# 01204211 Discrete Mathematics Lecture 8b: Modular arithmetic

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# Quick check 1

If a|m and b|m, can we say that ab|m? Prove this fact or provide a counter example.

### Quick check 2

If a|m, b|m, and  $a \neq b$  are both prime, can we say that ab|m? Prove this fact or provide a counter example.

#### Prime factorization

One useful fact that we use over and over again is the following.

# Unique Factorization (or Fundamental Theorem of Arithmetic)

Every integer greater than 1 can be written *uniquely* as a product of prime numbers (up to the order of factors).

#### Examples:

- $ightharpoonup 10 = 2 \cdot 5$
- **▶** 13 = 13
- $112 = 2 \cdot 2 \cdot 2 \cdot 2 \cdot 7 = 2^4 \cdot 7$

There are 3 clocks. At this moment, all three clocks ring at the same time. The first clock rings every 3 hours, the second clock rings every 4 hours, and the third clock rings every 10 hours. How long do you have to wait until you would hear all clocks ring a the same time again?

You have a large water container and two smaller buckets. The first bucket carries 3 litres of water and the second bucket carries 5 litres of water. Can you put exactly 1 litre of water in the water container?

You have a large water container and two smaller buckets. The first bucket carries 6 litres of water and the second bucket carries 15 litres of water. What is the minimum volume of water you can exactly put in the water container?

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In general if you have two buckets of volumes  $\boldsymbol{x}$  and  $\boldsymbol{y}$ , the amount that you can exactly make must be in the form of

$$ax + by$$
,

for some integers x and y. (Note that x and y may be negative.)

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for some integers x and y. (Note that x and y may be negative.) Do you see why the sum must be divisible by any common divisor of x and y?



### Useful fact

For any integer x and y, consider the term

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When the term is non-zero, it must be divisible by gcd(x,y), so it has to be at least gcd(x,y).

It turns out that you can actually attain that value, i.e., there exist a pair of integer  $\boldsymbol{a}$  and  $\boldsymbol{b}$  such that

$$a \cdot x + b \cdot y = gcd(x, y).$$

# Finding a and b: Extended Euclid Algorithm

We will modify the Euclid algorithm so that it also returns a and b together with  $\gcd(x,y)$ .

```
Algorithm Euclid(x,y):
if x \mod y == 0:
  return y,
else:
  g, a', b' = Euclid(y, x \mod y)
  a =
  b =
  return g, a, b
```

### Notes:

We have a' and b' such that

$$a' \cdot y + b' \cdot (x \bmod y) = g.$$

# Secret sharing

# Secret sharing scheme based on straight lines

What day is it today?

What day is it today? Thursday.

What day is it today? Thursday. What day is 3 days after today?

What day is it today? Thursday. What day is 3 days after today? Sunday.

What day is it today? Thursday. What day is 3 days after today? Sunday. What day is 20 days after today?

What day is it today? Thursday. What day is 3 days after today? Sunday. What day is 20 days after today? Wednesday.

What day is it today? Thursday. What day is 3 days after today? Sunday. What day is 20 days after today? Wednesday. What day is 10 days before today?

What day is it today? Thursday. What day is 3 days after today? Sunday. What day is 20 days after today? Wednesday. What day is 10 days before today? Monday.

Suppose that it is  $1\ {\rm o'clock}.$ 

Suppose that it is 1 o'clock. What time is the next 5 hours?

Suppose that it is 1 o'clock. What time is the next 5 hours? 6 o'clock.

Suppose that it is 1 o'clock. What time is the next 5 hours? 6 o'clock. What time is the next 10 hours?

Suppose that it is 1 o'clock.

What time is the next 5 hours? 6 o'clock.

What time is the next 10 hours? 11 o'clock.

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Suppose that it is 1 o'clock. What time is the next 5 hours? 6 o'clock. What time is the next 10 hours? 11 o'clock. What time is the next 20 hours? 9 o'clock.

As in the days of weeks and clocks examples (and also as the modulo in RSA algorithm in our experiment), when working under modular arithmetic, we start with a  $\frac{\text{modulus}}{m}$ .

