

Task-Based Programming

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Based on the Atmega328PB datasheet

Important Date for Mini Projects

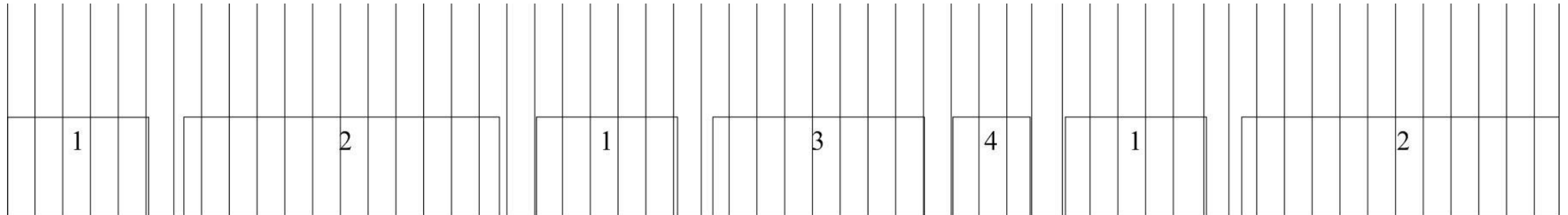
Date	To do List
4/8	<ul style="list-style-type: none">- Form a team with two students- Pick up one RedBot kit
4/10	<ul style="list-style-type: none">- Complete to form a team with two students- Practice basic RedBot functions
4/24	RedBot demo and racing: <ul style="list-style-type: none">- Task based programming code- Straight line and S curve line- Three fastest team will get 20% bonus points- Need to submit codes as well
5/6	3:30~ BNO055 based RedBot moving <ul style="list-style-type: none">- Three fastest team will get 20% bonus points- Need to submit codes as well

The Importance of Organizing Code

- The organized and neat code is highly preferable to inefficient and confusing code.
- As a general rule of thumb, the more that you write, the more debugging you will need to do later.
- We also want to reduce the amount of code sitting in our ISRs because 1) it's more difficult to debug, 2) Long ISRs may keep interrupts disabled so long that the MCU misses other events.
- We also want to cram more and more features/peripherals into our programs without hurting our performance.
- How can we do this???

Task-Based Programming

- Suppose that we have, 1. ADC, 2. Math Function #1, 3. UART, 4. Math Function #2
- Each task takes a certain amount of time to finish, so it's more efficient if we periodically space them out and repeat them at specific intervals.



Task Based Programming

- Ex. According to the datasheet, after the initial conversion, ADC will take ~ 13 ADC Clock cycles. This gives you insight as to how much time the 'task' should take. Doing this estimation for all tasks can give us a rough idea of how to schedule all of our features.
- Suppose that each task, 1, 2, 3, and 4 have software timer variables,

```
int ADC_Timer  
int Math_1_Timer  
int UART_Timer  
int Math_2_Timer
```

Task Based Programming

- Let's say that our ISR occurs every 1mS, ADC takes 2mS, Math_1 takes 4mS, UART takes 3mS, and Math_2 takes 5mS. Initializing our timers with 1mS in between them,

```
int ADC_Timer      = 18; // Waits 0mS
int Math_1_Timer   = 21; // Waits 3mS
int UART_Timer     = 26; // Waits 8mS
int Math_2_Timer   = 30; // Waits 12mS
```

- The whole cycle, repeating our four functions will take 18mS. With a 1mS ISR we decrement our variables at this frequency,

```
ISR(TIMERN_COMPN_Vect)
{
    if(ADC_Timer > 0)      ADC_Timer--;
    if(Math_1_Timer > 0)   Math_1_Timer--;
    if(UART_Timer > 0)     UART_Timer--;
    if(Math_2_Timer > 0)   Math_2_Timer--;
}
```

Task Based Programming

- The while(1) loop can now keep track of the status of our timer/counters. Once they've been decremented to zero, we reset the timers and then call our functions. Since we've staggered them already, they can be reset to the same values. You can space them out whatever you like. However, remember that the timer ISR will keep counting during the execution of your functions, and you may have two timers == 0 at the same time!!

```
while(1)
{
    if(ADC_Timer == 0{
        ADC_Timer=18; ADC();
    }

    if(Math_1_Timer == 0){
        Math_1_Timer=18; Math1();
    }
    if(UART_Timer == 0){
        UART_Timer=18;
        UART();
    }
    if(Math_2_Timer == 0){
        Math_2_Timer=18;
        Math2();
    }
}
```

Task Based Programming

- In summary, Task Based Programming allows us to organize and schedule many functions without reducing our performance.
- It also reduces the amount of code we keep in our timer ISR, and provides a basis for us to add things like tickets, variable scheduling times, permissions/priority (the basics of an operating system).

