

NATIONAL UNIVERSITY OF SINGAPORE  
SCHOOL OF COMPUTING  
SEMESTER I: 2012–2013  
EXAMINATION FOR  
CS3230 – DESIGN AND ANALYSIS OF ALGORITHMS  
November 2012 – Time Allowed 2 Hours

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**INSTRUCTIONS TO CANDIDATES**

1. This examination paper consists of FOUR (4) questions and comprises THREE (3) printed pages (including this page).
2. Answer **ALL** questions.
3. This is an **Open Book** examination.

## Question 1. (10 marks)

We are given two sets of natural numbers (all greater than or equal to 1)  $A = \{a_1, \dots, a_n\}$  and  $B = \{b_1, \dots, b_n\}$ . All elements in  $A$  are distinct. We say that  $B$  covers  $A$  if for every element  $a_i \in A$  there is an element  $b_j \in B$  such that  $\sqrt{b_j} \leq a_i \leq b_j^2$ . We are supposed to determine if  $B$  covers  $A$  in time  $O(n \log n)$ .

- (a) (4 marks) Write idea of an algorithm for this task.
- (b) (2 marks) Write the pseudocode.
- (c) (2 marks) Argue the correctness of your algorithm.
- (d) (2 marks) Argue that the running time is as desired.

## Question 2. (10 marks)

We are given an  $n$  node labeled tree  $T$  of depth  $d$  and degree at most 4. Each node  $v$  of  $T$  is labeled by a real number  $l(v)$ . A node  $v$  of  $T$  is a local-maximum if  $l(v)$  is greater than or equal to  $l(w)$  for all nodes  $w$  that are joined to  $v$  by an edge. The labels of the nodes are hidden. We have access to a function  $f$  such that for all pairs of nodes  $(v_1, v_2)$ ,  $f(v_1, v_2) = 1$  if  $l(v_1) \geq l(v_2)$  and 0 otherwise.

We are supposed to determine a local maximum using  $O(d)$  invocations of  $f$ .

- (a) (5 marks) Write idea of an algorithm for this task. (Hint: You may start by comparing the label of the root with its children.)
- (b) (2 marks) Write the pseudocode.
- (c) (3 marks) Argue that the number of invocations of  $f$  is  $O(d)$ .

## Question 3. (10 marks)

We are given a set  $S$  of  $n$  pairs of real numbers  $S = \{(a_1, b_1), \dots, (a_n, b_n)\}$  such that  $a_1 \leq a_2 \leq \dots \leq a_n$ . We say that a pair of indices  $(i, j)$  is compatible if  $|a_j - a_i| \geq |b_j - b_i|$ . We are required to select a set  $T \subseteq \{1, 2, \dots, n\}$ , containing  $n$ , of as large size as possible such that for any  $i, j \in T$  the pair  $(i, j)$  is compatible.

- (a) (5 marks) Give idea of a dynamic-programming algorithm running in time  $O(n^2)$  to find the size of optimal  $T$ . Justify the correctness of the recursion relation you obtain.
- (b) (3 marks) Write the pseudocode.
- (c) (2 marks) Argue that the running time is  $O(n^2)$ .

## Question 4. (10 marks)

For NP-complete problems, sometimes we wonder if special cases of them are easy to solve. We know that Vertex-Cover is an NP-complete problem. We may wonder if the problem becomes easy to solve in graphs of small degree. It turns out the vertex cover is a hard problem even for graphs of small degree. Consider the problem Degree-utmost-4-vertex-cover defined as follows.

Given a graph  $G = (V, E)$  with degree at most 4 and a number  $k$ , does this graph have a vertex cover of size at most  $k$ ?

Show that Degree-utmost-4-vertex-cover is NP-complete.

END of QUESTIONS