# 01204211 Discrete Mathematics Lecture 2a: Quantifiers ∀ ∃

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morassiona "predicate"

## Review (1)

$$5+7=12$$
  $P \Rightarrow P \Rightarrow P$ 

- A proposition is a statement which is either true or false.
- ightharpoonup We can use variables to stand for propositions, e.g., P= "today is Tuesday".
- We can use connectives to combine variables to get propositional forms.
  - **Conjunction:**  $P \wedge Q$  ("P and Q"),
  - **Disjunction:**  $P \checkmark Q$  ("P or Q"), and
  - **Negation:**  $\neg P$  ("not P")

  - **Equivalence:**  $P(\Leftrightarrow)Q$  ("P if and only if Q")



```
Review (2): Testing primes
```

6065 [x] tlook

Consider the following code.

```
Algorithm CheckPrime2(n): // Input: an integer n
if n <= 1:
    return False
let s = square root of n
i = 2
while i <= s:
    if n is divisible by i:
        return False
    i = i + 1
return True
```

How fast can it run? Note that  $s=\sqrt{n}$ ; therefore, it takes time approximately proportional to  $\sqrt{n}$  to run.

Ok, it should be faster. But is it correct?

## The goals

Let's recall what we are trying to do.

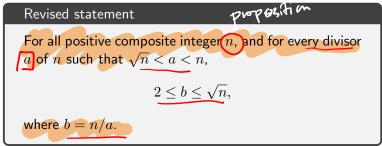
Original goal: To show that Algorithm CheckPrime2 is correct.

**Current (sub) goal:** Consider a positive composite n and its positive divisor a, where  $a > \sqrt{n}$ . Let b = n/a. We want to show that  $2 \le b \le \sqrt{n}$ .

[25 <u>b</u> 5 <u>m</u>]

## The (sub) goal

- ▶ Current (sub) goal: Consider a positive composite n and its positive divisor a, where  $a > \sqrt{n}$ . Let b = n/a. We want to show that  $2 \le b \le \sqrt{n}$ .
- ▶ We can be more specific about what values of *n* and *b* that we want to consider.



Note that this revised statement is now "quantified," that is, every variable in the statement has specific scope. Now the statement is either true or false.



#### **Predicates**



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#### **Predicates**

- In many cases, the statement we are interested in contains variables.
- For example, "x is even," "p is prime," or "s is a student."
- As we previously did with propositions, we can use <u>variables</u> to represent these statements. E.g.,

friend(x)

- let  $\underline{E(x)} \equiv "x \text{ is even"},$
- let  $\underline{P(y)} \equiv "y \text{ is prime, and}$
- let  $\overline{S(w)} \equiv "w \text{ is a student.}$

We call E(x), P(y), and S(w) predicates (You can think of predicates as statements that may be true of false depending on the values of its <u>variables</u>.)

## Quantifiers (1)

"5 1] 1100 of "-fake

- As we note before, these predicates are not propositions. But if we know the values of their variables, then they becomes propositions. For example, if we let x=5, then E(5) is a proposition which is false. Also, P(7) is true.
- Since the truth values of predicates depend on the assignments of their variables, we can put *quantifiers* to specify the scopes of these variables and how to interprete the truth values of the predicates over these values.

# Quantifiers (2): universal quantifiers ∀



▶ Let  $A = \{2, 4, 6, 8\}$ .



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$$(\forall x \in A)E(x).$$

The quantifier  $\forall$  is called a universal quantifier. (We usually pronounce "for all x", or "for every x.")



- Again, let  $A = \{2, 4, 6, 8\}$ .
- Note that  $\underline{P(2)}$  is true. This means that  $\underline{P(y)}$  is true for some  $y \in A$ .

```
There
exists y , n
such that P(y)
```

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▶ The quantifier  $\exists$  is called an existential quantifier. (We usually pronounce "for some x", or "there exists x.")

When the universe A is clear, we can leave it out and just write  $\forall x E(x)$  or  $\exists y P(y)$ .

