01204211 Discrete Mathematics Lecture 8b: Modular arithmetic

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September 29, 2022

Quick check 1

If a|m and b|m, can we say that ab|m? Prove this fact or provide a counter example.

Quick check 2

If a|m, b|m, and $a \neq b$ are both prime, can we say that ab|m? Prove this fact or provide a counter example.

Prime factorization

One useful fact that we use over and over again is the following.

Unique Factorization (or Fundamental Theorem of Arithmetic)

Every integer greater than 1 can be written *uniquely* as a product of prime numbers (up to the order of factors).

Examples:

- $ightharpoonup 10 = 2 \cdot 5$
- **▶** 13 = 13
- $112 = 2 \cdot 2 \cdot 2 \cdot 2 \cdot 7 = 2^4 \cdot 7$

There are 3 clocks. At this moment, all three clocks ring at the same time. The first clock rings every 3 hours, the second clock rings every 4 hours, and the third clock rings every 10 hours. How long do you have to wait until you would hear all clocks ring a the same time again?

You have a large water container and two smaller buckets. The first bucket carries 3 litres of water and the second bucket carries 5 litres of water.

Can you put exactly 1 litre of water in the water container?

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What is the minimum volume of water you can exactly put in the water container?

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In general if you have two buckets of volumes \boldsymbol{x} and \boldsymbol{y} , the amount that you can exactly make must be in the form of

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for some integers x and y. (Note that x and y may be negative.)

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In general if you have two buckets of volumes \boldsymbol{x} and \boldsymbol{y} , the amount that you can exactly make must be in the form of

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for some integers x and y. (Note that x and y may be negative.) Do you see why the sum must be divisible by any common divisor of x and y?

Useful fact

For any integer \boldsymbol{x} and \boldsymbol{y} , consider the term

$$a \cdot x + b \cdot y$$
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for some integer \boldsymbol{a} and \boldsymbol{b} .

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For any integer x and y, consider the term

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for some integer a and b.

When the term is non-zero, it must be divisible by gcd(x,y), so it has to be at least gcd(x,y).

It turns out that you can actually attain that value, i.e., there exist a pair of integer a and b such that

$$a \cdot x + b \cdot y = \gcd(x, y).$$



Finding a and b: Extended Euclid Algorithm

We will modify the Euclid algorithm so that it also returns a and b together with gcd(x,y).

```
Algorithm Euclid(x,y):
  if x \mod y == 0:
    return y,
  else:
    g, a', b' = Euclid(y, x mod y)
    a =
    b =
    return g, a, b
```

Notes:

We have a' and b' such that

$$a' \cdot y + b' \cdot (x \bmod y) = g.$$

What day is it today?

What day is it today? Thursday.

What day is it today? Thursday. What day is 3 days after today?

What day is it today? Thursday. What day is 3 days after today? Sunday.

What day is it today? Thursday. What day is 3 days after today? Sunday. What day is 20 days after today?

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What day is it today? Thursday.
What day is 3 days after today? Sunday.
What day is 20 days after today? Wednesday.
What day is 10 days before today? Monday.

Suppose that it is $1\ {\rm o'clock}.$

Suppose that it is 1 o'clock. What time is the next 5 hours?

Suppose that it is 1 o'clock. What time is the next 5 hours? 6 o'clock.

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Note that when you view integers under the lense of modulus 7, these numbers

$$\dots, -19, -12, -5, 2, 9, 16, 23, \dots$$

are essentially the same.



Properties (1)

 $a \bmod m = b \bmod m$, if and only if m|a - b.

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Proof.

 (\Rightarrow) Let $r = a \mod m$. We can write

$$a = qm + r$$
,

and

$$b = pm + r,$$

for some integers q and p. Thus, we have

$$a - b = qm + r - pm - r = (q - p)m.$$

Therefore m|a-b.



Properties (2)

- $(a+b) \bmod m = ((a \bmod m) + (b \bmod m)) \bmod m$
- $(a-b) \bmod m = ((a \bmod m) (b \bmod m)) \bmod m$
- $(a \cdot b) \bmod m = ((a \bmod m) \cdot (b \bmod m)) \bmod m$

Congruences

Definition (congruences)

For an integer m>0, if integers a and b are such that

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we write

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We also have that

$$a \equiv b \pmod{m} \Leftrightarrow m|(a-b)$$



Congruences: properties (1)

- (reflexivity) $a \equiv a \pmod{m}.$
- (symmetry) $a \equiv b \pmod{m}$ implies $b \equiv a \pmod{m}$.
- ▶ (transitiviey) $a \equiv b \pmod{m}$ and $b \equiv c \pmod{m}$ implies $a \equiv c \pmod{m}$.

If we have that

$$a \equiv b \pmod{m}$$
,

and

$$c \equiv d \pmod{m}$$
,

then

- $ightharpoonup a + c \equiv b + d \pmod{m}$
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What is missing here?

If we have that

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We can pretty much think of this "congruence" as a normal equation.

What is missing here? Division!

