

ECE 6504

Embedded and Controls Systems Security: Stack Smashing V

Reading:

T:kBouncer, ROPGuard, ROPecker, and Evaluating the Effectiveness of Current Anti-ROP Defenses
N:Security Based on Physical Unclonability and Disorder



“Learn from the mistakes of others. You can never live long enough to make them all yourself.”

–Groucho Marx

stack protection methods:

1. make the overflow difficult to exploit
2. make overflow noticeable
3. make it impossible

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→
assuming failure

detect ROP activity

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assuming failure

detect ROP activity

- ↓
1. fixed load addresses
 2. shared libraries incompatible with ASLR
 3. base address of a DLL can be either calculated via leaked pointer or brute-forced

an ideal ROP detector:

1. operates w/binaries (no source needed)
2. w/o debug symbols
3. does not insert code into calls
4. low overhead (time, memory)

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	ROP Types	No Source Code	No Binary Rewriting	Run-time Efficiency
DROP [1]	Ret-based	✓	X	X
ROPDefender [2]	Ret-based	✓	X	X
ROPGuard [3]	Ret-based	✓	X	✓
Return-less Kernel [4]	Ret-based	X	✓	✓
CFLocking [5]	All	X	✓	✓
G-Free [6]	All	X	✓	✓
ILR [7]	All	✓	X	✓
Binary Stirring [8]	All	✓	X	✓
IPR [9]	All	✓	X	✓
CCFIR [10]	All	✓	X	✓
KBouncer [11]	All	✓	X	✓
ROPecker	All	✓	✓	✓

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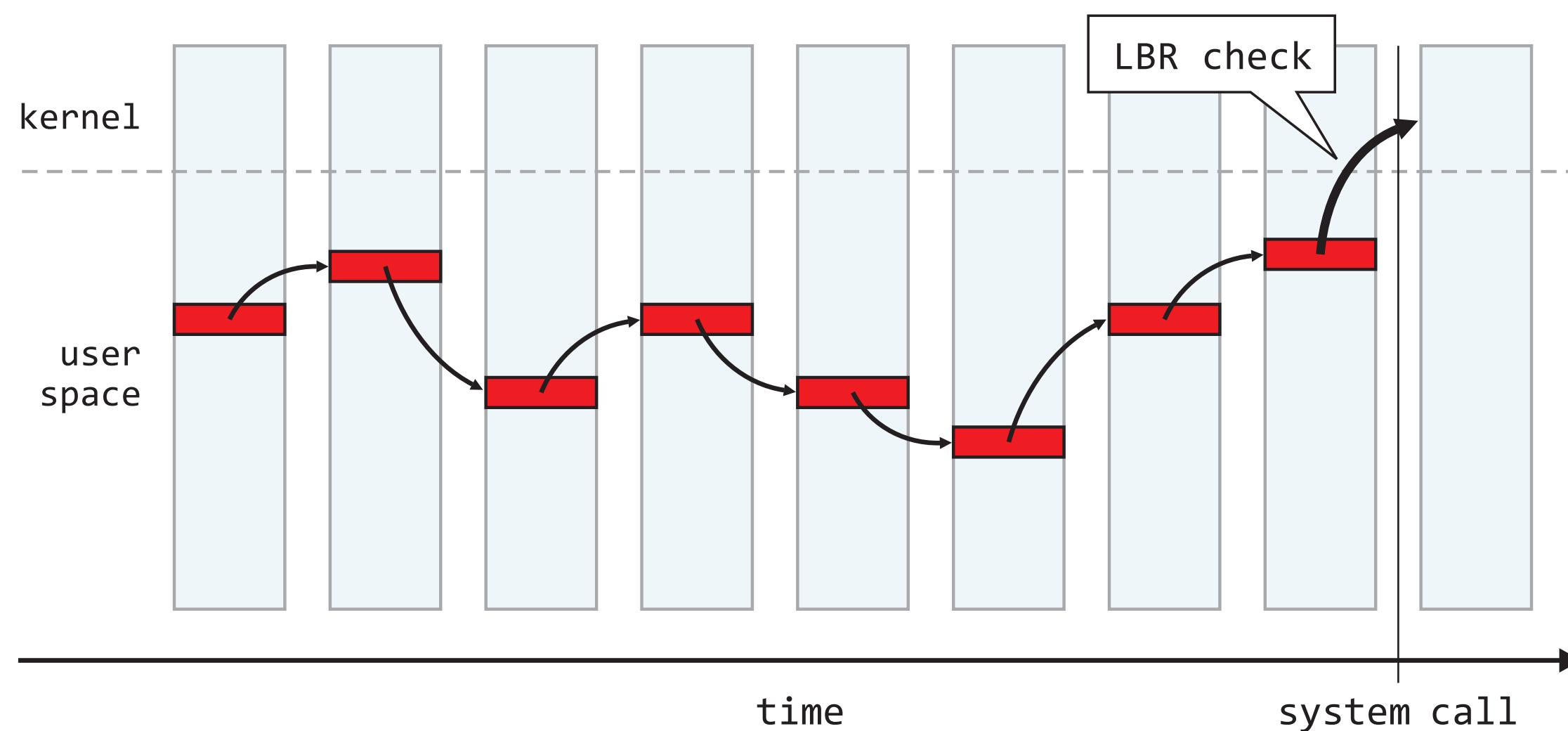
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system call anatomy (UNIX)

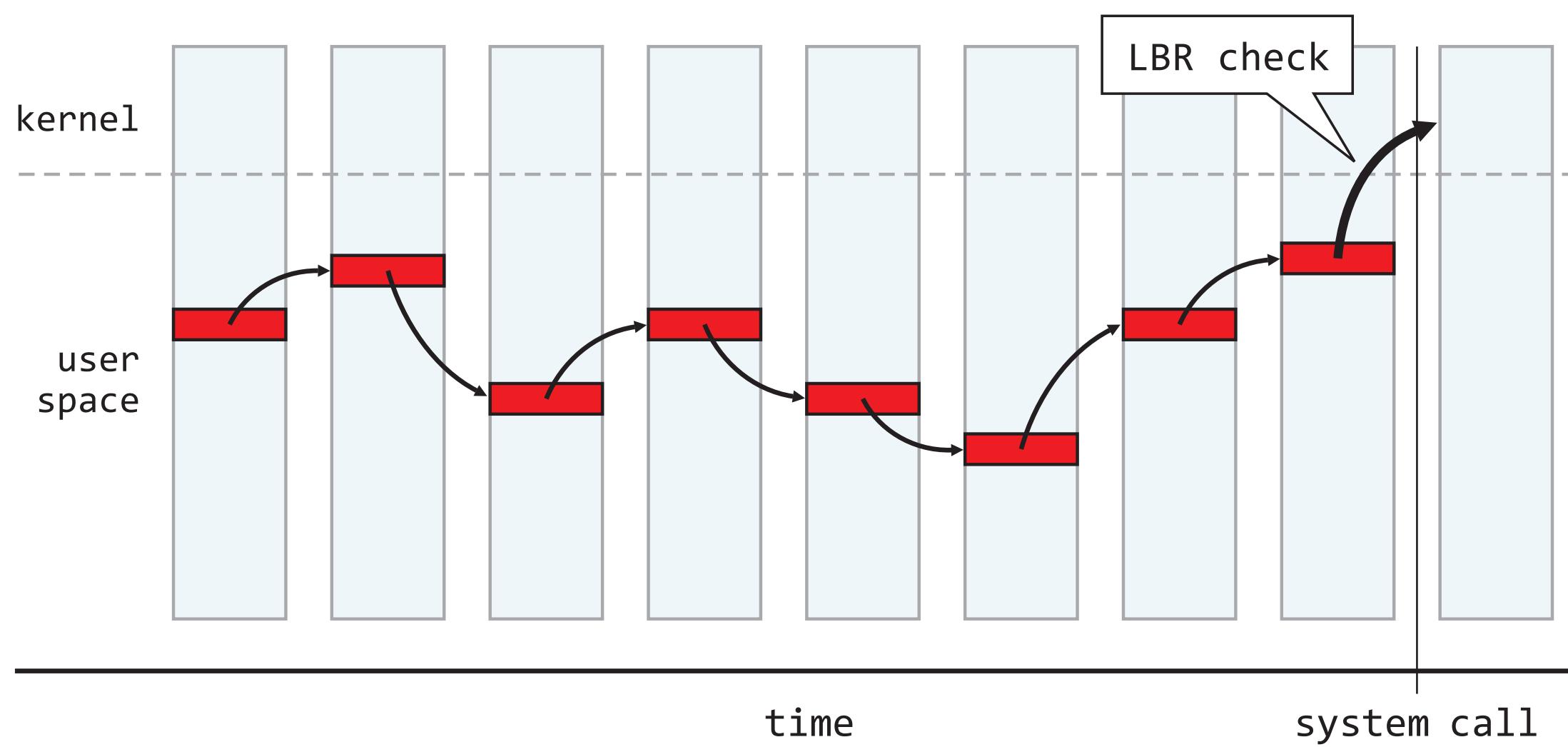
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2. wrapper function for call
(exposed directly)

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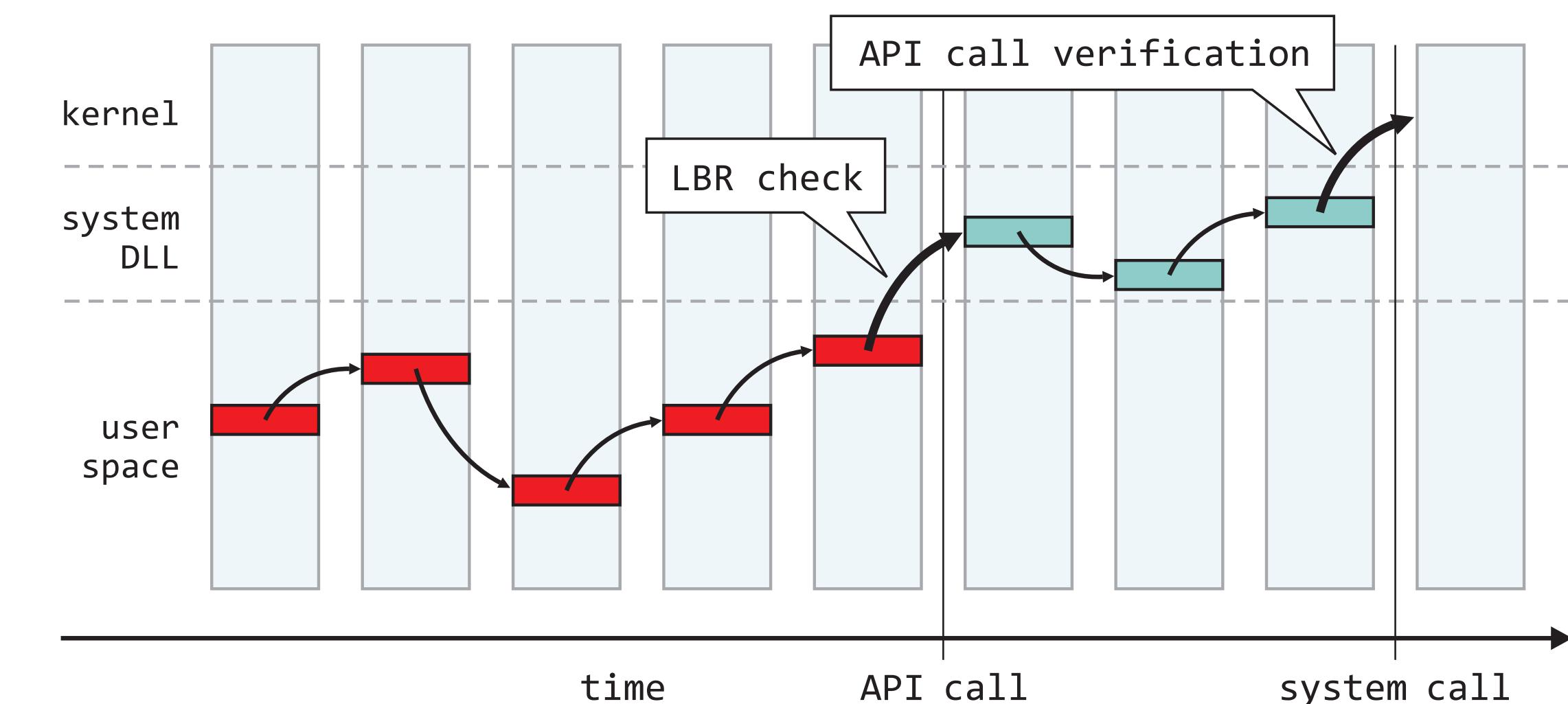
system call anatomy (Windows)

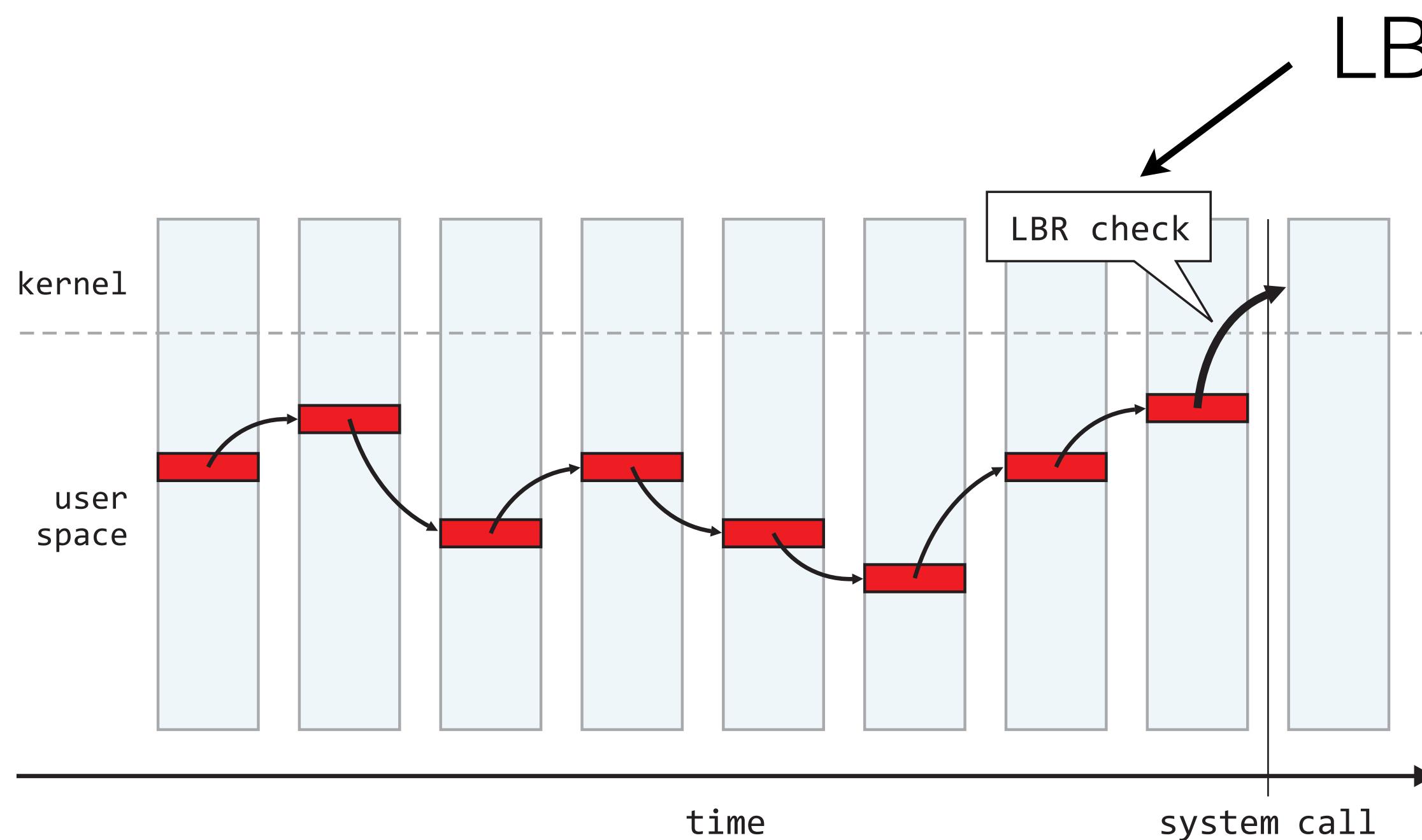
1. indirect branches
2. Windows API
3. Native API (not exposed)



system call anatomy (UNIX)

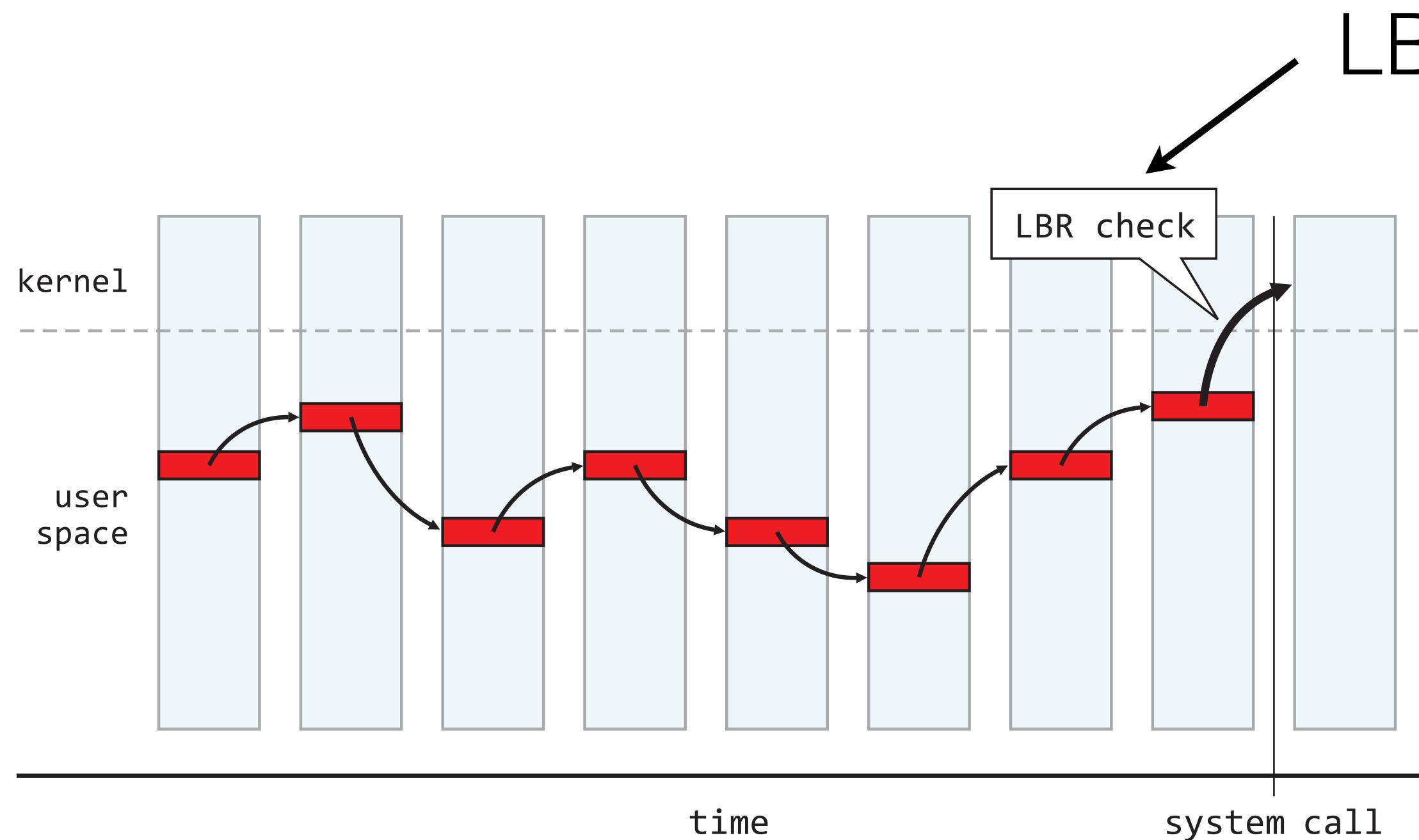
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LBR:

1. last branch record
2. hardware debug
3. last 16 branches
 - a. address of
 - b. target of
4. kernel-mode only

