

# Recursion

Thumrongsak Kosiyatrakul  
tkosiyat@cs.pitt.edu

- A Java method can call any other public Java method
  - `main()` is just a method itself, and we have called other methods from it
  - Thus, a method should be able to call itself
    - We call this a **Recursive Call**
  - At first thought this seems odd or even impossible
    - Why should we want to do this?
  - However, it will be very useful in a lot of different programming approaches

- Before we look at the programming in detail, let's try to get the idea down, using math
- Some mathematical functions are in fact defined recursively
  - Example: Factorial

$$n! = n \times (n - 1)!$$

- Note that the function is fedined in terms of itself, but with an important change:
  - The **recursive call** is smaller in size ( $n - 1$ ) than the original call  $n$
  - This is vital to recursion being viable
- Let's trace  $4!$  on the board in this way to see what happens
  - But let's stop when we reach  $-3!$

- What we are missing in the previous slide is a condition that allows the recursion to stop
  - Every recursive algorithm must have some terminating condition, to keep it from recursing **forever**
  - We call this the **base case**
- What is the base case for factorial?

$$0! = 1$$

- This now allows us to complete our algorithm:

$$n! = \begin{cases} 1 & \text{if } n = 0 \\ n \times (n - 1)! & \text{Otherwise} \end{cases}$$

- Three important rules for any recursive algorithm:
  - ① There must be some **recursive case**, in which the algorithm **calls itself**
  - ② There must be some **base case**, in which no recursive call is made
  - ③ The **recursive calls must lead** eventually **to the base case**
    - Usually by “reducing” the problem size in some way
- Do not forget these rules!!!

# More Recursion

- Let's look at another example:
  - Calculating an integer power of another integer

$$m^n = m \times m^{n-1} \quad \text{where } n > 0$$

- Do not forget about the **base case**

$$m^n = 1 \quad \text{where } n = 0$$

- The actions we take are slightly different from factorial, but the basic idea is similar
- Trace this on board
  - Note how **ifrst call made is last call to complete**
  - This is important in the implementation of recursion

# Implementing Recursion

- So, how do we implement recursion?
  - Luckily, the computer code is very similar to the mathematical functions
  - Consider factorial below:
    - Note that the recursive call is made within the return statement
    - This is fine – return is done **after** call completes

```
public static int factorial(int n)
{
    if(n <= 1)
        return 1;
    else
        return n * factorial(n - 1);
}
```

# Implementing Recursion

- How does recursion actually work?
  - Each time a method is called, an **activation record (AR)** is allocated for it
    - This consists of memory for the parameters and local variables used in the method
  - Each new activation record is placed on the top of the **run-time stack**
  - When a method terminates, its activation record is removed from the top of the run-time stack
  - Thus, the first activation record placed onto the stack is the last one removed



# Activation Record and Run-Time Stack

```
n = 4
n <= 1? No
return 4 * factorial(3); =
```

**factorial(4)**

# Activation Record and Run-Time Stack

```
n = 3
n <= 1? No
return 3 * factorial(2); =
```

**factorial(3)**

```
n = 4
n <= 1? No
return 4 * factorial(3); =
```

**factorial(4)**

# Activation Record and Run-Time Stack

n = 2  
n <= 1? No

?

return 2 \* factorial(1); = ?

**factorial(2)**

n = 3  
n <= 1? No

?

return 3 \* factorial(2); = ?

**factorial(3)**

n = 4  
n <= 1? No

?

return 4 \* factorial(3); = ?

**factorial(4)**

# Activation Record and Run-Time Stack

```
n = 1  
n <= 1? Yes  
return 1;
```

= 1

**factorial(4)**

```
n = 2  
n <= 1? No  
return 2 * factorial(1);
```

?

= ?

**factorial(2)**

```
n = 3  
n <= 1? No  
return 3 * factorial(2);
```

?

= ?

**factorial(3)**

```
n = 4  
n <= 1? No  
return 4 * factorial(3);
```

?

= ?

**factorial(4)**

# Activation Record and Run-Time Stack

```
n = 1  
n <= 1? Yes  
return 1;
```

= 1

**factorial(4)**

```
n = 2  
n <= 1? No  
return 2 * factorial(1);
```

1

= ?