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$$3 \cdot 4 =$$

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$$2-6 = -4 \mod 7 = 3 \mod 7 = 3.$$

Note that when you view integers under the lense of modulus 7, these numbers

$$\dots, -19, -12, -5, 2, 9, 16, 23, \dots$$

are essentially the same.



# Properties (1)

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#### Proof.

 $(\Rightarrow)$  Let  $r = a \mod m$ . We can write

$$a = qm + r,$$

and

$$b = pm + r,$$

for some integers q and p. Thus, we have

$$a - b = qm + r - pm - r = (q - p)m.$$

Therefore m|a-b.

(⇐) Exercise.



# Properties (2)

- $(a+b) \bmod m = ((a \bmod m) + (b \bmod m)) \bmod m$
- $(a-b) \bmod m = ((a \bmod m) (b \bmod m)) \bmod m$
- $(a \cdot b) \bmod m = ((a \bmod m) \cdot (b \bmod m)) \bmod m$

## Congruences

#### Definition (congruences)

For an integer m>0, if integers a and b are such that

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we write

$$a \equiv b \pmod{m}$$
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We also have that

$$a \equiv b \pmod{m} \Leftrightarrow m|(a-b)$$



# Congruences: properties (1)

- (reflexivity)  $a \equiv a \pmod{m}.$
- (symmetry)  $a \equiv b \pmod{m}$  implies  $b \equiv a \pmod{m}$ .
- ▶ (transitiviey)  $a \equiv b \pmod{m}$  and  $b \equiv c \pmod{m}$  implies  $a \equiv c \pmod{m}$ .

If we have that

$$a \equiv b \pmod{m}$$
,

and

$$c \equiv d \pmod{m}$$
,

then

- $a + c \equiv b + d \pmod{m}$
- $a c \equiv b d \pmod{m}$
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What is missing here?

If we have that

$$a \equiv b \pmod{m}$$
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and

$$c \equiv d \pmod{m}$$
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then

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We can pretty much think of this "congruence" as a normal equation.

What is missing here? Division!

Also, we wish we can do "cancellation", i.e., if

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Also, we wish we can do "cancellation", i.e., if

$$xa \equiv xb \pmod{m}$$
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then  $a \equiv b \pmod{m}$ . BUT THIS IS NOT ALWAYS TRUE. Let's see the following example:

$$2 \cdot 1 \equiv 2 \cdot 3 \pmod{4},$$

but

$$1 \not\equiv 3 \pmod{4}$$
.

### Multiplications as functions

Let's view multiplication by 2 as a function, i.e., let  $f(x) = 2 \cdot x \mod 4$ .

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Which functions have inverses?

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We are looking to a number x such that 2=5x. How can we do that? By dividing on both sides with 5:

$$2/5 = 5x/5 = x,$$

or equivalently, by multiplying with  $(1/5) = 5^{-1}$ :

$$2 \cdot 5^{-1} = 5x \cdot 5^{-1} = x \cdot 5 \cdot 5^{-1} = x \cdot 1 = x.$$

Here  $5^{-1}$  is a multiplicative inverse of 5.



# Multiplicative inverses (modular arithmetic)

You can do the same thing in modular arithmetic. Let the modulus be  $m=7. \ \mathrm{Note}$  that

$$5 \cdot 3 \equiv 15 \equiv 1 \pmod{7}.$$

Therefore,  $5^{-1} \equiv 3 \pmod{7}$ .

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To find 2/5, we can view our goal as to find the value of x such that

$$2 \equiv 5x \pmod{7}$$
.

We can multiply both sides with  $5^{-1} \equiv 3$  to get

$$2 \cdot 5^{-1} \equiv 2 \cdot 3 \equiv 6 \equiv 5^{-1} \cdot 5x \equiv x \pmod{7}.$$

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Let's check:

$$5 \cdot 6 \equiv 30 \equiv 2 \pmod{7}$$
,

as requied.



