

Algorithm Correctness

(or how to prove programs are correct)

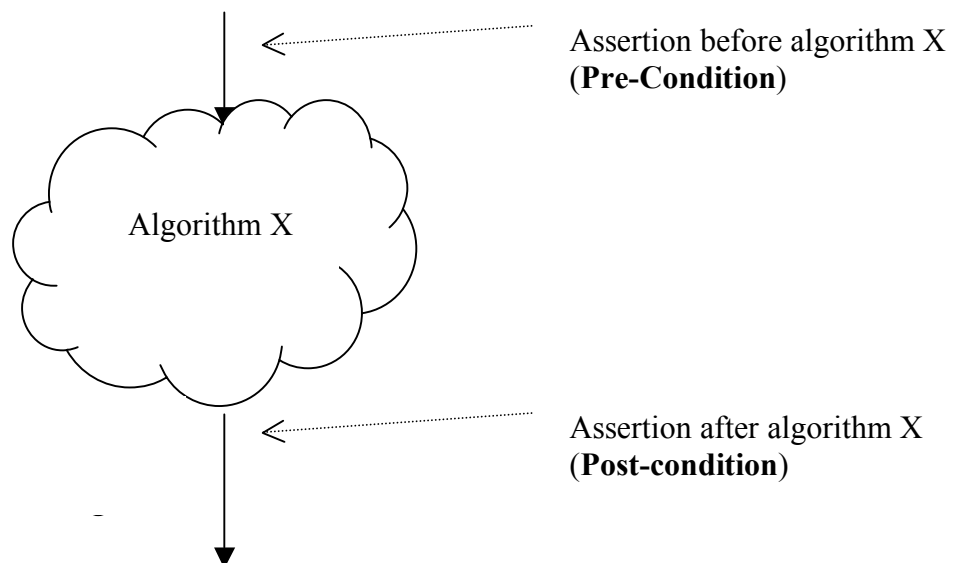


Pre-condition of the algorithm: a predicate that describes the initial state (before execution)

Post-condition of the algorithm: a predicate that describes the final state (after execution)

The algorithm is correct if it can be proved that if the pre-condition is true, the post-condition must be true.

Pre-condition \Rightarrow Post-condition



ex 1:

```
program mystery
begin
.
.
end
```

Pre-condition: $a = 3 \quad b = 5$

Post-condition: $a = 5 \quad b = 3$

ex 2:

```
x = 2
z = x + y
if ( y > 0 )
    z = z + 1
else
    z = 0
```

Pre-condition: $y = 6$

Post-condition: $z = 9$

Programs may be sliced into sequential parts and assertions inserted

```
program mystery2
begin
```

Segment 1

Segment 2

Segment 3

⋮

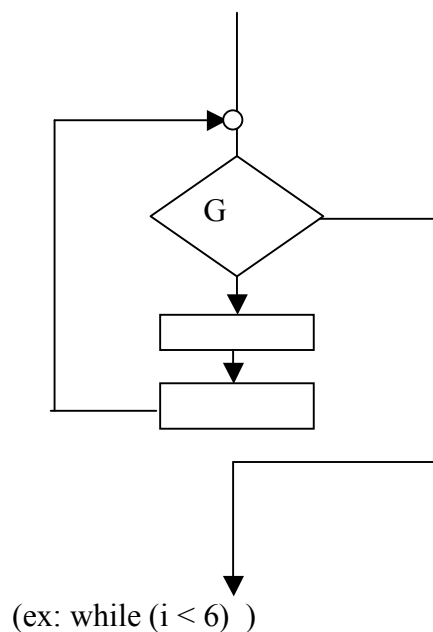
Segment n

```
end
```

The loop invariant $I(n)$ is a predicate where the variable, n , represents the number of times the loop has iterated.

What about loops?

Consider a while loop with Pre-Condition and Post-Condition and guard G .



To prove that the loop is “correct” do the following:

1. Write a **loop invariant** that “captures” the purpose of the loop; call it $I(n)$. (*Difficult!!*)
 $I(n)$ is a predicate that is true after n iterations of the loop.
2. Show that the pre-condition implies the loop invariant $I(0)$. (*Basis Step*)
3. Show that if $I(k)$ and guard G are true, then $I(k+1)$ is true. (*Inductive Step*)
In other words,
if $I(n)$ is true for $n = k$ and the guard G is still true, then it follows that $I(n)$ is true for $n = k+1$.
4. Show that guard G eventually becomes false (looping stops).
5. If the loop iterates a maximum of N times, show that $I(N)$ implies the post-condition.

Example 1

Pre-condition: m is a non-negative integer, x is real, $\text{Count} = 0$, $\text{Product} = 0$

while ($\text{Count} \neq m$)


{ $\text{Product} = \text{Product} + x$;

$\text{Count} = \text{Count} + 1$;

}

What's guard G?

Post-condition: $\text{Count} = m$ and $\text{Product} = m \cdot x$

-
1. **Write loop invariant $I(n)$:** (Remember that n represents the number of times the loop has iterated) \Rightarrow 

The loop invariant can be a conjunction (AND)

2. Show that the pre-condition implies that $I(0)$ is true.

3. **Inductive Step:** Show that if $I(k)$ and G are true, $I(k+1)$ is true.

Assume G : and $I(k)$: $\text{Product} = k \cdot x$ and $\text{Count} = k$ are true.

(these are $\text{Product}_{\text{old}}$ and $\text{Count}_{\text{old}}$)

Show $I(k+1)$: $\text{Product} = (k+1) \cdot x$ and $\text{Count} = k + 1$ is true.

(these are $\text{Product}_{\text{new}}$ and $\text{Count}_{\text{new}}$)

4. Show that G will eventually be false.

5. Show that the post-condition follows when the loop stops.

Example 2.

Pre-condition: $\text{Max} = \text{A}[0]; \text{Count} = 0$

```
while ( Count < 10)
{ Count = Count + 1;
  if ( A[Count] > Max )
    Max = A[Count];
}
```

Post-condition: Max contains the largest value in $\text{A}[0].. \text{A}[10]; \text{Count} = 10$

- 1) Write loop invariant $I(n)$:
- 2) Show that the pre-condition implies that $I(0)$ is true.
- 3) Show $I(k)$ and G imply $I(k+1)$.
 $I(k)$:
 $I(k+1)$:
- 4) Show that G will eventually be false.
- 5) Show the post-condition follows when the loop stops