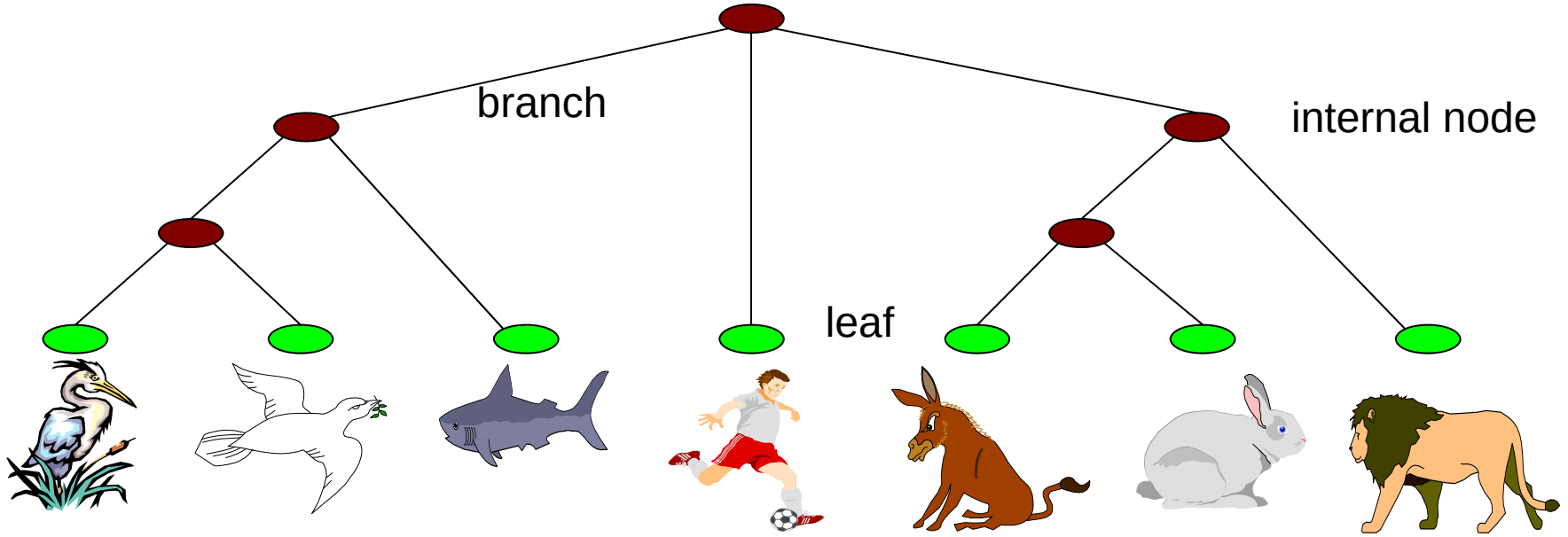
第4章：进化树构建的概率方法

- 问题介绍
- 进化树构建方法的概率方法

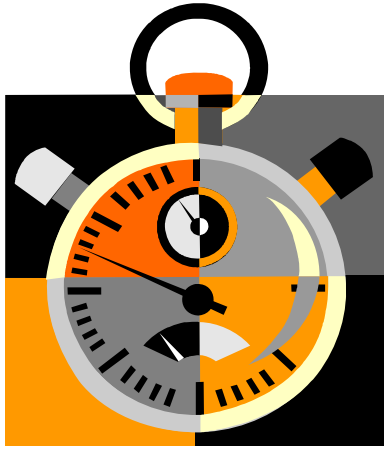
部分Slides修改自University of Basel的Michael Springmann课程“CS302 Seminar Life Science Informatics”的讲义

Phylogenetic Tree



- Topology: bifurcating
  - Leaves - *1…N*
  - Internal nodes *N+1…2N-2*
- Branch length

Molecular Clock Hypothesis



- Amount of genetic difference between sequences is a function of time since separation.
- Rate of molecular change is constant (enough) to predict times of divergence

Likelihood of a Tree

- Given:
  - n aligned sequences M= $X_1,\ldots,X_n$
  - A tree T, leaves labeled with $X_1,\ldots,X_n$
- Reconstruction t*:
  - Labeling of internal nodes
  - Branch lengths

Goal: Find optimal reconstruction t* : One maximizing the likelihood P(M|T, t*)

