

# **STRATEGO:**

# **A Model-Based Decision**

## **Procedure:**

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# Programming Language History

## 1940년 이전 [편집]

초기의 프로그래밍 언어는 현대의 컴퓨터에서 거슬러 올라간다. 초기 1801년에 발명된 [지카드식 문작기](#)는 자동으로 장식 패턴을 발생시켰다. 1842년~1843년의 9개월의 기간에 걸쳐 [에이다 러브레이스](#)는 이를 자세하게 일련의 주석으로 추가했는데, 이는 일부 역사가들이 세계 최초의 컴퓨터 부호들은 이들을 응용하여 만들어졌다.

## 1940년대 [편집]

이 시기에 개발된 일부 중요 언어는 다음을 포함한다:

- 1943년 - [Plankalkül](#)
- 1943년 - [애니악 코딩 시스템](#)<sup>[2]</sup>
- 1949년~1954년 - 나중에 [유니박](#)으로 발전함.<sup>[3]</sup>

## 1950년대 ~ 1960년대 [편집]

1950년대에는 다음의 세 가지 현대의 프로그래밍 언어가 설계되었다:

- [포트란](#) (1954년)
- [리스프](#) (1958년)
- [코볼](#) (1959년)

그 밖의 주요 언어들은 다음과 같다:

- 1951년 - 지역 어셈블리어(Regional Assembly Language)
- 1952년 - [오토코드](#)
- 1954년 - [IPL](#) (리스프의 선구자)
- 1955년 - [FLOW-MATIC](#) (코볼의 선구자)
- 1957년 - [COMTRAN](#) (코볼의 선구자)
- 1958년 - [알골 58](#)
- 1959년 - [FACT](#) (코볼의 선구자)
- 1959년 - [RPG](#)
- 1962년 - [APL](#)
- 1962년 - [시뮬라](#)
- 1962년 - [SNOBOL](#)
- 1963년 - [CPL](#) (C의 선구자)
- 1964년 - [BASIC](#)
- 1964년 - [PL/I](#)
- 1967년 - [BCPL](#) (C의 선구자)

## 1968년대 ~ 1978년대 [편집]

1960년대 말에서 1970년대 말의 기간 동안 다음의 다섯 가지 각이 기간 동안 개발된 주요 언어들은 다음과 같다:

- [시뮬라](#)(Simula, 1960년대 말 발명)
- [C](#)
- [스몰토크](#) (1970년대 중반)
- [프롤로그](#) (1972년)
- [ML](#)

그 밖의 주요 언어로는 다음과 같다.

- 1968년 - [로고](#)
- 1969년 - [B](#) (C의 선구자)
- 1970년 - [파스칼](#)
- 1970년 - [포스](#) (Forth)
- 1975년 - [스ქ](#) (Scheme)
- 1978년 - [SQL](#) (처음에는 쿼리 언어일 뿐이었으나, 훗날 프로세서로 기간 동안 개발된 주요 언어들은 다음과 같다:

## 1980년대 [편집]

1980년대 말에서 1990년대 말의 기간 동안 개발된 주요 언어들은 다음과 같다:

- 1980년 - [C++](#) (처음 이름은 C with classes였으나 1983년 7월 이름이 C++로 변경)
- 1983년 - [에이다](#)
- 1984년 - [커먼 리스프](#)
- 1984년 - [MATLAB](#)
- 1985년 - [에펠\(Eiffel\)](#)
- 1986년 - [오브젝티브-C](#)
- 1986년 - [얼랭](#)
- 1987년 - [펄](#)
- 1988년 - [Tcl](#)
- 1988년 - [매스매틱\(Mathematica\)](#)
- 1989년 - [FL](#) (Backus);

## 1990년대 [편집]

1990년대 말에서 2000년대 말의 기간 동안 개발된 주요 언어들은 다음과 같다:

- 1990년 - [하스켈](#)
- 1991년 - [파이썬](#)
- 1991년 - [비주얼 베이직](#)
- 1991년 - [HTML](#) (마크업 언어)
- 1993년 - [루비](#)
- 1993년 - [루아](#)
- 1994년 - [CLOS](#) (ANSI 커먼 리스프의 일부)
- 1995년 - [자바](#)
- 1995년 - [델파이 \(오브젝트 파스칼\)](#)
- 1995년 - [자바스크립트](#)
- 1995년 - [PHP](#)
- 1996년 - [WebDNA](#)
- 1997년 - [Rebol](#)
- 1999년 - [D](#)

## 현재의 경향 [편집]

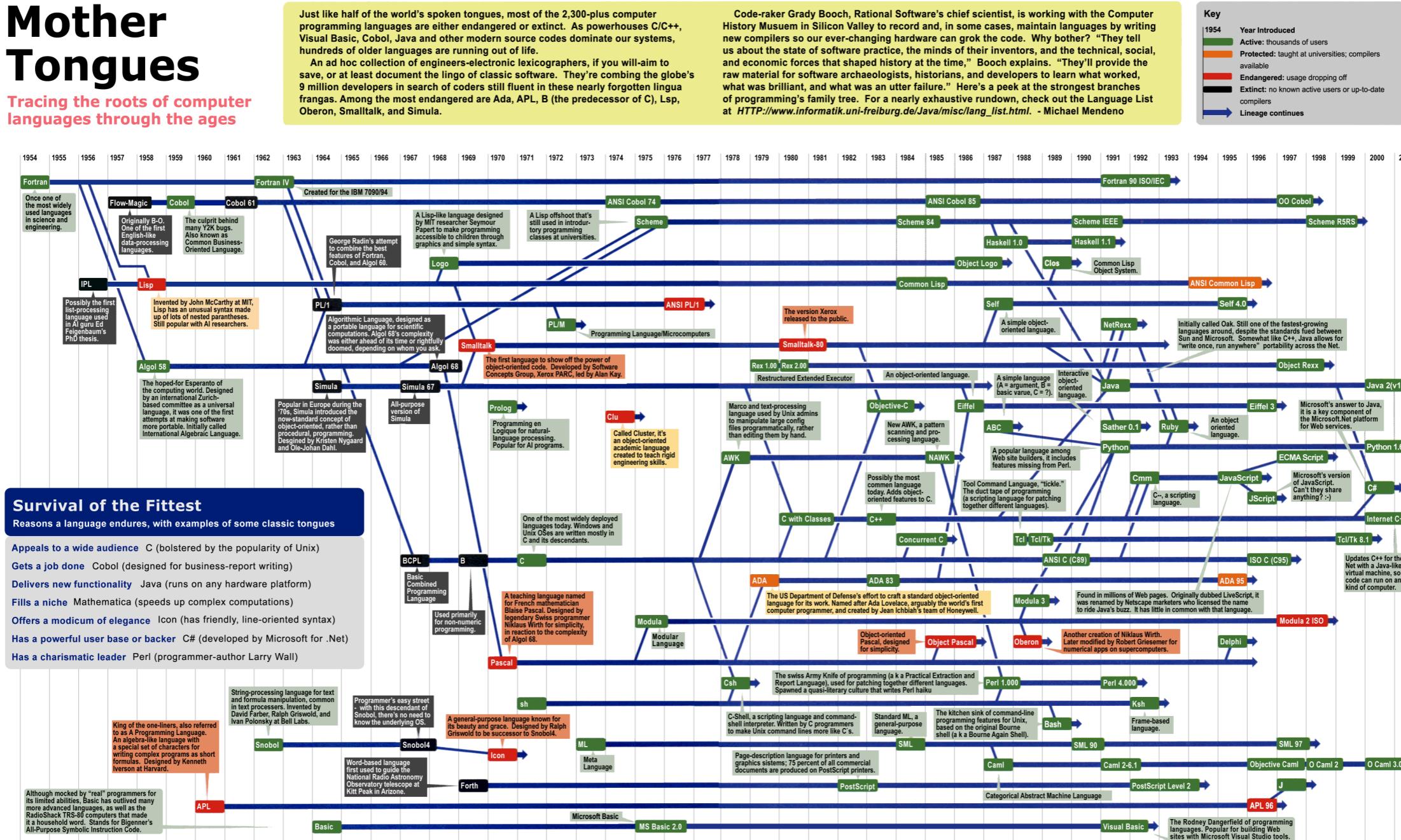
2000년 이후로 개발된 주요 언어들은 다음과 같다:

- 2000년 - [액션스크립트](#)
- 2001년 - [C#](#)
- 2001년 - [비주얼 베이직 닷넷](#)
- 2002년 - [F#](#)
- 2003년 - [그루비](#)
- 2003년 - [스칼라](#)
- 2003년 - [팩터](#)
- 2007년 - [Clojure](#)
- 2009년 - [고](#)
- 2011년 - [다트](#)
- 2014년 - [스위프트](#)
- 2015년 - [러스트](#)

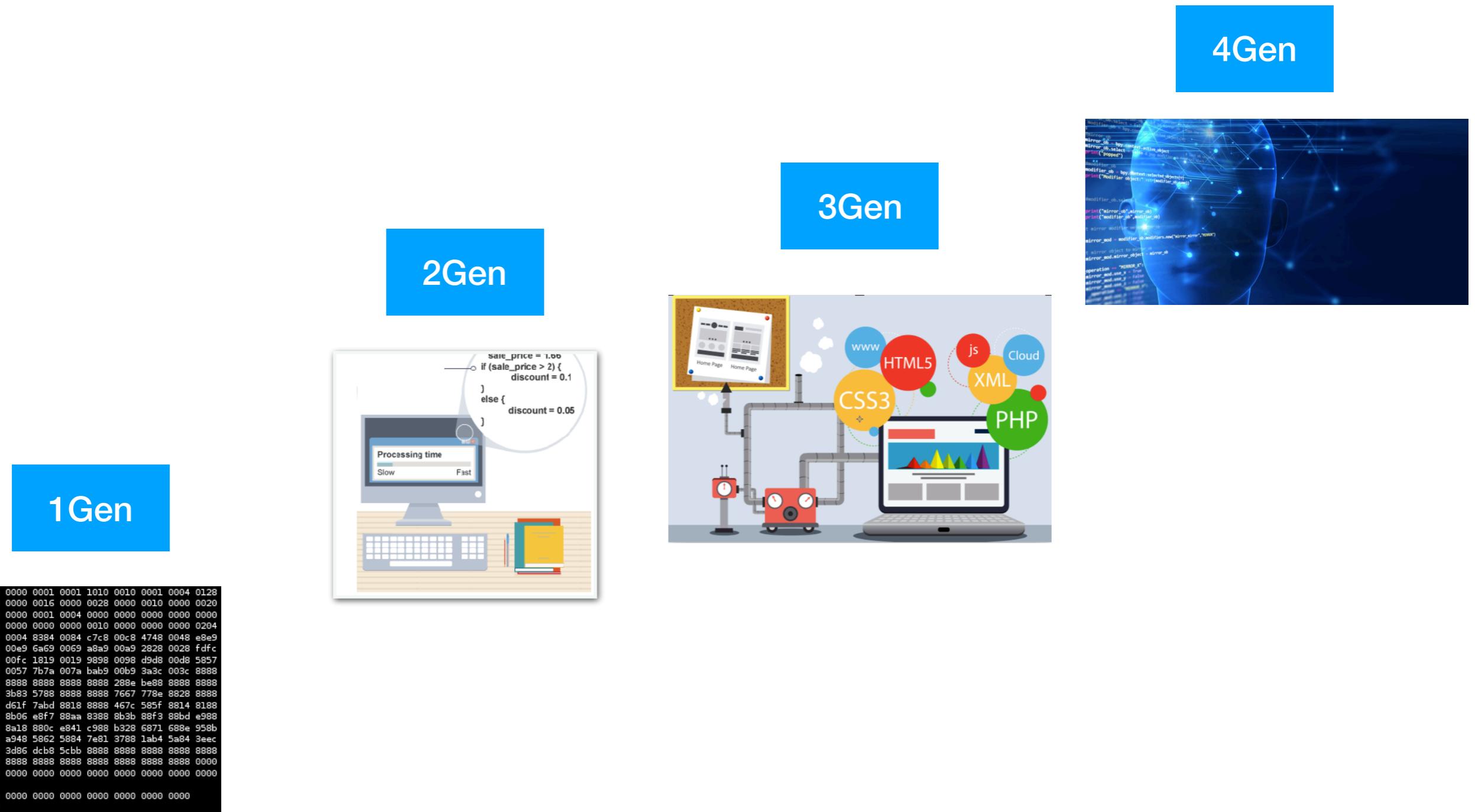
# Programming Language History

## Mother Tongues

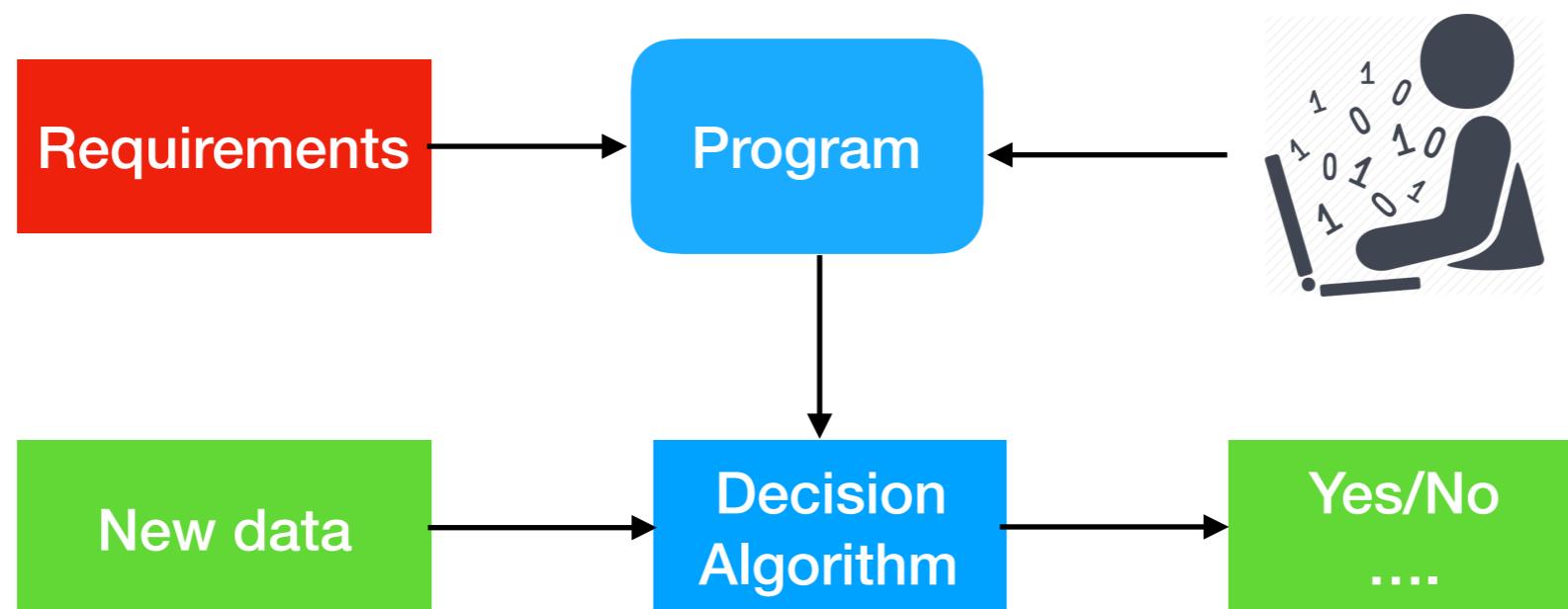
Tracing the roots of computer languages through the ages



