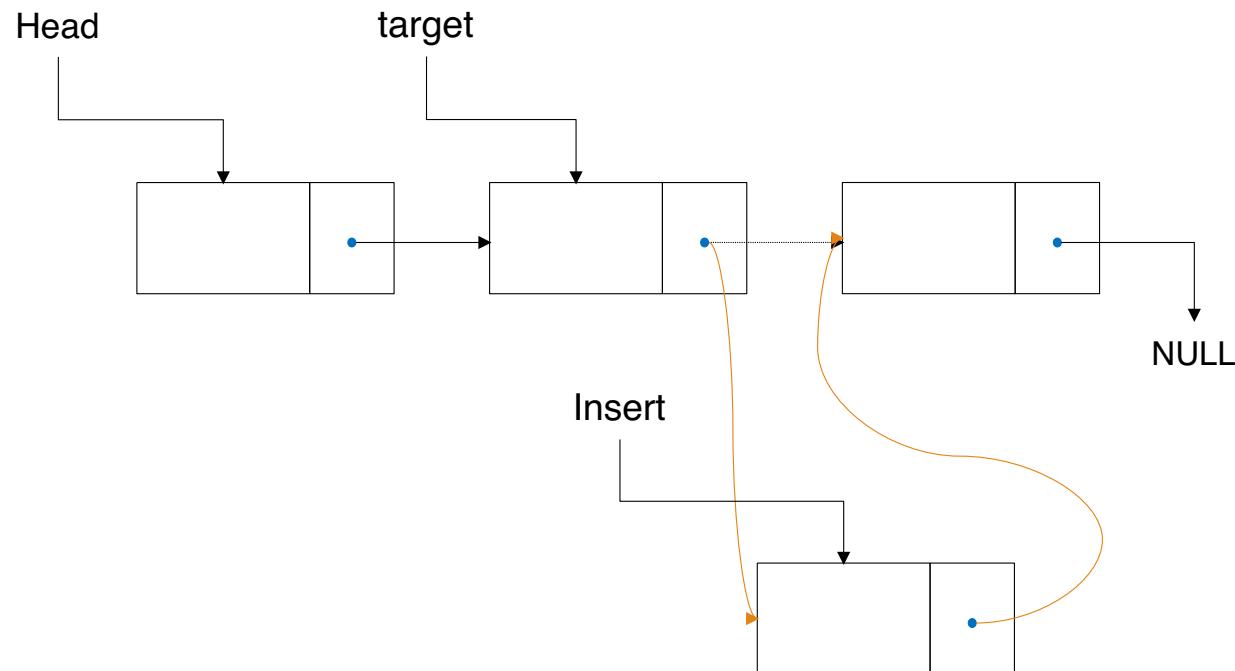
3. Update

- Target: Update data of a found node.
- Insert
 - At head
 - At tail
 - In middle
- Delete
 - Remove by value
 - Remove by position
- Special cases: head/tail

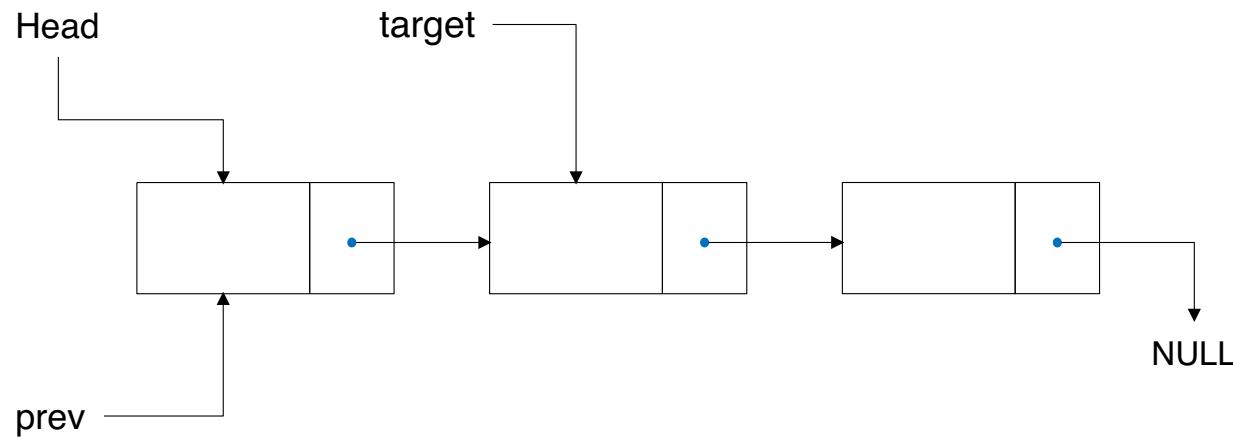
Linked List: Insert_after_Target()