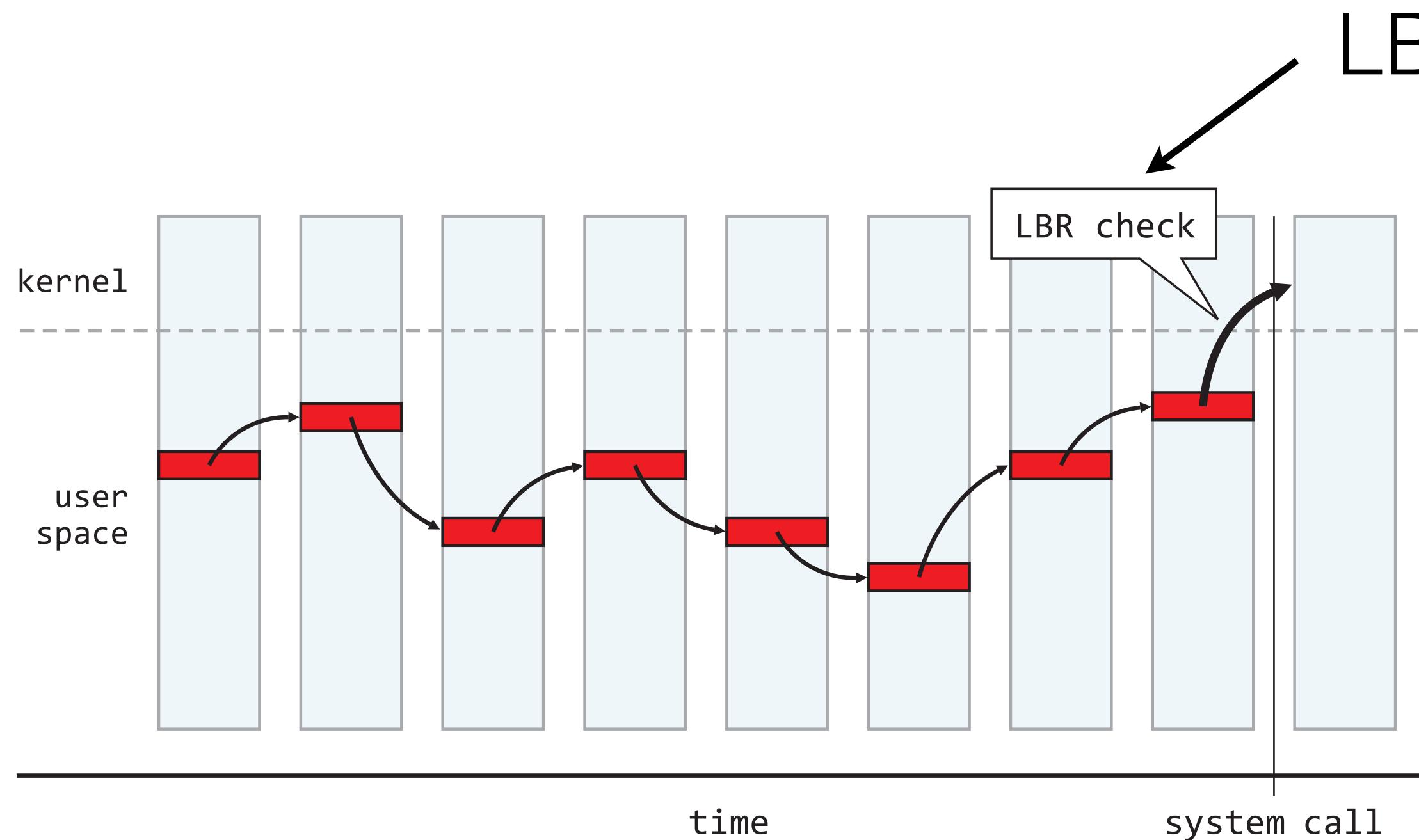


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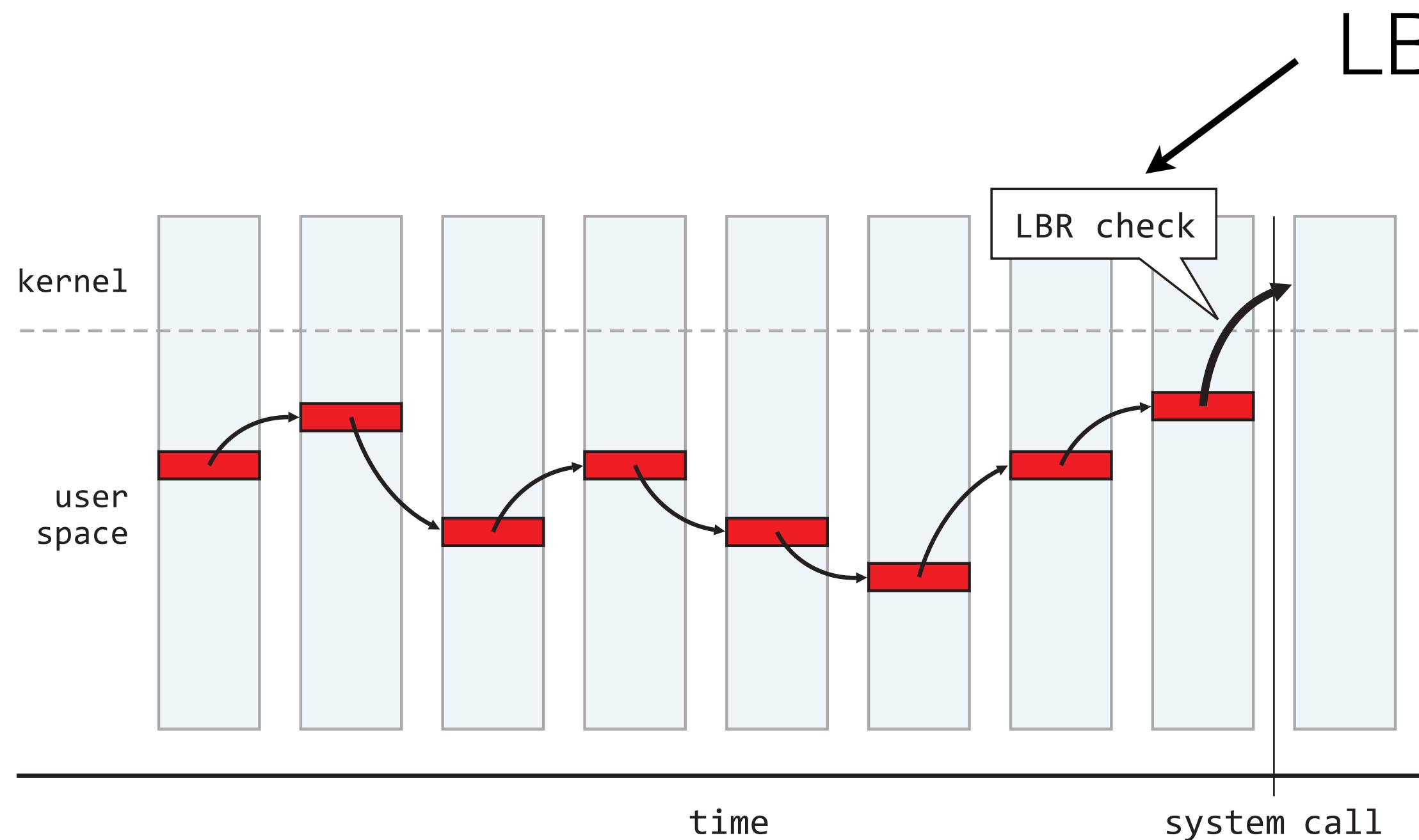
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non-malicious code does too



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use LBR to distinguish

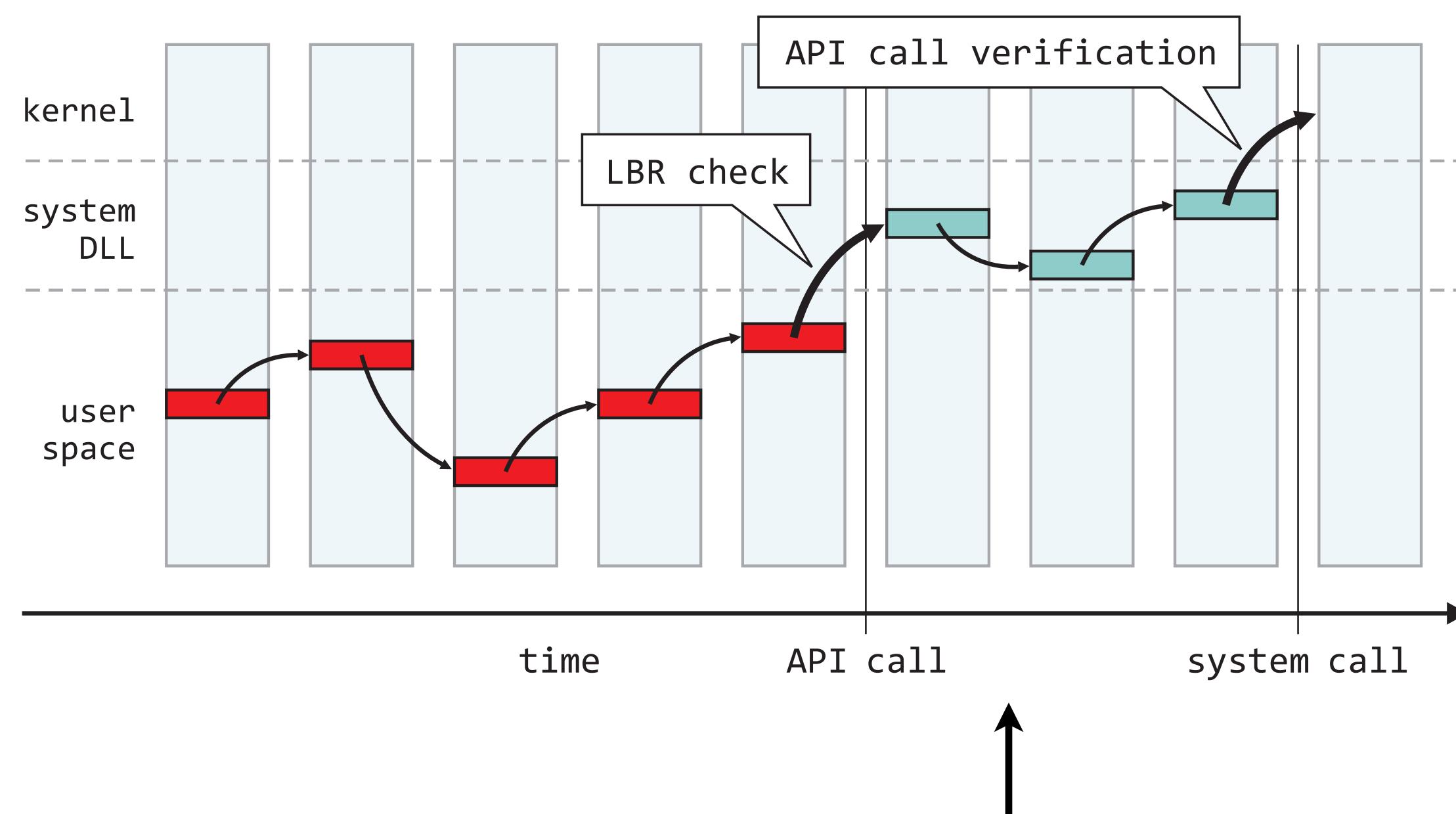


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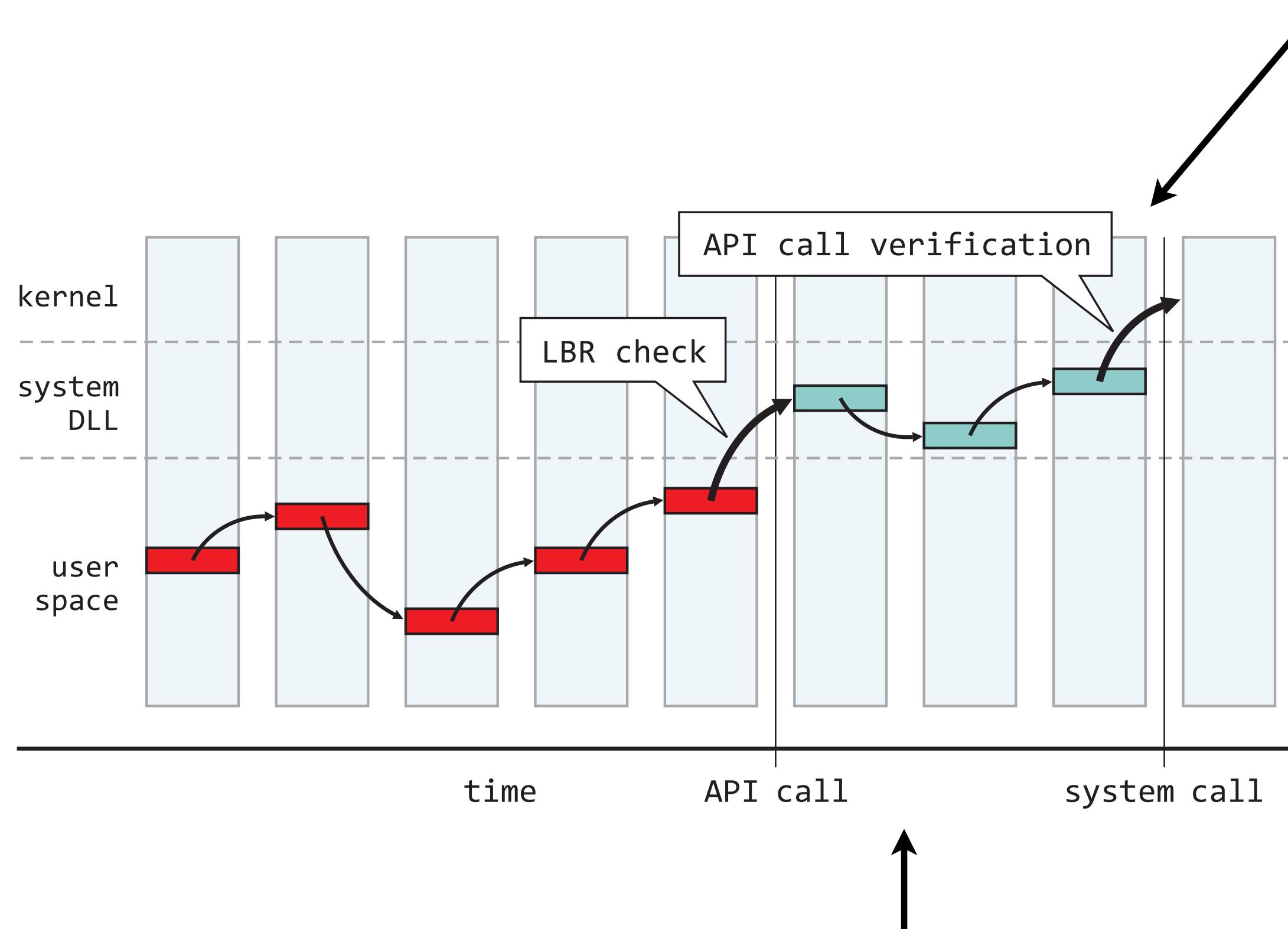


exploit code:

1. does not _(usually) invoke system call directly
2. detection code runs at _(select) API calls

if it does:

easy to verify that API
hasn't made call



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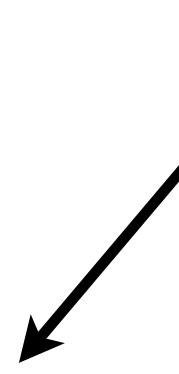
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1. illegal returns
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violation of calling convention



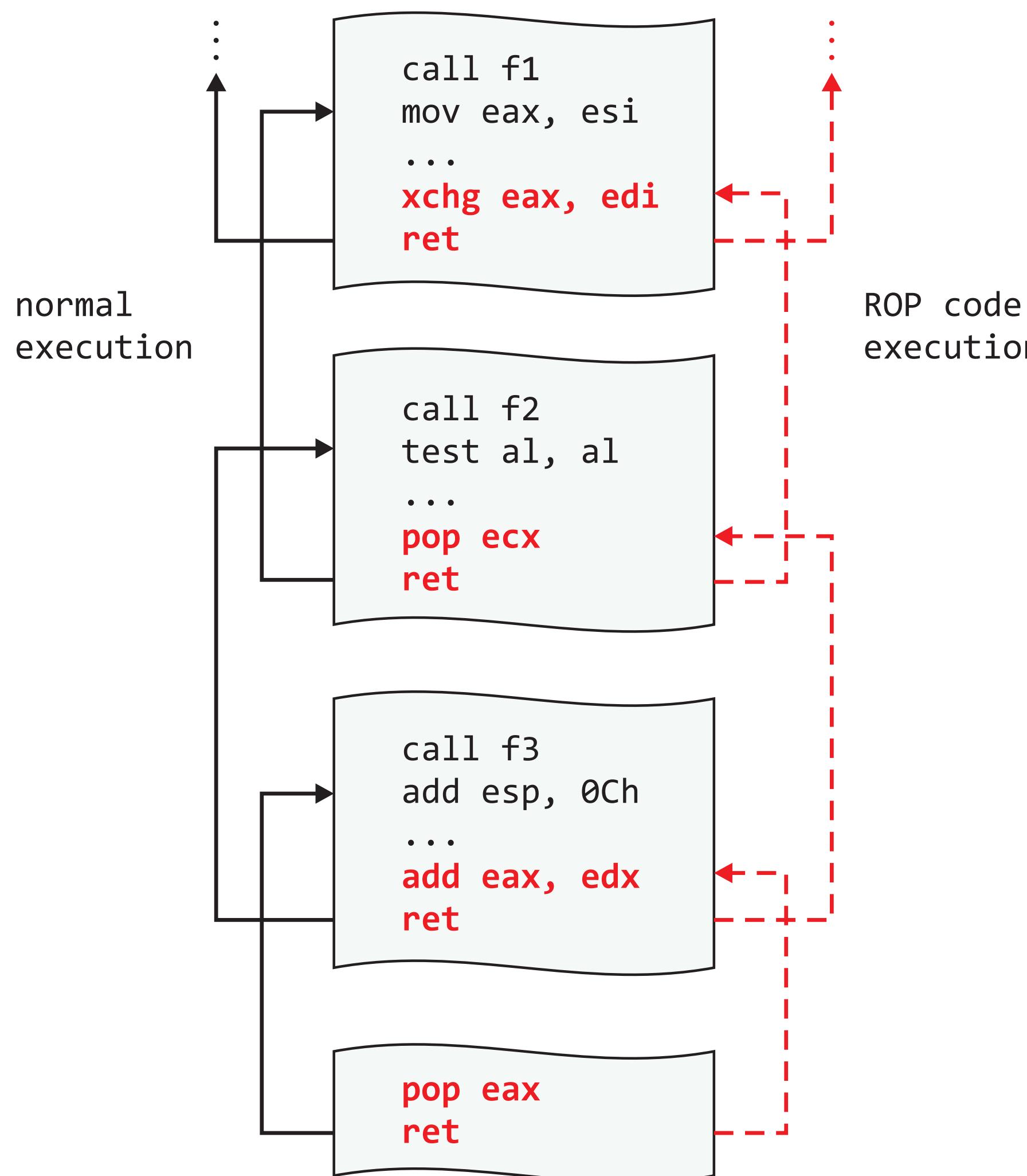
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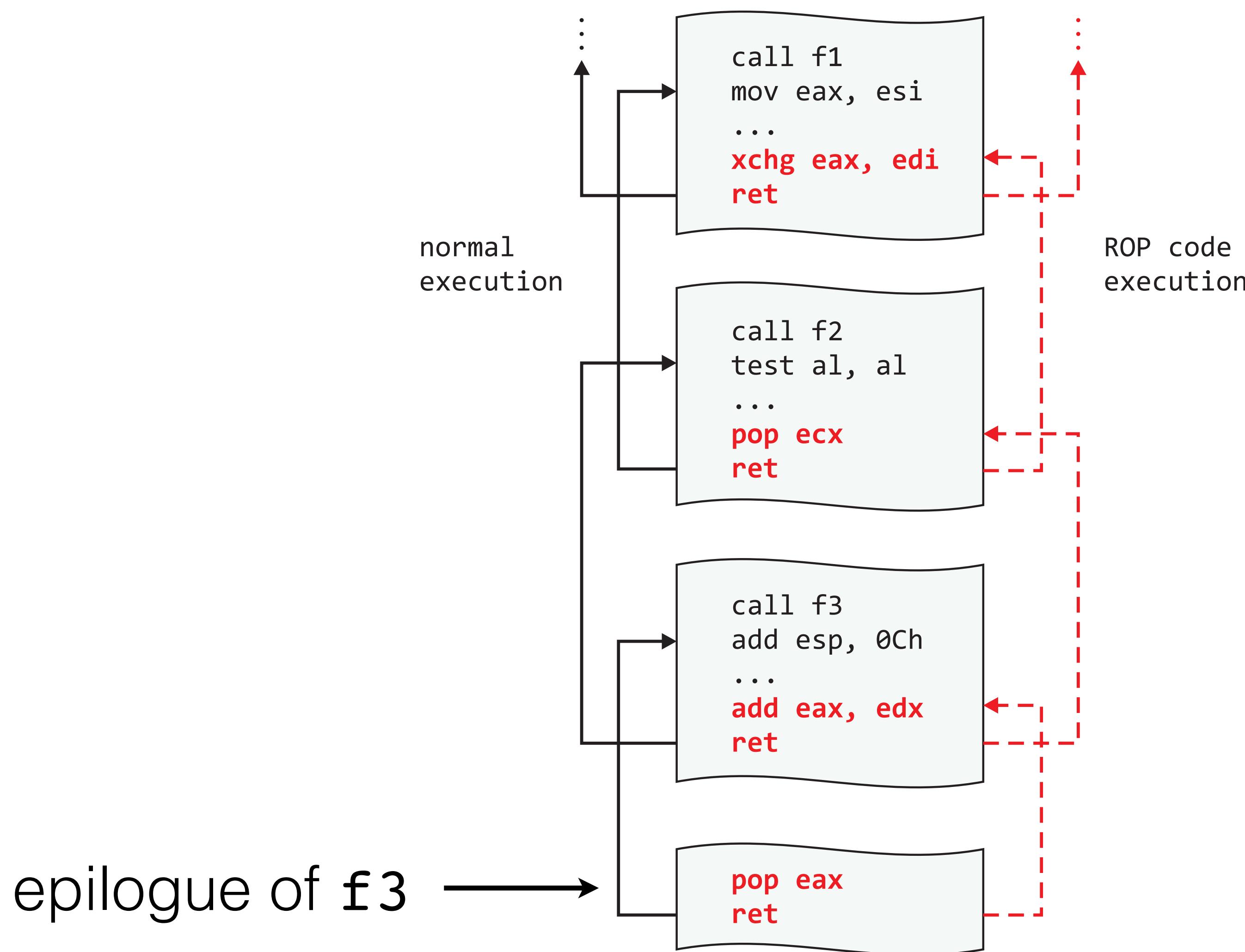
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too many indirect calls

