# CS3230 : Design and Analysis of Algorithms (Fall 2014) Tutorial Set #1

[For discussion during Week 3]

S-Problem Due: Friday, 22-Aug, before noon.

**OUT:** 19-Aug-2014 **Tutorials:** Tue & Wed, 26-27 Aug 2014

S-Problem Due: Friday, 22-Aug, before Noon.

## **IMPORTANT:** Read "Remarks about Homework" – also applies to tutorials.

## Prepare your answers to all the D-Problems in every tutorial set.

When preparing to present your answers,

- Think of a CLEAR EXPLANATION
- Illustrate with a good worked example;
- Describe the main ideas,
- Can you sketch why the solution works;
- Give analysis of running time, if appropriate
- Can you think of other (perhaps simpler) solutions?

In CS3230, you learn to develop high-level abstractions when describing algorithms. Try not to speak in ML/AL (machine/assembly language) or "for (j=0; j < n; j++) do". Instead give names to your sets (of objects or things or data structures), talk about Depth-First Search, Binary Search, traverse the graph, sort the set, use a priority queue, etc. You are no longer in CS1010, CS1020, CS2010 or CS2020. Speak with greater sophistication, and at a higher level of abstraction.

#### Remember:

- You can **freely quote** standard algorithms and data structures covered in the lectures (and including from pre-requisites) modules, textbook. Explain **any modifications** you make to them, and how they may affect the running time.
- There is no need to copy/re-prove algorithms or theorems already cover already;
- Unless otherwise specified, you are expected to **prove (justify)** your results. All logarithms are base 2, unless otherwise stated.

### Examples:

- a. Use Quicksort to sort the array X[1..n] in increasing order;
- b. Organize the set S as a Max-Heap (array-based);
- c. Run a post-order traversal of the tree T, and at each node, the processing of the node is ...
- d. Run Dijkstra's algorithm for single-source shortest path from vertex w on graph G=(V, E).
- e. Do <some-std-alg Q>, but with the following modifications: blah, blah, blah....
- f. By the Handshaking Lemma,  $(d_1 + d_2 + d_2 + \ldots + d_n) = 2e$  (OK, if you still don't know the Handshaking Lemma, Google it and learn it.)

Of course, it is your responsibility to ensure that the algorithm that you quote ACTUALLY solves your problem.

Routine Practice Problems -- do not turn these in -- but make sure you know how to do them.

- **R1.** (a) Show, by definition, that f(n) = O(n), where f(n) = 819n.
  - (a') Show that  $f(n) = O(n^2)$ ,  $f(n) = O(n^{819})$ .
  - **(b)** Show, by definition, that  $g(n) = O(n^2)$ , where  $g(n) = 17n^2$ .
  - **(b')** Show that  $g(n) = O(n^3)$ ,  $g(n) = O(n^{17})$ .
  - (c) Show, by definition, that  $h(n) = O(n^2)$ , where  $h(n) = 17n^2 + 819n$ .
  - (c') Show that  $h(n) = O(n^3)$ ,  $h(n) = O(n^{836})$ .
  - (d) Show, by definition, that  $k(n) = O(n^2)$ , where  $k(n) = 17n^2 + 37n(\lg n)$
  - (d') Show that  $g(n) = O(n^3)$ ,  $g(n) = O(n^{54})$ .
- **R2.** (a) Show, by definition, that  $f(n) = \Theta(n)$ , where f(n) = 819n.

(a') Is 
$$f(n) = \Theta(n^2)$$
? Is  $f(n) = \Theta(n^{819})$ ?

