

## 第 2 周阅读资料

### The Proof of Theory 2.4.

We will show that a subset  $P$  of  $\mathbb{R}^n$  is a polytope if and only if it is a bounded polyhedron. This might be intuitively clear, but a strictly mathematical proof requires some work.

We first give a definition. Let  $P$  be a convex set. A point  $z \in P$  is called a *vertex* of  $P$  if  $z$  is *not* a convex combination of two other points in  $P$ . That is, there do not exist points  $x, y$  in  $P$  and a  $\lambda$  with  $0 < \lambda < 1$  such that  $x \neq z, y \neq z$  and  $z = \lambda x + (1 - \lambda)y$ .

To characterize vertices we introduce the following notation. Let  $P = \{x \mid Ax \leq b\}$  be a polyhedron and let  $z \in P$ . Then  $A_z$  is the submatrix of  $A$  consisting of those rows  $a_i$  of  $A$  for which  $a_i z = b_i$ .

Then we can show:

**Theorem 2.2.** *Let  $P = \{x \mid Ax \leq b\}$  be a polyhedron in  $\mathbb{R}^n$  and let  $z \in P$ . Then  $z$  is a vertex of  $P$  if and only if  $\text{rank}(A_z) = n$ .*

**Proof.** *Necessity.* Let  $z$  be a vertex of  $P$  and suppose  $\text{rank}(A_z) < n$ . Then there exists a vector  $c \neq 0$  such that  $A_z c = 0$ . Since  $a_i z < b_i$  for every  $a_i$  that does not occur in  $A_z$ , there exists a  $\delta > 0$  such that:

$$(12) \quad a_i(z + \delta c) \leq b_i \text{ and } a_i(z - \delta c) \leq b_i$$

for every row  $a_i$  of  $A$  not occurring in  $A_z$ . Since  $A_z c = 0$  and  $Az \leq b$  it follows that

$$(13) \quad A(z + \delta c) \leq b \text{ and } A(z - \delta c) \leq b.$$

So  $z + \delta c$  and  $z - \delta c$  belong to  $P$ . Since  $z$  is a convex combination of these two vectors, this contradicts the fact that  $z$  is a vertex of  $P$ .

*Sufficiency.* Suppose  $\text{rank}(A_z) = n$  while  $z$  is not a vertex of  $P$ . Then there exist points  $x$  and  $y$  in  $P$  such that  $x \neq z \neq y$  and  $z = \frac{1}{2}(x + y)$ . Then for every row  $a_i$  of  $A_z$ :

$$(14) \quad \begin{aligned} a_i x \leq b_i = a_i z &\implies a_i(x - z) \leq 0, \text{ and} \\ a_i y \leq b_i = a_i z &\implies a_i(y - z) \leq 0. \end{aligned}$$

Since  $y - z = -(x - z)$ , this implies that  $a_i(x - z) = 0$ . Hence  $A_z(x - z) = 0$ . Since  $x - z \neq 0$ , this contradicts the fact that  $\text{rank}(A_z) = n$ . ■

Theorem 2.2 implies that a polyhedron has only a finite number of vertices: For each two different vertices  $z$  and  $z'$  one has  $A_z \neq A_{z'}$ , since  $A_z x = b_z$  has only one solution, namely  $x = z$  (where  $b_z$  denotes the part of  $b$  corresponding to  $A_z$ ). Since there exist at most  $2^m$  collections of subrows of  $A$ ,  $P$  has at most  $2^m$  vertices.

From Theorem 2.2 we derive:

**Theorem 2.3.** *Let  $P$  be a bounded polyhedron, with vertices  $x_1, \dots, x_t$ . Then  $P = \text{conv.hull}\{x_1, \dots, x_t\}$ .*

**Proof.** Clearly

$$(15) \quad \text{conv.hull}\{x_1, \dots, x_t\} \subseteq P,$$

since  $x_1, \dots, x_t$  belong to  $P$  and since  $P$  is convex.

The reverse inclusion amounts to:

$$(16) \quad \text{if } z \in P \text{ then } z \in \text{conv.hull}\{x_1, \dots, x_t\}.$$

We show (16) by induction on  $n - \text{rank}(A_z)$ .

If  $n - \text{rank}(A_z) = 0$ , then  $\text{rank}(A_z) = n$ , and hence, by Theorem 2.2,  $z$  itself is a vertex of  $P$ . So  $z \in \text{conv.hull}\{x_1, \dots, x_t\}$ .

If  $n - \text{rank}(A_z) > 0$ , then there exists a vector  $c \neq 0$  such that  $A_z c = 0$ . Define

$$(17) \quad \begin{aligned} \mu_0 &:= \max\{\mu \mid z + \mu c \in P\}, \\ \nu_0 &:= \max\{\nu \mid z - \nu c \in P\}. \end{aligned}$$

These numbers exist since  $P$  is compact. Let  $x := z + \mu_0 c$  and  $y := z - \nu_0 c$ .

Now

$$(18) \quad \mu_0 = \min\left\{\frac{b_i - a_i z}{a_i c} \mid a_i \text{ is a row of } A; a_i c > 0\right\}.$$

This follows from the fact that  $\mu_0$  is the largest  $\mu$  such that  $a_i(z + \mu c) \leq b_i$  for each  $i = 1, \dots, m$ . That is, it is the largest  $\mu$  such that

$$(19) \quad \mu \leq \frac{b_i - a_i z}{a_i c}$$

for every  $i$  with  $a_i c > 0$ .