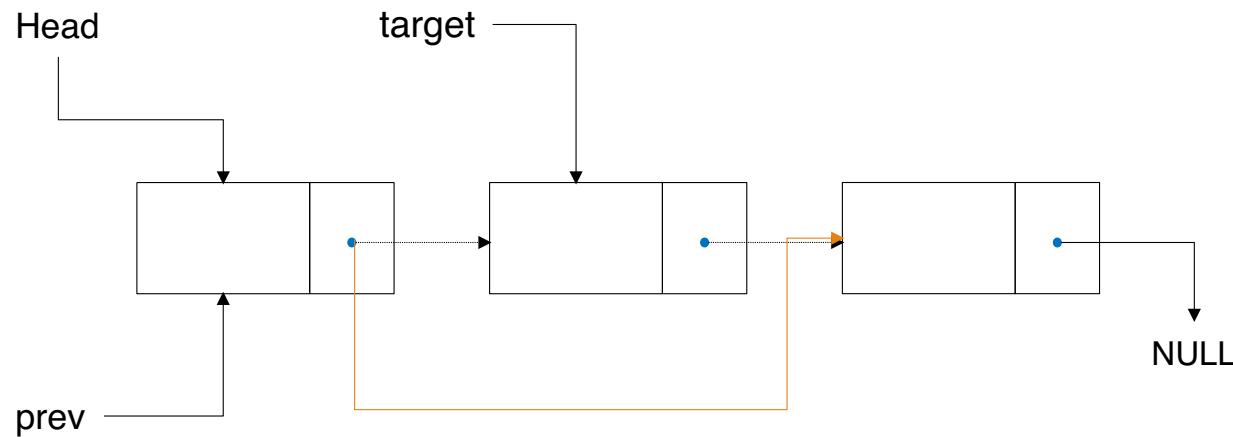
Linked List: Insert_after_Target()



Linked List: Delete()



Linked List: Delete()



Variation & Overhead

Extra pointer for quick operations: head & tail

Operations (ADTs)

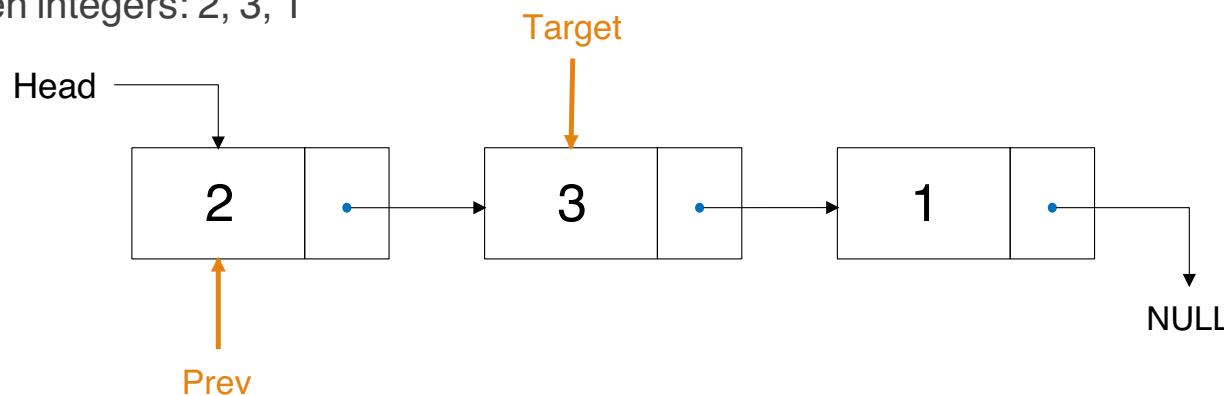
- Insert before
- Insert after
- Insert_first
- Insert_last
- Delete_by_value
- Delete_by_position
- Delete_first
- Delete_last



