#### - Complete Binary Tree

#### - Max-Heap

n = 6  
h = 2  
$$2^{h}$$
 = 4  
 $2^{h+1}$  - 1 = 7

### **Build Max-Heap in-place Iteratively**

- There two ways you can build the heap
- Both are iterative and in-place

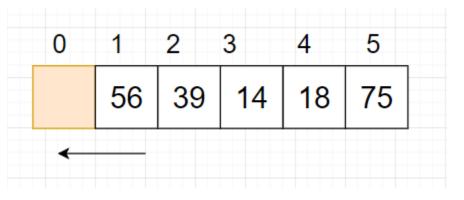
An **in-place algorithm** is an algorithm which transforms input using no auxiliary **data structure**. (Can use O(1) temporary variables).

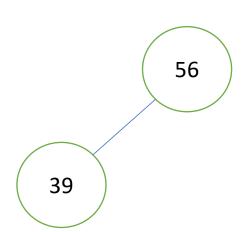
Top-down: O(nlog n).

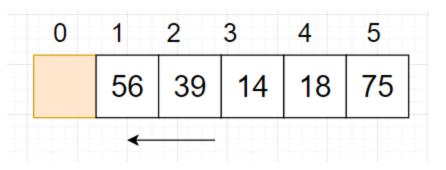
Bottom-up: O(n)

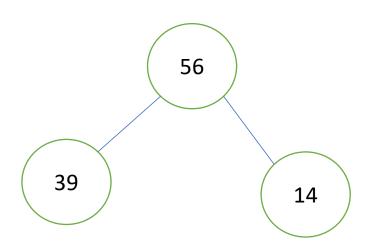
```
build_MaxHeap_TopDown(A, n)
  for (i <- 1 to n)
       upHeap(A, i)
                                       upHeap(A, i)
                                       j <- i
                                       while (j > 1 \& A[j/2] < A[j])
                                           swap(A[i], A[i/2])
                                            i < -i/2
```

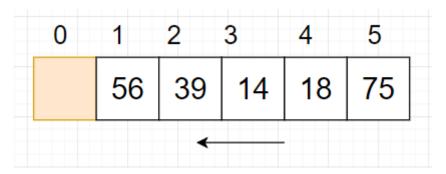


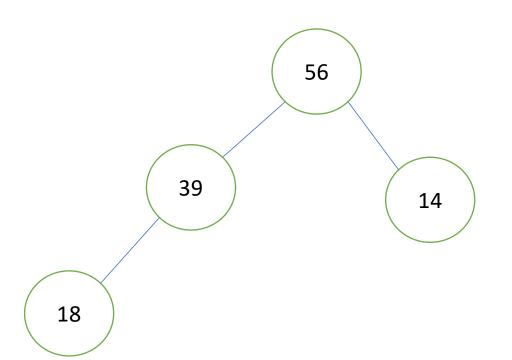


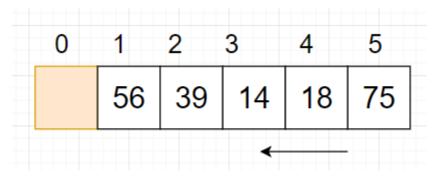


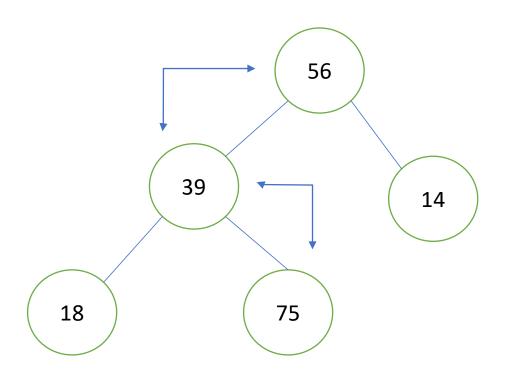


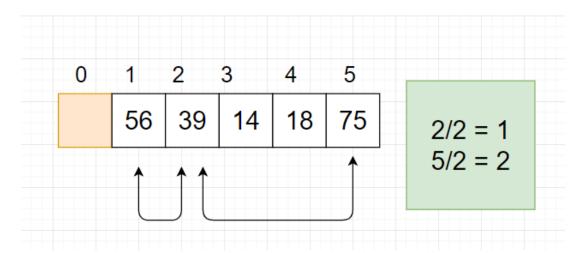


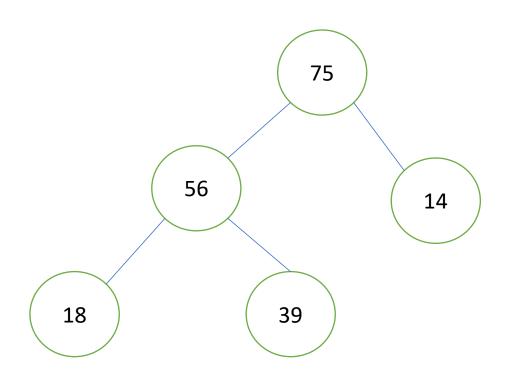










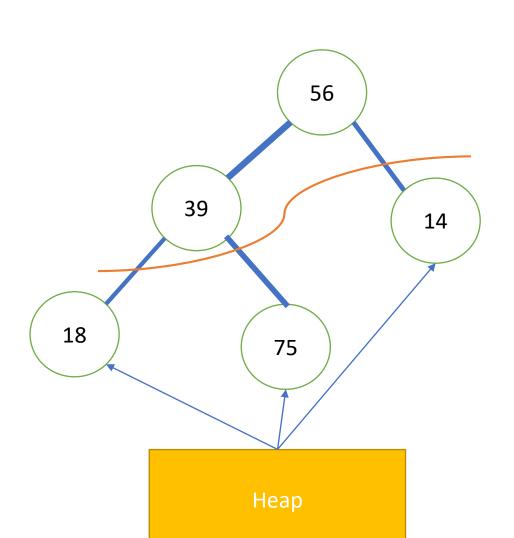


#### **Heap Sort**

There are two phases.