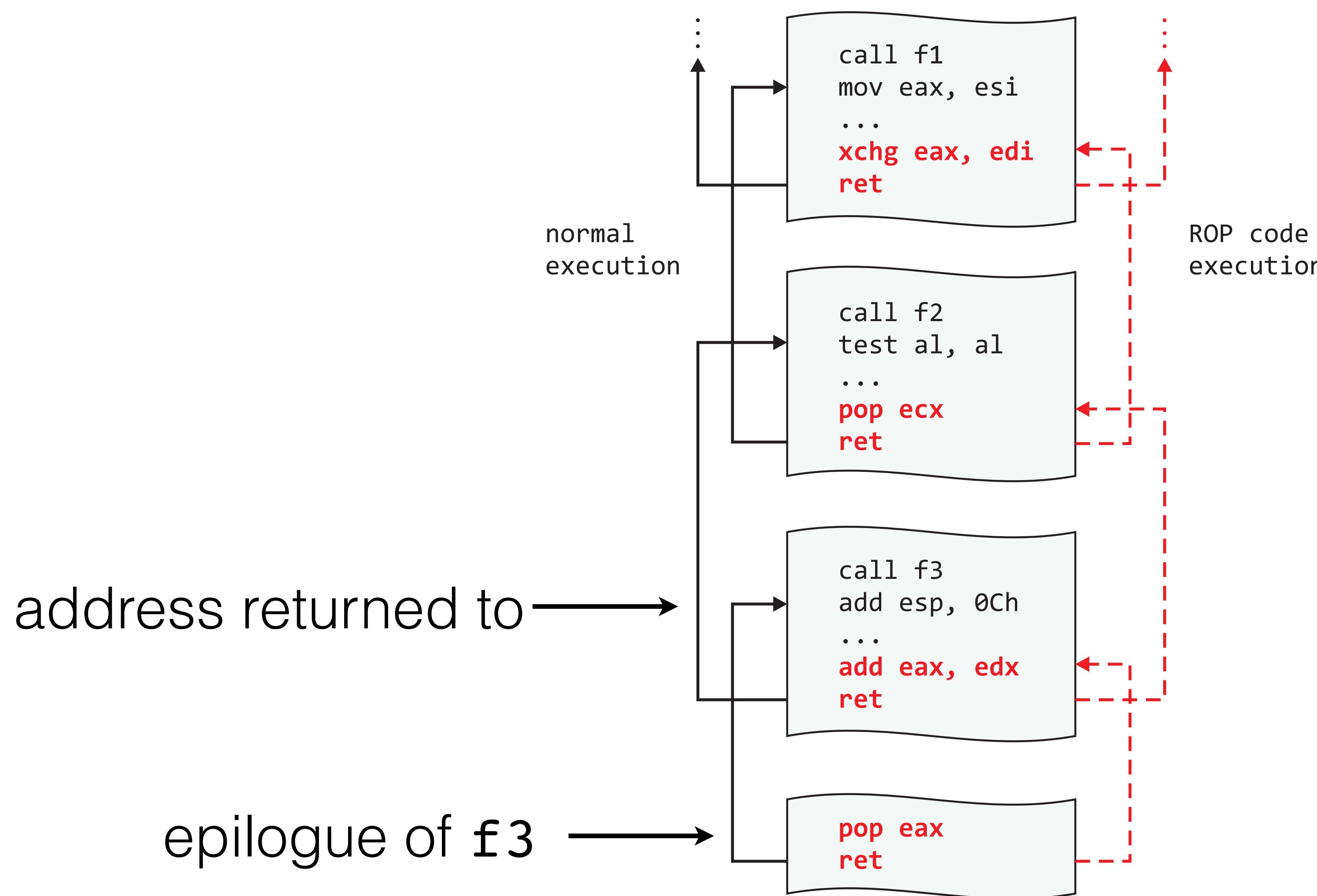
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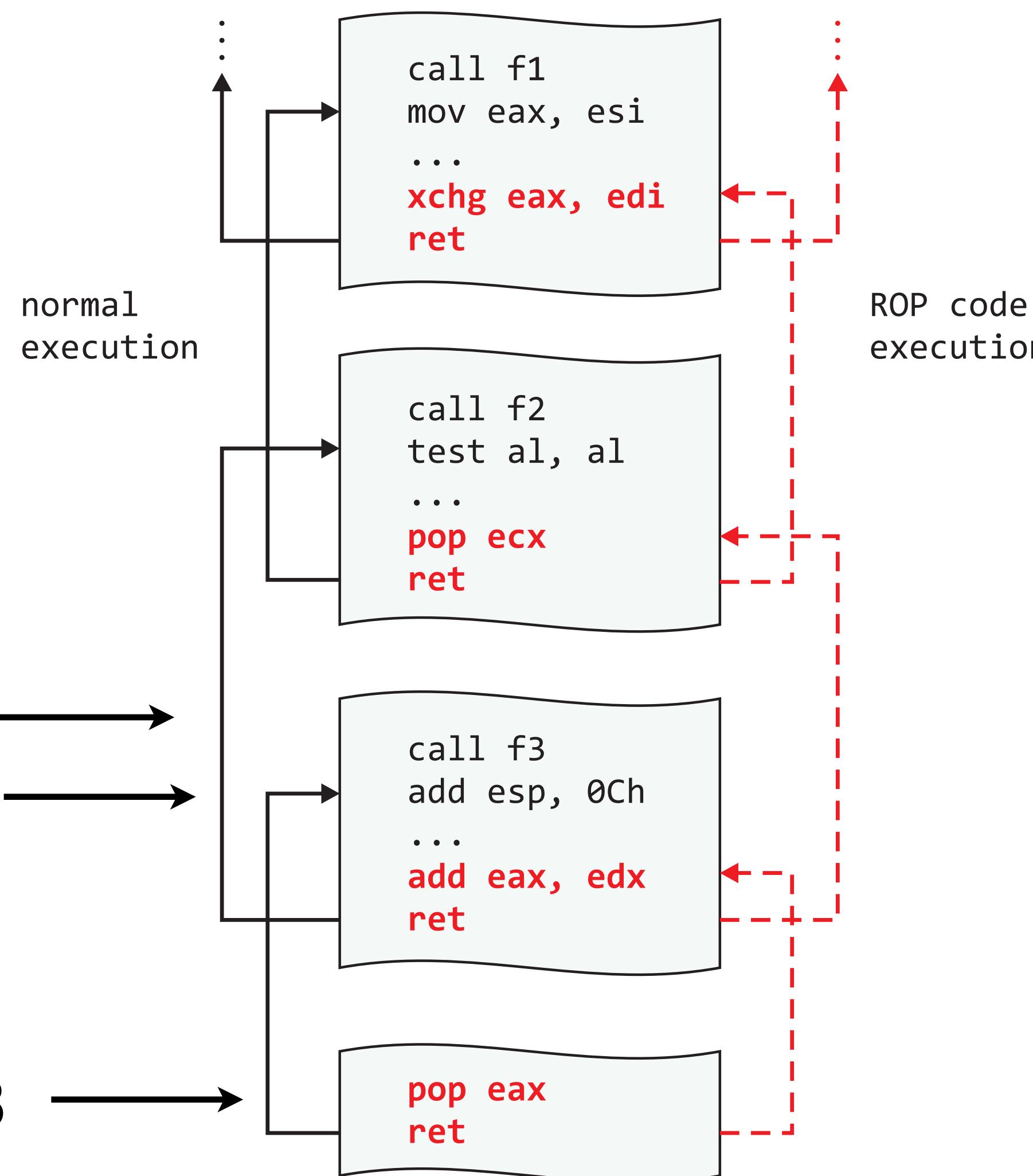


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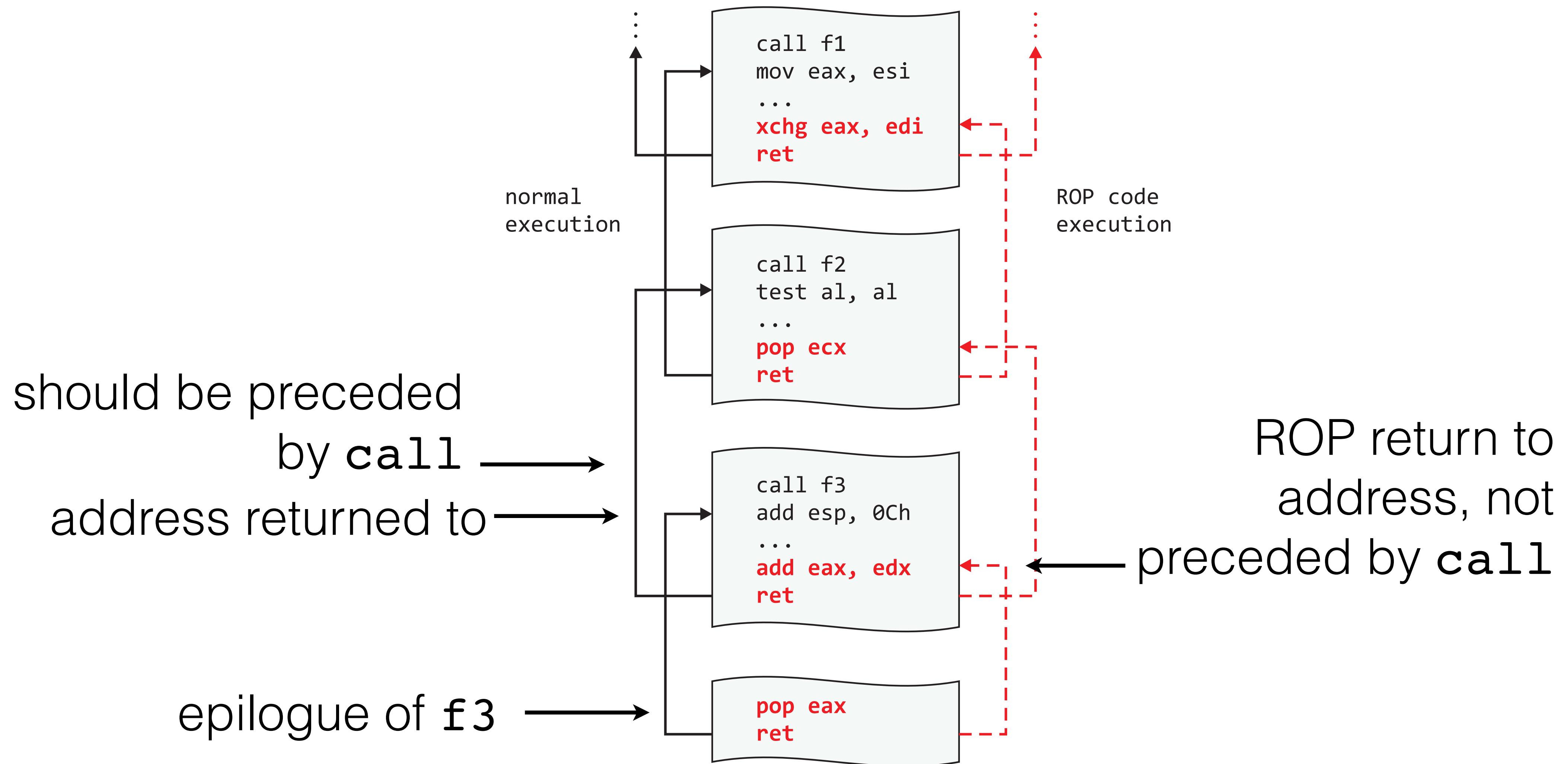


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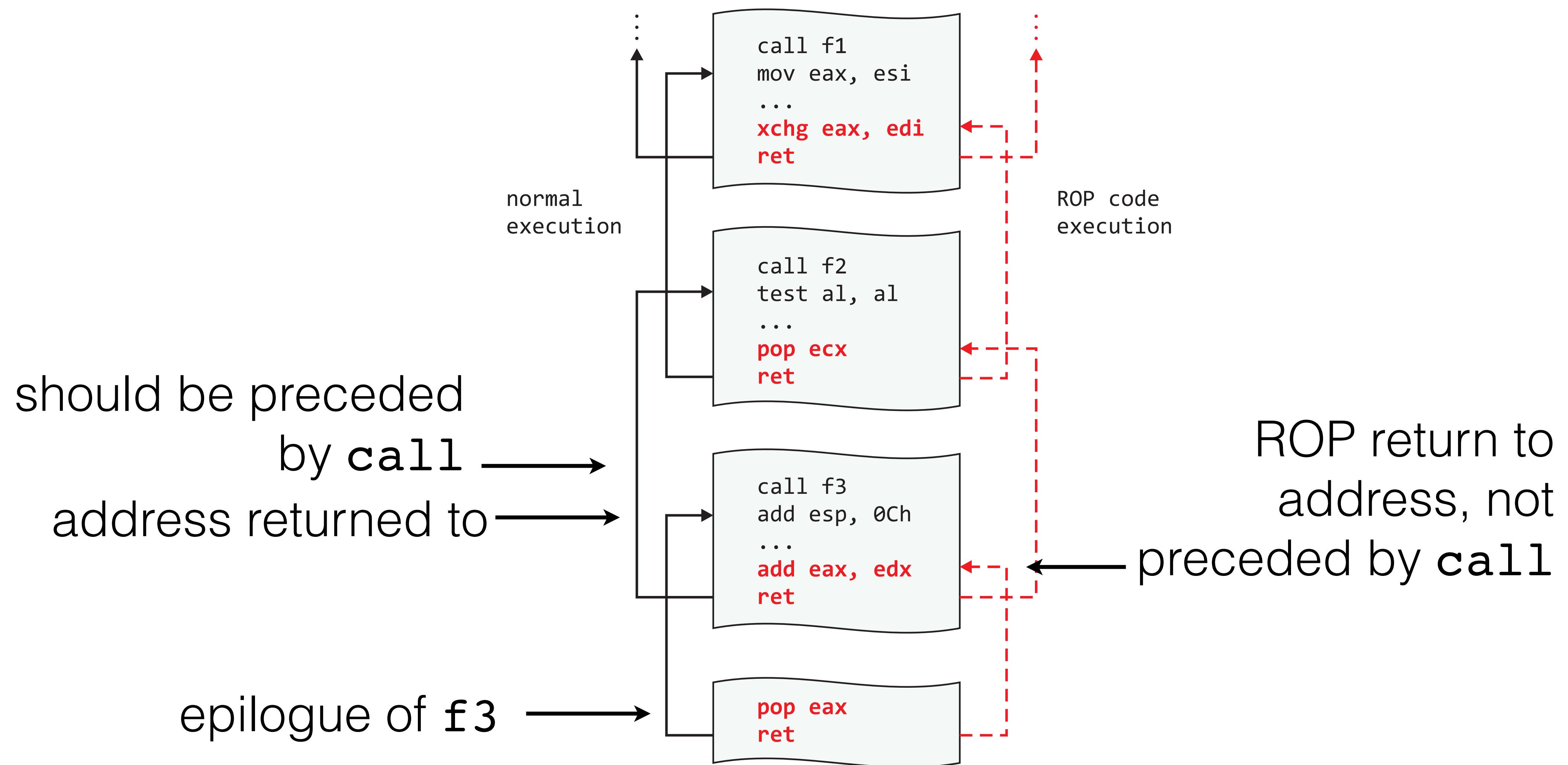
should be preceded
by **call** →
address returned to →
epilogue of f3 →



ROP detection: illegal returns



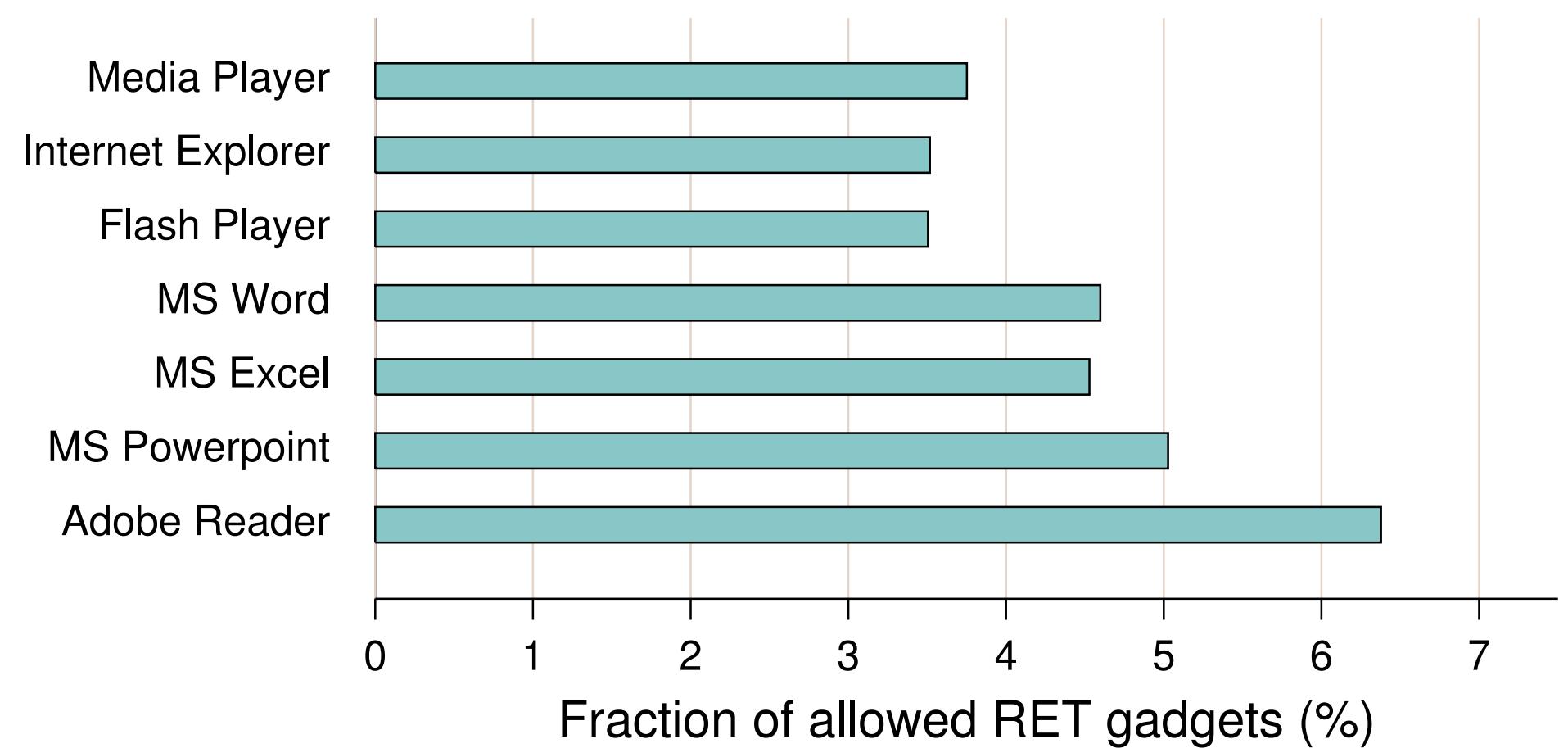
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Q: what about gadgets that begin at `call`?

