

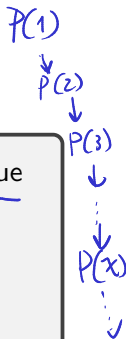
01204211 Discrete Mathematics

Lecture 4b: Mathematical Induction 2

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Review: Mathematical Induction



Suppose that you want to prove that property $P(n)$ is true for every natural number n .

Suppose that we can prove the following two facts:

→ **Base case:** $P(1)$

→ **Inductive step:** For any $k \geq 1$, $P(k) \Rightarrow P(k+1)$

induction

The **Principle of Mathematical Induction** states that $P(n)$ is true for every natural number n .

The assumption $P(k)$ in the inductive step is usually referred to as the **Induction Hypothesis**.

Example 1

Base case $P(1)$

$$1^2 = \frac{1}{6}(1+1)(2+1)$$

$$* P(2): 1^2 + 2^2 = 5 \quad \frac{2}{6}(3)(5)$$

Theorem: For every natural number n ,
 $\sum_{i=1}^n i^2 = \frac{n}{6}(n+1)(2n+1)$

Proof: We prove by induction. The property that we want to prove $P(n)$ is " $\sum_{i=1}^n i^2 = \frac{n}{6}(n+1)(2n+1)$."

Inductive step Assume in $P(k)$ and $P(k+1)$ is true for $k \geq 1$

Induction hypothesis: $P(k): \sum_{i=1}^k i^2 = \frac{k}{6}(k+1)(2k+1)$

Assume $P(k+1): \sum_{i=1}^{k+1} i^2 = \frac{(k+1)}{6}(k+2)(2k+3)$

$$\begin{aligned} \text{Then } \sum_{i=1}^{k+1} i^2 &= \left(\sum_{i=1}^k i^2 \right) + (k+1)^2 = \frac{k}{6}(k+1)(2k+1) + (k+1)^2 \\ &= \left(\frac{k+1}{6} \right) (k(2k+1) + 6(k+1)) = \left(\frac{k+1}{6} \right) (2k^2 + k + 6k + 6) \\ &= \left(\frac{k+1}{6} \right) (k+2)(2k+3) \quad \text{which is } P(k+1) \end{aligned}$$

By Principle of M.I. we have $P(n)$ is true for $n \geq 1$

Example 1

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Proof: We prove by induction. The property that we want to prove $P(n)$ is " $\sum_{i=1}^n i^2 = \frac{n}{6}(n+1)(2n+1)$."

Base case: We can plug in $n = 1$ to check that $P(1)$ is true:
 $1^2 = \frac{1}{6}(1+1)(2 \cdot 1 + 1)$.

Inductive step: We assume that $P(k)$ is true for $k \geq 1$ and show that $P(k+1)$ is true.

We first assume the Induction Hypothesis $P(k)$:

$$\sum_{i=1}^k i^2 = \frac{k}{6}(k+1)(2k+1)$$

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Example 1 (cont.)

Let's show $P(k+1)$. We write $\sum_{i=1}^{k+1} i^2 = \left(\sum_{i=1}^k i^2\right) + (k+1)^2$.

Using the Induction Hypothesis, we know that this is equal to

$$\begin{aligned}(k/6)(k+1)(2k+1) + (k+1)^2 &= \frac{(k+1)}{6}(k(2k+1) + 6(k+1)) \\ &\quad \text{(In this step, we factor out } (k+1)/6\text{)} \\ &= \frac{(k+1)}{6}(2k^2 + 7k + 6) \\ &= \frac{(k+1)}{6}((k+1) + 1)(2(k+1) + 1).\end{aligned}$$

This implies $P(k+1)$ as required.

From the Principle of Mathematical Induction, this implies that $P(n)$ is true for every natural number n . ■

Not an example (1)

$P(n) =$
set ของวัว ใน n
ที่มีสีเหมือนกัน n ตัว
คือทุกตัว $n=3$ สีดำดำดำ

Theorem 1

For any set of cows, all cows have the same color.

Proof.

We prove by induction on the size n of the set of cows.

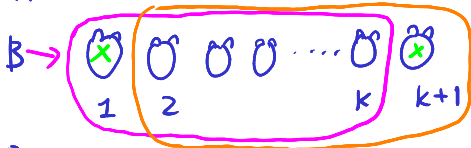
Base case: มีวัว $n=1$ ✓

Inductive step: ๑: มีวัว $n > 1$. $P(k) \Rightarrow P(k+1)$ ✓

พิจารณาเซตของวัวที่มี $k+1$ ตัว ให้ $B = A - \{\text{วัวตัวที่ } k+1\}$

มีวัว
 $|B| = k$, จาก $P(k)$

ทุกตัวใน B
มีสีเหมือนกัน



ให้ $C = A - \{\text{วัวตัวที่ } 1\}$
มีวัว $|C| = k$,
ทุกตัวใน C มี
สีเหมือนกัน

วัวตัวที่ 1 มีสีเดียวกับวัวตัวที่ 2 เพราะ \in เซต B และ
วัวตัวที่ k+1 มีสีเดียวกับวัวตัวที่ 2 (อยู่ใน C) \Rightarrow วัว 1 กับ วัว $k+1$ มี
สีเหมือนกัน

Not an example (1)

Theorem 1

For any set of cows, all cows have the same color.

Proof.

We prove by induction on the size n of the set of cows.

Base case: For $n = 1$, clearly for any set of a single cow, every cow in the set has the same color.

Not an example (1)

Theorem 1

For any set of cows, all cows have the same color.

Proof.

We prove by induction on the size n of the set of cows.

