

### Recursion

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### In this class

• How recursion works in program?

#### Prob 1. Factorial Function

Two statements for factorial function

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \cdot (n-1) \cdot (n-2) \cdots 3 \cdot 2 \cdot 1 & \text{if } n \ge 1. \end{cases}$$

$$n = 0$$

$$n \ge 1.$$

$$n! = \begin{cases} 1 \\ n \cdot (n-1)! \end{cases}$$

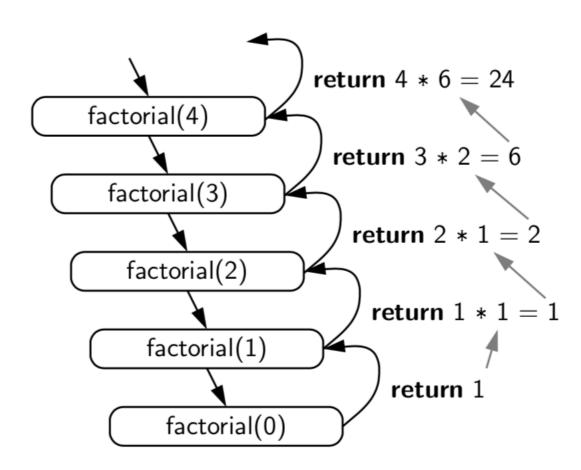
```
if n = 0 if n \ge 1.
```

Non-Recursion Program

```
1  def factorial(n):
2   if n == 0:
3    return 1
4   else:
5   return n * factorial(n-1)
```

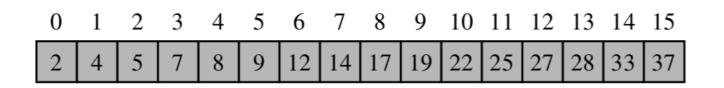
### Recursion

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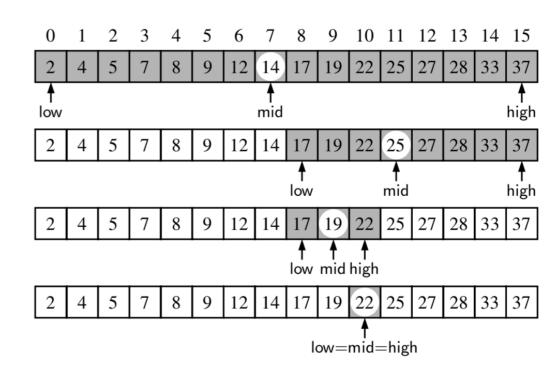


## Prob 2. Binary Search

- To locate a target value within a stored sequence
- Ex) Search 22



$$\mathsf{mid} = \lfloor (\mathsf{low} + \mathsf{high})/2 \rfloor$$
.



## Binary Search

```
def binary_search(data, target, low, high):
      """Return True if target is found in indicated portion of a Python list.
 3
      The search only considers the portion from data[low] to data[high] inclusive.
      if low > high:
        return False
                                                        # interval is empty; no match
      else:
        mid = (low + high) // 2
9
        if target == data[mid]:
10
                                                       # found a match
          return True
11
                                                                                                 9 10 11 12 13 14 15
        elif target < data[mid]:</pre>
12
                                                                                         12 [14] 17 | 19 | 22
                                                                                                                 33 | 37
           # recur on the portion left of the middle
13
          return binary_search(data, target, low, mid -1)
14
                                                                      low
                                                                                            mid
15
        else:
                                                                                   8
                                                                                      9 | 12 | 14 | 17 | 19 | 22
                                                                                                                 33 37
           # recur on the portion right of the middle
16
                                                                                                        mid
```

**return** binary\_search(data, target, mid + 1, high)

17

high

high

9 | 12 | 14 | 17 | 19 | 22 | 25 | 27 | 28 | 33 | 37

9 | 12 | 14 | 17 | 19 | 22 | 25 | 27 | 28 | 33 | 37

low=mid=high

low mid high

## Complexity Analysis of B-Search in Recursion

**Proposition 4.2:** The binary search algorithm runs in  $O(\log n)$  time for a sorted sequence with n elements.

**Justification:** To prove this claim, a crucial fact is that with each recursive call the number of candidate entries still to be searched is given by the value

$$high - low + 1$$
.

Moreover, the number of remaining candidates is reduced by at least one half with each recursive call. Specifically, from the definition of mid, the number of remaining candidates is either

$$(\mathsf{mid}-1)-\mathsf{low}+1 = \left\lfloor \frac{\mathsf{low}+\mathsf{high}}{2} \right\rfloor - \mathsf{low} \leq \frac{\mathsf{high}-\mathsf{low}+1}{2}$$

or

$$\mathsf{high} - (\mathsf{mid} + 1) + 1 = \mathsf{high} - \left | \frac{\mathsf{low} + \mathsf{high}}{2} \right | \leq \frac{\mathsf{high} - \mathsf{low} + 1}{2}.$$

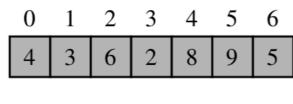
Initially, the number of candidates is n; after the first call in a binary search, it is at most n/2; after the second call, it is at most n/4; and so on. In general, after the  $j^{th}$  call in a binary search, the number of candidate entries remaining is at most  $n/2^j$ . In the worst case (an unsuccessful search), the recursive calls stop when there are no more candidate entries. Hence, the maximum number of recursive calls performed, is the smallest integer r such that

$$\frac{n}{2^r} < 1.$$

In other words (recalling that we omit a logarithm's base when it is 2),  $r > \log n$ . Thus, we have  $r = \lfloor \log n \rfloor + 1$ ,

which implies that binary search runs in  $O(\log n)$  time.

# Prob 3. Reversing Sequence





```
5 9 8 2 6 3 4
```

## Prob 5. Power Comp.

```
def power(x, n):
                                                                                                                                                    x^{n} = x \cdot x^{n-1} for n > 0.
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                 """ Compute the value x**n for integer n."""
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                  if n == 0:
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     4 return 1
power(x,n) = \begin{cases} 1 & \text{if } n = 0 \\ x \cdot power(x,n-1) & \text{otherwise.} \end{cases} for the second second
```

6 **return**  $\times * power(x, n-1)$ 

## Program Assignments

Build a non-recursive program for factorial function

### In next class

Array