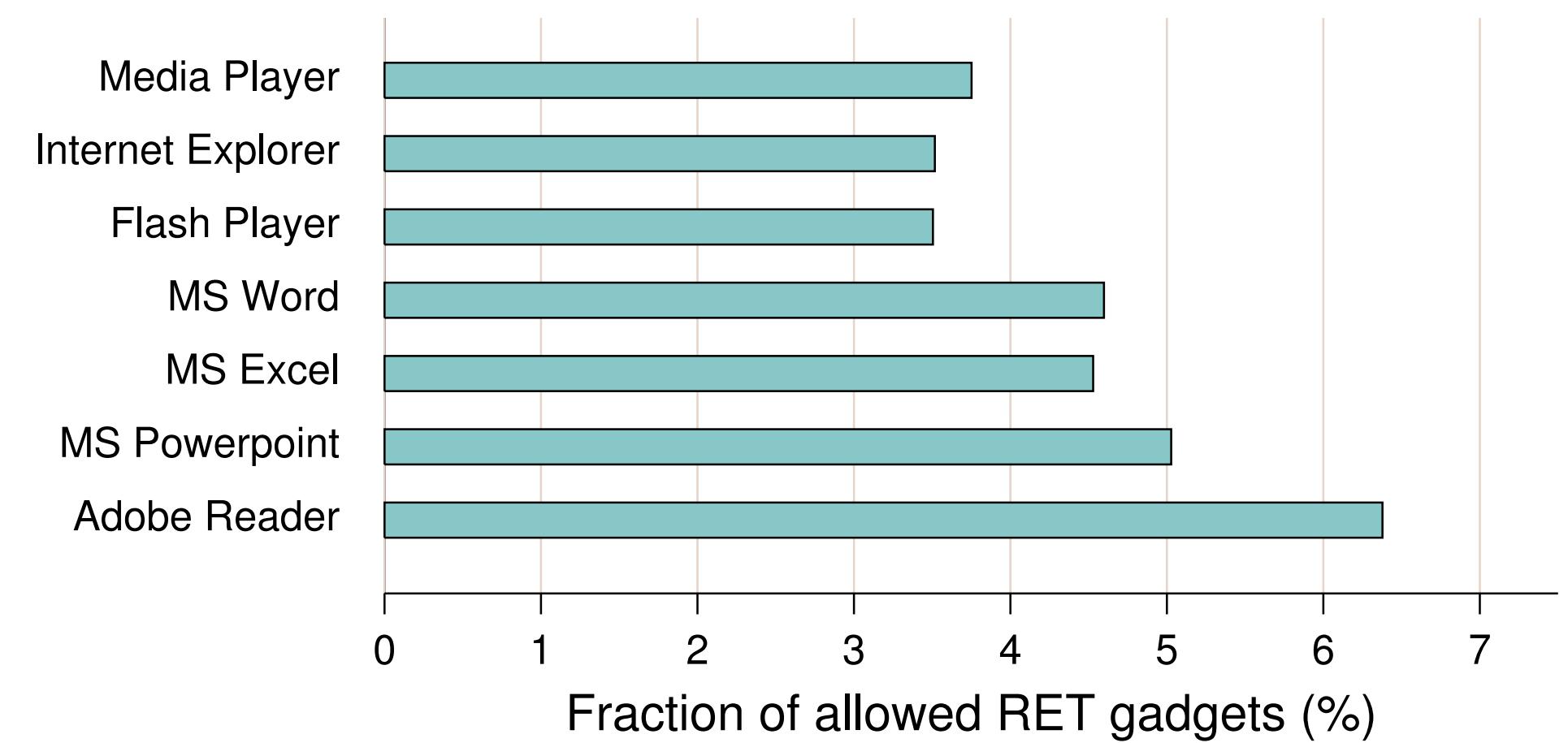
prevalence of **call**-preceded gadgets:

1. **call**
2. up-to 20 instructions
3. **ret**



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- ↑
1. thought to be too few
2. non-sense instructions should affect ROP
program state

ROP detection: gadget chaining

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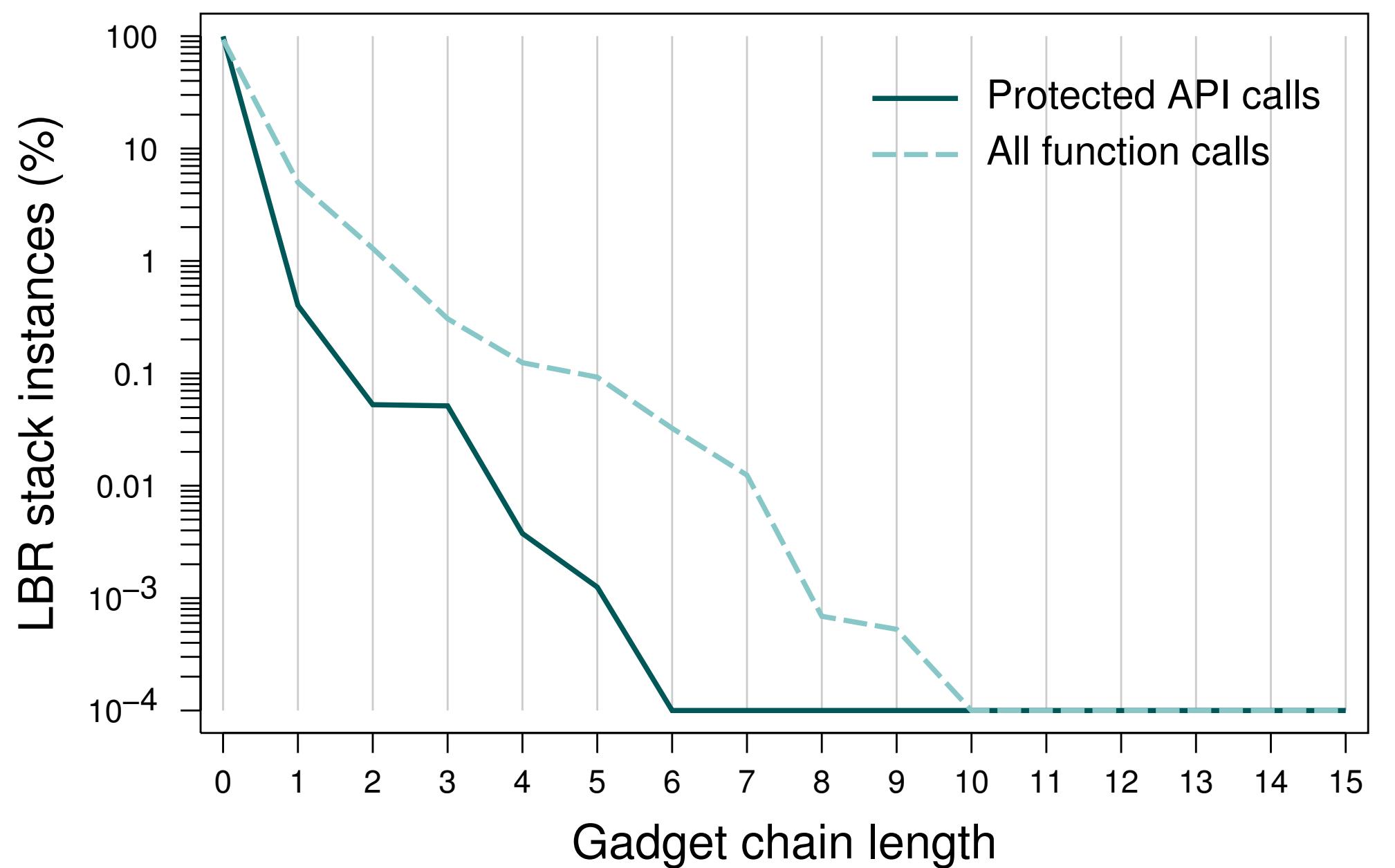


sequence of indirect calls, with fewer than 20
instructions between call and return

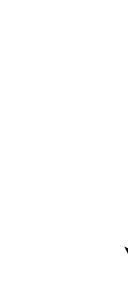
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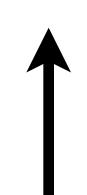
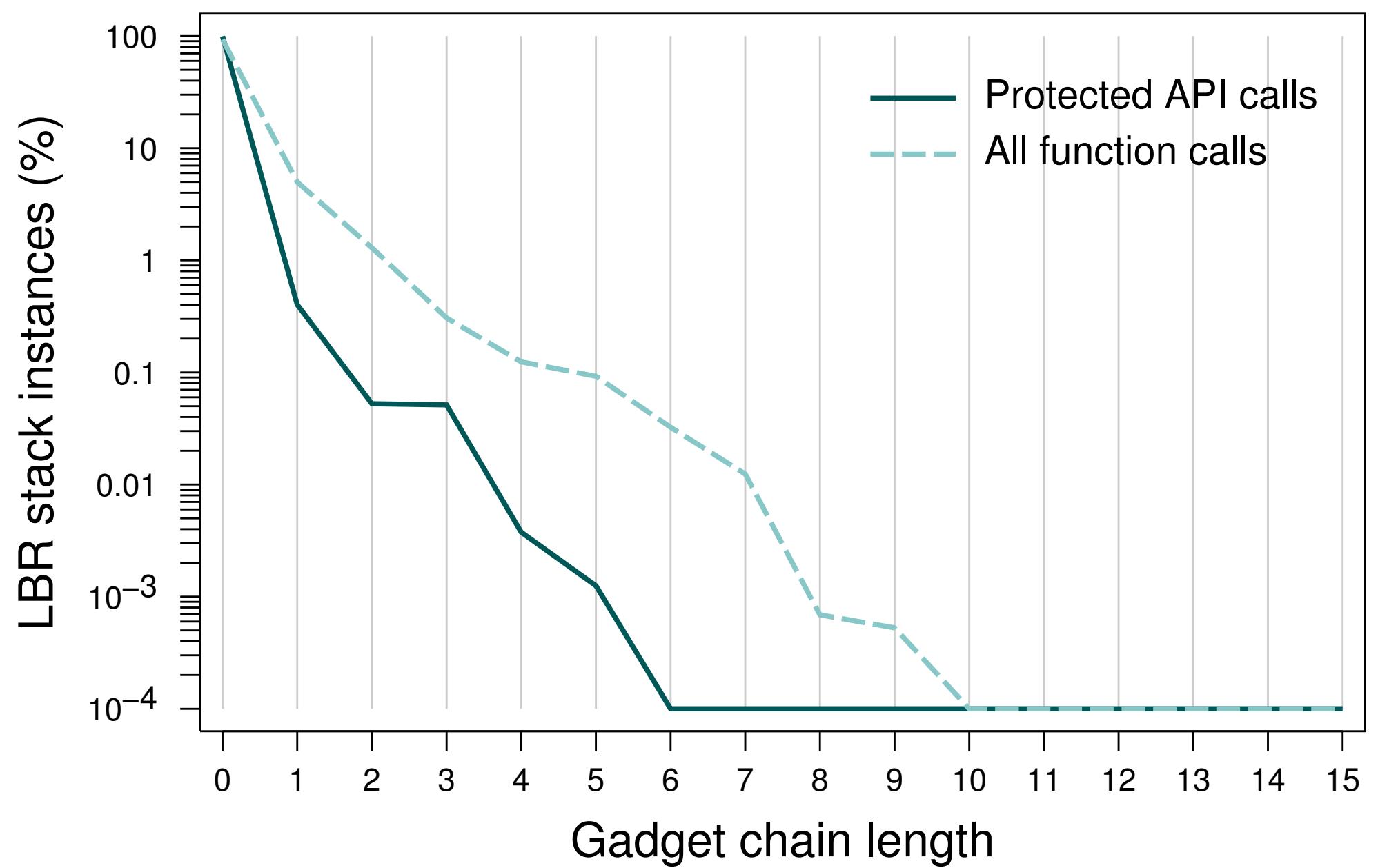
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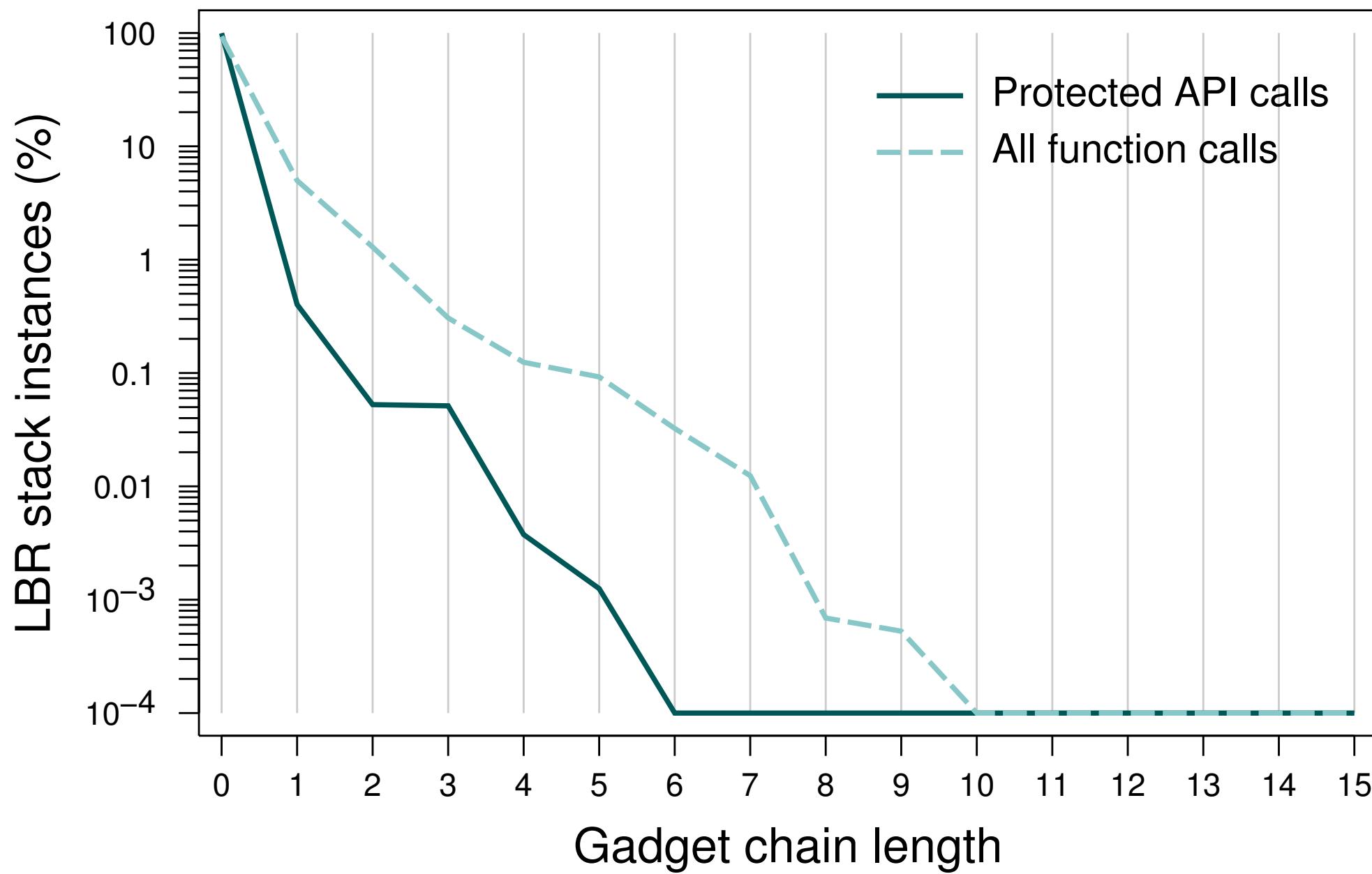


most APIs won't
exhaust LBR buffer

ROP detection: gadget chaining



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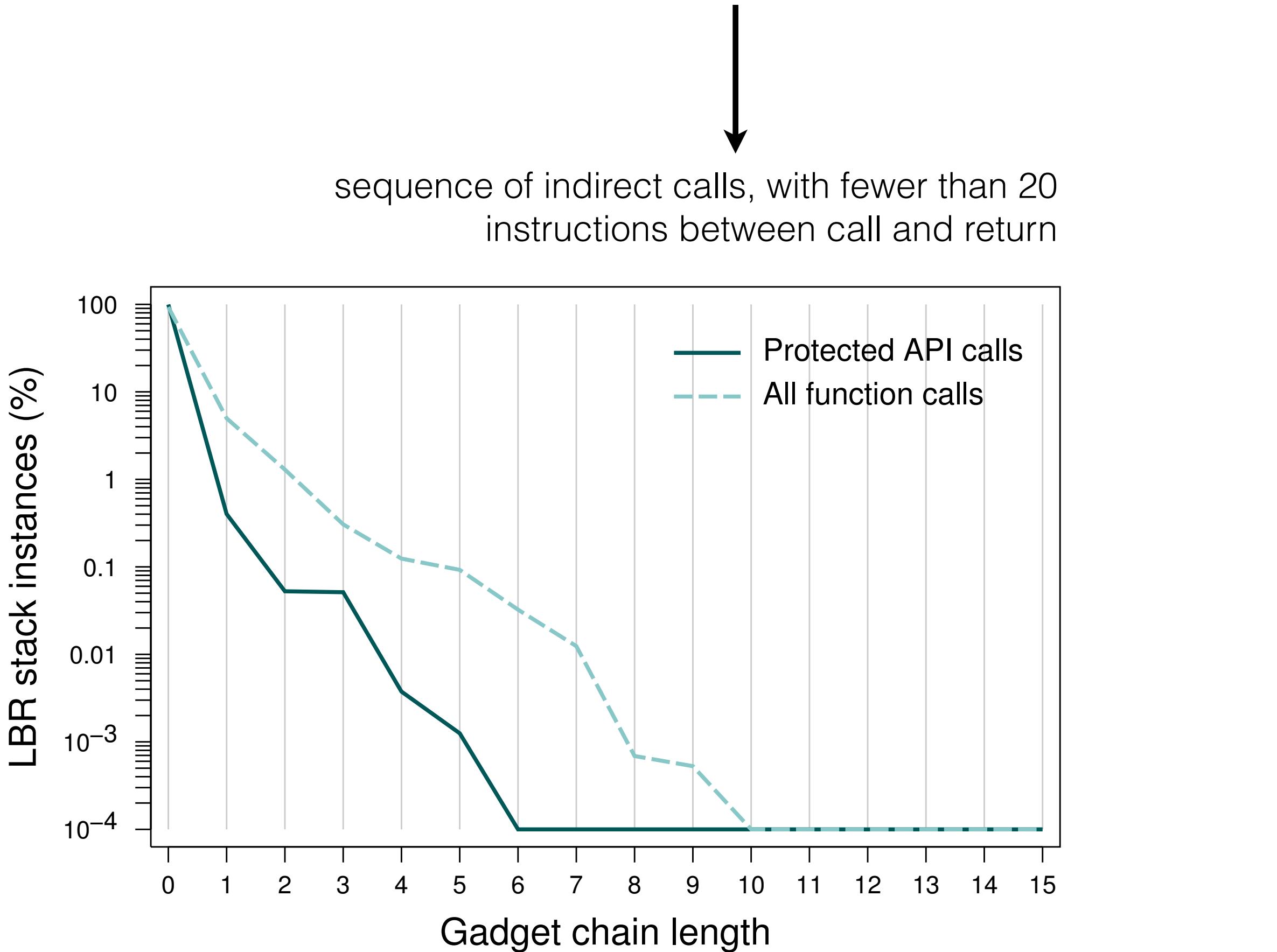
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but, gadgets

