컴퓨터 구조 5번째 과제

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[Review]

7.3) I/O module의 major function

Control and timing, processor communication, device communication, data buffering and error detection.

7.5) difference between memory-mapped I/O and isolated I/O

Memory-mapped I/O: there is a single address space for memory location and I/O devices. The processor treats the status and data registers of I/O modules as memory locations and uses the same machine instructions to access both memory and I/O devices.

Isolated I/O: The command line specifies whether the address refers to a memory location or an I/O device. The full range of addresses may be available for both. Address space for I/O is isolated from that for memory.

[Problems]

7.1)

For the first I/O instruction format, we have 8 bits for port addressing, therefore, the number of port addressed can be calculated as 2^8= 256 ports.

For the second instruction format, we have 16 bits for addressing, therefore, the number of ports addressed can be calculated as 2^16= 65536 ports.

Opcode selection allows us to change between the first and second instruction format, an opcode allows one input or output operation at a time.

7.2)

In direct addressing mode, an instruction can address up to $2^16 = 64$ K ports. In indirect addressing mode, the port address resides in a 16-bit registers, so again, the instruction can address up to $2^16 = 64$ K ports.

7.3)

2^16= 64KB

7.6)

- a. The printing rate is slowed to 5 cps.
- b. The situation must be treated differently with input devices such as the keyboard. It is necessary to scan the buffer at a rate of at least once per 60 ms. Otherwise, there is the risk of overwriting characters in the buffer.

7.9)

- a. The processor scans the keyboard 10 times per second. In 8 hours, the number of times the keykoard is scanned is $10 \times 60 \times 60 \times 8 = 288,000$.
- b. Only 60 visits would be required. The reduction is 1-(60/288000) = 0.999 -> 99%

7.11)

- a. 8KB/s \rightarrow 8 x 1024 = 8192B/s \rightarrow for 1B= 1/8192sec \rightarrow 1B/122 μ s(밀리세컨드). If each interrupt comsumes 122 μ s, then the fraction of processor time consumed is 100/122= 0.8196 = 0.82
- b. In this case, the time interval between interrupts is $16 \times 122 = 1952$ µs. Each interrupt now requires 100 µs for the first character plus the time for transferring each remaining character, which adds up to $8 \times 15 = 120$ µs, for a total of 220 µs. The fraction of processor time consumed is 220/1952 = 0.11
- c. each interrupt requires 100 µs for the first byte

plus 2 x 15 = 30 μ s for the remaining bytes, for a total of 130 μ s. The fraction of processor time consumed is 130/1952= 0.06