# Programming Paradigm Evolution

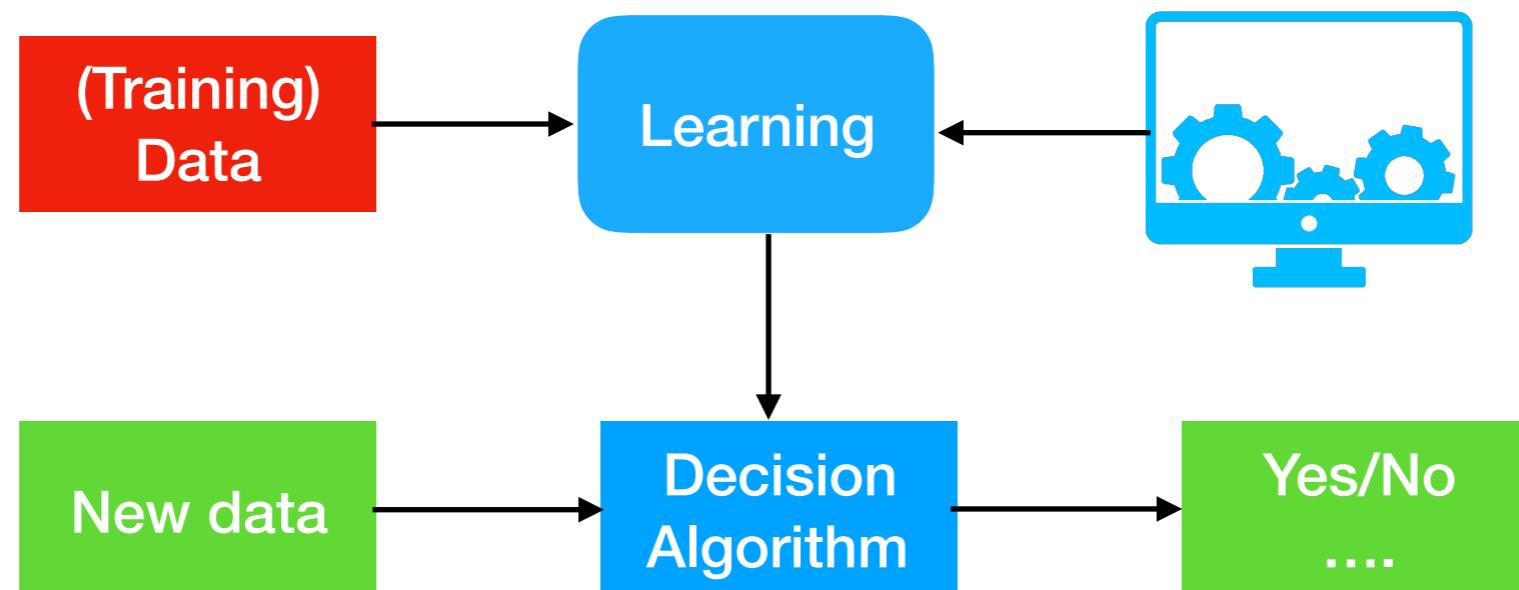


# Programming



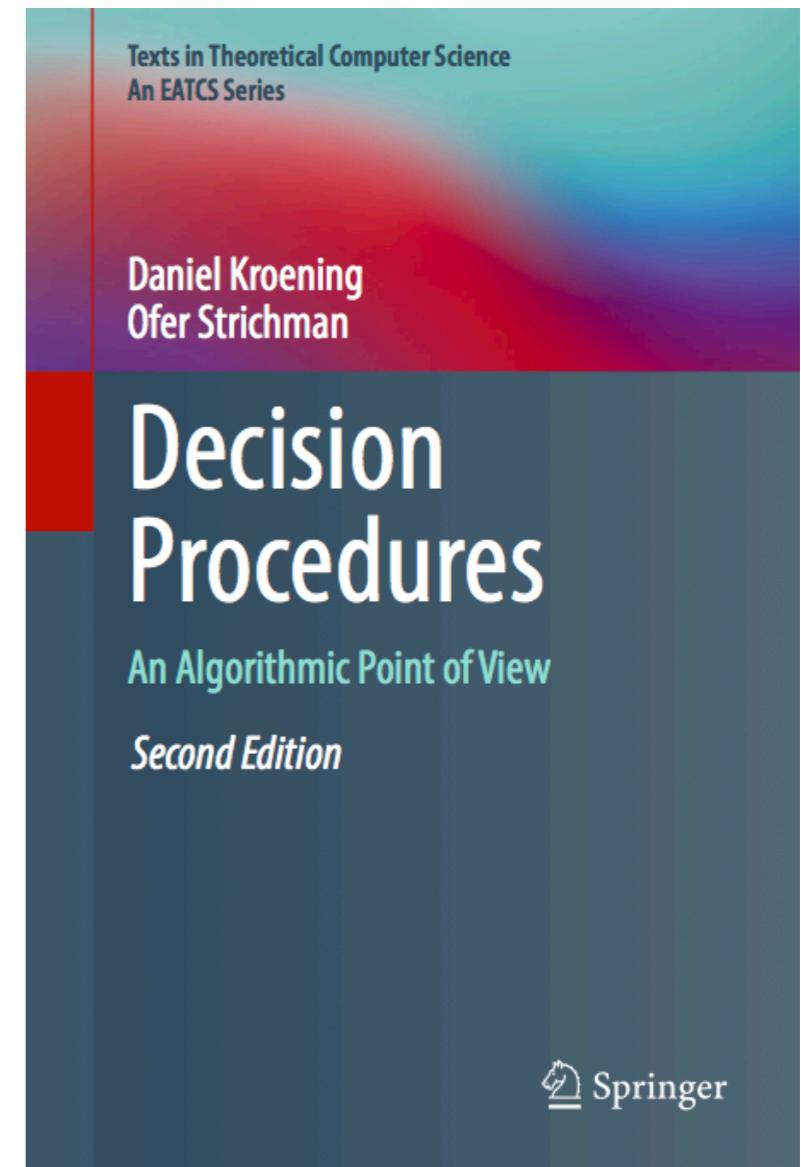
# Machine Learning

- Learning to decision



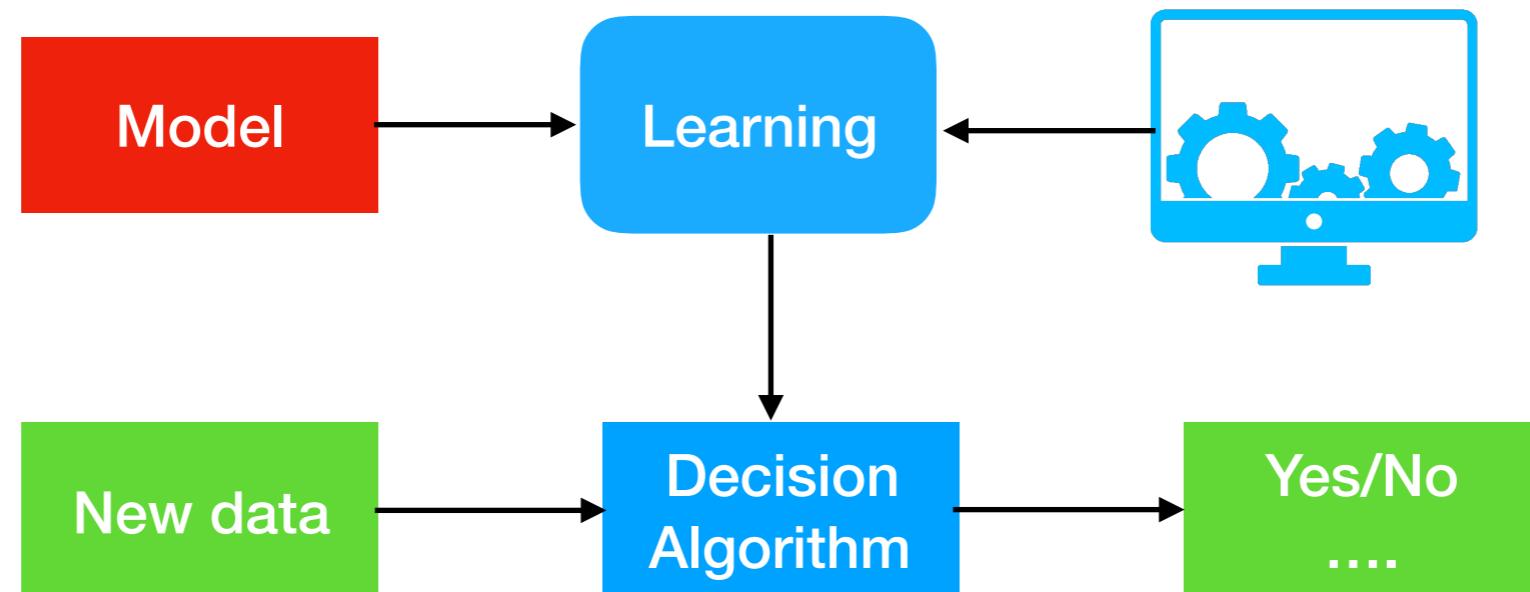
# Decision Procedures

- Satisfiability (SAT) - Propositional Logic
- Satisfiability Modulo Theory (SMT) - First Ordered Logic ...
- Integer Linear Programming (ILP) - Mathematical Optimization for integer variables
- Constrained Programming (CP) - No objective
- Machine Learning (ML)
- Model-based (Machine) Learning (MML)



# Model-based Machine Learning (MML)

- Instead of training data, a model is given



# Model-based Machine Learning (MML)

- No need of BIG training data
- Avoid UNSAFE/INSECURE training data
  - Ex) Pingpong → Tenis
- Need experience of building system models

# Stratego

- Strategy learning based on
  - Game theory