Image credit: <https://uxwing.com/idea-icon/>

Node Exchange in Linked List

Given integers: 2, 3, 1

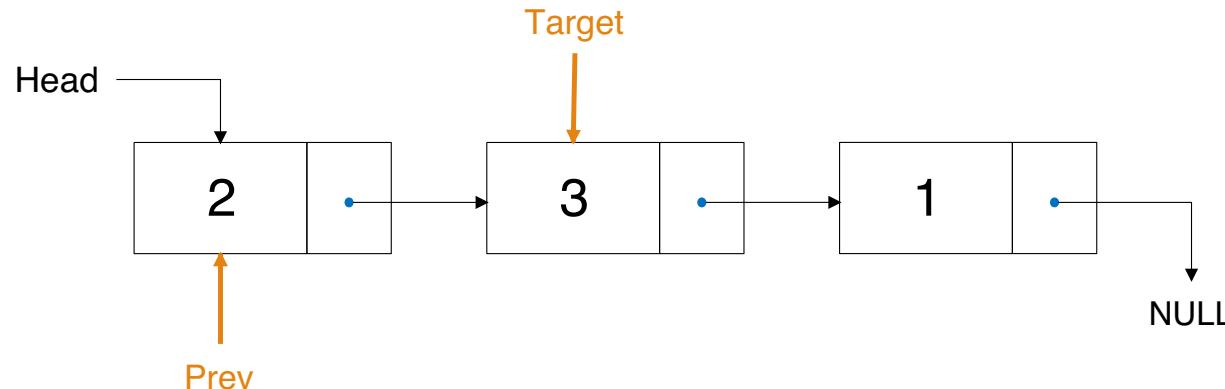


Question: Exchange Prev and Target

Prev: Prev
 Prev->next

Target: Target
 Target->next

Wrong

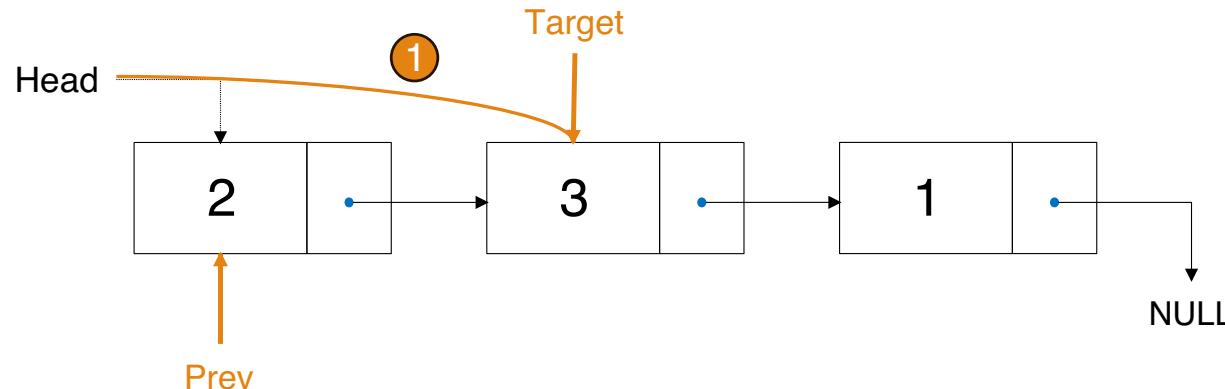


Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Target->next = Prev
3. Prev->next = Target->next (Prev)

Wrong

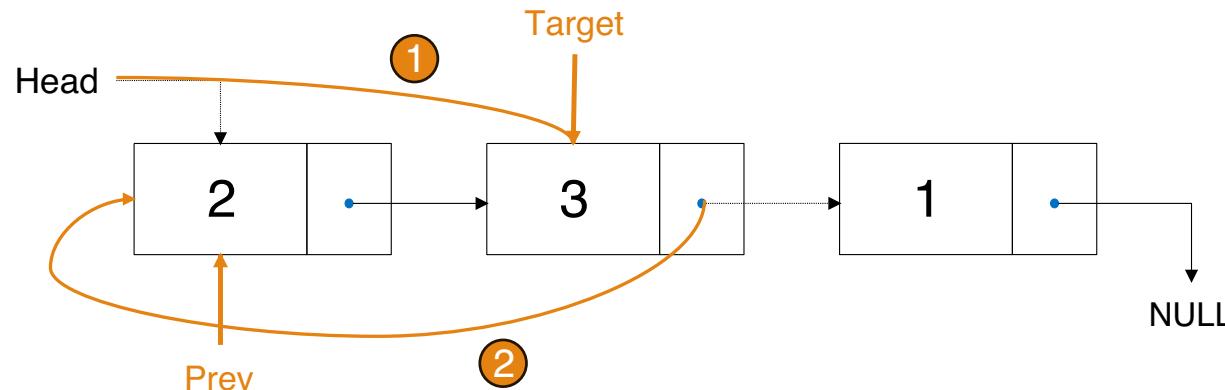


Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Target->next = Prev
3. Prev->next = Target->next (Prev)