require many more
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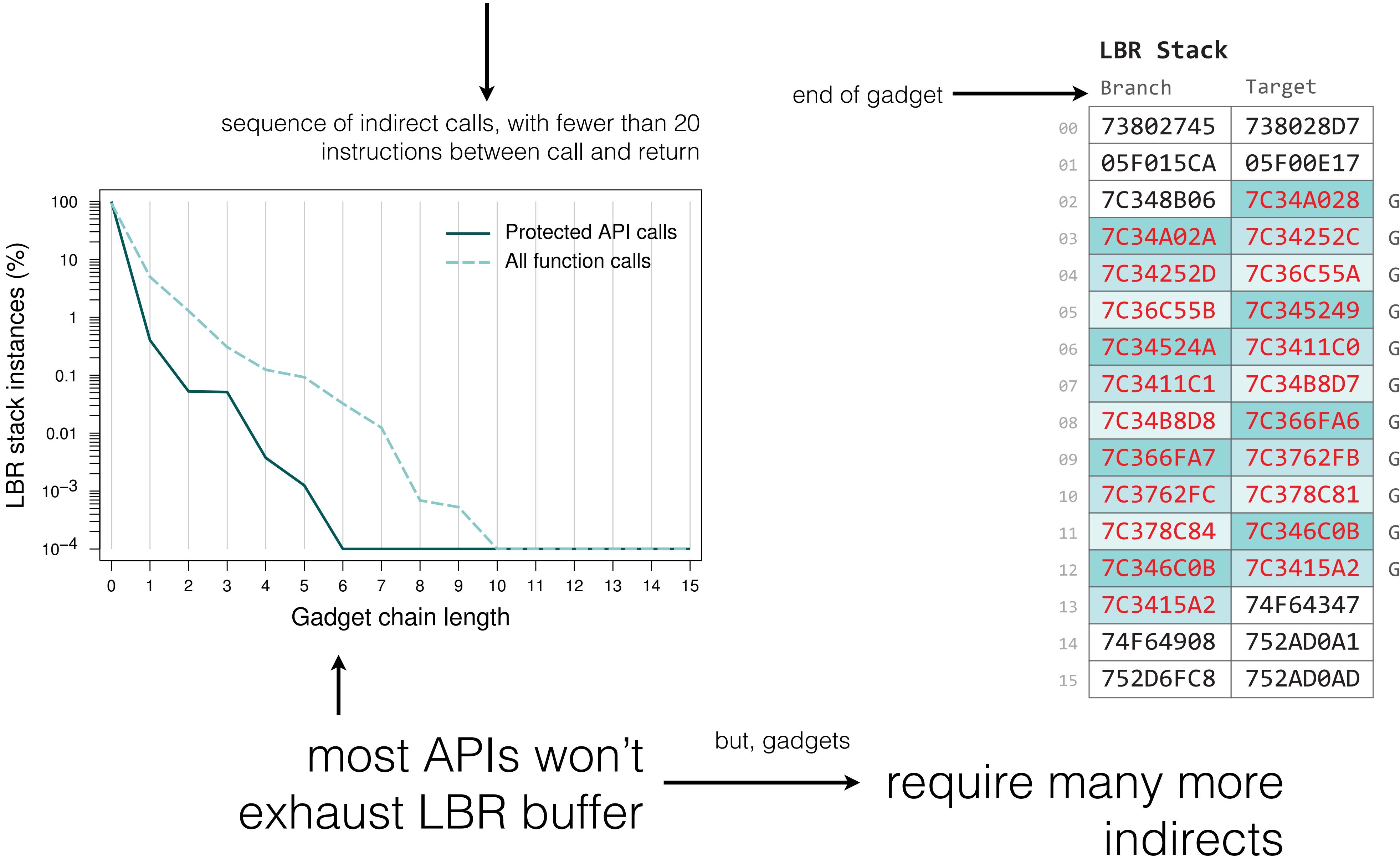
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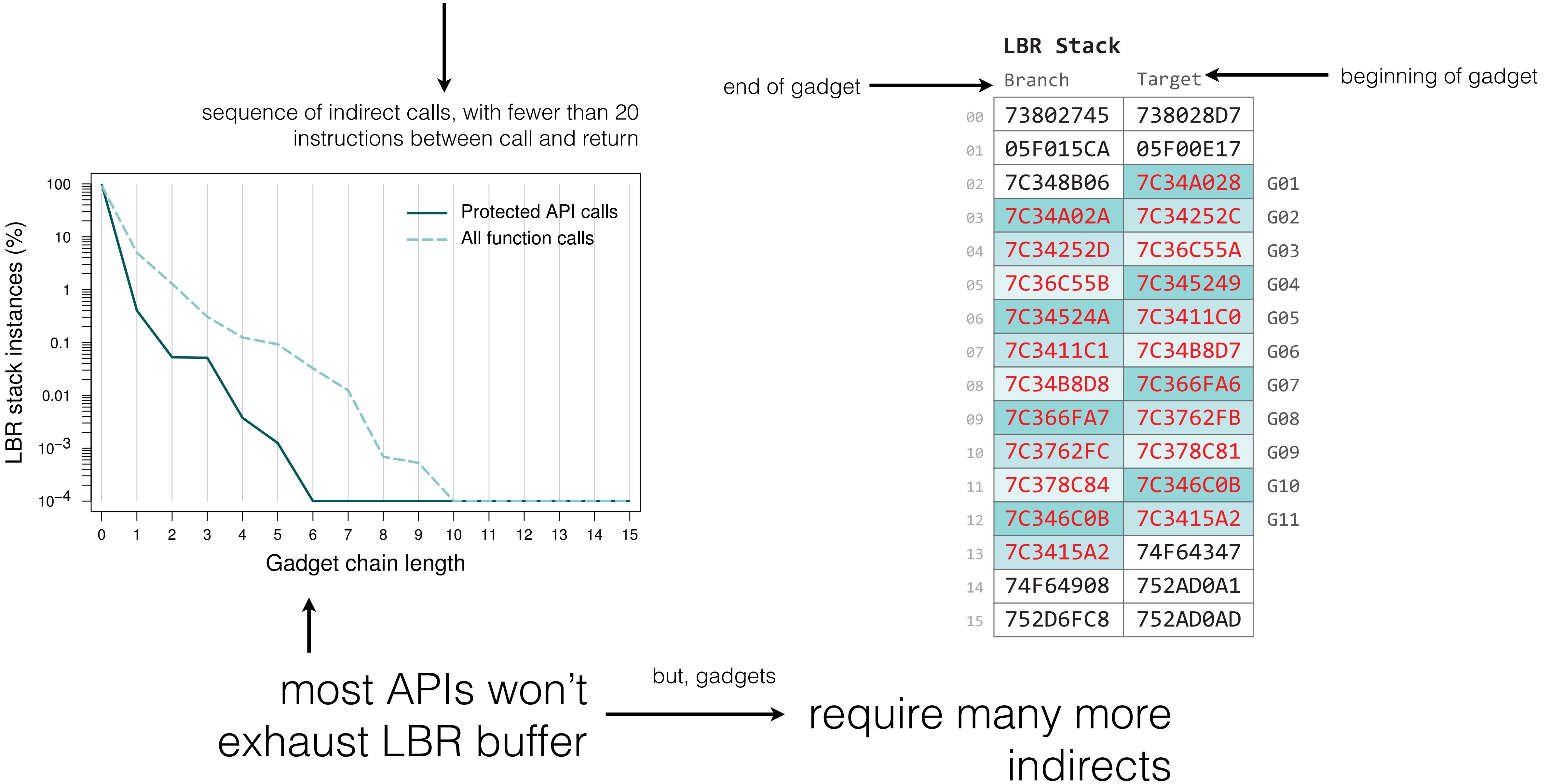
LBR Stack

Branch	Target
00	73802745 738028D7
01	05F015CA 05F00E17
02	7C348B06 7C34A028
03	7C34A02A 7C34252C
04	7C34252D 7C36C55A
05	7C36C55B 7C345249
06	7C34524A 7C3411C0
07	7C3411C1 7C34B8D7
08	7C34B8D8 7C366FA6
09	7C366FA7 7C3762FB
10	7C3762FC 7C378C81
11	7C378C84 7C346C0B
12	7C346C0B 7C3415A2
13	7C3415A2 74F64347
14	74F64908 752AD0A1
15	752D6FC8 752AD0AD

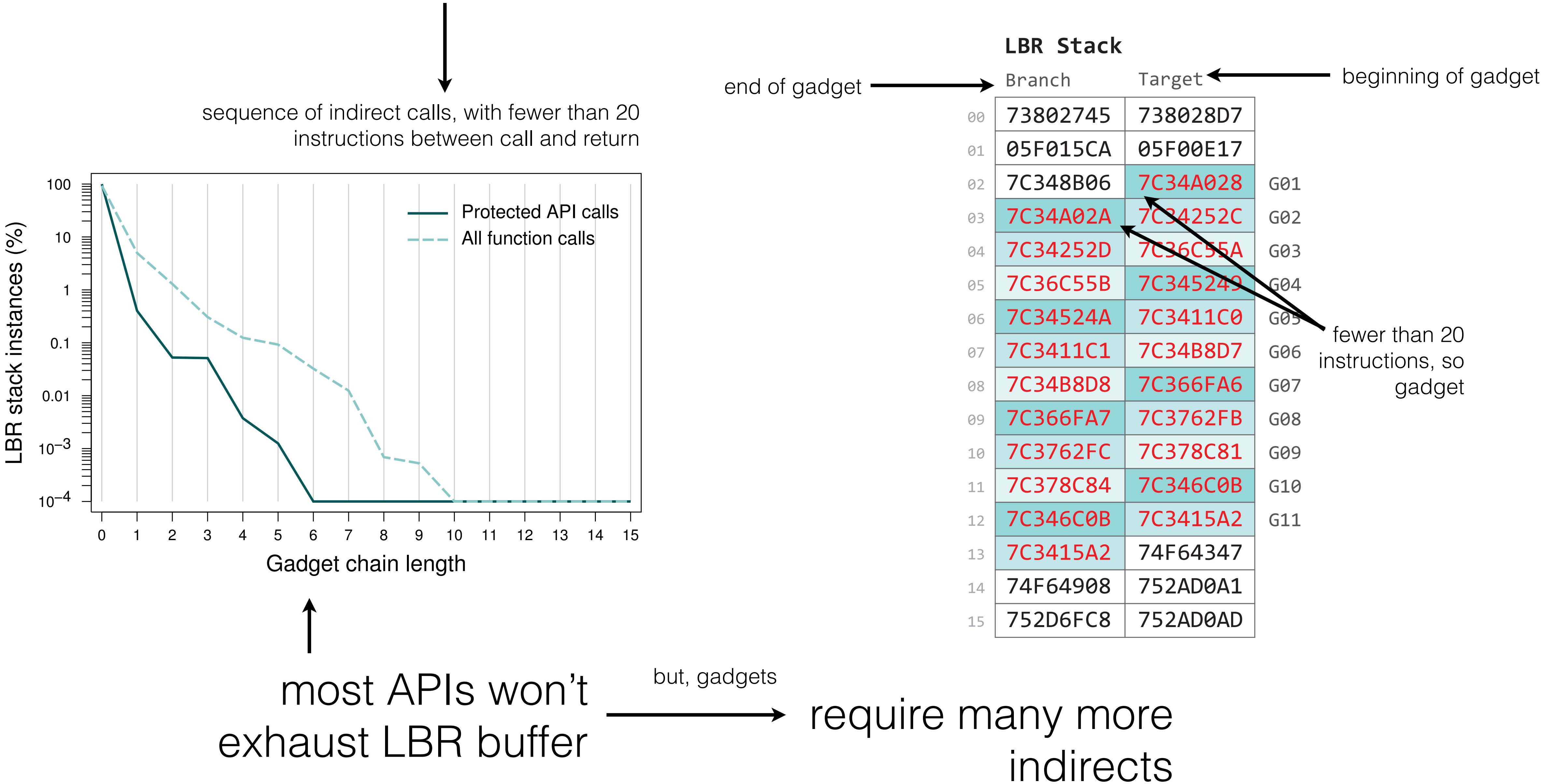
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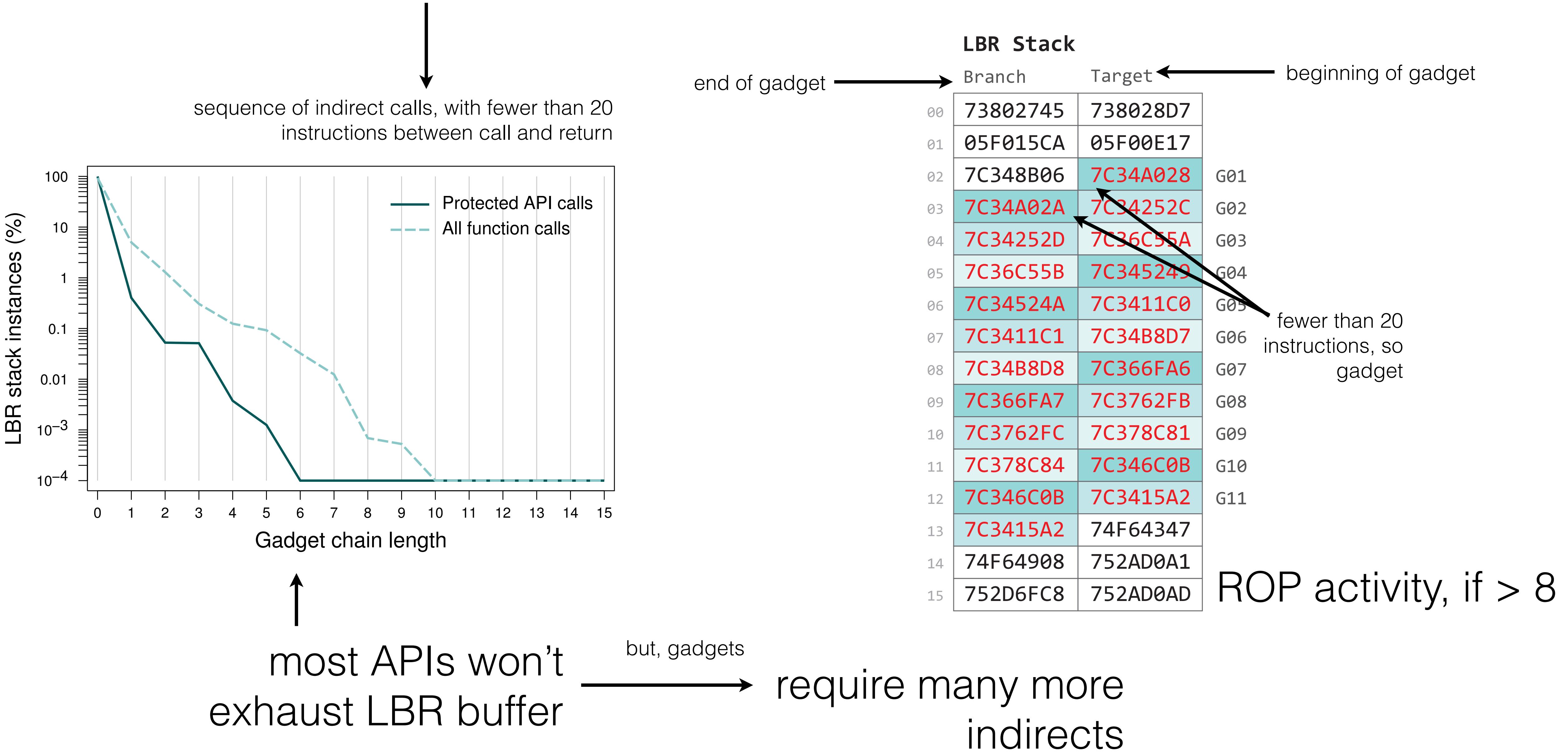
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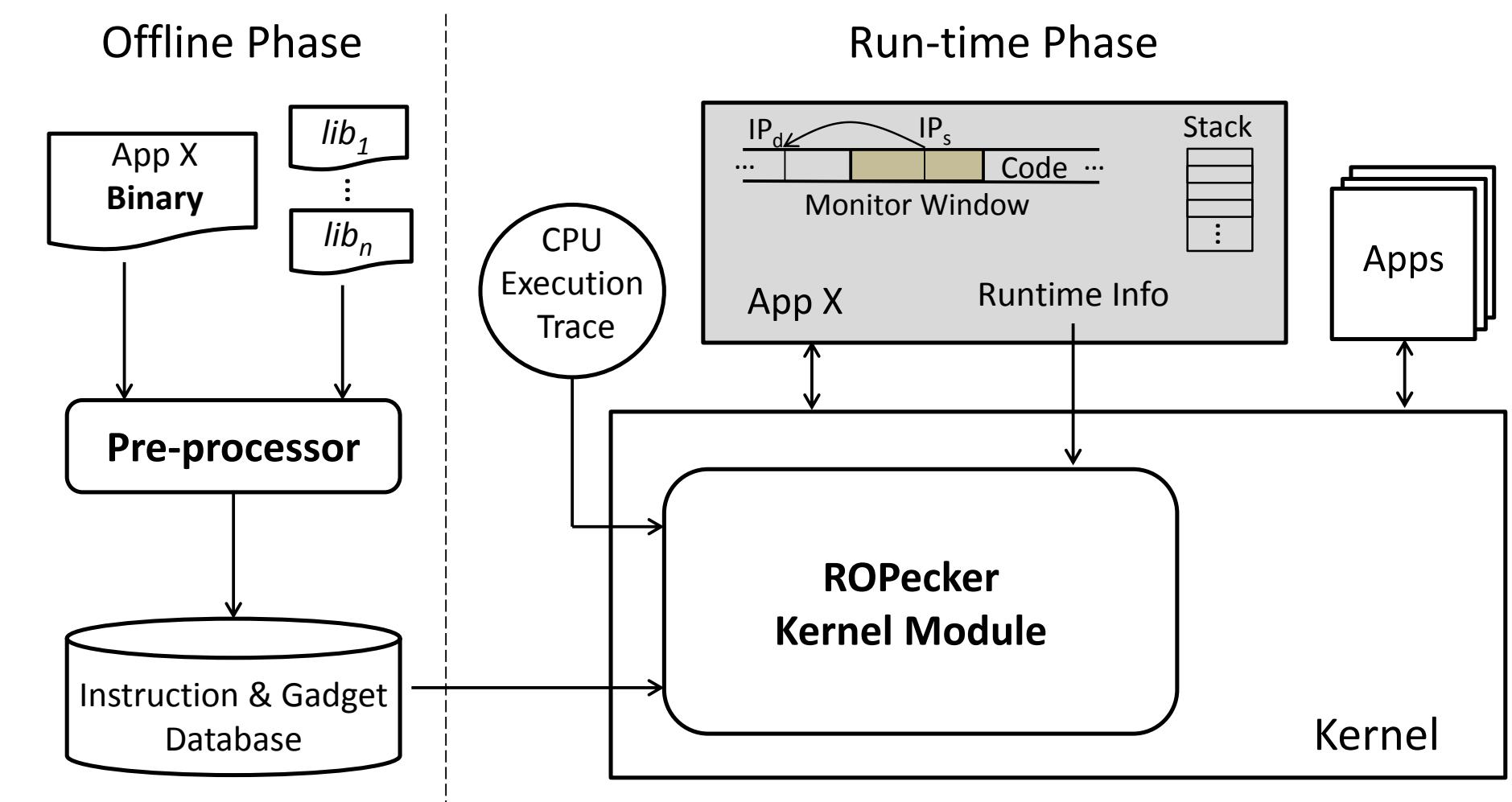


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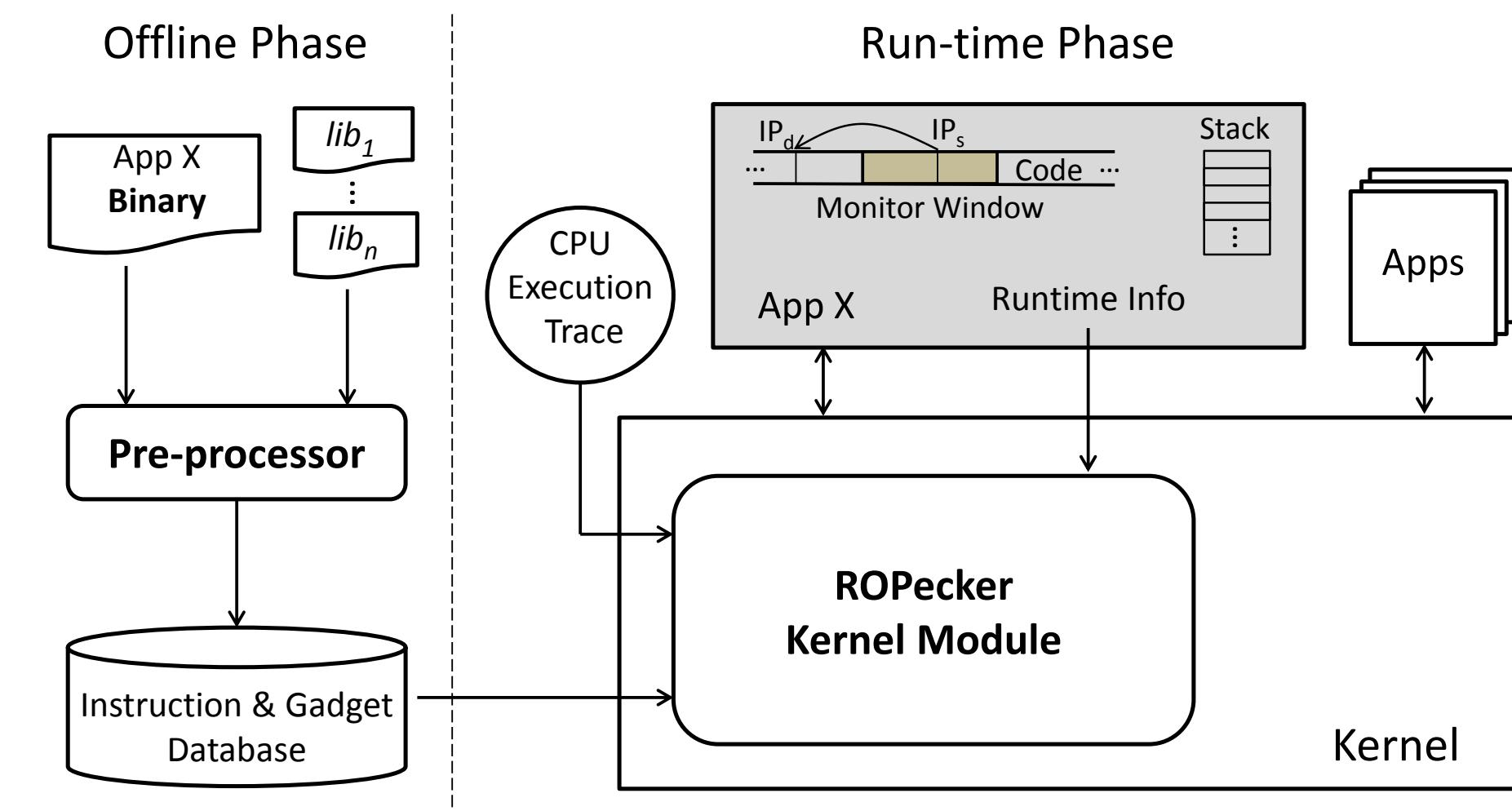
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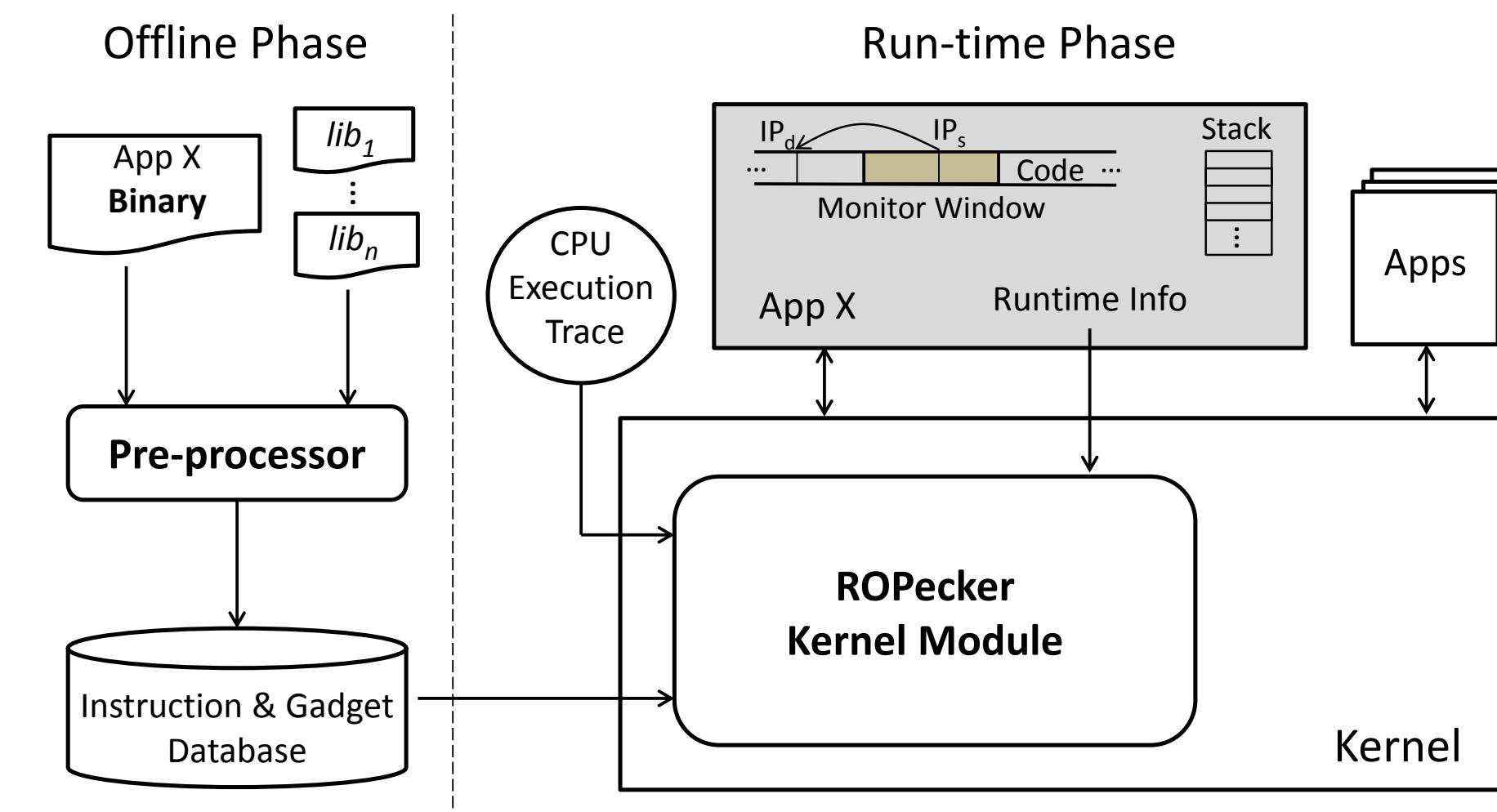
1. uses LBR, as kBouncer; past ROP activity