Probabilistic Methods

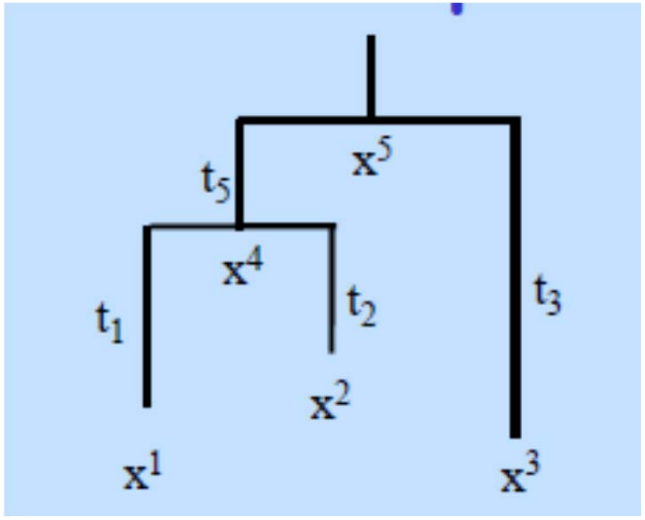
- The phylogenetic tree represents a generative probabilistic model (like HMMs) for the observed sequences.
- Background probabilities: q(a)
- Mutation probabilities: P(a|b, t)
- Models for evolutionary mutations
  - Jukes Cantor
  - Kimura 2-parameter model
- Such models are used to derive the probabilities

Probabilistic Model

- Assumptions:
  - Each character is independent
  - The branching is a Markov process: The probability that a node x has a specific label is only a function of the parent node y and the branch length t between them
  - The probabilities P(x|y,t) are known

Example

- Given then tree



$$P(x_1, x_2, x_3, x_4, x_5 | T, t^*)$$
$$= P(x_1|x_4, t_1) P(x_2|x_4, t_2) P(x_3|x_5, t_3) P(x_4|x_5, t_5)$$

连续时间马氏链

- 随机变量族 $\{\xi_t(\omega), t \geq 0\}$ 称为连续时间马氏链， 若 ξ_n 是离散的(其状态空间至多是可数集，即有限或者与自然数一一对应)，而且对于 $\forall m \geq 0, \forall s > s_1 \geq \cdots \geq s_m \geq 0$ 及任意状态序列 $i, j, i_1, \cdots, i_m$， 都有

$$\begin{aligned} & P(\xi_{t+s} = j \,|\, \xi_s = i, \cdots, \xi_{s_1} = i_1, \cdots, \xi_{s_m} = i_m) \\ =& P(\xi_{t+s} = j \,|\, \xi_s = i) \end{aligned}$$

转移概率矩阵

- 定义：对连续时间马氏链 $\{\xi_t(\omega), t \geq 0\}$，记

$$p_{ij}(s,t) = P(\xi_t = j | \xi_s = i), \forall t \geq s$$
$$P(s,t) = (p_{ij}(s,t))$$

称无穷矩阵P(s,t)为转移概率矩阵

- 易证转移概率满足Chapman-Kolmogorov方程

$$p_{ij}(s,u) = \sum_k p_{ik}(s,t) p_{kj}(t,u)$$

时齐的连续时间马氏链

- 定义：如果连续时间马氏链的转移概率矩阵P(s,s+t)与起始时间s无关，称此时的马氏链为时齐的。那么此时的转移概率矩阵就可用一个时间参数来标度

$$P(t-s) = P(s,t)$$

- 相应地，chapman-kolmogorov方程简化为

$$P(t+s) = P(s)P(t)$$

连续时间马氏链的有限维分布

- 类似于离散时间马氏链，连续时间马氏链的有限维分布有初始分布和转移概率唯一确定。设 $\mu_i(0) = P(\xi_0 = i)$, 转移概率矩阵P(t), 那么对 $\forall 0 < t_1 < t_2 < \cdots < t_n,$

$$
\begin{aligned}
&P(\xi_0 = i_0, \xi_{t_1} = i_1, \cdots, \xi_{t_n} = i_n) \\
&= \mu_{i_0}(0) p_{i_0 i_1}(t_1) p_{i_1 i_2}(t_2 - t_1) \cdots p_{i_{n-1} i_n}(t_n - t_{n-1})
\end{aligned}
$$

绝对概率

- 令 $\mu_i(t) = P(\xi_t = i)$, 称向量

$$\mu(t) = \begin{pmatrix} \mu_1(t) \\ \vdots \\ \mu_i(t) \\ \vdots \end{pmatrix}$$