Wrong

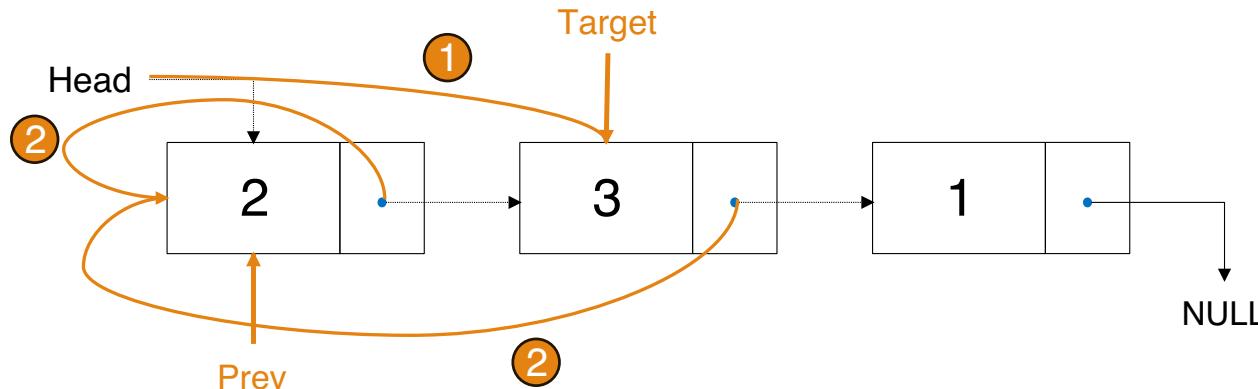


Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Target->next = Prev
3. Prev->next = Target->next (Prev)

Wrong

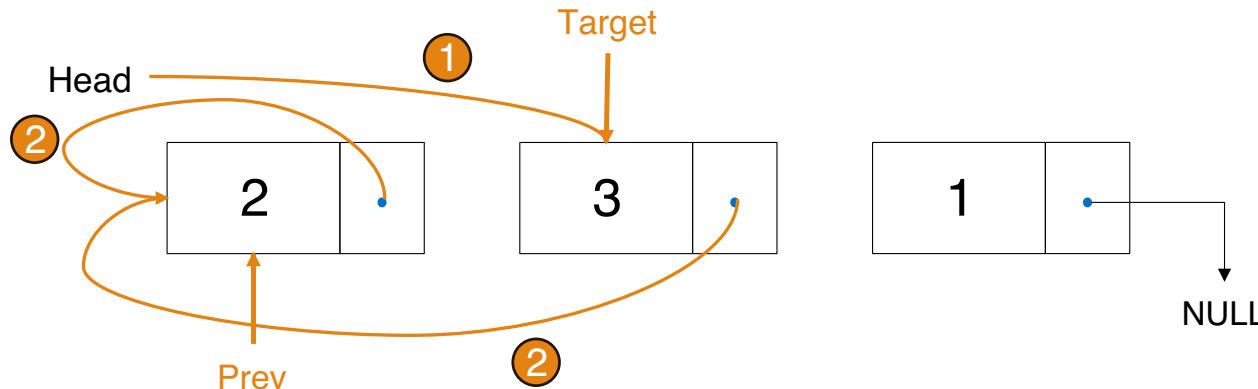


Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Target->next = Prev
3. Prev->next = Target->next (Prev)

Wrong (Final)



