01204211 Discrete Mathematics
Lecture 8b: Modular arithmetic

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Quick check 1

If $a|m$ and $b|m$, can we say that $ab|m$? Prove this fact or provide a counter example.

Quick check 2

relatively prime

$\gcd(a,b) = 1$

If $a|m$, $b|m$, and $a \neq b$ are both prime, can we say that $ab|m$? Prove this fact or provide a counter example.

Prime factorization

One useful fact that we use over and over again is the following.

Unique Factorization (or Fundamental Theorem of Arithmetic)

Every integer greater than 1 can be written *uniquely* as a product of prime numbers (up to the order of factors).

Examples:

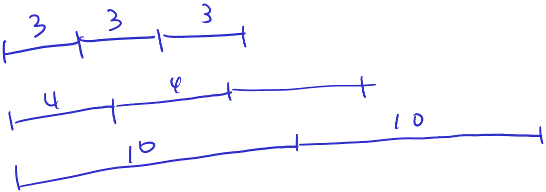
- $10 = 2 \cdot 5$
- $13 = 13$
- $112 = 2 \cdot 2 \cdot 2 \cdot 2 \cdot 7 = 2^4 \cdot 7$

$3 \cdot 2 \cdot 2 \cdot 5 = 60$

$3 \mid x$, $4 \mid x$, $10 \mid x$

$4 = 2 \cdot 2$, $10 = 2 \cdot 5$

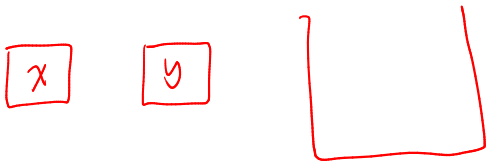
There are 3 clocks. At this moment, all three clocks ring at the same time. The first clock rings every 3 hours, the second clock rings every 4 hours, and the third clock rings every 10 hours. How long do you have to wait until you would hear all clocks ring a the same time again?



18 15

You have a large water container and two smaller buckets. The first bucket carries 3 litres of water and the second bucket carries 5 litres of water.
Can you put exactly 1 litre of water in the water container?

You have a large water container and two smaller buckets. The first bucket carries 6 litres of water and the second bucket carries 15 litres of water.
What is the minimum volume of water you can exactly put in the water container?



ปริมาตรที่น่าใส้

$ax + by$

a, b เป็นจำนวนเต็ม.

You have a large water container and two smaller buckets. The first bucket carries 6 litres of water and the second bucket carries 15 litres of water.
What is the minimum volume of water you can exactly put in the water container?

In general if you have two buckets of volumes x and y, the amount that you can exactly make must be in the form of

$$ax + by, \qquad = 1 \quad ?$$

for some integers x and y. (Note that x and y may be negative.)

You have a large water container and two smaller buckets. The first bucket carries 6 litres of water and the second bucket carries 15 litres of water.
What is the minimum volume of water you can exactly put in the water container?

In general if you have two buckets of volumes x and y, the amount that you can exactly make must be in the form of

$$ax + by,$$

for some integers x and y. (Note that x and y may be negative.)
Do you see why the sum must be divisible by any common divisor of x and y?

Useful fact

For any integer x and y, consider the term

$$a \cdot x + b \cdot y,$$

for some integer a and b.

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For any integer x and y, consider the term

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for some integer a and b.

When the term is non-zero, it must be divisible by $gcd(x, y)$, so it has to be at least $gcd(x, y)$.

It turns out that you can actually attain that value, i.e., there exist a pair of integer a and b such that

$$a \cdot x + b \cdot y = gcd(x, y).$$

Finding a and b: Extended Euclid Algorithm

We will modify the Euclid algorithm so that it also returns a and b together with $gcd(x, y)$.

```
Algorithm Euclid(x,y):
  if x mod y == 0:
    return y, 0, 1
  else:
    g, a', b' = Euclid(y, x mod y)
    a = b'
    b = a' - b' · (⌊x/y⌋)
    return g, a, b
```