### Multiplicative inverse modulo m

#### Definition

The multiplicative inverse modulo m of a, denoted by  $a^{-1}$ , is an integer such that

$$a \cdot a^{-1} \equiv 1 \pmod{m}$$
.

### Multiplicative inverse modulo 11

Let's try to figure out multiplicative inverse of every integer modulo  $11.\,$ 

	c. y c.	o ligare out
a	$a^{-1}$	$\pmod{11}$
1		
2 3		
4		
5		
6		
7		
8 9		
9		
10		

## Example: secret sharing

- ▶ Think of a secret number  $m \in \{0, 1, ..., 10\}$ .
- ▶ Pick a random number  $a \in \{1, 2, ..., 10\}$ .
- ▶ Your straight line function  $f(x) = (ax + m) \mod 11$ .
- We will generate 3 points from f and give them to 3 of your friends, each with only 1 point. Pick 3 numbers  $x_1, x_2, x_3$  from  $\{1, 2, \dots, 10\}$ .
- Let's compute

$$(x_1, f(x_1)), (x_2, f(x_2)), (x_3, f(x_3)).$$

ightharpoonup Give them to 3 of your friends and challenge them to form a group of 2 people and figure out your number m.



#### Theorem 1

An integer a has a multiplicative inverse modulo m iff  $\gcd(a,m)=1$ .

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#### Proof.

 $(\Leftarrow)$  Recall that there exist integers x and y such that

$$x \cdot a + y \cdot m = \gcd(a, m) = 1.$$

Thus,  $(x \cdot a + y \cdot m) \mod m = x \cdot a \mod m = 1 \mod m$ , i.e.,  $x \cdot a \equiv 1 \pmod m$ . Therefore x is the inverse.

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 $(\Rightarrow)$  Let r=gcd(a,m). Suppose that b is the multiplicative inverse of a modulo m, i.e., we have that

$$b \cdot a \equiv 1 \pmod{m}$$
,

Thus,  $ba \mod m = 1 \mod m = 1$ , i.e., there exists an integer q such that

$$ba = qm + 1,$$

or ba-qm=1. However, r since r|a and r|m, r also divides bd-qm and 1. But it  $r \not| 1$  because r>1 and we have the contradiction.

### Examples: division in modular arithmetic

Since the requirement for an existance of  $a^{-1}$  modulo m is that gcd(a,m)=1, if we let m be a prime number, every a which is not a multiple of m has an inverse. Can you solve this equation?

$$4x + 9 \equiv 0 \pmod{11}.$$

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$$4x + 9 \equiv 0 \pmod{11}.$$

We can even perform gaussian elimination (which is very useful later):

$$2x + y \equiv 3 \pmod{7}$$
$$x + 3y \equiv 5 \pmod{7}$$

# Public-key cryptography

## **RSA**

### Quick recap: RSA

- Private key: (e, n), Public key: (d, n)
- ▶ Encryption  $E(m) = m^e \mod n$ , Decryption:  $D(w) = w^d \mod n$ .
- ▶ Goal: Select e, d, n such that  $D(E(m)) = m^{ed} \mod n = m$ .

## Quick recap: RSA

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- ▶ Encryption  $E(m) = m^e \mod n$ , Decryption:  $D(w) = w^d \mod n$ .
- ▶ Goal: Select e, d, n such that  $D(E(m)) = m^{ed} \mod n = m$ .
- Pick two primes p and q. Let n = pq.
- Pick e (usually a small number)
- ▶ Pick d such that  $d = e^{-1} \pmod{(p-1)(q-1)}$ , i.e.,  $ed \equiv 1 \pmod{(p-1)(q-1)}$ , or

$$ed = k \cdot (p-1)(q-1) + 1,$$

for some integer k.

▶ What is  $m^{ed} \mod n$ ?

#### What's next?

- We will prove Fermat's Little Theorem and show how to efficiently test if a number is prime.
- ▶ We will also use Fermat's Little Theorem to prove the correctness of RSA.
- Modular arithmetic is also key to our usage of polynomials to perform secret sharing and error correcting codes, because now we can do Gaussian elimination using only integers.