# Lab #3: Memory Lab

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# Before getting started

- Do not directly fork the repository!
  - Your code will be visible to others
  - Refer to README on how to clone the repository
  - No pull requests either!

### **Important Dates**

- 9 Apr. Lab Hand-out Session (Today!)
- 16 Apr. Live Q&A Session (1)
- 23 Apr. Live Q&A Session (2)
- 29 Apr. 23:59 Submission Deadline

- Questions about the lab will be conducted through github issue.
   Feel free to post questions (except your code!)
- All Sessions are totally optional.

### **Before the Presentation::**

- All contents on this slide are sourced from the README file.
- For detailed information, please refer to README.
  - https://github.com/SNU-ARC/2024 spring sysprog Lab3/blob/main/README.md

# Goal of this Lab (1/2)

Implement a dynamic memory manager.

#### \$ mm\_driver ./test/demo.dmas

```
Processing './tests/demo.dmas'...
Configuration:
  script file:
                         ./tests/demo.dmas
  configuration:
    implementation:
                         memory manager
    mode:
                         correctness
   freelist policy:
                         implicit
    data segment size:
                         0x2000000 (33554432)
Action: m 1 16
Action: m 2 50
Action: m 3 47
Action: m 4 48
Action: m 5 32
Action: m 6 1000
Action: m 7 10000
Action: v
```

```
mm check ---
ds heap start:
                        0x7fc0f9279000
ds heap brk:
                        0x7fc0f9289000
heap start:
                        0x7fc0f9279020
heap end:
                        0x7fc0f9288fe0
free list policy:
                        Implicit
                                                             0), status: allocated
initial sentinel:
                        0x7fc0f9279018: size:
                                                             0), status: allocated
end sentinel:
                        0x7fc0f9288fe0: size:
                    offset size (hex) size (dec)
 address
                                                      payload status
 0x7fc0f9279020
                       0x0
                                  0x20
                                                           16 allocated
                                  0x60
                                                           80 allocated
 0x7fc0f9279040
                      0x20
 0x7fc0f92790a0
                      0x80
                                  0x40
                                                              allocated
 0x7fc0f92790e0
                      0xc0
                                  0x40
                                                              allocated
 0x7fc0f9279120
                     0x100
                                  0x40
                                                              allocated
 0x7fc0f9279160
                                 0x400
                                              1024
                                                               allocated
                     0x140
 0x7fc0f9279560
                     0x540
                                 0x2720
                                              10016
                                                               allocated
 0x7fc0f927bc80
                    0x2c60
                                0xd360
                                              54112
                                                        54096 free
Block structure coherent.
```

# Goal of this Lab (2/2)

#### You will learn

- how a dynamic memory manager works
- how to implement a dynamic memory manager
- how to work with macros in C
- how to work with function pointers in C
- how to debug code
- that writing a dynamic memory manager is simple in theory and difficult in practice.

# Environment setup (1/3)

- You can get skeleton code and test bench from git repo
  - o git clone https://github.com/SNU-ARC/2024\_spring\_sysprog\_Lab3.git

# **Environment setup (2/3)<Optional>**

- If you want to keep your own repository, keep the lab's visibility to private. Otherwise, others would see your work.
- Changing visibility
  - After cloning the repository, you should change the push remote URL to your own repository.
  - 1. Create an empty repository that you're going to manage (again, keep it private)
  - 2. Copy the url of that repository
  - 3. On your terminal in the cloned directory, type git remote set-url --push origin <repo url>
  - 4. Check with git remote -v if the push URL has changed to yours while the fetch URL remains the same (this repo)

# Environment setup (3/3)

The handout contains the following files and directories.