**factorial(2)**

```
n = 3  
n <= 1? No  
return 3 * factorial(2);
```

?

= ?

**factorial(3)**

```
n = 4  
n <= 1? No  
return 4 * factorial(3);
```

?

= ?

**factorial(4)**

# Activation Record and Run-Time Stack

n = 2  
n <= 1? No

1

return 2 \* factorial(1); = 2

factorial(2)

n = 3  
n <= 1? No

?

return 3 \* factorial(2); = ?

factorial(3)

n = 4  
n <= 1? No

?

return 4 \* factorial(3); = ?

factorial(4)

# Activation Record and Run-Time Stack

`n = 2`  
`n <= 1? No`  
`return 2 * factorial(1);` = `2`

1

`factorial(2)`

`n = 3`  
`n <= 1? No`  
`return 3 * factorial(2);` = `?`

2

`factorial(3)`

`n = 4`  
`n <= 1? No`  
`return 4 * factorial(3);` = `?`

?

`factorial(4)`

# Activation Record and Run-Time Stack

```
n = 3
n <= 1? No
return 3 * factorial(2); = 6
```

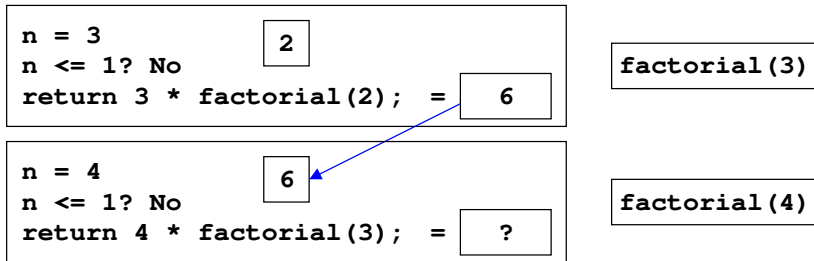
**factorial(3)**

```
n = 4
n <= 1? No
return 4 * factorial(3); = ?
```

**factorial(4)**



# Activation Record and Run-Time Stack



# Activation Record and Run-Time Stack

```
n = 4  
n <= 1? No  
return 4 * factorial(3); = 24
```

6

**factorial(4)**

# Activation Record and Run-Time Stack

24 → To caller

# Recursion vs Iteration

- Some recursive algorithms can also be easily implemented **with loop**
  - Both factorial and power can easily be done in this way
  - When possible, it is usually better to use iteration, since we do not have the overhead of the run-time stack (that we just saw on the previous slides)
- Other recursive algorithms are very difficult to do any other way (Tower of Hanoi in the textbook)
- You will see more about recursion in CS 0445

- Consider again **functional abstraction**
  - User of a method does not need to know how it is implemented
  - However, often recursive methods require more parameters than equivalent iterative methods
    - Extra parameters enable the testing for base cases
    - This can be problematic if the methods are part of an interface, which specifies the method header
  - We can get around this by using an additional, non-recursive method
    - The additional method satisfies the required header
    - It then calls the recursive method, adding any extra needed parameters

# More Recursion

- For example, consider a method to reverse an array of `int`
  - The header might be something like

```
public static void reverse(int[] data)
```

- However, to implement this recursively, we need extra parameters to keep track of the logical beginning and end of the array
- These extra parameters can be added in a call to the recursive method

```
public static void reverse(int[] data)
{
    recursive_reverse(data, 0, data.length - 1);
}
```

- Let's look at the `recursive_reverse()`

# recursive\_reverse()

- To reverse an array of size 6, we need to perform the following:
  - ① Swap data between index 0 and 5
  - ② Swap data between index 1 and 4
  - ③ Swap data between index 2 and 3
- In other words, to reverse an array
  - ① Swap data between the firstIndex and the lastIndex
  - ② Reverse the same array but
    - change the firstIndex to firstIndex + 1
    - change the lastIndex to lastIndex - 1
- This is why the signature of the recursive\_reverse() method is

```
public static void recursive_reverse(int[] data,  
                                     int firstIndex,  
                                     int lastIndex)
```