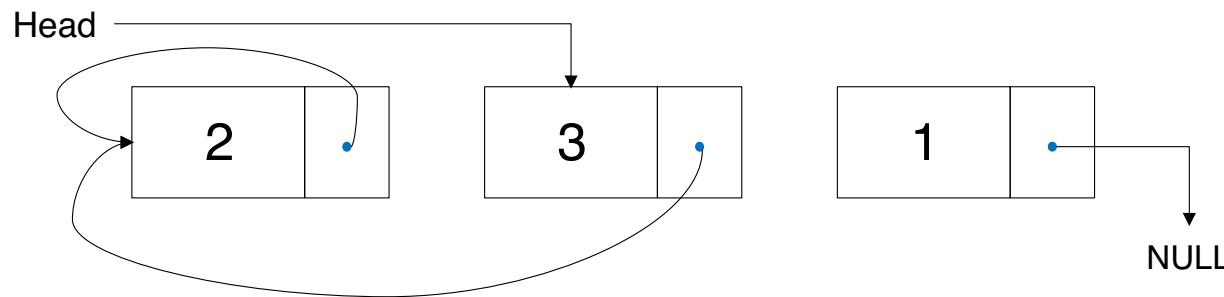
Prev: Prev
 Prev->next

Target: Target
 Target->next

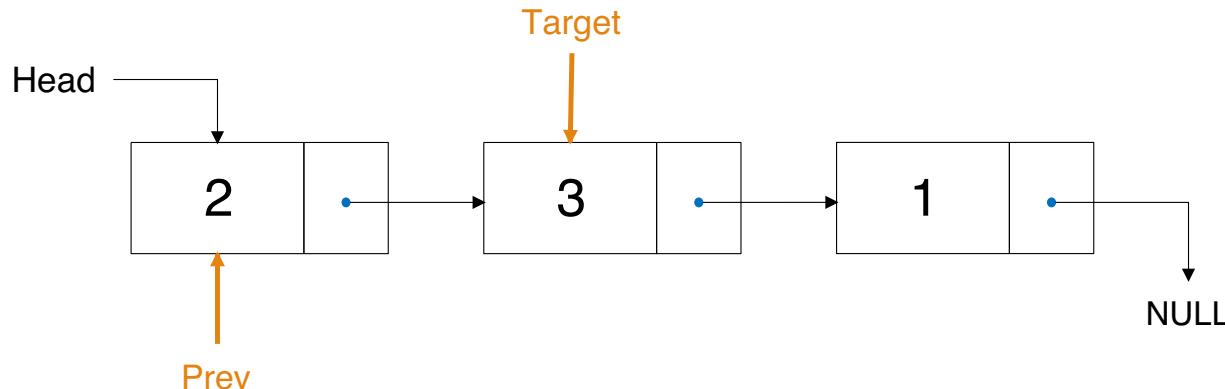
Steps

1. Head = Prev Head = Target
2. Target->next = Prev
3. Prev->next = Target->next (Prev)

Wrong (Final)



Correct



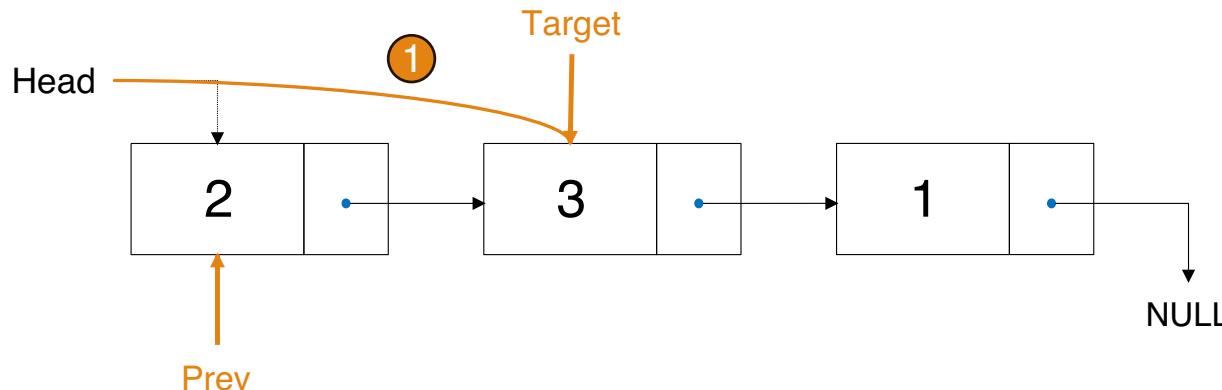
Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps

1. Head = Prev \sqcap Head = Target
2. Prev->next = Target->next
3. Target->next = Prev

Correct

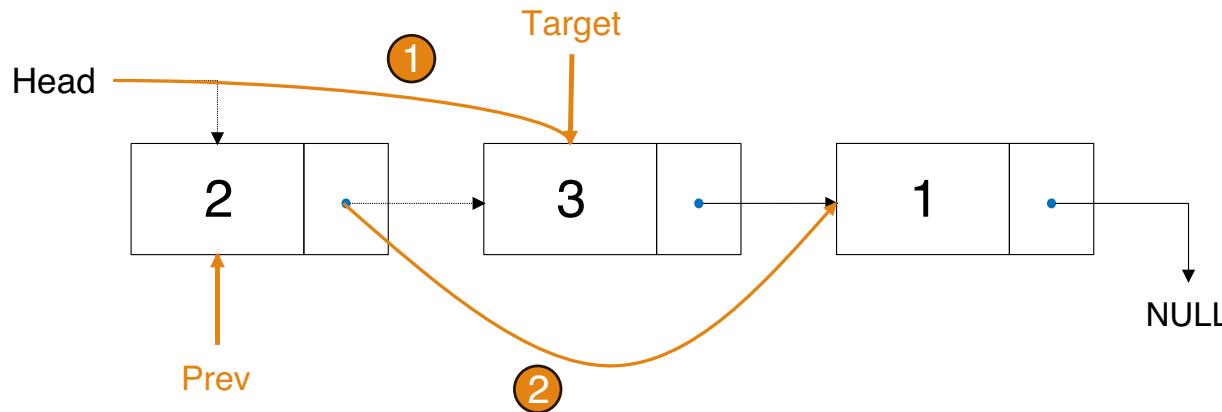


Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Prev->next = Target->next
3. Target->next = Prev