  - Model checking
- Generation, optimization, comparison as well as consequence and performance exploration of  
**strategies for stochastic priced timed games**

[<http://people.cs.aau.dk/~marius/stratego/intro.html>]

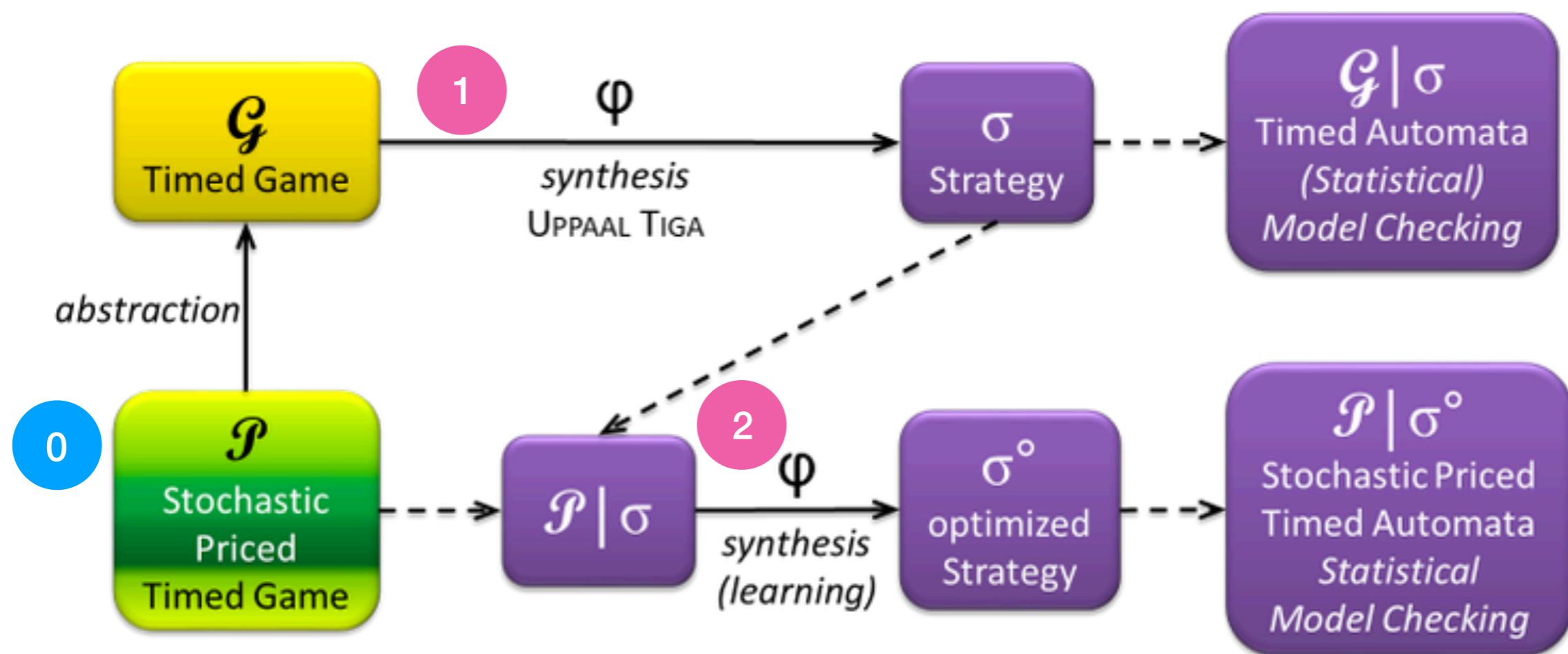
# Terminology

- A **game** is a mathematical model of a system consisting of several players (processes) which have independent objectives, often they are competing and sometimes opposing and even conflicting.
- A **strategy** is a prescription of one player's actions (transitions) for any possible situation (system state) eventually leading to achieving the player's goal.
  - **Deterministic strategy** specifies a single action per state and
  - **Non-deterministic strategy** may specify several alternatives.
  - **Fully permissive strategy** specifies all alternative actions leading to the goal.