## The main goal

Let's try to be more specific about our main goal:

Algorithm CheckPrime2 is correct.

- ► Can we re-write this statement so that the input/output of the algorithm are explicit?
- Note that the set of its input n is an integer. Thus, we are interested in every  $n \in \mathbb{Z}$ , where  $\mathbb{Z}$  denote the set of all integers.
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where  $C(n) \equiv$  "CheckPrime2(n) returns True", and  $P(n) \equiv$  "n is a prime."

## Quantified propositions with more than one variables



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Fake  $\forall x \forall y (x = y)$ 

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When you have many quantifiers, we can interprete the statement by nesting the quantifiers. E.g,

$$\exists x \forall y P(x,y) \equiv \exists x (\forall y (P(x,y))).$$
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Also note that usually,  $\exists x \forall y P(x,y) \not\equiv \forall y \exists x P(x,y)$ .



#### Quick check 4

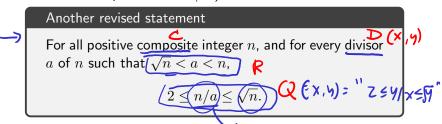
We will consider the universe to be "everything" Consider the following statements. Define appropriate predicates and rewrite them using the defined predicates and quantifiers. (Note: the predicates may have more than one variables.)

- ► Every human must die.  $\forall x (h uman(x) \Rightarrow die(x))$
- Some animal eats other animals.  $\exists x \exists y (A(x) \land A(y) \land E(x))$
- If a student works hard, that student will be successful.  $\Lambda O(x_1)$
- Everyone has someone that care about him or her.

$$\forall_x (\text{student}(x) \land \text{wh}(x) \Rightarrow \text{successful}(x))$$

#### Quick check 5

Let's consider the current subgoal. (Note that in this version, variable b is replaced with n/a.)



Define all required predicates and describe a quantified proposition equivalent to the revised statement above.

$$\forall n \forall a [C(n) \land D(a,n) \land R(a,n) \Rightarrow Q(n/a,n)]$$

# Negations of quantified propositions (1)

Let consider a set of positive integers  $\mathbb{Z}^+$  as our universe. Let predicate  $P(x) \equiv \text{``}x \text{ is a prime number.''}$ 

Consider this proposition

$$(\forall x \in \mathbb{Z}^+)P(x).$$

How can we show that this is false?

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When showing that a universally quantified proposition is false, we need to show "one" counter example. In this case, since P(4) is false,  $\forall x P(x)$  is false.

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When showing that a universally quantified proposition is false, we need to show "one" counter example. In this case, since P(4) is false,  $\forall x P(x)$  is false.

This way of disproving a statement is equivalent to showing that

$$(\exists x)(\neg P(x)).$$

## Negations of quantified propositions (2)

Let consider a set of positive integers  $\mathbb{Z}^+$  as our universe. Let predicate  $Q(x) \equiv$  "if x > 2, then  $x^2 \leq 2x$ ." Consider this proposition

$$(\exists x \in \mathbb{Z}^+)Q(x).$$

How can we show that this is false?

$$(\forall x) \neg Q(x)$$



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$$(\exists x \in \mathbb{Z}^+)Q(x).$$

How can we show that this is false? When showing that an existential quantified proposition is false, we need to show that Q(x) is false for every possible values of x. In this case, since  $x^2 = x \cdot x > 2 \cdot x$  for every x > 2, we have that  $(\exists x)Q(x)$  is false.

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## Negations of quantified propositions (3)

#### Thus, the following equivalences:

$$\neg (\forall x P(x)) \equiv \exists x (\neg P(x))$$

$$\neg (\exists x P(x)) \equiv \forall x (\neg P(x))$$

$$\neg (\exists x P(x)) \equiv \forall x (\neg P(x))$$

#### Quick check 6

Consider the following statements with the quantified propositions that you have written previously. Write down their negations in quantified propositional forms, and then translate them back to English sentences.

- Every human must die.
- Some animal eats other animals.
- If a student works hard, that student will be successful.
- Everyone has someone that care about him or her.