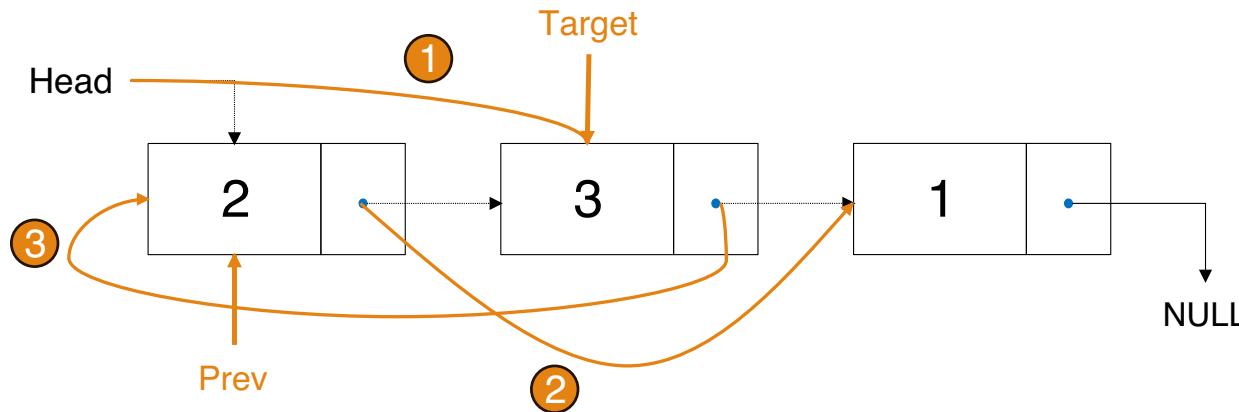
Correct



Prev: Prev
 Prev->next
Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Prev->next = Target->next
3. Target->next = Prev

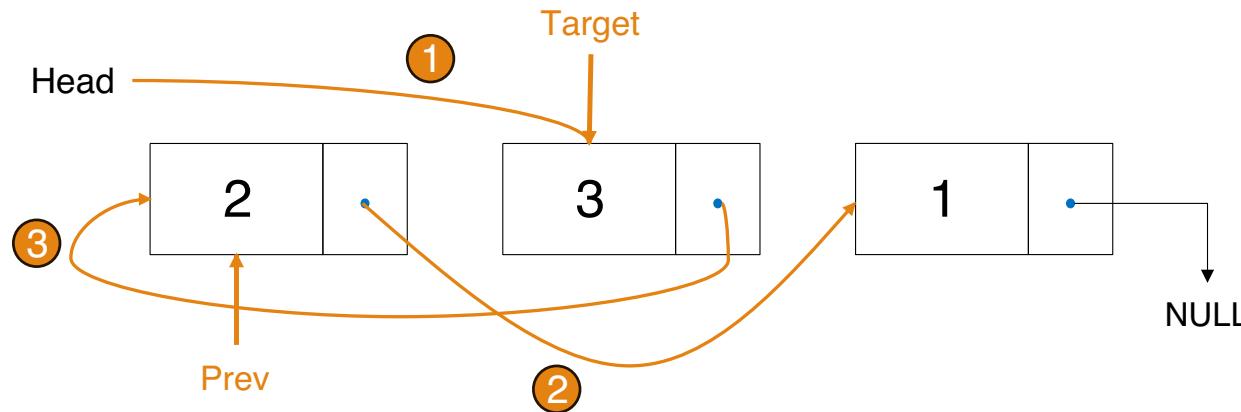
Correct



Prev: Prev
 Prev->next
Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Prev->next = Target->next
3. Target->next = Prev

Correct (Final)