$g = gcd(y, x \bmod y) = gcd(x, y)$

จาก recursive call

$$a' \cdot y + b' (x \bmod y) = g \quad -(1)$$

จาก (1) จะได้: $x \bmod y = x - \lfloor \frac{x}{y} \rfloor \cdot y$

จะได้ว่า $a' \cdot y + b' \left[x - \lfloor \frac{x}{y} \rfloor \cdot y \right] = g$

นั่นคือ $b' \cdot x + \left[a' - b' \lfloor \frac{x}{y} \rfloor \right] \cdot y = g$

ที่ $g = gcd(x, y)$ จะได้: $ax + by = g$

Notes:

We have a' and b' such that

$$a' \cdot y + b' \cdot (x \bmod y) = g.$$

Thm: สมมติ จำนวนเต็ม x, y มี จำนวนเต็ม a, b ที่ $ax + by = gcd(x, y)$

$$a' \cdot y + b' \left(x - \left\lfloor \frac{x}{y} \right\rfloor \cdot y \right) = g$$

$$a' \cdot y + b' x - b' \left\lfloor \frac{x}{y} \right\rfloor \cdot y = g$$

$$\underbrace{b'}_{a} x + \underbrace{\left(a' - b' \left\lfloor \frac{x}{y} \right\rfloor \right)}_{b} \cdot y = g$$

Secret sharing

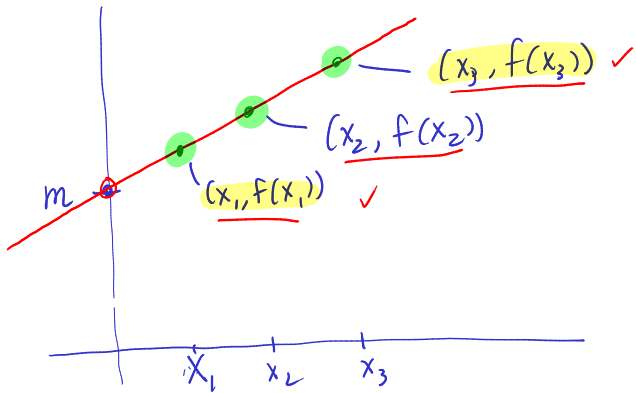
ความลับ

แบ่ง กุญแจ 3 คน

ใครที่ 2 คนขึ้นไป สามารถ สร้าง ความลับกลับมาได้

1 คนขึ้นไป ไม่ รู้อะไรเลยเกี่ยวกับ ความลับ.

Secret sharing scheme based on straight lines



$$ax_1 + m = f(x_1)$$

$$ax_3 + m = f(x_3)$$

Days

What day is it today?

Days

What day is it today? Thursday.

Days

What day is it today? Thursday.
What day is 3 days after today?

Days

What day is it today? Thursday.
What day is 3 days after today? Sunday.

Days

What day is it today? Thursday.
What day is 3 days after today? Sunday.
What day is 20 days after today?

Days

What day is it today? Thursday.
What day is 3 days after today? Sunday.
What day is 20 days after today? Wednesday.

Days

What day is it today? Thursday.
What day is 3 days after today? Sunday.
What day is 20 days after today? Wednesday.
What day is 10 days before today?

Days

What day is it today? Thursday.
What day is 3 days after today? Sunday.
What day is 20 days after today? Wednesday.
What day is 10 days before today? Monday.

Clocks

Suppose that it is 1 o'clock.

Clocks

Suppose that it is 1 o'clock.
What time is the next 5 hours?

Clocks

Suppose that it is 1 o'clock.
What time is the next 5 hours? 6 o'clock.

Clocks

Suppose that it is 1 o'clock.
What time is the next 5 hours? 6 o'clock.
What time is the next 10 hours?

Clocks

Suppose that it is 1 o'clock.
What time is the next 5 hours? 6 o'clock.
What time is the next 10 hours? 11 o'clock.

Clocks

Suppose that it is 1 o'clock.
What time is the next 5 hours? 6 o'clock.
What time is the next 10 hours? 11 o'clock.
What time is the next 20 hours?

Clocks

Suppose that it is 1 o'clock.
What time is the next 5 hours? 6 o'clock.
What time is the next 10 hours? 11 o'clock.
What time is the next 20 hours? 9 o'clock.