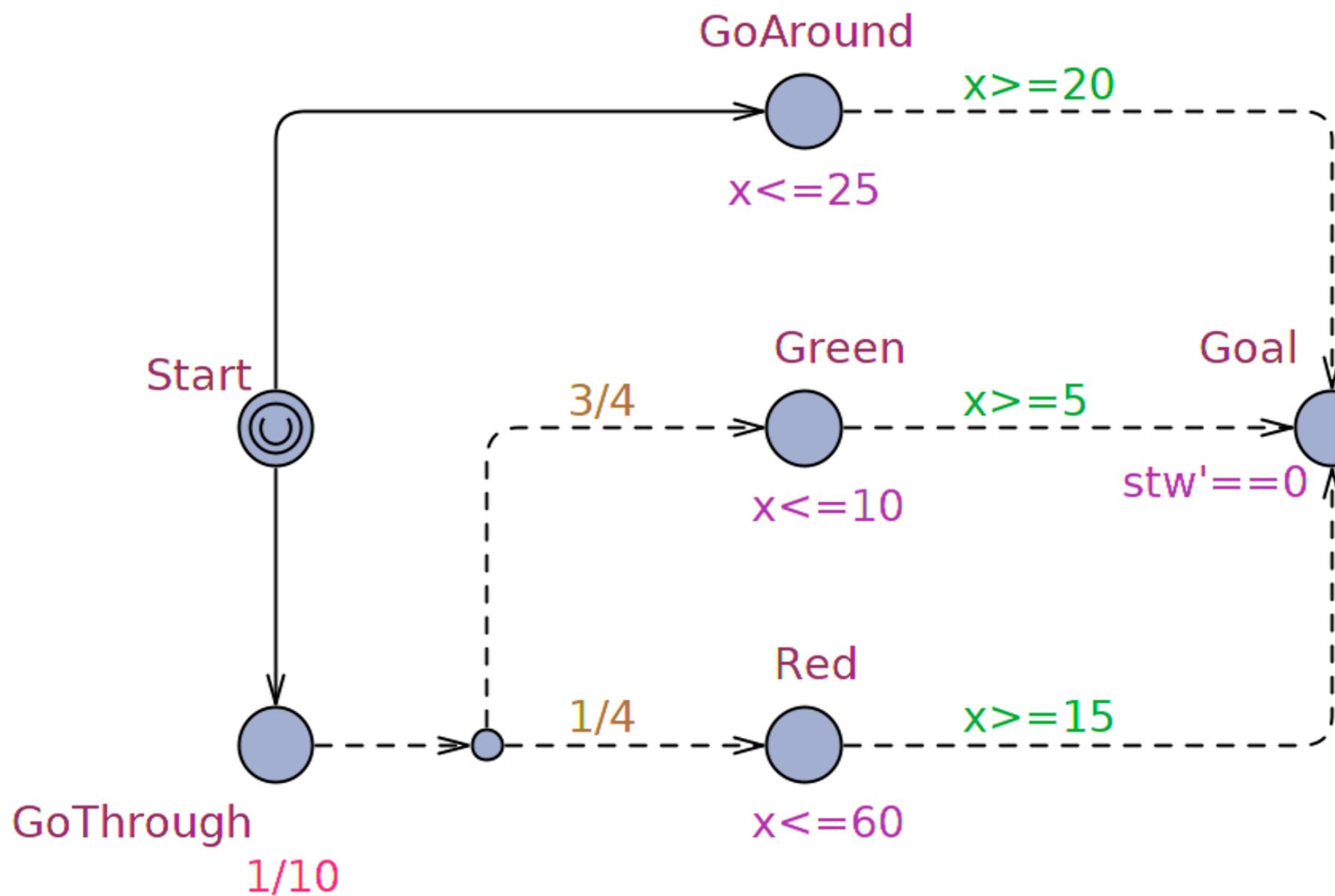
# Terminology

- Uppaal **timed game** is a two-player game specified by a network of timed automata with two types of edges:
  - Solid and dashed, corresponding to the player- and opponent-chosen transitions (Uppaal)
- A **stochastic timed game** is a probabilistic extension to a timed game.
- A **stochastic priced timed game** is a further extension with continuous price expressions which allow estimating the distribution of cost in a timed game.

# Overview

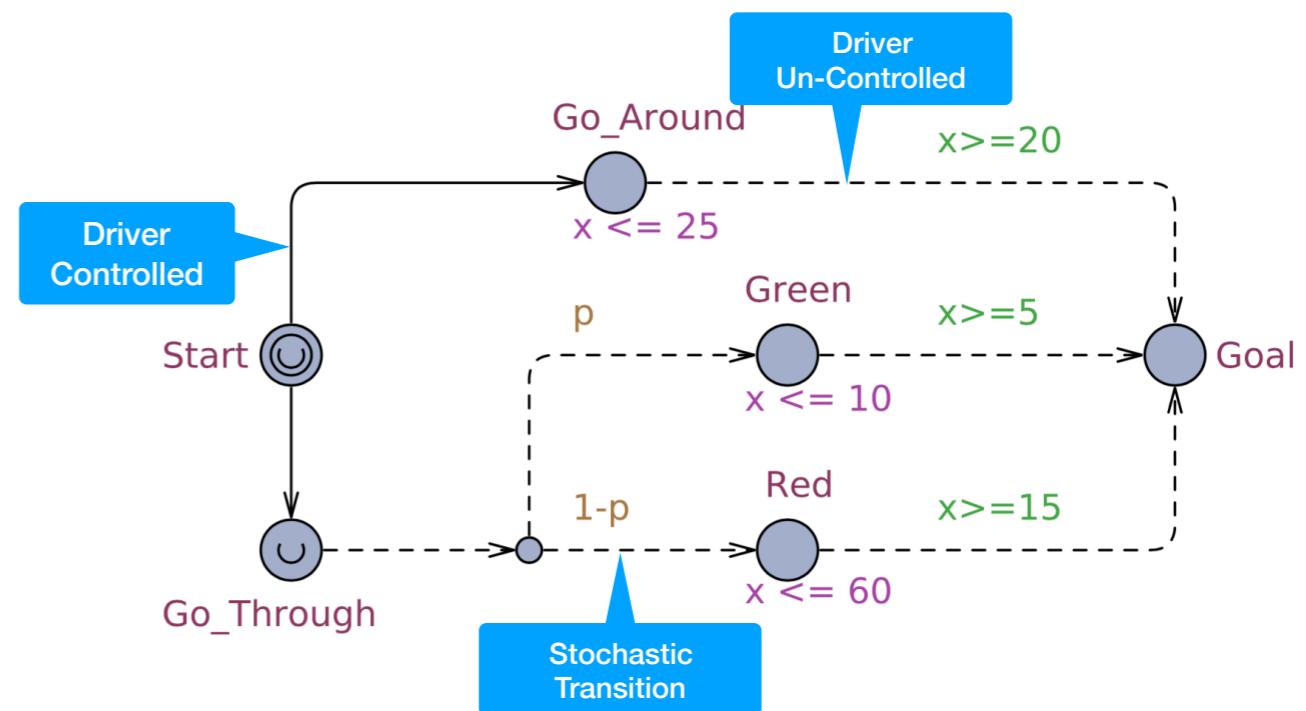


# Example



# 0. Construction (Constraint) Model

- Controlled transition
- Uncontrolled transition



# Formal Analysis

- $A<> (\text{carpassing.Goal} \&\& x \leq 25)$
- $\Pr[\leq 100] (<> \text{carpassing.Goal})$
- $E[\leq 100; 500] (\text{max: stw})$