为t时刻的绝对概率

转移速率矩阵(Q矩阵)

- 马氏链中时齐马氏链的多步转移矩阵是最小的转移矩阵P的次方。
- 连续情况下对应于无穷小生成元，但此时需要加条件
- 定义：设 $P(t) \to I, t \to 0,$ 若转移概率矩阵在0点的导数存在，记为$P'(0) = Q$, 称之为转移概率矩阵的转移速率矩阵(简称为Q矩阵)。

有限状态情形

- 当状态数有限时，如Q矩阵满足 $\sum_j q_{ij} = 0$ (即所谓保守)，此时转移概率矩阵由Q矩阵唯一确定

$$P(t) = e^{Qt} = I + \frac{Qt}{1!} + \frac{Q^2t^2}{2!} + \cdots + \frac{Q^kt^k}{k!} + \cdots$$

- 通常转移速率矩阵Q已知，我们可以通过指数矩阵给出转移概率矩阵P(t).

Kolmogorov方程和Master方程

- Kolmogorov向前方程(Fokker-Planck方程)：

$$P'(t) = P(t)Q, P(0) = I$$

- Kolmogorov向后方程：

$$P'(t) = QP(t), P(0) = I$$

- Master方程

$$(\mu_1'(t), \mu_2'(t), \cdots,) = (\mu_1(t), \mu_2(t), \cdots,)Q$$

其中 $\mu_k(t) = P(\xi_t = k)$ 为绝对概率

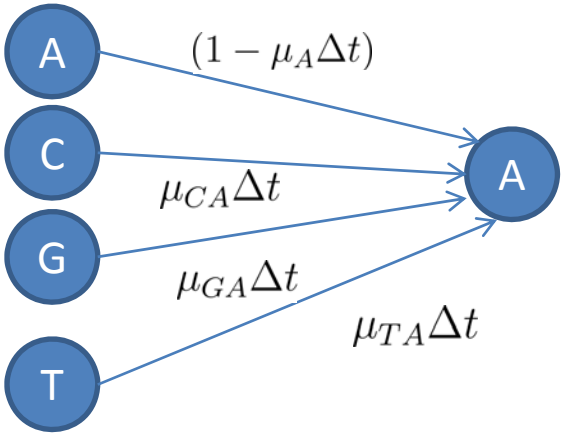
Molecular Evolution

Q: How can we model evolution on nucleotide level? (ignore gaps, focus on substitutions)

A: Consider what happens at a specific position for small time interval Δt

- p(t) = vector of probabilities of {A,C,G,T} at time t
- μ_{AC} = rate of transition from A to C per unit time
- $\mu_A = \mu_{AC} + \mu_{AG} + \mu_{AT}$ rate of transition out of A

Molecular Evolution



t 时刻　　t+Δt 时刻

$$p_A(t+\Delta t) = p_A(t)(1-\mu_A\Delta t) + p_C(t)\mu_{CA}\Delta t + p_G(t)\mu_{GA}\Delta t + p_T(t)\mu_{TA}\Delta t$$

Molecular Evolution

- 同理可得

$$p_C(t+\Delta t) = p_C(t)(1-\mu_C\Delta t) + p_A(t)\mu_{AC}\Delta t + p_G(t)\mu_{GC}\Delta t + p_T(t)\mu_{TC}\Delta t$$

$$p_G(t+\Delta t) = p_G(t)(1-\mu_G\Delta t) + p_A(t)\mu_{AG}\Delta t + p_C(t)\mu_{CG}\Delta t + p_T(t)\mu_{TG}\Delta t$$

$$p_T(t+\Delta t) = p_T(t)(1-\mu_T\Delta t) + p_A(t)\mu_{AT}\Delta t + p_C(t)\mu_{CT}\Delta t + p_G(t)\mu_{GT}\Delta t$$

Molecular Evolution

In vector notation, we get

$$p(t + \Delta t) = p(t) + P(t)Q\Delta t$$

where Q is the substitution rate matrix

$$Q = \begin{pmatrix} -\mu_A & \mu_{AG} & \mu_{AC} & \mu_{AT} \\ \mu_{GA} & -\mu_G & \mu_{GC} & \mu_{GT} \\ \mu_{CA} & \mu_{CG} & -\mu_C & \mu_{CT} \\ \mu_{TA} & \mu_{TG} & \mu_{TC} & -\mu_T \end{pmatrix}$$