Modular arithmetic

As in the days of weeks and clocks examples (and also as the modulo in RSA algorithm in our experiment), when working under modular arithmetic, we start with a **modulus** m.

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$$4 + 5 = 9 \bmod m = 2.$$

Or

$$3 \cdot 4 =$$

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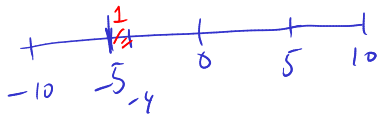

Or

$$3 \cdot 4 = 12 \bmod m = 5.$$

Or

$$2 - 6 = -4 \bmod 7 = 3$$

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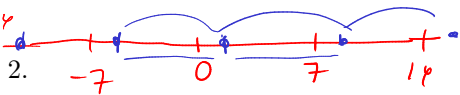

Or

$$3 \cdot 4 = 12 \bmod m = 5.$$

Or

$$2 - 6 = -4 \bmod 7 = 3 \bmod 7 = 3.$$

Note that when you view integers under the lense of modulus 7, these numbers

$$\ldots, -19, -12, -5, 2, 9, 16, 23, \ldots.$$

are essentially **the same**.

Properties (1)

$a \bmod m = b \bmod m$, if and only if $m|a - b$.

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Proof.

($\Rightarrow$) Let $r = a \bmod m$. We can write

$$a = qm + r,$$

and

$$b = pm + r,$$

for some integers q and p. Thus, we have

$$a - b = qm + r - pm - r = (q - p)m.$$

Therefore $m|a-b$.

($\Leftarrow$) Exercise.

□

Properties (2)

$1011002157 \mod 11177$

192141124210112

- $(a + b) \bmod m = ((a \bmod m) + (b \bmod m)) \bmod m$
- $(a - b) \bmod m = ((a \bmod m) - (b \bmod m)) \bmod m$
- $(a \cdot b) \bmod m = ((a \bmod m) \cdot (b \bmod m)) \bmod m$

▷ mod คำนวณแบ่งทางได้

Congruences

Definition (congruences)

For an integer $m > 0$, if integers a and b are such that

$$a \bmod m = b \bmod m,$$

we write

$$a \equiv b \pmod{m}.$$

$a \equiv b \pmod{m}$

$c \equiv d \pmod{m}$ $\Rightarrow a + c \equiv b + d \pmod{m}$

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We also have that

$$a \equiv b \pmod{m} \quad \Leftrightarrow \quad m|(a-b)$$

Congruences: properties (1)

- (reflexivity)
  $a \equiv a \pmod{m}$.
- (symmetry)
  $a \equiv b \pmod{m}$ implies $b \equiv a \pmod{m}$.
- (transitiviey)
  $a \equiv b \pmod{m}$ and $b \equiv c \pmod{m}$ implies $a \equiv c \pmod{m}$.

Congruences: properties (2) – operations

If we have that

congruence.

$$a \equiv b \pmod{m},$$

and

$$c \equiv d \pmod{m},$$

then

- $a + c \equiv b + d \pmod{m}$
- $a - c \equiv b - d \pmod{m}$
- $ac \equiv bd \pmod{m}$

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What is missing here?
Division!

Also, we wish we can do "cancellation", i.e., if

$$xa \equiv xb \pmod{m},$$

then $a \equiv b \pmod{m}$. **BUT THIS IS NOT ALWAYS TRUE.**

$m = 15$

① หา x, a, b $= 6, 5, 10$

ที่ $xa \equiv xb \pmod{15}$

แต่ $a \not\equiv b \pmod{15}$

$m = 11$

② หา x, a, b

$xa \equiv xb \pmod{11}$

$a \not\equiv b \pmod{11}$

Also, we wish we can do "cancellation", i.e., if

$$xa \equiv xb \pmod{m},$$

then $a \equiv b \pmod{m}$. **BUT THIS IS NOT ALWAYS TRUE.**
Let's see the following example:

$$2 \cdot 1 \equiv 2 \cdot 3 \pmod{4},$$

but

$$1 \not\equiv 3 \pmod{4}.$$

Multiplications as functions

Let's view multiplication by 2 as a function, i.e., let $f(x) = 2 \cdot x \bmod 4$.