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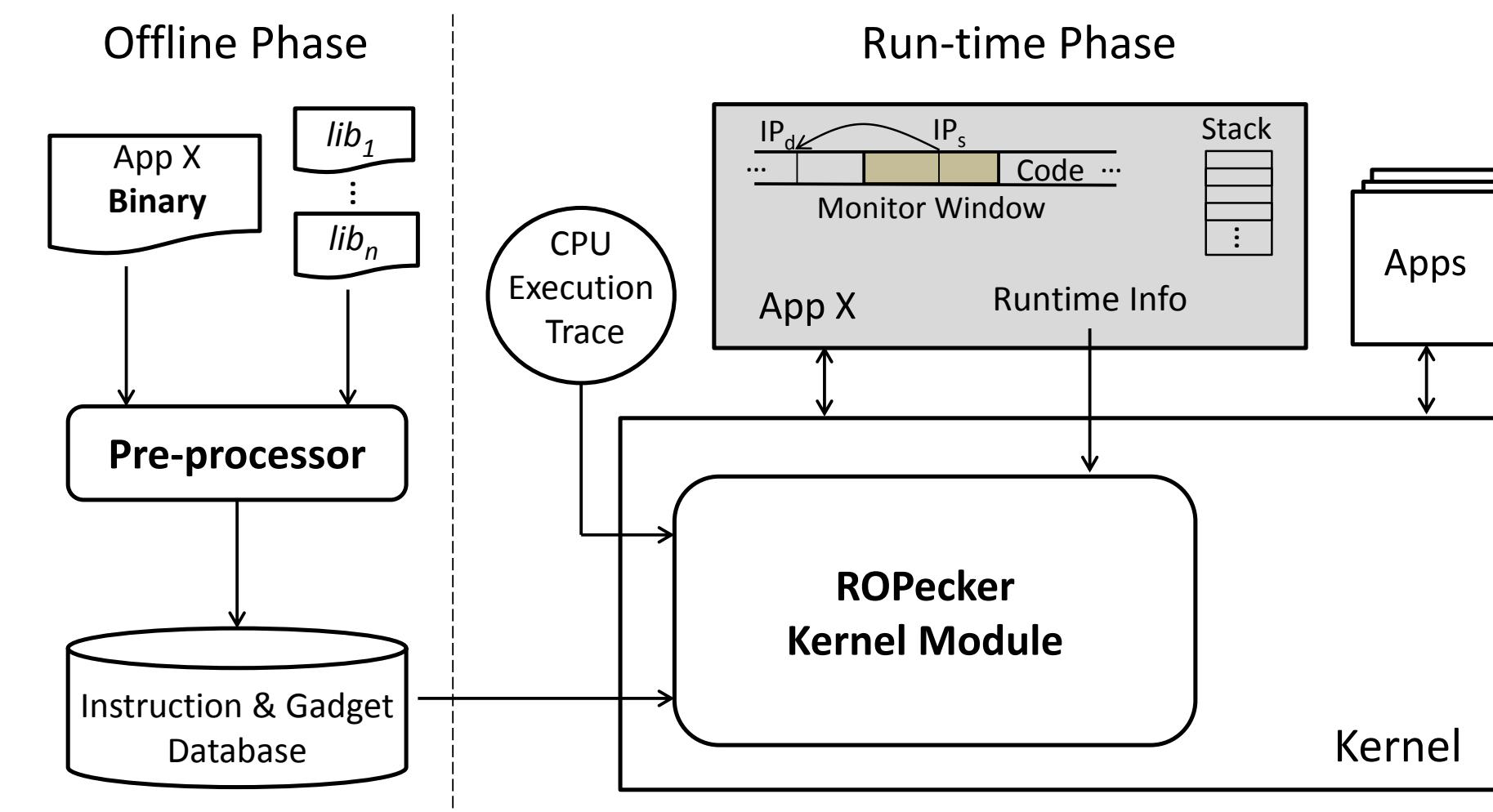
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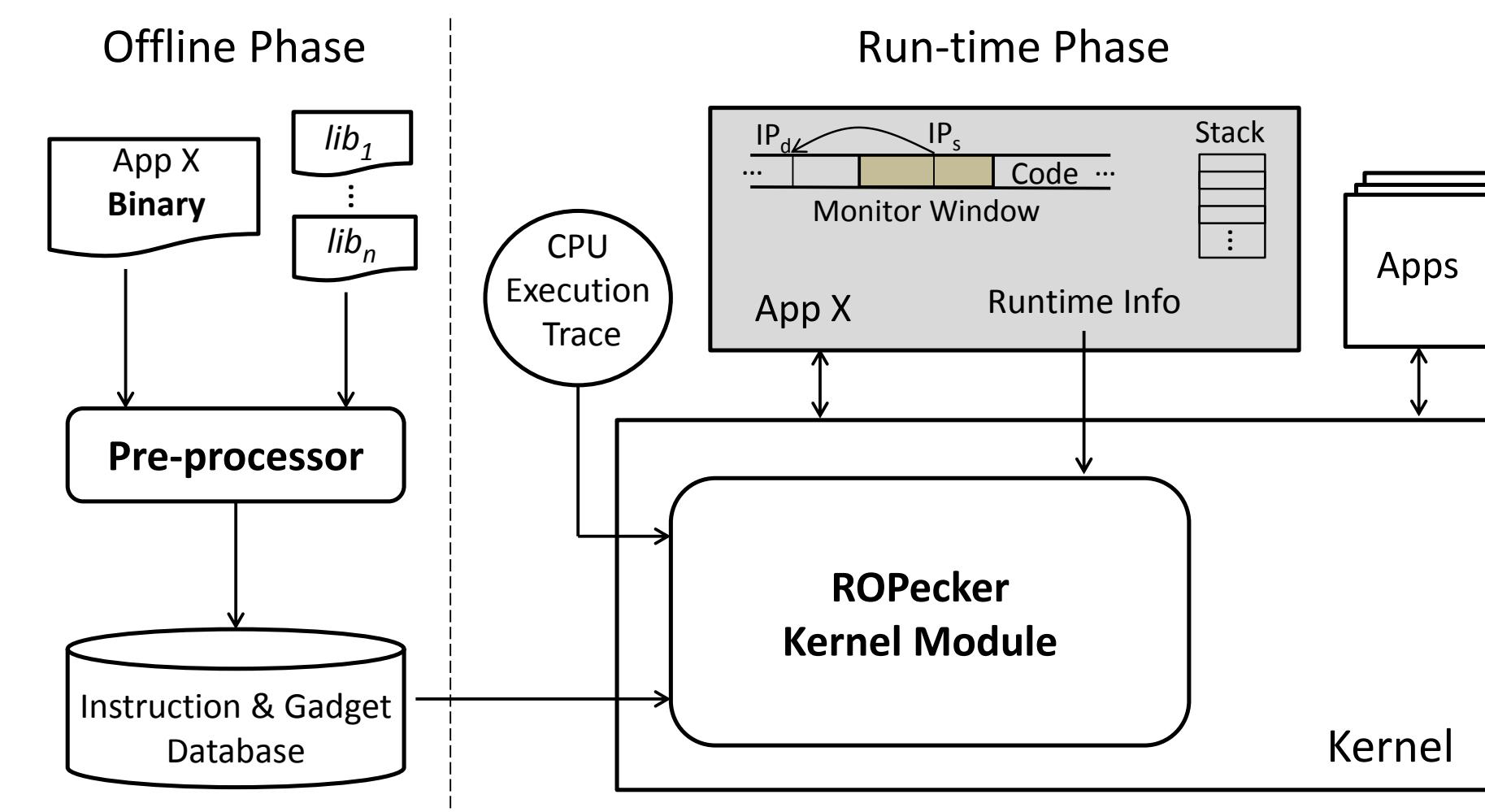
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3. decreases availability of gadgets



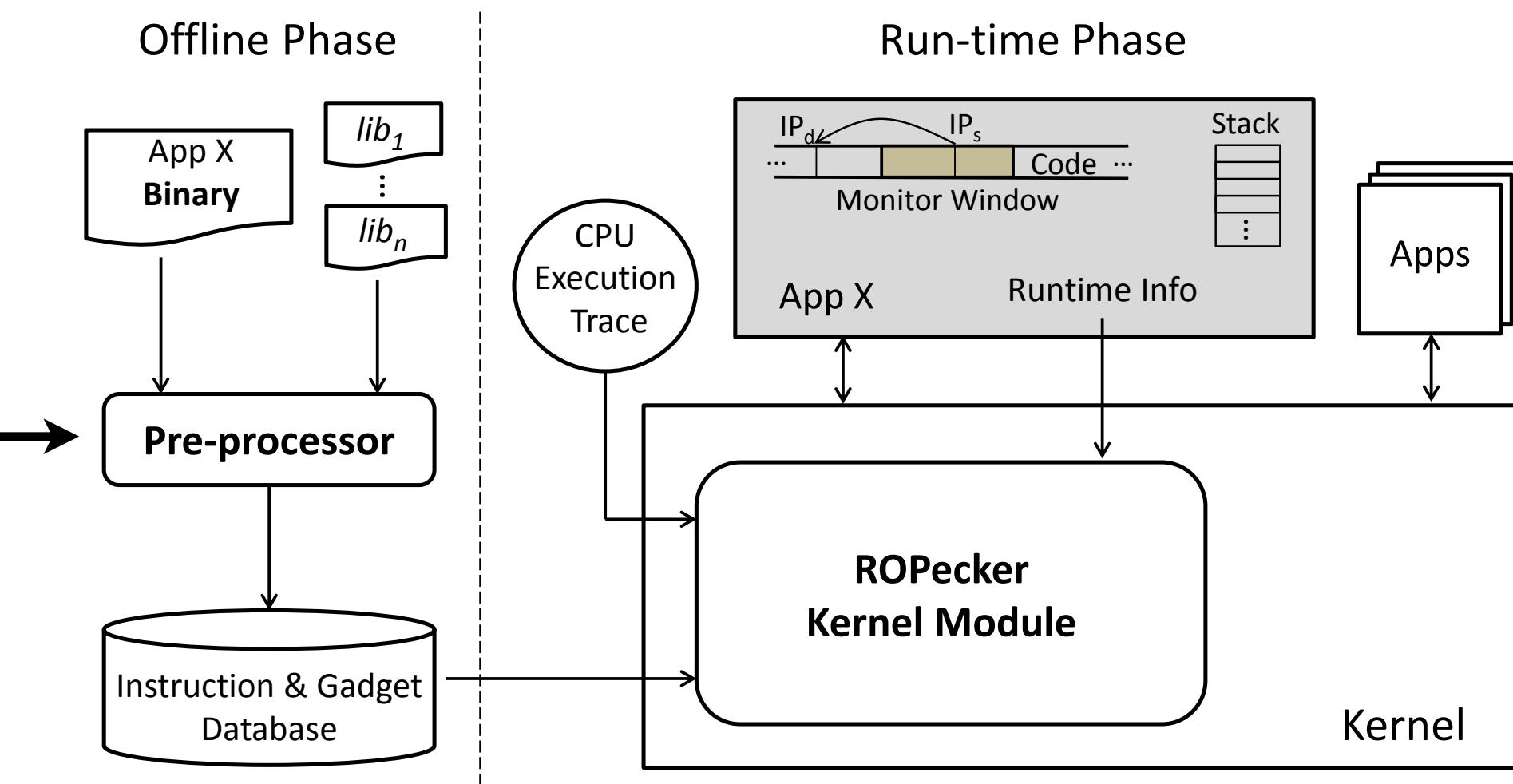
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1. uses LBR, as kBouncer; past ROP activity
2. emulates instructions; future ROP activity
3. decreases availability of gadgets
4. identifies potential gadgets



identify potential gadgets (six instructions)
via disassembly:
1. application
2. shared library

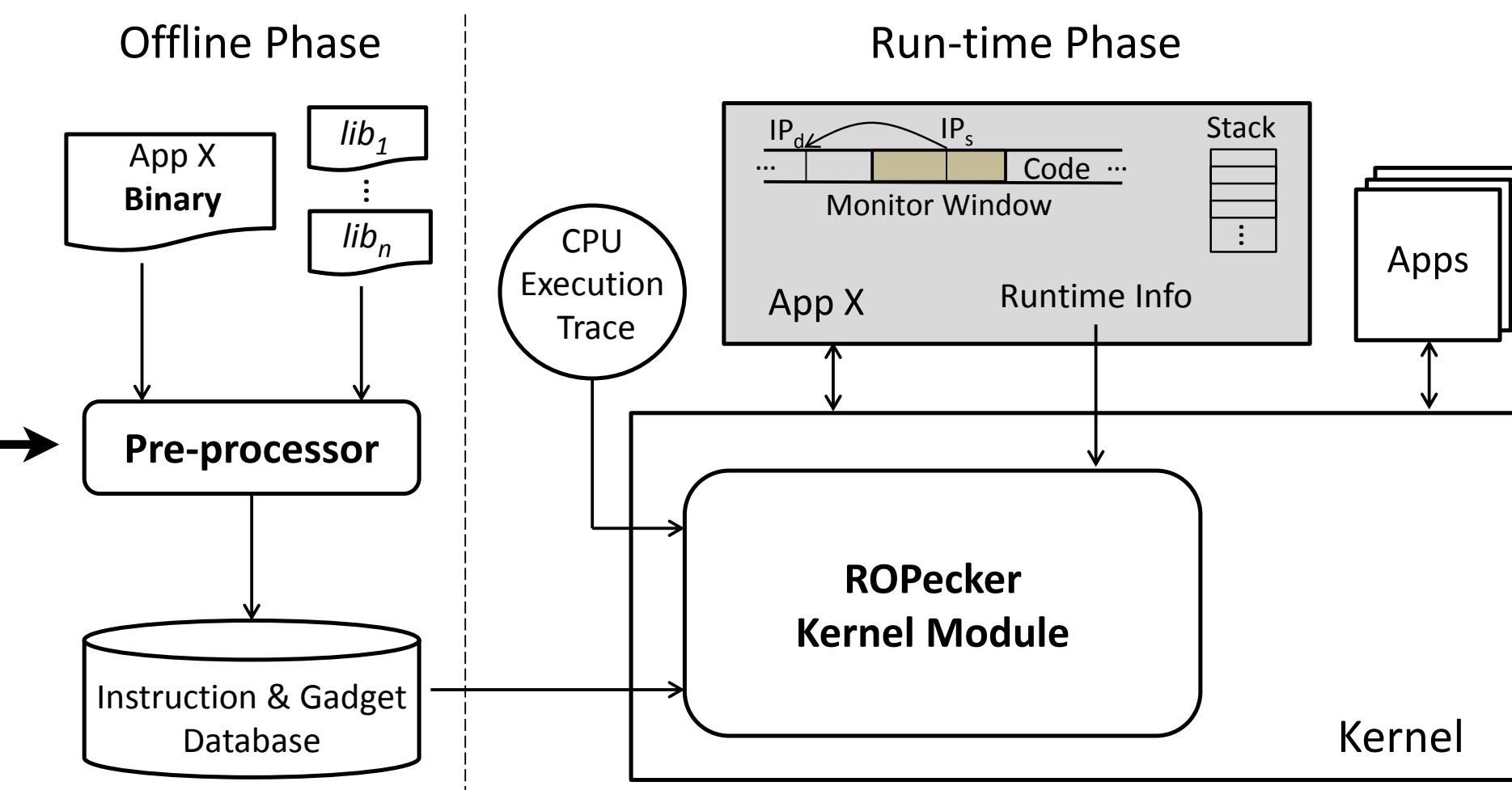


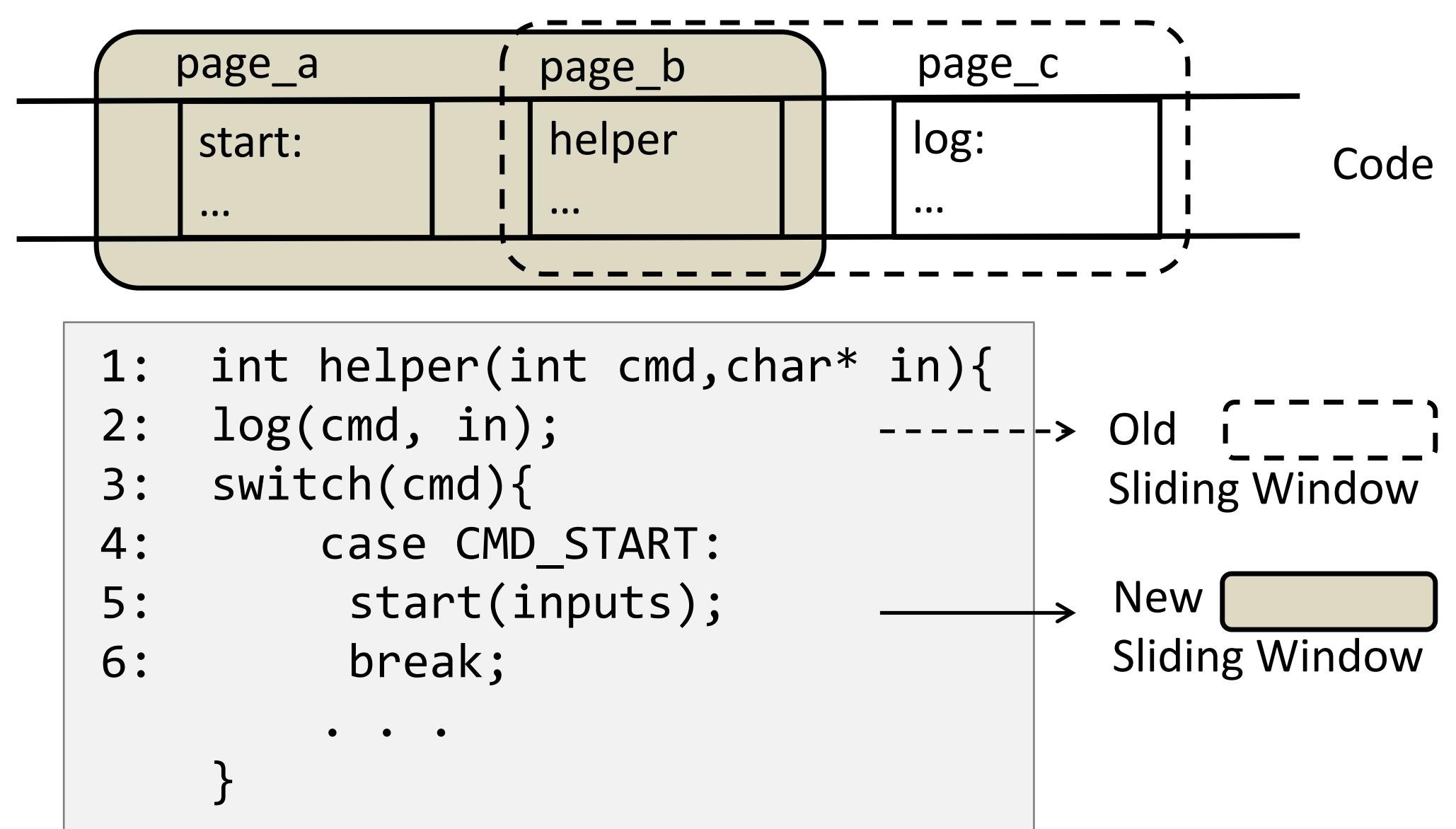
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if addresses of gadget
used, we assume ROP
activity