Molecular Evolution

- 写成列向量形式

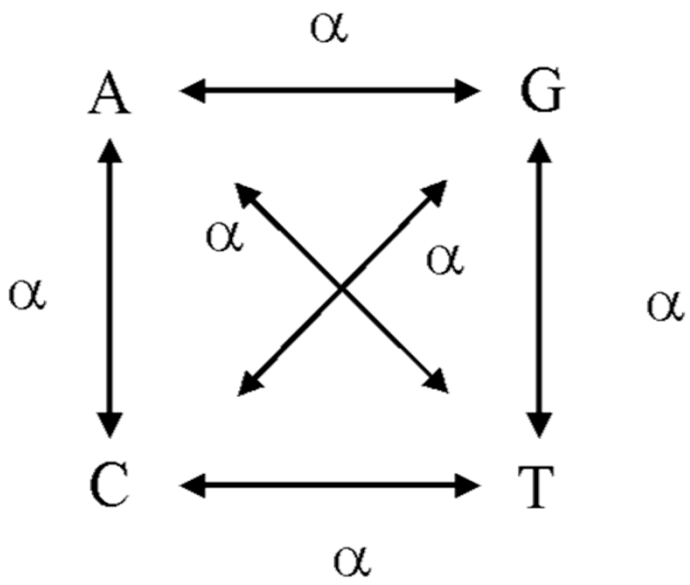
$$\vec{p}(t+\Delta t) = \vec{p}(t) + Q^T \vec{p}(t)$$

- 取转置后即为Master方程
- 其中Q为转移速率矩阵

$$Q = \begin{pmatrix} -\mu_A & \mu_{AG} & \mu_{AC} & \mu_{AT} \\ \mu_{GA} & -\mu_G & \mu_{GC} & \mu_{GT} \\ \mu_{CA} & \mu_{CG} & -\mu_C & \mu_{CT} \\ \mu_{TA} & \mu_{TG} & \mu_{TC} & -\mu_T \end{pmatrix}$$

Jukes Cantor model

- Mutation occurs at a constant rate
- Each nucleotide is equally likely to mutate into any other nucleotide with rate a.



$$Q = \begin{pmatrix} -3\alpha & \alpha & \alpha & \alpha \\ \alpha & -3\alpha & \alpha & \alpha \\ \alpha & \alpha & -3\alpha & \alpha \\ \alpha & \alpha & \alpha & -3\alpha \end{pmatrix}$$

Substitution Matrix

- 由对称性，可设

$$P(t) = \begin{pmatrix} \gamma(t) & s(t) & s(t) & s(t) \\ s(t) & \gamma(t) & s(t) & s(t) \\ s(t) & s(t) & \gamma(t) & s(t) \\ s(t) & s(t) & s(t) & \gamma(t) \end{pmatrix}$$

- 由Kolmogorov后向方程

$$\frac{dP(t)}{d(t)} = QP(t)$$

Substitution Matrix

- 可得方程

$$\begin{cases} \dfrac{d\gamma(t)}{dt} = -3\alpha\gamma(t) + 3\alpha s(t) \\ \dfrac{ds(t)}{dt} = -\alpha s(t) + \alpha\gamma(t) \end{cases}$$

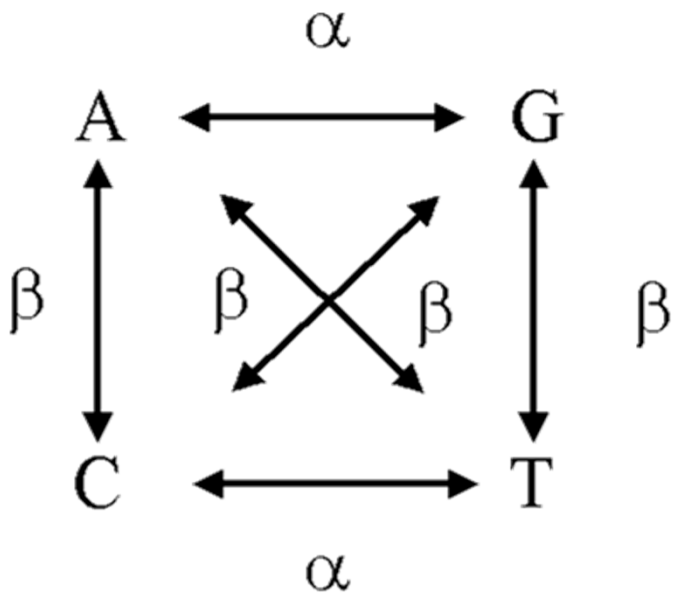
- 容易求得

$$\gamma(t) = \frac{1}{4}(1 + 3e^{-4\alpha t})$$

$$s(t) = \frac{1}{4}(1 - e^{-4\alpha t})$$

Kimura 2-parameter Model

- Allows a different rate for transitions and transversions.