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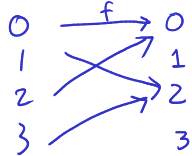
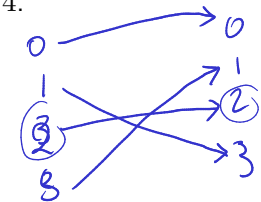


Let's also see $g(x) = 3 \cdot x \bmod 4$.



$3 \cdot x \equiv 2 \pmod 4$

$x \equiv 2 \pmod 4$

$3 \cdot y \equiv 1 \pmod 4$

$y \equiv 3 \pmod 4$

Multiplications as functions

Let's view multiplication by 2 as a function, i.e., let $f(x) = 2 \cdot x \bmod 4$.

Let's also see $g(x) = 3 \cdot x \bmod 4$.

Which functions have inverses? f, g

Multiplicative inverses (standard arithmetic)

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We are looking to a number x such that $2 = 5x$. How can we do that?

$5x = 2$
$x = 2/5$

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By dividing on both sides with 5:

$$2/5 = 5x/5 = x,$$

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By dividing on both sides with 5:

$$2/5 = 5x/5 = x,$$

or equivalently, by multiplying with $(1/5) = 5^{-1}$:

$$2 \cdot 5^{-1} = 5x \cdot 5^{-1} = x \cdot 5 \cdot 5^{-1} = x \cdot 1 = x.$$

Here 5^{-1} is a multiplicative inverse of 5.

Multiplicative inverses (modular arithmetic)

You can do the same thing in modular arithmetic. Let the modulus be $m = 7$. Note that

$$5 \cdot 3 \equiv 15 \equiv 1 \pmod{7}.$$

Therefore, $5^{-1} \equiv 3 \pmod{7}$.

$5y \equiv 4 \pmod{7}$

$y \equiv 5^{-1} \cdot 5y \equiv 3 \cdot 4 \equiv 5 \pmod{7}$

Multiplicative inverses (modular arithmetic)

You can do the same thing in modular arithmetic. Let the modulus be $m = 7$. Note that

$$5 \cdot 3 \equiv 15 \equiv 1 \pmod 7.$$

$2/5 \equiv 6 \pmod 7$

Therefore, $5^{-1} \equiv 3 \pmod 7$.

To find $2/5$, we can view our goal as to find the value of x such that

$$2 \equiv 5x \pmod 7.$$

We can multiply both sides with $5^{-1} \equiv 3$ to get

$$2 \cdot 5^{-1} \equiv 2 \cdot 3 \equiv 6 \equiv 5^{-1} \cdot 5x \equiv x \pmod 7.$$

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To find $2/5$, we can view our goal as to find the value of x such that

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We can multiply both sides with $5^{-1} \equiv 3$ to get

$$2 \cdot 5^{-1} \equiv 2 \cdot 3 \equiv 6 \equiv 5^{-1} \cdot 5x \equiv x \pmod{7}.$$

Let's check:

$$5 \cdot 6 \equiv 30 \equiv 2 \pmod{7},$$

as requied.

Multiplicative inverse modulo m

$a^{-1} \pmod m$

Definition

The multiplicative inverse modulo m of a, denoted by a^{-1}, is an integer such that

$$a \cdot a^{-1} \equiv 1 \pmod m.$$

Multiplicative inverse modulo 11

Let's try to figure out multiplicative inverse of every integer modulo 11.

a	$a^{-1} \pmod{11}$
1	1
2	6
3	4
4	3
5	9
6	2
7	8
8	7
9	5
10	10

$7x \equiv 6 \pmod{11}$

$x \equiv 8 \cdot 7x \equiv 8 \cdot 6 \equiv 4 \pmod{11}$

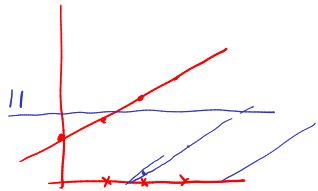
$7a + m \equiv 9 \pmod{11}$

$3a + m = 4 \pmod{11}$

$4a \equiv 5 \pmod{11}$

$a \equiv 3 \cdot 4a \equiv 3 \cdot 5 \equiv 15 \equiv 4 \pmod{11}$

Example: secret sharing



- Think of a secret number $m \in \{0, 1, \ldots, 10\}$.
- Pick a random number $a \in \{1, 2, \ldots, 10\}$.
- Your straight line function $f(x) = (ax + m) \bmod 11$.
- We will generate 3 points from f and give them to 3 of your friends, each with only 1 point. Pick 3 numbers x_1, x_2, x_3 from $\{1, 2, \ldots, 10\}$.
- Let's compute