Example: How are the tasks scheduled?

```
while (1)
{
    if (task1_timer == 0)           // if task1_timer is not already equal to 0,
                                    // it is being decremented every 1 millisecond
                                    // during a timer ISR
    {
        task1_timer = t1;
        task1();                   // task1 takes m1 milliseconds
    }

    if (task2_timer == 0)           // if task2_timer is not already equal to 0,
                                    // it is being decremented every 1 millisecond
                                    // during a timer ISR
    {
        task2_timer = t2;
        task2();                   // task2 takes m2 milliseconds
    }
}
```

Example Cont'd

- Suppose $t1=5$, $m1=1$, $t2=10$, and $m2=15$
- What is the frequency $f1$ in Hz at which `task1()` is called?
- What is the frequency $f2$ in Hz at which `task2()` is called?
- Answer:
 - Since both `task1_timer` and `task2_timer` are decremented to 0 during the execution of `task2()`, `task1()` and `task2()` alternate.
 - Therefore, $f1=f2 = 1$ every 16 ms which is equal to $1000/16$ Hz.

Example Cont'd

- Suppose $t_1=20$, $m_1=1$, $t_2=10$, and $m_2=15$
- What is the frequency f_1 in Hz at which `task1()` is called?
- What is the average frequency f_2 in Hz at which `task2()` is called?

Answer:

- Since `task2_timer` is decremented to 0 during the execution of `task2()`, `task2()` is called as often as possible.
- When it is `task1()`'s turn to be executed, it takes more than one and less than two executions of `task2_timer` to get `task1()` decremented to 0.
- Therefore, the execution pattern converges to a repetition of `task2()` (takes 15 ms), `task2()` (takes 15 ms), `task1()` (takes 1 ms) giving
 - a frequency $f_1=1000/31$ Hz and
 - an average frequency $f_2=2 * 1000/31$.

Example Cont'd

- Suppose $t_1=20$, $m_1=1$, $t_2=25$, and $m_2=15$
- What is the frequency f_1 in Hz at which $\text{task1}()$ is called?
- What is the frequency f_2 in Hz at which $\text{task2}()$ is called?
- Answer:
 - During the time that $\text{task2}()$ is executed (which takes 15 ms), task1_timer (which initial value is 20) is decremented to a value $v \leq 5$.
 - The MCU will be idle for v ms after which task2_timer is decremented to $25-15-v$ and task1_timer just turned into 0.
 - So, after v ms $\text{task1}()$ is executed taking 1 ms during which task1_timer reduces to 19 and task2_timer reduces by 1 to $9-v$.
 - The MCU will be idle for another $9-v$ ms after which task1_timer is equal to $10+v$ and task2_timer just turned into 0.
 - Now $\text{task2}()$ is executed (which takes 15 ms) after which task1_timer is equal to 0 and task2_timer is equal to 10.
 - The same argument is now repeated for $v=0$ showing that the execution pattern converges to a repetition of $\text{task2}()$ (takes 15 ms), $\text{task1}()$ (takes 1 ms), idle time (takes 9 ms) giving
 - a frequency $f_1 = f_2 = 1000/25$ Hz.

Example Cont'd

- Suppose $t_1=4$, $m_1=1$, $t_2=8$, and $m_2=4$.
- Assume initially $\text{task1_timer} = 0$ and $\text{task2_timer} = t_2$
- What is the average frequency f_1 in Hz at which $\text{task1}()$ is called?
- What is the average frequency f_2 in Hz at which $\text{task2}()$ is called?
- Answer:
 - Task 1 executes during the intervals $[12n, 12n+1]$, $[12n+5, 12n+6]$, for integers $n \geq 0$.
 - Task 2 executes during intervals $[12n+8, 12n+12]$ for integers $n \geq 0$.
 - This gives frequencies $f_1 = 1000 \cdot 2 / 12$ Hz and $f_2 = 1000 / 12$ Hz.