## 01204211 Discrete Mathematics Lecture 5a: Mathematical Induction 4

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#### Review: Mathematical Induction

Suppose that you want to prove that property P(n) is true for every natural number n.

Suppose that we can prove the following two facts:

Base case: P(1)

**Inductive step:** For any  $k \ge 1$ ,  $P(k) \Rightarrow P(k+1)$ 

The **Principle of Mathematical Induction** states that P(n) is true for every natural number n.

The assumption P(k) in the inductive step is usually referred to as the Induction Hypothesis.

## Review: Strengtening the Induction Hypothesis

To prove the following theorem with induction:

#### Theorem 1

For any integer 
$$n \ge 1$$
,  $\frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \cdots + \frac{1}{n^2} < 2$ .

You need to prove the following "stronger" version:

#### Theorem 2

For any integer 
$$n \ge 1$$
,  $\frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} + \dots + \frac{1}{n^2} \le 2 - \frac{1}{n}$ .

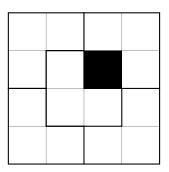


#### A Lesson learned

- Is a stronger statement easier to prove?
- ▶ In this case, the statement is indeed stronger, but the induction hypothesis gets stronger as well. Sometimes, this works out nicely.

## L-shaped tiles $(1)^1$

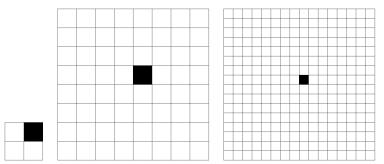
A 4x4 area with a hole in the middle can be tiled with L-shaped tiles.



¹This section is from Berkeley CS70 lecture notes. <□ > <♂ > < ≧ > < ≧ > < ≥ < ∞ < ○ <

# L-shaped tiles (2)

This is true for 2x2 area, 8x8 area, even 16x16 area.



This motivates us to try to prove that it is possible to use L-shaped tiles to tile a  $2^n \times 2^n$  area.

#### Theorem 3

For integer  $n \ge 1$ , an area of size  $2^n \times 2^n$  with one hole in the middle can be tiled with L-shaped tiles.

Proof: We prove by induction on n.

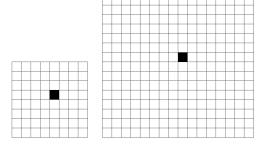
Base case: For  $n=1,\,2^1\times 2^1$  area with a hole in the middle can be tiled.

**Inductive step:** Assume that for  $k \ge 1$ , an  $2^k \times 2^k$  area with a hole in the middle can be tiled. We shall prove the statement for n=k+1, i.e., that an  $2^{k+1} \times 2^{k+1}$  area with one hole in the middle can be tiled.

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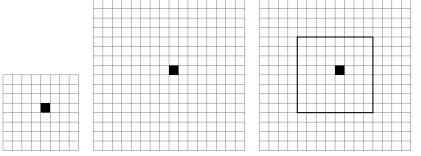
Proof: (cont.)

Let's see the Induction Hypothesis and the goal:



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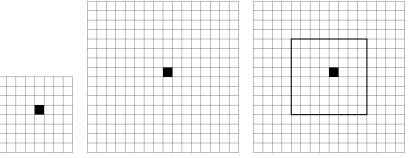
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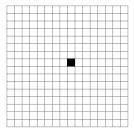
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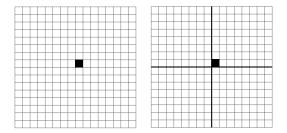
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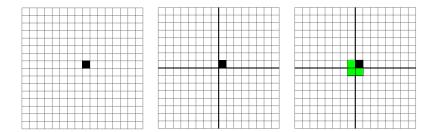
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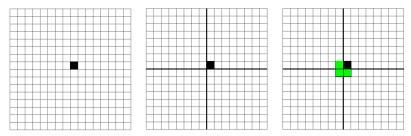


With the current form of the Induction Hypothesis, this is probably the way to use it. But it seems hard to go further with this approach....

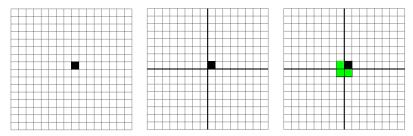




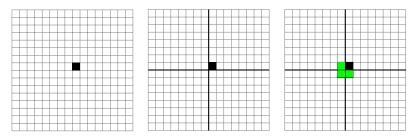




The last step seems nice, because it shows how we can solve the problem in the  $2^{k+1}\times 2^{k+1}$  area with 4 problems in the  $2^k\times 2^k$  areas.

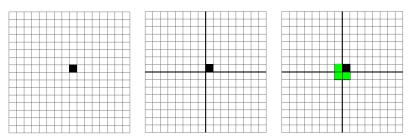


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Current Inductive Hypothesis: Assume that for  $k \ge 1$ , an  $2^k \times 2^k$  area with "a hole in the middle" can be tiled.



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**A Stronger Inductive Hypothesis:** Assume that for  $k \geq 1$ , an  $2^k \times 2^k$  area with one hole can be tiled.

Theorem: For integer  $n \ge 1$ , an area of size  $2^n \times 2^n$  with one hole can be tiled with L-shaped tiles.

Proof: We prove by induction on n.

Base case: For n=1,  $2^1 \times 2^1$  area with one hole can be tiled;

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**Base case:** For n=1,  $2^1 \times 2^1$  area with one hole can be tiled; there are 4 cases shown below.





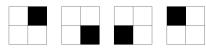




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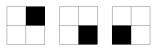


**Inductive step:** Assume that for  $k \geq 1$ , an  $2^k \times 2^k$  area with one hole can be tiled. We shall prove the statement for n=k+1, i.e., that an  $2^{k+1} \times 2^{k+1}$  area with one hole can be tiled.

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#### Theorem 4

If P(1) and for any integer  $k \ge 1$ ,  $P(k) \Rightarrow P(k+1)$ , then P(n) for all natural number n.

Proof.

<sup>&</sup>lt;sup>2</sup>This section is from Berkeley CS70 lecture notes.

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Therefore, P(n) is true for every positive integer n. Is this proof correct?

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### The Well-Ordering Property

The proof of the Principle of Mathematical Induction depends on the following axiom of natural numbers  $\mathbb{N}$ :

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▶ Previously, we use the well-ordering property of natural numbers to prove the Principle of Mathematical Induction, but it turns out that we can use the induction to prove the well-ordering property as well. Therefore, we can take one as an axiom, and use it to prove the other.