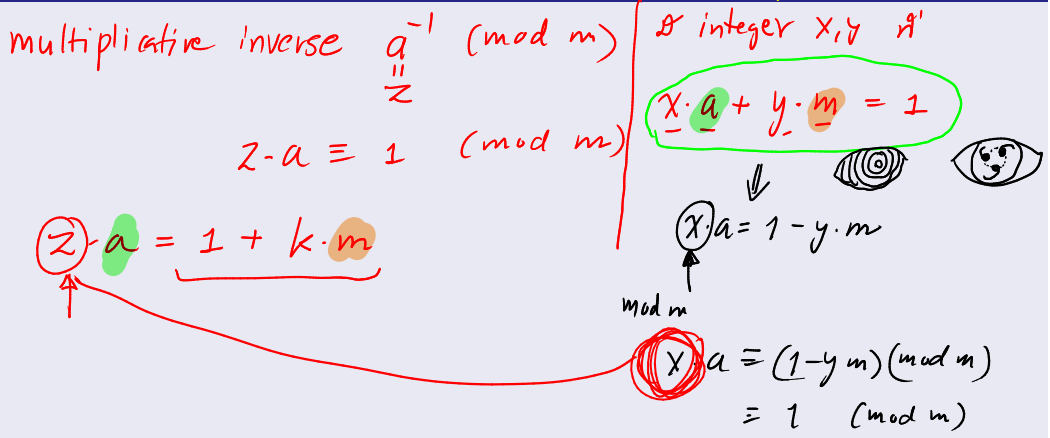
$$(x_1, f(x_1)), \quad (x_2, f(x_2)), \quad (x_3, f(x_3)).$$

- Give them to 3 of your friends and challenge them to form a group of 2 people and figure out your number m.

Theorem 1

An integer a has a multiplicative inverse modulo m iff $gcd(a, m) = 1$.

Proof.



Theorem 1

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Proof.

($\Leftarrow$) Recall that there exist integers x and y such that

$$x \cdot a + y \cdot m = gcd(a, m) = 1.$$

Thus, $(x \cdot a + y \cdot m) \bmod m = x \cdot a \bmod m = 1 \bmod m$, i.e., $x \cdot a \equiv 1 \pmod m$. Therefore x is the inverse.

0

ex $z^{-1} \pmod{17}$

$gcd(17, 2)$

$= gcd(2, 17 \bmod 2)$

1

$0 \cdot 2 + 1 \cdot 1 = 1$

a' b'

$1 \cdot 17 + \left(0 - 1 \cdot \left\lfloor \frac{17}{2} \right\rfloor\right) \cdot 2$

$1 \cdot 17 - 8 \cdot 2 = 1$

$-8 \cdot 2 \equiv 1 \pmod{17}$

```
Algorithm Euclid(x,y):
  if x mod y == 0:
    return y,   0   ,  1
  else:
    g, a', b' = Euclid(y, x mod y)
    a = b'
    b = a' - b'·(⌊x/y⌋)
    return g, a, b
```

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($\Leftarrow$) Recall that there exist integers x and y such that

$$x \cdot a + y \cdot m = gcd(a, m) = 1.$$

Thus, $(x \cdot a + y \cdot m) \bmod m = x \cdot a \bmod m = 1 \bmod m$, i.e., $x \cdot a \equiv 1 \pmod{m}$. Therefore x is the inverse.
($\Rightarrow$) Let $r = gcd(a, m)$. Suppose that b is the multiplicative inverse of a modulo m, i.e., we have that

$$b \cdot a \equiv 1 \pmod{m},$$

Thus, $ba \bmod m = 1 \bmod m = 1$, i.e., there exists an integer q such that

$$ba = qm + 1,$$

or $ba - qm = 1$. However, r since $r|a$ and $r|m$, r also divides $bd - qm$ and 1. But it $r \not| 1$ because $r > 1$ and we have the contradiction. □