# 1. Check Strategy (Objective)

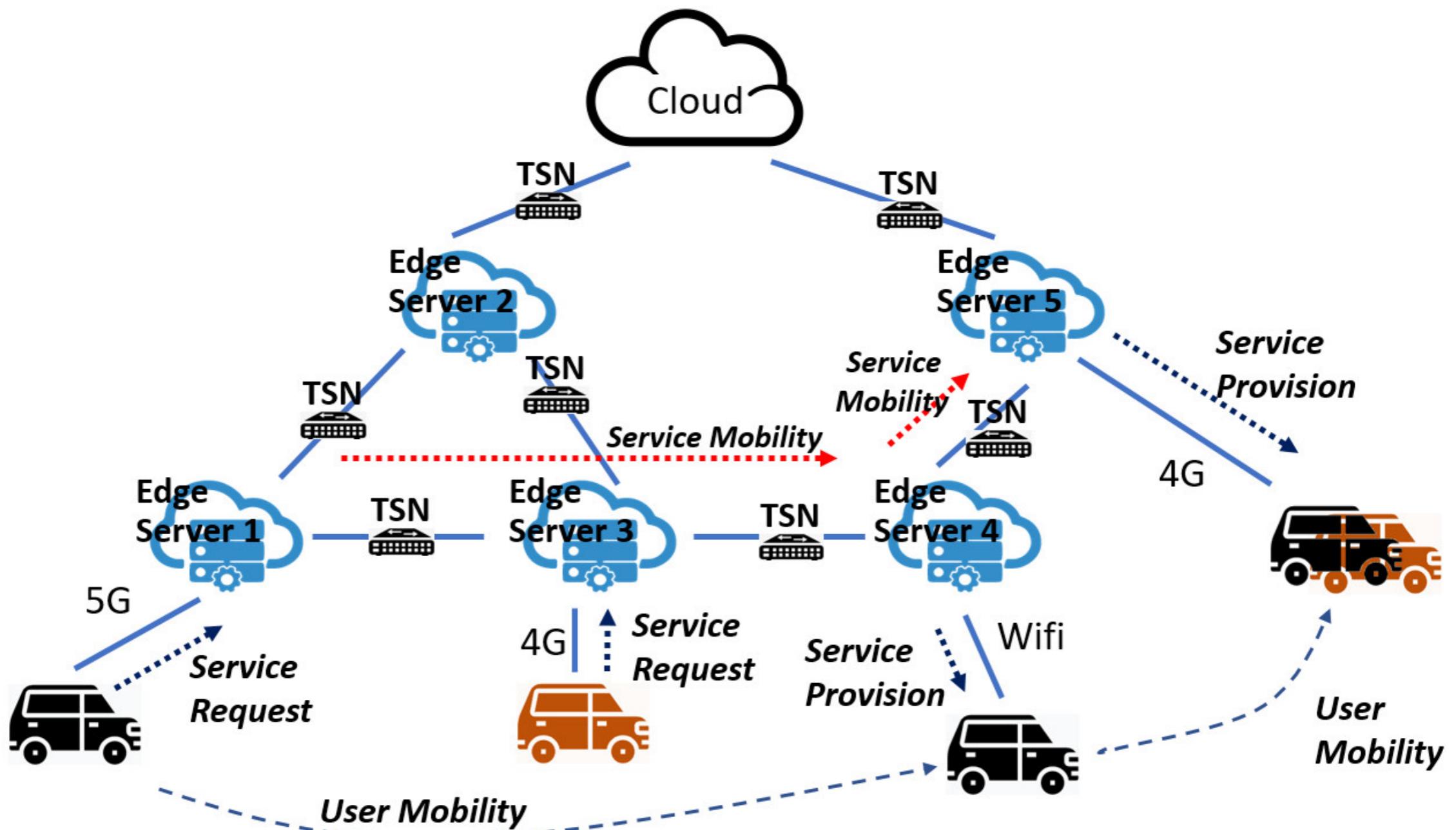
- **control: A<> carpassing.Goal && x <= 25**
  - Is it possible to choose a passage to pass the city within 25 minutes no matter which way is taken?
  - Answer: possible (satisfied).
- **strategy GoFast = control: A<> carpassing.Goal && x <= 25**
  - Is it possible to choose a passage to pass the city within 25 minutes no matter which way is taken?
  - Answer: possible (satisfied) and the tool computes fully permissive strategy named "GoFast".

# 2. Learning

- **A<> (carpassing.Goal && x<=25) under GoFast**
  - Is it possible to pass the city within 25 minutes under decisions of GoFast strategy?
  - Answer: yes (satisfied).
- **E[<=100; 500] (max: stw) under GoFast**
  - Estimate the time (needed to pass the city, where the clock time is stopped).

# **Use Case: Cloud-Based Real-Time SW Update**

# Cloud-based Real-Time SW Update



# Problem Description (Informal)

- For vehicles using connectivity to infrastructure through edges on paths and roads,
  - Assign SW app updates to edges such that an objective, such as minimized traveling time for a given update, is satisfied.

# Our Approach

- Model-based learning
  - Set up a model of the system and properties of the system to hold,
  - Find a strategy, i.e., a solution, that meets properties,
  - Verify that the system under the control guided by the strategy always holds given properties

# Assumptions

- An edge is defined by a range to cover, where a vehicle completes SW update at a specific edge,
- A vehicle has a different driving speed over edges, but determining driving speed of a vehicle over an edge is out of this work scope,
- An SW update is completed at one edge (one-to-one mapping)

# Model of System

Let  $m$  be a memory quantity that an SW update of a vehicle requires,  $m$  available quantity of memory on an edge, and  $c$  the computation time of a vehicle for processing each unit of memory.

A route is defined by a sequence of paths,

A path is defined by  $p = (e, tt_{best}, tt_{worst})$ , where  $e$  is an edge on the path, and  $tt_{best}$  and  $tt_{worst}$  are, respectively, the best-case traveling time and the worst-case traveling time to cross the edge.

$ut_i$  is a time that elapses to complete all the updates of apps for a vehicle  $v_i$ .

# Model of System

The underlying system comprises a set of edges and vehicles traveling over edges,

$E$  is a set of edges which relay between a cloud server and vehicles and formally defined by

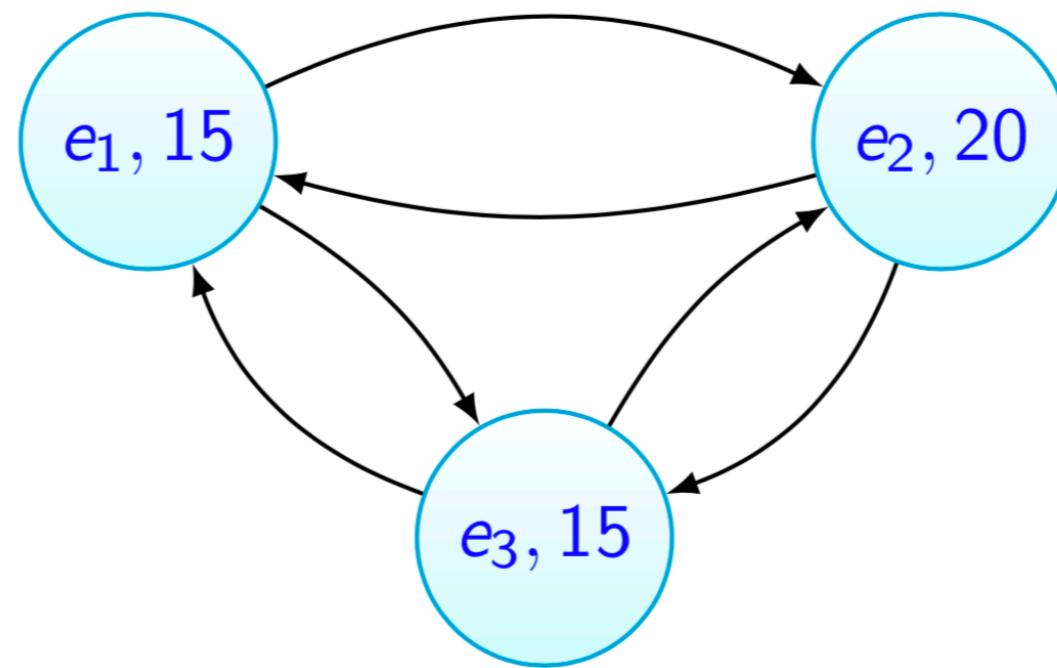
$E = \{e_k | e_k = (m_k, 2^E)\}$  where  $2^E$  is a subset of  $E$  which can be reached from  $e_k$ .

$V$  is a set of vehicles and defined by

$$V = \{v_i | v_i = (r_i, u_i, c_i)\}$$

where  $r_i$  is a route,  $u_i = \{m_j\}$  where  $m_j$  is a required memory capacity for the SW update of  $j$ , and  $c_i$  a processing time for update.

# Running Case



$$E = \{e_1 = (15, \{e_2, e_3\}), e_2 = (20, \{e_1, e_3\}), e_3 = (15, \{e_1, e_2\})\}$$

$$V = \{v_1 = ((e_1, 10, 15) \rightarrow (e_2, 14, 15) \rightarrow (e_1, 10, 15), \{3, 4\}, 2), \\ v_2 = ((e_2, 24, 35) \rightarrow (e_1, 10, 15) \rightarrow (e_3, 21, 25), \{1, 4, 2\}, 1), \\ v_3 = ((e_1, 10, 15) \rightarrow (e_2, 25, 36) \rightarrow (e_3, 21, 26), \{5, 6\}, 2)\}$$

# Constraints

Let  $upd_{i,j}$  be a predicate that indicates whether the update  $i$  of app for a vehicle  $j$  is over.