Let the minimum (18) be attained by  $i_0$ . So for  $i_0$  we have equality in (18). Therefore

$$(20) \quad \begin{aligned} \text{(i)} \quad A_z x &= A_z z + \mu_0 A_z c = A_z z, \\ \text{(ii)} \quad a_{i_0} x &= a_{i_0}(z + \mu_0 c) = b_{i_0}. \end{aligned}$$

So  $A_x$  contains all rows in  $A_z$ , and moreover it contains row  $a_{i_0}$ . Now  $A_z c = 0$  while  $a_{i_0} c \neq 0$ . This implies  $\text{rank}(A_x) > \text{rank}(A_z)$ . So by our induction hypothesis,  $x$  belongs to  $\text{conv.hull}\{x_1, \dots, x_t\}$ . Similarly,  $y$  belongs to  $\text{conv.hull}\{x_1, \dots, x_t\}$ . Therefore, as  $z$  is a convex combination of  $x$  and  $y$ ,  $z$  belongs to  $\text{conv.hull}\{x_1, \dots, x_t\}$ . ■

As a direct consequence we have:

**Corollary 2.3a.** *Each bounded polyhedron is a polytope.*

**Proof.** Directly from Theorem 2.3. ■

Conversely:

**Theorem 2.4.** *Each polytope is a bounded polyhedron.*

**Proof.** Let  $P$  be a polytope in  $\mathbb{R}^n$ , say

$$(21) \quad P = \text{conv.hull}\{x_1, \dots, x_t\}.$$

We may assume that  $t \geq 1$ . We prove the theorem by induction on  $n$ . Clearly,  $P$  is bounded.

If  $P$  is contained in some affine hyperplane, the theorem follows from the induction hypothesis.

So we may assume that  $P$  is not contained in any affine hyperplane. It implies that the vectors  $x_2 - x_1, \dots, x_t - x_1$  span  $\mathbb{R}^n$ . It follows that there exist a vector  $x_0$  in  $P$  and a real  $r > 0$  such that the ball  $B(x_0, r)$  is contained in  $P$ .

Without loss of generality,  $x_0 = 0$ . Define  $P^*$  by

$$(22) \quad P^* := \{y \in \mathbb{R}^n \mid x^T y \leq 1 \text{ for each } x \in P\}.$$

Then  $P^*$  is a polyhedron, as

$$(23) \quad P^* = \{y \in \mathbb{R}^n \mid x_j^T y \leq 1 \text{ for } j = 1, \dots, t\}.$$

This follows from the fact that if  $y$  belongs to the right hand set in (23) and  $x \in P$  then  $x = \lambda_1 x_1 + \dots + \lambda_t x_t$  for certain  $\lambda_1, \dots, \lambda_t \geq 0$  with  $\lambda_1 + \dots + \lambda_t = 1$ , implying

$$(24) \quad x^T y = \sum_{j=1}^t \lambda_j x_j^T y \leq \sum_{j=1}^t \lambda_j = 1.$$

So  $y$  belongs to  $P^*$ .

Moreover,  $P^*$  is bounded, since for each  $y \neq 0$  in  $P^*$  one has that  $x := r \cdot \|y\|^{-1} \cdot y$  belongs to  $B(0, r)$  and hence to  $P$ . Therefore,  $x^T y \leq 1$ , and hence

$$(25) \quad \|y\| = (x^T y)/r \leq 1/r.$$

So  $P^* \subseteq B(0, 1/r)$ .

This proves that  $P^*$  is a bounded polyhedron. By Corollary 2.3a,  $P^*$  is a polytope. So there exist vectors  $y_1, \dots, y_s$  in  $\mathbb{R}^n$  such that

$$(26) \quad P^* = \text{conv.hull}\{y_1, \dots, y_s\}.$$

We show:

$$(27) \quad P = \{x \in \mathbb{R}^n \mid y_j^T x \leq 1 \text{ for all } j = 1, \dots, s\}.$$

This implies that  $P$  is a polyhedron.

To see the inclusion  $\subseteq$  in (27), it suffices to show that each of the vectors  $x_i$  belongs to the right hand side in (27). This follows directly from the fact that for each  $j = 1, \dots, s$ ,  $y_j^T x_i = x_i^T y_j \leq 1$ , since  $y_j$  belongs to  $P^*$ .

To see the inclusion  $\supseteq$  in (25), let  $x \in \mathbb{R}^n$  be such that  $y_j^T x \leq 1$  for all  $j = 1, \dots, s$ . Suppose  $x \notin P$ . Then there exists a hyperplane separating  $x$  and  $P$ . That is, there exist a vector  $c \neq 0$  in  $\mathbb{R}^n$  and a  $\delta \in \mathbb{R}$  such that  $c^T x' < \delta$  for each  $x' \in P$ , while  $c^T x > \delta$ . As  $0 \in P$ ,  $\delta > 0$ . So we may assume  $\delta = 1$ . Hence  $c \in P^*$ . So there exist  $\mu_1, \dots, \mu_s \geq 0$  such that  $c = \mu_1 y_1 + \dots + \mu_s y_s$  and  $\mu_1 + \dots + \mu_s = 1$ . This gives the contradiction:

$$(28) \quad 1 < c^T x = \sum_{j=1}^s \mu_j y_j^T x \leq \sum_{j=1}^s \mu_j = 1. \quad \blacksquare$$