

# CS3230

## Tutorial 3

1. Draw a decision tree to show that 4 numbers can be sorted using at most 5 comparisons.
2. Given an array  $A$  of integers and a number  $k$ , give an efficient algorithm to decide if there exist two elements  $A[i]$  and  $A[j]$ ,  $i \neq j$ , such that  $A[i] + A[j] = k$ . Give time complexity bound of your algorithm.
3. Consider the following modification of the partition algorithm done in class. Show that it works correctly.

Partition( $A, i, j$ )

Assumption:  $i < j$

1. Let  $m = i + 1$ ,  $n = j$ ;
  2. While  $m \leq n$ , do {
    3. While  $A[m] < A[i]$  and  $m \leq n$  do {  $m = m + 1$  }
    4. While  $A[n] \geq A[i]$  and  $n \geq m$  do {  $n = n - 1$  }
    5. If  $n > m$ , then swap( $A[m], A[n]$ ).}
  6. swap( $A[i], A[n]$ )
  7. Return  $n$ .
- End

4. Given as input a sorted array  $A$ , containing  $n$  elements, and two numbers  $\ell$  and  $u$  (where  $\ell \leq u$ ). Give an algorithm to find how many numbers are there in the array which are between  $\ell$  and  $u$  (both inclusive). That is, find the number of  $x$  in the array  $A$  such that  $\ell \leq x \leq u$ .

What is the time complexity of your algorithm.

5. Suppose we are given an array of  $n$  integers between 1 to  $m$  (both inclusive). Preprocess the array such that one can answer the following query in constant time:

How many numbers are there in the array which are between  $\ell$  and  $u$  (both inclusive), where  $1 \leq \ell \leq u \leq m$ .

What is the time complexity of your preprocessing algorithm? Try to make it linear in  $m$  and  $n$ .