Examples: division in modular arithmetic

Since the requirement for an existance of a^{-1} modulo m is that $gcd(a, m) = 1$, if we let m be a prime number, every a which is not a multiple of m has an inverse.
Can you solve this equation?

$$4x + 9 \equiv 0 \pmod{11}.$$

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Since the requirement for an existance of a^{-1} modulo m is that $gcd(a, m) = 1$, if we let m be a prime number, every a which is not a multiple of m has an inverse.
Can you solve this equation?

$$4x + 9 \equiv 0 \pmod{11}.$$

We can even perform gaussian elimination (*which is very useful later*):

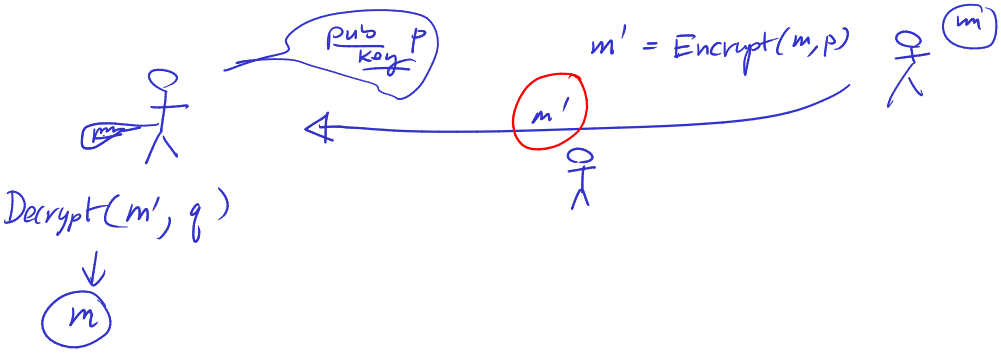
$$\begin{cases} 2x + y \equiv 3 \pmod{7} \\ x + 3y \equiv 5 \pmod{7} \end{cases}$$

Public-key cryptography

~~private-key cryptography~~

public key

private key



RSA Rivest Shamir Adelman

จน. เฉพาะ; ยิ่งใหญ่ p, q — 2048 บิต

n

pub: (e, n)

private: (d, n)

หา จน. เพื่อ d, e

- เลือก e (65537)
- หา d ที่ $d \cdot e \mod (p-1)(q-1) = 1$

↳ multiplicative inverse ของ e mod $(p-1)(q-1)$

def: $E(m) = m^d \mod n$

$D(k) = k^e \mod n$

Goal: message m

$(m^d \bmod n)^e \bmod n$

$= (m^d)^e \bmod n$

$= m$

Quick recap: RSA

- Private key: (d, n), Public key: (e, n)
- Encryption $E(m) = m^e \bmod n$, Decryption: $D(w) = w^d \bmod n$.
- Goal: Select e, d, n such that $D(E(m)) = m^{ed} \bmod n = m$.

Quick recap: RSA

- Private key: (d, n), Public key: (e, n)
- Encryption $E(m) = m^e \bmod n$, Decryption: $D(w) = w^d \bmod n$.
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- Pick two primes p and q. Let $n = pq$.
- Pick e (usually a small number)
- Pick d such that $d = e^{-1} \pmod{(p-1)(q-1)}$, i.e., $ed \equiv 1 \pmod{(p-1)(q-1)}$, or

$$ed = k \cdot (p-1)(q-1) + 1,$$

  for some integer k.
- What is $m^{ed} \bmod n$?

$m^{ed} \equiv m^{k\cdot(p-1)(q-1)+1} \pmod n$

$\equiv m^{k(p-1)(q-1)} \cdot m \pmod n$

$1 \equiv$

What's next?

- We will prove Fermat's Little Theorem and show how to efficiently test if a number is prime.
- We will also use Fermat's Little Theorem to prove the correctness of RSA.
- Modular arithmetic is also key to our usage of polynomials to perform secret sharing and error correcting codes, because now we can do Gaussian elimination using only integers.