- **(b)** Show, by definition, that  $g(n) = \Theta(n^2)$ , where  $g(n) = 17n^2$ .
- **(b')** Is  $g(n) = \Theta(n)$ ? Is  $g(n) = \Theta(n^3)$ ? Is  $g(n) = \Theta(n^{17})$ ?
- (c) Show, by definition, that  $h(n) = \Theta(n^2)$ , where  $h(n) = 17n^2 + 819n$ .
- (c') Is  $h(n) = \Theta(n)$ ? Is  $h(n) = \Theta(n)$
- Is  $h(n) = \Theta(n^3)$ ? Is  $g(n) = \Theta(n^{836})$ ?
- (d) Show, by definition, that  $k(n) = \Theta(n^2)$ , where  $k(n) = 17n^2 + 37n(\lg n)$
- (d') Is  $k(n) = \Theta(n \lg n)$ ? Is  $k(n) = \Theta(n^3)$ ? Is  $k(n) = \Theta(n^{54})$ ?
- **R3.** From **T1-R1** and **T1-R2**, can you see the big different between O-notation and  $\Theta$ -notation?

**S-Problems:** (To do and submit by Friday, 22-Aug, before noon.) Solve this S-problem and submit for grading.

# S1: (Two Important Processes in CS)

## [Repeated-halving]

Start with a number n. Repeatedly "divide by two (throw away the remainder)" until we reach 0. How many steps will you take? Let h(n) be the number of steps.

# [Repeated-doubling]

Start with the number 1. Repeatedly *multiply by two* until we get a number greater than or equal to n. How *many steps* will you take? Let d(n) be the number of steps.

- (a) [In a nice table, list the value of h(n) d(n) for n = 1-25, 31, 32, 33, 63, 64, 65, 100, 127, 128, 129, 100, 1023, 1024, 1025,  $10^6$ ,  $10^9$ .]
- **(b)** Write the repeated-halving and repeated-doubling processes in pseudo-code.
- (c) Write down any relationship you see between the two processes? between h(n) and d(n)?
- (d) Give an exact mathematical formulae for h(n) and d(n). [Hint: As a self-check, you can test the formula out yourself.]

#### **D-Problems:**

Solve these D-problems and prepare to discuss them in tutorial class. You will call upon one of you to present your solution or your best attempt at a solution. Your solution presentation does NOT need to be fully correct, given your best attempt. The TA will help clarify and correct any issues or errors.

- **D1. Discuss solution to HW0-S1.** (Hint: Compare with problem T1-S1.)
- D2. Discuss solution to HW0-S3.
- D3. [Proving  $O, \Omega, \Theta$  by definition]

Let 
$$f(n) = 14n^2 - 6n + 707$$
  $g(n) = 19n^2 - 707n(\lg n) + 8n - 30.$ 

Prove the following by using the definitions of O,  $\Omega$ ,  $\Theta$ . Namely, find the respective constants c,  $c_1$ ,  $c_2$ , and the positive integer  $n_0$ .

(i) 
$$f(n) = O(n^2)$$
 (ii)  $f(n) = \Omega(n^2)$ 

(ii) 
$$f(n) = \Omega(n^2)$$

(iii) 
$$f(n) = \Theta(n^2)$$

(iv) 
$$g(n) = O(n^2)$$
 (v)  $g(n) = \Omega(n^2)$  (vi)  $g(n) = \Theta(n^2)$ 

(v) 
$$g(n) = \Omega(n^2)$$

(vi) 
$$g(n) = \Theta(n^2)$$

- Prove the results in T1-D3(a) above by making use of some of the following short-cut methods: sum-rule or product-rule, by polynomial rule, L'Hopital's rule (limits), etc. (Find the method that is easiest for you.)
- **D4.** [Analysis of Algorithm Ace.] Analyze the following algorithm (called Ace) and give a good estimate of its running time (in  $\Theta$  notation).

## Procedure Ace(N):

1. Set A[k]=1 for all k=2,3,...,N

// Initialize A to all "primes"

- 2. for  $(p := 2; p \le N; p++)$  do begin
- 3. if (A[p] = 1) then Print p;
- 4. for  $(j := 2p; j \le N; j := j + p)$  do A[j] := 0;

// mark j as "not prime"

**Note:** The running time is NOT  $\Theta(n^2)$ .

[Hint: The loop in Step 4, ask how many times does it loop?]