Phase I: Build the heap. In-place, bottom-up, iteratively

Phase II: Sort. In-place, iteratively

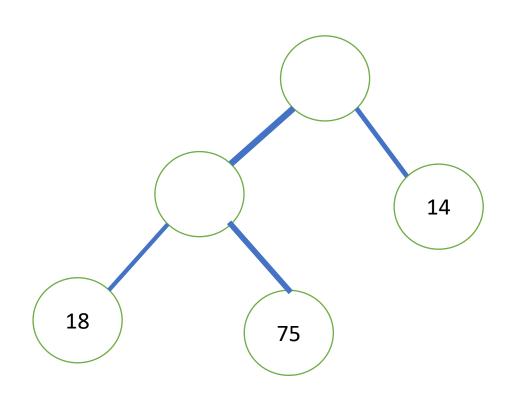


- $\lceil n/2 \rceil$  = 3 leaves are already heaps.
- There are  $\lfloor n/2 \rfloor = 2$  internal nodes.

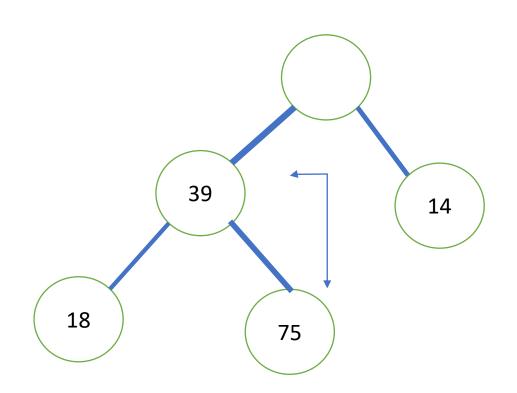
```
build_MaxHeap_BottomUp(A, n)
    for (i <- [n/2] to 1) downHeap(A, i)

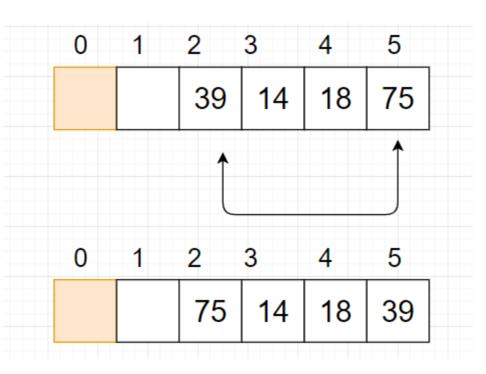
downHeap(A, i)
    j <- i
    k <- maxChildIndex(A, j)
    while (k!=0)
        swap(A[j], A[k])
        j <- k
        k <- maxChildIndex(A, j)</pre>
```

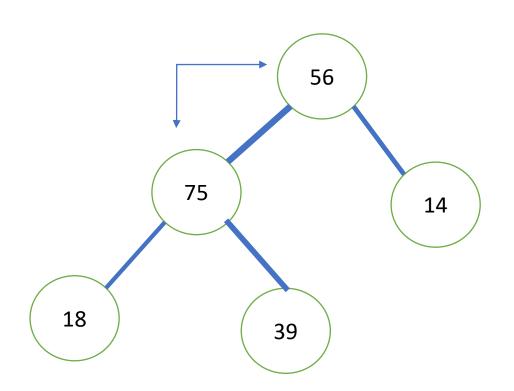
Note: maxChildIndex returns the index of a child of A[j] with maximum value among A[j], A[2j], A[2j+1]. If there is no such child it returns 0.

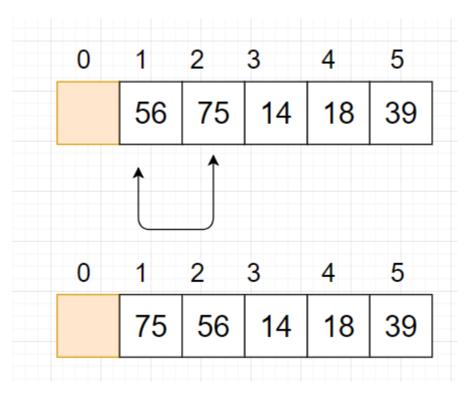


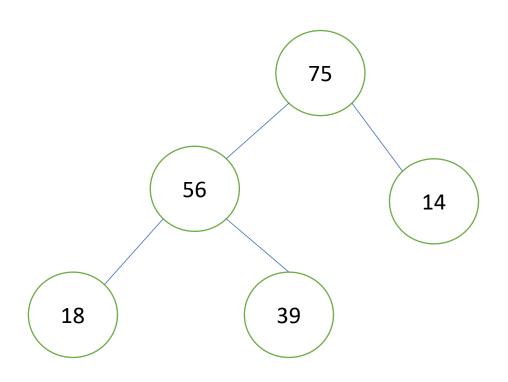
0	1	2	3	4	5
			14	18	75

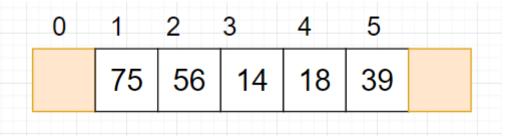












Put the root to sorted list, place the last element in the heap to the root and rebuild the heap