$$Q = \begin{pmatrix} -\alpha - 2\beta & \beta & \alpha & \beta \\ \beta & -\alpha - 2\beta & \beta & \alpha \\ \alpha & \beta & -\alpha - 2\beta & \beta \\ \beta & \alpha & \beta & -\alpha - 2\beta \end{pmatrix}$$

Substitution Matrix

- 由对称性，可设转移概率矩阵

$$P(t) = \begin{pmatrix} \gamma(t) & s(t) & u(t) & s(t) \\ s(t) & \gamma(t) & s(t) & u(t) \\ u(t) & s(t) & \gamma(t) & s(t) \\ s(t) & u(t) & s(t) & \gamma(t) \end{pmatrix}$$

- 由Kolmogorov后向方程

$$\frac{dP(t)}{d(t)} = QP(t)$$

Substitution Matrix

- 可得方程

$$\begin{cases} \dfrac{d\gamma(t)}{dt} = -(2\beta+\alpha)\gamma(t) + 2\beta s(t) + \alpha u(t) \\ \dfrac{ds(t)}{dt} = -2\beta s(t) + \beta\gamma(t) + \beta u(t) \\ \dfrac{du(t)}{dt} = -(2\beta+\alpha)u(t) + 2\beta s(t) + \alpha\gamma(t) \end{cases}$$

- 容易求得

$$\begin{cases} s(t) = \dfrac{1}{4}(1 - e^{-4\beta t}) \\ s(t) = \dfrac{1}{4}(1 + e^{-4\beta t} - 2e^{-2(\alpha+\beta)t}) \\ \gamma(t) = 1 - 2s(t) - u(t) \end{cases}$$

Substitution Matrix: General Case

- 对于对称矩阵Q可以对角化, 即存在正交矩阵U, 和特征值 $\lambda_1 \geq \cdots \geq \lambda_n$, 使得

$$Q = U^T diag\{\lambda_1, \cdots, \lambda_n\} U$$

- 于是

$$P(t) = U^T diag\{e^{\lambda_1 t}, \cdots, e^{\lambda_n t}\} U$$

PAM矩阵

- Point accepted mutation (Dayhoff et al 1978)
- Given an tree of protein family, the frequence matrix A_{ab} counting the occurrence of an “a” in the ancestral sequence was replaced by a “b” in the descendant.
- Estimate the conditional probability p(b|a)

$$P(b|a) = B_{a,b} = \frac{A_{ab}}{\sum_c A_{ac}}$$

PAM矩阵

- Scaling B

$$C_{ab} = \sigma B_{ab}, C_{aa} = \sigma B_{aa} + (1 - \sigma)$$

- Such that the expected number of substitution is 1%, i.e.

$$\sum_{ab} q_a q_b C_{ab} = 0.01$$

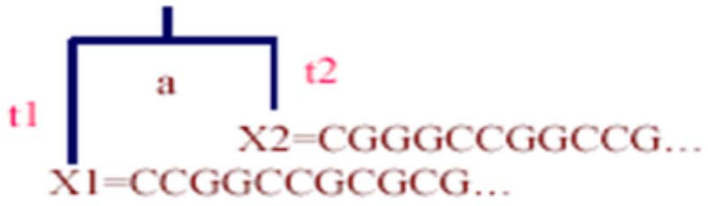
- Then the PAM(1) matrix is given by

$$S(1) = (C_{ab})$$

Calculating the Likelihood for Ungapped Alignments

$$P(X^1, X^2|T, t_1, t_2) = \prod_{u=1}^{N} P(X_u^1, X_u^2|T, t_1, t_2)$$

$$P(X_u^1, X_u^2|T, t_1, t_2) = \sum_a q_a P(X_u^1, |a, t_1) P(X_u^2|a, t_2)$$



- 假设突变符合JC model, 等初始概率 $q_A = q_C = q_G = q_T = \frac{1}{4}$

$$P(C, C|T, t_1, t_2) = q_c\gamma(t_1)\gamma(t_2) + q_G s(t_1)s(t_2) + q_a s(t_1)s(t_2) + q_T s(t_1)s(t_2)$$
$$= \frac{1}{3}(r(t_1)r(t_1) + 3S(t_1)S(t_2))$$
$$P(C, G|T, t_1, t_2) = P(G, C|T, t) = \frac{1}{4}(\gamma(t_1)s(t_1) + s(t_1)\gamma(t_2) + 2s(t_1)s(t_2))$$
$$P(X^1, X^2|T, t_1, t_2) = 16^{-(n_1+n_2)}(1 + 3e^{-4\alpha(t_1+t_2)})^{n_1}(1 - e^{-4\alpha(t_1+t_2)})^{n_2}$$