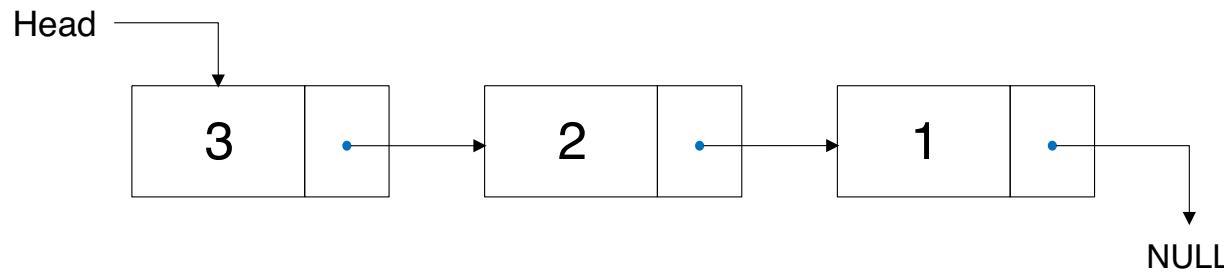


Prev: Prev
 Prev->next

Target: Target
 Target->next

Steps
1. Head = Prev Head = Target
2. Prev->next = Target->next
3. Target->next = Prev

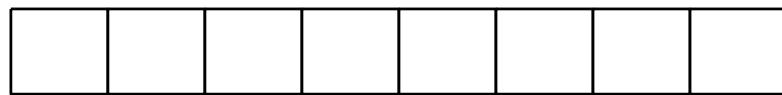
Correct (Final)



Node Exchange in Linked List

Correct vs. Wrong: pointer dependence

Think: the Integer Array, Sorting



front

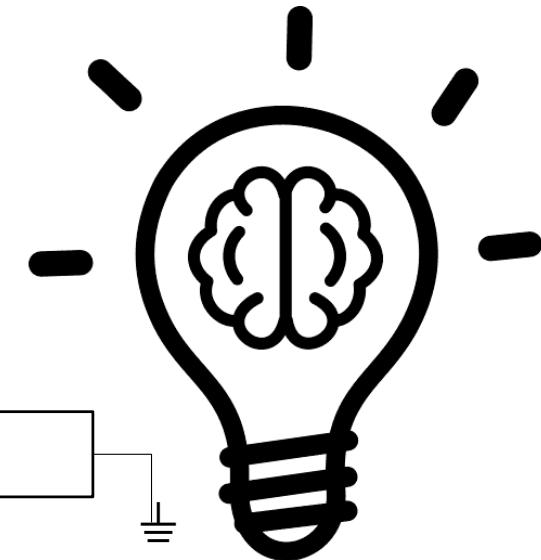
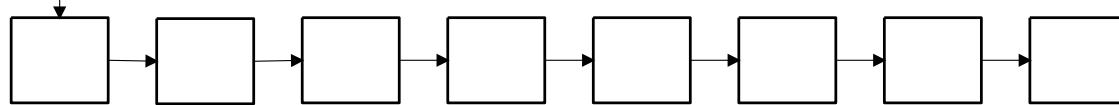
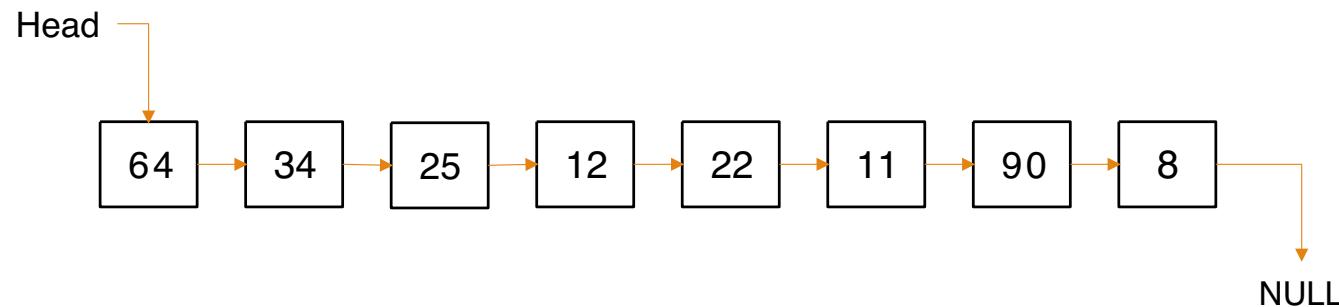


Image credit: <https://uxwing.com/idea-icon/>

Linked List Sorting

Think: the Integer Linked List

64	34	25	12	22	11	90	8
----	----	----	----	----	----	----	---



Selection Sort - Array

```
procedure selectionSort(A[1..n]):  
    for i from 1 to n-1:  
        minIndex = i  
        for j from i+1 to n:  
            if A[j] < A[minIndex]:  
                minIndex = j  
        swap A[i] and A[minIndex]
```