File/Directory	Description
doc/	Doxygen instructions, configuration file, and auto-generated documentation. Run make doc to generate.
driver/	Pre-compiled modules required to link your implementation to the mm_driver test program
reference/	Reference implementation
src/	Source code of the lab. You will modify memmgr.c/h, mm_test.c
tests/	Test allocation/deallocation sequences to test your allocator
README.md	this file
Makefile	Makefile driver program

File	Description
blocklist.c/h	Implementation of a list to manage allocated blocks for debugging/verification purposes. <b>Do not modify!</b>
datasec.c/h	Implementation of the data segment. <b>Do not modify!</b>
nulldriver.c/h	Implementation of an empty allocator that does nothing. Useful to measure overhead. <b>Do not modify!</b>
memmgr.c/h	The dynamic memory manager. A skeleton is provided. Implement your solution by editing the C file.
mm_test.c	A simple test program to test your implementation step-by-step.

### **Overview**

- Handle memory allocation and management in memmgr.c
  - Implement dynamic memory allocation and management functions such as malloc, realloc and free.
  - Implement both implicit and explicit free lists to track and allocate free blocks using the best fit policy.
  - For explicit free list, use the LIFO(Last-In-First-Out) policy.

### **Overview**

#### We will use a simulated heap

- Real heap is not directly manipulable
- o dataseg.c will give a simulated heap space
- memmgr.c will manage that heap space given by dataseg.c

```
File: mm test.c
   user-level process. After initializing the data
   segment and the heap, mm malloc/free/calloc/
   realloc can be used as in libc.
1. ds allocate() |
                        sequence of mm malloc(),
                            mm free(), mm calloc(),
                               and mm realloc()
            2. mm init()
           custom memory manager. Manages the heap
         File: memmgr.c
                           ds sbrk()
   custom data segment implementation. Manages the
   data segment to be used by our allocator.
 File: dataseg.c
```

- 64-bit operator
  - One word = 8 bytes
- Minimal Block size = 32 bytes
  - All block addresses should be aligned to 32 bytes
- Each block should have boundary tags
  - Header & Footer
  - Bit 0 of each boundary tag should indicate the status of the block
     (1: allocated, 0: free)

#### mm\_malloc(size)

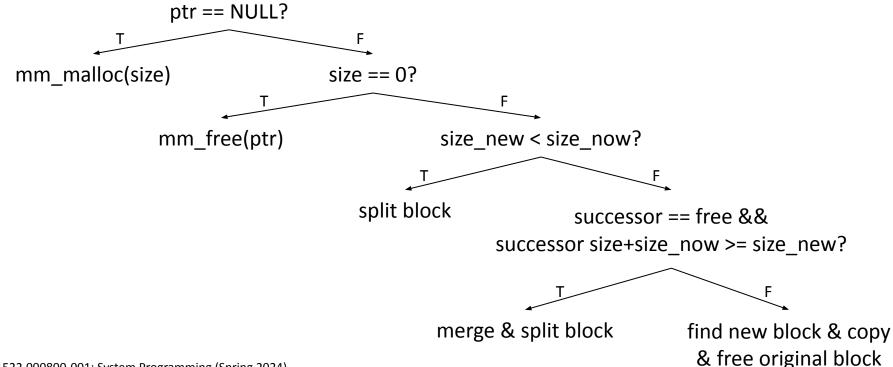
- Returns a pointer to an allocated payload block of at least size bytes.
- The entire allocated block must lie within the heap region and must not overlap with any other block.

#### mm\_free(ptr)

- Deallocate the block pointed to by ptr that was returned by an earlier call to mm\_malloc(), mm\_calloc(), or mm\_realloc() and has not yet been freed.
- When the caller attempts to release a memory block that has already been freed, an error message is generated.

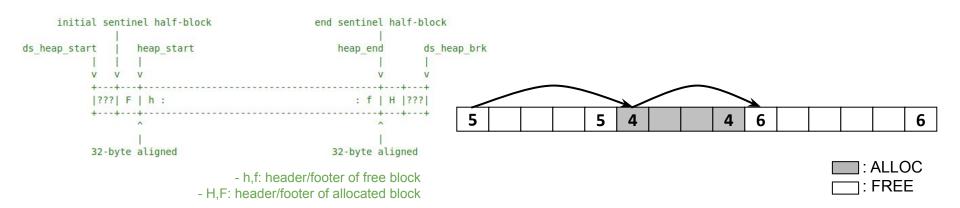
- mm\_realloc(ptr, size)
  - If (ptr == NULL), then mm malloc(size).
  - o If (size == 0), then mm\_free(ptr).
  - If (ptr != NULL) then it must point to a valid allocated block.
    - if size(new) < size pointed by ptr(current)
      - the block should be split into allocated block and a new free block.
    - if size(new) > size pointed by ptr(current)
      - if successor block is free and large enough when merged, increase the block size and fix the remaining free block.
      - if not, assign a new block by the allocation policy(best-fit) and copy the original contents. Afterwards, the original block must be freed.