Also, we wish we can do "cancellation", i.e., if

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Also, we wish we can do "cancellation", i.e., if

$$xa \equiv xb \pmod{m}$$
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then $a \equiv b \pmod{m}$. BUT THIS IS NOT ALWAYS TRUE. Let's see the following example:

$$2 \cdot 1 \equiv 2 \cdot 3 \pmod{4},$$

but

$$1 \not\equiv 3 \pmod{4}$$
.

Multiplications as functions

Let's view multiplication by 2 as a function, i.e., let $f(x) = 2 \cdot x \mod 4$.

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Which functions have inverses?

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We are looking to a number x such that 2=5x. How can we do that?

By dividing on both sides with 5:

$$2/5 = 5x/5 = x,$$

or equivalently, by multiplying with $(1/5) = 5^{-1}$:

$$2 \cdot 5^{-1} = 5x \cdot 5^{-1} = x \cdot 5 \cdot 5^{-1} = x \cdot 1 = x.$$

Here 5^{-1} is a multiplicative inverse of 5.

Multiplicative inverses (modular arithmetic)

You can do the same thing in modular arithmetic. Let the modulus be $m=7. \ \mathrm{Note}$ that

$$5 \cdot 3 \equiv 15 \equiv 1 \pmod{7}.$$

Therefore, $5^{-1} \equiv 3 \pmod{7}$.

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To find 2/5, we can view our goal as to find the value of \boldsymbol{x} such that

$$2 \equiv 5x \pmod{7}$$
.

We can multiply both sides with $5^{-1} \equiv 3$ to get

$$2 \cdot 5^{-1} \equiv 2 \cdot 3 \equiv 6 \equiv 5^{-1} \cdot 5x \equiv x \pmod{7}.$$

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Let's check:

$$5 \cdot 6 \equiv 30 \equiv 2 \pmod{7}$$
,

as requied.



Multiplicative inverse modulo m

Definition

The multiplicative inverse modulo m of a, denoted by a^{-1} , is an integer such that

$$a \cdot a^{-1} \equiv 1 \pmod{m}$$
.

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An integer a has a multiplicative inverse modulo m iff $\gcd(a,m)=1$.

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Proof.

 (\Leftarrow) Recall that there exist integers x and y such that

$$x \cdot a + y \cdot m = \gcd(a, m) = 1.$$

Thus, $(x\cdot a+y\cdot m) \bmod m=x\cdot a \bmod m=1 \bmod m$, i.e., $x\cdot a\equiv 1 \pmod m$. Therefore x is the inverse.

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Proof.

 (\Leftarrow) Recall that there exist integers x and y such that

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Thus, $(x \cdot a + y \cdot m) \mod m = x \cdot a \mod m = 1 \mod m$, i.e., $x \cdot a \equiv 1 \pmod m$. Therefore x is the inverse.

 (\Rightarrow) Let r=gcd(a,m). Suppose that b is the multiplicative inverse of a modulo m, i.e., we have that

$$b \cdot a \equiv 1 \pmod{m}$$
,

Thus, $ba \mod m = 1 \mod m = 1$, i.e., there exists an integer q such that

$$ba = qm + 1,$$

or ba-qm=1. However, r since r|a and r|m, r also divides bd-qm and 1.

But it $r \not| 1$ because r > 1 and we have the contradiction.



Examples: division in modular arithmetic

Since the requirement for an existance of a^{-1} modulo m is that gcd(a,m)=1, if we let m be a prime number, every a which is not a multiple of m has an inverse.

Can you solve this equation?

$$4x + 9 \equiv 0 \pmod{11}.$$

Examples: division in modular arithmetic

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Can you solve this equation?

$$4x + 9 \equiv 0 \pmod{11}.$$

We can even perform gaussian elimination (which is very useful later):

$$2x + y \equiv 3 \pmod{7}$$
$$x + 3y \equiv 5 \pmod{7}$$

Quick recap: RSA

- Private key: (e, n), Public key: (d, n)
- ▶ Encryption $E(m) = m^e \mod n$, Decryption: $D(w) = w^d \mod n$.
- ▶ Goal: Select e, d, n such that $D(E(m)) = m^{ed} \mod n = m$.

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- Private key: (e, n), Public key: (d, n)
- ▶ Encryption $E(m) = m^e \mod n$, Decryption: $D(w) = w^d \mod n$.
- ▶ Goal: Select e, d, n such that $D(E(m)) = m^{ed} \mod n = m$.
- Pick two primes p and q. Let n = pq.
- Pick e (usually a small number)
- Pick d such that $d=e^{-1} \pmod{(p-1)(q-1)}$, i.e., $ed \equiv 1 \pmod{(p-1)(q-1)}$, or

$$ed = k \cdot (p-1)(q-1) + 1,$$

for some integer k.

▶ What is $m^{ed} \mod n$?

What's next?

- ► We will prove Fermat's Little Theorem and show how to efficiently test if a number is prime.
- We will also use Fermat's Little Theorem to prove the correctness of RSA.
- Modular arithmetic is also key to our usage of polynomials to perform secret sharing and error correcting codes, because now we can do Gaussian elimination using only integers.