Selection Sort - Array

```
procedure selectionSort(A[1..n]):  
    for i from 1 to n-1:  
        minIndex = i  
        for j from i+1 to n:  
            if A[j] < A[minIndex]:  
                minIndex = j  
        swap A[i] and A[minIndex]
```

A0	64	34	25	12	22	11	90	8
A1	8	34	25	12	22	11	90	64
A2	8	11	25	12	22	34	90	64
A3	8	11	12	25	22	34	90	64
A4	8	11	12	22	25	34	90	64
A5	8	11	12	22	25	34	90	64
A6	8	11	12	22	25	34	90	64
A7	8	11	12	22	25	34	64	90

Selection Sort - Linked List

Linked List

- Swap the value
- Swap the pointer

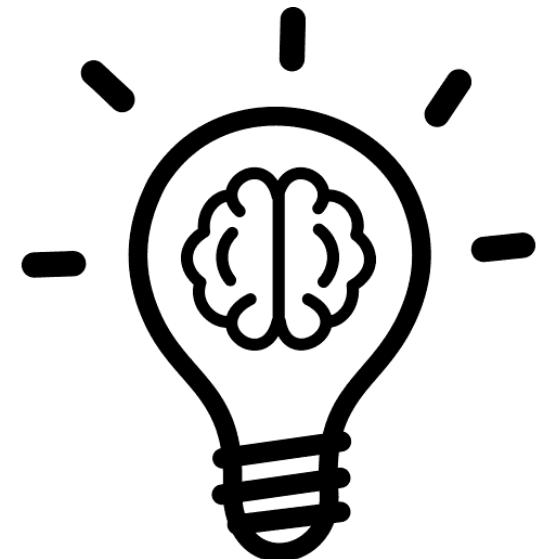


Image credit: <https://uxwing.com/idea-icon/>

Selection Sort - Linked List

Swap the value

```
void selectionSortValue(Node* head){  
    for(Node* i=head; i && i->next; i=i->next){  
        Node* minNode=i;  
        for(Node* j=i->next; j; j=j->next){  
            if(j->val < minNode->val) minNode=j;  
        }  
        if(minNode!=i){  
            int t=i->val; i->val=minNode->val; minNode->val=t;  
        }  
    }  
}
```

Selection Sort - Linked List

Swap the pointer

```
Node* selectionSortPointer(Node* head){  
    if(!head) return head;  
    Node dummy = {0, head};  
    Node* sortedTail = &dummy;  
  
    while(sortedTail->next){  
        Node* minPrev = sortedTail;  
        for(Node* p=sortedTail->next; p && p->next; p=p->next){  
            if(p->next->val < minPrev->next->val) minPrev = p;  
        }  
        Node* minNode = minPrev->next;  
        // detach minNode  
        minPrev->next = minNode->next;  
        // insert after sortedTail  
        minNode->next = sortedTail->next;  
        sortedTail->next = minNode;  
        sortedTail = sortedTail->next;  
    }  
    return dummy.next;  
}
```

Bonus: AI Prompt for Studying “Linked List” (1)

I understand basic singly linked lists and can create, traverse, and add nodes. Now I want to advance to intermediate level. Please help me with:

1. Advanced List Types :

- Doubly linked lists - when and why to use them
- Circular linked lists - practical applications
- Singly vs doubly vs circular - trade-offs and comparisons

2. Complex Operations :

- Insertion at any position (beginning, middle, end)
- Deletion operations (by value, by position, specific cases)
- Searching and finding nodes efficiently
- Reversing a linked list (iterative and recursive approaches)

Bonus: AI Prompt for Studying “Linked List” (2)

3. Algorithm Implementation:

- Merge two sorted linked lists
- Detect cycles in linked lists (Floyd's cycle detection)
- Find middle element efficiently
- Remove duplicates from sorted/unordered lists

4. Memory Management & Error Handling :

- Proper memory allocation and deallocation
- Handling edge cases (empty lists, single nodes)
- Defensive programming techniques
- Memory leak detection and prevention

Bonus: AI Prompt for Studying “Linked List” (3)

5. Performance Analysis:

- Time complexity of different operations
- Space complexity considerations
- When to use linked lists vs arrays vs vectors
- Performance trade-offs in real applications

6. Advanced Techniques:

- Using dummy/sentinel nodes to simplify code
- Two-pointer techniques (slow/fast pointers)
- Recursive vs iterative approaches
- Generic linked lists using templates

Bonus: AI Prompt for Studying “Linked List” (4)

7. Practical Applications:

- Implementing stacks and queues using linked lists
- Undo functionality in applications
- Hash table collision resolution with chaining
- Music playlist or browser history implementation