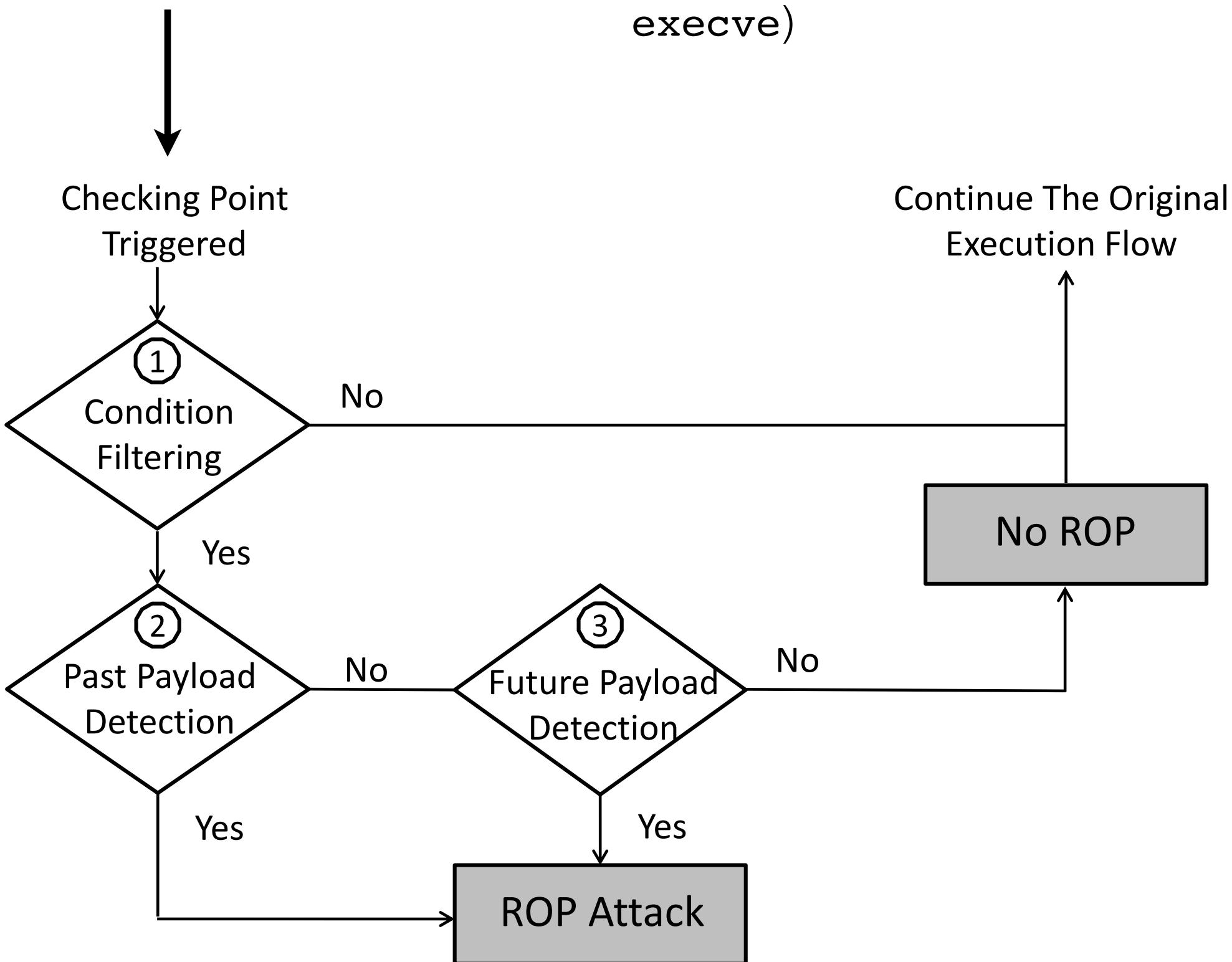




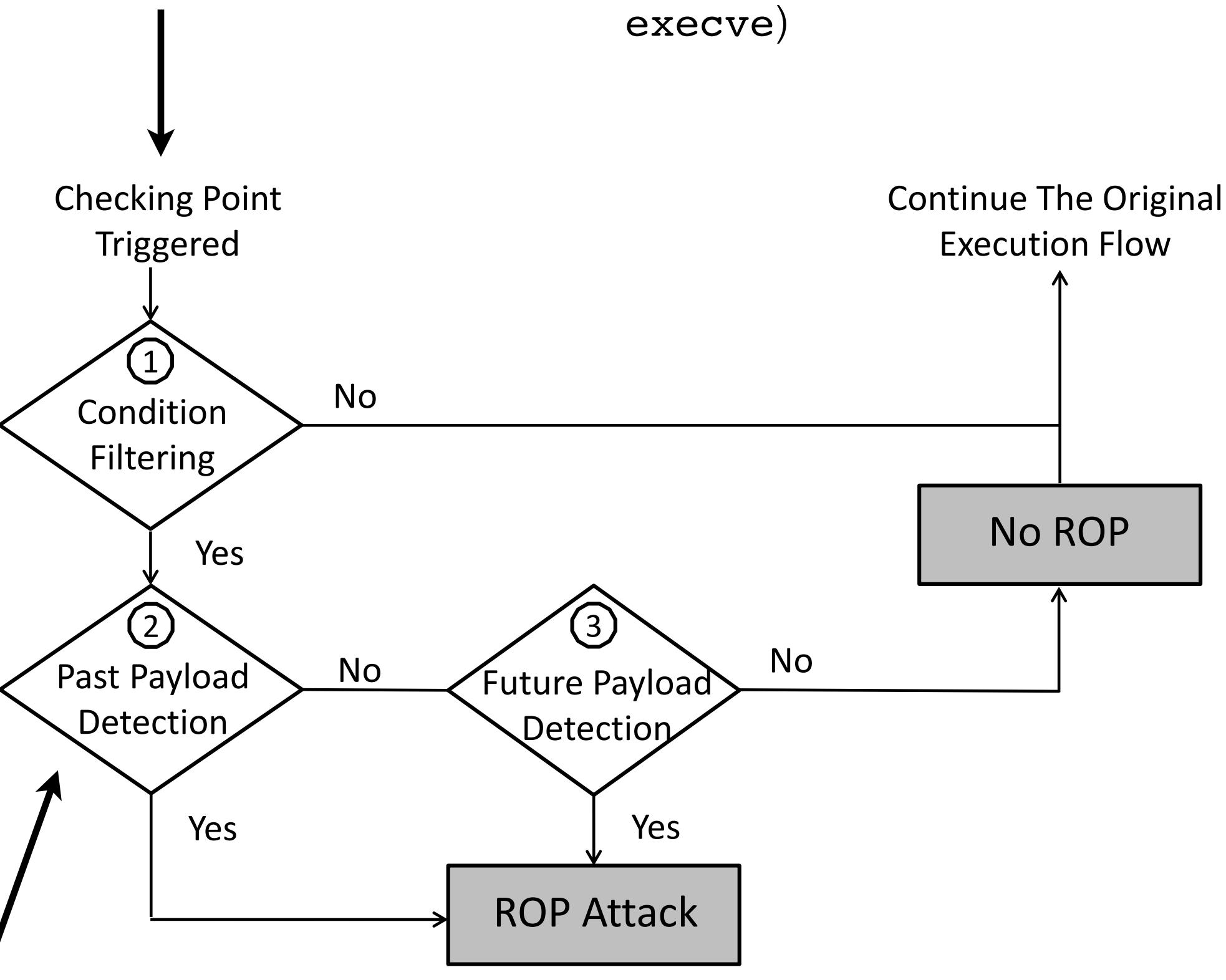
sliding window:

1. all pages set NX
2. for each call:
 - a. enable X for next N bytes/pages
 - b. disable X for previous N bytes/pages
3. 8 KB or 16 KB

1. page fault from execution in NX region
2. call to protect system call (`mprotect`, `mmap2`, `execve`)

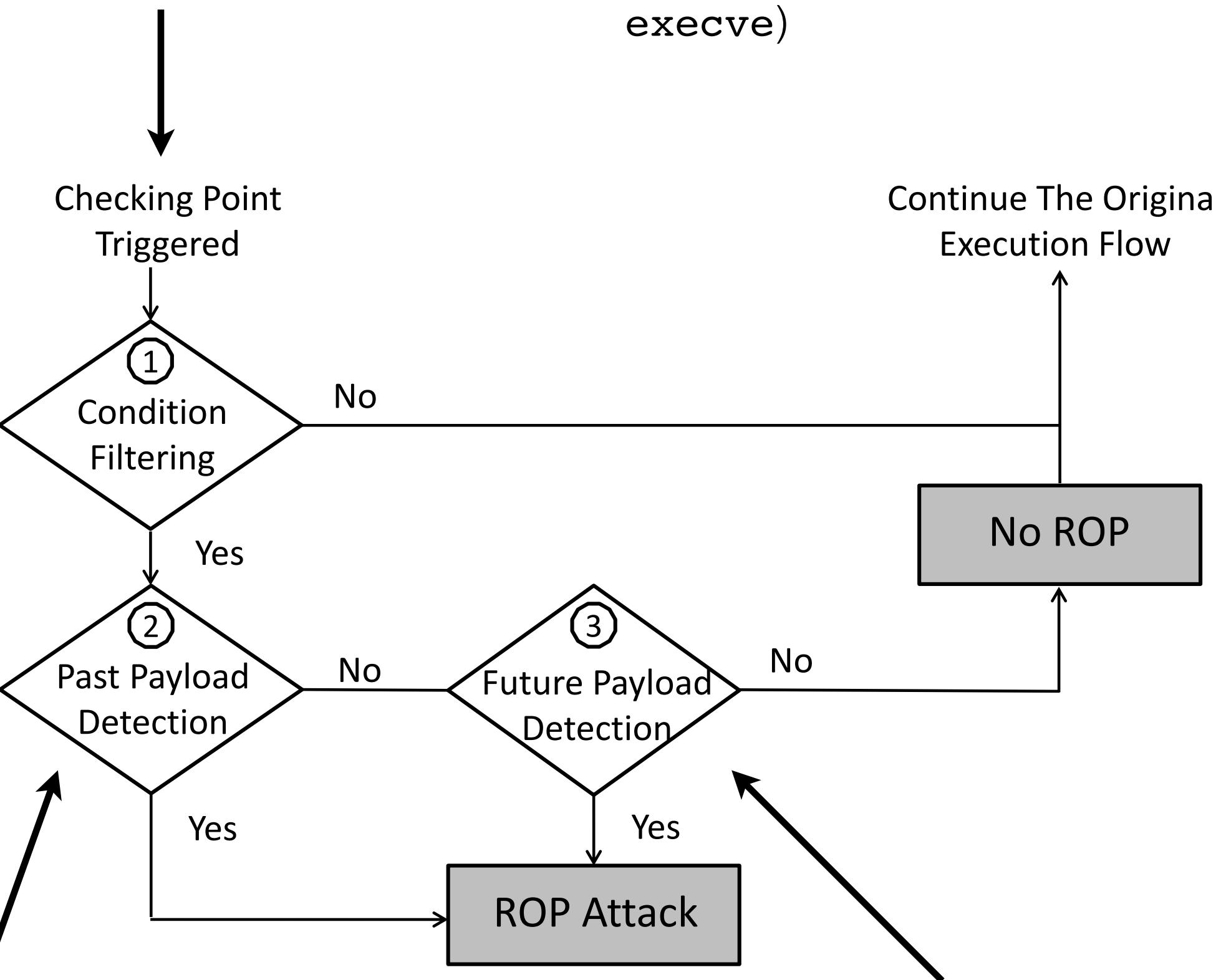


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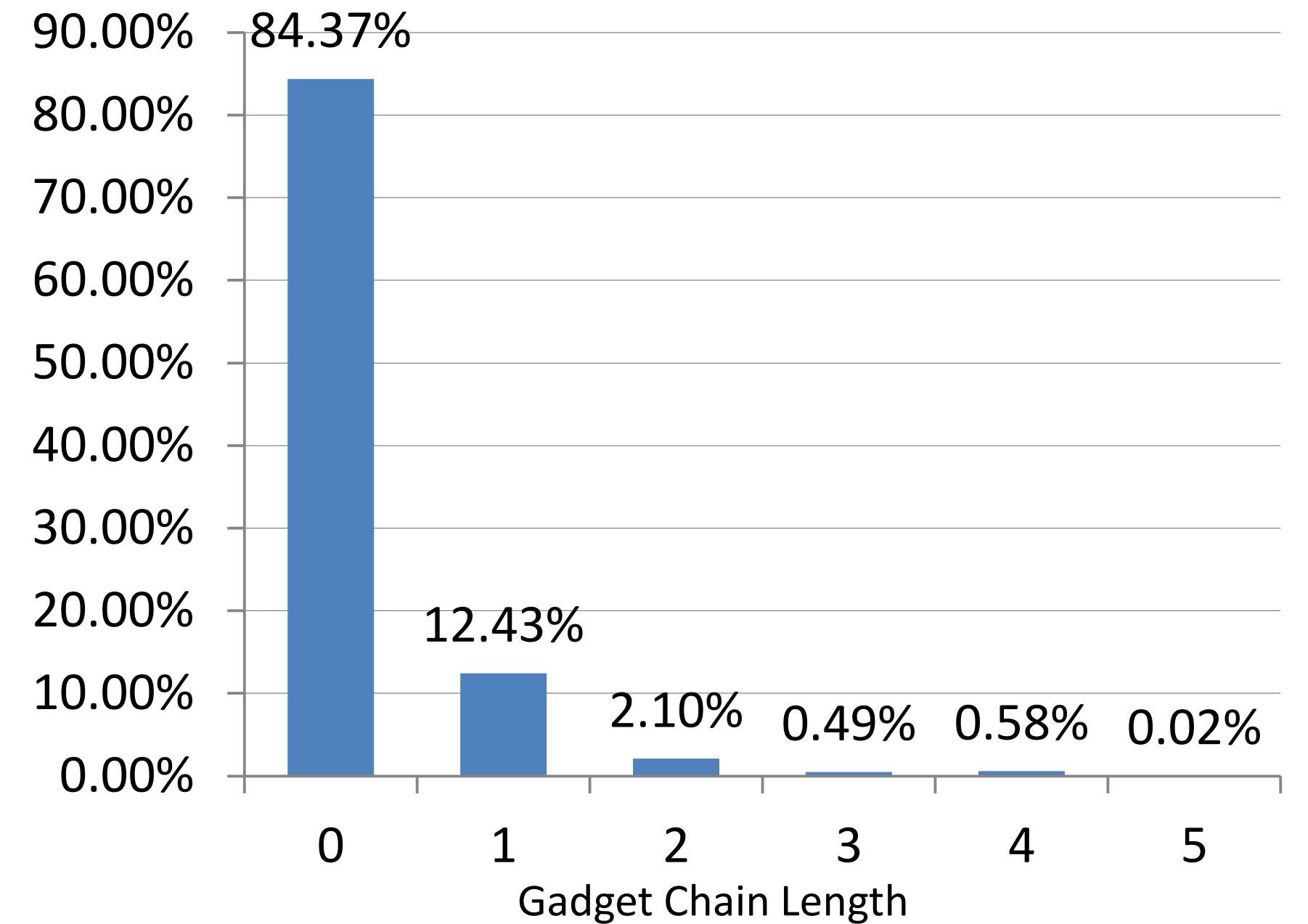
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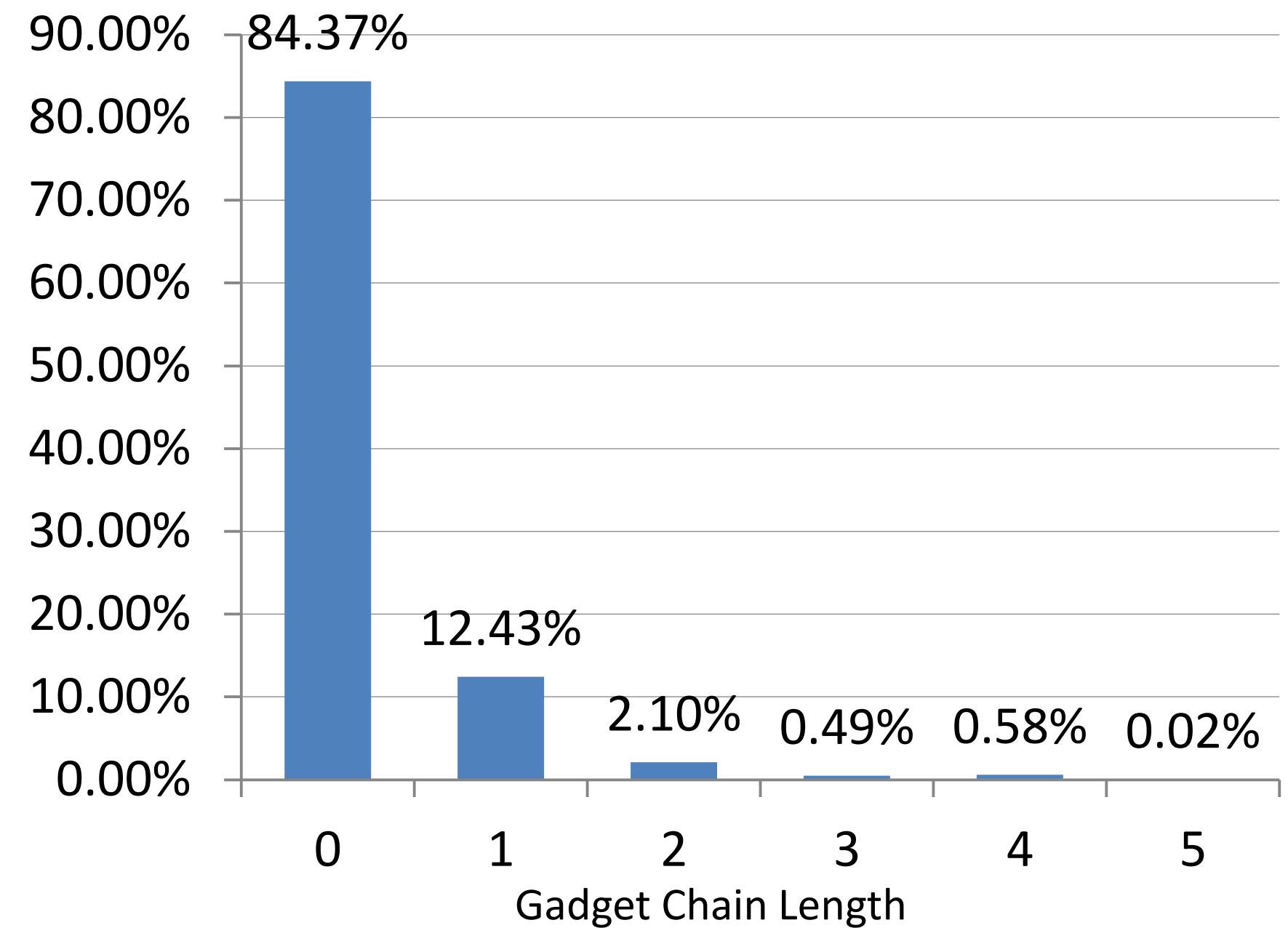


1. for indirect calls after context switch
2. count branches to potential gadgets

1. gadget chain too small?
2. emulate call stack, do we keep returning to suspected gadgets?



length of (potential) gadget chains in non-malicious code
17 popular Linux tools, 12 benchmark tools of SPEC INT2006, and 3 large binaries



if count from (2,3) > 11—16, have
ROP activity



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ROPGuard

Fratrić, Ivan. "ROPGuard: Runtime prevention of return-oriented programming attacks." (2012).