mm\_realloc(ptr, size)

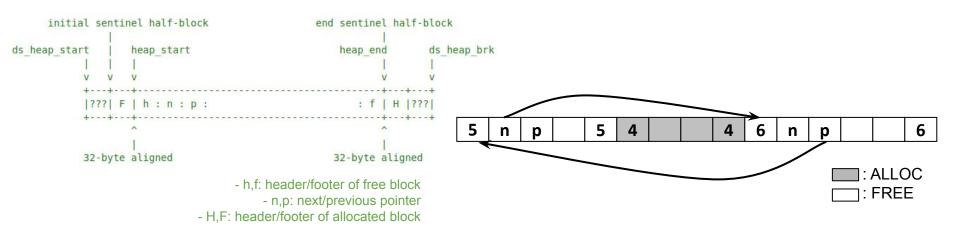


- Best-fit policy
  - Select the smallest available block that is large enough
  - Aims to reduce fragmentation and optimize space utilization

- Support different free list managements Implicit free list
  - Utilize the spaces within the memory itself to track free blocks



- Support different free list managements Explicit free list
  - Create a linked list of free blocks
  - Each block points to the next and previous block

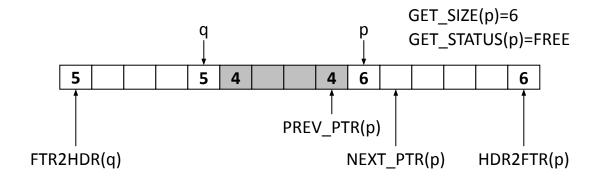


# Code and Test (1/3)

- The skeleton provides some useful tools such as
  - macros for easy data manipulation
  - mm\_check for checking and dumping heap status
  - functions for logging information
- You have to implement the following two parts:
  - o in mm\_init(), mm\_malloc(), mm\_realloc(), mm\_free()
    - Main functions to allocate and manage blocks.
    - Functions should be compatible with both implicit and explicit free list.
  - o in bf\_get\_free\_block\_implicit(), bf\_get\_free\_block\_explicit()
    - Track and allocate a free block based on the best fit policy within each free list method.
    - For peak utilization, search for free blocks using the above principles.

# Code and Test (1/3)

Macro examples



# Code and Test (2/3)

- mm\_test
  - An interactive testing tool to check your implementation
  - Reference mm\_test is given, but do not take the outputs literally.
  - Build with \$ make mm\_test

```
Select freelist policy.

(i) implicit list
(e) explicit list
(q) quit
Your selection:

(m) malloc
(f) free
(c) check heap
(l) set log level
(q) quit
Your selection:
```

```
example) m 16 \rightarrow m 130 \rightarrow m 31 \rightarrow m 5 \rightarrow m 1023 \rightarrow f
 (2) Block list
      0] 0x7f8f1a33f028: size:
                                                  16), status: allocated
                                                 130). status: allocated
      1] 0x7f8f1a33f048: size:
         0x7f8f1a33f0e8: size:
                                                  31), status: allocated
                                        5 (
                                                   5), status: allocated
      3| 0x7f8f1a33f128: size:
                                               1023), status: allocated
      41 0x7f8f1a33f148: size:
                                                           in this case
Enter index/indices of blocks to free:
                                                           we choose 2
  ds heap start:
                        0x7f8f1a33f000
  ds heap brk:
                        0x7f8f1a34f000
  heap start:
                        0x7f8f1a33f020
                        0x7f8f1a34efe0
  heap end:
  free list policy:
                        Implicit
  initial sentinel:
                        0x7f8f1a33f018: size:
                                                           0), status: allocated
  end sentinel:
                        0x7f8f1a34efe0: size:
                                                           0), status: allocated
                           size (hex) size (dec)
    address
                     offset
                                                    payload status
    0x7f8f1a33f020
                       0x0
                                  0x20
                                               32
                                                         16 allocated
    0x7f8f1a33f040
                       0x20
                                  0xa0
                                                        144 allocated
    0x7f8f1a33f0e0
                       0xc0
                                  0x40
                                               64
                                                            free
    0x7f8f1a33f120
                      0x100
                                  0x20
                                               32
                                                         16
                                                            allocated
    0x7f8f1a33f140
                      0x120
                                 0x420
                                             1056
                                                      1040
                                                            allocated
    0x7f8f1a33f560
                      0x540
                                0xfa80
                                            64128
                                                      64112 free
  Block structure coherent.
```