The constraints that we have: Given an edge  $k$ ,

$$\sum_{i \in V} \sum_{j \in u_i} upd_{i,j,k} \times m_j \leq \overline{m_k}$$

Given an update  $j$  and a vehicle  $i$ ,  $upd_{i,j,k}=1$

$$\sum_{k \in E} upd_{i,j,k} = 1$$

# Objectives

Given  $E$  and  $V$ , a scheduling strategy that assigns an update of a vehicle to an edge is computed such that

$$\min \sum_{i \in V} ut_i$$

# Edge and Vehicle Specifications

Listing 1: Specification of edges and vehicles in Uppaal models

```
// Edge's capacities
memsz_t curmemsz[eid_t] = {15, 20, 15};

//Vehicle's routes and duration for each edge
const paths_t v1path = {{1,10,15}, {2,14,15}, {1,10,15}, nullpath} ;
const paths_t v2path = {{2,24,35}, {1,10,15}, {3,21,25}, nullpath} ;
const paths_t v3path = {{1,10,15}, {2,25,36}, {3,21,26}, nullpath} ;

//Vehicle update requirements
upd_t v1upds[updid_t] = {{3, 0},{4, 0}, noneupd};
upd_t v2upds[updid_t] = {{1, 0},{4, 0}, {2, 0}};
upd_t v3upds[updid_t] = {{5, 0},{6, 0}, noneupd};

Tony = Veh(1, v1upds, 2, v1path, curmemsz);
Nany = Veh(2, v2upds, 1, v2path, curmemsz);
Jony = Veh(3, v3upds, 2, v3path, curmemsz);
```

# Vehicle Model

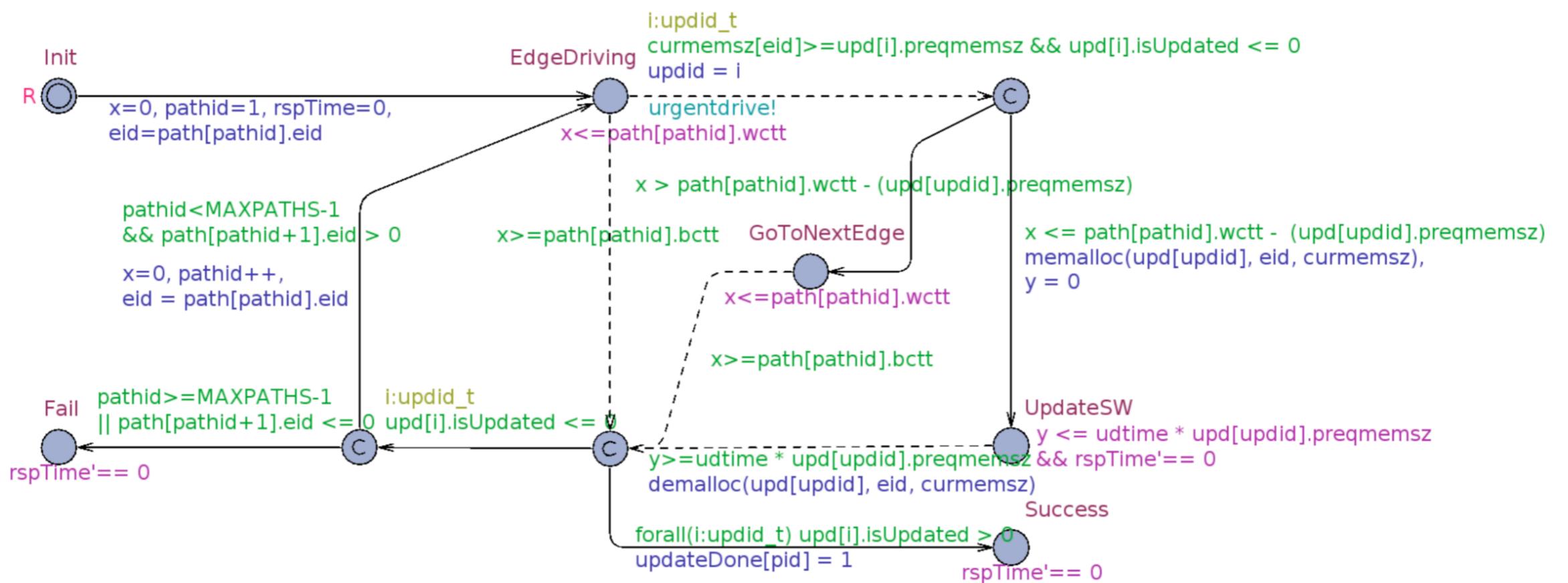


Figure: Vehicle model in stochastic timed automata

# Queries (Symbolic)

Check if all updates of 3 vehicle finishes within 100 time units, ( $\phi$ )

A < $\rightarrow$  updateDone[1] > 0 && updateDone[2] > 0 && updateDone[3] > 0 && time <= 100

Search for a strategy ( $\sigma$ ), named  $\text{BoundedUT}$ , that all updates for 3 vehicles are over within 100 time units.

strategy  $\text{BoundedUT}$  = control : A < $\rightarrow$  updateDone[1] > 0 && updateDone[2] > 0 && updateDone[3] > 0 && time <= 100

Check if all updates of 3 vehicle finishes within 100 time units under the strategy  $\text{BoundedUT}$ , ( $P \mid \sigma \models \phi$ )

A < $\rightarrow$  updateDone[1] > 0 && updateDone[2] > 0 && updateDone[3] > 0 && time <= 100 under  $\text{BoundedUT}$

# Queries (Statistical)

Search for a strategy that all updates for 3 vehicles are over such that the update time for the vehicles are minimized,