runtime checks when critical functions are called:

1. ensure stack pointer stays within area
(prevent stack pivot attack)
2. address of function would be on the top of the stack if called
using return otherwise not
3. instructions preceded by call; return address executable; call
must be related to function
4. valid stack: return address just below frame pointer
5. stack pointer should always be resolvable from machine code
6. function should not be allowed to make code sections
executable; don't load dlls from some places (network)

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kBouncer assumptions:

1. long gadgets unlikely
2. can't string together non-malicious indirects that exhaust LBR

strategy: last 14 indirect branches should be benign

1. identify: LBR-flushing function
 - a. legitimate function
 - b. many legitimate indirects (call-preceded)



strategy: last 14 indirect branches should be benign

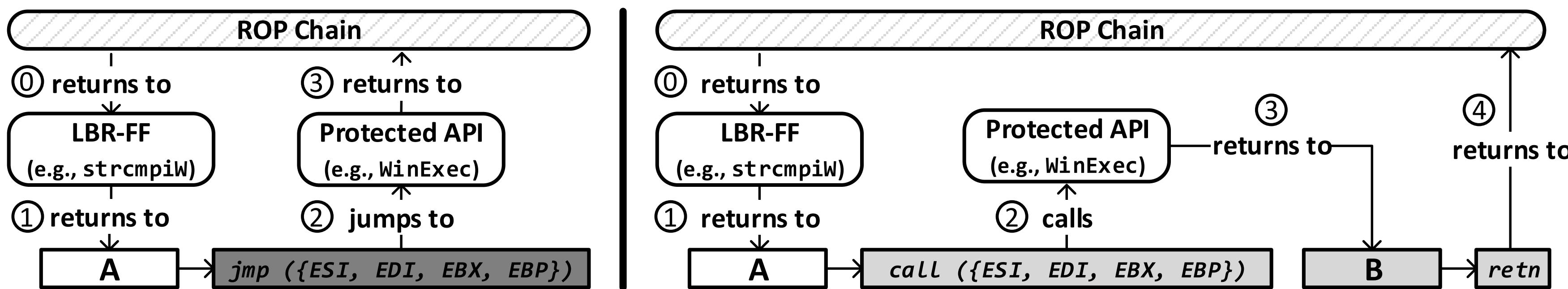
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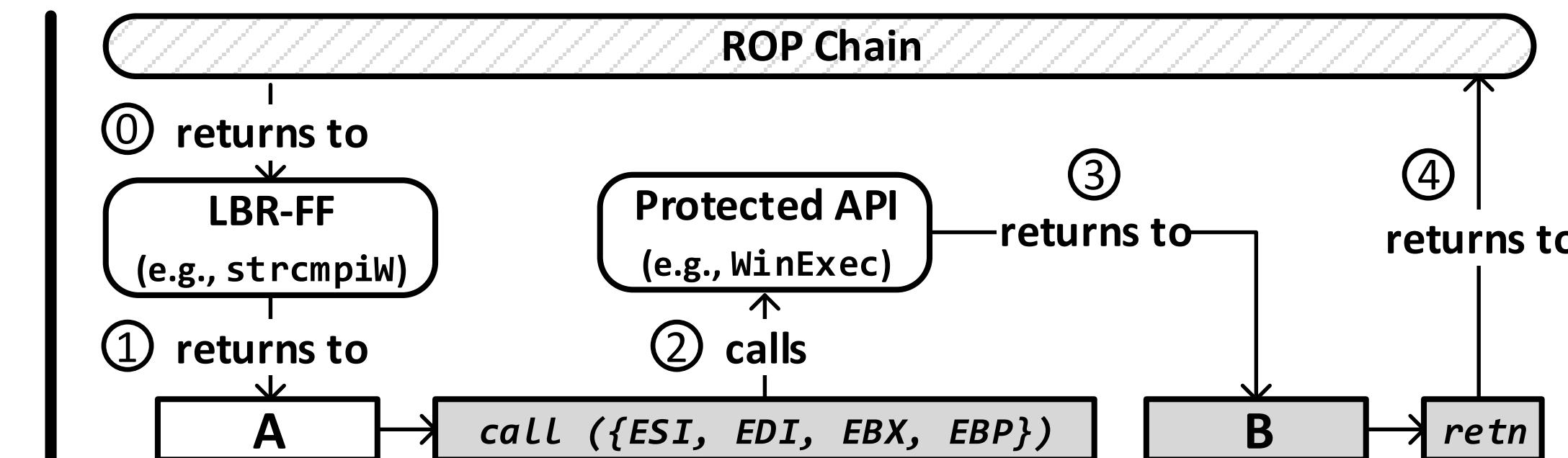
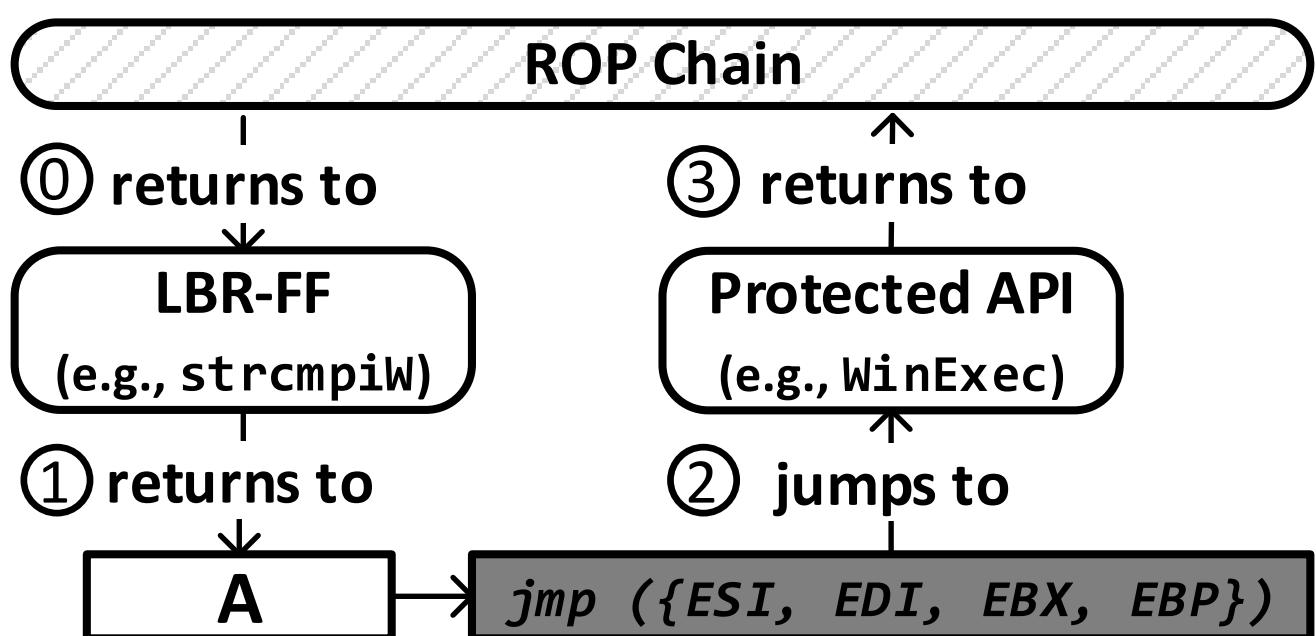
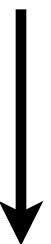
2. identify: call-preceded gadget that doesn't affect ROP state

i-jump-gadget
<code>call <anything></code>
A

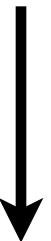
i-call-gadget
<code>call <anything></code>
A
<code>call ({ESI, EDI, EBX, EBP})</code>
B
<code>retn</code>



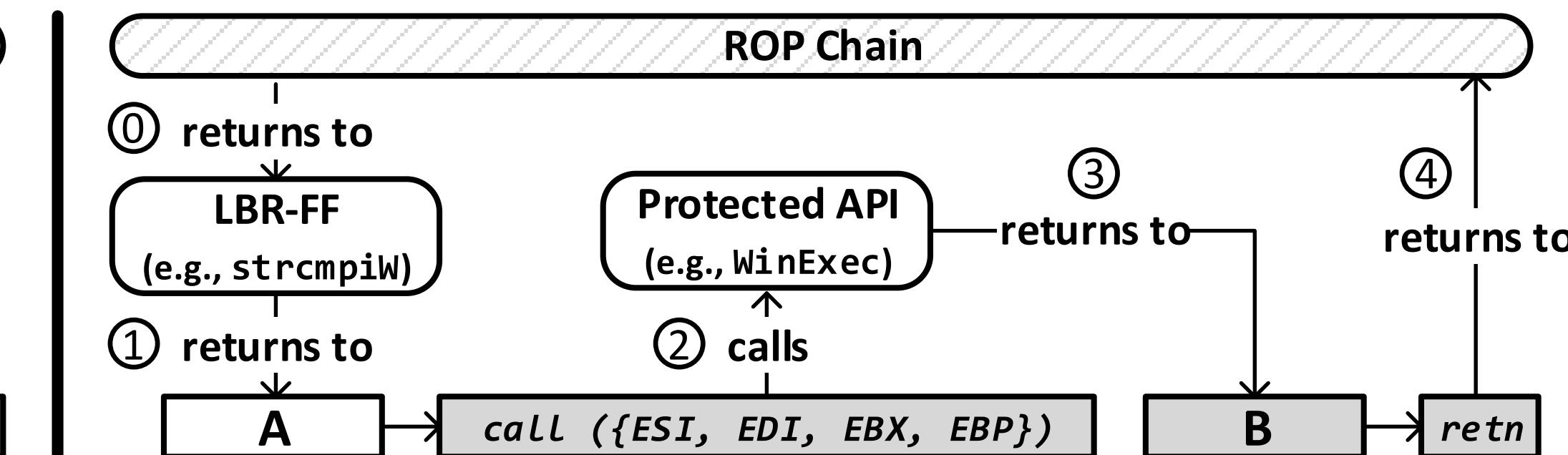
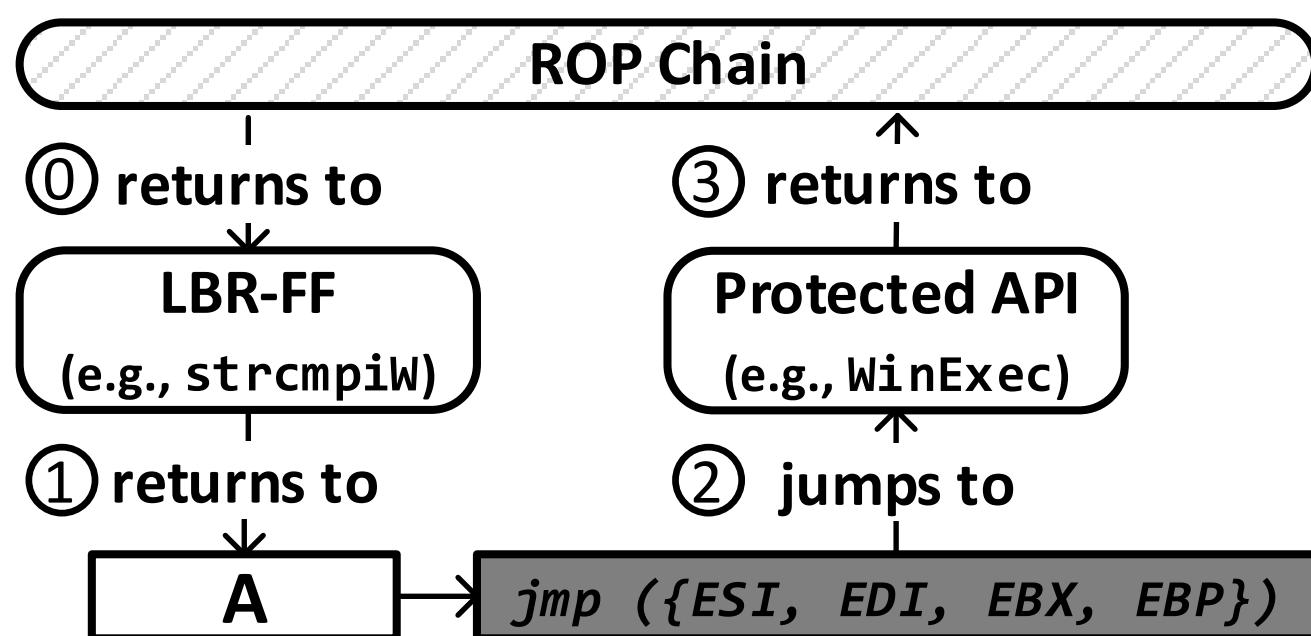
setup vars/args



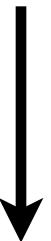
setup vars/args



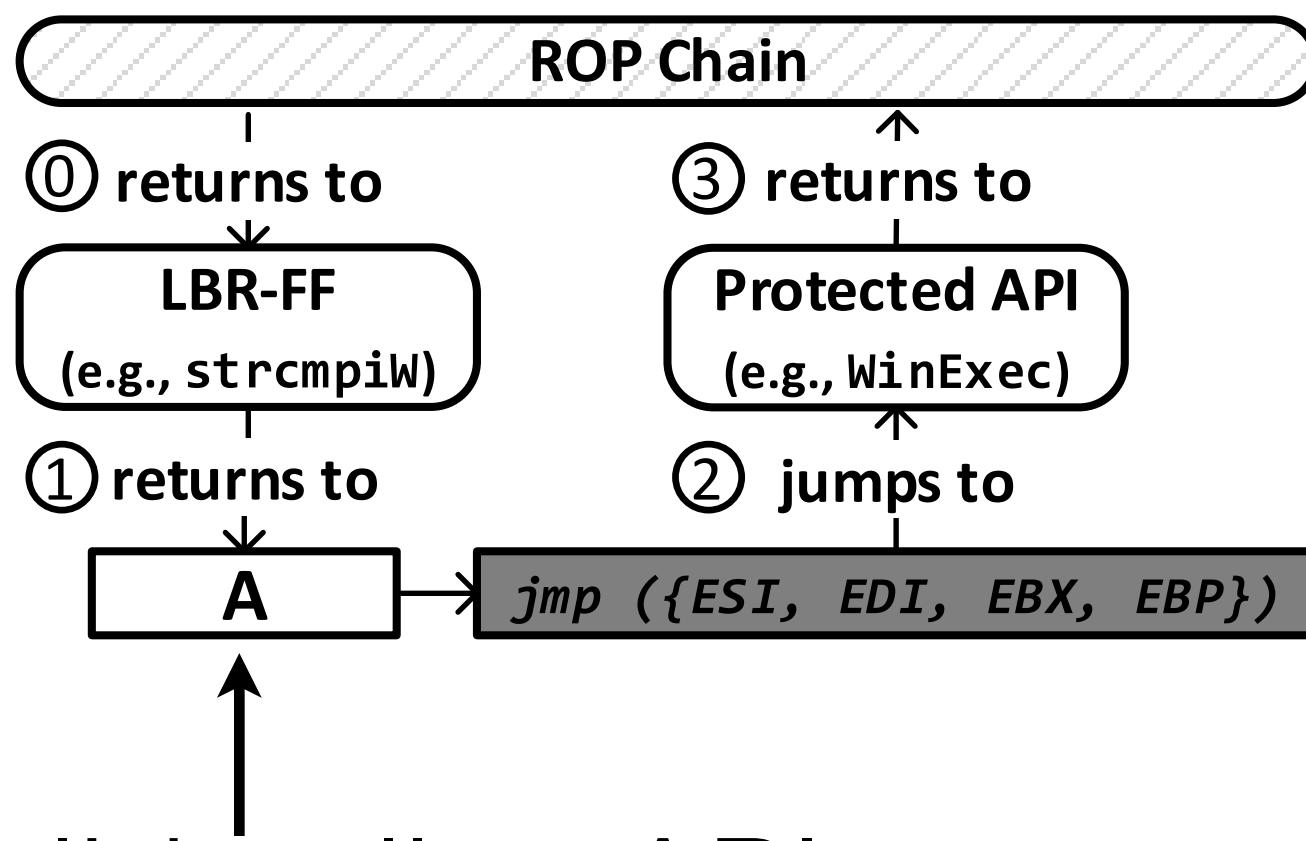
flush LBR →



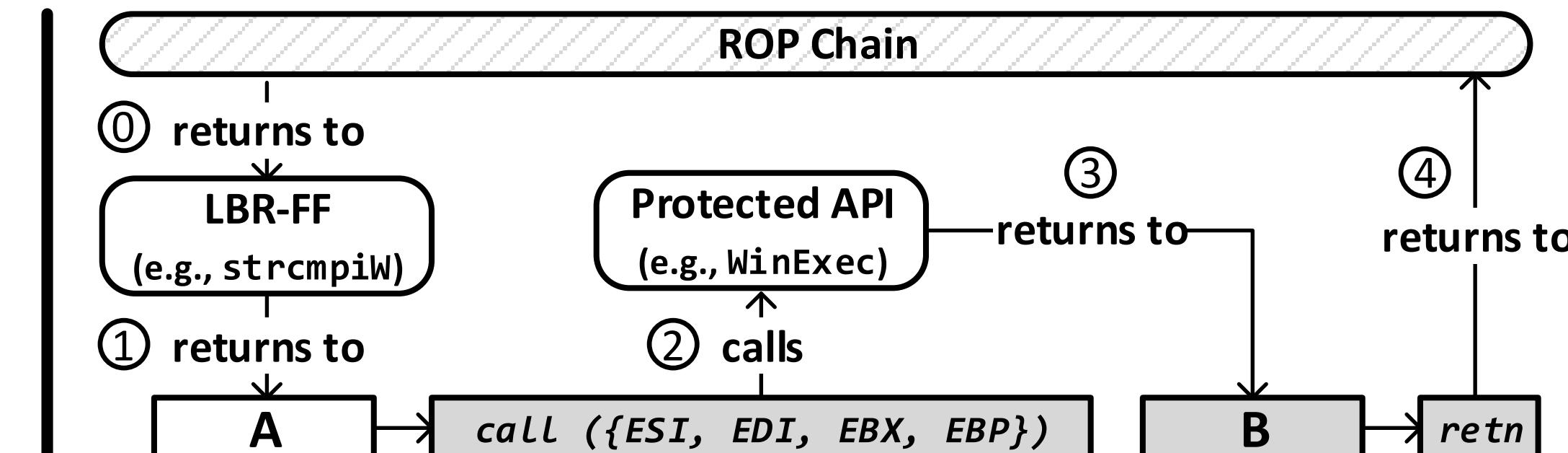
setup vars/args



flush LBR →



for valid call to API



ROPecker vulnerabilities:

1. long gadgets

ROPecker vulnerabilities:

1. long gadgets
2. gadget that contains direct branch

ROPecker vulnerabilities:

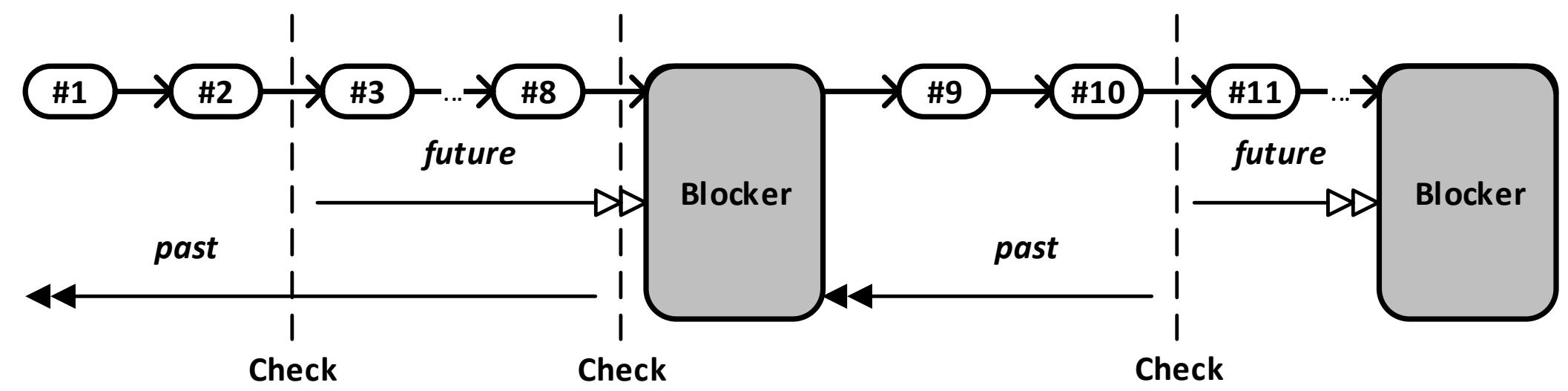
1. long gadgets
2. gadget that contains direct branch
to another gadget

ROPecker vulnerabilities:

1. long gadgets
2. gadget that contains direct branch
to another gadget
3. gadgets within windows

ROPecker vulnerabilities:

1. long gadgets
2. gadget that contains direct branch to another gadget
3. gadgets within windows



blocker: end of detection

1. smallest window size: one page
2. direct jumps to gadgets
3. long gadgets