Base case: For $n = 1$, clearly for any set of a single cow, every cow in the set has the same color.

Inductive step: Suppose that for every set of size k of cows, all cows in the set have the same color.

We will show that every set of size $k + 1$ of cows, all cows in this set have the same color.

Not an example (2)

Inductive step (cont.): Consider set A of $k + 1$ cows.

Not an example (2)

Inductive step (cont.): Consider set A of $k + 1$ cows.

พิจารณาเซต A ของวัวที่มี $k+1$ ตัว ให้ $B = A - \{\text{วัวตัวที่ } k+1\}$

ให้ $|B| = k$, จาก $P(k)$

พิจารณาเซต B มีสีเดียวกัน

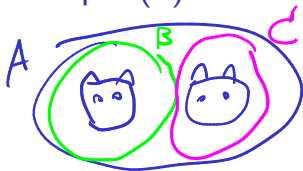
ให้ $C = A - \{\text{วัวตัวที่ } 1\}$

ให้ $|C| = k$,
พิจารณาเซต C มีสีเดียวกัน

วัวตัวที่ 1 มีสีเดียวกับวัวตัวที่ 2 เพราะ $2 \in B$ และ 1 และ 2 มีสีเดียวกัน
วัวตัวที่ $k+1$ มีสีเดียวกับวัวตัวที่ 2 ($2 \in C$) \Rightarrow วัว 1 วัว 2 และ วัว $k+1$ มีสีเดียวกัน

Because we have established that the base case and the inductive step is true, we can conclude that for any set of cows, all cows have the same color. ■

Not an example (3)



$$P(1) \not\Rightarrow P(2) \Rightarrow P(3) \Rightarrow \dots$$

$$\underline{P(k) \Rightarrow (P(k+1))}$$

Clearly the following theorem cannot be true.

เมื่อ $k=1$

Theorem 2

For any set of cows, all cows have the same color.

What is wrong with its proof based on mathematical induction?

พิจารณาเซตของวัวที่มี $k+1$ ตัว $B = A - \{\text{วัวตัวที่ } k+1\}$

เมื่อ $|B| = k$, จาก $P(k)$

วัวทุกตัวใน B มีสีเดียวกัน

ใน B → $k+1$

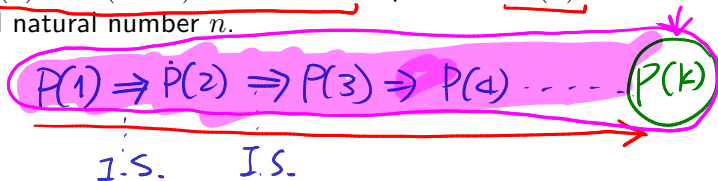
นั่น $C = A - \{\text{วัวตัวที่ } 1\}$

เมื่อ $|C| = k$,
วัวทุกตัวใน C มีสีเดียวกัน

วัวตัวที่ 1 มีสีเดียวกับวัวตัวที่ 2 (เพราะอยู่ใน B) และวัวตัวที่ 2 มีสีเดียวกับวัวตัวที่ $k+1$ (เพราะอยู่ใน C) \Rightarrow วัวตัวที่ 1 มีสีเดียวกับวัวตัวที่ $k+1$

Unused facts

- Let's informally think about how proving $P(1)$ and $P(k) \Rightarrow P(k+1)$ for all $k \geq 1$ implies that $P(n)$ is true for all natural number n .



$$P(k) \Rightarrow P(k+1)$$

Unused facts

- ▶ Let's informally think about how proving $P(1)$ and $P(k) \Rightarrow P(k+1)$ for all $k \geq 1$ implies that $P(n)$ is true for all natural number n .
- ▶ One may notice that when we prove a statement $P(n)$ for all natural number n by induction, during the inductive step where we want to show $P(k+1)$ from $P(k)$, we usually have that $P(1), P(2), \dots, P(k)$ is true at hands as well.

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- ▶ Then why don't we use them as well?

Strong Mathematical Induction

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Strong Mathematical Induction

$$\underline{R} \Rightarrow S$$

Strong Induction

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Suppose that we can prove the following two facts: **Base case:** $P(1)$

Inductive step: For any $\underline{k \geq 1}$,

$$\underbrace{P(1) \wedge P(2) \wedge \cdots \wedge P(k)}_{\text{Inductive Hypothesis}} \Rightarrow \underline{P(k+1)}.$$

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Then $P(n)$ is true for every natural number n .

Example 2

Theorem: For any integer $n \geq 4$, one can use only 2-baht coins and 3-baht coins to obtain exactly n baht.

Proof: We prove by strong induction on n .

Base cases: For $n = 4$, we can use two 2-baht coins. For $n = 5$, we can use one 2-baht coin and one 3-baht coin.

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Inductive step: Assume that for $k \geq 5$, we can obtain exactly ℓ baht, for $4 \leq \ell \leq k$, using only 2-baht and 3-baht coins. We will show how to obtain a set of $k + 1$ baht.

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Since $k \geq 5$, we have that $k - 1 \geq 4$. Therefore from the Induction Hypothesis, we can use only 2-baht coins and 3-baht coins to form a set of coins of total value $k - 1$ baht.

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From the Principle of Strong Mathematical Induction, we conclude that the theorem is true. ■

Is strong induction more powerful?

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- ▶ In fact, if you can prove that $P(n)$ is true for all natural number n with strong induction, you can always prove it with mathematical induction.
- ▶ Hint: Let $Q(n) = P(1) \wedge P(2) \wedge \cdots \wedge P(n)$.