**strategy GoFast = minE(time)[<= 100] : <> updateDone[1] > 0 &&**  
**updateDone[2] > 0 && updateDone[3] > 0**

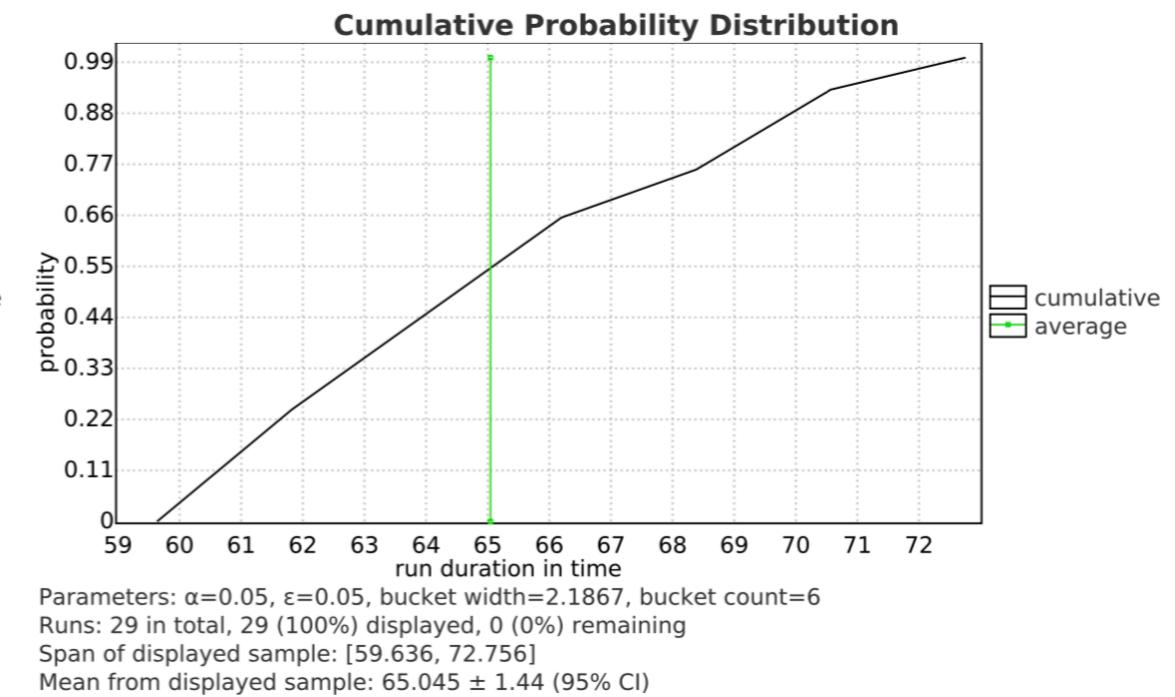
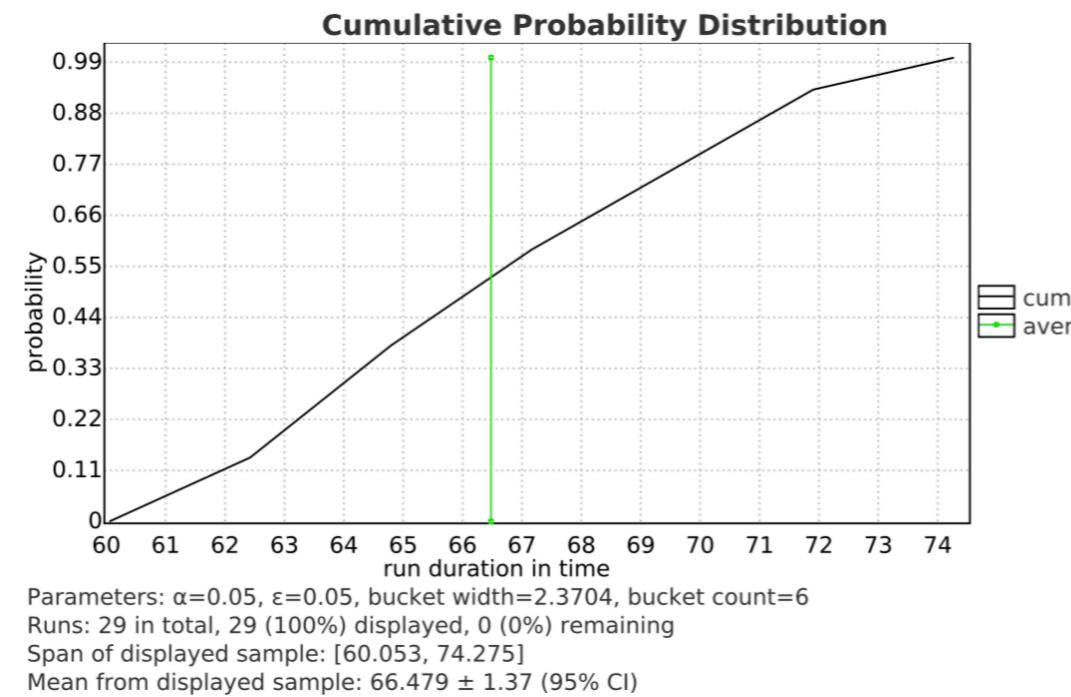
Check the possibility that all update of 3 vehicles are over

**Pr[<= 100](<> updateDone[1] > 0 && updateDone[2] > 0 &&**  
**updateDone[3] > 0)**

Check the possibility that all update of 3 vehicles are over under the minimum response time

**Pr[<= 100](<> updateDone[1] > 0 && updateDone[2] > 0 && updateDone[3] > 0)**  
**under GoFast**

# Cumulative Probability Distributions

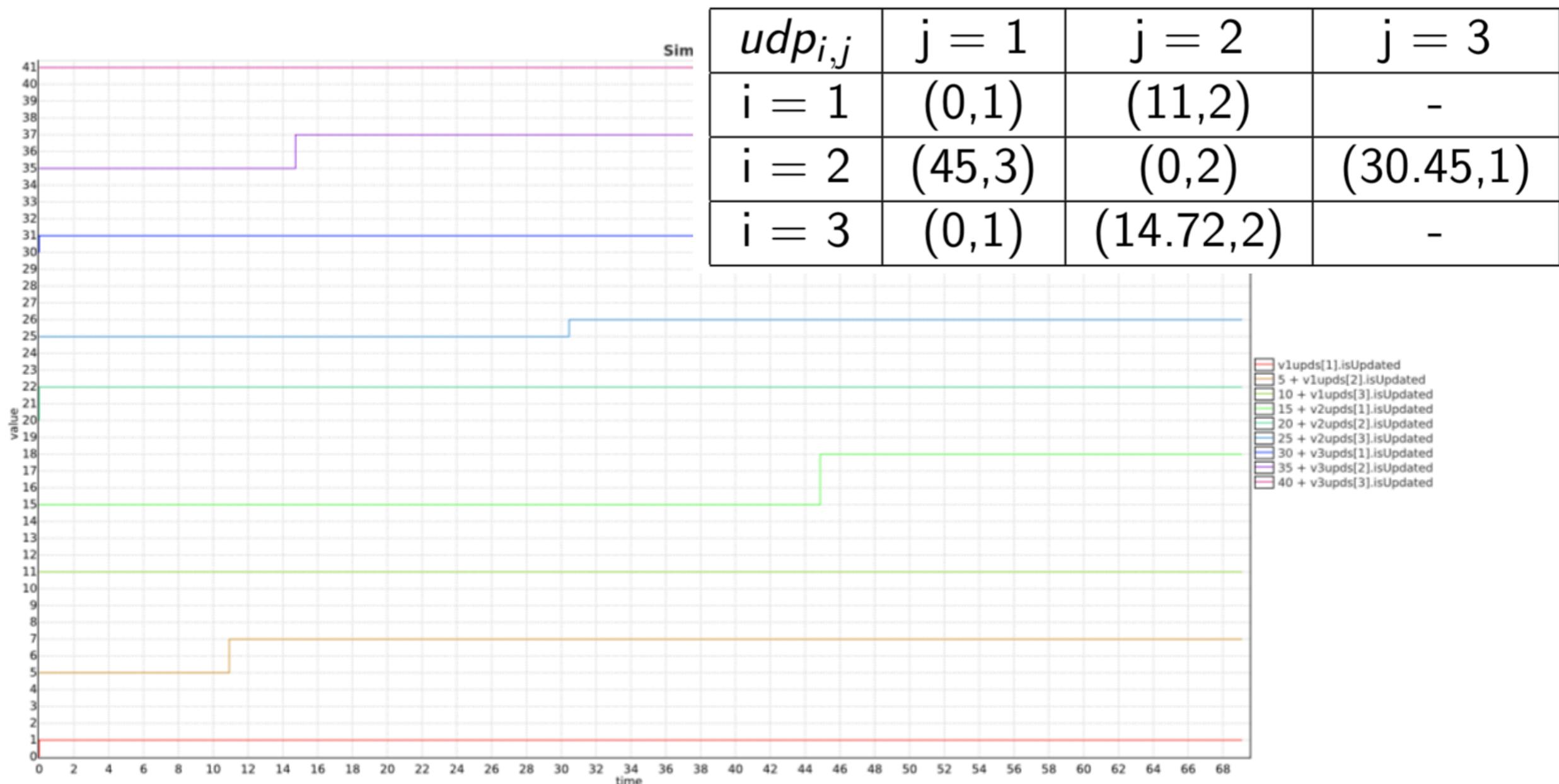


The model not under strategy finished in average 66.479 time units in between [60.053, 74.275], but the model not under strategy finished in average 65.045 time units in between [59.636, 72.756]

# Maximum SW Update Time

- The maximum SW update time:
  - Simulation-based estimation : 65.1763
    - $E[<= 1000;1000]$  (max : sysrsp time)
  - Simulation-based estimation under the strategy control: 65.026
    - $E[<= 1000;1000]$  (max : sysrsp time) under GoFast

# SW Update Strategy with Mim-Update-Time Strategy



# Conclusions

- Machine learning is a next generation of software programming
- Decision procedure is ALSO one of machine learning techniques,
  - Instead of data, Stratego generates a constrained behavior trace (data) from a model, which is guaranteed by controller
  - Create a control algorithm that guarantees user-defined properties.