# Code and Test (3/3)

- mm\_driver
  - Driver program that issues allocations and free requests from script
  - Can print stats at the end
  - Build with \$ make mm\_driver

```
$ mm_driver --help
```

```
Syntax: mm_driver [--dssize <size>] [--implementation <impl>] [--mode <mode>] [--policy <policy>] [--statfile <file>] [--help] <script(s)>
  --dssize <size>
                            set size of datasegment to <size>
  --implementation <impl>
                            select implementation from one of
                                            vour implementation
                                            C standard library
                                            empty implementation (to measure overhead)
  --mode <mode>
                             set mode to one of
                              performance minimize output and measure performance
                              correctness perform extra correctness checks
                                            print lots of output on what's going on
  --policy <policy>
                             set freelist policy to
                              implicit
                                            implicit freelist policy
                                            explicit freelist policy
                            Note: the policy setting only has an effect on the mmemmgr implementation
  --statfile <file>
                            write statistics in CSV format to <file>
  --help
                            this screen
  <script(s)>
                            one or more .dmas scripts
Note: settings given on the command line override settings in the script files.
```

# Code and Test (3/3)

- mm\_driver
  - If you want to own test, you can copy \*.dmas and change test sequence.
  - You can change the policy by '\$ mm\_driver --policy <policy>' or by changing 'heap <policy>' in \*.dmas

#### \$ mm\_driver ./test/demo.dmas

```
Statistics:
                        17
 Actions:
    malloc:
    calloc:
    realloc:
    free:
  Utilization:
    payload:
                        30 (
                                   48) bytes
   heap size:
                     10000 (
                                65536) bytes
   utilization:
                         0.1%
   #sbrk():
                          1 times
  Performance:
    total time:
                          0.000212 sec
   throughput:
                        79.97 kops/sec
```

- test sequence in directory 'test'
  - o 100K.dmas
    - 100K allocation for performance comparison of explicit and implicit free list.
  - o alloc.dmas
    - 2048 malloc for normal operation without error.
  - o demo.dmas
    - Normal allocation for checking performance.
  - o Is.dmas
    - Memory sequence of 1s -R

# **Grading Policy**

- Test bench : 80 %
  - Hidden test cases will be used along with the given example tests
- Report : 20 %
  - Briefly explain your code including following information
    - How and when to increase/decrease the heap size
    - How to implement an explicit free list
    - (If any) Explain the added macros
  - Show and analyse the performance difference between implicit and explicit free list
- For late submission:
  - A deduction of 20%p per 24 hours

# Submission (via eTL)

- Write-up
  - Briefly describe your implementation.
  - Filename: [student\_id].pdf (example: 2024-12345.pdf)
  - Please submit it in pdf format. Other formats are not accepted.
- Compress your source code and write-up into a single file
  - Compress memmgr.c and your report
  - Filename should be [student id].tar (example: 2024-12345.tar).
  - Please submit it in tar format. Other formats are not accepted.
  - Refer to README.md for submission instructions.
- Submission deadline: by 23:59 on April 29, 2024

# **Question?**