其中n1是匹配数，n2是不匹配数目.

Calculating the Likelihood for Ungapped Alignments

- n sequences of length N, site u=1...N
- Given a rooted tree contains *2n - 1* nodes*, 1... n* being the leaf nodes, n+1 ... 2n-1 non-leaf, tree lengths t1, …, t_{2n-1}.
- Let a(i) denote the ancestor of node a^i

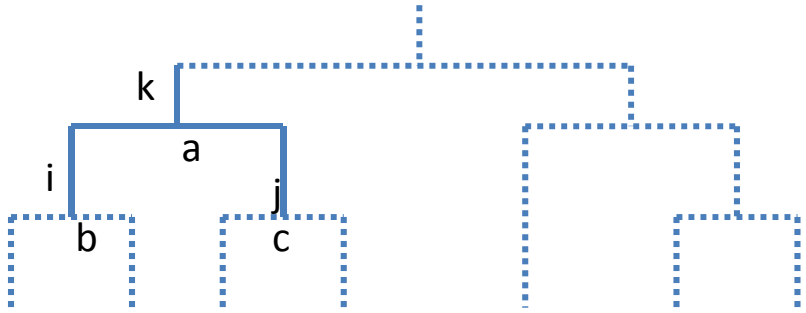
$$P(x^1, \cdots, x^n | T, t) = \prod_{u=1}^{N} P(x_u^1, \cdots, x_u^n | T, t)$$

$$P(x_u^1, \cdots, x_u^n | T, t) = \sum_{a^{n+1}, \cdots, a^{2n-1}} q_{a^{2n-1}} \prod_{i=n+1}^{2n-2} P(a^i | a^{\alpha(i)}, t_i)$$

$$\times \prod_{i=1}^{n} P(x_u^i | a^{\alpha(i)}, t_i)$$

Felsenstein's Recursive Algorithm

- Let $P(L_k|a)$ denote the probability of all the leafs below node k given that the residue at k is a.
- Then we compute $P(L_k|a)$ from the probabilities $P(L_i | b)$ and $P(L_j |c)$ for all b and c, where i and j are the daughter nodes of k.



Felsenstein's Recursive Algorithm

- Initialization: set k=2n-1
- Recursion: Compute $P(L_k \mid a)$ for all a as follows:
  - If k is leaf node: $P(L_k|a)=1$ only if $a = x_u^k$.
  - If k is not a leaf node:
    - Compute $P(L_i|a)$, $P(L_j|a)$ for all a at the daughter nodes i,j , and set $P(L_k|a) = \sum_{bc} P(b|a,t_i)P(L_i|b)P(c|a,t_j)P(L_j|c)$
- Temination: Likelihood at site u,

$$P(x_u|T,t) = \sum_a P(L_{2n-1}|a)q_a$$

Reversibility & Independence of Root Position

- The score of the optimal tree is independent of the root position if and only if:
  - the substitution matrix is **multiplicative**
  - the substitution matrix is **reversible**
- A substitititution matrix is reversible if for all a,b and t:

$$P(b|a,t)q_a = P(a|b,t)q_b$$

Maximum Likelihood (ML)

- Score each tree by
  - Assumption of independent positions “m”
- Branch lengths *t* can be optimized
  - Gradient Ascent
  - EM
- We look for the highest scoring tree
  - Exhaustive
  - Sampling methods (Metropolis)

Computational Problem

- Such procedures are computationally expensive!
- Computation of optimal parameters, per candidate, requires non-trivial optimization step.
- Spend non-negligible computation on a candidate, even if it is a low scoring one.
- In practice, such learning procedures can only consider small sets of candidate structures

参考文献

- S. Durbin, S. Eddy, A. Krogh and G. Mitchison. Biological Sequence Analysis—Probabilistic Models of Proteins and Nucleic Acids. 1998, Cambridge University Press.